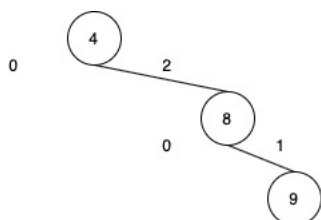
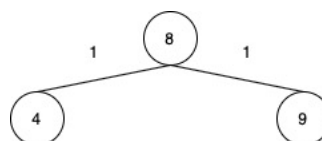


Question 1

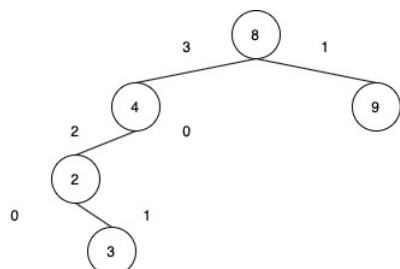
a)



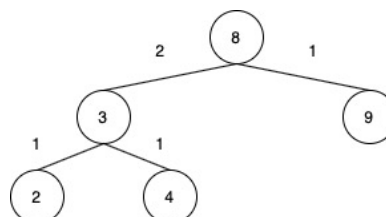
After insert until node 9, the tree is right-right imbalanced, thus, left rotation on node 9 is used



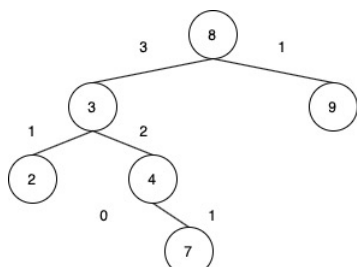
After update the height, the tree is balanced again



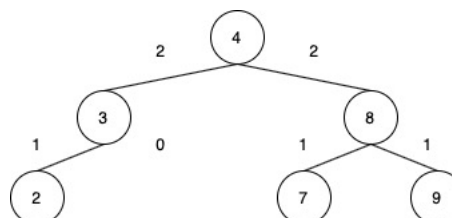
After insert until node 3, the tree is left-right imbalanced, thus, double rotation is used



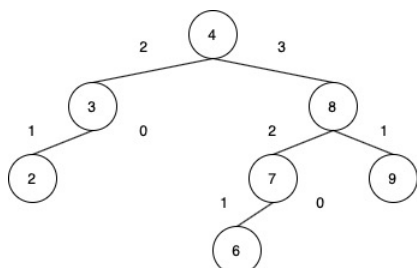
After update the height, the tree is balanced again



After insert until node 7, the tree is left-right imbalanced, thus, double rotation is used

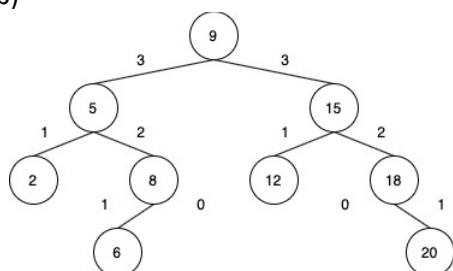


After update the height, the tree is balanced again

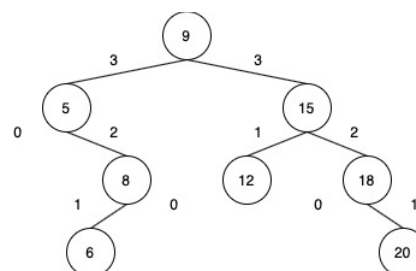


The insertion is finished

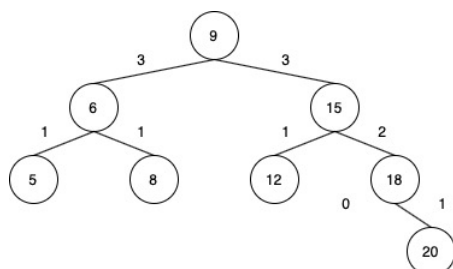
b)



Delete node 2 in the AVL tree



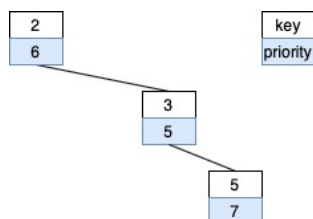
The AVL tree suffers from right-left imbalance, thus, double rotation is used



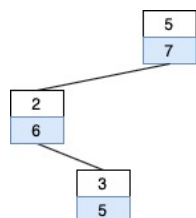
After update the height, the AVL tree is balanced again
 The deletion is finished

Question 2

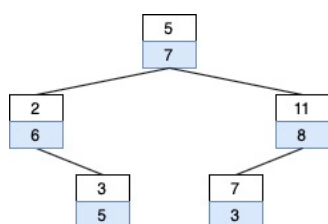
a)



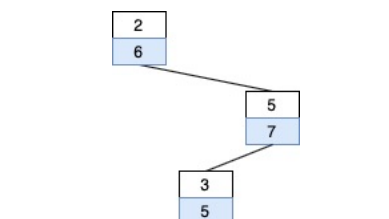
After insert until node 5 with priority 7, it violates the max tree property, thus, move node 5 upward



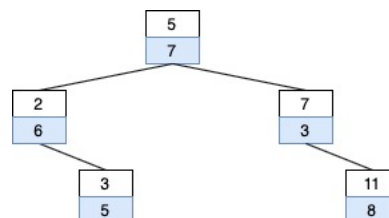
After rotation, there is no more violation



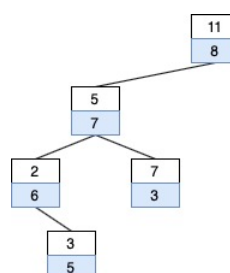
After rotation, it still violates the max tree property, thus, move node 11 upward again



After rotation, it still violates the max tree property, thus, move node 5 upward again

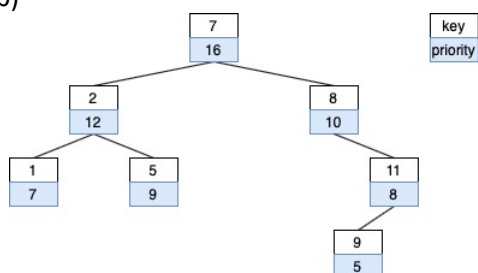


After insert until node 11 with priority 8, it violates the max tree property, thus, move node 11 upward

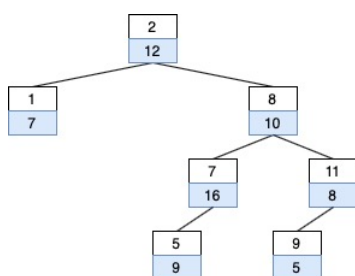


After rotation, there is no more violation
The insertion is finished

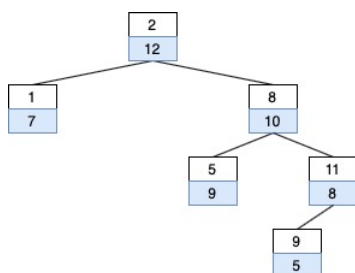
b)



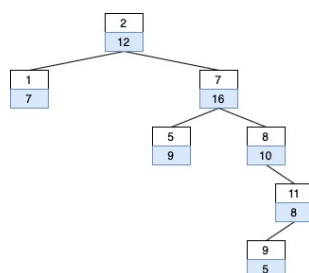
Attempt to delete node 7 with priority 16, we rotate node 2, which has a larger priority, upward



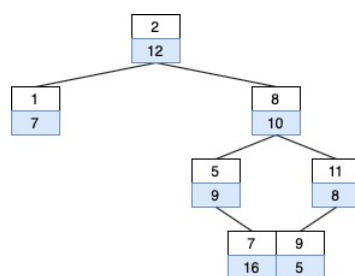
We rotate node 5 with priority 9 upward



The deletion is finished



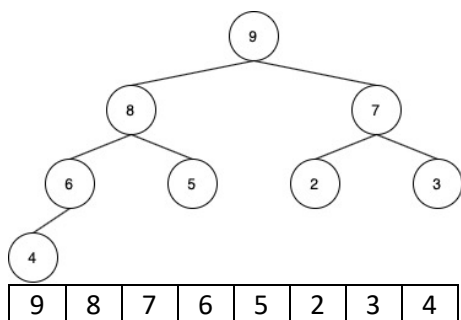
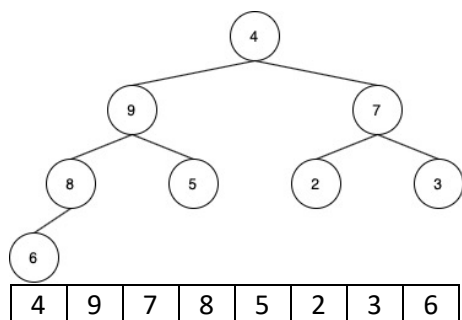
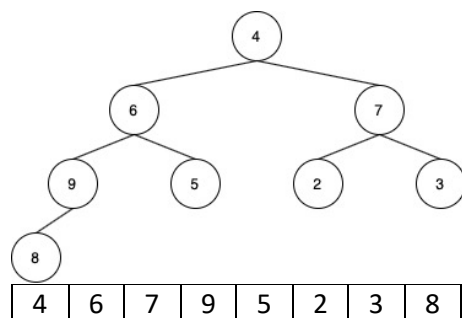
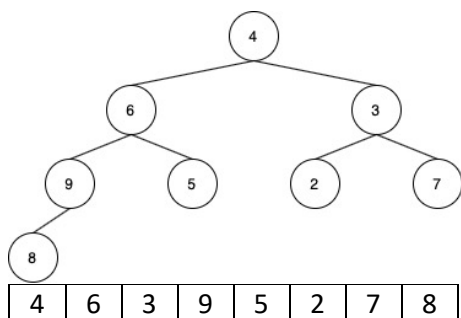
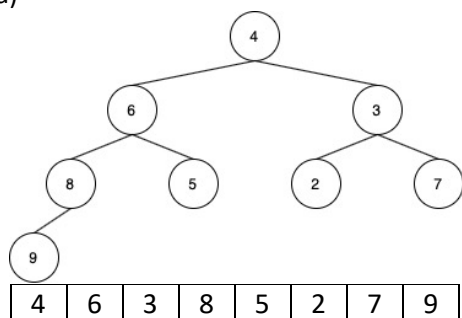
Again, we rotate node 8 upward as it has the larger priority



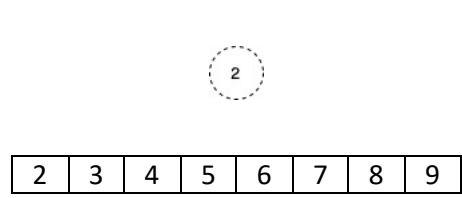
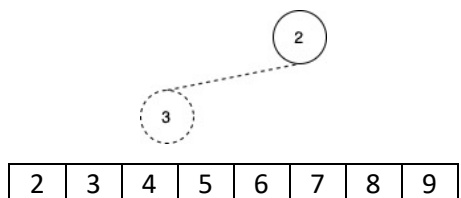
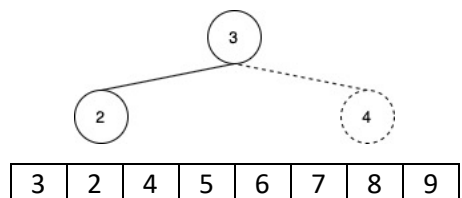
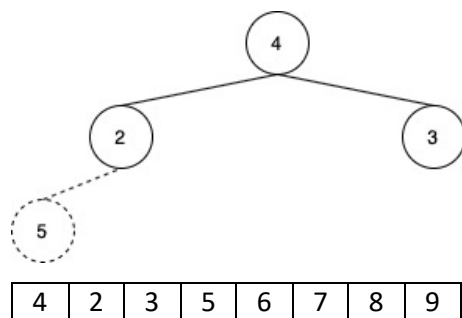
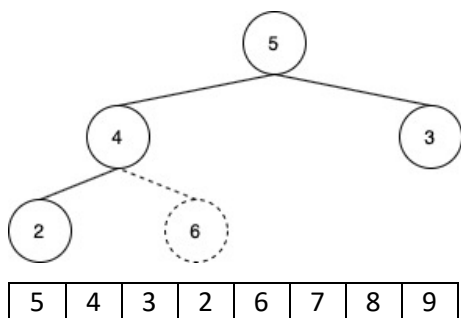
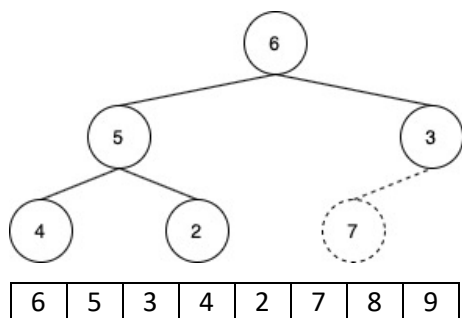
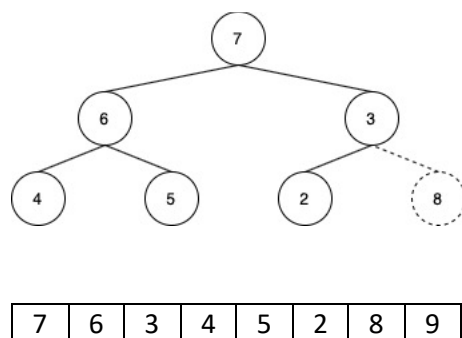
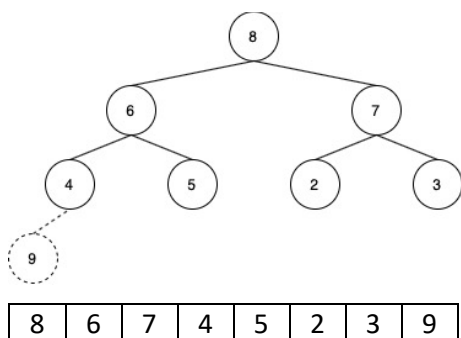
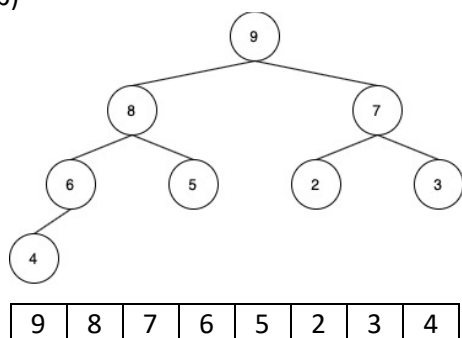
Node 7 became a leaf node, and we remove it

Question 3

a)

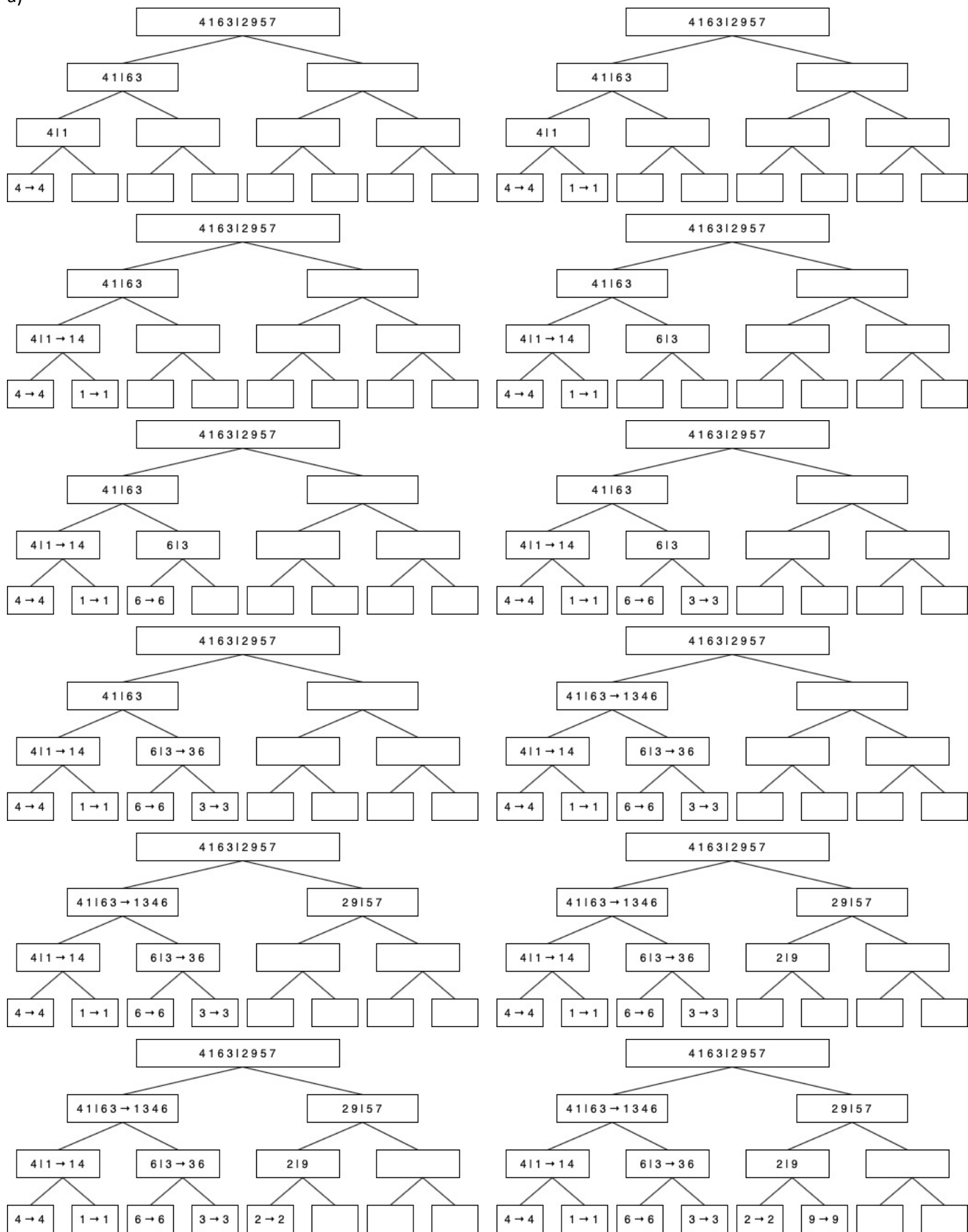


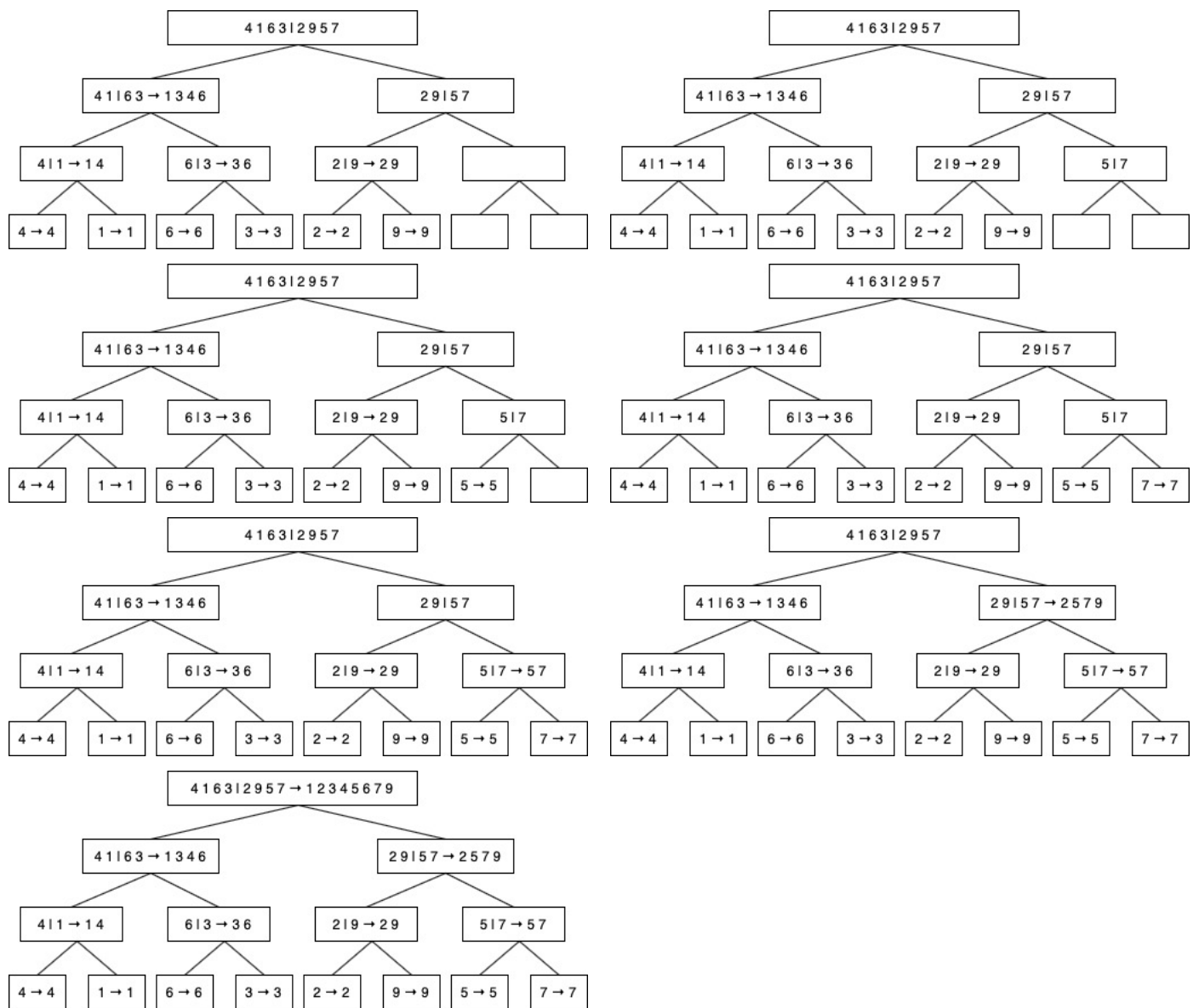
b)



Question 4

a)





b)

1	3	6	8	2	5	7	9
---	---	---	---	---	---	---	---

left = 0

left = 1

left = 1

left = 2

left = 2

left = 3

left = 3

left = 4

secondleft = 4

secondleft = 4

secondleft = 5

secondleft = 5

secondleft = 6

secondleft = 6

secondleft = 7

secondleft = 7

Temporary								
1								
1	2							
1	2	3						
1	2	3	5					
1	2	3	5	6				
1	2	3	5	6	7			
1	2	3	5	6	7	8		
1	2	3	5	6	7	8	9	

c)

6	3	5	2	7	4	1	8
6	3	8	2	7	4	1	5
3	6	8	2	7	4	1	5
3	2	8	6	7	4	1	5
3	2	4	6	7	8	1	5
3	2	4	1	7	8	6	5
3	2	4	1	5	8	6	7

Swap the pivot data and the last data

nextsmallpos = 0

nextsmallpos = 1

nextsmallpos = 2

nextsmallpos = 3

nextsmallpos = 4

Swap the last data and the data in *nextsmallpos*

Question 5

Array A									
6	1	2	1	3	4	6	1	3	2

counter					
3	2	2	1	0	2
3	5	7	8	8	10

Array B									
				2					
				2		3			
		1		2		3			
		1		2		3			6
		1		2		3	4		6
		1		2	3	3	4		6
	1	1		2	3	3	4		6
	1	1	2	2	3	3	4		6
1	1	1	2	2	3	3	4		6
1	1	1	2	2	3	3	4	6	6

counter					
3	4	7	8	8	10
3	4	6	8	8	10
2	4	6	8	8	10
2	4	6	8	8	9
2	4	6	7	8	9
2	4	5	7	8	9
1	4	5	7	8	9
1	3	5	7	8	9
0	3	5	7	8	9
0	3	5	7	8	8

Question 6

By counting sort, it builds up an array A for counter. The sort cost (pre-processing time) is $O(n + k)$ for integers are ranged in $[1, k]$. The number of integers fall within a range of $[a, b]$ is $A[b] - A[a]$, such algorithm takes $O(1)$ time

Question 7

a)

4650		4650		3918		3152		2465
7391		9830		9830		7391		3152
5479	Sort on	7391	Sort on	4650	Sort on	2465	Sort on	3918
3152	digit 1	3152	digit 2	3152	digit 3	5479	digit 4	4650
2465		2465		2465		4650		5479
3918		3918		5479		9830		7391
9830		5479		7391		3918		9830

b) We can represent numbers in range of $[1, n^3 - 1]$ with digit of base n , such that $a \cdot n^2 + b \cdot n + c$, where $a, b, c \in [1, n^3 - 1]$. By radix sort, $k = n$ and $d = c$, such that $O(d \cdot (n + k)) = O(c \cdot (n + n)) = O(n)$

Question 8

a) The expected number of the basic operation `random(left, right)` executed from Lines 1 to 7

$$E(X) = 2 \cdot \sum_{i=1}^{\infty} \frac{i}{3^i} = 6, \text{ given that } Pr(X = 3i) = \frac{1}{3^i}$$

For `partition` takes $O(n)$ and others take constant time

Thus, the expected time complexity for Lines 1 to 7 is $O(n)$

$$\begin{aligned} b) \quad T(n) &\leq cn + \frac{1}{n} \sum_{i=1}^n \max(T(i-1), T(n-i)) \\ &\leq cn + \frac{2}{n} \sum_{i=n/2}^{n-1} T(i) \\ &\leq cn + \frac{2}{n} \left[T(n-1) + T(n-2) + \dots + T\left(\frac{n}{2}\right) \right] \\ &\leq 4nc \end{aligned}$$

\therefore the time complexity is $O(n)$