嵌入式作業系統實作 Embedded OS Implementation

PA_1

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[PART I] Task Control Block Linked List [20%]

• The screenshot results. (10%)

- A report that describes your implementation (please attach the screenshot of the code and MARK the modified part). (10%)
- 1. 上半部顯示 task 被創建時的 link,這段程式碼加在 OS_TCBInit 下,因為在執行 OSTaskCreateExt 建立 Task 內會執行到 OS_TCBInit 這個 function,OS_TCBInit 會初始化 task 的 OS_TCB 的數值,所以可以在這裡 抓到 task 被建立時的 link 數值。

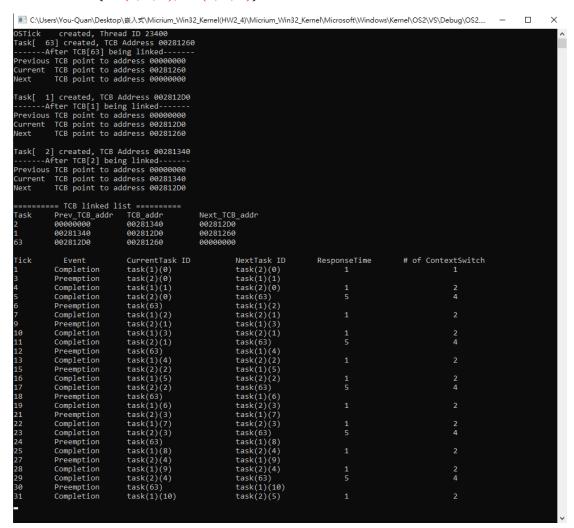
```
//OS_TCBInit會初始化task的OS_TCB所需要的變數
         //所以在這個地方print作業所需要的數值
         if (ptcb->OSTCBId == OS_TASK_IDLE_ID) //假如是idle_task的話就顯示優先權
            printf("Task[ \ \%d] \ created, \ TCB \ Address \ \%p\n", \ ptcb->OSTCBPrio, \ ptcb);
            printf("-----After TCB[%d] being linked-----\n", ptcb->OSTCBPrio);
            printf("Previous TCB point to address %p\n", ptcb->OSTCBPrev);
            printf("Current\t TCB point to address %p\n", ptcb);
            printf("Next\t TCB point to address %p\n", ptcb->OSTCBNext);
            printf("\n");
Ιė
            printf("Task[ %d] created, TCB Address %p\n", ptcb->OSTCBId, ptcb);
            printf("Previous TCB point to address %p\n", ptcb->OSTCBPrev);
            printf("Current\t TCB point to address %p\n", ptcb);
            printf("Next\t TCB point to address %p\n", ptcb->OSTCBNext);
            printf("\n");
```

2. 下半部顯示 linked list 的表,取 OSTCBList 做開頭,使用一個 while 一直指向下一個 link 的 OS_TCB, 然後顯示出所有 task 的 link 情況。

[PART II] RM Scheduler Implementation [80%]

• The screenshot results (with the given format) of four task sets. (Time ticks 0-30 or miss deadline). (40%)

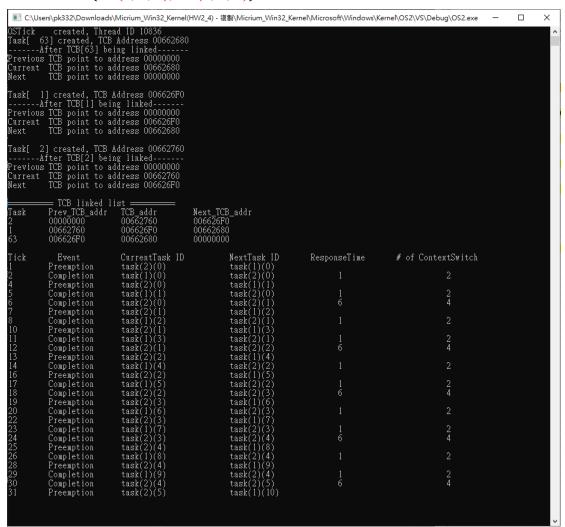
Task set $1 = \{\tau 1 \ (0, 1, 3), \tau 2 \ (0, 3, 6)\}$



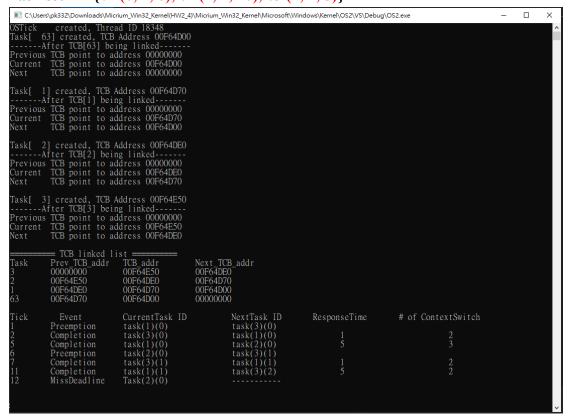
Task set $2 = {\tau 1 (0, 8, 15), \tau 2 (0, 2, 5)}$

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OSTick	created, Thr	ead ID 22624				^
		B Address 00AF1260				
	-After TCB[63] b	eing linked				
		address 00000000				
		address 00AF1260				
Next	TCB point to	address 00000000				
		Address 00AF12D0				
		ing linked				
		address 000000000 address 00AF12D0				
Next		address 00AF1200				
Next	ICB POINT TO	auuress 00AF1200				
Task[21 created, TCB	Address 00AF1340				
		ing linked				
		address 00000000				
		address 00AF1340				
Next	TCB point to	address 00AF12D0				
ž						
	==== TCB linked					
Task	Prev_TCB_addr		Next_TCB_addr			
2	00000000	00AF1340	00AF12D0			
1 63	00AF1340	00AF12D0	00AF1260			
03	00AF12D0	00AF1260	00000000			
Tick	Event	CurrentTask ID	NextTask ID	ResponseTime	# of ContextSwitch	
2	Completion	task(2)(0)	task(1)(0)	2	1	
5 7	Preemption	task(1)(0)	task(2)(1)			
7	Completion	task(2)(1)	task(1)(0)			
10	Preemption	task(1)(0)	task(2)(2)			
12	Completion	task(2)(2)	task(1)(0)	2		
14	Completion	task(1)(0)	task(63)	14		
15	Preemption	task(63)	task(2)(3)			
17	Completion	task(2)(3)	task(1)(1)	2		
20 22	Preemption	task(1)(1)	task(2)(4)	2	2	
25	Completion Preemption	task(2)(4) task(1)(1)	task(1)(1) task(2)(5)	2	Z	
27	Completion	task(1)(1) task(2)(5)	task(2)(3) task(1)(1)	2	2	
29	Completion	task(1)(1)	task(63)	14	6	
30	Preemption	task(63)	task(2)(6)	3. 1		
32	Completion	task(2)(6)	task(1)(2)	2	2	
						V .

Task set $3 = \{\tau 1 (1, 1, 3), \tau 2 (0, 4, 6)\}$



Task set $4 = \{\tau 1 \ (0, 4, 6), \tau 2 \ (2, 2, 10), \tau 3 \ (1, 1, 5)\}$



• A report that describes your implementation (please attach the screenshot of the code and MARK the modified part). (40%)

實現方式條列簡述說明:

1. 在OS_TCB中新增變數

INT32U begin_ready_time; //紀錄task變ready的時間點

INT32U response; //反應時間 INT32U arrival; //到達時間 INT32U execution; //執行時間

INT32U period; //週期

INT32U job_id; //工作次數

- 2. 在OSTimeTick中去察看目前的ReadyTable有哪些Task正在Ready狀態,把所有正在Ready且不是目前正在執行的task的所有task的response都加一,因為task正在Ready卻沒有被執行表示被延後了一個tick,response一開始初始值是週期內的執行時間,所以task工作執行只要沒有執行滿response的時間就會一直卡在while迴圈內。
- 3. Task執行完會執行OS_Sched()去切換給下個Task,所以在OS_Sched()中顯示 完成的字串。
- 4. 當TimeTick中斷產生之後會進OSIntExit()把低優先權的task中斷給高優先權的task,所以在這裡顯示中斷狀態的字串。
- 5. 當有task已經快要完成時有可能會被其他優先權高的task給搶占,所以會在 OSIntExit()增加判斷程式迴避掉那次的context switch,讓快要做完的 task先完成它的工作。
- 6. 當有一個task完成了一個周期內的工作後要做下個週期的工作時因為不會進入OS_Sched()做context switch,所以在OSTimeDly(0)時表示task的反應時間等於週期,在OSTimeDly內去print完成的字串。
- 7. 判斷是否有task已經miss deadline,只要在OSTimeTick中去判斷reponse是 否大於period,表示說反應時間已經大於週期就是miss deadline。
- 8. 以下是程式的截圖以及內部程式碼中較詳細的註解說明

Task set 1	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30
Task1 (0,1,3)	0			1			2			3			4			5			6			7			8			9			10
Task2 (0,3,6)			0					1						2						3						4					5
Scheduling Result	0		0	1	0		2		1	3	1		4		2	5	2		6		3	7	3		8		4	9	4		10
Task set 2		t																													
Task1 (0,8,15)					0															1											2
Task2 (0,2,5)	0					1					2					3					4					5				6	
Scheduling Result	ult 0			0			1		0		- 3	2		0			3		1			4		1			5		1		6
Task set 3																															
Task1 (1,1,3)			0		1			2			3			4			5			6			7			8			9		
Task2 (0,4,6)	0		0						1						2						3						4				5
Scheduling Result	0		0	0	1	0	1	2		1	3	1	2	4		2	5	2	3	6		3	7	3	4	8		4	9	4	5
Task set 4																															
Task1 (0,4,6)			0						1						2						3						4				5
Task2 (2,2,10)				0										1	Г									2							
Task3 (1,1,5)			0				1	L				2					3					4					5				
Scheduling Result	0		0	0		0	1			1		2			2		3	0		3		4	3	1		4	5		4	2	

Main.c

```
static woid task2(woid* p_arg);
   * LOCAL GLOBAL VARIABLES
   static OS_STK StartupTaskStk[APP_CFG_STARTUP_TASK_STK_SIZE];
   #define TASK_STACKSIZE 2048
static OS_STK Task1_STK[TASK_STACKSIZE];
   static OS_STK Task2_STK[TASK_STACKSIZE];
static OS_STK Task3_STK[TASK_STACKSIZE];
#define TASK1_ID 1 //定義task1的id、priority、arrival、execution、period
   #define TASK1_PRIORITY 2
   #define TASK1_ARRIVAL 0
   #define TASK1_EXECUTION 4
   #define TASK1_PERIOD 6
#define TASK2_ID 2 //定義task2的id、priority、arrival、execution、period #define TASK2_PRIORITY 3
  #define TASK2_ARRIVAL 2
#define TASK2_EXECUTION 2
   #define TASK2_PERIOD 10
■ #define TASK3_ID 3 //定義task3的id、priority、arrival、execution、period
#define TASK3_PRIORITY 1
#define TASK3_ARRIVAL 1
   #define TASK3_EXECUTION 1
#define TASK3_PERIOD 5
```

```
□#ifdef TASK1_ID //假如task1有被定義的話 就創建task1工作
      OSTaskCreateExt(task1,
          &Task1_STK[TASK_STACKSIZE - 1u],
          TASK1_PRIORITY,
          TASK1_ID,
          &Task1_STK[Ou],
          TASK_STACKSIZE,
          (OS_TASK_OPT_STK_CHK | OS_TASK_OPT_STK_CLR),
          TASK1_ARRIVAL,
          TASK1_EXECUTION,
          TASK1_PERIOD);
  #endif // TASK1_ID
□#ifdef TASK2_ID //假如task2有被定義的話 就創建task2工作
      OSTaskCreateExt(task2,
          0,
          &Task2_STK[TASK_STACKSIZE - 1u],
          TASK2_PRIORITY,
          TASK2_ID,
          &Task2_STK[Ou],
          TASK_STACKSIZE,
          (OS_TASK_OPT_STK_CHK | OS_TASK_OPT_STK_CLR),
          TASK2_ARRIVAL,
          TASK2_EXECUTION,
          TASK2_PERIOD);
  #endif // TASK2_ID
白#ifdef TASK3_ID //假如task3有被定義的話 就創建task3工作
      OSTaskCreateExt(task3,
          &Task3_STK[TASK_STACKSIZE - 1u],
          TASK3_PRIORITY,
          TASK3_ID,
          &Task3_STK[Ou],
          TASK_STACKSIZE,
          Ou,
          (OS_TASK_OPT_STK_CHK | OS_TASK_OPT_STK_CLR),
          TASK3_ARRIVAL,
          TASK3_EXECUTION,
          TASK3_PERIOD);
      OSTimeSet(0); //重新歸零timetick
```

ucos_ii.h

os_task.c

os core. c

```
INT8U OS_TCBInit (INT8U
                    OS_STK *ptos,
                    OS_STK *pbos,
                    INT16U
                    INT32U
                            stk_size,
                            *pext,
                    INT16U opt,
                    INT32U
                     INT32U
ΙÞ
                             period) //在OS_TCBInit內新增變數 arrival、exe
                     INT32U
     OS_TCB
               *ptcb;
 ⊟#if OS_CRITICAL_METHOD == 3u
      OS_CPU_SR cpu_sr = Ou;
  #endif
 ⊨#if OS_TASK_REG_TBL_SIZE > Ou
      INT8U
  #endif
 □#if OS_TASK_CREATE_EXT_EN > Ou
 □#if defined(OS_TLS_TBL_SIZE) && (OS_TLS_TBL_SIZE > Ou)
  #endif
  #endif
      OS_ENTER_CRITICAL();
      ptcb = OSTCBFreeList;
      if (ptcb != (OS_TCB *)0) {
          OSTCBFreeList
                                  = ptcb->OSTCBNext;
                                                          /* Update poin
          OS_EXIT_CRITICAL();
          ptcb->begin_ready_time = arrival; //初始化在OS_TCB內新增的變數
          ptcb->response = execution;
          ptcb->arrival = arrival;
          ptcb->execution = execution;
          ptcb->period = period;
          ptcb->job_id = 0;
```

```
⊟void OS_Sched (void)
 ⊟#if OS CRITICAL METHOD == 3u
           OS_CPU_SR cpu_sr = Ou;
                      if (OSLockNesting == Ou) {
                              OS_SchedNew();
                               OS_SchedNew();
OSTCBHighRdy = OSTCBPrioTbl[OSPrioHighRdy];
OSTCBHighRdy = OSTCBPrioChr) { /* No Ctx Sw if current task is highest rdy
                                         if (OSTimeGet() > OSTCBCur->arrival)
                                                  //task完成之後會執行OS_Sched所以在這裡print完成
//假設是task完成後切入idle task,print的方式會不
                                                  if(OSTCBHighRdv->OSTCBPrio == OS TASK IDLE PRIO)
                                                           printf("%d\t Completion\t task(%d)(%d)\t\t task(%d) \t\t %d\t\t %d\n", OSTimeGet(),
                                                                   OSTCBCur->OSTCBId, OSTCBCur->job_id,
                                                                    OS_TASK_IDLE_PRIO,
                                                                    OSTimeGet() - OSTCBCur->begin_ready_time, OSTCBCur->OSTCBCtxSwCtr);
                                                           printf(\mbox{\em sk}(\mbox{\em sk})(\mbox{\em sk})\mbox{\em sk}(\mbox{\em sk}(\mbox{\em sk})\mbox{\em sk}(\mbox{\em sk})\mbox{\em sk}(\mbox{\em sk})\mbox{\em sk}(\mbox{\em sk}(\mbox{\em sk})\mbox{\em sk}(\mbox{\em sk}(\mbox{\em sk})\mbox{\em sk}(\mbox{\em sk}(\mbox{\em sk})\mbox{\em sk}(\mbox{\em sk}(\mbox{\em sk})\mbox{\em sk}(\mbox{
                                                                    {\tt OSTCBHighRdy->OSTCBId,\ OSTCBHighRdy->job\_id,}\\
                                                                    OSTimeGet() - OSTCBCur->begin_ready_time, OSTCBCur->OSTCBCtxSwCtr);
                                                  OSTCBCur->begin_ready_time = OSTimeGet();
                                                  OSTCBCur->response = OSTCBCur->execution;
                                                  OSTCBCur->job_id += 1;
⊟#if OS TASK PROFILE EN > Ou
                                             OSTCBHighRdy->OSTCBCtxSwCtr++;
   ⊟void OSIntExit (void)
 =#if OS_CRITICAL_METHOD == 3u
     #endif
              if (OSRunning == OS_TRUE) {
   OS_ENTER_CRITICAL();
                       if (OSIntNesting > Ou) {
                        if (OSIntNesting == Ou) {
]þ
                                if (OSLockNesting == Ou) {
                                         if (OSTimeGet() - OSTCBCur->begin_ready_time != OSTCBCur->response)
                                                  OS_SchedNew();
                                                  OS_SchedNew();
OSTCBHighRdy = OSTCBPrioTbl[OSPrioHighRdy];

/* No Ctx Sw if current task is highest rdy */
                                                            //假設是idle task被搶占,print的方式會不一樣
if (OSTCBCur->OSTCBPrio == OS_TASK_IDLE_PRIO)
                                                                    printf("%d\t Preemption\t task(%d) \t\t task(%d)(%d)\n", OSTimeGet(),
    OS_TASK_IDLE_PRIO,
                                                                              OSTCBHighRdy->OSTCBId, OSTCBHighRdy->job_id);
                                                                     printf("\%d\t Preemption\t task(\%d)(\%d)\t task(\%d)(\%d)\n", \ OSTimeGet(),
                                                                              OSTCBCur->OSTCBId, OSTCBCur->job_id,
                                                                              {\tt OSTCBHighRdy->OSTCBId}, \ {\tt OSTCBHighRdy->job\_id});
 ⊞#if OS_TASK_PROFILE_EN > Ou
                                                            OSTCBHighRdy->OSTCBCtxSwCtr++;
```

os_time.c

```
□void OSTimeDly (INT32U ticks)
□#if OS_CRITICAL_METHOD == 3u
     OS_CPU_SR cpu_sr = Ou;
      if (OSIntNesting > Ou) {
                                                 /* See if trying to call from an ISR
      if (OSLockNesting > Ou) {
                                                  /* 0 means no delay!
      if (ticks > Ou) {
          OS_ENTER_CRITICAL();
                      = OSTCBCur->OSTCBY;
          OSRdyTbl[y] &= (OS_PRIO)~OSTCBCur->OSTCBBitX;
          OS_TRACE_TASK_SUSPENDED(OSTCBCur);
          if (OSRdyTbl[y] == Ou) {
             OSRdyGrp &= (OS_PRIO)~OSTCBCur->OSTCBBitY;
          OSTCBCur->OSTCBDly = ticks;
          OS_TRACE_TASK_DLY(ticks);
          OS_EXIT_CRITICAL();
          OS_Sched();
      //當發生執行 OSTimeDly(0) 的情況表示反應時間剛好是週期時間
//task已經完成這次週期的工作要做下個週期的工作
₽
if (OSTCBCur->period == OSTCBCur->response)
          printf("%d\t Completion\t task(%d)(%d)\t\t task(%d)(%d)\t\t %d\t\t\ %d\n", OSTimeGet(),
             OSTCBCur->OSTCBId, OSTCBCur->job_id,
              OSTCBCur->OSTCBId, OSTCBCur->job_id + 1,
              OSTimeGet() - OSTCBCur->begin_ready_time, OSTCBCur->OSTCBCtxSwCtr);
          //初始化task的OS_TCB變數參數
          OSTCBCur->begin_ready_time = OSTimeGet();
          OSTCBCur->response = OSTCBCur->execution;
          OSTCBCur->OSTCBCtxSwCtr = 0;
          OSTCBCur->job_id += 1;
```

os_cpu_c. c

```
624 [:
625 □≠if (OS_MSG_TRACE > Ou)
626 ■ [: //OS_Printf("Task[%3.1d] created, Thread ID %5.0d\n", p_tcb->OSTCBPrio, p_stk->ThreadID); //因為作業不需要print這行 所以註解掉這行
627   #endif
```

os_cpu_c.c \ main.c \ os_time.c \ ucos_ii.h \ os_task.c \ os_core.c