注册表回调是一个比较监控注册表读写的回调,它的能量非常大,一个回调能实现在 SSDT 上 HOOK 十几个 API 的效果。部分游戏保护还会在注册表回调上做功夫,监控 service 键的子键,实现双层拦截驱动加载(在映像回调那里还有一层)。而在卡巴斯基等 HIPS 类软件里,则用来监控自启动等键值。

注册表回调在 XP 系统上貌似是一个数组,但是从 WINDOWS 2003 开始,就变成了一个链表。这个链表的头称为 CallbackListHead,可在 CmUnRegisterCallback 中找到:

```
lkd> uf CmUnRegisterCallback
nt!CmUnRegisterCallback:
fffff800`01cba790 48894c2408
                                          qword ptr [rsp+8], rcx
                                  mov
fffff800`01cba795 53
                                          rbx
                                  push
fffff800`01cba796 56
                                  push
                                          rsi
fffff800`01cba797 57
                                  push
                                          rdi
ffffff800`01cba798 4154
                                          r12
                                  push
ffffff800`01cba79a 4155
                                  push
                                          r13
ffffff800`01cba79c 4156
                                  push
                                          r14
fffff800`01cba79e 4157
                                          r15
                                  push
fffff800`01cba7a0 4883ec60
                                  sub
                                          rsp, 60h
fffff800`01cba7a4 41bc0d0000c0
                                          r12d, 0C000000Dh
                                  mov
fffff800`01cba7aa 4489a424b0000000 mov
                                           dword ptr [rsp+0B0h], r12d
fffff800`01cba7b2 33db
                                  xor
                                          ebx, ebx
fffff800`01cba7b4 48895c2448
                                          qword ptr [rsp+48h], rbx
                                  mov
fffff800`01cba7b9 33c0
                                          eax, eax
                                  xor
fffff800`01cba7bb 4889442450
                                          qword ptr [rsp+50h], rax
                                  mov
fffff800`01cba7c0 4889442458
                                          qword ptr [rsp+58h], rax
                                  mov
fffff800`01cba7c5 448d6b01
                                          r13d, [rbx+1]
                                  1ea
fffff800`01cba7c9 458afd
                                          r15b, r13b
                                  mov
fffff800`01cba7cc 4488ac24a8000000 mov
                                           byte ptr [rsp+0A8h], r13b
fffff800`01cba7d4 440f20c7
                                          rdi, cr8
                                  mov
fffff800`01cba7d8 450f22c5
                                          cr8, r13
fffff800`01cba7dc f00fba351b6adcff00 lock btr dword ptr [nt!CallbackUnregisterLock
(fffff800`01a81200)], 0
ffffff800`01cba7e5 720c
                                          nt!CmUnRegisterCallback+0x63 (fffff800`01cba7f3)
                                  jb
//省略中间无关代码
nt!CmUnRegisterCallback+0xc6:
fffff800`01cba856 4533c0
                                          r8d, r8d
                                  xor
fffff800`01cba859 488d542420
                                  lea
                                          rdx, [rsp+20h]
fffff800`01cba85e 488d0d6b69dcff lea
                                          rcx, [nt!CallbackListHead (fffff800`01a811d0)]
fffff800`01cba865 e8a261e5ff
                                  call
                                          nt!CmListGetNextElement (fffff800`01b10a0c)
fffff800`01cba86a 488bf8
                                          rdi, rax
                                  mov
fffff800`01cba86d 4889442428
                                          qword ptr [rsp+28h], rax
                                  mov
fffff800`01cba872 483bc3
                                  cmp
                                          rax, rbx
```

搜索的过程跟寻找进程、线程、映像的数组类似,根据 lea REG,XXX 来定位。不过为了更加精准,这次我采用两次 lea REG,XXX 来定位。:

```
ULONG64 FindCmpCallbackAfterXP()
{
    ULONG64
                      uiAddress=0;
    PUCHAR
                      pCheckArea=NULL, i=0, j=0, StartAddress=0, EndAddress=0;
    ULONG64
                      dwCheckAddr=0;
    UNICODE STRING unstrFunc;
    UCHAR b1=0,b2=0,b3=0;
    ULONG templong=0,QuadPart=0xfffff800;
    RtlInitUnicodeString(&unstrFunc, L"CmUnRegisterCallback");
    pCheckArea = (UCHAR*)MmGetSystemRoutineAddress (&unstrFunc);
    if (!pCheckArea)
         KdPrint(("MmGetSystemRoutineAddress failed."));
         return 0;
    StartAddress = (PUCHAR)pCheckArea;
    EndAddress = (PUCHAR)pCheckArea + PAGE_SIZE;
    for(i=StartAddress;i<EndAddress;i++)</pre>
         if( MmIsAddressValid(i) && MmIsAddressValid(i+1) && MmIsAddressValid(i+2) )
         {
             b1=*i;
             b2=*(i+1);
             b3=*(i+2);
             if( b1==0x48 \&\& b2==0x8d \&\& b3==0x0d ) //488d0d(lea rcx,)
                  j=i-5;
                  b1=*j;
                  b2=*(j+1);
                  b3=*(j+2);
                  if( b1==0x48 && b2==0x8d && b3==0x54 ) //488d54(lea rdx,)
                       memcpy(&templong,i+3,4);
                       uiAddress = MakeLong64ByLong32(templong) + (ULONGLONG)i + 7;
                       return uiAddress;
                  }
             }
         }
```

```
return 0;
}
```

定位完毕之后,就是枚举链表了。注册表回调是一个"结构体链表",类似于 EPROCESS,它的定义如下:

我们只关心两个值,一个是 Cookie,一个是 Function。前者可以理解成注册表回调的"句柄"(用 CmUnRegisterCallback 注销回调传入的就是这个 Cookie),后者是回调函数的地址。代码如下:

```
ULONG CountCmpCallbackAfterXP(ULONG64* pPspLINotifyRoutine)
    ULONG
                           sum = 0;
    ULONG64
                           dwNotifyItemAddr;
    ULONG64*
                      pNotifyFun;
    ULONG64*
                      baseNotifyAddr;
    ULONG64
                           dwNotifyFun;
                      cmpCookie;
    LARGE_INTEGER
    PLIST ENTRY
                           notifyList;
    PCM_NOTIFY_ENTRY notify;
    dwNotifyItemAddr = *pPspLINotifyRoutine;
    notifyList = (LIST_ENTRY *)dwNotifyItemAddr;
    do
         notify = (CM_NOTIFY_ENTRY *)notifyList;
         if (MmIsAddressValid(notify))
                  (MmIsAddressValid((PVOID) (notify->Function)) && notify->Function
{
                  DbgPrint("[CmCallback]Function=%p\tCookie=%p",
(PVOID) (notify->Function), (PVOID) (notify->Cookie. QuadPart));
                  //notify->Function=(ULONG64)MyRegistryCallback;
                  sum ++:
             }
```

```
notifyList = notifyList->Flink;
} while ( notifyList != ((LIST_ENTRY*)(*pPspLINotifyRoutine)) );
return sum;
}
```

执行效果类似于:

```
11.94157314 [MY FUNCTION]: FFFFF88003C23008
11.94157600 [MY COOKIE]: 01CEF9B0165BD342
11.94158268 CmCallbackListHead: FFFFF800040C91D0
11.94158459 [CmCallback]Function=FFFFF88005ED2CB8 Cookie=01CEF9B0165BD341
11.94158554 [CmCallback]Function=FFFFF88003C23008 Cookie=01CEF9B0165BD342
```

不过需要注意的是,干净的 WIN7X64 系统是没有注册表回调的。为了体现枚举效果,可以在测试驱动前运行 WIN64AST。对付注册表回调有三种方法(老三套): 1.直接使用 CmUnRegisterCallback 把回调注销; 2.把链表中记录的回调地址修改为自定义的空函数的回调地址; 3.直接在目标回调地址上写一个 RET,使其不执行任何代码就返回。第三种方法没有针对性,可以用于对付任何回调函数。DisableFunctionWithReturnValue 用来对付有返回值的回调函数,DisableFunctionWithoutReturnValue 用于对付无返回值的回调函数。

```
KIRQL WPOFFx64()
    KIRQL irql=KeRaiseIrqlToDpcLevel();
    UINT64 cr0= readcr0();
    __writecr0(cr0);
    _disable();
    return irql;
}
void WPONx64(KIRQL irql)
    UINT64 cr0= readcr0();
    cr0 = 0x10000;
    enable();
    __writecr0(cr0);
    KeLowerIrql(irql);
VOID DisableFunctionWithReturnValue(PVOID Address)
    KIRQL irql;
    CHAR patchCode[] = "\x33\xC0\xC3"; //xor eax,eax + ret
    if(MmIsAddressValid(Address))
        irql=WPOFFx64();
        memcpy(Address,patchCode,3);
```

```
WPONx64(irql);
}

VOID DisableFunctionWithoutReturnValue(PVOID Address)

{
    KIRQL irql;
    if(MmIsAddressValid(Address))
    {
        irql=WPOFFx64();
        RtIFillMemory(Address,1,0xC3);
        WPONx64(irql);
    }
}
```

等大家深入学习之后,会发现以上写两个函数纯属多此一举,在此先卖个关子不解释。