

Homework 1 (Due Jan 18, 2023)

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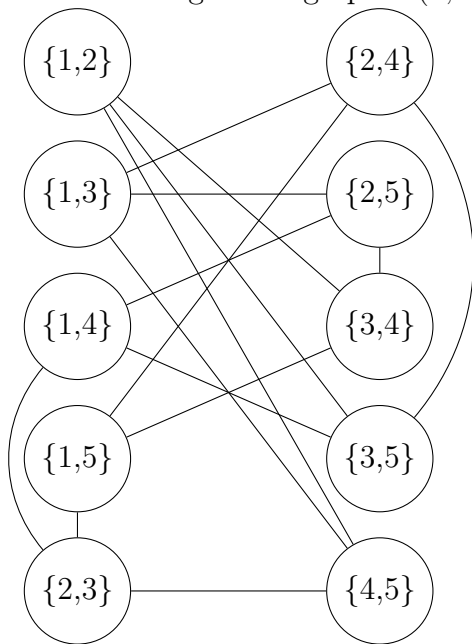
MATH 575 - Discrete Mathematics II - Spring 2023

April 24, 2023

Justify all of your answers completely.

- For positive integers n and k , consider the graph $G(n, k)$ which is defined as follows: the vertex set of $G(n, k)$ is the set of subsets of $[n]$ of size k , and two vertices are connected by an edge in $G(n, k)$ if and only if the corresponding subsets are disjoint.¹

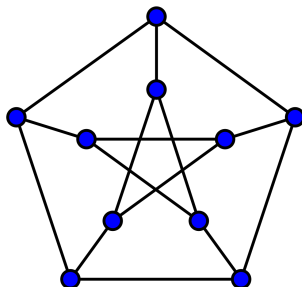
(a) Give a drawing of the graph $G(5, 2)$.



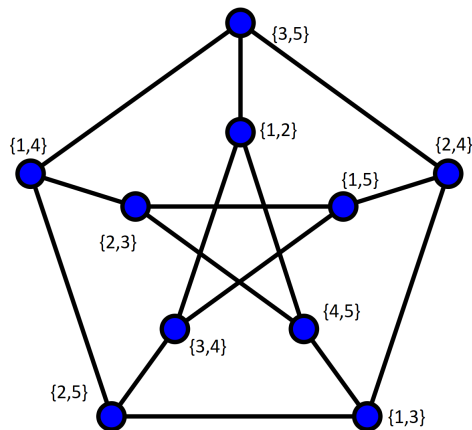
¹E.g., $V(G(4, 2)) = \{\{1, 2\}, \{1, 3\}, \{1, 4\}, \{2, 3\}, \{2, 4\}, \{3, 4\}\}$ and

$$E(G(4, 2)) = \{\{\{1, 2\}, \{3, 4\}\}, \quad \{\{1, 3\}, \{2, 4\}\}, \quad \{\{1, 4\}, \{2, 3\}\}\}$$

- (b) Let G be the graph drawn below. Show that $G(5,2)$ is isomorphic to G by relabelling the vertices of G in the drawing below.

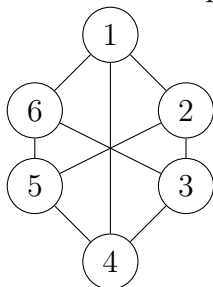


Here is the labeled graph

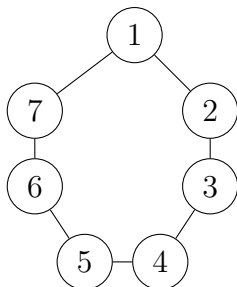


2. A graph is called **k -regular** if every vertex has degree k .

- (a) Draw an example of a 3-regular graph on 6 vertices.



- (b) Draw two nonisomorphic 2-regular graphs on 7 vertices.



- (c) Prove that if k is odd, then there does not exist a k -regular graph with an odd number of vertices.

Proof. Assume we have a graph, G , that is k -regular, where k is odd. Assume towards contradiction that G has an odd number of vertices, namely n . By Handshaking lemma, summing up the degrees of the vertices will be an even number, but in G , it will be nk . This is an odd number multiplied by an odd

BOOM, A CONTRADICTION!!!

number, which is not even. ■

3. Prove that every graph G must contain a pair of vertices with the same degree.

Proof. Let v_1, v_2, \dots, v_n be the vertices of G , and d_i be the degree of v_i . There are n different possible numbers d_i can be, namely $0, \dots, n-1$. If a vertex has a degree of $n-1$, then it is neighbors to every other vertex, making another vertex having a degree of 0 impossible. So a graph with every degree is impossible, meaning a graph can at most have $n-1$ different d_i 's. So by PHP, matching n vertices to $n-1$ degrees must result in 2 vertices having the same degree. ■

4. Let G be a graph with $V(G) = \{v_1, \dots, v_n\}$. Recall that the *adjacency matrix* of G is the $n \times n$ matrix A such that $A_{ij} = 1$ if $v_i v_j \in E(G)$ and $A_{ij} = 0$ otherwise. Use induction to prove that for all integers $k \geq 1$, the (i, j) -entry of A^k is equal to the number of v_i, v_j -walks of length k in G .

Proof. Let us induct on the degree of adjacency matrix A .

Base case: $k = 1$

Obvious by the definition of adjacency matrix.

Induction step: Assume that A is an adjacency matrix of G and that the (i, j) -entry of A^k is equal to the number of v_i, v_j -walks of length k in G . Let us right multiply A to get $A^k \cdot A = A^{k+1}$. Looking at the (i, j) -entry of A^{k+1} , we get it from the dot product of the i th row of A^k and the j th column of A . Each element in the i th row of A^k is the number of walks from vertex i to each intermediate vertex q , and those entries will be multiplied by 0 if that intermediate vertex q is not adjacent to j . This means those walks of length k starting at vertex i that cannot reach j by adding one more edge to the end are not added. The walks of length k that do reach j after adding an edge to the end of the walk are added, since the number they are being multiplied by is 1 in the j th column of A . So we are left with the number of walks from vertex i to an intermediate vertex q (walk of length k) to vertex j , which gives a walk of length $k+1$. ■

5. Let G be an n -vertex graph with degree sequence (d_1, d_2, \dots, d_n) .

- (a) What is the degree sequence of \overline{G} ?
 $(n - 1 - d_1, n - 1 - d_2, \dots, n - 1 - d_n)$ is the sequence.
- (b) A graph G is called *self-complementary* if it is isomorphic to its complement. Prove that if G is self-complementary, then either n or $n - 1$ is divisible by 4.

Proof. Assume G is self-complementary. This means that $|E(G)|$ is half of the total number of edges for the complete graph with the same number of vertices, since $G \cup \overline{G} = K_n$. So $|E(K_n)| = \frac{n(n-1)}{2} \implies |E(G)| = \frac{n(n-1)}{4}$. And since $|E(G)|$ is a whole number, 4 must divide n or $n - 1$. ■

- (c) Show that for all n divisible by 4, there exists a self-complementary graph on n vertices.²

Proof. The rule set for enumerated vertices: $v_i v_j$ is an edge in the graph

- i. if $v_i \equiv 1 \pmod{4} \wedge (v_j \equiv 1 \pmod{4} \vee v_j \equiv 2 \pmod{4})$
- ii. if $v_i \equiv 2 \pmod{4} \wedge (v_j \equiv 1 \pmod{4} \vee v_j \equiv 3 \pmod{4})$
- iii. if $v_i \equiv 3 \pmod{4} \wedge (v_j \equiv 2 \pmod{4} \vee v_j \equiv 0 \pmod{4})$
- iv. if $v_i \equiv 0 \pmod{4} \wedge (v_j \equiv 3 \pmod{4} \vee v_j \equiv 0 \pmod{4})$

gives a self complementary graph on $4k$ vertices.

We will show the above statement by inducting on $n = 4k$.

Base case: **$k = 1$**

The P_4 graph is self-complementary.



$k = 2$

Layering P_4 graphs with the rule set results in self-complementary graphs.



Induction Step: Assume we have a simple graph, G , such that there are $4k$ vertices $\{v_1, \dots, v_{4k}\}$, with the rules stated in the beginning, where $k \geq 2$. (This is the pattern shown in base case $k = 2$)

Let us add 4 vertices to the graph and apply the rule to those vertices. We know from the base case $k = 2$ that the subgraph of this new "layer" and any other layer will be self-complementary. Since the vertices in the new layer are self complementary with every other layer, and the other layers were already self-complementary in total, this new graph with $4(k + 1)$ vertices is also self complementary. ■

²Hint: generalize the structure of the path P_4 .

- (d) Show that for all n such that $n-1$ is divisible by 4, there exists a self-complementary graph on n vertices.³

In part (c), we constructed a graph with $4k$ vertices that was self-complementary. Let us add a single vertex to our graph. It is now not self-complementary. To make it self complementary, we must connect the vertex to the "outer" vertices in each "row". This makes the graph self-complementary, since the vertices our new vertex is adjacent to in the complement graph are now the "outer" vertices (i.e. 2 and 3 become the outer vertices), and we still have half the total possible edges. So now we have a self-complementary construction for every graph with n vertices where $4|(n-1)$.

³Hint: add a vertex to a construction in part (c).