CSE 120: Principles of Operating Systems
Lecture 13: Protection

Prof. Joseph Pasquale
University of California, San Diego
March 4, 2019

Protection

- Processes access resources
- Resources are shared, need to be protected
  - from processes without permission
  - from improper access by a process
- What is the right protection model?
- What are the mechanisms?

The Kernel Enforces Protection

- To protect resources, have kernel "own" them
  - kernel can then allow access (temporarily)
- To access a resource, a process must ask for it
  - kernel can test whether access should be given
- Once a process is given access
  - kernel can prevent others for gaining access
  - kernel may/may not be able to take away access
- This assumes the kernel operates correctly

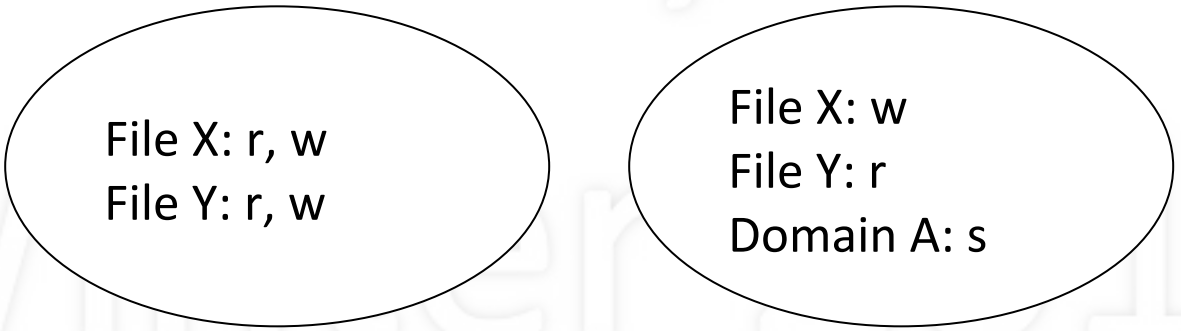
Protecting the Kernel

- The kernel itself must be protected!
- Mechanisms
  - Memory protection
  - Protected mode of operation: kernel vs. user
  - Clock interrupt, so kernel eventually gets control
- Notice, mechanisms are hardware supported
- Protected kernel can protect other resources

Goals Supported by Kernel

- Allow range of permissions
- Allow user to set/get them
- Be fast/simple for common case
- Support user expressing complex permissions

A Formal Model of Protection



- Protection: how to limit access to a resource
- Resource: object that requires protection
- Domain: set of (resource, permission) pairs
- Process: accesses resources within domain

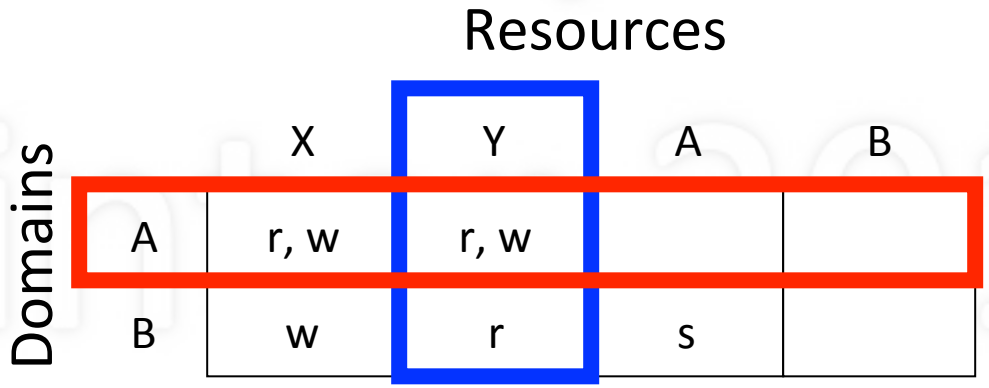
Protection Matrix

Resources

Domains	X	Y	A	B
A	r, w	r, w		
B	w	r	s	

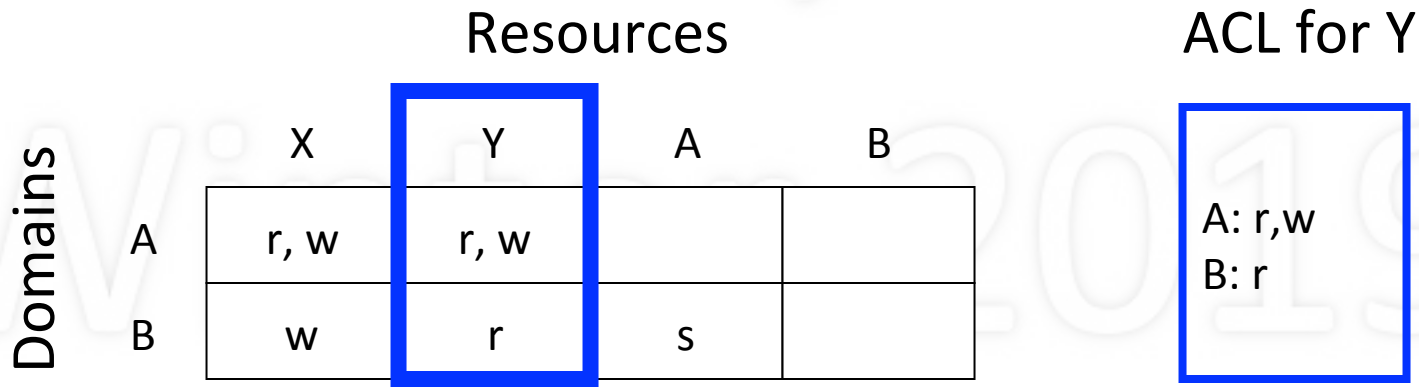
- Can describe all domains as a matrix
  - Rows are domains
  - Columns are resources
  - Matrix entry [d, r] contains permissions/rights

Efficient Representations



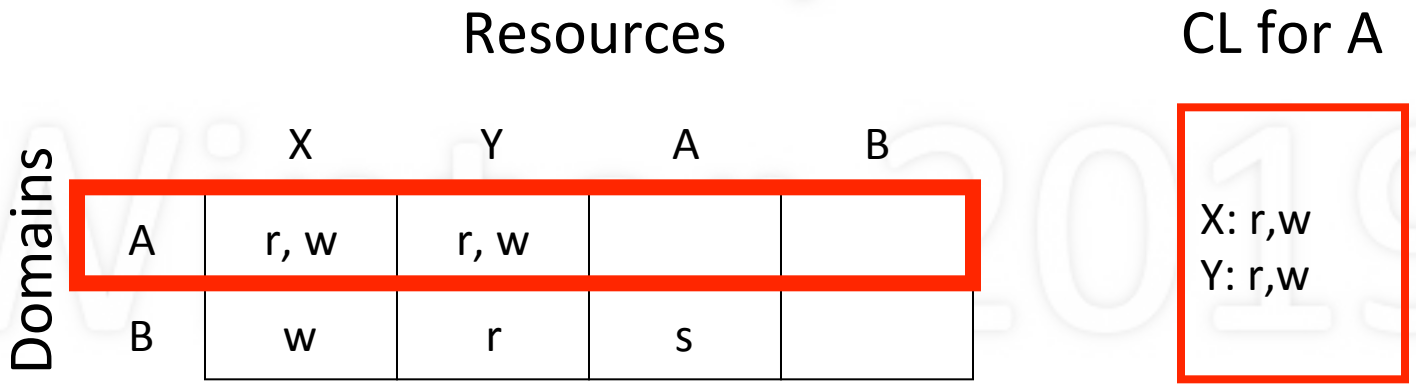
- Access Control Lists
  - For each resource, list (domain, permissions) pairs
- Capability Lists
  - For each domain, list (resource, permissions) pairs

Access Control Lists



- ACL is associated with resource
- Like a registry: if name is on list, ok to access
- Can be inefficient: must lookup on each access
- Revocation is easy; just remove from list

Capability Lists



- Capability list associated with each domain
- Like key/ticket: if you have it, you get access
- Efficient: on access, just produce capability
- Hard to revoke

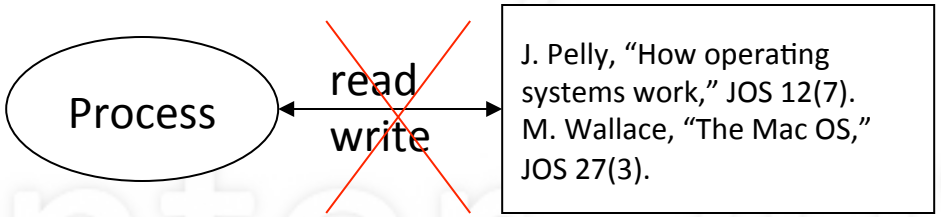
UNIX Protection

- Associated with each file is set of permissions
  - Permission bits r/w/x for owner, group, world
  - Limited form of access control list
- Protection domain: UID (user account ID) + …
  - A process is always in some domain
- When process opens file, check permission
- If ok, provide process with a capability
  - Future operations then carried out efficiently

Extending Protection in UNIX

- For common case, r/w/x for o/g/w adequate
- For special cases, can extend via user program
- SETUID mechanism: causes domain switch
- If executable file has SETUID bit set
  - Process runs in domain of owner (of executable)
  - Therefore, it runs with all the rights of the owner

Example Use of SETUID bit

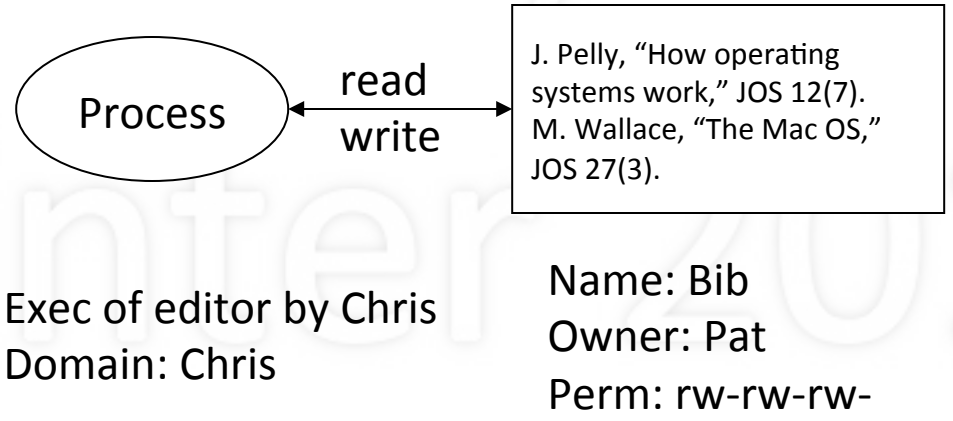


Exec of editor by Chris
Domain: Chris

Name: Bib
Owner: Pat
Perm: rw-------

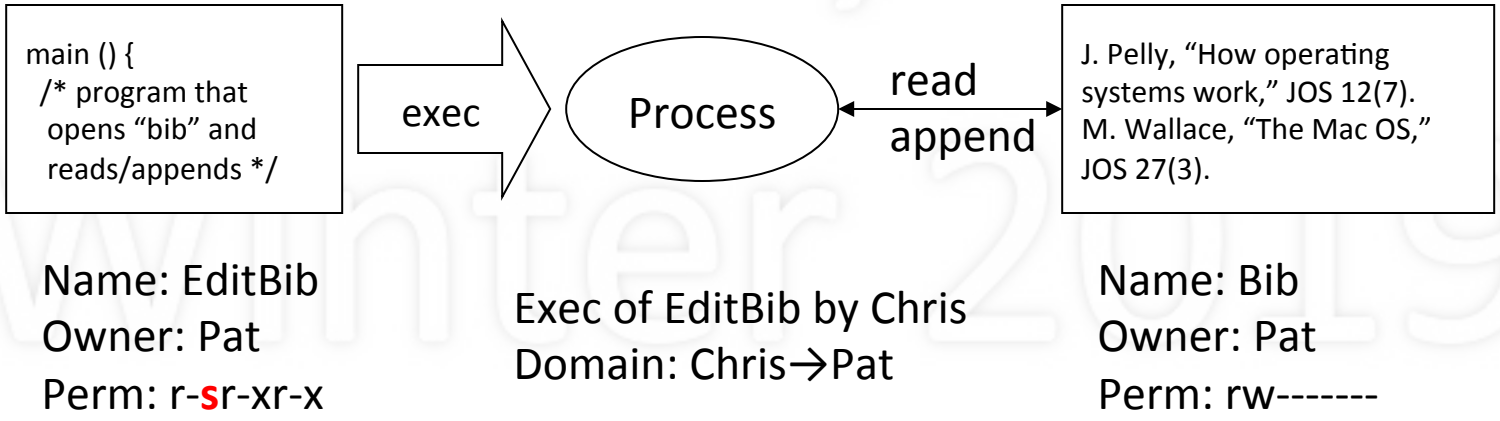
- Pat has a file "Bib" of bibliographic references
- Chris wants to read *and add entries*
- But, Chris lacks permissions (only Pat can r/w)
- Pat wishes to allow append access, but how?

Example Use of SETUID bit



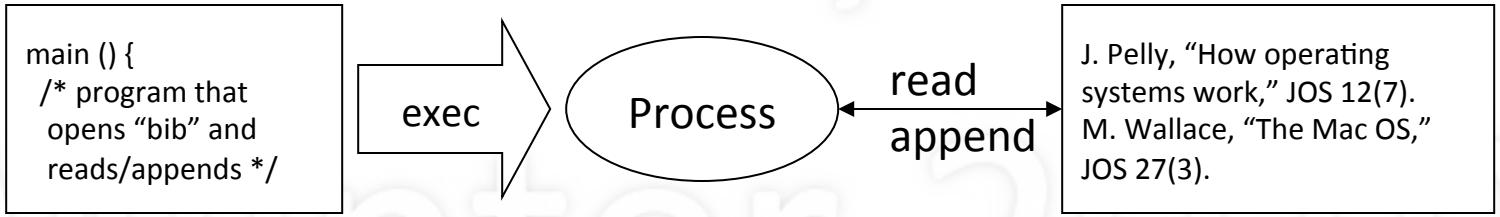
- Pat can set permissions so Chris can r/w
- But this gives Chris too much power
  - Chris can modify the file arbitrarily
  - Pat would like to give Chris only append access

Example Use of SETUID bit



- Pat provides program: only reads/appends
- Sets permissions
  - of program: execute (for Chris), and SETUID on
  - of Bib file: read/write only for Pat, not Chris

Example Use of SETUID bit



Name: EditBib
Owner: Pat
Perm: r-**s**r-xr-x

Exec of EditBib by Chris
Domain: Chris→Pat

Name: Bib
Owner: Pat
Perm: rw-------

- When Chris executes EditBib, runs as Pat
  - SETUID causes a domain switch to Pat's domain
- Since Pat has r/w access, file can be modified
- Limited to what program does: read/append

Summary

- Protection
- Formal Model
- Access Control Lists
- Capability Lists

Textbook

- Chapter 17 (on Protection)
  - Lecture-related: 17.1-17.2, 17.4-17.7, 17.13
  - Recommended: 17.3, 17.8-17.12

CSE 120, UCSD Winter 2019

Review & Research

- What is meant by "protection" in operating systems?*
- What are the reasons require protection?*
- In what way does the kernel enforce protection?**
- How is the kernel itself protected?**
- Why is hardware support needed to protect the kernel?**

R&R

- What are the goals for the kernel to support protection?*
- In the formal model of protection, what is the definition of each of the following: protection, resource, domain, process?

R&R

- What is a protection matrix?*
- What is represented by the matrix rows?
- What is represented by the matrix columns?
- What is contained in a matrix cell?
- What is the “s” right/permission?*
- Why is the protection matrix inefficient to represent as full matrix?**

R&R

- What is an access control list (ACL)?
- In what way is an ACL like a registry?*
- In what way is inefficiency a disadvantage of ACLs?*
- What is revocation, and why is it an advantage of ACLs?*

R&R

- What is a capability list?
- In what way is a capability like a key/ticket?*
- In what way is efficiency an advantage of a capability?*
- Why is revocation difficult when using capability lists?*

R&R

- What is the design for protection in UNIX?
- What is a protection domain in UNIX?*
- How is the protection model integrated with the UNIX file system?**
- How does UNIX use ACLs and capabilities?**
- What is the SETUID mechanism in UNIX?**

R&R

- If an executable file has the SETUID bit set, how does this affect a process that executes that file?**
- On slide 13, what happens if Chris tries to read the Bib file, and why?*
- What would happen if Pat sets the permission bits to rw-rw-rw, how would it affect your previous answer?*

R&R

- On slide 15, what is the purpose of the EditBib program?*
- Why is it that Pat would be the owner of this file (rather than Chris)?**
- Why is it that the SETUID bit must be set for this file (and not for the Bib file)?**
- Why would it not be good to change the permissions of Bib to rw-rw-rw-?*

R&R

- When Chris's process executes the EditBib file, what does it mean for the process to run in Pat's domain?*
- By running in Pat's domain, doesn't this mean that Chris can now impersonate Pat: why or why not?**