

#### **Product Specification**

#### Z8400/Z84C00 NMOS/CMOS Z80° CPU Central Processing Unit

#### **FEATURES**

The extensive instruction set contains 158 instructions, including the 8080A instruction set as a subset.

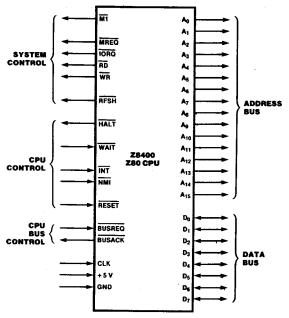
- NMOS version for low cost high performance solutions, CMOS version for high performance low power designs.
- NMOS Z0840004 4 MHz, Z0840006 6.17 MHz, Z0840008 - 8 MHz.
- CMOS Z84C0006 DC to 6.17 MHz, Z84C008 DC to 8 MHz, Z84C0010 - DC to 10 MHz, Z84C0020 - DC -20 MHz
- 6 MHz version can be operated at 6.144 MHz clock.

- The Z80 microprocessors and associated family of peripherals can be linked by a vectored interrupt system. This system can be daisy-chained to allow implementation of a priority interrupt scheme.
- Duplicate set of both general-purpose and flag registers.
- Two sixteen-bit index registers.
- Three modes of maskable interrupts:

  Mode 0—8080A similar;

  Mode 1—Non-Z80 environment, location 38H;

  Mode 2—Z80 family peripherals, vectored interrupts.
- On-chip dynamic memory refresh counter.





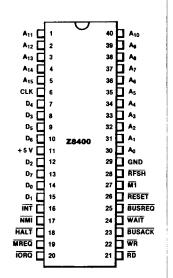
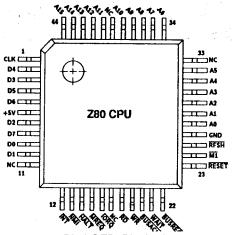
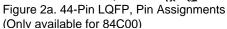


Figure 2. 40-pin Dual-In-Line (DIP), Pin Assignments





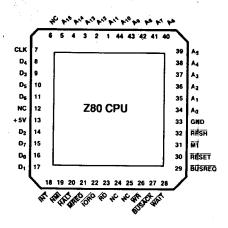


Figure 2b. 44-Pin Chip Carrier Pin Assignments

#### **GENERAL DESCRIPTION**

The CPUs are fourth-generation enhanced microprocessors with exceptional computational power. They offer higher system throughput and more efficient memory utilization than comparable second- and third-generation microprocessors. The internal registers contain 208 bits of read/write memory that are accessible to the programmer. These registers include two sets of six general-purpose registers which may be used individually as either 8-bit registers or as 16-bit register pairs. In addition, there are two sets of accumulator and flag registers. A group of "Exchange" instructions makes either set of main or alternate registers accessible to the programmer. The alternate set allows operation in foreground-background mode or it may be reserved for very fast interrupt response.

The CPU also contains a Stack Pointer, Program Counter, two index registers, a Refresh register (counter), and an Interrupt register. The CPU is easy to incorporate into a system since it requires only a single +5V power source. All output signals are fully decoded and timed to control standard memory or peripheral circuits; the CPU is supported by an extensive family of peripheral controllers. The internal block diagram (Figure 3) shows the primary functions of the processors. Subsequent text provides more detail on the I/O controller family, registers, instruction set, interrupts and daisy chaining, and CPU timing.

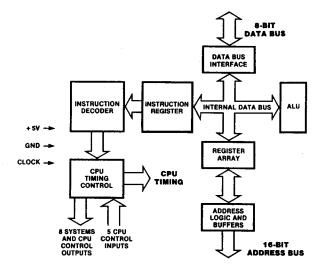


Figure 3. Z80C CPU Block Diagram

#### **CPU REGISTERS**

Figure 4 shows three groups of registers within the CPU. The first group consists of duplicate sets of 8-bit registers: a principal set and an alternate set [designated by ' (prime), e.g., A']. Both sets consist of the Accumulator register, the Flag register, and six general-purpose registers. Transfer of data between these duplicate sets of registers is accomplished by use of "Exchange" instructions. The result is faster response to interrupts and easy, efficient implementation of such versatile programming techniques

as background-foreground data processing. The second set of registers consists of six registers with assigned functions. These are the I (Interrupt register), the R (Refresh register), the IX and IY (Index registers), the SP (Stack Pointer), and the PC (Program Counter). The third group consists of two interrupt status flip-flops, plus an additional pair of flip-flops which assists in identifying the interrupt mode at any particular time. Table 1 provides further information on these registers.

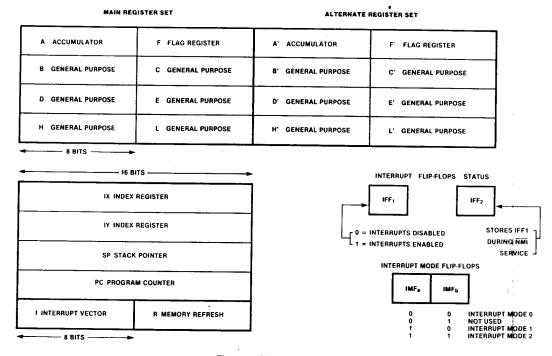


Figure 4. CPU Registers

#### **INTERRUPTS: GENERAL OPERATION**

The CPU accepts two interrupt input signals:  $\overline{\text{NMI}}$  and  $\overline{\text{INT}}$ . The  $\overline{\text{NMI}}$  is a non-maskable interrupt and has the highest priority.  $\overline{\text{INT}}$  is a lower priority interrupt and it requires that interrupts be enabled in software in order to operate.  $\overline{\text{INT}}$  can be connected to multiple peripheral devices in a wired-OR configuration.

The Z80 has a single response mode for interrupt service on the non-maskable interrupt. The maskable interrupt, INT, has three programmable response modes available. These are:

- Mode 0 similar to the 8080 microprocessor.
- Mode 1 Peripheral Interrupt service, for use with non-8080/Z80 systems.

Mode 2 - a vectored interrupt scheme, usually daisychained, for use with the Z80 Family and compatible peripheral devices.

The CPU services interrupts by sampling the  $\overline{\text{NMI}}$  and  $\overline{\text{INT}}$  signals at the rising edge of the last clock of an instruction. Further interrupt service processing depends upon the type of interrupt that was detected. Details on interrupt responses are shown in the CPU Timing Section.

Non-Maskable Interrupt (NMI). The nonmaskable interrupt cannot be disabled by program control and therefore will be accepted at all times by the CPU. NMI is usually reserved for servicing only the highest priority type interrupts, such as that for orderly shutdown after power

Table 1. Z80C CPU Registers

	Register	Size (Bits)	Remarks
A, A'	Accumulator	8	Stores an operand or the results of an operation.
F, F'	Flags	8	See Instruction Set.
B, B'	General Purpose	8	Can be used separately or as a 16-bit register with C.
C, C'	General Purpose	8	Can be used separately or as a 16-bit register with C.
D, D'	General Purpose	8	Can be used separately or as a 16-bit register with E.
E, E'	General Purpose	8	Can be used separately or as a 16-bit register with E.
H, H'	General Purpose	8	Can be used separately or as a 16-bit register with L.
L, L'	General Purpose	8	Can be used separately or as a 16-bit register with L.
			Note: The (B,C), (D,E), and (H,L) sets are combined as follows:  B — High byte C — Low byte  D — High byte E — Low byte  H — High byte L — Low byte
	Interrupt Register	8	Stores upper eight bits of memory address for vectored interrupt processing.
R	Refresh Register	8	Provides user-transparent dynamic memory refresh. Automatically incremented and placed on the address bus during each instruction fetch cycle.
IX	Index Register	16	Used for indexed addressing.
IY .	Index Register	16	Used for indexed addressing
SP	Stack Pointer	16	Holds address of the top of the stack. See Push or Pop in instruction set.
PC	Program Counter	16	Holds address of next instruction.
IFF <sub>1</sub> -IFF <sub>2</sub>	Interrupt Enable	Flip-Flops	Set or reset to indicate interrupt status (see Figure 4).
IMFa-IMFb	Interrupt Mode	Flip-Flops	Reflect Interrupt mode (see Figure 4).

failure has been detected. After recognition of the \$\overline{NMI}\$ signal (providing \$\overline{BUSREQ}\$ is not active), the CPU jumps to restart location 0066H. Normally, software starting at this address contains the interrupt service routine.

Maskable Interrupt (INT). Regardless of the interrupt mode set by the user, the CPU response to a maskable interrupt input follows a common timing cycle. After the interrupt has been detected by the CPU (provided that interrupts are enabled and BUSREQ is not active) a special interrupt processing cycle begins. This is a special fetch (M1) cycle in which IORQ becomes active rather than MREQ, as in a normal M1 cycle. In addition, this special M1 cycle is automatically extended by two WAIT states, to allow for the time required to acknowledge the interrupt request.

**Mode 0 Interrupt Operation.** This mode is similar to the 8080 microprocessor interrupt service procedures. The interrupting device places an instruction on the data bus. This is normally a Restart instruction, which will initiate a call

to the selected one of eight restart locations in page zero of memory. Unlike the 8080, the Z80 CPU responds to the Call instruction with only one interrupt acknowledge cycle followed by two memory read cycles.

**Mode 1 Interrupt Operation.** Mode 1 operation is very similar to that for the NMI. The principal difference is that the Mode 1 interrupt has only one restart location, 0038H.

Mode 2 Interrupt Operation. This interrupt mode has been designed to most effectively utilize the capabilities of the Z80 microprocessor and its associated peripheral family. The interrupting peripheral device selects the starting address of the interrupt service routine. It does this by placing an 8-bit vector on the data bus during the interrupt acknowledge cycle. The CPU forms a pointer using this byte as the lower 8 bits and the contents of the I register as the upper 8 bits. This points to an entry in a table of addresses for interrupt service routines. The CPU then jumps to the routine at that

address. This flexibility in selecting the interrupt service routine address allows the peripheral device to use several different types of service routines. These routines may be located at any available location in memory. Since the interrupting device supplies the low-order byte of the 2-byte vector, bit 0 (A<sub>a</sub>) must be a zero.

Interrupt Enable/Disable Operation. Two flip-flops, IFF1 and IFF2, referred to in the register description, are used to signal the CPU interrupt status. Operation of the two flip-flops is described in Table 2. For more details, refer to the Z80 CPU Technical Manual (03-0029-01) and Z80 Assembly Language Programming Manual (03-0002-01).

Table 2. State of Flip-Flops

Action	IFF <sub>1</sub>	IFF <sub>2</sub>	Comments
CPU Reset	0	0	Maskable interrupt
DI instruction execution	0	0	Maskable interrupt INT disabled
El instruction execution	1	1	Maskable interrupt
LD A,I instruction execution	•	•	IFF <sub>2</sub> → Parity flag
LD A,R instruction execution	•	•	IFF <sub>2</sub> → Parity flag
Accept NMI	0	•	Maskable interrupt
RETN instruction execution	IFF <sub>2</sub>	•	IFF <sub>2</sub> → IFF <sub>1</sub> at completion of an NMI service routine.

#### **INSTRUCTION SET**

The microprocessor has one of the most powerful and versatile instruction sets available in any 8-bit microprocessor. It includes such unique operations as a block move for fast, efficient data transfers within memory, or between memory and I/O. It also allows operations on any bit in any location in memory.

The following is a summary of the instruction set which shows the assembly language mnemonic, the operation, the flag status, and gives comments on each instruction. For an explanation of flag notations and symbols for mnemonic tables, see the Symbolic Notations section which follows these tables. The Z80 CPU Technical Manual (03-0029-01). the Programmer's Reference Guide (03-0012-03), and Assembly Language Programming Manual (03-0002-01) contain significantly more details for programming use.

The instructions are divided into the following categories: ☐ 8-bit loads □ 16-bit loads ☐ Exchanges, block transfers, and searches □ 8-bit arithmetic and logic operations ☐ General-purpose arithmetic and CPU control □ 16-bit arithmetic operations □ Rotates and shifts

- ☐ Bit set, reset, and test operations
- □ Jumps
- □ Calls, returns, and restarts
- □ Input and output operations

A variety of addressing modes are implemented to permit efficient and fast data transfer between various registers, memory locations, and input/output devices. These addressing modes include:

- □ Immediate
- □ Immediate extended
- □ Modified page zero
- □ Relative
- □ Extended
- □ Indexed
- □ Register
- □ Register indirect
- □ Implied
- □ Bit

#### **PIN DESCRIPTIONS**

**A<sub>0</sub>-A<sub>15</sub>.** Address Bus (output, active High, 3-state). A<sub>0</sub>-A<sub>15</sub> form a 16-bit address bus. The Address Bus provides the address for memory data bus exchanges (up to 64K bytes) and for I/O device exchanges.

BUSACK. Bus Acknowledge (output, active Low). Bus Acknowledge indicates to the requesting device that the CPU address bus, data bus, and control signals MREQ, IORQ, RD, and WR have entered their high-impedance states. The external circuitry can now control these lines.

BUSREQ. Bus Request (input, active Low). Bus Request has a higher priority than NMI and is always recognized at the end of the current machine cycle. BUSREQ forces the CPU address bus, data bus, and control signals MREQ, IORQ, RD, and WR to go to a high-impedance state so that other devices can control these lines. BUSREQ is normally wired-OR and requires an external pullup for these applications. Extended BUSREQ periods due to extensive DMA operations can prevent the CPU from properly refreshing dynamic RAMs.

**D<sub>0</sub>-D<sub>7</sub>.** Data Bus (input/output, active High, 3-state). D<sub>0</sub>-D<sub>7</sub> constitute an 8-bit bidirectional data bus, used for data exchanges with memory and I/O.

HALT. Halt State (output, active Low). HALT indicates that the CPU has executed a Halt instruction and is awaiting either a nonmaskable or a maskable interrupt (with the mask enabled) before operation can resume. While halted, the CPU executes NOPs to maintain memory refresh.

**INT.** Interrupt Request (input, active Low). Interrupt Request is generated by I/O devices. The CPU honors a request at the end of the current instruction if the internal software-controlled interrupt enable flip-flop (IFF) is enabled. INT is normally wired-OR and requires an external pullup for these applications.

IORQ. Input/Output Request (output, active Low, 3-state). IORQ indicates that the lower half of the address bus holds a valid I/O address for an I/O read or write operation. IORQ is also generated concurrently with M1 during an interrupt acknowledge cycle to indicate that an interrupt response vector can be placed on the data bus.

M1. Machine Cycle One (output, active Low). M1, together with MREQ, indicates that the current machine cycle is the opcode fetch cycle of an instruction execution. M1, together with IORQ, indicates an interrupt acknowledge cycle.

MREQ. Memory Request (output, active Low, 3-state). MREQ indicates that the address bus holds a valid address for a memory read or memory write operation.

NMI. Non-Maskable Interrupt (input, negative edgetriggered). NMI has a higher priority than INT. NMI is always recognized at the end of the current instruction, independent of the status of the interrupt enable flip-flop, and automatically forces the CPU to restart at location 0066H.

RD. Read (output, active Low, 3-state). RD indicates that the CPU wants to read data from memory or an I/O device. The addressed I/O device or memory should use this signal to gate data onto the CPU data bus.

**RESET.** Reset (input, active Low). RESET initializes the CPU as follows: it resets the interrupt enable flip-flop, clears the PC and Registers I and R, and sets the interrupt status to Mode 0. During reset time, the address and data bus go to a high-impedance state, and all control output signals go to the inactive state. Note that RESET must be active for a minimum of three full clock cycles before the reset operation is complete.

**RFSH.** Refresh (output, active Low). RFSH, together with MREQ, indicates that the lower seven bits of the system's address bus can be used as a refresh address to the system's dynamic memories.

WAIT. Wait (input, active Low). WAIT indicates to the CPU that the addressed memory or I/O devices are not ready for a data transfer. The CPU continues to enter a Wait state as long as this signal is active. Extended WAIT periods can prevent the CPU from properly refreshing dynamic memory.

WR. Write (output, active Low, 3-state). WR indicates that the CPU data bus holds valid data to be stored at the addressed memory or I/O location.

PS017801-0602

#### **8-BIT LOAD GROUP**

	Symbolic				Fk	ngs					Opcod	•		No. of	No. of M	No. of T		
Mnemonic	Operation	S	Z		H		P/V	N	C	76	543	210	Hex	Bytes	Cycles	States	Com	ments
LD r, r'	r ← r'	•	•			Х	•	•	•	01	r	r'		1	1	4	r, r'	Reg.
LD r, n	r ← n	•	•	Χ	•	Х	•	•	•	00	r	110		2	2	7	000	В
											<b>←</b> n→						001	С
LD r, (HL)	r ← (HL)	•	•	Χ	•	Х	•	•	•	01	r	110		1	2	7	010	D
LD r, (IX + d)	r ← (IX + d)	•	•	Х	•	Χ	•	•	•	11	011	101	DD	3	5	19	011	Ε
										01	r	110					100	н
											<b>~</b> d→						101	L
LD r, (IY + d)	$r \leftarrow (IY + d)$	•	•	Х	•	Х	•	•	•	11	111	101	FD	3	5	19	111	Α
										01	r	110						
											<b>←</b> d→							
LD (HL), r	(HL) ← r	•	•	Х	•	Χ	•	•	•	01	110	ſ		1	2	7		
LD (IX + d), r	(IX+d) <del>←</del> r	•	•	X	•	Χ	•	•	•	11	011	101	DD	3	5	19		
										01	110	r						
											<b>←</b> d→							
LD (IY + d), r	(IY+d) <del></del> r	•	•	Х	•	Х	•	•	•	11	111	101	FD	3	5	19		
		•								01	110	r						
											<b>←</b> d→							
LD (HL), n	(HL) ← n	•	•	Х	•	Х	•	•	•	00	110	110	36	2	3	10		
											←n→							
LD (IX + d), n	(IX + d) ← n	•	•	Х	•	Х	•	•	•	11	011	101	DD	4	5	19		
										00	110	110	36					
											←d →							
											←n→							

### 8-BIT LOAD GROUP (Continued)

	Symbolic					ags					Opcod			No. of	No. of M	No. of T	
Mnemonic	Operation	S	Z		Н		PΛ	/ N	С	76	543	210	Hex	Bytes	Cycles	States	Comments
LD (IY + d), n	(IY+d) <del>←</del> n	•	•	Х	•	Х	•	•	•	11	111	101	FD	4	5	19	
										00	110	110	36				
											<b>←</b> d→						
											←n→						
LD A, (BC)	A ← (BC)	•	•	Χ	•	Х	•	•	•	00	001	010	OA	1	2	7	
LD A, (DE)	A ← (DE)	•	•	Χ	•	Х	٠	•	•	00	011	010	1A	1	2	7	
LD A, (nn)	A ← (nn)	•	•	Х	•	Х	•	•	•	00	111	010	3A	3	4	13	
											<b>←</b> n→						
											<b>←</b> n→						
LD (BC), A	(BC) ← A	•	•	Х	•	Х	•	•	•	00	000	010	02	1	2	7	
LD (DE), A	(DE) ← A	•	•	Х	•	Χ	•	•	•	00	010	010	12	1	2	7	
LD (nn), A	(nn) ← A	•	•	Х	•	Х	•	•	•	00	110	010	32	3	4	13	
											<b>←</b> n →						
_											<b>←</b> n→						
LD A, I	A←I	#	<b>‡</b>	Х	0	Х	IFF	0	•	11	101	101	ED	2	2	9	
										01	010	111	57				
LDA, R	A←R	<b>‡</b>	<b>‡</b>	X	0	Х	IFF	0	•	11	101	101	ED	2	2	9	
										01	011	111	5F				
_D I, A	1 <del></del> A	•	•	Х	•	Х	•	•	•	11	101	101	ED	2	2	9	
	_									01	000	111	47				
₋DR, A	R←A	•	•	X	•	Х	•	•	•	11	101	101	ED	2	2	9	
										01	<b>0</b> 01	111	4F				

NOTE: IFF, the content of the interrupt enable flip-flop, (IFF2), is copied into the P/V flag.

#### **16-BIT LOAD GROUP**

Mnemonic	Symbolic Operation	s	z		FI: H	ags	P/V	N	С		Opcod 543		Hex	No. of Bytes	No. of M Cycles	No. of T States	Con	ımentı
LD dd, nn	dd ← nn	•	•	X	•	X	•	•	•	00	dd0	001		3	3	10	dd	Pair
											←n→						00	BC
											<b>←</b> n →						01	DE
LD IX, nn	IX ← nn	•	•	X	•	Х	•	•	•	11	011	101	DD	4	4	14	10	HL
										00	100	001	21				11	SP
											<b>←</b> n→							
. =											←n→							
LD IY, nn	IY ← nn	•	•	Х	•	Х	•	•	•	11	111	101	FD	4	4	14		
										00	100	001	21					
											+n→							
											<b>←</b> n →							
LD HL, (nn)	H ← (nn + 1)	•	•	Х	•	Х	•	•	•	00	101	010	2A	3	5	16		
	L ← (nn)										←n→							
S Date of the state of the stat											←n →							
LD dd, (nn)	dd <sub>H</sub> ← (nn + 1)	•	•	Х	•	X	•	•	•	11		101	ED	4	6	20		
	dd <sub>L</sub> ← (nn)									01	dd1	011						
											+n→							
											<b>←</b> n →							

NOTE:  $(PAIR)_H$ ,  $(PAIR)_L$  refer to high order and low order eight bits of the register pair respectively, e.g.,  $BC_L = C$ ,  $AF_H = A$ .

### 16-BIT LOAD GROUP (Continued)

Mnemonic	Symbolic Operation	s	z		Fla H	gs	P/V	N	С	76	Opcod 543		Hex	No. of Bytes	No. of M Cycles	No. of T States	Comme	118
DIX, (nn)	IX <sub>H</sub> ← (nn + 1)	•	•	×	•	χ.	•	•	•	11	011	101	DD	4	6	20		+
<b>-</b> 17 (1 (1 (1)	IX <sub>I</sub> ← (nn)			^		^				00	101		2A		·			
										••	+n→	• • •						į
											+n→							
D IY, (nn)	IY <sub>H</sub> ← (nn + 1)			x		Х				11	. 111		FD	4	6	20		
J 11, (1.1.)	IY <sub>L</sub> ← (nn)			^		•				00	101		2A	,	·			
	115 (111)									00	+n→	0.0						
											+n→							
O (nn), HL	(nn + 1) ← H			Х		х				00	100	010	22	3	5	16		
٠, ١،١٠), ١،١٠	(nn)+-L	-	•	^	•	^	•		•	00	+n→	010		J	J			
	(111)										+n→							
D (nn), dd	(nn + 1) ← dd <sub>H</sub>			¥		х				11	101	101	ED	4	6	20		
- (rii), uu	(nn) ← dd <sub>L</sub>	•	٠	^	•	^	-	•	•	01	dd0		LD	7	J	20		
	(iii) · uu[									U	+ n →	011						
											+n→							
O (nn), IX	(nn + 1) ← IX <sub>H</sub>			¥		х				11	011	101	DD	4	6	20		i
- (ι ιι η, ι∧	(nn) ← IX <sub>L</sub>	-	•	^	•	^	•	•	•	00	100		22	7	J	20		İ
	(111) - 12(									w	+n→	010	~~					
											+n→							
O (nn), IY	(nn+1) ← IY <sub>H</sub>			¥		X ·				11	111	101	FD	4	6	20		
J (1111), 11	(nn) ← IY <sub>L</sub>	•	•	^	•	^	•	•	•	00	100		22	7	Ū	20		
	(iii) · II[									00	+n→	010	22					
											+n→							1
D SP. HL	SP - HL	_		х		х		_		11	111	001	F9	1	1	6		i
O SP. IX	4SP + IX	-	•	x	•	x	•		•	11	011	101	DD	2	2	10		
J UF, IA	10F 1-IA	•	•	^	•	^	•	•	•	11	111	001	F9	-	-			
D SP, IY	SP + IY			х		х				11	111	101	FD	2	2	10		
5 51,11	G( - 11	-	•	^	•	^	-	-	-	11	111	001	F9	-	•		qq Pa	air
USH qq	(SP - 2) ← qq <sub>L</sub>			¥		х				11	qq0	101		1	3	11	00 B	-
	(SP - 1) ← qq <sub>H</sub>	•	٠	^	•	^	•	•	•	1,1	440			•	•	•••	00 DI	
	(SP → SP - 2																10 H	- 1
USH IX	(SP-2) + IXL			¥		х				11	011	101	DD	2	4	15	11 AF	- 1
00111A	(SP - 1) ← IX <sub>H</sub>	•	•	^	•	^	•	•	•	11	100	101	E5	-	7	.0	^	
	SP→SP-2									- 11	,00	101	LU					-
USHIY	SP-2) ← IY <sub>L</sub>			¥		х				11	111	101	FD	2	4	15		i
OOMII		•	•	^	•	^	-	•	-	11	100	101	E5	_	7			
	$(SP-1) \leftarrow IY_H$ $SP \rightarrow SP-2$									11	100	101	ES					
OP oc	or → or - 2 qq <sub>H</sub> ← (SP + 1)			v	_	х				11	qq0	001		1	3	10		
OP qq		•	•	^	•	^	•	•	•	11	440	001		'	3	.0		
	qqL ← (SP) SP → SP +2																	
OD IV			_	v	_	v	_	_	_	11	011	101	DD	2	4	14		
OP IX	IX <sub>H</sub> + (SP + 1)	•	•	^	•	Х	•	•	•	11	011 100	001	E1	2	4	1-4		
	IX <sub>L</sub> ← (SP)									11	100	w	E1					
	SP → SP +2					v		_			444	404				14		
OP IY	IY <sub>H</sub> ← (SP + 1)	•	•	X	•	X	. •	•	•	11	111	101	FD	2	4	14		
	IY <sub>L</sub> ← (SP)									11	100	001	E1					
	SP → SP + 2																	

NOTE: (PAIR)<sub>H</sub>, (PAIR)<sub>L</sub> refer to high order and low order eight bits of the register pair respectively, e.g., BC<sub>L</sub> = C, AF<sub>H</sub> = A.

### EXCHANGE, BLOCK TRANSFER, BLOCK SEARCH GROUPS

Mnemonic	Symbolic Operation	s	z		FI	ngs		/ N	С	76	Opcoc 543	ie 210	Hex	No. of Bytes	No. of M Cycles	No. of T States	Comments
EX DE, HL	DE ++ HL	•	•	х	•	x	•	•	•	11	101	011	EB	1	1	4	
EX AF, AF'	AF ↔ AF'	•	•	Х	•	Х	•	•	•	00	001	000	08	1	1	4	
EXX	BC ++ BC' DE ++ DE' HL ++ HL'	•	•	X	•	X	•	•	•	11	011	001	D9	1	1	4	Register bank and auxiliary register bank exchange
EX (SP), HL	H ++ (SP + 1) L ++ (SP)	•	•	X	•	X	•	•	•	11	100	011	E3	1	5	19	o o name
EX (SP), IX	IX <sub>H</sub> ++ (SP + 1) IX <sub>1</sub> ++ (SP)	•	•	X	•	X	•	•	•	11 11	011 100	101 011	DD E3	2	6	23	
EX (SP), IY	IYH ++ (SP+1)	•	•	х	•	х	•		•	11	111	101	FD	2	6	23	
	IYL ↔ (SP)						<b>①</b>			11	100	011	E3		_		
LDI	(DE) ← (HL) DE ← DE + 1 HL ← HL + 1 BC ← BC - 1	•	•	X	0	X	•	0	•	11 10	101 100	101 000	ED A0	2	4	16	Load (HL) into (DE), increment the pointers and decrement the byte counter
							<b>②</b>										(BC)
LDIR	(DE) ← (HL)	•	•	Х	0	X	0	0	•	11	101	101	ED	2	5	21	If BC ≠ 0
	DE ← DE + 1 HL ← HL + 1 BC ← BC − 1 Repeat until BC = 0									10	110	000	BO	2	4	16	If BC = 0
							①							÷			
LDD	(DE) ← (HL) DE ← DE – 1 HL ← HL – 1 BC ← BC – 1	•	•	X	0	X		0	•	11 10	101 101	101 000	ED A8	2	4	16	
LDDR	(DE) <del></del> (HL)	_	_	x	0	х	Õ	^	_		404	404		•	_	•	<b>****</b>
LDON	DE + DE - 1 HL + HL - 1 BC + BC - 1 Repeat until BC = 0	•	•					U	•	11	101	101 000	ED B8	2 2	5 4	21 16	If BC ≠ 0 If BC = 0
CPI	A - (HL) HL ← HL + 1 BC ← BC - 1	<b>‡</b>	<b>③</b> ‡	x	*	x	<b>①</b>	1	•	11 10	101 100	101 001	ED A1	2	4	16	

NOTE:

(1) P/V flag is 0 if the result of BC - 1 = 0, otherwise P/V = 1.

(2) P/V flag is 0 only at completion of instruction.

(3) Z flag is 1 if A = HL, otherwise Z = 0.

#### EXCHANGE, BLOCK TRANSFER, BLOCK SEARCH GROUPS (Continued)

	Symbolic	_	_			ıgs			_		Opcod			No. of	No. of M		
Mnemonic	Operation	5	Z		Н		P/V	N	C	76	543	210	Hex	Bytes	Cycles	States	Comments
			3				1			•							
CPIR	A – (HL)	‡	#	X	<b>‡</b>	X	ŧ	1	•	11	101	101	ED	2	5	21	If BC ≠ 0 and A ≠ (HL)
•	HL ← HL + 1 BC ← BC − 1 Repeat until A = (HL) or									10	110	001	B1	2	4	16	If BC = 0 or A = (HL)
	BC = 0		3				①										
CPD	A - (HL) HL ← HL - 1 BC ← BC - 1	*	•	X	*	X	•	1	•	11 10	101 101	101 001	ED A9	2	4	16	
CPDR	A – (HL)	<b>‡</b>	③ •	X	<b>‡</b>	X	0	1	•	11	101	101	ED	2	5	21	If BC ≠ 0 and A ≠ (HL)
	HL ← HL − 1 BC ← BC − 1 Repeat until A = (HL) or BC = 0									10	111	001	В9	2	4	16	If BC = 0 or A = (HL)

NOTE: 

P/V flag is 0 if the result of BC - 1 = 0, otherwise P/V = 1.

P/V flag is 0 only at completion of instruction.

Triangle is 1 if A = (HL), otherwise Z = 0.

#### 8-BIT ARITHMETIC AND LOGICAL GROUP

Mnemonic	Symbolic Operation	s	z		Fle H	gs	P/V	N	С	76	Opcod 543	9 210	Hex	No. of Bytes	No. of M Cycles	No. of T States	Com	ments
ADD A, r	A←A+r	*	<b>‡</b>	Х	<b>‡</b>	X	٧	0	<b>‡</b>	10	000	ſ		1	1	4	r	Reg.
ADD A, n	A ← A+n	#	#	Х	<b>‡</b>	Х	٧	0		11	000	110		2	2	7	000	В
											←n→						001	C
																	010	D
ADD A, (HL)	A + A+(HL)	<b>‡</b>	<b>‡</b>	Х	<b>‡</b>	Х	٧	0	<b>‡</b>	10	000	110		1	2	7	011	E
ADD A, (IX + c	d) A←A + (IX + d)	#		Х	<b>‡</b>	Х	٧	0	<b>‡</b>	11	011	101	DD	3	5	19	100	H
										10	000	110					101	L
											<b>-</b> d→						111	A
ADD A, (IY+c	d) A ← A + (IY + d)	<b>‡</b>	\$	Х	<b>‡</b>	Х	٧	0	<b>‡</b>	11	111	101	FD	3	5	19		
										10	000	110						
											<b>-</b> d→							
ADC A, s	A - A+s+CY	<b>‡</b>	<b>‡</b>	Χ	<b>‡</b>	Х	٧	0	#		001						s is a	ny of r, n
SUB s	A ← A – s	<b>‡</b>	<b>‡</b>	X	<b>‡</b>	Х	٧	1	\$		010						(HL),	(IX+d),
SBC A, s	A - A-s-CY	<b>‡</b>	<b>‡</b>	Χ	<b>‡</b>	Х	٧	1	<b>‡</b>		011						(IY+	d) as
ANDs	A ← A > s	<b>‡</b>	<b>‡</b>	X	1	Х	Ρ	0	0		100						show	n for AC
OR s	A ← A > s	<b>‡</b>	<b>‡</b>	X	0	Х	Ρ	0	0		110						instru	ction. T
XOR s	A - Aes	<b>‡</b>	<b>‡</b>	Х	0	Х	Ρ	0	0		101						indica	ated bits
CP s	A-s	<b>‡</b>	<b>‡</b>	Х	<b>‡</b>	Х	٧	1	<b>‡</b>		111						repla	ce the
																	000	] in the
																	ADD	set abo

PS017801-0602

#### 8-BIT ARITHMETIC AND LOGICAL GROUP (Continued)

	Symbolic					ngs					Орсос			No. of	No. of M	No. of T	
Mnemonic	Operation	8	Z		H		PΛ	N	С	76	543	210	Hex	Bytes	Cycles	States	Comments
INC r	r+r+1	<b>‡</b>	\$	х	‡	Х	٧	0	•	00	r	100		1	1	4	· · · · · · · · · · · · · · · · · · ·
INC (HL)	(HL) ←																
	(HL) + 1	<b>‡</b>	<b>‡</b>	Х	<b>‡</b>	Х	٧	0	•	00	110	100		1	3	11	
INC (IX+d)	(IX + d) ←	<b>‡</b>		X	<b>‡</b>	Х	٧	0	•	11	011	101	DD	3	6	23	
	(IX + d) + 1									00	110	100					
											<b>-</b> -d-	•					
INC (IY+d)	(IY + d) ←	#		X	<b>‡</b>	Х	٧	0	•	11	111	101	FD	3	6	23	
	(IY+d)+1									00	110	100					
											<b>-</b> d-	•					
DEC m	m ← m – 1	#	<b>*</b>	X	<b>‡</b>	X	٧	1	•			101					

NOTE: m is any of r, (HL), (IX+d), (IY+d) as shown for INC. DEC same format and states as INC. Replace 100 with 101 in opcode.

### GENERAL-PURPOSE ARITHMETIC AND CPU CONTROL GROUPS

	Symbolic				FI	ngs				(	Орсос	le		No. of	No. of M	No. of T	
Mnemonic	Operation	\$	Z		H	_	PΛ	/ N	C	76		210	Hex	Bytes	Cycles	States	Comments
DAA	<b>@</b>	<b>‡</b>	<b>‡</b>	X	*	X	Р	•	<b>‡</b>	00	100	111	27	1	1	4	Decimal adjust
CPL	A <del>-</del> A	•	•	· X	1	X	•	1	•	00	101	111	2F	1	1	4	Complement accumulator (one's complement).
NEG	A - 0 - A	<b>‡</b>	\$	X	<b>‡</b>	Х	٧	1	<b>‡</b>	11	101	101	ED	2	2	8	Negate acc.
										01	000	100	44				(two's complement).
CCF	CY ← CY	•	•	X	X	X	•	0	‡	00	111	111	3F	1	. 1	4	Complement carry flag.
SCF	CY ← 1	•	•	X	0	X	•	0	1	00	110	111	37	1	1	4	Set carry flag.
NOP	No operation	•	•	Х	•	Х	•	•	•	00	000	000	00	1	1	4	
HALT	CPU halted	•	•	Х	•	Х	•	•	•	01	110	110	76	1	1	4	
DI ★	IFF ← 0	•	•	Х	•	Х	•	•	•	11	110	011	F3	1	1	4	
El ★	IFF ← 1	•	٠	Х	•	Х	•	•	•	11	111	011	FB	1	1	4	:
IM 0	Set interrupt	•	•	Х	•	Х	•	•	•	11	101	101	ΕD	2	2	8	
	mode 0									01	000	110	46				
IM 1	Set interrupt mode 1	•	•	X	•	X	•	•	•	11	101	101	ED	2	2	8	
IM 2	Set interrupt			х		х	_			01 11	010 101	110 101	56 ED	•	•	•	:
	mode 2	•		^	•	^	•	•	•	01	011	110	ED 5E	2	2	8	!

NOTES: @ converts accumulator content into packed BCD following add or subtract with packed BCD operands. IFF indicates the interrupt enable flip-flop.

CY indicates the carry flip-flop.

\* indicates interrupts are not sampled at the end of EI or DI.

#### **16-BIT ARITHMETIC GROUP**

Mnemonic	Symbolic Operation	s	z		Fla	ngs	P/V	N	С		Opcod 543		Hex	No. of Bytes	No. of M Cycles	No. of T States	Com	merits
														•				
ADD HL, ss	HL ← HL+ss	•	•	Х	Х	Х	•	0	ŧ	00	ssi	001		1	3	1,1	ss	Reg
																	00	B¢
ADC HL, ss	HL←							_						_			01	D₿
	HL+ss+CY	ŧ	ŧ	Х	Х	Х	٧	O	ŧ	11	101	101	ED	2	4	15	10	HĻ
SBC HL, ss	HL←									01	ss1	010					11	SP
3	HL-ss-CY	<b>‡</b>	<b>‡</b>	X	Х	Х	٧	1	<b>‡</b>	11	101	101	ED	2	4	15		
										01	ss0	010						
ADD IX, pp	IX ← IX + pp	•	•	Χ	Х	Х	•	0	<b>‡</b>	11	011	101	DD	2	4	15	pp	Reg
										01	pp1	001					00	В¢
																	01	DE
																	10	IX
																	11	SP
ADD IY, rr	$IY \leftarrow IY + rr$	•	•	Χ	Х	Х	•	0	<b>‡</b>	11	111	101	FD	2	4	15	rr	Reg.
										00	rr1	001					00	В¢
INC ss	ss ss + 1	•	•	Х	•	Х	•	•	•	00	ss0	011		1	1	6	01	D₿
INC IX	IX + IX + 1	•	•	Х	•	Х	•	•	•	11	011	101	DD	2	2	10	10	ΙY
										00	100	011	23				11	SP
INC IY	IY ← IY + 1	•	•	Х	•	Х	•	•	•	11	111	101	FD	2	2	10		
										00	100	011	23					
DEC ss	ss - ss - 1	•	•	Х	•	Х	•	•	•	00	ss1	011		1	1	6		
DEC IX	IX ← IX – 1	•	•	X	•	Х	•	•	•	11	011	101	DD	2	2	10		
										00	101	011	2B					
DEC IY	IY ← IY – 1	•	•	х	•	х	•	•	•	11	111	101	FD	2	2	10		
										00	101	011	28		_			

#### **ROTATE AND SHIFT GROUP**

Mnemo	Symbolic onle Operation	s	z		Fla	gs		N	С		Opcod 543		Hex	No. of Bytes	No. of M Cycles	No. of T States	Comments
RLCA	CY = 7 = 0 =	•	•	x	0	x	•	0	<b>‡</b>	00	000	111	07	1	1	4	Rotate left circular
RLA	CY - 7 - 0	•	•	x	0	x	•	0	<b>‡</b>	00	010	111	<b>17</b>	1	1	4	accumulator. Rotate left accumulator.
RRCA	7 <b>0 CY</b>	•	•	x	0	X	•	0	<b>‡</b>	00	001	111	0F	1	1	4	Rotate right circular
RRA	7 — 0 CY	•	•	x	0	X	•	0	<b>‡</b>	00	011	111	1F	1	1	4	accumulator. Rotate right accumulator.

### ROTATE AND SHIFT GROUP (Continued)

Maar -	Symbolic	_	_		FI	age			_		Opcod			No. of	No. of M	No. of T	
	nic Operation	S	Z		Н		PΛ	/ N	<u> </u>	76	543	210	Hex	Bytes	Cycles	States	Comments
RLCr		<b>‡</b>	<b>‡</b>	x	0	x	P	0	• ‡	11 00	001	011 r	СВ	2	2	8	Rotate left circular register r.
RLC (HL	.) [ev]- 7	‡ 	<b>‡</b>	X	0	X	P	0	<b>‡</b>	11 00	001 000	011 110	СВ	2	4	15	r Reg 000 B
RLC (IX -		.‡	<b>‡</b>	X	0	X	P	0	*	11 11	011 001 ← d →	101 011	DD CB	4	6	23	001 C 010 D 011 E 001 H
RLC (IY +	+ d) }	<b>*</b>	<b>‡</b>	x	0	x	P	0		11	111	101	FD	4	6	23	101 L 111 A
iL m	$m = r_i(HL_i(IX + d))$	] ; ,(1Y+	‡ d)	x	0	x	Р	0	•	00	001 ← d → 000 010		СВ				Instruction format and states are as shown for
RCm	m = r, (HL), (IX + d)	‡ ),(IY+	<b>‡</b> d)	x	0	x	Ρ	0	*		001						RLCs. To form new opcode replace 000 or RLCs with
lR m	m = r, (HL), (iX + d)	•	‡ d)	x	0	x	Р	0	<b>‡</b>		011						shown code.
LA m	$CY \longrightarrow 7 \longrightarrow 0 \longrightarrow 0$ $m = r_1(HL), (IX + d)$	-	<b>‡</b> ď)	X	0	X	P	0	<b>‡</b>		100						
RA m	$m = r_i(HL), (IX + d)$			X	0	X	Ρ	0	*	. 2.	101						
RLm	$0+7 \rightarrow 0 \rightarrow CY$ m=r,(HL),(IX+d)			X	0	X	P	0	*		111						:
רט [	7-4 3-0 7-4 3-0 A (HL)		<b>‡</b>	x	0	x	Р	0	•	11 01	101 101	101 111	ED 6F		5		Rotate digit left and right between the accumu- lator and
RD [	74 30 74 30 A (HL)	*	<b>‡</b>	x	0	x	P	0	•	11 01	101 100	101 111	ED 67	2	5	18	location (HL). The content of the upper half of the accumulator is unaffected.

### BIT SET, RESET AND TEST GROUP

Mnemonic	Symbolic Operation	8	z		Fla H	gs	P/V	N	С	76	Opcod 543		Hex	No. of Bytes	No. of M Cycles	No. of T States	Con	nments
BIT b, r	Z←rb	х	<b>‡</b>	Х	1	х	х	0	•	11	001	011	СВ	2	2	8	r	Reg.
										01	b	ſ					000	В
BIT b, (HL)	Z ← (HL) <sub>b</sub>	Х	<b>‡</b>	Х	1	Х	Х	0	•	11	001	011	CB	2	3	12	001	С
										01	b	110					010	D
BIT b,(IX + d)b	$Z \leftarrow (IX + d)_b$	X	<b>‡</b>	X	1	X	X	0	•	11	011	101	DD	4	5	20	011	E
										11	001	011	CB				100	Н
											<b>-</b> d-	•					101	L
										01	b	110					111	Α
																	b	Bit Tested
BIT b, $(IY + d)_b$	Z ← (IY+d) <sub>b</sub>	X	<b>‡</b>	X	1	Х	X	0	•	11	111	101	FD	4	5	20	000	0
										11	001	011	CB				001	1
											<b>-</b> d→	•					010	2
										01	b	110					011	3
SET b, r	r <sub>b</sub> ←1	•	•	X	•	Х	•	•	. •	11	001	011	CB	2	2	8	100	4
										11	b	r					101	5
SET b, (HL)	(HL) <sub>b</sub> ← 1	•	•	X	•	X	•	•	•	11	001	011	CB	2	4	15	110	6
										11	b	110					111	7
SET b, $(1X + d)$	(IX+d) <sub>b</sub> <del>-</del> 1	•	•	X	•	X	•	•	•	11	011	101	DD	4	6	23		
		-								11	001	011	CB					
											-d-	•						
										11	b	110						
SET b, (IY+d)	$(iY+d)_b \leftarrow 1$	•	•	X	•	Х	•	•	•	11	111	101	FD	4	6	23		
										11	001	011	CB					
											+d →	•						
										11	b	110						
RES b, m	m <sub>b</sub> ← 0	•	•	X	•	X	•	•	•	10							To fo	kw usiA
	m≡r, (HL),														•			ode replace
	(IX+d), $(IY+d)$			•														of SET b, s
									•									10 Flags
																	and	
																		s for SET
																	instr	uction.

NOTE: The notation  $m_b$  indicates location  $m_s$  bit b (0 to 7).

#### **JUMP GROUP**

Mnemonic	Symbolic Operation	s	z		FI	<b>ag</b> s		۷N	С		Opco 543	<b>ie</b> 210	Hex	No. of Bytes	No. of M Cycles	No. of T States	•	nments
JP nn	PC ← nn	•	•	х	•	Х	•	•	•	11	000	011	СЗ	3	3	10	œ	Condition
											←n-	•					000	NZ (non-zero)
											<b>←</b> n~	•					001	Z (zero)
JP cc, nn	If condition cc		•	Х	•	Χ	•	•	•	11	œ	010		3	3	10	010	NC (non-carry)
	is true PC←nn,										←n-	•					011	C (carry)
	otherwise										<b>←</b> n-	•					100	PO (parity odd)
_	continue																101	PE (parity even)
JR e	PC ← PC+e	•	•	Х	•	Х	•	•	•	00		000	18	2	3	12	110	P (sign positive)
	_									•	-e-2						111	M (sign negiative
JR C, e	#C=0,	•	•	X	•	X	•	•	•	00		000	38	2	2	7	If cor	ndition not met.
	continue									•	-e-2	<b>→</b>						
	HC=1,													2	3	12	If cor	ndition is met.
	PC ← PC+e															•		
JR NC, e	IFC=1,	•	•	Х	•	Х	•	•	•	00			30	2	2	7	If cor	ndition not met.
	continue									•	-e-2	<b>→</b>						
	If C=0,													2	3	12	If cor	ndition is met.
JP Z, e	PC ← PC+e	_	_	v		v								_	_	_		
	continue	•	•	Х	•	X	•	•	•	00		000	28	2	2	7	If cor	ndition not met.
	If Z = 1,									•	-e-2	-		•		40		and the second
	PC←PC+e													2	3	12	IT CO	ndition is met.
	IfZ=1.	_	_	x	_	х	_		_	00	100	000	20	2	2	7	16	
	continue	٠	Ť	^	•	^	٠	•	•		-e-2		20	2	2	′	II COL	ndition not met.
	If Z = 0.									•	6-2			2	3	12	lf aar	ndition is met.
	PC + PC+e													2	3	12	II COI	idition is met.
	PC + HL			¥		Y	•			11	101	001	'E9	1	1	4		
, ,	PC+IX	•					•			11	011	101	DD	2	2	8		
. ()				^		^				11	101	001	E9	-	2	Ü		
JP (IY)	PC ← IY			x		x	•			11	111	101	FD	2	2	8		
,				^	-	^	-	-	-	11	101	001	E9	-	-	U		
DJNZ, e	B <b>←</b> B – 1			x	•	x		•		00		000	10	2	2	8	If B =	n
=	If B = 0,			••							-e-2		••	~	-	Ū	., 5	-
	continue																	
	lf B≠0,													2	3	13	If B≠	0.
	PC ← PC+e													_	-			

NOTES: e represents the extension in the relative addressing mode.
e is a signal two's complement number in the range < - 126, 129 >.
e - 2 in the opcode provides an effective address of pc + e as PC is incremented by 2 prior to the addition of e.

#### **CALL AND RETURN GROUP**

Mnemonic	Symbolic Operation	s	z		Fia H	igs		/N	С	76	Opcod 543		Hex	No. of Bytes	No. of M Cycles	No. of T States	Com	ments
CALL nn	(SP-1)←PC <sub>H</sub>	•	•	х	•	Х	•	•	•	11	001	101	CD	3	5	17		
	(SP-2)←PC <sub>L</sub> PC ← nn.										+n→							
CALL cc nr	PC ← nn, If condition			¥		х				11	←n→ cc	100		3	3	10	If co.is	s false.
O/1LL 00,111	cc is false	_	-	^	-	^		-		••	+ n →			Ü	·			5 Ka300.
	continue, otherwise										+-n-			3	5	17	If oc is	s true.
	same as CALL nn																	
RET	PC <sub>L</sub> ← (SP) PC <sub>H</sub> ←(SP+1)	•	•	x	•	X	•	•	•	11	<b>0</b> 01	001	C9	1	3	10		
RET ∞	If condition cc is false	•	•	X	•	X	•	•	•	11	cc	000		1	1	5	If cc is	s false.
	continue,													/1	3	11	If oc is	s true.
	same as RET																	Condition
																		NZ (non-zero)
																		Z (zero)
																		NC (non-carry)
RETI	Return from	•	•	Х	•	Х	•	•	•	11	101	101	ED	2	4	14		C (carry)
	interrupt									01	001	101	4D					PO (parity odd)
RETN <sup>1</sup>	Return from	•	•	Х	•	Х	•	•	•	11	101	101	ED	2	4	14		PE (parity even)
	non-maskable									01	000	101	45				110	P (sign positive)
	interrupt																	·M (sign negative)
RST p	(SP-1)←PCH	•	•	X	•	Х	•	•	•	11	t	111		1	3	11	t	P
	(SP-2)←PC <sub>L</sub>																000	
	PC <sub>H</sub> ← 0																	08H
	PC <sub>L</sub> ← p																-	10H
																	011	18H
																	100	20H
																		28H
																	110	30H
																	111	38H

NOTE: ¹RETN loads IFF2 → IFF1

#### **INPUT AND OUTPUT GROUP**

Mnemonic	Symbolic Operation	S	Z		FI	age		۷N	C	76	Opcod 543	le 210	Hex	No. of Bytes	No. of M Cycles	No. of T States	Comments
IN A, (n)	A ← (n)	•	٠.	x	•	X	•	•	•	11	011	01	DB	2	3	11	n to A <sub>0</sub> ~ A <sub>7</sub>
											<b>←</b> n-	•					Acc. to A <sub>8</sub> ~ A <sub>15</sub>
IN r, (C)	r +- (C)	<b>‡</b>	#	Х	<b>‡</b>	Х	Ρ	0	•	11	101	101	ED	2	3	12	C to Ao ~ A <sub>7</sub>
	if $r = 110$ only									01	r	000					B to A <sub>8</sub> ~ A <sub>15</sub>
	the flags will																
	be affected																
			①	)													
INI	(HL) ← (C)	Х	ŧ	Х	Х	Х	Х	1	х	11	101	101	ED	2	4	16	C to A <sub>0</sub> ~ A <sub>7</sub>
	B ← B – 1									10	100	010	A2				B to A <sub>8</sub> ~ A <sub>15</sub>
	HL+HL+1		2	)													0 10
INIR	(HL) ← (C)	X	1	Х	Х	X	Х	1	х	11	101	101	ED	2	5	21	C to A <sub>0</sub> ~ A <sub>7</sub>
	B ← B – 1									10	110	010	B2		(If B≠0)		B to A <sub>8</sub> ~ A <sub>15</sub>
	HL ← HL+1													2	4	16	- 10 1 10
	Repeat until								s						(If B = 0)		
	B=0														·,		
			1	)													
IND	(HL) ← (C)	Х	Ť	х	х	Х	Х	1	х	11	101	101	ΕD	2	4	16	C to A <sub>0</sub> ~ A <sub>7</sub>
	B ← B – 1									10	101	010	AA				B to A <sub>8</sub> ~ A <sub>15</sub>
	HL+HL-1		<b>②</b>	ı													- 101 6 1113
INDR	(HL) ← (C)	Х	$\stackrel{\smile}{1}$	х	х	Х	х	1	х	11	101	101	ED	2	5	21	C to A <sub>0</sub> ~ A <sub>7</sub>
	B - B-1									10	111	010	BA		(If B≠0)		B to A <sub>8</sub> ~ A <sub>15</sub>
	HL+HL-1													2	4	16	- 101 0 1110
	Repeat until													_	(If B = 0)		
	B=0														<b>(</b> )		
OUT (n), A	(n) <del>-</del> A	•	•	Х	•	X	•	•.	•	11	010	011	D3	2	3	11	n to A <sub>0</sub> ~ A <sub>7</sub>
-											<b>+</b> n→						Acc. to A <sub>8</sub> ~ A <sub>15</sub>
OUT (C), r	(C) ← r	•	•	X	•	Х	•	•	•	11	101	101	ED	2	3	12	C to Ao ~ A7
										01	r	001					B to A <sub>8</sub> ~ A <sub>15</sub>
			1														•
OUTI	(C) ← (HL)	X	#	X	X	X	X	1	Х	11	101	101	ED	2 -	4	16	C to A <sub>0</sub> ~ A <sub>7</sub>
	B ← B – 1									10	100	011	A3				B to A <sub>8</sub> ~ A <sub>15</sub>
	HL←HL+1		2														•
OTIR	(C) + (HL)		1	X	Х	X	Х	1	Х	11	101	101	ED	2	5	21	C to A <sub>0</sub> ~ A <sub>7</sub>
	B ← B – 1									10	110	011	<b>B</b> 3		(If B≠0)		B to A <sub>8</sub> ~ A <sub>15</sub>
	HL+HL+1													2	4	16	0 .0
	Repeat until														(If $B = 0$ )		
	B=0														·/		
			①														•
DTUC	(C) ← (HL)	X	*	Х	X	X	X	1	Х	11	101	101	ED	2	4	16	C to $A_0 \sim A_7$
	B ← B – 1									10	101	011	AB				B to A <sub>B</sub> ~ A <sub>15</sub>
	HL ← HL – 1																5 10
			@														
OTOR	(C) ← (HL)		$\tilde{1}$	х	Х	Х	Х	1	Х	11	101	101	ED	2	5	21	C to A <sub>0</sub> ~ A <sub>7</sub>
	B ← B – 1									10	111	011			(If B≠0)		B to A <sub>8</sub> ~ A <sub>15</sub>
	HL+HL-1													2	4	16	00
	Repeat until														(If B = 0)	• •	
	B=0														···/		

NOTES: ① If the result of B – 1 is zero, the Z flag is set; otherwise it is reset.
② Z flag is set upon instruction completion only.

#### **SUMMARY OF FLAG OPERATION**

	D <sub>7</sub>							Do	
Instructions	s	Z		Н		P/V	N	C	Comments
ADD A, s; ADC A, s	<b>‡</b>	#	X	<b>‡</b>	X	٧	0	<b>‡</b>	8-bit add or add with carry.
SUB s; SBC A, s; CP s; NEG	<b>‡</b>	#	X	<b>‡</b>	Х	٧	1	<b>‡</b>	bit subtract, subtract with carry, compare and negate accumulator.
AND s	<b>‡</b>	#	Х	1	Х	Ρ	0	0	Logical operation.
OR s, XOR s	<b>‡</b>	#	Х	0	Х	Ρ	0	0	Logical operation.
INCs	<b>‡</b>	<b>‡.</b>	Х	<b>‡</b>	Х	٧	0	•	8-bit increment.
DEC s	<b>‡</b>	<b>‡</b>	Х	<b>\$</b>	Х	٧	1	•	8-bit decrement.
ADD DD, as	•	•	Х	Х	Х	•	0	<b>‡</b>	16-bit add.
ADC HL, ss	<b>‡</b>	<b>‡</b>	Х	Х	Х	٧	0	<b>‡</b>	16-bit add with carry.
SBC HL. ss			Х	Х	Х	٧	1	<b>‡</b>	16-bit subtract with carry.
RLA; RLCA; RRA; RRCA	•	•	Х	0	Х	•	0	<b>‡</b>	Rotate accumulator.
RL m; RLC m; RR m; RRC m; SLA m; SRA m; SRL m	<b>‡</b>	<b>‡</b>	Х	0	X	P	0	<b>‡</b>	Rotate and shift locations.
RLD: RRD		#	Х	0	Χ	Р	0	•	Rotate digit left and right.
DAA			X	ŧ	X	P	•	<b>‡</b>	Decimal adjust accumulator.
CPL	•	•	X	1	X	•	1	•	Complement accumulator.
SCF	•	•	X	0	Х	•	0	1	Set carry.
CCF	•	•	X	X	X	•	Ö	<b>‡</b>	Complement carry.
IN r (C)			X	0	X	Ρ	0	•	Input register indirect.
INI; IND; OUTI; OUTD	X	į	X	X	Х	Х	1	•	Block input and output, $Z = 1$ if $B \neq 0$ , otherwise $Z = 0$ .
INIR; INDR; OTIR; OTDR	X	1	Х	Х	Х	Х	1	•	Block input and output. $Z = 1$ if $B \neq 0$ , otherwise $Z = 0$ .
LDI; LDD	X	X	X	0	X	#	0	•	Block transfer instructions. P/V = 1 if BC ≠ 0, otherwise P/V = 0.
LDIR; LDDR	X	X	Х	Ô	X	Ó	0	•	Block transfer instructions. $P/V = 1$ if $BC \neq 0$ , otherwise $P/V = 0$
CPI; CPIR; CPD; CPDR	X	<b>‡</b>	X	X	X	<b>‡</b>	1	•	Block search instructions. $Z = 1$ if $A = (HL)$ , otherwise $Z = 0$ . $P/V = 1$ if $BC \neq 0$ , otherwise $P/V = 0$ .
LD A; I, LD A, R	<b>‡</b>	#	Х	0	X	IFF	0	•	IFF, the content of the interrupt enable flip-flop, (IFF <sub>2</sub> ), is copied into the P/V flag.
BIT b, s	X	#	Х	1	Χ	Х	0	•	The state of bit b of location s is copied into the Z flag.

#### **SYMBOLIC NOTATION**

ymbol	Operation	Symbol	Operation
S	Sign flag, S = 1 if the MSB of the result is 1.	<b>‡</b>	The flag is affected according to the result of the
Z	Zero flag. $Z = 1$ if the result of the operation is 0.		operation.
PN	Parity or overflow flag. Parity (P) and overflow (V)	•	The flag is unchanged by the operation.
	share the same flag. Logical operations affect	0	The flag is reset by the operation.
	this flag with the parity of the result while	1	The flag is set by the operation.
	arithmetic operations affect this flag with the	X	The flag is indeterminate.
	overflow of the result. If P/V holds parity: P/V = 1	V	P/V flag affected according to the overflow result
	if the result of the operation is even; PN = 0 if		of the operation.
	result is odd. If P/V holds overflow, P/V = 1 if the result of the operation produced an overflow. If	Р	P/V flag affected according to the parity result of the operation.
	P/V does not hold overflow, $P/V = 0$ .	r	Any one o the CPU registers A, B, C, D, E, H, L.
Н*	Half-carry flag. H = 1 if the add or subtract operation produced a carry into, or borrow from,	s .	Any 8-bit location for all the addressing modes allowed for the particular instruction.
	bit 4 of the accumulator.	SS	Any 16-bit location for all the addressing modes
N*	Add/Subtract flag. N = 1 if the previous		allowed for that instruction.
	operation was a subtract.	ü	Any one of the two index registers IX or IY.
С	Carry/Link flag. C = 1 if the operation produced	R	Refresh counter.
_	a carry from the MSB of the operand or result.	n	8-bit value in range < 0, 255 >.
		nn	16-bit value in range < 0, 65535 >.

<sup>\*</sup>H and N flags are used in conjunction with the decimal adjust instruction (DAA) to properly correct the result into packed BCD format following addition or subtraction usin. perands with packed BCD format.

#### **CPU TIMING**

The Z80 CPU executes instructions by proceeding through a specific sequence of operations:

- Memory read or write
- I/O device read or write
- Interrupt acknowledge

The basic clock period is referred to as a T time or cycle, and three or more T cycles make up a machine cycle (M1, M2 or M3 for instance). Machine cycles can be extended either by the CPU automatically inserting one or more Wait states or by the insertion of one or more Wait states by the user.

**Instruction Opcode Fetch.** The CPU places the contents of the Program Counter (PC) on the address bus at the start of the cycle (Figure 5). Approximately one-half clock cycle later, MREQ goes active. When active, RD indicates that the memory data can be enabled onto the CPU data bus.

The CPU samples the  $\overline{WAIT}$  input with the falling edge of clock state  $T_2$ . During clock states  $T_3$  and  $T_4$  of an  $\overline{M1}$  cycle, dynamic RAM refresh can occur while the CPU starts decoding and executing the instruction. When the Refresh Control signal becomes active, refreshing of dynamic memory can take place.

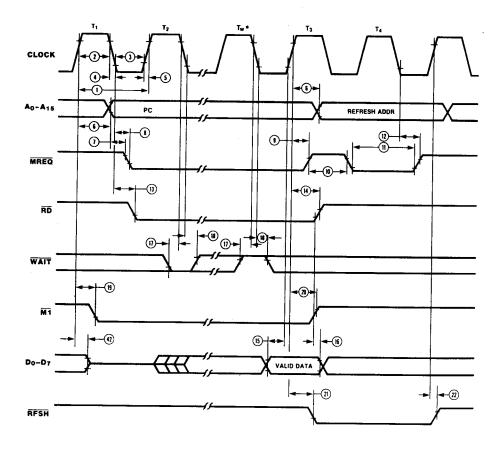


Figure 5. Instruction Opcode Fetch

Memory Read or Write Cycles. Figure 6 shows the timing of memory read or write cycles other than an opcode fetch (M1) cycle. The MREQ and RD signals function exactly as in the fetch cycle. In a memory write cycle, MREQ also

becomes active when the address bus is stable. The  $\overline{WR}$  line is active when the data bus is stable, so that it can be used directly as an  $R\overline{W}$  pulse to most semiconductor memories.

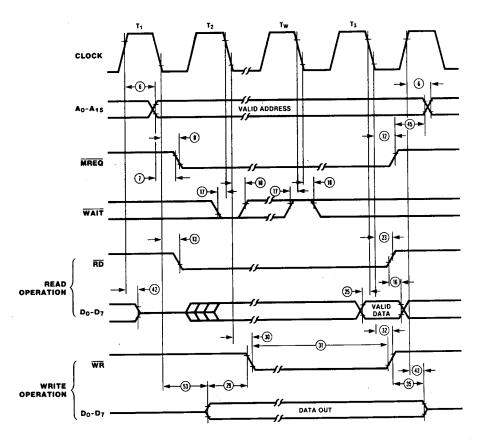
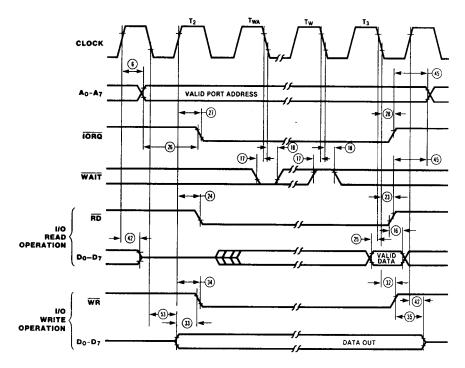


Figure 6. Memory Read or Write Cycles

Input or Output Cycles. Figure 7 shows the timing for an I/O read or I/O write operation. During I/O operations, the CPU automatically inserts a single Wait state ( $T_{WA}$ ). This

extra Wait state allows sufficient time for an 1/O port to decode the address from the port address lines.

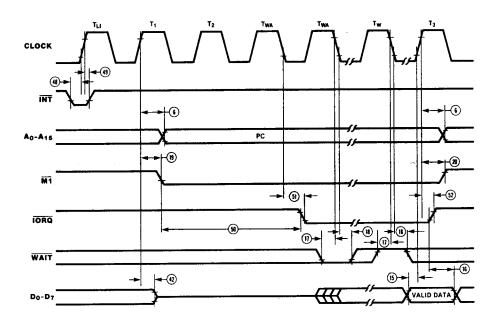


T<sub>WA</sub> = One wait cycle automatically inserted by CPU.

Figure 7. Input or Output Cycles

Interrupt Request/Acknowledge Cycle. The CPU samples the interrupt signal with the rising edge of the last clock cycle at the end of any instruction (Figure 8). When an interrupt is accepted, a special  $\overline{\text{M1}}$  cycle is generated.

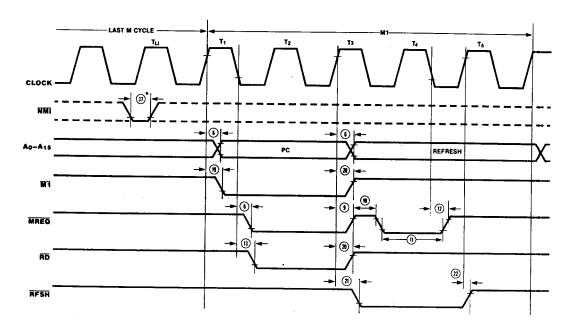
During this  $\overline{\text{M1}}$  cycle,  $\overline{\text{IORQ}}$  becomes active (instead of  $\overline{\text{MREQ}}$ ) to indicate that the interrupting device can place an 8-bit vector on the data bus. The CPU automatically adds two Wait states to this cycle.



27

Non-Maskable Interrupt Request Cycle. NMI is sampled at the same time as the maskable interrupt input INT but has higher priority and cannot be disabled under software control. The subsequent timing is similar to that of a normal

memory read operation except that data put on the bus by the memory is ignored. The CPU instead executes a restart (RST) operation and jumps to the  $\overline{\text{NMI}}$  service routine located at address 0066H (Figure 9).

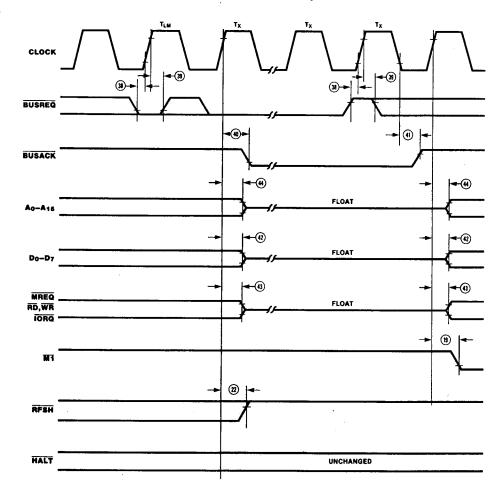


<sup>\*</sup>Although NMI is an asynchronous input, to guarantee its being recognized on the following machine cycle, NMI's falling edge must occur no later than the rising edge of the clock cycle preceding the last state of any instruction cycle (T<sub>LI</sub>).

Figure 9. Non-Maskable Interrupt Request Operation

Bus Request/Acknowledge Cycle. The CPU samples BUSREQ with the rising edge of the last clock period of any machine cycle (Figure 10). If BUSREQ is active, the CPU sets its address, data, and MREQ, IORQ, RD, and WR lines

to a high-impedance state with the rising edge of the next clock pulse. At that time, any external device can take control of these lines, usually to transfer data between memory and I/O devices.

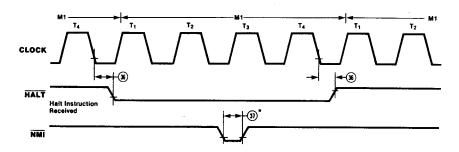


NOTES: 1)  $T_{LM}$  = Last state of any M cycle. 2)  $T_X$  = An arbitrary clock cycle used by requesting device.

Figure 10. BUS Request/Acknowledge Cycle

Halt Acknowledge Cycle. When the CPU receives a HALT instruction, it executes NOP states until either an INT or NMI input is received. When in the Halt state, the HALT output is

active and remains so until an interrupt is received (Figure 11). INT will also force a Halt exit.



\*Although NMI is an asynchronous input, to guarantee its being recognized on the following machine cycle, NMI's falling edge must occur no later than the rising edge of the clock cycle preceding the last state of any instruction cycle (T<sub>L1</sub>).

Figure 11. Halt Acknowledge

Reset Cycle. RESET must be active for at least three clock cycles for the CPU to properly accept it. As long as RESET remains active, the address and data buses float, and the control outputs are inactive. Once RESET goes inactive, two

internal T cycles are consumed before the CPU resumes normal processing operation. RESET clears the PC register, so the first opcode fetch will be to location 0000H (Figure 12).

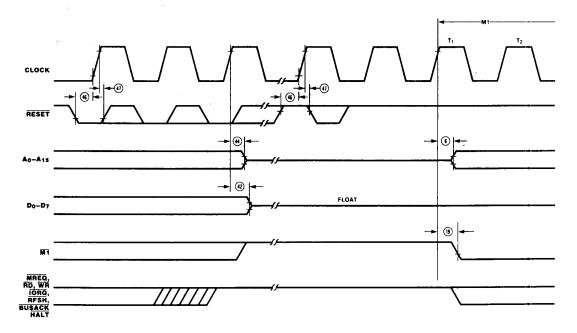


Figure 12. Reset Cycle

Power-Down mode of operation (Only applies to CMOS Z80 CPU).

 $\textbf{CMOSZ80}\,\textbf{CPU}\,\textbf{supports}\,\textbf{Power-Down}\,\textbf{mode}\,\textbf{of}\,\textbf{operation}.$ 

This mode is also referred to as the "standby mode", and supply current for the CPU goes down as low as 10 uA (Where specified as lcc<sub>2</sub>).

**Power-Down Acknowledge Cycle.** When the clock input to the CPU is stopped at either a High or Low level, the CPU stops its operation and maintains all registers and control signals. However,  $I_{cc2}$  (standby supply current) is guaranteed only when the system clock is stopped at a Low

level during  $T_4$  of the machine cycle following the execution of the HALT instruction. The timing diagram for the power-down function, when implemented with the HALT **instruction, is shown in Figure 13.** 

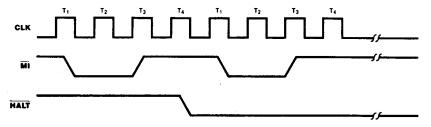


Figure 13. Power-Down Acknowledge

**Power-Down Release Cycle.** The system clock must be supplied to the CPU to release the power-down state. When the system clock is supplied to the CLK input, the CPU restarts operations from the point at which the power-down state was implemented.

The timing diagrams for the release from power-down mode are shown in Figure 14.

#### NOTES:

- When the external oscillator has been stopped to enter the power-down state, some warm-up time may be required to obtain a stable clock for the release.
- 2) When the HALT instruction is executed to enter the power-down state, the CPU will also enter the Halt state. An interrupt signal (either NMI or INT) or a RESET signal must be applied to the CPU after the system clock is supplied in order to release the power-down state.

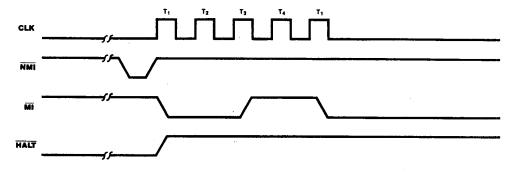


Figure 14a.

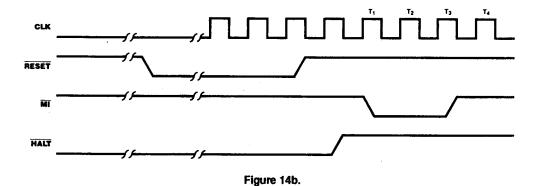


Figure 14c.

Figure 13. Power-Down Release

#### **ABSOLUTE MAXIMUM RATINGS**

Voltage on V <sub>CC</sub> with respect to	V <sub>SS</sub> 0.3V to +7V
Voltages on all inputs with resp	pect
to V <sub>SS</sub>	0.3V to V <sub>CC</sub> + 0.3V
Operating Ambient	
Temperature	See Ordering Information
Storage Temperature	65°C to +150°C

Stresses greater than those listed under Absolute Maximum Ratings may cause permanent damage to the device. This is a stress rating only; operation of the device at any condition above those indicated in the operational sections of these specifications is not implied. Exposure to absolute maximum rating conditions for extended periods may affect device reliability.

#### STANDARD TEST CONDITIONS

The DC Characteristics and capacitance sections below apply for the following standard test conditions, unless otherwise noted. All voltages are referenced to GND (0V). Positive current flows into the referenced pin.

Available operating temperature ranges are:

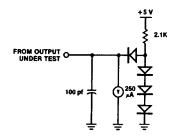
■ S = 0°C to +70°C Voltage Supply Range:

NMOS: +4.75V ≤ VCC ≤ +5.25V CMOS: +4.50V ≤ VCC ≤ +5.50V

■ E=  $-40^{\circ}$ C to  $100^{\circ}$ C, +4.50V  $\leq$  VCC  $\leq$  +5.50V

All ac parameters assume a load capacitance of 100 pf. Add 10 ns delay for each 50 pf increase in load up to a maximum of 200 pf for the data bus and 100 pf for address and control lines. AC timing measurements are referenced to 1.5 volts (except for clock, which is referenced to the 10% and 90% points).

The Ordering Information section lists temperature ranges and product numbers. Package drawings are in the Package Information section. Refer to the Literature List for additional documentation.



### DC CHARACTERISTICS (Z84C00/CMOS Z80 CPU)

Symbol	Parameter	Min	Max	Unit	Condition
V <sub>ILC</sub>	Clock Input Low Voltage	-0.3	0.45	٧	
VIHC	Clock Input High Voltage	V <sub>CC</sub> 6	V <sub>CC</sub> +.3	٧	
$V_{IL}$	Input Low Voltage	-0.3	0.8	٧	
V <sub>IH</sub>	Input High Voltage	2.2	Vcc	V	
V <sub>OL</sub>	Output Low Voltage		0.4	٧	$I_{OL} = 2.0  \text{mA}$
V <sub>OH1</sub>	Output High Voltage	2.4		٧	$I_{OH} = -1.6  \text{mA}$
V <sub>OH2</sub>	Output High Voltage	V <sub>CC</sub> - 0.8		٧	$I_{OH} = -250 \mu\text{A}$
lcc <sub>1</sub>	Power Supply Current 4 MHz 6 MHz 8 MHz 10 MHz 20 MHz		20 30 40 50	mA mA mA	$V_{CC} = 5V$ $V_{IH} = V_{CC} - 0.2V$ $V_{IL} = 0.2V$
Icc <sub>2</sub>	Standby Supply Current		100	mΑ μΑ	$V_{\infty} = 5V$ $V_{CC} = 5V$
					CLK = (0) $V_{IH} = V_{CC} - 0.2V$ $V_{IL} = 0.2V$
ILI	Input Leakage Current	-10	10	μΑ	$V_{IN} = 0.4 \text{ to } V_{CC}$
ILO	3-State Output Leakage Current in Float	-10	10 <sup>2</sup>	μΑ	$V_{OUT} = 0.4$ to $V_{CC}$

#### **CAPACITANCE**

Symbol	Parameter	Min	Max	Unit
C <sub>CLOCK</sub>	Clock Capacitance		10	prf
C <sub>IN</sub>	Input Capacitance		5	pf
C <sub>OUT</sub>	Output Capacitance		15	pif

T<sub>A</sub> = 25°C, f = 1 MHz. Unmeasured pins returned to ground.

<sup>1.</sup> Measurements made with outputs floating.
2. A<sub>15</sub>·A<sub>0</sub>, D<sub>7</sub>·D<sub>0</sub>, MREQ, IORQ, RD, and WR.
3. I<sub>CC<sub>2</sub></sub> standby supply current is guaranteed only when the supplied clock is stopped at a low level during T<sub>4</sub> of the machine cycle immediately following the execution of a HALT instruction.

### AC CHARACTERISTICS† (Z84C00/CMOS Z80 CPU)

 $V_{cc}$ =5.0V  $\pm$  10%, unless otherwise specified

			Z84	C0004	*Z84	C0006	784	C0008	Z840	C0010	784	C0020[1]	Unit	Note
No	Symbol	Parameter		Max		Max		Max	Min	Max	Min		OH	:40(6
1	TcC	Clock Cycle time	250	, DC	162	DC	125	DC	100*	DC	50*	DC	nS	
2	TwCh	Clock Pulse width (high)	110	DC	65	DC	55	DC	40	DC	20	DC	nS	
3	TwCi	Clock Pulse width (low)		DC	65	DC	55	DC	40	DC	20	DC	nS	
4	TfC	Clock Fall time		30		20		10		10		10	nS	
5	TrC	Clock Rise time		30		20		10		10		10	nS	
6	TdCr(A)	Address vaild from Clock Rise	i	110		90		80		65		57	nS	[2]
7	TdA(MREQf)	Address valid to /MREQ Fall	65*		35*		20*		5*		-15*		nS	
8	TdCf(MREQf)	Clock Fail to MREQ Fail delay		85		70		60		55		40	nS	
9	TdCr(MREQr)	Clock Rise to /MREQ Rise delay		85		70		60		55		40	nS	
10	TwMREQh	/MREQ pulse width (High)	110*		65*		45**		30*		10*		nS	[3]
	TwMREQI	/MREQ pulse width (low)	220*		132*		100*		75*		25*		nS	[3]
		Clock Fall to MREQ Rise delay		85		70		60		55		40	nS	• •
	TdCf(RDf)	Clock Fall to /RD Fall delay		95		80		70		<b>6</b> 5		40	nS	
	TdCr(RDr)	Clock Rise to /RD Rise delay		85		70		60		55		40	nS	
15	TsD(Cr)	Data setup time to Clock Rise	35		30		30		25		12		nS	
	ThD(RDr)	Data hold time after /RD Rise	0		0		0		0		0		nS	
	TsWAIT(Cf)	WAIT setup time to Clock Fall	70		60		50		20		7.5		nS	
	ThWAIT(Cf)	WAIT hold time after Clock Fall	10		10		10		10		10		nS	
	TdCr(M1f)	Clock Rise to /M1 Fall delay		100		80	•	70		65		<b>4</b> 5	nS	
20	TdCr(M1r)	Clock Rise to /M1 Rise delay		100		80		70		<b>6</b> 5		<b>4</b> 5	nS	
	TdCr(RFSHf)	Clock Rise to /RFSH Fall delay		130		110		95		80		60	nS	
	TdCr(RFSHr)	Clock Rise to /RFSH Rise delay		120		100		85		80		60	nS	
	TdCf(RDr)	Clock Fall to /RD Rise delay		85		70		60		55		40	nS	
	TdCr(RDf)	Clock Rise to /RD Fall delay		85		70		60		<b>5</b> 5		40	nS	
25	TsD(Cf)	Data setup to Clock Fall during												
		M2, M3, M4 or M5 cycles	50		40		30		25		12		nS	
26	TdA(IORQf)	Address stable prior to /IORQ Fall	180*		107*		75*		50*		0*		nS	
27	TdCr(IORQf)	Clock Rise to /IORQ Fall delay		75		<b>65</b> .		55		50		40	nS	
28	TdCf(IORQr)	Clock Fall to /IORQ Rise delay		85		70		60		55		40 '	nS	
29	TdD(WRf)Mw	Data stable prior to /WR Fall	80*		22*		5*		40 <b>*</b>		-10*		nS	
30	TdCf(WRf)	Clock Fall to /WR Fall delay		80		70	···	60		55		40	nS	
31	TwWR	MR pulse width	220*		132*		100*		75*		25*		nS	
32	TdCf(WRr)	Clock Fall to MR Rise delay		80		70		60		55		40	nS	
33	TdD(WRf)IO	Data stable prior to /WR Fall	-10*		-55*		-55*		-10*		-30*		nS	
34	TdCr(WRf)	Clock Rise to /WR Fall delay		<b>6</b> 5		60		60		50		40	nS	
35	TdWRr(D)	Data stable from MR Rise	60*		30*		15*		10*		0*		nS	
<b>36</b> '	TdCf(HALT)	Clock Fall to /HALT 'L' or 'H'		300		260		225		90		70	nS	
	TwnM!	/NMI pulse width	80		60		60		60		60		nS	
8	TsBUSREQ	/BUSREQ setup time	50		50		40		30		15		nS	
(	(Cr)	to Clock Rise												

<sup>\*</sup>For clock periods other than the minimums shown, calculate parameters using the table on the following page. Calculated values above assumed TrC = TfC = 20 ns.

<sup>†</sup>Units in nanoseconds (ns). †† For loading ≥ 50 pf. Decrease width by 10 ns for each additional 50 pf...

<sup>\*\*4</sup> MHz CMOS Z80 is obsoleted and replaced by 6 MHz

#### AC CHARACTERISTICS<sup>†</sup> (Z84C00/CMOS Z80 CPU; Continued)

 $V_{\infty}$ =5.0V ± 10%, unless otherwise specified

			Z84C0004" Z84C0006		Z840	Z84C0008		Z84C0010		Z84C0020[1]		Note		
No	Symbol	Parameter	Min	Max	Min	Max	Min	Max	Min	Max		Мах		
39	ThBUSREQ	/BUSREQ hold time	10		10		10		10		10		nS	
	(Cr)	after Clock Rise												
40	TdCr	Clock Rise to /BASACK		100		90		80		75		40	nS	
	(BUSACKI)	Fall delay												
41	TdCf	Clock Fall to /BASACK		100		90		80		75		40	nS	
	(BUSACKr)	Rise delay												
42	TdCr(Dz)	Clock Rise to Data float delay		90		80		70		<b>65</b>		40	nS	
43	TdCr(CTz)	Clock Rise to Control Outputs												
		Float Delay (/MREQ, /IORQ,												
		/RD and /WR)		80		70		60		65		40	nS	
44	TdCr(Az)	Clock Rise to Address		90		80		70		75		40	n\$	
		float delay												
45	TdCTr(A)	Address Hold time from /MREQ,	80*		35*		20*		20*		0*		nS	
		/IORQ, /RD or /WR												
46	TsRESET(Cr)	/RESET to Clock Rise setup time	60		60		45		40		15		nS	
47	ThRESET(Cr)	/RESET to Clock Rise Hold time	10		10		10		10		10		nS	
48	TsINTf(Cr)	/INT Fall to Clock Rise	80		70		55		50		15		nS	
		Setup Time												
49	ThINTr(Cr)	/INT Rise to Clock Rise	10		10		10		10		10		nS	
		Hold Time												
50	TdM1f	/M1 Fall to /IORQ Fall delay	565	,	359	,	270*	,	220	•	100	*	nS	
	(IORQf)													
51	TdCf(IORQf)	/Clock Fall to /IORQ Fall delay		<b>8</b> 5		70		60		55		45	пS	
52	TdCf(IORQr)	Clock Rise to /IORQ Rise delay		85		70		60		55		<b>4</b> 5	пS	
53	TdCf(D)	Clock Fall to Data Valid delay		150		130		115		110		<b>7</b> 5	nS	

- Notes:
  For Clock periods other than the minimum shown, calculate parameters using the following table.
- Calculated values above assumed TrC = TfC = maximum.
  \*\* 4 MHz CMOS Z80 is obsoleted and replaced by 6 MHz

- [1] Z84C0020 parameters are guuaranteed with 50pF load Capacitance.
  [2] If Capacitive Load is other than 50pF, please use Figure 1. to calculate the value.
  [3] Increasing delay by 10nS for each 50pF increase in loading, 200pF max for data lines, and 100pF for control lines.

#### **FOOTNOTES TO AC CHARACTERISTICS**

No	Symbol	Parameter	Z84C0004°	Z84C0006	Z84C0008	Z84C0010	Z84C0020
1	TcC	TwCh + TwCl + TrC + TfC					
7	TdA(MREQf)	TwCh + TfC	-65	-50	-45	-45	-45
10	TwMREQh	TwCh + TfC	-20	-20	-20	-20	-20
11	TwMREQI	TcC	-30	-30	-25	-25	-25
26	TdA(IORQf)	TcC	-70	-55	-50	-50	-50
29	TdD(WRf)	TcC	-170	-140 .	-120	-60	-60
31	TwWR	TcC /	-30	-30	-25	-25	-25
33	TdD(WRf)	TwCl + TrC	-140	-140	-120	-60	-60
35	TdWRr(D)	TwCl + TrC	-70	<b>-5</b> 5	-50	-40	-25
45	TdCTr(A)	TwCl + TrC	-50	-50	-45	-30	-30
50	TdM1f(IORQf)	2TcC + TwCh + TfC	-65	-50	-45	-30	-30
C Test	Conditions: V <sub>IH</sub> = 2.0 V <sub>II</sub> = 0.8		V <sub>IHC</sub> =	V <sub>CC</sub> -0.6 V 0.45 V	FLOAT = 1	±0.5 V	

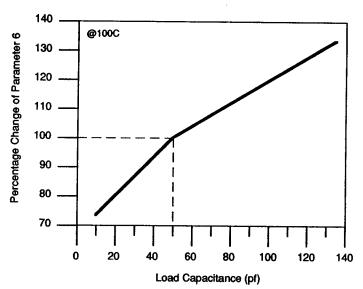


Figure 1. Address Delay Characteristics (Parameter 6)

#### DC CHARACTERISTICS (Z8400/NMOS Z80 CPU)

All parameters are tested unless otherwise noted.

Symbol	Parameter	Min	Max	Unit	Test Condition
V <sub>ILC</sub>	Clock Input Low Voltage	-0.3	0.45	v	
V <sub>IHC</sub>	Clock Input High Voltage	V <sub>CC</sub> 6	V <sub>CC</sub> +.3	٧	
V <sub>IL</sub>	Input Low Voltage	- 0.3	0.8	V	
V <sub>IH</sub>	Input High Voltage	2.0 <sup>1</sup>	Vcc	V	
VOL	Output Low Voltage		0.4	V	$I_{OL} = 2.0  \text{mA}$
V <sub>OH</sub>	Output High Voltage	2.4 <sup>1</sup>		٧ .	I <sub>OH</sub> = -250 μA
lcc.	Power Supply Current	•	200	mA	Note 3
lLi	Input Leakage Current		10	μΑ	$V_{IN} = 0$ to $V_{CC}$
LO	3-State Output Leakage Current in Float	-10	10 <sup>2</sup>	μA	V <sub>OUT</sub> = 0.4 to V <sub>C</sub> (

For military grade parts, refer to the Z80 Military Electrical Specification.
 A<sub>15</sub>-A<sub>0</sub>. D<sub>7</sub>-D<sub>0</sub>, MREQ, IORO, RD, and WR.
 Measurements made with outputs floating.

#### **CAPACITANCE**

Guaranteed by design and characterization.

Symbol	Parameter	Min	Max	Unit
C <sub>CLOCK</sub>	Clock Capacitance		35	pf
C <sub>IN</sub>	Input Capacitance	•	5	pf
C <sub>OUT</sub>	Output Capacitance		15	pf

NOTES:

T<sub>A</sub> = 25°C, f = 1 MHz.
Unmeasured pins returned to ground.

### AC CHARACTERISTICS<sup>†</sup> (Z8400/NMOS Z80 CPU)

			<b>Z08400</b> 04		<b>Z08</b> 4	0006	Z084	8000
Number	Symbol	Parameter	Min	Max	Min	Max	Min	Max
1	TcC	Clock Cycle Time	250*		162*		125*	
2	TwCh	Clock Pulse Width (High)	110	2000	<b>6</b> 5	2000	55	2000
3	TwCl	Clock Pulse Width (Low)	110	2000	65	2000	55	2000
4	TfC	Clock Fall Time		30		20		10
5	TrC	Clock Rise Time		30		20		10
6	TdCr(A)	Clock † to Address Valid Delay		110		90		80
7	TdA(MREQf)	Address Valid to MREQ ↓ Delay	65*		35*		20*	
8	TdCf(MREQf)	Clock I to MREQ I Delay		85		70		60
9	TdCr(MREQr)	Clock f to MREQ f Delay		85		70		60
10	TwMREQh	MREQ Pulse Width (High)	110**	Ħ	65**	Ħ	45*1	+
11	Twmreqi	MREQ Pulse Width (Low)	220*	Ħ	135*1	<del>i</del>	100*1	+
12	TdCf(MREQr)	Clock I to MREQ ↑ Delay		85		70		60
13	TdCf(RDf)	Clock I to RD I Delay		95		80		70
14	TdCr(RDr)	Clock † to RD † Delay		85		70		60
15	TsD(Cr)	Data Setup Time to Clock †	35		30		30	
16	ThD(RDr)	Data Hold Time to RD †		0		0		0
17	TsWAIT(Cf)	WAIT Setup Time to Clock ↓	70		60		50	
18	ThWAIT(Cf)	WAIT Hold Time after Clock ↓		0		0		0
19	TdCr(M1f)	Clock † to M1 ↓ Delay		100		80		70
20	TdCr(M1r)	Clock † to M1 † Delay		100		80		70
21	TdCr(RFSHf)	Clock ↑ to RFSH ↓ Delay		130		110		95
22	TdCr(RFSHr)	Clock to RFSH t Delay		120		100		85
23	TdCf(RDr)	Clock I to RD ↑ Delay		85		70		60
24	TdCr(RDf)	Clock † to RD ↓ Delay		85		70		60
25	TsD(Cf)	Data Setup to Clock ↓ during M <sub>2</sub> , M <sub>3</sub> , M <sub>4</sub> , or M <sub>5</sub> Cycles	50		40		30	
26	TdA(IORQf)	Address Stable prior to IORQ ↓	180*		110*		75*	
27	TdCr(IORQf)	Clock † to IORQ ↓ Delay		75		65		55
28	TdCf(IORQr)	Clock  to IORQ ↑ Delay		85		70		.60
29	TdD(WRf)	Data Stable prior to WR ↓	80*		25*		5*	
30	TdCf(WRf)	Clock ∮ to WR ∮ Delay		80		70		60
31	TwWR	WR Pulse Width	220*		135*		100*	
32	TdCf(WRr)	Clock ↓ to WR↑ Delay		80		70		60
· 33	TdD(WRf)	Data Stable prior to WR ↓	<b>−10</b> *		-55*		55*	
34	TdCr(WRf)	Clock † to WR ↓ Delay		65		60		55
35	TdWRr(D)	Data Stable from WR †	60*		30*		15*	
36	TdCf(HALT)	Clock ↓ to HALT ↑ or ↓		300		260		225
37	TwNMI	NMI Pulse Width	80		70		60*	
38	TsBUSREQ(Cr)	BUSREQ Setup Time to Clock †	50		50		40	

<sup>\*</sup>For clock periods other than the minimums shown, calculate parameters using the table on the following page. Calculated values above assumed TrC = TrC = 20 ns.
†Units in nanoseconds (ns).

<sup>#</sup> For loading  $\geq$  50 pf., Decrease width by 10 ns for each additional 50 pf.

### AC CHARACTERISTICS† (Z8400/NMOS Z80 CPU; Continued)

			<b>Z08</b> 4	Z0840006		Z0840008		
Number	Symbol	Parameter	Min	Max	Min	Max	Min	Max
39	ThBUSREQ(Cr)	BUSREQ Hold Time after Clock †	0		0		0	<del>-</del>
40	TdCr(BUSACKf)	Clock † to BUSACK ↓ Delay		100		90		80
41	TdCf(BUSACKr)	Clock   to BUSACK   Delay		100		90		80
42	TdCr(Dz)	Clock ↑ to Data Float Delay		90		80		70
43	TdCr(CTz)	Clock † to Control Outputs Float Delay (MREQ, IORQ, RD, and WR)		80		70		60
44	TdCr(Az)	Clock † to Address Float Delay		90		80		70
45	TdCTr(A)	MREQ t, IORQ t, RD t, and WR t to Address Hold Time	. 80*		35*		20*	
46	TsRESET(Cr)	RESET to Clock † Setup Time	60		60		45	
47	ThRESET(Cr)	RESET to Clock † Hold Time		0		0		0
48	TsINTf(Cr)	INT to Clock † Setup Time	80		70		55	
49	ThINTr(Cr)	INT to Clock † Hold Time		0		0		0
50	TdM1f(IORQf)	M1 ↓ to IORQ ↓ Delay	565*		365*		270*	
51	TdCf(lORQf)	Clock I to IORQ I Delay		85		70		60
52	TdCf(IORQr)	Clock † IORQ † Delay		85		70		60
53	TdCf(D)	Clock		150		130		115

<sup>\*</sup>For clock periods other than the minimums shown, calculate parameters using the following table. Calculated values above assumed TrC = TrC = 20 ns. †Units in nanoseconds (ns).

#### **FOOTNOTES TO AC CHARACTERISTICS**

Number	Symbol	General Parameter	Z0840004	Z0840006	Z084D008
1	TcC	TwCh + TwCl + TrC + TfC	714		
7	TdA(MREQf)	TwCh + TfC	- 65	-50	- 45
10	TwMREQh	TwCh + TfC	- 20	-20	-20
11	TwMREQI	TcC	- 30	-30	- 25
26	TdA(IORQf)	TcC	- 70	55	-50
29	TdD(WRf)	TcC	- 170	140·	- 120
31	TwWR	TcC	- 30	-30	- 25
33	TdD(WRf)	TwCl + TrC	- 140	- 140	120
35	TdWRr(D)	TwCl + TrC	- 70	- 55	50
45	TdCTr(A)	TwCl + TrC	- 50	- 50	45
50	TdM1f(IORQf)	2TcC + TwCh + TfC	- 65	- 50	45

$$V_{OH} = 1.5 V$$
  
 $V_{OL} = 1.5 V$   
FLOAT = ±0.5 V

AC Test Conditions: V<sub>IH</sub> = 2.0 V V<sub>IL</sub> = 0.8 V

 $V_{IHC} = V_{CC} - 0.6 V$   $V_{ILC} = 0.45 V$ 

# **Customer Support**

For answers to technical questions about the product, documentation, or any other issues with Zilog's offerings, please visit Zilog's Knowledge Base at <a href="http://www.zilog.com/kb">http://www.zilog.com/kb</a>.

For any comments, detail technical questions, or reporting problems, please visit Zilog's Technical Support at <a href="http://support.zilog.com">http://support.zilog.com</a>.

PS017801-0602 Customer Support

## **Mouser Electronics**

**Authorized Distributor** 

Click to View Pricing, Inventory, Delivery & Lifecycle Information:

### ZiLOG:

```
        Z84C0008VEC
        Z84C0008FEC
        Z84C0006FEC
        Z84C0020FEC
        Z84C0006VEG
        Z84C0006PEG
        Z84C0010VEG

        Z84C0010PEG
        Z84C0020PEG
        Z84C0006AEG
        Z84C0008AEG
        Z84C0010AEG
        Z84C0020AEG
        Z84C0020VEG

        Z84C0020VEG
        Z84C0010PEC
        Z84C0010VEC
        Z84C0010FEC
        Z84C0006PEC
        Z84C0006VEC

        Z84C0020VEC
        Z84C0010PEC
        Z84C0006FEG
        Z84C0010FEG
        Z84C0020PEC
        Z84C0008FEG
```