

Prolog

Need to learn

Prolog database of facts and rules

Inference engine of Prolog

- Backtrack search (Depth-First Search)

Logical variables

- Different from variables in most languages
- Can be assigned with any data type

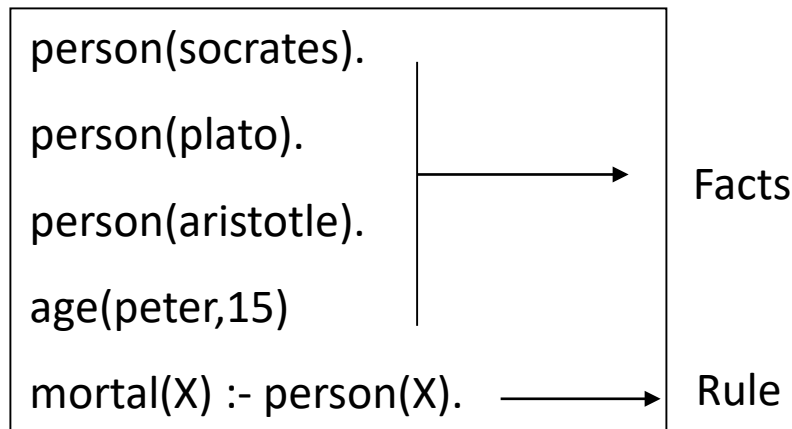
Unification

- Built-in pattern matcher

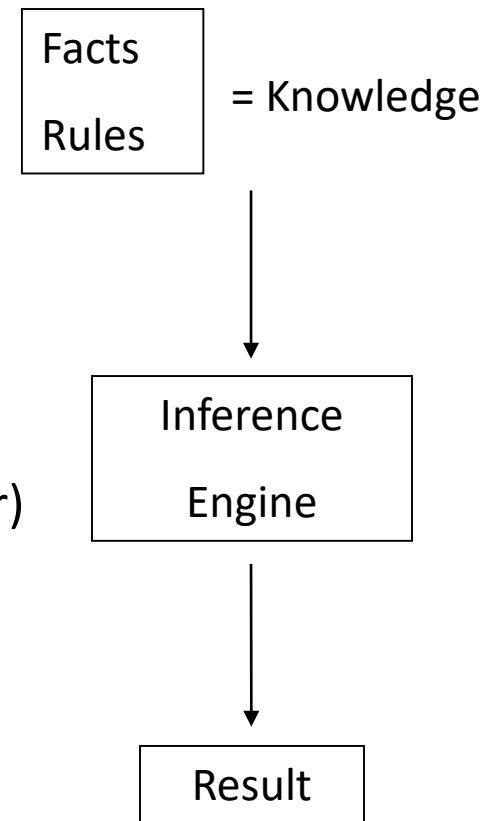
Prolog Program

Database

- A set of clauses (sentences)
 - Facts or rules
 - All are predicates
 - Small units, similar to subroutine call
 - Return truth values



Prolog =
(Interpreter)



Prolog

Prolog Listener

- Listen to queries from user

Query

- ?- mortal(X).
- Return X = socrates.

Arity

- Number of arguments in a predicate
- mortal(X): arity = 1
- parent(peter, mary): arity = 2
- Represented as mortal/1, parent/2

Predicates with different arity

- Different
- name/2 \neq name/3

```
person(socrates).  
person(plato).  
person(aristotle).  
mortal(X) :- person(X).
```

Syntax of Prolog

Syntax

`pred(arg1, arg2, ... argN).`

- `pred`: name of predicate
- `arg1... : arguments`
- `N`: arity
- `..`: syntactic end of Prolog clause
- E.g. `parent(peter, mary)` ← a wrong syntax

Predicate of arity 0

- `pred.`

Data types

Integer

Atom

- Text constant
- Begin with a **lowercase** letter

Variable

- Begins with an **uppercase** letter, or **underscore** (_)

Structure

- Complex terms, list

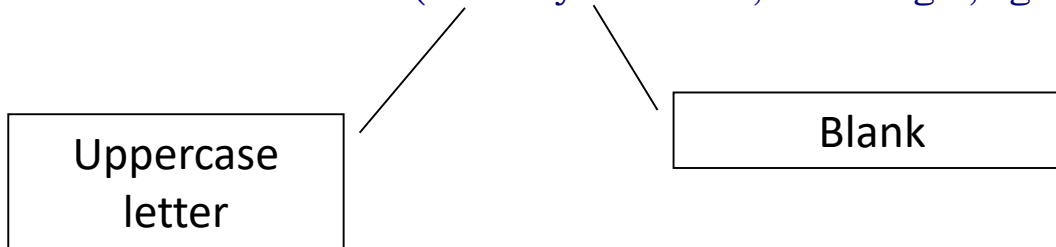
Facts

Data or information of Prolog program

E.g. customer/3

`customer('John Jones', boston, good_credit).`

`customer('Sally Smith', chicago, good_credit).`



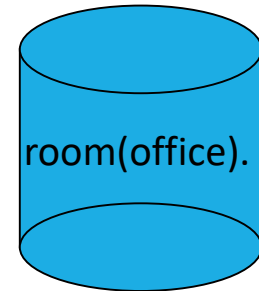
customer	Name	Place	Credit
1	'John Jones'	boston	good
2

Simple Queries

Prolog Queries

Work by pattern matching

- ?- room(X).
- Check if Query = facts or rules in DB



Query = goal

- If there is a fact that matches the goal
 - Query succeeds → Responds 'yes'
 - Otherwise → Responds 'no'

Unification

Process of pattern matching

- For computer
- $\text{room}(X) = \text{room}(\text{office})$.

Rules of unification

- Predicate name in the goal and the one in DB are the same
- Both predicates have same arity
 - Same number of arguments
- All of the arguments are the same
 - Variables can be instantiated
 - Constants cannot be instantiated

Bindings of Variables

When a logical variable is assigned

- A value /a structure /a term
- Logical variable is bound
- E.g. $A = 1$
 - A is bound to the value of 1
 - “Binding of A” = 1
- Variable once bound, cannot change value
 - E.g. ?- $A=1, A=2$.
 - Fail
 - Can unbind A by backtracking automatically

Bindings of Variables

```
?-room(X).  
X= kitchen
```

```
?-room(X), X = 1.  
fail
```

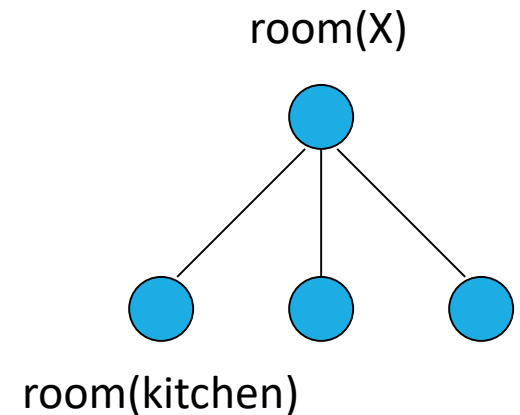
Unbind value by backtracking

- Use (;) to find alternatives
- no → no more answer

```
?-room(X).  
X= kitchen;  
X= office;  
X= hall;  
X= 'dinning room';  
X= cellar;  
no
```

Prolog program

```
room(kitchen).  
room(office).  
room(hall).  
room('dining room').  
room(cellar).
```



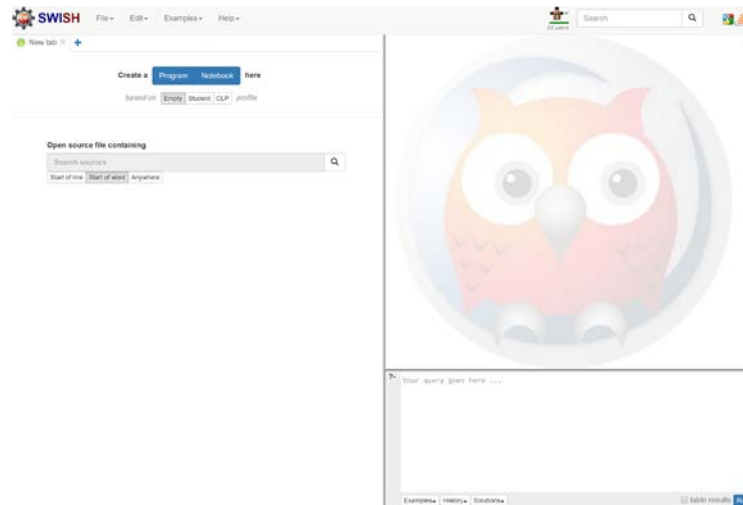
To use Prolog

Online

- <http://swish.swi-prolog.org/>

Download

- <http://www.swi-prolog.org/download/stable>



To use Prolog

Unix: type 'pl'.

Windows: double click the icon

Writing prolog program

- Consult the program
 - ?- consult('a.pl').
 - ?- consult(a).
- Ask your questions

If program is changed

- Consult the program again

Exercise 1

Family database


- Create a Prolog file called “family.pl”
- Add members of family
 - male/1, female/1
 - male(dennis). female(diana).
- Add predicate that records parent-child relationship
 - parent(diana, dennis).

Queries

Consult program “family.pl”

Ask queries

- ?- male(dennis).
- ?- male(X).
- ?- parent(X, dennis).
- ?- parent(dennis, X).
- ?- parent(diana, _).
- ?- parent(X, Y).



Don't care

Exercise 2

Family database

- Write queries
 - Confirm a parent relationship
 - Find someone's parent
 - Find someone's children
 - List all parent-children
 - Semi-colon “;”
 - Predicate “fail”
- Compound queries to find family relationships
 - Father, mother, sons, daughters, grandmothers, grandfathers

Prolog Rules

Prolog Rules

Stored query

Syntax

- head :- body.
- where_food(X,Y) :- location(X,Y), edible(X).

head

- Predicate definition (like a fact)

:-

- Neck symbol, read as "if"

body

- One or more goals
- Simple or compound query

Example

?- where_food(X, kitchen).

- X = apple ;
- X = crackers ;
- no

?- where_food(Thing, 'dining room').

- no

?- where_food(apple, kitchen).

- yes

```
location(desk, office).  
location(apple, kitchen).  
location	flashlight, desk).  
location('washing machine', cellar).  
location(nani, 'washing machine').  
location(broccoli, kitchen).  
location(crackers, kitchen).  
location(computer, office).  
  
edible(apple).  
edible(crackers).  
  
tastes_yucky(broccoli).  
  
here(kitchen).
```

Multiple Rules

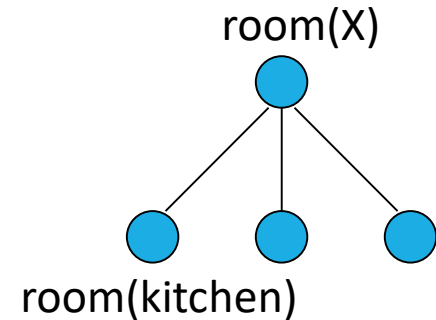
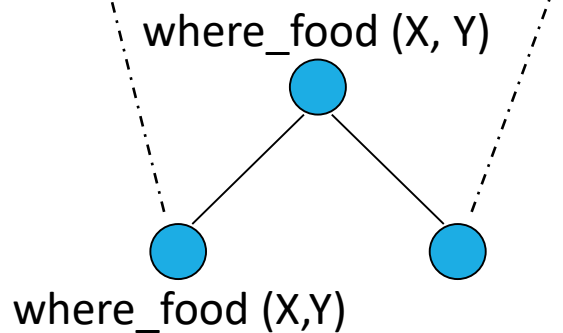
Have multiple facts

Have multiple rules

- Defining a predicate

For example

- `where_food(X,Y) :- location(X,Y), edible(X).`
- `where_food(X,Y) :- location(X,Y), tastes_yucky(X).`



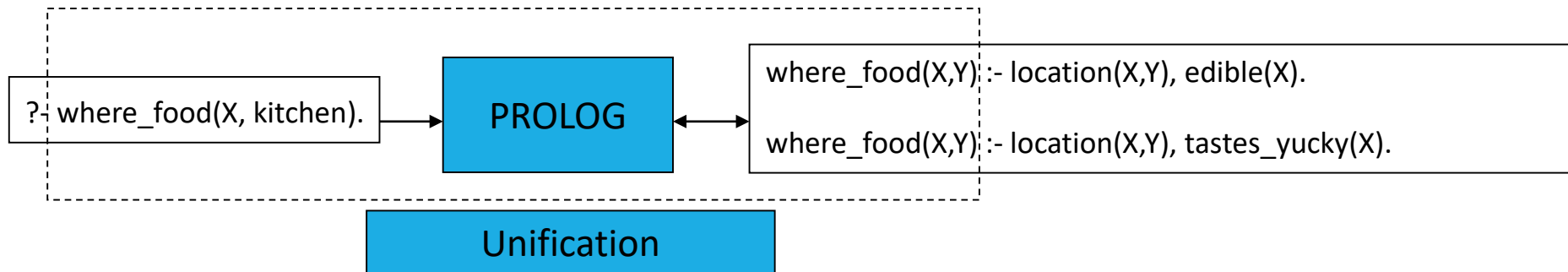
```
?- where_food(X, kitchen).  
X = apple ;  
X = crackers ;  
X = broccoli ;  
no
```

How Rules Work?

Unification

- Unify goal pattern with head of the rule
- Succeeds
 - Initiate a new query
 - With goals in body of the rule

```
location(desk, office).  
location(apple, kitchen).  
location	flashlight, desk).  
location('washing machine', cellar).  
location(nani, 'washing machine').  
location(broccoli, kitchen).  
location(crackers, kitchen).  
location(computer, office).  
  
edible(apple).  
edible(crackers).  
  
tastes_yucky(broccoli).  
  
here(kitchen).
```



Using Rules

Solve the problem of door/2

- Define a two-way predicate

`connect(X,Y):- door(X,Y).`

`connect(X,Y):- door(Y,X).`

Means
'OR'

`door(office, kitchen).`

`?- connect(kitchen, office).`

yes

`?- connect(office, kitchen).`

yes

Exercises

Family database

- Build a rule for siblings
 - `siblings(X,Y) :- parent(P,X), parent(P,Y).`
 - Allow an individual as his/her own sibling
- Fix the problem
 - With a built-in predicate “Not Equal” \neq
 - `siblings(X,Y) :- parent(P,X), parent(P,Y), X \neq Y.`
- Use sibling to build rules
 - Brothers, sisters, uncles, aunts
 - `brothers(X,Y) :- siblings(X,Y), male(X).`
 - `sisters(X,Y) :- siblings(X,Y), female(X).`
 - `uncles(X,Y) :- parent(P,Y), brothers(X,P).`
 - `aunts(X,Y) :- parent(P,Y), sisters(X,P).`

Exercises

Define married/2 by using facts spouse/2

- married(X,Y) :- spouse(X,Y).
- married(X,Y) :- spouse(Y,X).

Define uncles & aunts using married/2 further

- uncles(X,Y):- parent(P,Y), brothers(X,P).
- uncles(X,Y):- aunts(A,Y), married(X,A).
- aunts(X,Y):- parent(P,Y), sisters(X,P).
- aunts(X,Y):- uncles(U,Y), married(X,U).

Arithmetic & Logic

Arithmetic

Evaluation of arithmetic expression

- Built-in predicate: 'is'
- X is <arithmetic expression>
- E.g. $\text{Tmp} = X$, X is $\text{Tmp} + 1$.

Variable X

- Set to value of the arithmetic expression

Examples

Use Prolog as calculator

- ?- X is 2 + 2.
- X = 4
- ?- X is 3 * 4 + 2.
- X = 14

Parentheses clarify precedence

- ?- X is 3 * (4 + 2).
- X = 18
- ?- X is (8 / 4) / 2.
- X = 1

Comparison

Prolog provides a number of operators

- $X > Y$
- $X < Y$
- $X \geq Y$
- $X \leq Y$

?- X is 2 + 2, X > 3.

X = 4

?- X is 2 + 2, 3 >= X.

no

?- 3+4 > 3*2.

yes

Recursion

Recursion

Ability for a unit of code to call itself

- Repeatedly if necessary
- Prolog
 - A predicate contains a goal that refers to itself

At least two parts

- Boundary / termination condition
- Recursive case

Recursive Definition

Boundary condition

- Simple case that we know to be true

Recursive case

- Simplify the problem
 - First remove a layer of complexity
 - Then call itself
- At each level
 - Check boundary condition
 - If it is reached, recursion ends
 - Otherwise, recursion continues

Example

Flashlight is in desk, and desk is in office.

?- location flashlight, office).

no

is_contained_in/2

- Dig through layers of nested things

Boundary condition

- T1 is directly located in T2
- is_contained_in(T1,T2) :- location(T1,T2).

Simplify & recur

- T1 is contained in **X**, some intermediate thing, which is located in T2
- is_contained_in(T1,T2) :- location(X,T2), is_contained_in(T1,X).

```
location(desk, office).  
location(apple, kitchen).  
location	flashlight, desk).  
location('washing machine', cellar).  
location(nani, 'washing machine').  
location(broccoli, kitchen).  
location(crackers, kitchen).  
location(computer, office).
```

Exercises

Family database

- Use recursion to write ancestor/2

```
ancestor(X,Y) :- parent(X,Y).
ancestor(X,Y) :- parent(P,Y), ancestor(X, P).
```
- Use ancestor/2
 - Find all of a person's ancestors

```
all_ancestor(X) :- ancestor(Y, X), write(Y), nl, fail.
```
 - Find all of a person's descendants

```
all_descendent(X) :- ancestor(X, Y), write(Y), nl, fail.
```

Unification

Unification

Built-in pattern-matching algorithm

- Make two items identical

Unification process

- Variable & any term

$$A = abc$$

$$A = f(a,b)$$

- Constant & constant

$$a = a$$

$$abc = abc$$

- Structure & structure

$$f(a,g(b,c)) = f(a,X)$$

Explicit Unification “=

Built-in predicate =/2

- Succeed if two arguments unify
- $\text{arg1} = \text{arg2}$
- =(arg1, arg2)

Warning

- Do not cause arithmetic evaluation
 - Evaluation is done by “is/2”

Example without Variable

?- a = a.

yes

?- a = b.

no

?- location(apple, kitchen) = location(apple, kitchen).

yes

?- location(apple, kitchen) = location(pear, kitchen).

no

?- a(b,c(d,e(f,g))) = a(b,c(d,e(f,g))).

yes

?- a(b,c(d,e(f,g))) = a(b,c(d,e(g,f))).

no

Example for a Variable & a Constant

?- X = a.

X = a

?- 4 = Y.

Y = 4

?- location(apple, kitchen) = location(apple, X).

X = kitchen

?- location(X,Y) = location(apple, kitchen).

X = apple

Y = kitchen

?- location(apple, X) = location(Y, kitchen).

X = kitchen

Y = apple

Example for Variables

?- X = Y.

X = _01

Y = _01

?- location(X, kitchen) = location(Y, kitchen).

X = _01

Y = _01

Bound Variables

?- X = Y, Y = hello.

X = hello

Y = hello

?- X = Y, a(Z) = a(Y), X =
hello.

X = hello

Y = hello

Z = hello

?- X = Y, Y = 3, write(X).

3

X = 3

Y = 3

?- X = Y, tastes_yucky(X),
write(Y).

broccoli

X = broccoli

Y = broccoli

Structures with Variables

?- X = a(b,c).

X = a(b,c)

?- a(b,X) = a(b,c(d,e)).

X = c(d,e)

?- a(b,X) = a(b,c(Y,e)).

X = c(_01,e)

Y = _01

?- a(b,X) = a(b,c(Y,e)), Y = hello.

X = c(hello, e)

Y = hello

?- food(X,Y) = Z, write(Z), nl,
tastes_yucky(X), edible(Y),
write(Z).

food(_01,_02)

food(broccoli, apple)

X = broccoli

Y = apple

Z = food(broccoli, apple)

Important to Note

If a value unified to a variable in later goal

- Conflicts with pattern set earlier, unification fails

?- $a(b,X) = a(b,c(Y,e))$, $X = \text{hello}$.

no

Second goal fails

- No value of X
- Allow hello to unify with $c(Y,e)$

?- $a(b,X) = a(b,c(Y,e))$, $X = c(\text{hello}, e)$.

$X = c(\text{hello}, e)$

$Y = \text{hello}$

Succeed

- Same structure can be unified

Important to Note

?- $a(X) = a(b,c)$.

no

?- $a(b,c,d) = a(X,X,d)$.

no

Fail

- Pattern asks that the first two arguments be the same

?- $a(c,X,X) = a(Y,Y,b)$.

no

No value of X and Y allow the two structures to unify

- $c = Y$
- $X = Y$
- $X = b$
- $\rightarrow c = b?$

Important to Note

Anonymous variable (`_`)

- Wildcard variable
- Do not bind to values
- Multiple occurrences do not imply equal values

```
?- a(c,X,X) = a(_,_ ,b).  
X = b
```

Implicit unification

- Prolog searches for the head of a clause
 - Match a goal pattern
 - `?- food(Z, kitchen).`

```
food(X, Y):- location(X, Y),  
edible(X).
```

Exercises

Predict results of unification

?- $a(b,c) = a(X,Y)$.

- $X = b, Y = c$

?- $a(X,c(d,X)) = a(2,c(d,Y))$.

- $X = 2, Y = 2$

?- $a(X,Y) = a(b(c,Y),Z)$.

- $X = b(c, _01), Y = _01, Z = _01$

?- $\text{tree}(\text{left}, \text{root}, \text{Right}) = \text{tree}(\text{left}, \text{root}, \text{tree}(a, b, \text{tree}(c, d, e)))$.

- $\text{Right} = \text{tree}(a, b, \text{tree}(c, d, e))$

List

List

Powerful data structure

- Hold and manipulate groups of things

List in Prolog

- Collection of terms
- Atoms, integers, structures, lists

Represented as

- [apple, broccoli, refrigerator]

List

Comparison

- Without list
 - `location(apple, kitchen).`
 - `location(broccoli, kitchen).`
 - `location(crackers, kitchen).`
- With list
 - `loc_list([apple, broccoli, crackers], kitchen).`

Special list

- Empty list
 - Nil represented as `[]`
 - `loc_list([], hall).`

List

Work as usual terms

loc_list([apple, broccoli, crackers], kitchen).

?- loc_list(X, kitchen).

X = [apple, broccoli, crackers]

?- [_,X,_] = [apples, broccoli, crackers].

X = broccoli

Have to know

- Number of items in list
- Order of items

Special Notation

[H | T]

- Allow reference to
 - First element of the list
 - List of remaining elements

H

- Head
- Bound to the first element of the list

T

- Tail
- Bound to list of remaining elements

Unification Using Lists

?- [a|[b,c,d]] = [a,b,c,d].

yes

Succeed

- 'a' is a head (an atom)
- [b,c,d] is the tail (a list)

?- [a|b,c,d] = [a,b,c,d].

no

Fail

- The tail is not a list

Examples

?- [H|T] = [apple, broccoli, refrigerator].

H = apple

T = [broccoli, refrigerator]

?- [H|T] = [a, b, c, d, e].

H = a

T = [b, c, d, e]

?- [H|T] = [apples, bananas].

H = apples

T = [bananas]

?- [H|T] = [a, [b,c,d]].

H = a

T = [[b, c, d]]

?- [H|T] = [apples].

H = apples

T = []

Tail is empty list

?- [H|T] = [].

no

Do not unify

- Empty list has no head

Multiple Heads

Before the bar (|)

- Can specify more than one element

?- [One, Two | T] = [apple,
sprouts, fridge, milk].

One = apple

Two = sprouts

T = [fridge, milk]

?- [X,Y|T] = [a|Z].

X = a

Y = _01

T = _03

Z = [_01 | _03]

?- [H|T] = [apple, Z].

H = apple

T = [_01]

Z = _01

List Predicate

member/2

- Check membership of a term

```
member(H,[H|T]).
```

```
member(X,[H|T]) :- member(X,T).
```

```
?- member(banana, [apple, broccoli, crackers]).
```

```
no
```

```
?- member(X, [apple, broccoli, crackers]).
```

```
X = apple ;
```

```
X = broccoli ;
```

```
X = crackers ;
```

```
no
```


append/3

Recursive Definition

Build lists from other lists

Split lists into separate pieces

?- append([a,b,c],[d,e,f],X).

X = [a,b,c,d,e,f]

Boundary condition

- $[] + X = X$

append([],X,X).

append([H|T1],X,[H|T2]) :- append(T1,X,T2).

Recursive case

- $[H|T1] + X = [H|T2]$
- $T2 = T1 + X$

$[1|[]] + [a,b] = [1|T2]$

$T2 = [] + [a,b] = [a,b]$

Exercises

Remove a given element from a list

`remove(H, [], []).`

`remove(H, [H|T], T).`

`remove(X, [H|T], [H|T2]) :- remove(X, T, T2).`

Get last element of a list

`last([E], E).`

`last([H|T], E) :- last(T, E).`

Count elements in a list

`len([], 0).`

`len([H|T], Count) :- len(T, NewCount), Count is 1 + NewCount.`

Exercises

?- [] = [H|T].

no

?- [a] = [H|T].

H = a, T = []

?- [apple,3,X,'What?'] = [A,B|Z].

A = apple, B = 3, Z = [_01, 'What?']

?- [[a,b,c],[d,e,f],[g,h,i]] = [H|T].

H = [a,b,c], T = [[d,e,f], [g,h,i]]

Control Structures

Tail Recursion

Traditional recursion

- $N! = (N - 1)! \cdot N$
- Reduce problem into smaller one
- Problem is not solved until $(N-1)!$ is known
- Stack is necessary to keep information

Iteration

- $N! = N \cdot (N - 1)!$
- No stack is necessary

Tail recursion

- Combine the advantages

Example

factorial_1(1,1).

factorial_1(N,F):-

N > 1,

NN is N - 1,

factorial_1(NN,FF),

F is N * FF.

factorial_2(1,F,F).

factorial_2(N,T,F)÷

N > 1,

TT is N * T,

NN is N - 1,

factorial_2(NN,TT,F).

Need a stack

Handle large amount of value
simultaneously