

COMP412 Computer Security

Lec 04 Block Cipher Operation

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2021/2022

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Modes of Operation

Electronic Code Book

Cipher Block Chaining Mode

Cipher Feedback Mode

Output Feedback Mode

Counter Mode

Feedback Characteristics of Modes

XTS-AES

Modes of Operation

- › Block cipher: operates on fixed length **b -bit input** to produce **b -bit ciphertext**
- › What about encrypting plaintext longer than b bits?
 - › Break plaintext into b -bit blocks (padding if necessary) and apply cipher on each block
- › Security issues arise: different **modes of operation** have been developed

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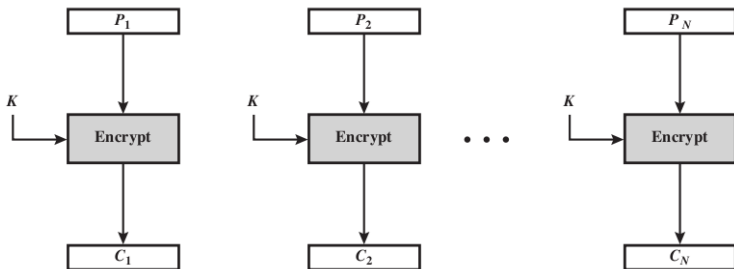
Output Feedback Mode

Counter Mode

Feedback Characteristics of Modes

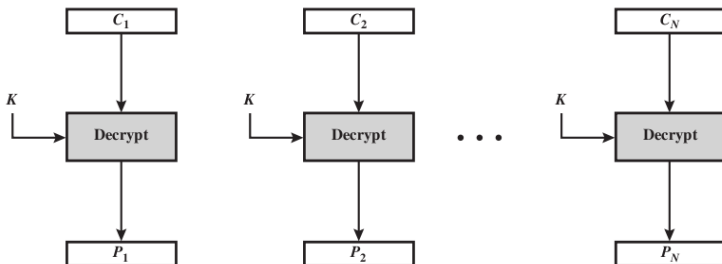
XTS-AES

Electronic Code Book (ECB) Encryption



- Each block: 64 bits
- Same key in each block
- Independently

Electronic Code Book (ECB) Decryption



Modes

ECB

CBC

CFB

OFB

CTR

Feedback

XTS-AES

Summary

- › Each block of 64 plaintext bits is encoded independently using same key
- › Typical applications: secure transmission of single values (e.g. encryption key)
- › Problem: with long message, repetition in plaintext may cause repetition in ciphertext

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Cipher Block Chaining Mode (CBC) Encryption

Modes

ECB

CBC

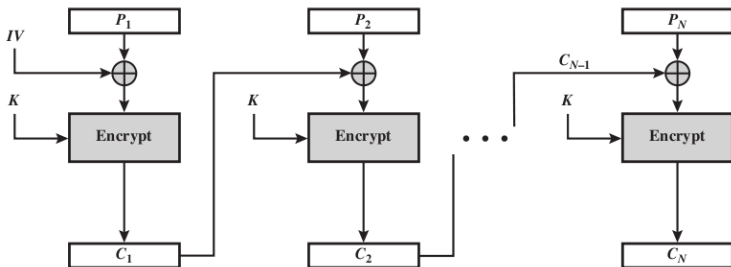
CFB

OFB

CTR

Feedback

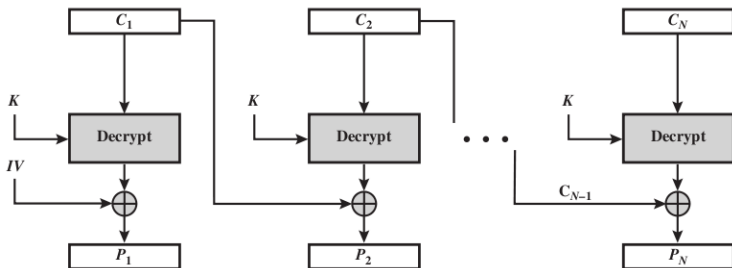
XTS-AES



- Each block: 64 bits
- Same key in each block
- Initialization Vector (IV) necessarily
- IV XORed P_1 , C_i XORed P_{i+1}

Question: *How to avoid the repetition in ciphertext for repetition in plaintext?*

Cipher Block Chaining Mode (CBC) Decryption



Modes

ECB

CBC

CFB

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XTS-AES

CBC Summary

- › Input to encryption algorithm is **XOR** of next 64-bits plaintext and preceding 64-bits ciphertext
- › Typical applications: General-purpose **block-oriented** transmission; authentication
- › Initialisation Vector (IV) must be known by sender/receiver, but secret from attacker
 - In particular, it must be impossible to predict the IV for any given plaintext;
 - For maximum security, IV should be protected against unauthorized changes.
 - E.g., send the IV using ECB encryption.

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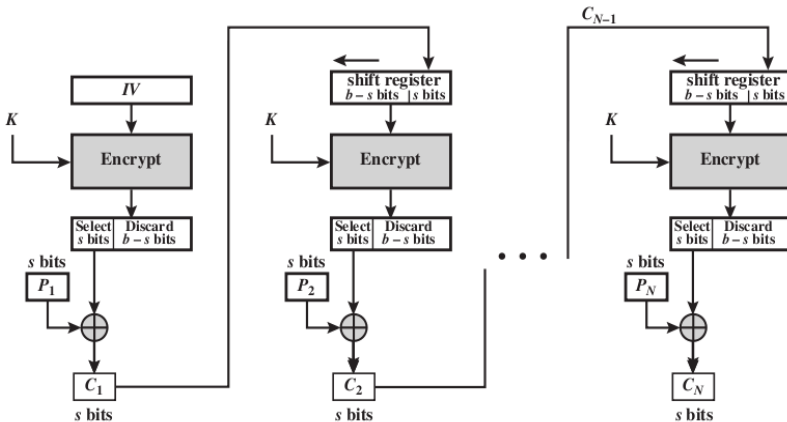
Output Feedback Mode

Counter Mode

Feedback Characteristics of Modes

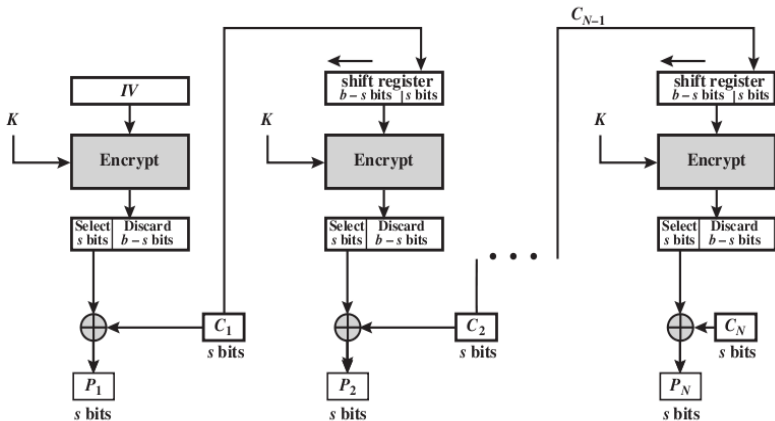
XTS-AES

Cipher Feedback Mode (CFB) Encryption



- Totally: **b** bits
- Each time process **s** bits
- Shift register applied each time
- Initialization Vector (IV) necessarily

Cipher Feedback Mode (CFB) Decryption



CFB Summary

- › Converts **block cipher** into **stream cipher**
 - › No need to pad message to integral number of blocks
 - › Operate in real-time: each character encrypted and transmitted immediately
- › Input processed **s** bits at a time
- › **Preceding ciphertext** used as input to cipher to produce pseudo-random output
- › XOR **output** with **plaintext** to produce ciphertext
- › Typical applications: General-purpose **stream-oriented** transmission; authentication

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Output Feedback Mode (OFB) Encryption

Modes

ECB

CBC

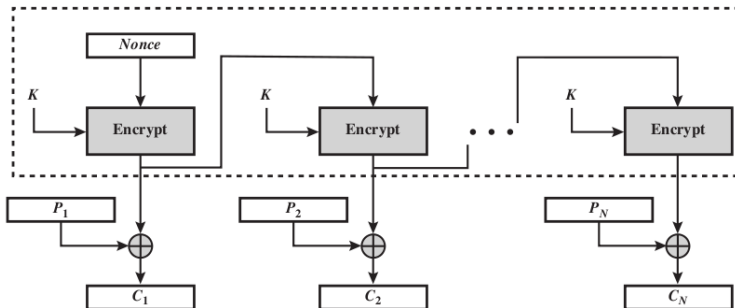
CFB

OFB

CTR

Feedback

XTS-AES



- Initialization Vector (IV) necessarily
 - IV must be a nonce,
 - must be unique to each execution of the encryption operation
 - Because each encryption output depends only on the key and the IV.

Output Feedback Mode (OFB) Decryption

Modes

ECB

CBC

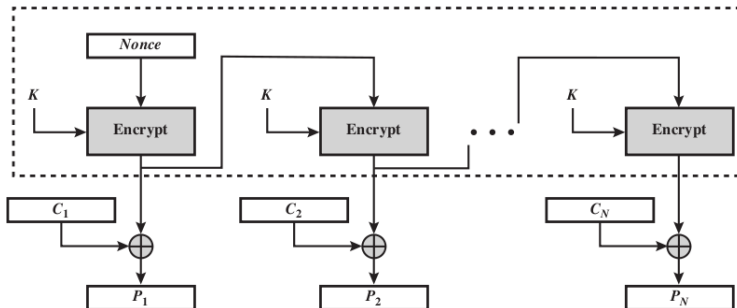
CFB

OFB

CTR

Feedback

XTS-AES



OFB Summary

- › Converts **block cipher** into **stream cipher**
 - OFB has structure of a typical stream cipher;
 - Distinction from the stream cipher is OFB encrypts a full block at a time; while many stream ciphers encrypt one byte at a time.
- › Similar to CFB, except input to encryption algorithm is **preceding encryption output**
- › Typical applications: **stream-oriented** transmission over **noisy** channels (e.g. satellite communications)
- › Advantage compared to CFB: bit errors do not propagate
- › Disadvantage: more vulnerable to message stream modification attack

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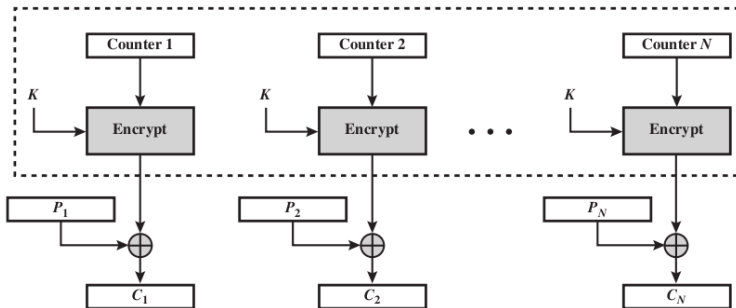
Output Feedback Mode

Counter Mode (CTR)

Feedback Characteristics of Modes

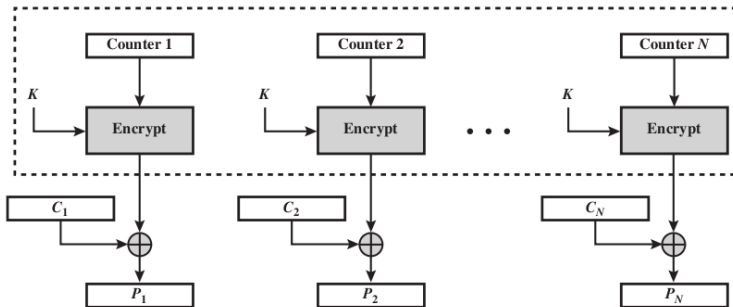
XTS-AES

Counter Mode (CTR) Encryption



- Initial Counter value must be a **nonce**
- All the counter values across all the messages should be **unique**.
- Same key in each block

Counter Mode (CTR) Decryption



Modes

ECB

CBC

CFB

OFB

CTR

Feedback

XTS-AES

CTR Summary

- › Converts **block cipher** into **stream cipher**
- › Each block of plaintext XORed with encrypted counter
- › Typical applications: General-purpose **block-oriented** transmission; useful for high speed requirements
- › Efficient hardware and software implementations
- › Simple and secure

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Feedback: CBC and CFB

Modes

ECB

CBC

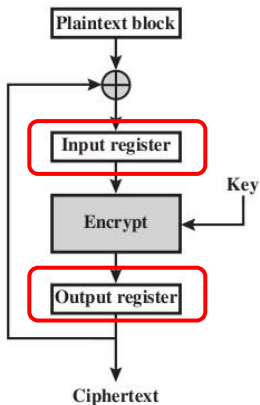
CFB

OFB

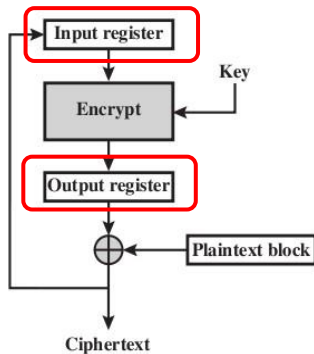
CTR

Feedback

XTS-AES



(a) Cipher block chaining (CBC) mode



(b) Cipher feedback (CFB) mode

Feedback: OFB and CTR

Block Cipher
Operation

Modes

ECB

CBC

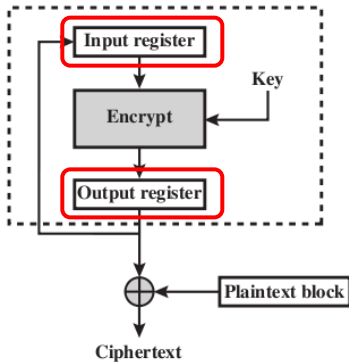
CFB

OFB

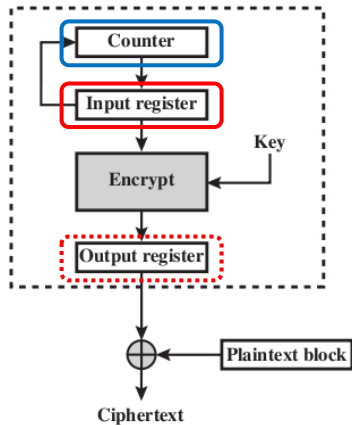
CTR

Feedback

XTS-AES



(c) Output feedback (OFB) mode



(d) Counter (CTR) mode

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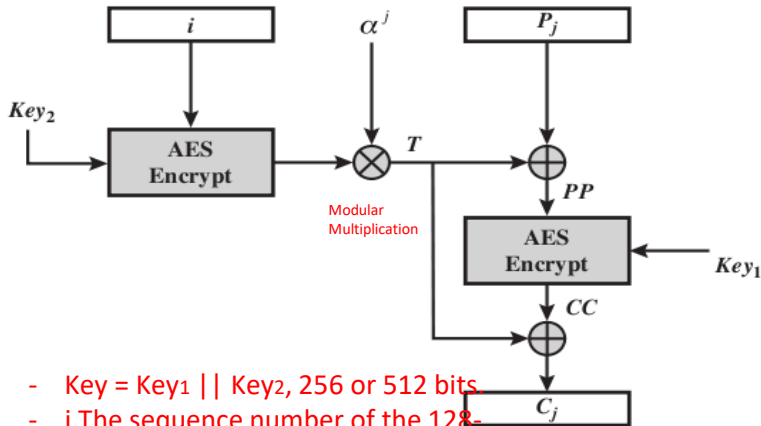
Output Feedback Mode

Counter Mode

Feedback Characteristics of Modes

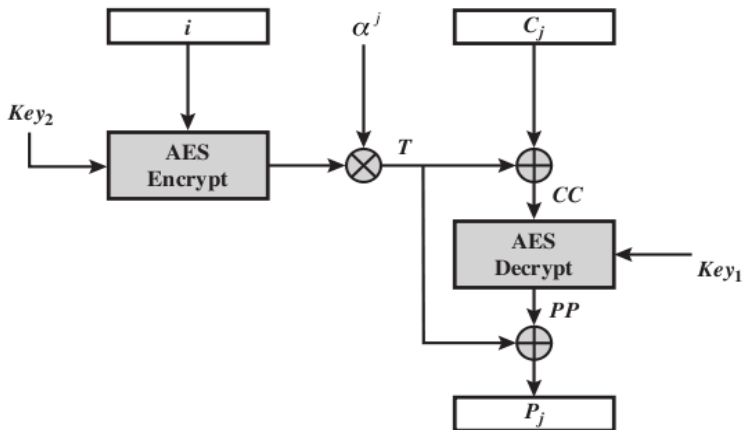
XTS-AES

XTS-AES Encryption of Single Block



- Key = $Key_1 || Key_2$, 256 or 512 bits.
- j The sequence number of the 128-bit block.
- P_j The j th block of plaintext
- i The value of the 128-bit tweak; a nonnegative integer.

XTS-AES Decryption of Single Block



Prove: Decrypted P = Plaintext P?

XTS-AES Encryption

Modes

ECB

CBC

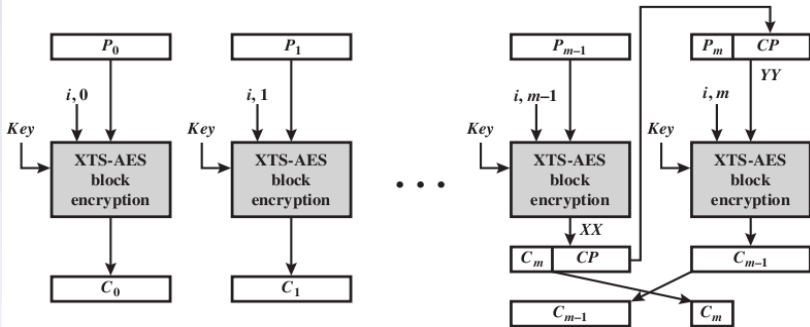
CFB

OFB

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Feedback

XTS-AES



XTS-AES Decryption

Modes

ECB

CBC

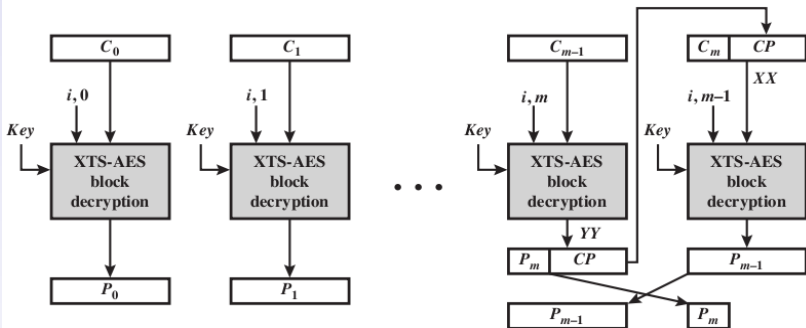
CFB

OFB

CTR

Feedback

XTS-AES



Encryption for Stored Data

- › XTS-AES designed for encrypting **stored data** (as opposed to **transmitted data**)
- › The P1619 standard was designed to specify the requirement for encrypting stored data.
 - “data at rest” differ somewhat from those for transmitted data.

Storage Encryption Requirements

- › The **ciphertext** is freely available for an attacker.
- › The data layout is not changed on the storage medium and in transit. The encrypted data must be the same size as the plaintext data.
- › **Data** are assessed in fixed sized blocks, independently from each other.
- › Encryption is performed in 16-byte blocks, independently from other blocks.
- › There are no other metadata used, except the location of the data blocks within the whole data set.
- › The same plaintext is encrypted to different ciphertexts at different locations, but always to the same ciphertext when written to the same location again.
- › A standard conformant device can be constructed for decryption of data encrypted by another standard conformant device.