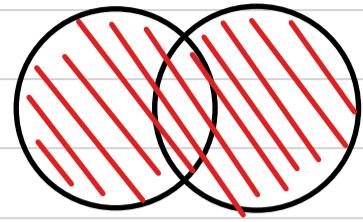
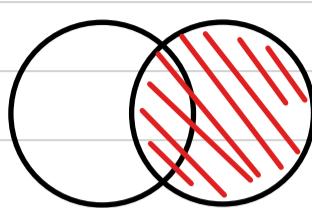
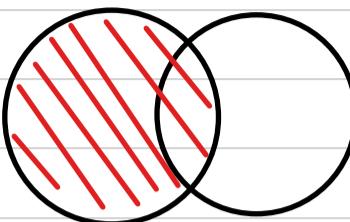
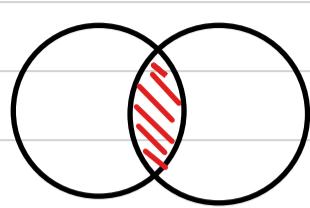


# Week 8

- o Inner join      left join      Right join      Full join



- o Multiple tables :

```
FROM actors, roles  
WHERE actors.id = roles.actor_id;
```

```
FROM actors AS a, roles AS r  
WHERE a.id = r.actor_id;
```

```
FROM actors AS a,  
movies AS m,  
roles AS r  
WHERE a.id = r.actor_id  
AND m.id = r.movie_id;
```

- o Mending pakai JOIN (Innerjoin = Join) :

```
FROM actors AS a  
JOIN roles AS r ON (a.id = r.actor_id);
```

```
FROM actors AS a  
JOIN roles AS r  
ON (a.id = r.actor_id)  
JOIN movies AS m  
ON (m.id = r.movie_id);
```

```
FROM actors_copy AS a  
NATURAL JOIN roles AS r  
NATURAL JOIN movies_copy AS m;
```

```
FROM actors AS a  
LEFT JOIN roles AS r  
ON (a.id = r.actor_id);
```

ON VS USING

```
FROM actors_copy AS a  
JOIN roles AS r USING (actor_id)  
JOIN movies_copy AS m USING (movie_id);
```

Muncul 1 kolom sama

Natural join gabung 2 table berdasarkan attribut/kolom nama yang sama & datatype. Kalau inner join gabung 2 table dispecified di ON clause.

Muncul 2/lebih

mysql> SELECT CONCAT(first_name, ' ', last_name) AS Name, r.role AS Role --> FROM actors AS a --> LEFT JOIN roles AS r ON (a.id = r.actor_id);	
Actor	Role
Nicholas Cage	Balthazar Blake
Nicholas Cage	Benjamin Franklin Gates
Dwayne Johnson	Jackson Bane
Dwayne Johnson	Derek Thompson
Dwayne Johnson	Hank Parsons
Diane Kruger	Abigail Chase
Jay Baruchel	David Stutler
Josh Hutcherson	Peeta Mellark
Josh Hutcherson	Sean Anderson
Teresa Palmer	Rebecca Barnes
Alfredo Quindug	Cato
Alexander Ludwig	Seth
Vanessa Hudgens	Kalliani
Ryan Sheckler	Mick Donnelly
Jennifer Lawrence	Katniss Everdeen
Ben Stiller	Larry Daley
AnnaSophia Robb	Sara
Dane Cook	Frankenbot
Dylan O'Brien	NULL
Thomas Brodie-Sangster	NULL

20 rows in set (0.00 sec)

o Ada juga CROSS JOIN yang memakai konsep perkalian Cartesius/matriks

o FULL JOIN = LEFT + RIGHT JOIN

```

④ FROM actors AS a
  JOIN roles AS r
  ON (a.id = r.actor_id)
  RIGHT JOIN movies AS m
  ON (m.id = r.movie_id);

```

Actor	Movie	Role
Nicholas Cage	The Sorcerer's Apprentice	Balthazar Blake
Jay Baruchel	The Sorcerer's Apprentice	David Stutler
Teresa Palmer	The Sorcerer's Apprentice	Rebecca Barnes
Josh Hutcherson	The Hunger Games	Peeta Mellark
Alexander Ludwig	The Hunger Games	Cato
Jennifer Lawrence	The Hunger Games	Katniss Everdeen
Nicholas Cage	National Treasure	Benjamin Franklin Gates
Diane Kruger	National Treasure	Abigail Chase
Dwayne Johnson	Race to Witch Mountain	Jack Bruno
Alexander Ludwig	Race to Witch Mountain	Seth
AnnaSophia Robb	Race to Witch Mountain	Sara
Dwayne Johnson	Tooth Fairy	Derek Thompson
Ryan Sheckler	Tooth Fairy	Mick Donnelly
Dwayne Johnson	Journey 2: The Mysterious Island	Hank Parsons
Josh Hutcherson	Journey 2: The Mysterious Island	Sean Anderson
Vanessa Hudgens	Journey 2: The Mysterious Island	Kailani
Ben Stiller	Night at the Museum: Secret of the Tomb	Larry Daley
Dan Stevens	Night at the Museum: Secret of the Tomb	Sir Lancelot
NULL	Project Almanac	NULL

- Gabung baris dalam kolom yang sama memakai UNION / ALL

```

④ SELECT first_name, last_name,
      gender, birth_date
    FROM actors
   WHERE gender = 'F'
UNION / ALL
SELECT first_name, last_name,
      gender, birth_date
    FROM actors
   WHERE birth_date > '1992-01-01';

```

```

mysql> SELECT first_name, last_name, gender, birth_date
-> FROM actors
-> WHERE gender = 'F'
-> UNION
-> SELECT first_name, last_name, gender, birth_date
-> FROM actors
-> WHERE birth_date > '1992-01-01';
+-----+-----+-----+-----+
| first_name | last_name | gender | birth_date |
+-----+-----+-----+-----+
| Diane      | Kruger    | F     | 1976-07-15 |
| Teresa     | Palmer    | F     | 1986-02-26 |
| Vanessa    | Hudgens   | F     | 1988-12-14 |
| Jennifer   | Lawrence  | F     | 1990-08-15 |
| AnnaSophia | Robb      | F     | 1993-12-08 |
| Josh       | Hutcherson| M     | 1992-10-12 |
| Alexander  | Ludwig    | M     | 1992-05-07 |
+-----+-----+-----+-----+
7 rows in set (0.00 sec)

mysql>

```

④ UNION = Yang sama ambil 1 saja  
 UNION ALL = Yang sama ambil semua yang ada

- LEFT JOIN = LEFT OUTER JOIN
- RIGHT JOIN = RIGHT OUTER JOIN

# Week 9

- o Multi-user kalau mau akses database itu secara concurrently
- o **Transactions** = An executing program that form logical unit of database processing.
- o Proses transaksi tentu memakan banyak akses database operations seperti insertion, deletion, modification, or retrieval operations.
- o Jika operasi database tidak update dan hanya retrieve di transaksi, maka dinamakan **read-only transaction**. Otherwise namanya **read-write transaction**
- o Basic operation :
  - o **read-item ( $X$ )** = Reads a database item  $X$  into a program variable
  - o **write-item ( $X$ )** = Writes the value of program variable  $X$  into database  $X$

o

$T_1$	$T_2$
<code>read_item(<math>X</math>); <math>X := X - N;</math></code>	
<code>write_item(<math>X</math>); read_item(<math>Y</math>);</code>	<code>read_item(<math>X</math>); <math>X := X + M;</math></code>
<code><math>Y := Y + N;</math> write_item(<math>Y</math>);</code>	<code>write_item(<math>X</math>);</code>

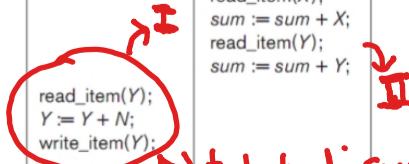
→  **$T_1$  masalah karena  $X$  overwrite**  
**Needs concurrency control**

o

$T_1$	$T_2$
<code>read_item(<math>X</math>); <math>X := X - N;</math> write_item(<math>X</math>);</code>	<code>read_item(<math>X</math>); <math>X := X + M;</math> write_item(<math>X</math>);</code>
<code>read_item(<math>Y</math>);</code>	

→ **Kalau  $T_1$  rusak maka harus kembali ke nilai awal tetapi  $T_2$  memegang nilai yang temporary, sehingga pembacaan selanjutnya akan salah**  
**Needs concurrency control**

o

$T_1$	$T_3$
<code>read_item(<math>X</math>); <math>X := X - N;</math> write_item(<math>X</math>);</code>  <b>I</b> <b>II</b>	<code>sum := 0; read_item(<math>A</math>); sum := sum + <math>A</math>;</code> $\vdots$ <code>read_item(<math>X</math>); sum := sum + <math>X</math>;</code> <code>read_item(<math>Y</math>); sum := sum + <math>Y</math>;</code>

→  **$T_3$  reads  $X$  after  $N$  subtracted and reads  $Y$  before  $N$  is added sehingga kalkulasi salah. Yang 1 update nilai, yang 1 lagi baca function sum sehingga ada perhitungan nilai sebelum / sesudah update**  
**Needs concurrency control**

**telat di sum**

- o Oleh karena itu, perlu recovery karena mungkin ada :

### 1. Computer failure (System crash)

↳ Hardware, software, or network error sewaktu lagi transaction

↳ Hardware crash are usually media failures seperti main memory failure

### 2. Transaction / System error

↳ Integer overflow, division by 0, erroneous parameter values, logical programming error

### 3. Local errors / exception conditions detected by transactions

↳ Data for transaction may not be found

↳ Insufficient account balance in banking database

### 4. Concurrency control enforcement

↳ Bisa abort transaksi untuk menyelesaikan deadlock dari berbagai transaksi.

### 5. Disk failure

↳ Bisa hilang data di disk karena read/write malfunction / head crash

### 6. Physical problems & catastrophes

↳ Fire, theft, sabotage, salah pasang, etc.

- o A transaction is an atomic unit of work that is either completed in its entirely or not done at all.

- o Recovery manager keeps track this perintah :

1. BEGIN TRANSACTION

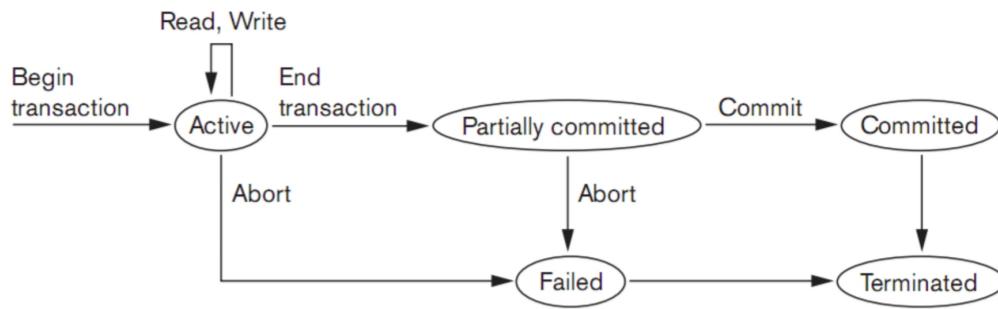
2. READ OR WRITE

3. END TRANSACTION

4. COMMIT TRANSACTION

5. ROLLBACK / ABORT

o



o

Properties of transactions : ACID

- o Atomicity : Transaction is atomic unit of processing and perform semua / go.
- o Consistency preservation : If its complete execution takes database from one consistent state to another
- o Isolation : Should appear as though it is executed in isolation from other transaction and jgn interfered with by any other transactions executing concurrently.
- o Durability/permanency : Change pke commit hrs persist di database dan must not be lost because of any failure.

o

Off autocommit = SET @@ AUTOCOMMIT=0

On autocommit = SET @@ AUTOCOMMIT=1

o

COMMIT = Make changes permanent

ROLLBACK = Undo changes

o

Undo only a part of current transaction by save point.

↳ SAVEPOINT savepoint\_name

↳ ROLLBACK TO SAVEPOINT savepoint\_name

# Week 10

- Database threat :- Loss of integrity (Protect from improper modification seperti CRUD)
  - Loss of availability (Make object available yang ada legitimate right)
  - Loss of confidentiality (Protect data from unauthorized disclosure)
- 2 tipe database security mekanisme ◦

## 1. Discretionary Security Mechanisms

↳ Grant privileges to user untuk CRUD

## 2. Mandatory Security Mechanisms

↳ Multilevel security by classifying data seperti hanya libat data classified at user's own / lower classification level

- Control measures ◦

### 1. Access control

↳ Include provision for restrict access ke database by creating user account & password to control the login

### 2. Inference Control

↳ Statistical database used to provide statistical information seperti perusahaan besar untuk cari informasi populasi bkn individu.

### 3. Flow Control

↳ Prevent information from flowing yg dpt reach unauthorized users

↳ Covert channels are channels that pathways for information to flow implicitly in ways that violate the security policy of an organization

## 4c Data Encryption

- ↳ Protect data yg dikirim lewat communication network
- ↳ Encoded using some algorithm

- Database Administrator /DBA is the central authority for managing a database system
- DBA has DBA account in DBMS, sometimes called a system or superuser account which provide powerful capabilities that are not made available to regular database accounts and users
- Types of action :-
  - Account Creation
  - Privilege Granting
  - Privilege Revocation
  - Security Level Assignment
- If any tampering with database is suspected, a database audit is performed, which consists of review log to examine all akses dan operasi applied to database during certain time period.
- When an illegal or unauthorized operation is found, the DBA can determine account number used to perform operasi
- View = Single table that is derived from other tables dan virtual disimpan ke database

○

```
CREATE USER 'username'@'hostname'  
IDENTIFIED BY 'password';
```

- Example

```
CREATE USER 'dennis'@'localhost'  
IDENTIFIED BY 'dennis';
```

- Show the contents of the users catalog view

```
SELECT host, user, password  
FROM mysql.user;
```

```
DROP USER 'username'@'hostname';
```

Example

```
DROP USER 'dennis'@'localhost';
```

```
RENAME USER 'username'@'hostname'  
TO 'new_username'@'new_hostname';
```

Example

```
RENAME USER 'dennis'@'localhost'  
TO 'dennisgunawan'@'localhost';
```

0

```
SET PASSWORD FOR 'username'@'hostname' = PASSWORD('password');
```

Example

```
SET PASSWORD FOR 'dennisgunawan'@'localhost'  
= PASSWORD('dennisgunawan');
```

0

```
GRANT privileges | ALL PRIVILEGES  
ON table_name | [database_name.] * | *.*  
TO grantees;
```

```
GRANT SELECT  
ON actors  
TO 'dennis'@'localhost';
```

```
GRANT SELECT  
ON entertainment.*  
TO 'dennis'@'localhost';
```

```
GRANT CREATE USER  
ON *.*  
TO 'dennis'@'localhost';
```

```
GRANT INSERT, UPDATE  
ON actors  
TO 'dennis'@'localhost';
```

```
GRANT ALL PRIVILEGES  
ON *  
TO 'dennis'@'localhost';
```

```
GRANT CREATE, ALTER, DROP  
ON *.*  
TO 'dennis'@'localhost';
```

```
GRANT UPDATE (year, rank)  
ON movies  
TO 'dennis'@'localhost';
```

```
GRANT CREATE, ALTER, DROP, CREATE VIEW  
ON entertainment.*  
TO 'dennis'@'localhost';
```

0

```
GRANT SELECT, INSERT  
ON actors  
TO 'dennis'@'localhost'  
WITH GRANT OPTION;
```

'dennis'@'localhost'

```
GRANT INSERT  
ON actors  
TO 'rico'@'localhost';
```

→ Agar bisa grant kasih orang

0

```
REVOKE privileges | ALL PRIVILEGES  
ON table_name | [database_name.] * | *.*  
FROM grantees;
```

REVOKE = Mencabut

```
REVOKE SELECT  
ON actors  
FROM 'dennis'@'localhost';
```

```
REVOKE GRANT OPTION  
ON actors  
FROM 'dennis'@'localhost';
```

0

```
CREATE [OR REPLACE] VIEW view_name [(column_list)]  
AS SELECT_statement  
[WITH CHECK OPTION];
```

Example

```
CREATE VIEW movie_production  
AS SELECT name, year  
FROM movies;
```

```
SELECT name, year  
FROM movie_production;
```

```
CREATE OR REPLACE VIEW good_movies  
AS SELECT name, rank  
FROM movies  
WHERE rank > 6.5;
```

```
SELECT name, rank  
FROM good_movies;
```

0

```
CREATE OR REPLACE VIEW good_movies  
AS SELECT name, rank  
FROM movies  
WHERE rank > 6.5  
WITH CHECK OPTION;
```

Untuk cek ada perubahan di update, insert, delete

0

```
DROP VIEW [IF EXISTS] view_name;
```

When IF EXISTS is specified, no error messages appear.

Example

```
DROP VIEW movie_production;
```

```
DROP VIEW IF EXISTS good_movies;
```

## Week 11

- o Functional dependency = Constraint between 2 sets of attributes from the database
- o  $X \rightarrow Y$  :- Y depends on X
  - Y is determined by X
  - Values of X determined by values of Y
  - Y is functionally dependent on X
- o Functional dependency = A property of semantics / meaning of the attributes
- o Employee(EmployeeID, EmployeeName)
  - EmployeeID  $\rightarrow$  EmployeeName
- Project(ProjectNumber, ProjectName, ProjectLocation)
  - ProjectNumber  $\rightarrow$  { ProjectName, ProjectLocation }
- o Normalisasi dibentuk oleh Codd (1972) yang berguna untuk :-
  - Minimizing redundancy
  - Minimizing insertion, deletion, update anomalies
- o Normal Forms
  - First Normal Form (1NF)
  - Second Normal Form (2NF)
  - Third Normal Form (3NF)
  - Boyce-Codd Normal Form (BCNF)
  - Fourth Normal Form (4NF)
  - Fifth Normal Form (5NF)
  - Domain-Key Normal Form (DKNF)
- o 1NF :- Disallow multivalue, composite attribute
- o Domain of attribute must include only atomic (Simple, indivisible) value.

## o Full Functional Dependency

A functional dependency  $X \rightarrow Y$  is a full functional dependency if  $Y$  is functionally dependent on  $X$  but **not functionally dependent on any proper subset of  $X$**

A functional dependency  $X \rightarrow Y$  is a full functional dependency if removal of any attribute  $A$  from  $X$  means that the dependency does not hold any more

- For any attribute  $A \in X$ ,  $(X - \{A\}) \rightarrow Y$  does not functionally determine  $Y$

## Partial Dependency

- A functional dependency  $X \rightarrow Y$  is a partial dependency if  $Y$  is functionally dependent on  $X$  and **also functionally dependent on any proper subset of  $X$**

- A functional dependency  $X \rightarrow Y$  is a partial dependency if some attribute  $A \in X$  can be removed from  $X$  and the dependency still holds
  - For some  $A \in X$ ,  $(X - \{A\}) \rightarrow Y$

## o WorksOn(EmployeeID, ProjectNumber, Hours, ProjectName)

- $\text{EmployeeID}, \text{ProjectNumber} \rightarrow \text{Hours}, \text{ProjectName}$
- $\text{ProjectNumber} \rightarrow \text{ProjectName}$

$\text{EmployeeID}, \text{ProjectNumber} \rightarrow \text{Hours}$  is a full functional dependency

- Neither  $\text{EmployeeID} \rightarrow \text{Hours}$  nor  $\text{ProjectNumber} \rightarrow \text{Hours}$  holds

$\text{EmployeeID}, \text{ProjectNumber} \rightarrow \text{ProjectName}$  is a partial dependency

- $\text{ProjectNumber} \rightarrow \text{ProjectName}$  holds

## o 2NF :- Harus 1NF

- Hilangkan partial functional dependencies
- Make sure to keep relation with original primary key and any attributes that are fully functionally dependent on it.

## o A functional dependency $X \rightarrow Y$ in a relation schema $R$ is a **transitive dependency** if there is a set of attributes $Z$ that is neither a candidate key nor a subset of any key of $R$ , and both $X \rightarrow Z$ and $Z \rightarrow Y$ hold.

## o Works-For(SSN, D#, DPhone)

- $\text{SSN} \rightarrow \text{D\#}, \text{DPhone}$
- $\text{D\#} \rightarrow \text{DPhone}$

The dependency  $\text{SSN} \rightarrow \text{DPhone}$  is transitive through  $\text{D\#}$

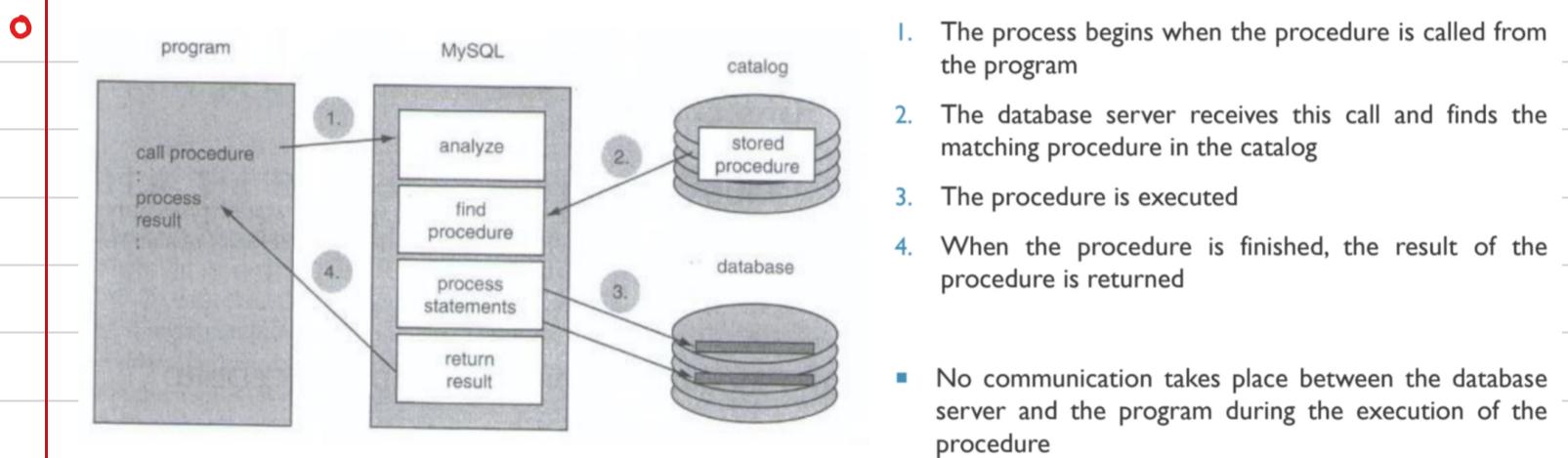
- Both the dependencies  $\text{SSN} \rightarrow \text{D\#}$  and  $\text{D\#} \rightarrow \text{DPhone}$  hold
- $\text{D\#}$  is neither a key itself nor a subset of the key of  $\text{Works-For}$

$\text{SSN} \rightarrow \text{D\#}$  is not a transitive dependency

- 3NF :- Hasus 2NF
  - Hilangkan transitive functional dependencies
  - Decompose & set up a relation that includes the nonkey attributes that functionally determines other nonkey attributes
- Higher NF eliminates more modification anomalies than lower NF

# Week 12

- o **Stored procedures** = Certain piece of code consist of declarative and procedural SQL statements stored in the catalog of a database that can activate by call program, trigger, or other stored procedure.
- o Each stored procedure consists of at least 3 parts:
  - Name
  - List of parameters
  - Body



```
CREATE PROCEDURE procedure_name
  ( [ [ IN | OUT | INOUT ] parameter_name data_type,
    [ IN | OUT | INOUT ] parameter_name data_type,
    ... ]
routine_body
```

- CREATE PROCEDURE delete\_actor ( IN actor\_id CHAR(5) )  
...
  - CREATE PROCEDURE show\_all\_movies ()  
...
- A stored procedure does not need parameters, but opening and closing brackets are still required

- o 3 type parameter:
  - **Input** → Data pass to stored procedure
  - **Output** → Store procedure use parameter when answer/result must be returned
  - **Input/Output** → Act as input & output

o

```
DECLARE variable_name, variable_name, ... data_type  
[ DEFAULT default_value ];
```

- `DECLARE score DECIMAL(5,2);`
- `DECLARE name VARCHAR(40);`
- `DECLARE number1, number2 INTEGER;`
- `DECLARE number1 INTEGER DEFAULT 100;`

o

```
SET variable_name { = | := } scalar_expression,  
variable_name { = | := } scalar_expression,  
...;
```

- `SET var1 = 1;`
- `SET var1 := 1;`
- `SET var1 = 1, var2 = var1;`

o

```
CALL [ database_name. ] stored_procedure_name  
([ parameter, parameter, ... ] );
```

- `CALL delete_actor ('A0001');`
- `CALL show_all_movies ();`

o

Untuk fetch data 1 row memakai `SELECT INTO`, kalau multiple row pakai `CURSOR`

o

```
SELECT ...  
INTO local_variable, local_variable, ...  
FROM ...  
WHERE ...  
GROUP BY ...  
HAVING ...  
ORDER BY ...  
LIMIT ...;
```

```
CREATE PROCEDURE count_actors ( OUT counter INTEGER )  
BEGIN  
    SELECT COUNT(*)  
    INTO counter  
    FROM actors;  
END  
  
CALL count_actors (@number_of_actors);  
  
SELECT @number_of_actors;
```

o

```
DROP PROCEDURE [IF EXISTS] [database_name.] procedure_name;
```

- `DROP PROCEDURE delete_actor;`
- `DROP PROCEDURE show_all_movies;`

# Week 13

o 4 special SQL statements with cursor :

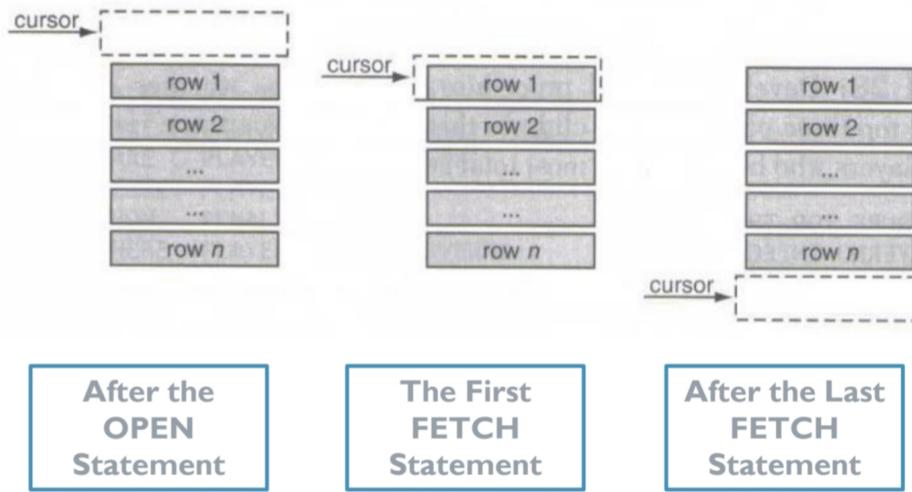
1. DECLARE CURSOR
2. OPEN CURSOR
3. FETCH CURSOR
4. CLOSE CURSOR

o

`FETCH cursor_name INTO local_variable, local_variable, ...;`

- We can use FETCH CURSOR statements to retrieve the created result row by row into the stored procedure
- At a certain moment, only one row from the result is visible: the current row
- It is as if an arrow points to one row from the result
- With the FETCH CURSOR statement, we move this cursor to the next row

o



o

Advantage of stored procedure :

- Jika specific set of update di database logically form a unit and if set itu use multiple application, better put di 1 procedure.
- Minimize network traffic (Hrs selesai dulu proceduresnya)
- Not depend on particular host language

o

STORED FUNCTIONS

VS

STORED PROCEDURES

Have input but no output parameter

No call

Ada return

Menyakon output parameter

Ada call  
No return

o

```
CREATE FUNCTION function_name  
  ([ parameter_name data_type,  
    parameter_name data_type,  
    ... ] )  
  RETURNS data_type  
routine_body
```

```
GRANT EXECUTE  
ON FUNCTION stored_function_name  
TO grantees  
[WITH GRANT OPTION];
```

```
DROP FUNCTION [IF EXISTS] [database_name.] function_name;
```

- DROP FUNCTION count\_actors;

o

**Trigger** = Piece of code consist of procedural & declarative statements stored in the catalog and activated by the database server if a specific operation is executed on database and only when certain condition holds

o

Trigger turn passive database server into active

o

Trigger add 3 main elements :

- Trigger moment
- Trigger event
- Trigger action

o

```
CREATE TRIGGER insert_actor  
  AFTER INSERT ON actors FOR EACH ROW  
BEGIN  
  INSERT INTO actors_copy(id,first_name,last_name,gender,birth_date)  
  VALUES(NEW.id,NEW.first_name,NEW.last_name,NEW.gender,NEW.birth_date);  
END
```

- NEW is specified in front of the column names
- If a row is inserted, it looks as if there is a table called NEW
- The column names of this NEW table are equal to those of the triggering table (those in which the new row appears)
- As a result of specifying NEW in front of the column names, the actor that is added to the actors table is used

o

```
CREATE TRIGGER delete_actor  
  BEFORE DELETE ON actors FOR EACH ROW  
BEGIN  
  INSERT INTO actors_history(id,first_name,last_name,gender,birth_date)  
  VALUES(OLD.id,OLD.first_name,OLD.last_name,OLD.gender,OLD.birth_date);  
END
```

- The keyword OLD is now specified instead of NEW
- After we remove a row, a table called OLD exists with column names that are equal to those of the triggering table, in which the removed row occurs
- When you update rows, the NEW and the OLD tables both exist
  - The row with the old values appears in the OLD table and the new row appears in the NEW table

o

```
DROP TRIGGER trigger_name
```