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SIMULATED ANNEALING



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DEFINITION

- **Simulated Annealing (SA)** adalah sebuah algoritma metaheuristik untuk memecahkan masalah-masalah optimasi, tepatnya menemukan global optimum dari sebuah fungsi
- Dikenalkan oleh Kirkpatrick et al. pada 1983 untuk menjawab Traveling Salesman Problem, SA terinspirasi dari perilaku atom pada **proses annealing atau pemanasan logam** pada bidang metalurgi yang dilakukan untuk mengubah properti fisika logam
- SA juga merupakan adaptasi dari Metropolis-Hastings oleh Metropolis et al. pada 1953, sebuah metode Monte Carlo untuk mendapatkan sampling

ANNEALING PROCESS



- Dalam proses annealing di metalurgi, sebuah logam **dipanaskan ke temperatur yang tinggi secara cepat**, lalu **didinginkan secara berkala**
- Pada temperatur yang tinggi, atom-atom logam juga bergerak dengan cepat dan saat temperatur sudah diturunkan, energi kinetik pada atom-atom pun juga berkurang
- Pada akhir proses annealing, atom-atom pada logam berada di sebuah crystalline structure atau **kondisi yang lebih teratur**, sehingga material logam lebih **lentur dan mudah dipakai atau dikerjakan**

WHY USE SA?

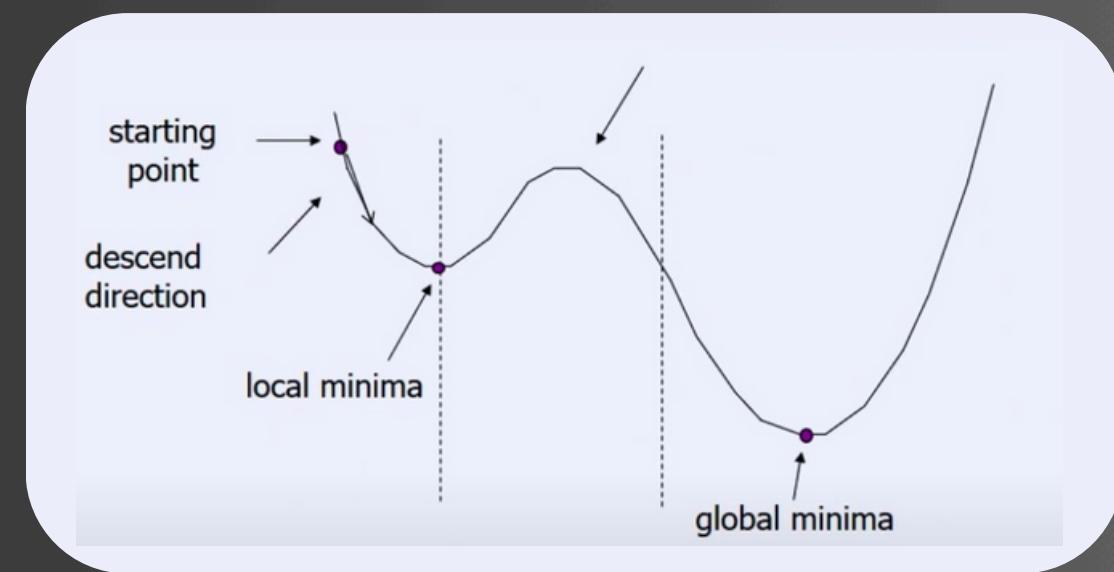
- Simpel dan fleksibel karena dapat diaplikasikan di banyak jenis masalah optimasi
- Dapat lolos dari local optimum dan menemukan global optimum secara baik karena mampu menerima solusi yang lebih buruk
- Tidak memerlukan kalkulasi matematika yang kompleks seperti informasi gradien pada fungsi objektif seperti gradient descent

APPLICATIONS

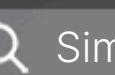
- VLSI Design
- Multiprocessor scheduling
- Crystal structure prediction
- Traveling salesman
- Vehicle routing
- Scheduling problems

HOWEVER...

- Cooling schedule perlu diatur secara baik: apabila terlalu cepat, SA berisiko konvergensi menuju local optimum; apabila terlalu lambat, SA dapat membutuhkan waktu konvergensi yang terlalu lama
- Memerlukan banyak iterasi untuk solusi yang baik, terutama apabila masalahnya kompleks
- Tidak memanfaatkan heuristik atau bimbingan yang problem-specific sehingga dapat lebih inefisien dibandingkan algoritma lain yang lebih dispesialisasi



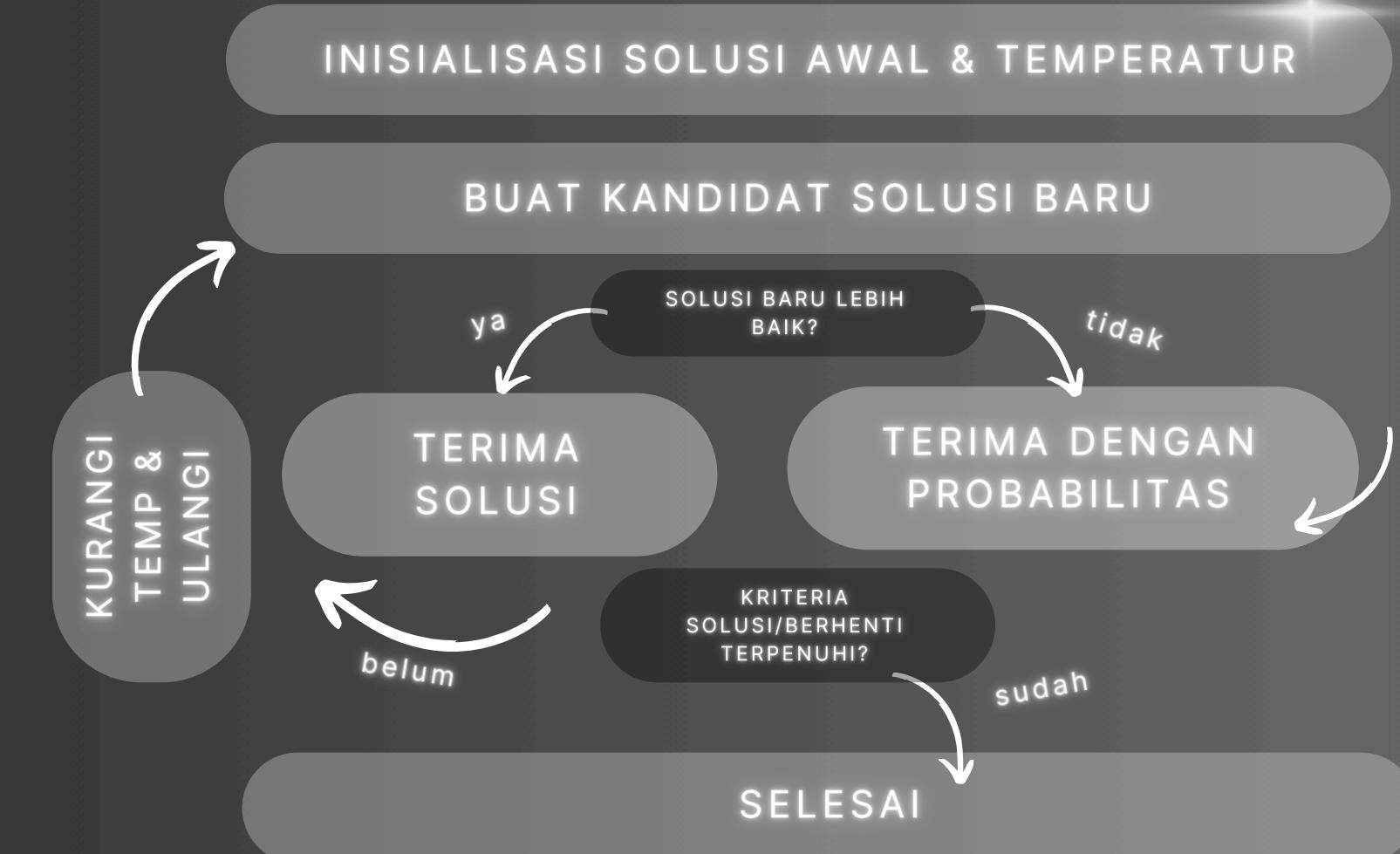
*Harus kabur
dari terjebaknya
di local minima



HOW IT WORKS

- Awalnya, algoritma dimulai dengan **temperatur tinggi** dan suatu konfigurasi molekul logam yang merepresentasikan **solusi awal random**
- Di setiap iterasi, **solusi baru akan dibuat** dengan membuat perubahan dari sebelumnya, dan **temperatur sistem akan perlakan diturunkan**
- Saat temperatur masih tinggi, solusi lebih buruk memiliki probabilitas lebih tinggi untuk diterima
- Saat temperatur sudah merendah, sistem akan lebih bercondong ke solusi yang lebih baik, layaknya hill-climbing klasikal
- Dalam kata lain, algoritma ini **rela menerima kondisi yang buruk di awal** demi mendapat yang **terbaik di akhir**
- Dengan demikian, semestinya suatu crystalline structure atau **solusi optimal akan didapat di akhir** algoritma
- Solusi global optima dapat diraih apabila cooling proses semakin melambat

OVERALL CONCEPT



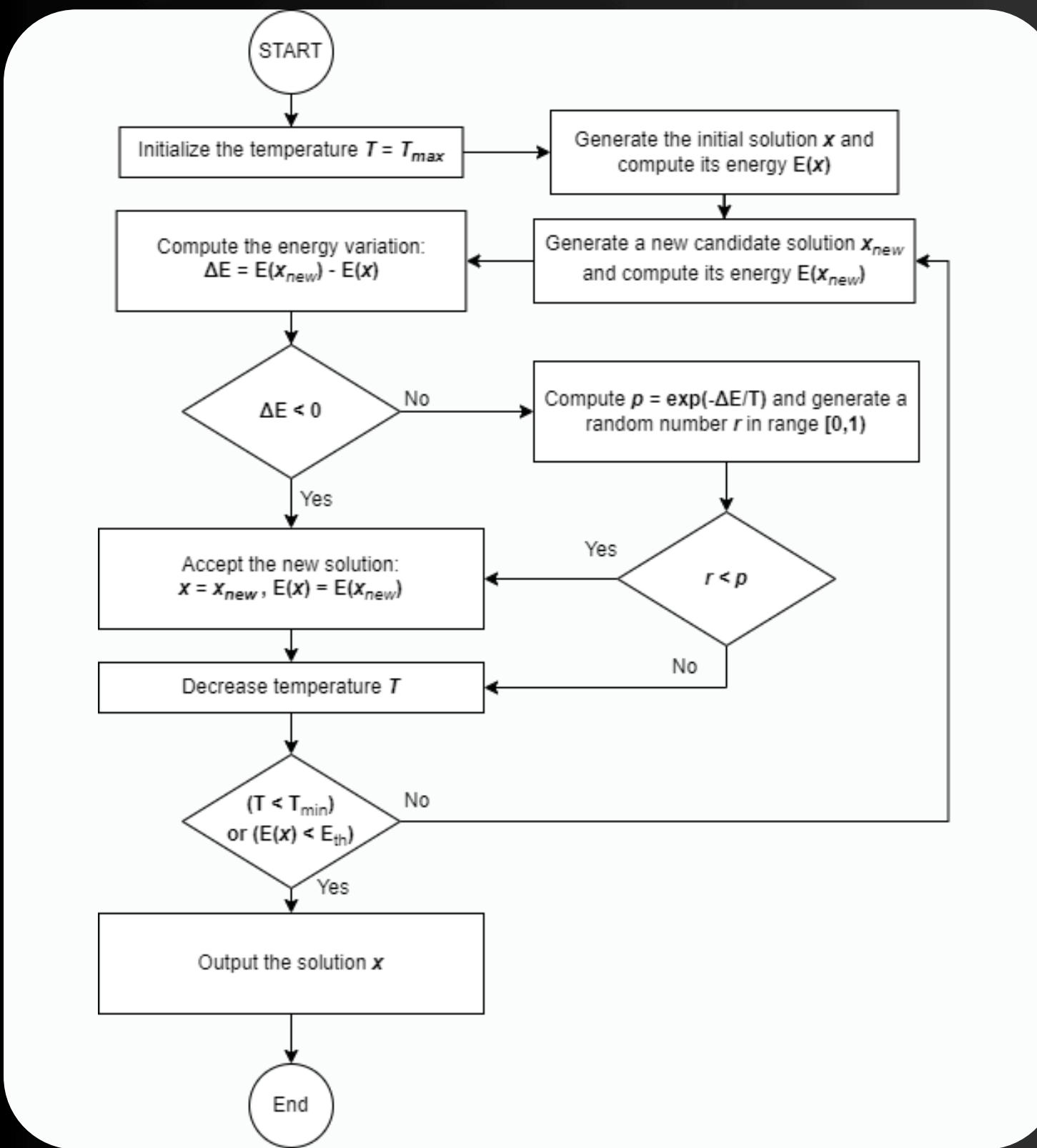
Physical Annealing	Simulated Annealing
Metal	Optimization Problem
Energy State	Cost Function
Temperature	Control Parameter
Crystalline Structure	The Optimal Solution

KEY COMPONENTS & STOPPING CRITERIA





PSEUDOCODE



Algorithm 1: Simulated Annealing Algorithm

```

Input:  $T_0, c, itmax$ 
Output:  $x_{opt}, f_{best}$ 
1 Bangkitkan  $x_0$ ;
2  $f_{best} = f(x_0)$ ;
3  $T = T_0$ ;
4  $it = 1$ ;
5 while kriteria penghentian belum dicapai do
6   if  $it < itmax$  then
7     Bangkitkan  $\Delta x$  berdistribusi Gaussian;
8     if  $f(x_0 + \Delta x) < f(x_0)$  then
9        $f_{best} = f(x_0 + \Delta x)$ ;
10       $x_0 = x_0 + \Delta x$ ;
11    else
12       $\Delta f = f(x_0 + \Delta x) - f(x_0)$ ;
13      Bangkitkan bilangan random  $r$ ;
14      if  $r < \exp\left(\frac{-\Delta f}{kT}\right)$  then
15         $x_0 = x_0 + \Delta x$ ;
16    $it = it + 1$ ;
17    $T = T * c$ ;

```

<https://github.com/oscarjiro/simulated-annealing-example>



PARAMETER TUNING

Parameter	Exploration	Exploitation
Initial Temperature (T_0)	Higher Temperature leads to more exploration	Lower Temperature accelerates exploitation
Cooling Schedule (c)	Slow Cooling promotes extensive exploration	Fast Cooling focuses on early exploitation
Number of Iteration (it_{max})	Longer run times provide more opportunities for exploration	Premature stopping focuses on immediate exploitation

Complex problem may require more exploration but simple problem will benefit more from exploitation

COOLING SCHEDULE

Schedule	Formula	Cooling Rate	Characteristics
Geometric Cooling	$T_{k+1} = \alpha \cdot T_k$	Moderate	Most common, balance between speed and exploration
Linear Cooling	$T_{k+1} = T_k - \beta$	Moderate to Fast	Constant decrease, simple but needs good parameter tuning
Exponential Cooling	$T_{k+1} = T_0 \cdot e^{-\lambda k}$	Fast	Faster cooling at the start, slows down later
Logarithmic Cooling	$T_{k+1} = \frac{T_0}{\log(k+1)}$	Slow	Slow cooling, allows extensive exploration
Quadratic Cooling	$T_{k+1} = \frac{T_0}{(k+1)^2}$	Very Fast	Rapid cooling, can lead to early convergence
Adaptive Cooling	Varies based on improvement	Varies (based on progress)	Flexible, adjusts cooling based on optimization progress

TIME COMPLEXITY $\rightarrow O(T * I)$

SPACE COMPLEXITY $\rightarrow O(M * N * K)$

Informations:

- I : Number of iterations
- T : Temperature
- N : Number of array
- M : Dimension
- K : Additional Space



DISTRIBUSI PROBABILITAS BOLTZMANN

$$P(E) = e^{\frac{-E}{k \times T}}$$

KRITERIA METROPOLIS

$$\Delta E = E_{i+1} - E_i = A$$

$$r \leq e^{\frac{-\Delta E}{k \times T}} = B$$

$$|x_{i+1} - x_i| < \varepsilon = C$$

Informations :

- k = Boltzmann constant
= 1.380649×10^{-23} J/K
- e = Euler's number
= 2.71828
- x_i = i^{th} vertex
- ΔE = Energy / Value differences
- r = Random value = $\text{randn}(0,1)$
- T = Temperature
- ε = Termination criterion
- σ = Standard deviation
- E_i = This solution & E_{i+1} = Neighborhood solution

ADDENDUM TO DETERMINE NEXT POINT

$$x = x_{i+1} + \sigma \left[\sum_{i=1}^n R_i - \frac{N}{2} \right]$$

$$\sigma = \frac{\max - \min}{6}$$

Informations :

- R_i = Random value
- N = Number of all element or R_i

COOLING SCHEDULE → STEPS

STEP-1: Choose initial point (x_0), termination criterion (ε). Set initial temperature to high value (T), and number of iterations (n)

STEP-2: Calculate neighboring point randomly

STEP-3: If $A < 0$, set $t = t + 1$
else create a random number $r = \text{rand}(0,1)$

If B true, set $t = t + 1$
else go to **STEP-2**

STEP-4: If C true and T is small. Return E
else if $(t \bmod n) = 0$, then lower T . Go to **STEP-2**

$$T_{new} = T_{old} \times c$$

c = Temperature factor



OPTIMIZING A FUNCTION

Himmelblau function:

$$f(\vec{x}) = (x_1^2 + x_2 - 11)^2 + (x_1 + x_2^2 - 7)^2 \quad \text{for } x_1, x_2 \in (0, 5)$$

I. Initialize parameter & Choose search domain (0, 5) for x_1 and x_2

Initial Point

Take mid-point : $\vec{x}_0 = (2.5, 2.5)$

Termination factor = Permissible error

$\varepsilon = 0.001$

Temperature factor Max Iteration

$c = 0.5$

$t = 2$

Initial Temp (T) * 0th Iteration

Search domain = (0,0) ; (5,0) ; (0,5) ; (5,5)

$$f(x_1, x_2) = (x_1^2 + x_2 - 11)^2 + (x_1 + x_2^2 - 7)^2$$

$$f(0,0) = (0+0-11)^2 + (0+0-7)^2 = 170$$

$$f(5,0) = (5^2+0-11)^2 + (5+0^2-7)^2 = 200$$

$$f(0,5) = (0^2+0-11)^2 + (0+5^2-7)^2 = 360$$

$$f(5,5) = (5^2+5-11)^2 + (5+5^2-7)^2 = 890$$

$$T = (170 + 200 + 360 + 890)/4 = 405$$

II. Generate value point for x_1 dan x_2 , random value, determine next point, & ΔE value

Value of Initial Point

$$f(x_0) = f(2.5; 2.5) = (2.5^2 + 2.5 - 11)^2 + (2.5 + 2.5^2 - 7)^2 = 8.125$$

Generate random value

Assume $\Delta_1 = 0.037$ & $\Delta_2 = -0.086$

Determine next point

$$x_1 = (2.5; 2.5) + (0.037; -0.086) = (2.537; 2.414)$$

$$f(x_1) = (2.537^2 + 2.414 - 11)^2 + (2.537 + 2.414^2 - 7)^2 = 6.482$$

ΔE value

$$\Delta E = f(x_1) - f(x_0) = 6.482 - 8.125 = -1.643$$

Because $\Delta E < 0$, then accept new value and increment

$$t=t+1=0+1=1$$

$$\|x_1 - x_0\| = \sqrt{(0.037)^2 + (-0.086)^2} = 0.101$$

Because $0.101 > 0.001$, then don't stop & because

$$t \bmod n = 1 \bmod 1 = 0, \text{ then } T = T/2 = 405/2 = 202.5$$

III. Repeat Steps 2 for next iteration (0 → 1)

Generate random value

Assume $\Delta_1 = -0.465$ & $\Delta_2 = -1.810$

Determine next point

$$x_2 = (2.537; 2.414) + (-0.465; -1.810) = (2.072; 0.604)$$

$$f(x_1) = (2.072^2 + 0.604 - 11)^2 + (2.072 + 0.604^2 - 7)^2 = 58.067$$

ΔE value

$$\Delta E = f(x_1) - f(x_0) = 58.067 - 6.482 = 51.585$$

Because $\Delta E > 0$, then don't reject straight way.

Instead, use Metropolis criterion to decide whether to accept or reject this worse point

$$P(E_2) = e^{\frac{-\Delta E}{T}} = e^{\frac{-51.585}{202.5}} = 0.775$$

Generate uniformly distributed number r

r = 0.649 then r < 0.775 = 0.649 < 0.775 means accept the point & increment t=t+1=1+1=2 (Stopping criteria). **6.482 to 58.067 (Worse Acceptance)**



OPTIMIZING A FUNCTION

$$Z = 2 \times x_1 + x_2^2 \quad \text{for } 0 \leq x_1, x_2 \leq 9$$

I. Initialize parameter & Choose
search domain (0, 9) for x_1 and x_2

Initial Temp (T) * Iterasi ke-0

Search domain = (4,6);(6,4);(4,8);(0,6);(5,7)

$$Z = f(x_1, x_2) = 2 \times x_1 + x_2^2$$

$$f(4,6) = 2 \times 4 + 6^2 = 44$$

$$f(6,4) = 2 \times 6 + 4^2 = 28$$

$$f(4,8) = 2 \times 4 + 8^2 = 72$$

$$f(0,6) = 2 \times 0 + 6^2 = 36$$

$$f(5,7) = 2 \times 5 + 7^2 = 59$$

$$T = (44 + 28 + 72 + 36 + 59)/5 = 47.9$$

Multiple T with suitable number to optimize the temperature for more maximize

$$T * 10 = 47.9 * 10 = 479 \approx 500$$

Temperature factor Max Iteration

$$c = 0.5$$

$$t = 2$$

II. Generate value of initial point for x_1 and x_2 ,
update T, & Generate next point

Value of Initial Point = (5, 7) & Update T * Iterasi 0

$$f(x_0) = f(5,7) = 2 \times 5 + 7^2 = 59$$

$$T_{\text{new}} = T_{\text{old}} * c = 500 * 0.5 = 250$$

Generate next point using formula

$$x_{\text{old}(1)} = 5 \text{ and } \sigma = (\text{max-min})/6 = (9-0)/6 = 1.5$$

Generate N and R_i :

- Assume N=4, $R_1=0.72$, $R_2=0.72$, $R_3=0.15$, $R_4=0.32$
- $\sum_{i=1}^4 R_i = 0.72 + 0.72 + 0.15 + 0.32 = 4.95$

$$\therefore x_{\text{new}(1)} = x_{\text{old}(1)} + \sigma * (\sum_{i=1}^4 R_i - N/2) = 5 + 1.5 * (4.95 - 4/2) = 4.95$$

$$x_{\text{old}(2)} = 5 \text{ and } \sigma = (\text{max-min})/6 = (9-0)/6 = 1.5$$

Generate N and R_i :

- Assume N=4, $R_1=0.32$, $R_2=0.5$, $R_3=0.8$, $R_4=0.4$
- $\sum_{i=1}^4 R_i = 0.32 + 0.5 + 0.8 + 0.4 = 2.02$

$$\therefore x_{\text{new}(2)} = x_{\text{old}(2)} + \sigma * (\sum_{i=1}^4 R_i - N/2) = 7 + 1.5 * (2.02 - 4/2) = 7.03$$

**∴ New point = ($x_{\text{old}(1)}$, $x_{\text{new}(2)}$) = (4.05; 7.03)
 $f(4.05; 7.03) = 2 \times 4.05 + 7.03^2 = 59.32$**

III. Check ΔE , Repeat Steps 2 for next iteration (0 → 1), Update T

ΔE value for iteration 0

$$\Delta E = 59.32 - 59 = 0.32 > 0 \rightarrow \text{Go for metropolis algorithm}$$

Metropolis Algorithm

$$P(E_2) = e^{\frac{-\Delta E}{T}} = e^{\frac{-0.32}{250}} = 0.99 \quad \text{Assume } r = 0.495, \text{ then } r < 0.99 = 0.495 < 0.99$$

Generate next point using formula

Using $x_{\text{old}(1)} = 4.05$ & $x_{\text{old}(2)} = 7.03$

For $x_{\text{new}(1)}$, assume N=4, $R_1=0.39$, $R_2=0.22$, $R_3=0.99$, $R_4=0.06$

For $x_{\text{new}(2)}$, assume N=4, $R_1=0.5$, $R_2=0.2$, $R_3=0.2$, $R_4=0.1$

For the same method like in Step 2, new point for iteration 1 is (3.54; 5.53)

$$\therefore f(3.54; 5.53) = 2 \times 3.54 + 5.53^2 = 39.46$$

Update T

$$T_{\text{new}} = T_{\text{old}} * c = 250 * 0.5 = 125$$

ΔE value for iteration 1

$$\Delta E = 39.46 - 59.32 = -19.82 < 0 \rightarrow \text{Accept value and return}$$

∴ 59 to 59.32 to 39.46 (Most optimal)



TRAVELING SALESMAN PROBLEM

	Kota 1	Kota 2	Kota 3	Kota 4	Kota 5
Kota 1	0	132	217	164	58
Kota 2	132	0	290	201	79
Kota 3	217	290	0	113	303
Kota 4	164	201	113	0	196
Kota 5	58	79	303	196	0

 $T_{old} = 1000$ $T_{akhir} = 15$ **Ambil rute pertama misalkan rute**

$$1-2-3-5-4-1 = 132+290+303+196+164 = 1085 = f_1$$

Konstruksi rute baru dengan menukar 2-3 menjadi 3-2

$$1-3-2-5-4-1 = 217+290+79+196+164 = 946 = f_2$$

∴ Karena $f_2 < f_1$, maka rute 2 diterima**Konstruksi rute baru dengan slide 5-4 ke kiri dua kali**

$$1-5-4-3-2-1 = 58+196+113+290+132 = 789 = f_3$$

∴ Karena $f_2 > f_3$, maka gunakan kriteria metropolis**Cek nilai P(E) dengan $r = 0.2$**

$$P(E) = e^{\frac{-E}{T}} = e^{\frac{-789}{1000}} = 0.4543$$

$r < 0.4543 = 0.2 < 0.4543$ (Accept solusi)
 $T_{new} = T_{old} * 0.01 = 1000 * 0.01 = 10 < T_{akhir}$
 $= 10 < 15$ (Proses berhenti)

∴ Solusi terbaik saat ini = 1-5-4-3-2-1 $*f = E$

STEPS IN TSP USING SA

STEP-1: Tentukan rute awal secara acak**STEP-2:** Bangkitkan rute baru dengan 3 metode (Beserta contohnya)

Misalkan ada rute dengan 7 kota yaitu 4-6-2-3-5-1-7

1. Metode membalik (Flip)
Rute baru menjadi 4-6-1-5-3-2-7
2. Metode menukar (Swap)
Rute baru menjadi 4-6-1-3-5-2-7
3. Metode menggeser (Slide)
Rute baru menjadi 4-6-3-5-1-2-7

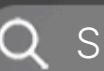
Cara menentukan antara ketiga metode ini dapat menggunakan **bilangan random r**

- $r \leq 0.33$ → Pilih metode swap
- $0.33 \leq r \leq 0.67$ → Pilih metode slide
- $r \geq 0.67$ → Pilih metode flip

STEP-3: Jika jarak total rute baru < rute awal, maka rute baru digunakan sebagai urutan terbaik

Jika jarak total rute baru > rute awal, maka gunakan kriteria metropolis untuk menolak / menerima rute baru

STEP-4: Jika iterasi maksimum tercapai, maka harus ada penurunan temperatur. Jika suhu sudah lebih kecil dari suhu awal, maka proses berhenti dan dapatkan rute terbaik di proses tersebut. Seiring temperatur turun, maka jumlah pasangan kota untuk ditukar berkurang / Pencarian semakin sempit dan didapatkan rute terbaik dengan jarak total minimum



TAMBAHAN

Pada penentuan next point terdapat beberapa cara dari patokan previous point

I. Dengan Penambahan bilangan random Δ

Sesuai dengan contoh sebelumnya pada bagian Himmelblau function yaitu :

- Titik awal = $(x_1, x_2) = (2.5; 2,5)$
- Misalkan $\Delta_1 = 0.037$ & $\Delta_2 = -0.086$
- Titik baru = $(2.5; 2,5) + (0.037; -0.086) = (2.537; 2.414)$

II. Dengan Akurasi ± 6

Misalkan titik awal = $(x_1, x_2) = (4, 5)$ dan random number = $(u_1, u_2) = (0.31; 0.57)$

Akurasi ± 6 berarti titik awal berada di rentang :

$$x_1 = (4-6, 4+6) = (-2, 10) \text{ dan } x_2 = (5-6, 5+6) = (-1, 11)$$

$$\text{Titik baru } x_1 = -2 + u_1 * (10 - (-2)) = -2 + 0.31 * 12 = 1.72$$

$$\text{Titik baru } x_2 = -1 + u_2 * (11 - (-1)) = -1 + 0.57 * 12 = 5.84$$

$$\therefore \text{Titik baru} = (1.72; 5.84)$$

Memakai rumus dan ada hubungannya
III. dengan standard deviation seperti contoh
pada fungsi Z sebelumnya



Links:

- https://webpages.iust.ac.ir/yaghini/Courses/AOR_891/05_Simulated%20Annealing_01.pdf
- <https://www.mathworks.com/help/gads/how-simulated-annealing-works.html>
- <https://www.eecs.uottawa.ca/ordal/papers/sander/node7.html>
- <https://saturncloud.io/glossary/simulated-annealing/>
- <https://sites.gatech.edu/omscs7641/2024/02/19/simulated-annealing-methods-and-real-world-applications/>
- https://algorithmafternoon.com/books/simulated_annealing/chapter02/
- <https://chatgpt.com/share/66fd2850-7a30-8008-8a90-15d72da2b2d8>
- <https://cs.stackexchange.com/questions/110238/why-is-the-usual-simulated-annealing-algorithm-using-a-boltzmann-probability>

Paper:

- https://www.researchgate.net/publication/225260290_The_Theory_and_Practice_of_Simulated_Annealing/link/0c960516543772adef000000/download?tp=eyJjb250ZXh0Ijp7ImZpcnN0UGFnZSI6InB1YmxpY2F0aW9uliwicGFnZSI6InB1YmxpY2F0aW9uln19

Video:

- <https://www.youtube.com/watch?v=Yo1F7Z6ha6s&t=471s>
- <https://www.youtube.com/watch?v=kra3EsSFt2A>
- <https://www.youtube.com/watch?v=A39YiEfXwec&t=379s>
- <https://www.youtube.com/watch?v=TC9WNwM2noM>
- <https://www.youtube.com/watch?v=D2zyF8zhSQ8&t=898s>
- <https://youtu.be/D2zyF8zhSQ8?si=5Dijjer4hRTLDmpl>