

# CS 373: Combinatorial Algorithms, Spring 1999

<http://www-courses.cs.uiuc.edu/~cs373>

## Homework 3 (due Thu. Mar. 11, 1999 by noon)

Name:	
Net ID:	Alias:

Everyone must do the problems marked ►. Problems marked ▷ are for 1-unit grad students and others who want extra credit. (There's no such thing as "partial extra credit"!) Unmarked problems are extra practice problems for your benefit, which will not be graded. Think of them as potential exam questions.

Hard problems are marked with a star; the bigger the star, the harder the problem.

**Note:** When a question asks you to "give/describe/present an algorithm", you need to do four things to receive full credit:

1. (New!) If not already done, model the problem appropriately. Often the problem is stated in real world terms; give a more rigorous description of the problem. This will help you figure out what is assumed (what you know and what is arbitrary, what operations are and are not allowed), and find the tools needed to solve the problem.
2. Design the most efficient algorithm possible. Significant partial credit will be given for less efficient algorithms, as long as they are still correct, well-presented, and correctly analyzed.
3. Describe your algorithm succinctly, using structured English/pseudocode. We don't want full-fledged compilable source code, but plain English exposition is usually not enough. Follow the examples given in the textbooks, lectures, homeworks, and handouts.
4. Justify the correctness of your algorithm, including termination if that is not obvious.
5. Analyze the time and space complexity of your algorithm.

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### Undergrad/.75U Grad/1U Grad Problems

#### ►1. Hashing

- (a) (2 pts) Consider an open-address hash table with uniform hashing and a load factor  $\alpha = 1/2$ . What is the expected number of probes in an unsuccessful search? Successful search?
- (b) (3 pts) Let the hash function for a table of size  $m$  be

$$h(x) = \lfloor A mx \rfloor \bmod m$$

where  $A = \hat{\phi} = \frac{\sqrt{5}-1}{2}$ . Show that this gives the best possible spread, i.e. if the  $x$  are hashed in order,  $x + 1$  will be hashed in the largest remaining contiguous interval.

## ►2. (5 pts) Euler Tour:

Given an **undirected** graph  $G = (V, E)$ , give an algorithm that finds a cycle in the graph that visits every edge *exactly* once, or says that it can't be done.

## ►3. (5 pts) Makefiles:

In order to facilitate recompiling programs from multiple source files when only a small number of files have been updated, there is a UNIX utility called 'make' that only recompiles those files that were changed, *and* any intermediate files in the compilation that depend on those changed. Design an algorithm to recompile only those necessary.

## ►4. (5 pts) Shortest Airplane Trip:

A person wants to fly from city A to city B in the shortest possible time. S/he turns to the traveling agent who knows all the departure and arrival times of all the flights on the planet. Give an algorithm that will allow the agent to choose an optimal route. Hint: rather than modify Dijkstra's algorithm, modify the data. The time is from departure to arrival at the destination, so it will include layover time (time waiting for a connecting flight).

## ►5. (9 pts, 3 each) Minimum Spanning Tree changes Suppose you have a graph G and an MST of that graph (i.e. the MST has already been constructed).

- (a) Give an algorithm to update the MST when an edge is added to G.
- (b) Give an algorithm to update the MST when an edge is deleted from G.
- (c) Give an algorithm to update the MST when a vertex (and possibly edges to it) is added to G.

## Only 1U Grad Problems

## ▷1. Nesting Envelopes

You are given an unlimited number of each of  $n$  different types of envelopes. The dimensions of envelope type  $i$  are  $x_i \times y_i$ . In nesting envelopes inside one another, you can place envelope A inside envelope B if and only if the dimensions A are *strictly smaller* than the dimensions of B. Design and analyze an algorithm to determine the largest number of envelopes that can be nested inside one another.

## Practice Problems

1. The incidence matrix of an undirected graph  $G = (V, E)$  is a  $|V| \times |E|$  matrix  $B = (b_{ij})$  such that

$$b_{ij} = \begin{cases} 1 & (i, j) \in E, \\ 0 & (i, j) \notin E. \end{cases}$$

- (a) Describe what all the entries of the matrix product  $BB^T$  represent ( $B^T$  is the matrix transpose). Justify.
- (b) Describe what all the entries of the matrix product  $B^TB$  represent. Justify.
- ★(c) Let  $C = BB^T - 2A$ . Let  $C'$  be  $C$  with the first row and column removed. Show that  $\det C'$  is the number of spanning trees. ( $A$  is the adjacency matrix of  $G$ , with zeroes on the diagonal).

2.  $O(V^2)$  Adjacency Matrix Algorithms

- (a) Give an  $O(V)$  algorithm to decide whether a directed graph contains a *sink* in an adjacency matrix representation. A sink is a vertex with in-degree  $V - 1$ .
- (b) An undirected graph is a scorpion if it has a vertex of degree 1 (the sting) connected to a vertex of degree two (the tail) connected to a vertex of degree  $V - 2$  (the body) connected to the other  $V - 3$  vertices (the feet). Some of the feet may be connected to other feet.  
Design an algorithm that decides whether a given adjacency matrix represents a scorpion by examining only  $O(V)$  of the entries.
- (c) Show that it is impossible to decide whether  $G$  has at least one edge in  $O(V)$  time.

3. Shortest Cycle:

Given an **undirected** graph  $G = (V, E)$ , and a weight function  $f : E \rightarrow \mathbf{R}$  on the **edges**, give an algorithm that finds (in time polynomial in  $V$  and  $E$ ) a cycle of smallest weight in  $G$ .

4. Longest Simple Path:

Let graph  $G = (V, E)$ ,  $|V| = n$ . A *simple path* of  $G$ , is a path that does not contain the same vertex twice. Use dynamic programming to design an algorithm (not polynomial time) to find a simple path of maximum length in  $G$ . Hint: It can be done in  $O(n^c 2^n)$  time, for some constant  $c$ .

5. Minimum Spanning Tree:

Suppose all edge weights in a graph  $G$  are equal. Give an algorithm to compute an MST.

6. Transitive reduction:

Give an algorithm to construct a *transitive reduction* of a directed graph  $G$ , i.e. a graph  $G^{TR}$  with the fewest edges (but with the same vertices) such that there is a path from  $a$  to  $b$  in  $G$  iff there is also such a path in  $G^{TR}$ .

7. (a) What is  $5^{2^2 9 5^0 + 2 3^4 1 + 1 7 3^2 + 1 1 2^3 + 5 1^4} \bmod 6$ ?
- (b) What is the capital of Nebraska? Hint: It is not Omaha. It is named after a famous president of the United States that was not George Washington. The distance from the Earth to the Moon averages roughly 384,000 km.