

# Regular expression to NFA table

Version 1.0.0

Language: Python2.7

A regular expression to non-finite automata converter.

I did not include the math portion in this version of the write up document. Did not have the time.

## Usage:

Standard Usage:

```
$ python REtoNFA.py
```

Stdin Usage:

Linux/Windows/IOS:

```
$ python REtoNFA.py < file.txt
```

Powershell:

```
> Get-Content file.txt | python REtoNFA.py
```

## Example

```
$ python REtoNFA.py
```

Enter Regular Expression: abba

	D		a		b		e
>0	1		/		/		
1	/		2		/		
2	/		3		/		
3	4		/		/		
*4	/		/		/		

```
$ python REtoNFA.py
```

Enter Regular Expression: aab+bbb

	D		a		b		e
>0	1		4		/		
1	2		/		/		
2	/		3		/		
*3	/		/		/		
4	/		5		/		
5	/		3		/		

```
$ python REtoNFA.py
```

Enter Regular Expression: a\*+b\*

	D		a		b		e
>0	0		0		1		
*1	/		/		/		

## Regular Expression Format.

Format: Any combination of accepted characters.

String	Represents
$\alpha\beta$	$\alpha$ then $\beta$
$\alpha+\beta$	$\alpha$ or $\beta$
$\alpha^*$	0 or more $\alpha$ 's
$\alpha(\rho)\beta$	$\rho$ has priority over $\beta$

**Note:** The output dictionary will have all transition values, even it they are not represented in the default print table.

**Accepted Characters:**

Note: This program will only display the transitions for a, b, and  $\epsilon$  in its default printout table.

This program will accept almost all characters (*See non-accepted characters for exceptions*).

Char	Represents
a-z, A-Z, 0-9, $\epsilon$	Literal Character
+	Or operation
*	Kleene Star
( )	Grouping

**Non-Accepted Characters:**

Some characters may break the program. Those characters include but are not limited to:

Char	Represents	Will Cause
\\$	Final Node in return dict	Unwanted Final Node Declaration
&	Epsilon transition in return dict	Unwanted Epsilon Transition
\t	End of RE	Premature Termination of program
$\alpha$ +	any string ending with a plus	Breaking of Program

## How it works:

At its core, this program is a simple case of parsing a string, identifying cases, and performing the proper function for each case.

After the program recieves a string in a proper format from the user, The string gets passed to the `parse()` function, where an end-tag( `\t` ) is appended to the end of the string and then is passed to the `parseRecur()` function.

Once in `parseRecur()` , the string is first checked if it only contains the terminal character, if yes then terminate. Otherwise the string is broken down into 3 pieces: \*  $\alpha$  = First character of string \*  $\beta$  = Rest of string \*  $\nu$  = Next character following  $\alpha$

Now,  $\alpha$  is checked to see if it is in `[+, (, ), *]` , if its one of these characters, perform the neccessary operation (*See math section for details of each operation*). Otherwise, append `{curpose: { $\alpha$ lpha: nextpos}}` to the final dictionary. return `parseRecur()` until the terminating character is found.

Then return the final dictionary of key-value pairs.

## Time Spent

- 108 Hours Total
  - Began working on April 27

- Figured out math on May 7
- Began rudament programming on May 8
- Completed program on May 13

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