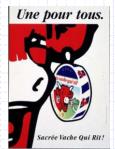
#### 3. Recursion

- · What is Recursion
- · What it is good for and what it is not good for
- What are the characteristics of recursion
- How is a recursive function executed
- How is recursive function call implemented.

#### What is Recursion

- Self referential (defined) in terms of itself)
- The laughing-cow (*la vache qui rit*) package shows a cow wearing laughing-cow packages as earrings, which show a cow wearing laughing-cow packages as earrings which ...



#### Other examples

#### A linked list is:

- a) empty, or
- b) has a head (first element) and a tail, which is a linked list

#### A tree is:

- a) empty, or
- b) has a root, and left and/or right (sub-) trees

#### **Factorial function**

Factorial 5, written 5!, is:  $5\times4\times3\times2\times1$ 

and 6! is

 $6 \times 5 \times 4 \times 3 \times 2 \times 1$ , so  $6 \times 5!$ 

Factorial function, for non-negative integers is:

a) 0! = 1

b) if n > 0, then  $n! = n \times (n - 1)!$ 

#### In C

```
if (n == 0) return 1;
  else return (n * factorial(n - 1));
```

Caution: inefficient

#### Useful recursion

- ◆To be useful the recursion must terminate, so there must be at least one non-recursive case
- such as: 0!
- as well as recursive cases.

such as: n \* (n - 1)!

#### Infinite recursion

```
void TellStory(){
printf("%s", "It was a dark and stormy night ");
printf("%s", "and the captain said to the mate ");
printf("%s", ": Tell us a story mate' ");
printf("%s", " and this is the story he told ...");
TellStory();
}
```

#### **Recursive Programming**

- Consider the problem of computing the sum of all the integers between 1 and any positive integer N
- This problem can be recursively defined as:

$$\begin{split} \sum_{i=1}^{N} i &= N + \sum_{i=1}^{N-1} i &= N + N - 1 + \sum_{i=1}^{N-2} i \\ &= N + N - 1 + N - 2 + \sum_{i=1}^{N-3} i \\ &M \end{split}$$

Recursion

#### **Recursive Programming**

#### **Recursive Programming**

- Note that just because we can use recursion to solve a problem, doesn't mean we should
- For instance, we usually would not use recursion to solve the sum of 1 to N problem, because the iterative version is easier to understand and more efficient
- However, for some problems, recursion provides an elegant solution, often cleaner than an iterative version
- You must carefully decide whether recursion is the correct technique for any problem

cursion

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#### **Indirect Recursion**

- A function invoking itself is considered to be direct recursion
- A function could invoke another function, which invokes another, etc., until eventually the original function is invoked again
- For example, function f1 could invoke f2, which invokes f3, which in turn invokes f1 again
- This is called indirect recursion, and requires all the same care as direct recursion
- ◆ It is often more difficult to trace and debug

cursion 1

#### Length of a list

- a) the length of an empty list is 0
- b) the length of a (non-empty) list is:
  - 1 + the length of the tail of the list

Recursion

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#### Length of a list in C

```
int length_v1 (node* p){ /* iteration */
    int countNodes = 0;
    while (p) do {
        countNodes++;
        p = p->next
    }
    return countNodes;
}

int length_v2(node* p){ /* recursion */
    if (p) return (1 + length_v2(p->next));
    else return 0;
}
```

#### Traversing a list: iterative

Traversing a (singly) linked list *iteratively* in the forward direction is easy:

```
void traverse (node* p){
  while (p){
    process(p->data); /* assume a process function */
    p = p->next;
  }
```

Traversing iteratively in the backward direction is **hard** (no pointers, so need to *stack* return pointers)

## Traversing a list: recursive, forward

Traversing a (singly) linked list *recursively* in the forward direction is easy:

```
void traverse (node* p){
   if (p){
      process(p->data);
      traverse(p->next);
   }
}
```

## Traversing a list: recursive, backward

Traversing a (singly) linked list recursively in the **backward** direction is also easy:

```
void reverseTraverse (node* p){
  if (p){
    reverseTraverse(p->next);
    process(p->data);
  }
}
```

#### How recursion works

- •When a function is called, its parameters, local variables and return address are stacked on the function-call stack
- Nested calls lead to deeper stacking.
- A call of a function to itself is just another nested call.

Recursion 17

#### When not to use recursion

- Don't use a recursive approach when a simple iterative approach is available
- Examples: searching, traversing and inserting in a list is easy to do iteratively
- Traversing a list backwards ('backtracking') is easy to do recursively but hard to do iteratively.

Recursion 18

```
When not to use recursion: example

Eibonacci Numbers:

fib<sub>0</sub> = 0

fib<sub>1</sub> = 1

fib<sub>n</sub> = fib<sub>n-1</sub> + fib<sub>n-2</sub>, for n > 0

int fib(n: integer){ /* doubly recursive */

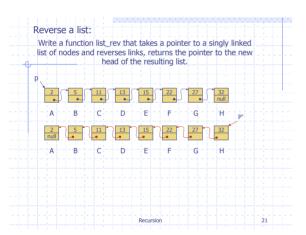
if (n == 0) return 0;
else if (n ==1) return 1;
return (fib(n - 1) + fib (n - 2));
}

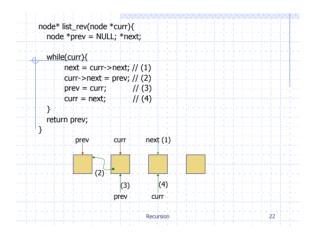
Very inefficient: values repeatedly calculated, then
'forgotten'
```

```
A Better way:

int fib (n: integer){ / *iterative */
    int i, x, y, z;

    i = 1; x = 1; y = 0;
    while (i != n) {
        z = x;
        i++;
        x = x + y;
        y = z;
    }
    return x;
}
```





```
node *list_rev_recursion(node *curr, node *prev) {
node *revHead;

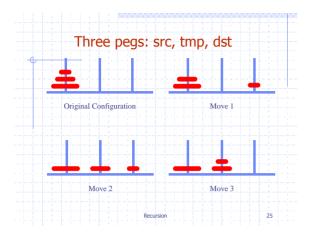
if (curr == NULL)
    revHead = prev;
    else {
        revHead = list_rev_recursion(curr->next, curr);
        curr->next = prev;
    }
    return revHead;
}

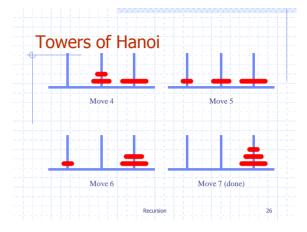
Initial method call should be
head = list_rev_recursion(head, NULL)
```

#### Towers of Hanoi

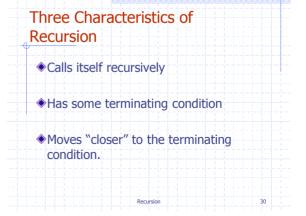
- The Towers of Hanoi is a puzzle made up of three vertical pegs and several disks that slide on the pegs
- The disks are of varying size, initially placed on one peg with the largest disk on the bottom with increasingly smaller ones on top
- The goal is to move all of the disks from one peg to another under the following rules:
  - Only one disk can be moved at a time
  - A bigger disk can never be placed on top of a smaller one

Recursion





# Towers of Hanoi An iterative solution to the Towers of Hanoi is quite complex A recursive solution is much shorter and more elegant if (n == 1) { (move one disk directly from src to dst) } else { (move a tower of n-1 disks from src to tmp) (move a tower of n-1 disk from tmp to dst) } Recursion 27



## Two Flavors of Recursion if (terminating condition) { do final actions } else { move one step closer to terminating condition recursive call(s) } -or if (!(terminating condition)) { move one step closer to terminating condition recursive call(s) }

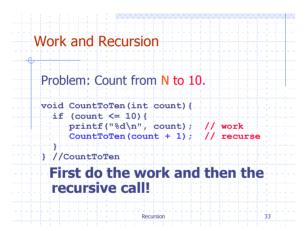
```
Tracing The Recursion

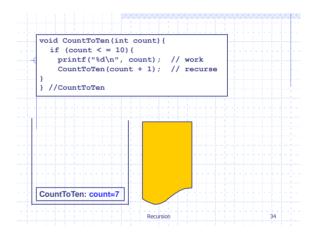
To keep track of recursive execution, do what a computer does: maintain information on an activation stack.

Each stack frame contains:

• Module identifier and variables
• Any unfinished business

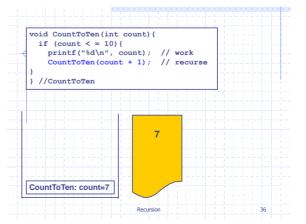
ModuleID: Data values Unfinished business
```





```
void CountToTen(int count){
  if (count < = 10) {
    printf("%d\n", count); // work
    CountToTen(count + 1); // recurse
}
} //CountToTen

CountToTen: count=7</pre>
```



```
void CountToTen(int count) {
   if (count <= 10) {
      printf("%d\n", count); // work
      CountToTen(count + 1); // recurse
   }
} //CountToTen

CountToTen: count=8
CountToTen: count=7</pre>
Recursion 37
```

