

AVL Trees

CSCI-2270

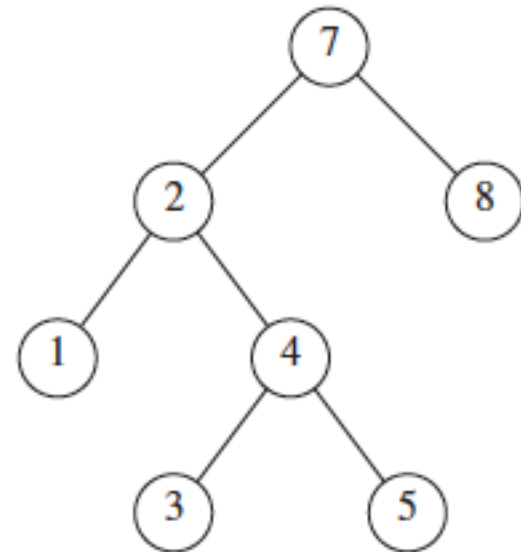
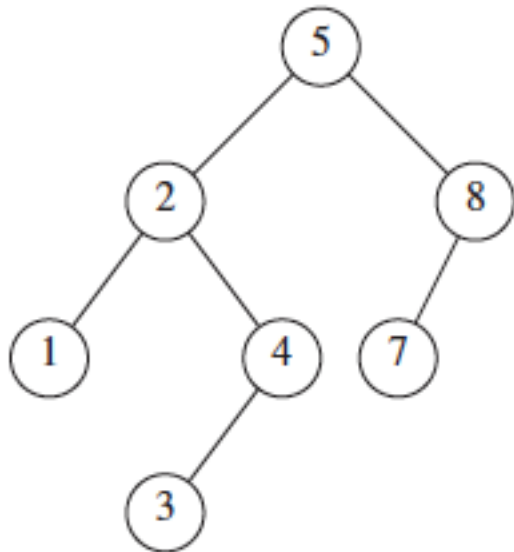
Elizabeth Boese



AVL

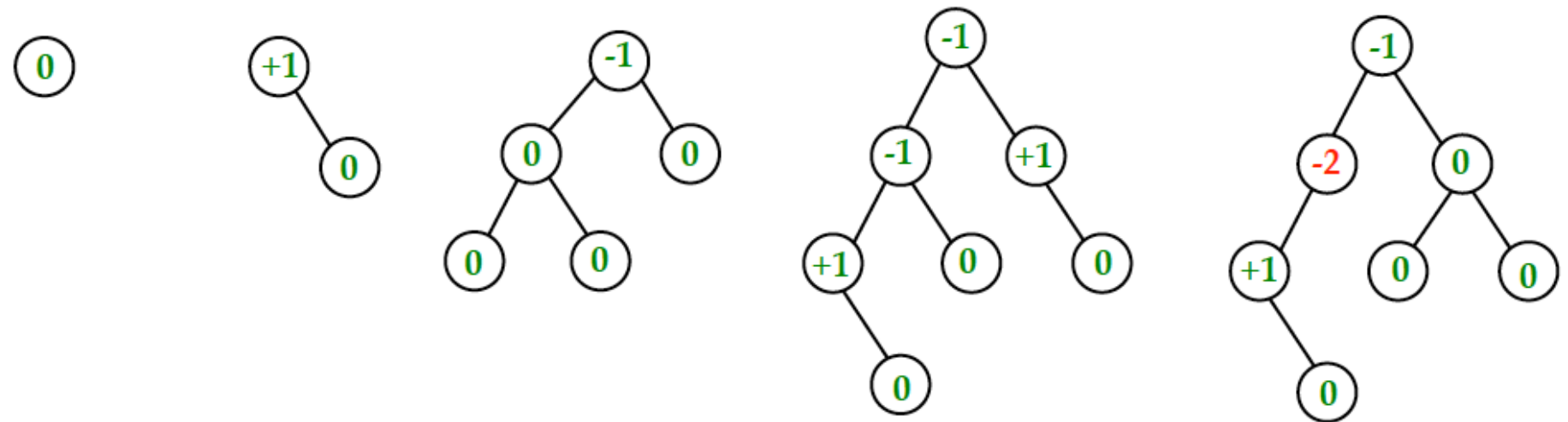
AVL (Adelson-Velskii and Landis)

- BST with balance
 - identical to a binary search tree, except that for every node in the tree, the height of the left and right subtrees can differ by at most 1.
- Which of the following is a binary tree? AVL tree?



AVL

- A binary tree is an AVL tree if $\text{balance}(u) \in \{-1, 0, +1\}$ for every node u
- I.e. the heights of $\text{LEFT}(u)$ and $\text{RIGHT}(u)$ are “about the same” for every node u .
- $\text{balance}(u) := \text{right_height}(u) - \text{left_height}(u)$



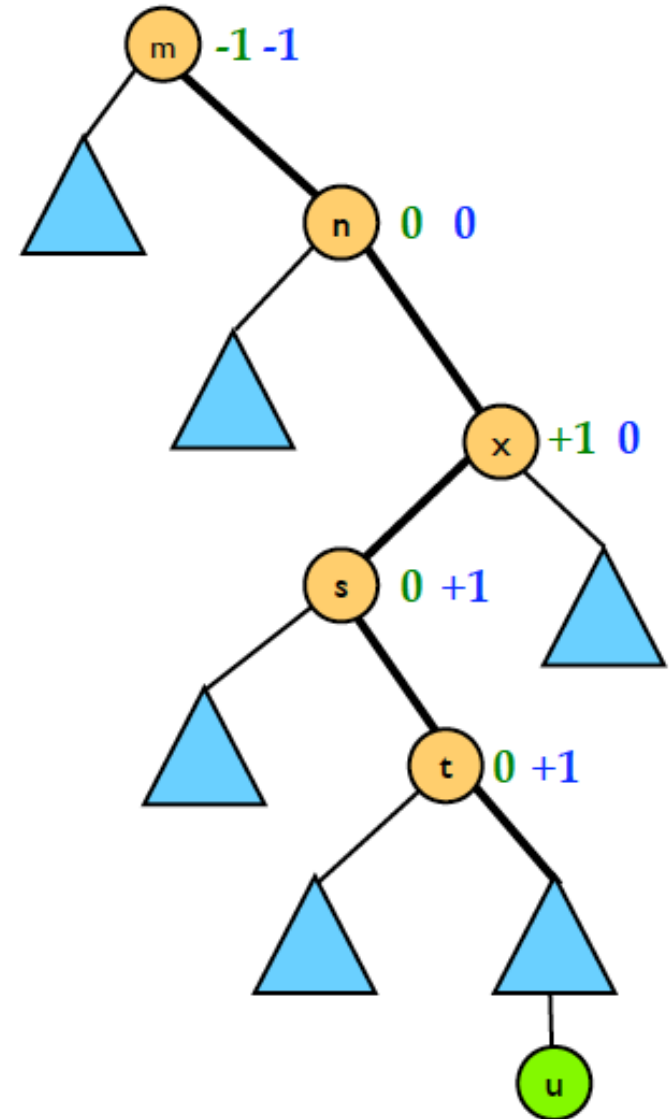
NOT an AVL tree

AVL Trees

- AVL tree with n nodes has height $O(\log n)$.
 - search will run in $O(\log n)$ time if AVL has binary search tree property.
- **insert, delete** can be implemented in $O(\log n)$ time.
- Good structure to implement dictionary or sorted set ADTs.

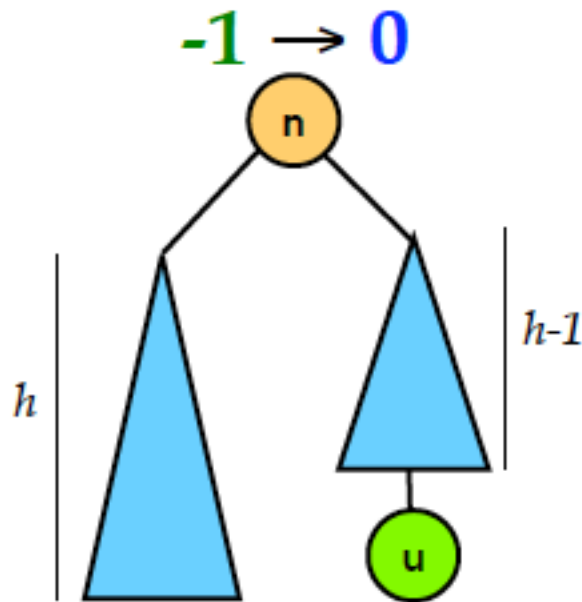
AVL Insert

1. First, do a standard BST insert: do a find and add node where you “fall off the tree.”
2. Walk insertion path back up to root, updating balances.
3. If node was added to the left subtree, *decrement* balance by 1, otherwise *increment* balance by 1. Stop when node’s height doesn’t change.
4. If a balance becomes +2 or -2, fix it.

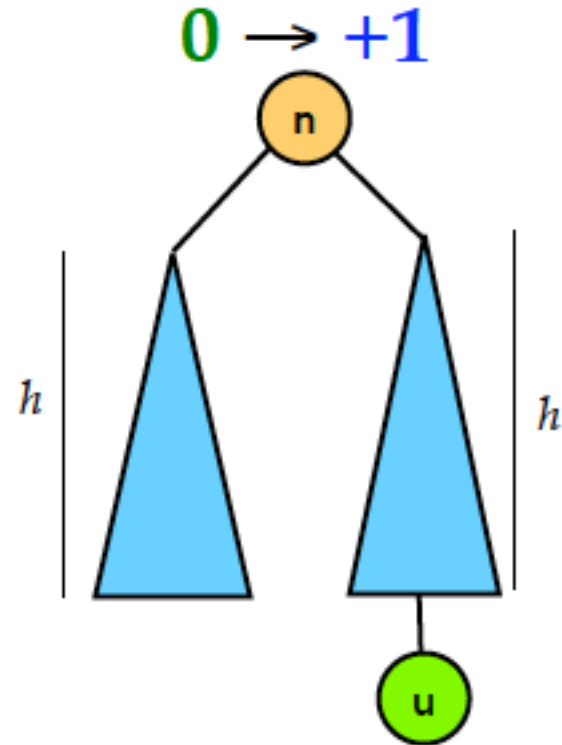


AVL Insert

Easy Cases



Node was added to
the shorter subtree

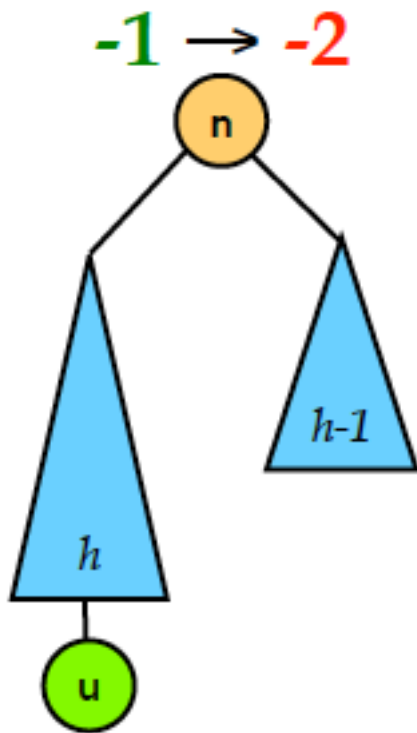


Subtrees were equal,
now slightly unbalanced

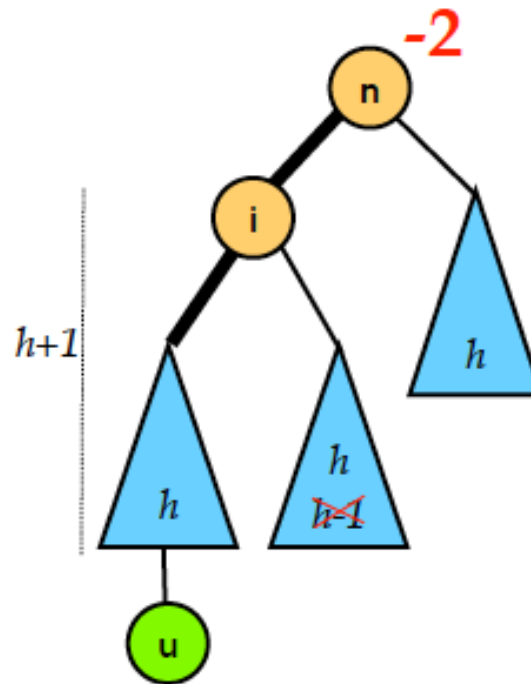
AVL Insert

What to do?

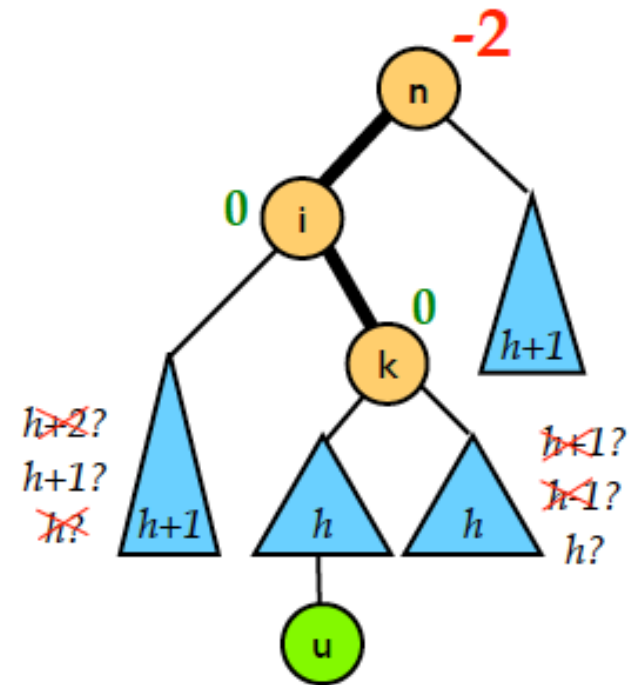
Two cases:



Suppose n is the lowest node that would become -2



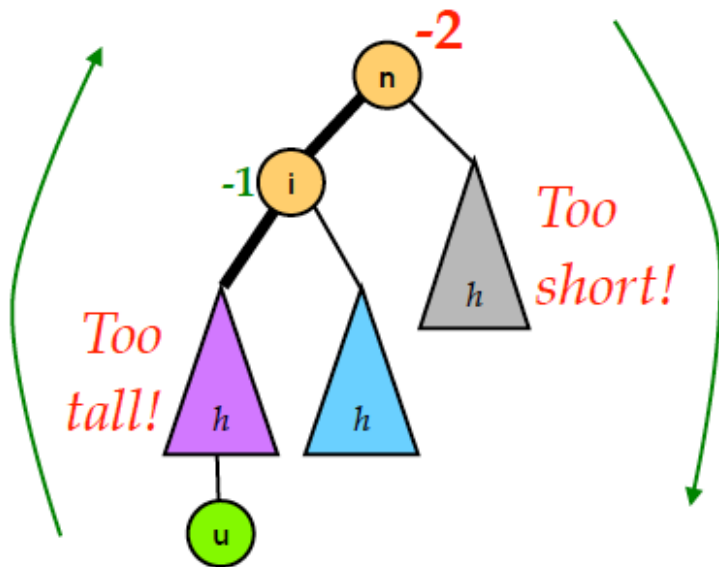
Left, Left



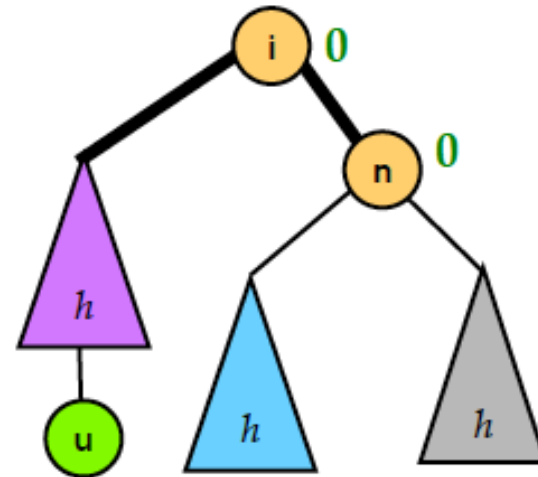
Left, Right


AVL Insert

Left, Left Case



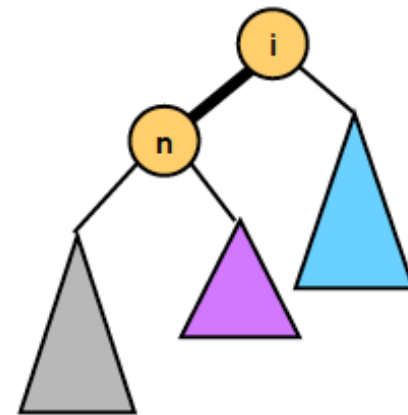
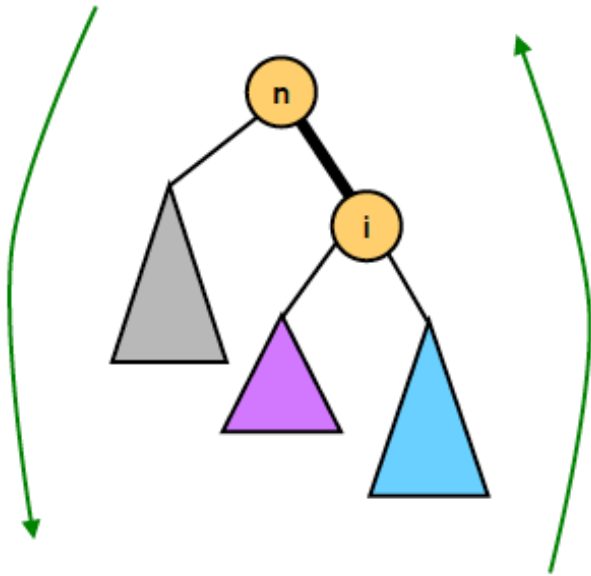
Right rotation
(aka clockwise rotation)



Why does  obey BST ordering?

AVL Insert

Symmetric Left Rotation

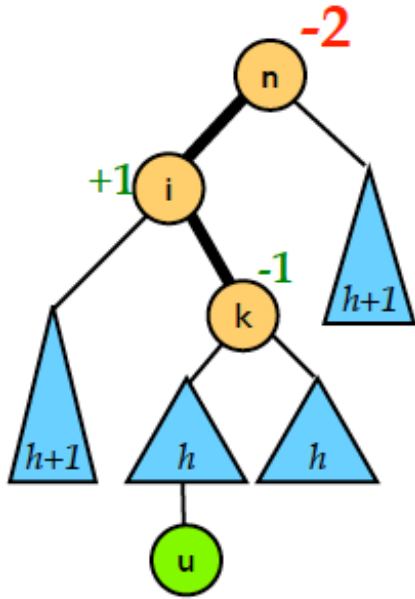


Left rotation
(aka counterclockwise rotation)

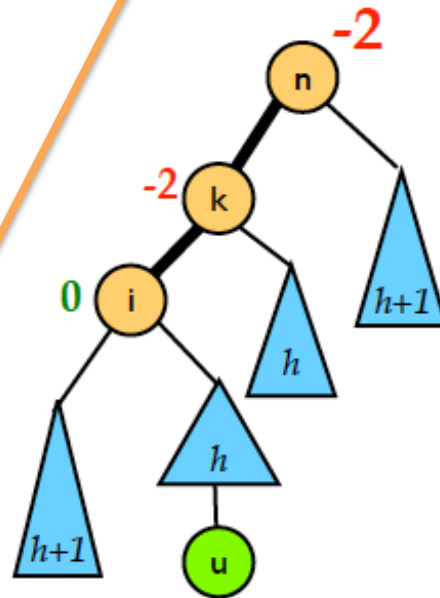
Only a constant # of pointers need to be updated for a rotation: $O(1)$ time

AVL Insert

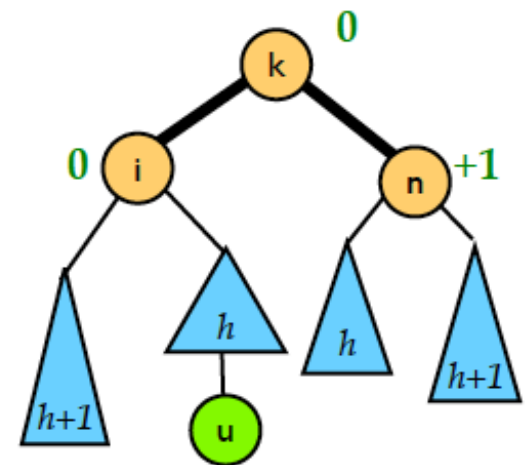
Left, Right Case



Left, Right



(1) Left rotation at i

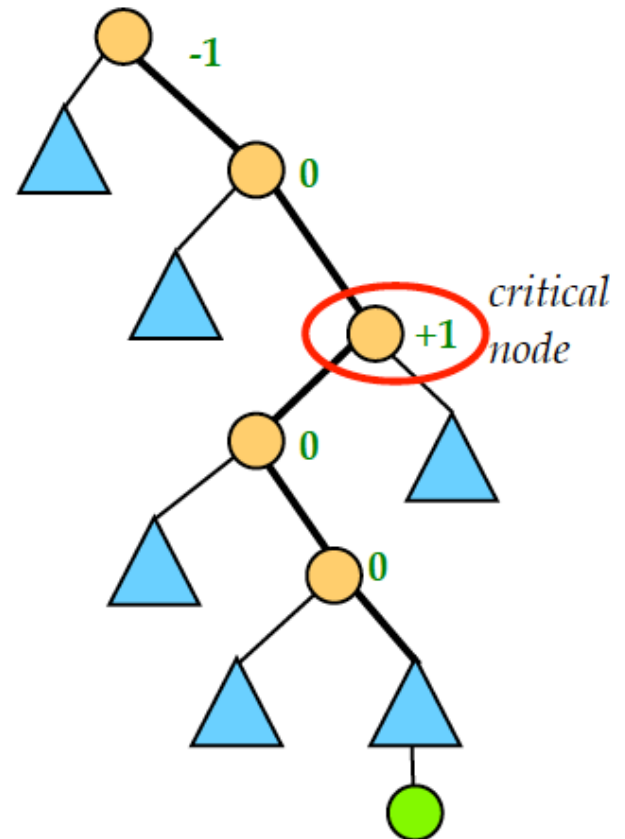


*(2) Then right rotation
at n*

AVL Insert

- Critical Node
 - node on the insertion path closest to the leaves with balance $\neq 0$

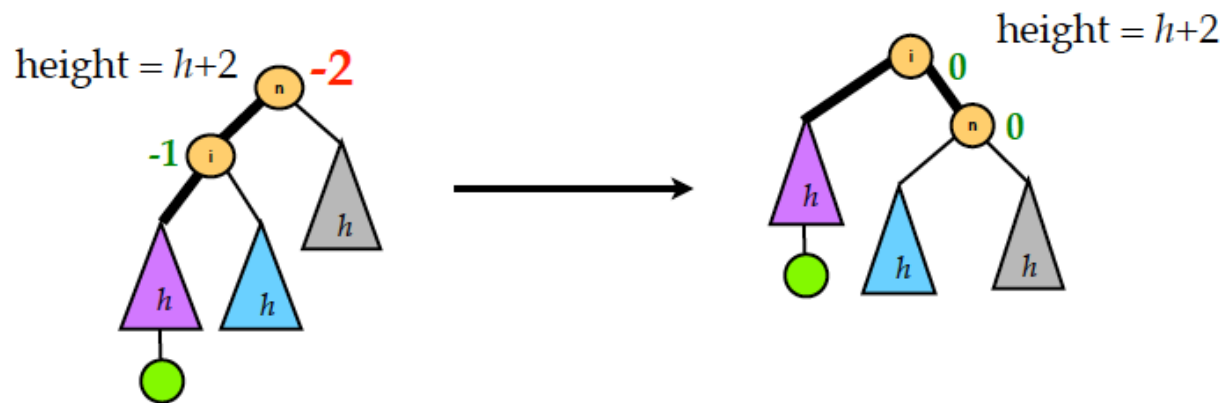
- Rotations leave subtree rooted at critical node balanced with unchanged height.



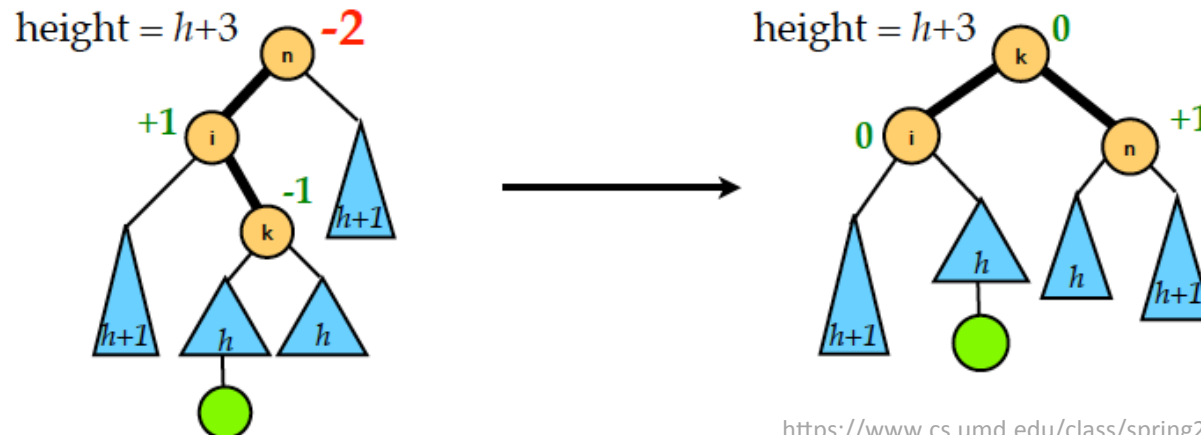
AVL Insert

- Rotations preserve height of critical subtree

Left, Left Case:



Left, Right Case:



AVL Insert

- Because height of critical subtree doesn't change, it can't effect the balance of any nodes higher up in the tree.
- We can stop processing once we process the critical node.
- Therefore, only one rotation will occur.
- Optimization:
 - on first pass down the tree to insert a node, remember the critical node (last node with non-zero balance)
 - Then, to adjust balances, start at critical node and rewalk the path down to inserted node.

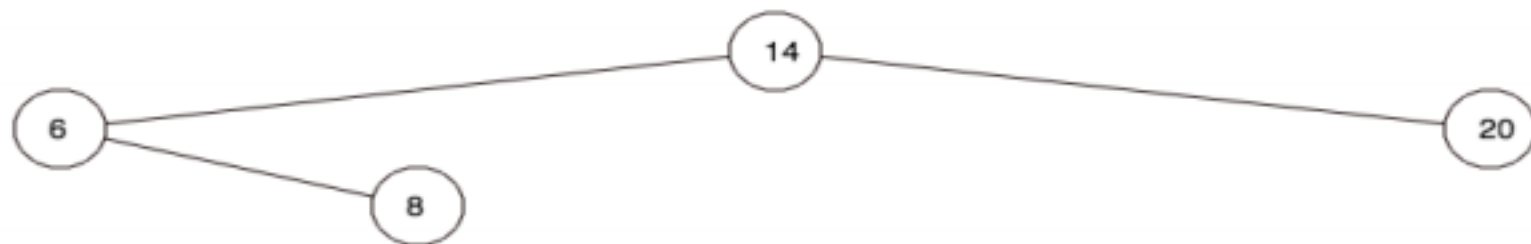
AVL Trees

- Nice Features:
 - Worst case $O(\log n)$ performance guarantee
 - Fairly simple to implement
- Problems:
 - Have to maintain extra balance factor storage at each node.

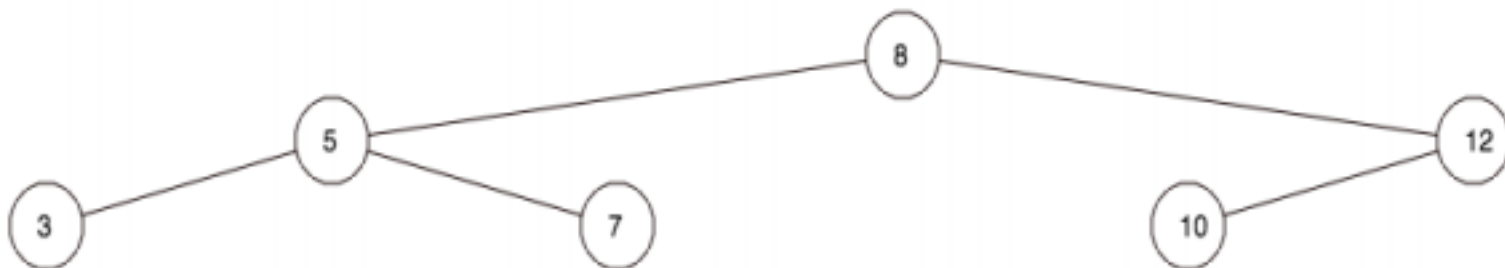


QUESTIONS TO PONDER

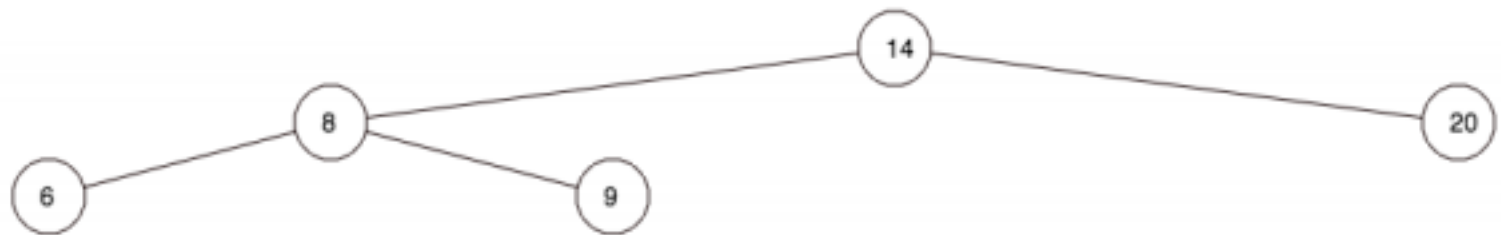
1. [2 Points] If 10 is inserted into the following AVL tree, draw the resulting tree.



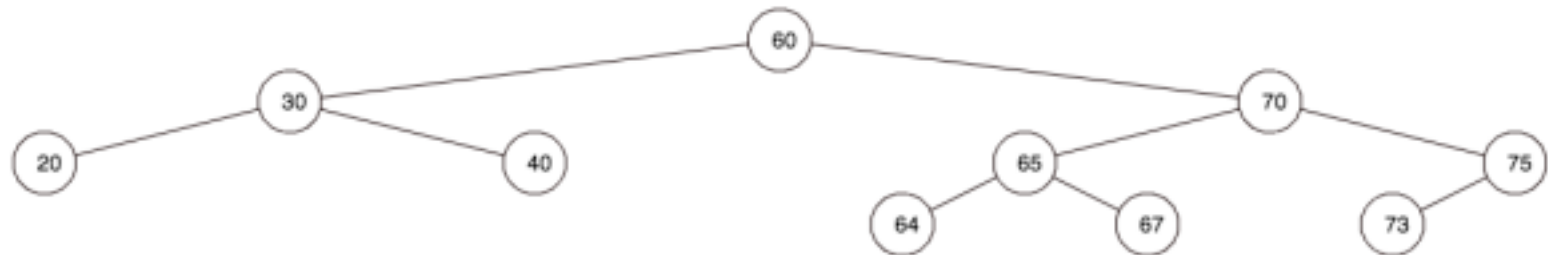
2. [2 Points] If 14 is inserted into the following AVL tree, draw the resulting tree.



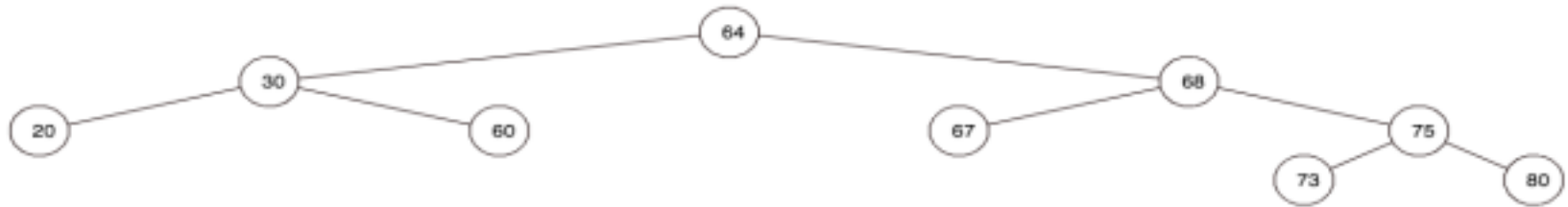
3. [2 Points] If 12 is inserted into the following AVL tree, draw the resulting tree.



4. [2 Points] If 68 is inserted into the following AVL tree, draw the resulting tree.

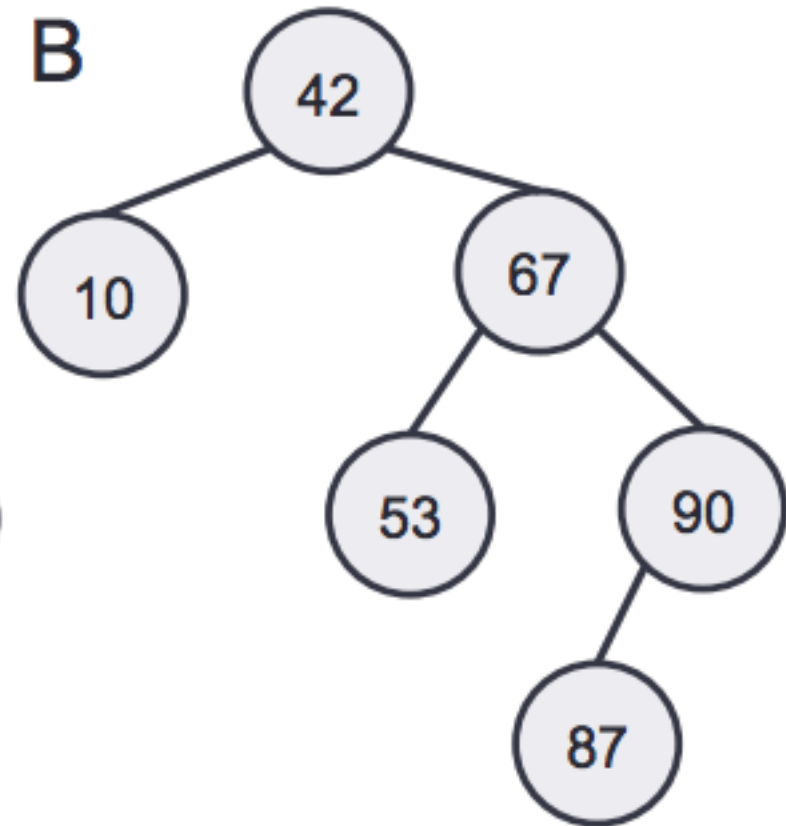
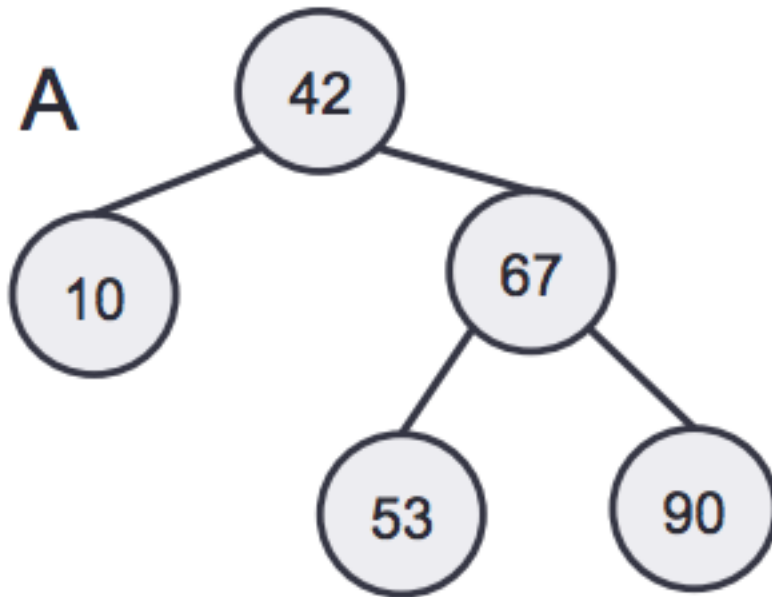


5. [2 Points] If 90 is inserted into the following AVL tree, draw the resulting tree.



Questions to Ponder

6. Which of the following is a balanced AVL tree?



Questions to Ponder

7.

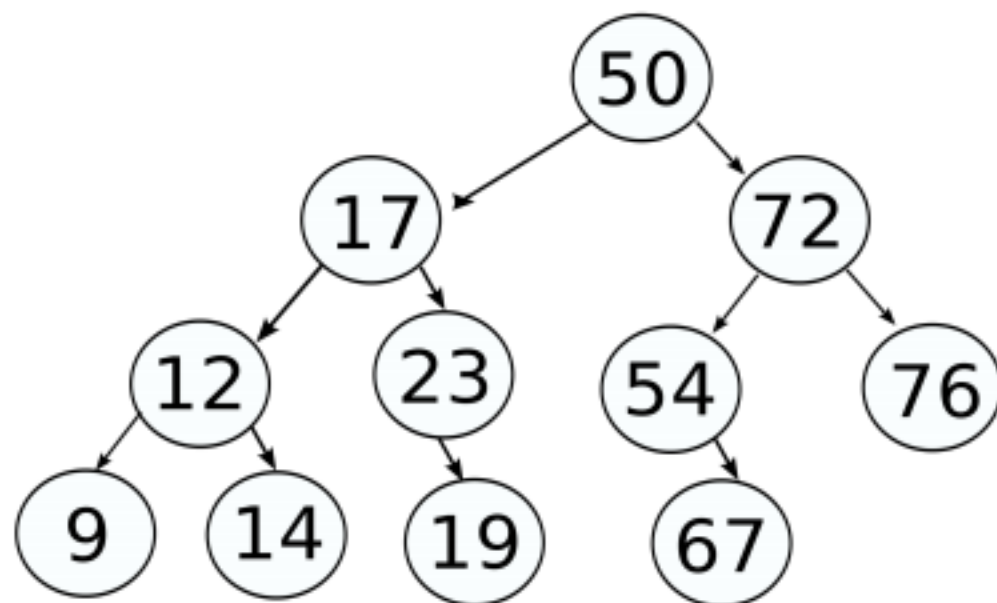
- (e) T F An AVL tree is balanced, therefore a median of all elements in the tree is always at the root or one of its two children.

Explain:

<https://courses.csail.mit.edu/6.006/oldquizzes/solutions/q1-f2008-sol.pdf>

8. What are the two invariants that an AVL tree must hold?

5. Given this tree structure:



- (a) Assuming the tree is a binary search tree, and not an AVL tree, draw the tree structure created by the following code. Hint: Don't try to do it all at once in your head. Draw each insertion/deletion as an entirely separate tree.

```
t->insert(71);  
t->delete(19);  
t->delete(23);  
t->delete(50);  
t->insert(77);  
t->insert(78);  
t->insert(79);  
t->insert(80);
```

- (b) Repeat part (a) assuming the tree is an AVL tree.