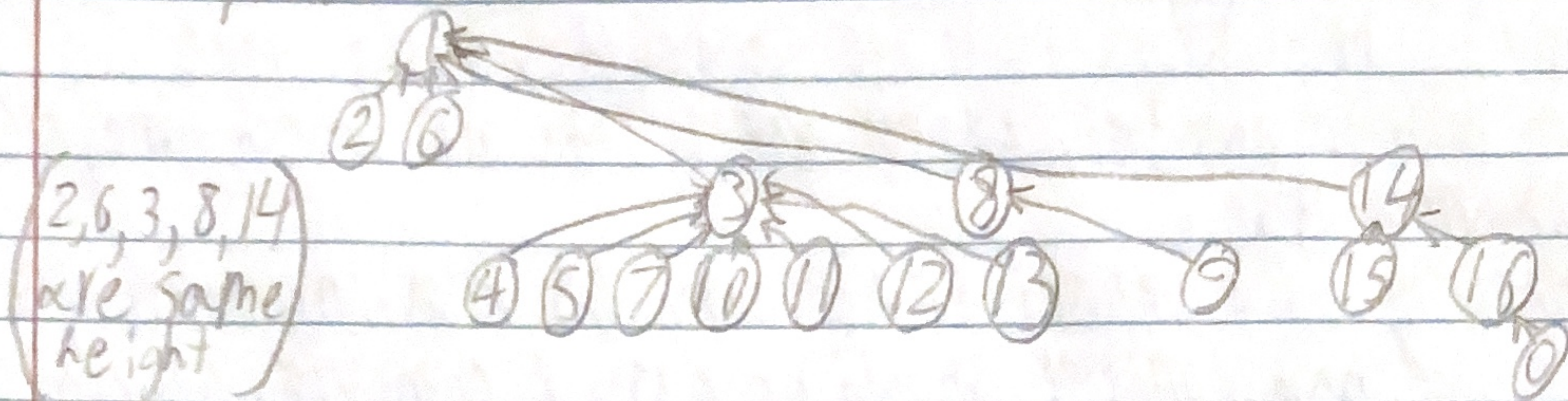
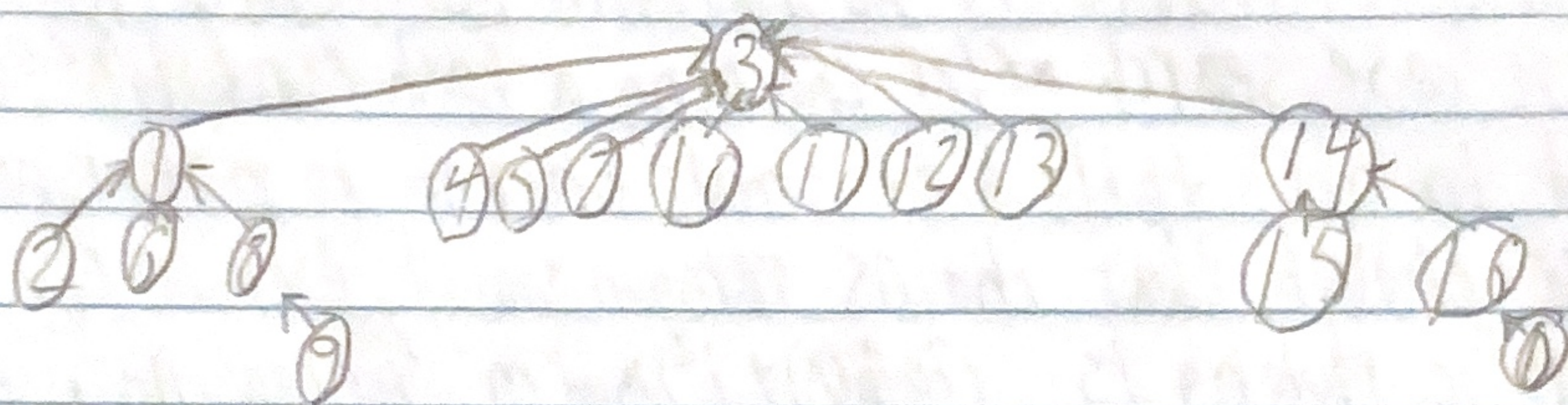


8.1) performing the sequence arbitrarily and by height both yield:

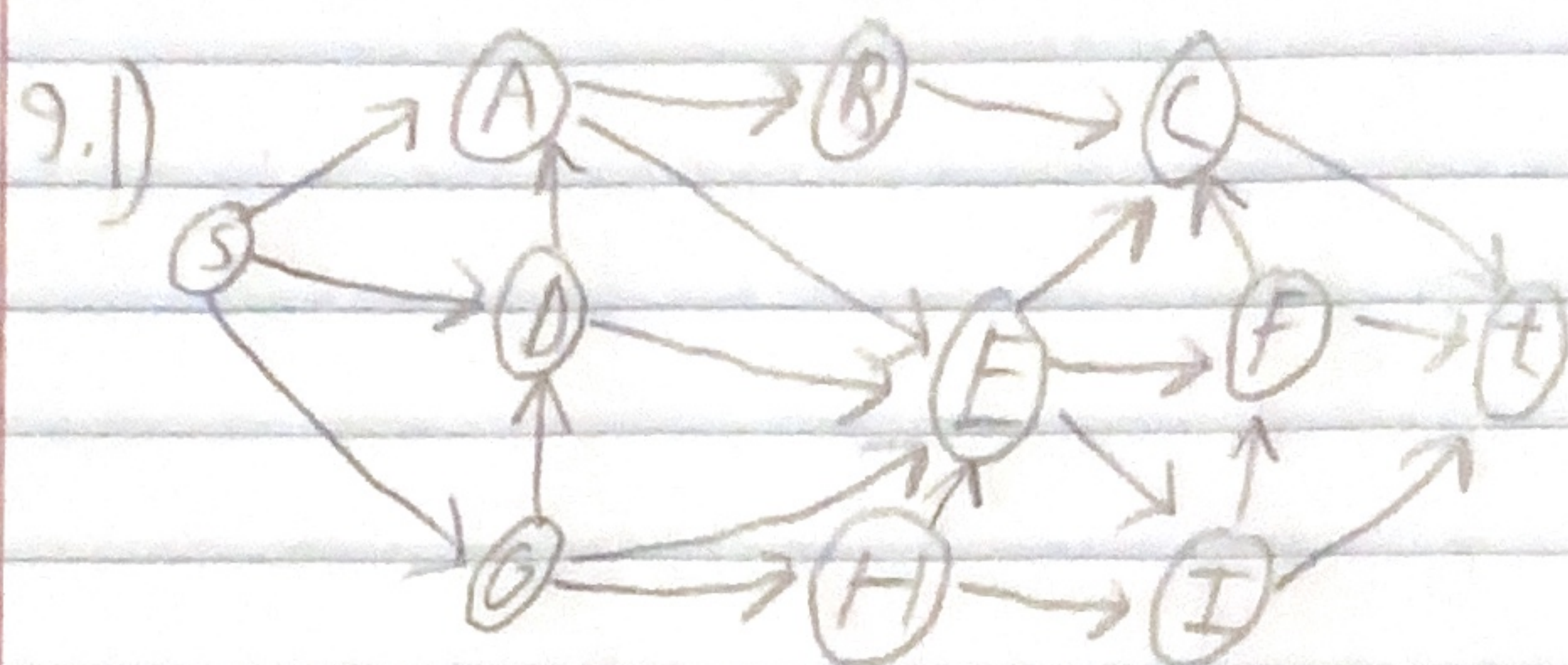


Performing the sequence by size yields:





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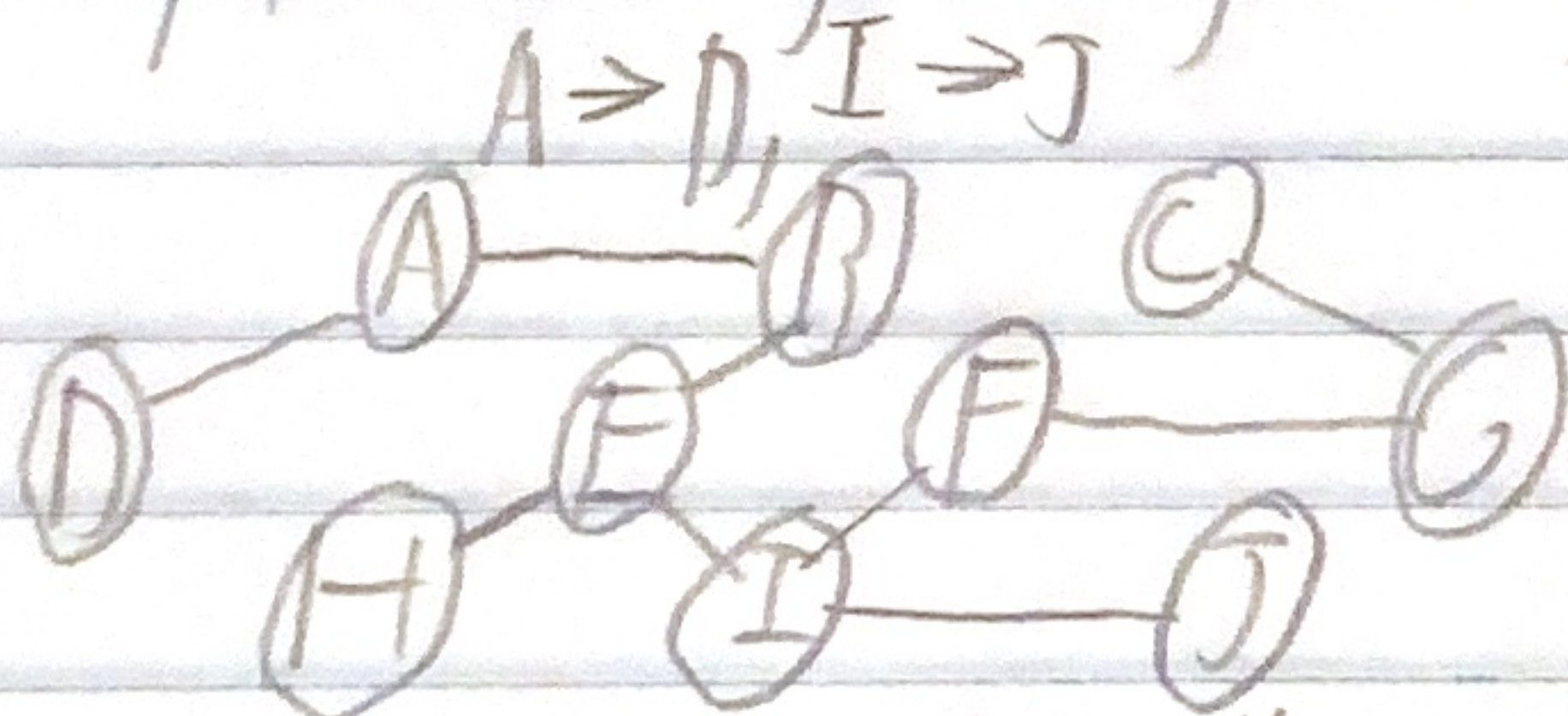


Topological ordering: s, G, D, H, A, B, E, I, F, C, t

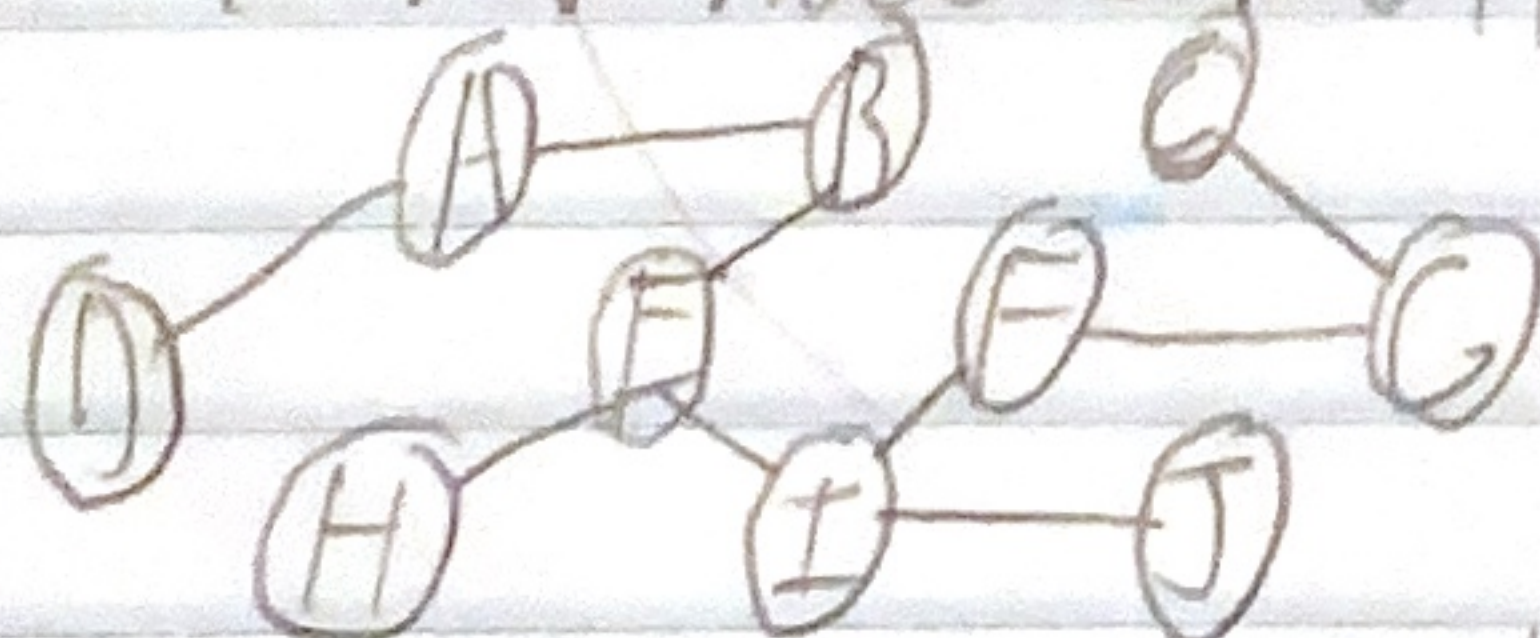
9.5 a) shortest weighted paths from A:  $A \rightarrow B, A \rightarrow C, A \rightarrow B \rightarrow G, A \rightarrow B \rightarrow G \rightarrow E, A \rightarrow B \rightarrow G \rightarrow E \rightarrow F, A \rightarrow C \rightarrow D$

b) shortest unweighted paths from B:  $B \rightarrow G, B \rightarrow E, B \rightarrow C, B \rightarrow E \rightarrow F, B \rightarrow C \rightarrow D, B \rightarrow C \rightarrow D \rightarrow A$

9.15 a) i. Using Kruskal's, add edges of smallest weight without creating a cycle:  $C \rightarrow G, E \rightarrow I, F \rightarrow G, E \rightarrow B, E \rightarrow H, F \rightarrow I, B \rightarrow A, A \rightarrow D, I \rightarrow J$



ii. using Prim's algorithm, start at node and branch off to least weighted edge without creating cycles:  $C \rightarrow G, G \rightarrow F, F \rightarrow I, E \rightarrow I, E \rightarrow H, E \rightarrow B, B \rightarrow A, I \rightarrow J, A \rightarrow D$ . If new node can offer a better path than before, take it.



b) This solution is not unique;  $B \rightarrow F$  could be used instead of  $F \rightarrow I$  and the spanning network would remain the same length.