

# Sorting Review

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Java

# Big O of Sorting Algorithms

Algorithm	Time Complexity (Best)	Time Complexity (Average)	Time Complexity (Worst)	Space Complexity
Bubble Sort	$O(n)$	$O(n^2)$	$O(n^2)$	$O(1)$
Insertion Sort	$O(n)$	$O(n^2)$	$O(n^2)$	$O(1)$
Selection Sort	$O(n^2)$	$O(n^2)$	$O(n^2)$	$O(1)$

# Examples of Runtime

Here are the running times of some operations we might perform on the phone book, from fastest to slowest:

- **$O(1)$  (in the worst case):** Given the page that a business's name is on and the business name, find the phone number.
- **$O(1)$  (in the average case):** Given the page that a person's name is on and their name, find the phone number.
- **$O(\log n)$ :** Given a person's name, find the phone number by picking a random point about halfway through the part of the book you haven't searched yet, then checking to see whether the person's name is at that point. Then repeat the process about halfway through the part of the book where the person's name lies. (This is a binary search for a person's name.)
- **$O(n)$ :** Find all people whose phone numbers contain the digit "5".
- **$O(n)$ :** Given a phone number, find the person or business with that number.
- **$O(n \log n)$ :** There was a mix-up at the printer's office, and our phone book had all its pages inserted in a random order. Fix the ordering so that it's correct by looking at the first name on each page and then putting that page in the appropriate spot in a new, empty phone book.

# Examples of Runtime

For the below examples, we're now at the printer's office. Phone books are waiting to be mailed to each resident or business, and there's a sticker on each phone book identifying where it should be mailed to. Every person or business gets one phone book.

- **$O(n \log n)$ :** We want to personalize the phone book, so we're going to find each person or business's name in their designated copy, then circle their name in the book and write a short thank-you note for their patronage.
- **$O(n^2)$ :** A mistake occurred at the office, and every entry in each of the phone books has an extra "0" at the end of the phone number. Take some white-out and remove each zero.
- **$O(n \cdot n!)$ :** We're ready to load the phonebooks onto the shipping dock. Unfortunately, the robot that was supposed to load the books has gone haywire: it's putting the books onto the truck in a random order! Even worse, it loads all the books onto the truck, then checks to see if they're in the right order, and if not, it unloads them and starts over. (This is the dreaded [bogo sort](#).)
- **$O(n^n)$ :** You fix the robot so that it's loading things correctly. The next day, one of your co-workers plays a prank on you and wires the loading dock robot to the automated printing systems. Every time the robot goes to load an original book, the factory printer makes a duplicate run of all the phonebooks! Fortunately, the robot's bug-detection systems are sophisticated enough that the robot doesn't try printing even more copies when it encounters a duplicate book for loading, but it still has to load every original and duplicate book that's been printed.

[Link to thread](#)

# Examples of Runtime

$O(\log N)$  basically means time goes up linearly while the  $n$  goes up exponentially. So if it takes 1 second to compute 10 elements, it will take 2 seconds to compute 100 elements, 3 seconds to compute 1000 elements, and so on.

It is  $O(\log n)$  when we do divide and conquer type of algorithms e.g binary search. Another example is quick sort where each time we divide the array into two parts and each time it takes  $O(N)$  time to find a pivot element. Hence it  $N O(\log N)$

[Link to thread](#)

## Non-Recursive Merge Sort

# Merge Sort Pseudo Code

```
void mergeSort(int A[], int first, int last)  
{  
    // find middle index of A  
    // sort the first half of A  
    // sort the second half of A  
    // merge the first and second halves of A  
}
```

**Let's look at merge...**

**/\***

**precondition: lists A and B are sorted in  
non-decreasing order**

**postcondition: list C contains all the  
values from lists A and  
B in nondecreasing  
order**

**\*/**

**void merge(int A[], int B[], int C[])**

List A

3	11	17	19	24	29	31	37
---	----	----	----	----	----	----	----

List B

1	4	5	15	18	25	27	36
---	---	---	----	----	----	----	----

How many elements does  
List C have?



# Pseudo Code for Merge

- A) List A is done, get value from List B**
- B) List B is done, get value from List A**
- C) Neither is done, if List A[i] < B[k],  
then get value from List A**
- D) Neither is done, if List B[k] <= List A[i]  
then get value from List B**

List A

3	11	17	19	24	29	31	37
---	----	----	----	----	----	----	----

List B

1	4	5	15	18	25	27	36
---	---	---	----	----	----	----	----

Let's count which rules  
we use...

# Merge Sort Pseudo Code

```
void mergeSort(int A[], int first, int last)  
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    // find middle index of A  
    // sort the first half of A  
    // sort the second half of A  
    // merge the first and second halves of A  
}
```

Complete the MergeSort Lab folder

Create the Merge method that merges 2 arrays into 1.