Graph problems!

Time-outs:

- Python: 5 seconds on Q1 through Q4, 10 seconds on Q5
- C++ (w/g++ -O3 flag): 2 seconds on Q1 through Q4, 4 seconds on Q5

You may solve the programming problems in either Python or C++. To use Python, name the file containing your code main.py. To use C++, name the file main.cpp. You are encouraged to test your code locally.

At any time, you can choose to run the tests that will be used to grade your solution. To do so, navigate to the Gradescope assignment for that particular homework question and upload **only** main.py or main.cpp. Gradescope will tell you how many test cases passed, and what error occurred on the ones that failed.

When you run the tests on Gradescope, debugging output will be printed to tell you which test cases you passed, which ones you failed, and what your total score currently is for that problem. If a test case failed, the output will specify a failure mode (e.g. timeout, compiler error, wrong answer, etc.).

The first few test cases for each problem are visible test cases. This means that if you fail them, Gradescope will tell you the the expected output and the output your code actually gave. Note that these outputs may not be listed in order. The rest of the test cases for each problem are hidden test cases, which means you will not be shown the input or expected output.

All assignments MUST be manually submitted to Gradescope in order to complete submission! There is no auto-submit feature enabled for any Gradescope coursework. There will be no exemptions granted if you forget to manually submit your assignment. If time permits, you can contact the course staff and they can check on the Gradescope platform and confirm your assignment submission status.

1. Establishing Communications

(100 pts)

Description

As each team of explorers from Oganesson Dynamics arrives at their destination, the first thing they want to do is setup a communications base. But where should they put it?

To decide, they build a graph of the area. Each node represents a work site, and each edge indicates two sites that are close enough to communicate if one of them has the communications base. In fact, pairs of sites can relay through the base to communicate, as long as they are both connected to it. It can be slow, but it works well.

You must determine where to put the base such that the MOST pairs of sites can communicate.

Input Format

- The first line contains N, the number of sites being analyzed, and M, the number of pairs (edges) that could potentially communicate.
- The next *M* lines each have two values identifying all pairs of sites that could communicate with a base station at one of them.

Constraints

- $1 \le N \le 200$
- $1 \le M \le 10,000$
- $0 \le \text{site IDs} < N$

Output Format

Output the count of how many PAIRS of sites would be able to communicate if the communications base were placed in the optimal position.

Sample Input

5 7

0 1

0 2

0 3

1 2

1 3

1 4

2 3

Sample Output

10

Explanation

If a communications node is placed at size 1, it will connect to 0, 2, 3, and 4.

This means that there are 10 pairs that can communicate. Node 1 can directly communicate with each of the four others PLUS it can relay communications between (0,2), (0,3), (0,4), (2,3), (2,4), and (3,4).

Sample Input

- 10 15
- 0 1
- 0 2
- 0 3
- 0 4
- 1 2
- 1 5
- 2 3
- 2 5
- 3 4
- 3 5
- 3 6
- 3 7
- 3 9
- 6 9
- 8 9

Sample Output

Sample Input

- 10 20
- 0 3
- 0 5
- 0 6
- 0 9
- 1 4
- 1 6
- 1 8
- 1 9
- 2 3
- 2 4
- 2 5
- 2 7
- 2 8
- 3 6
- 5 0
- 3 9
- 4 6
- 5 6
- 5 8
- 6 9 7 8

Sample Output

2. Expanding Communications

(100 pts)

Description

Some teams from Oganesson dynamics were sent with two communications bases. These bases must be connected to each other by an edge, but have the ability to relay messages through one another, connecting any of their neighbors and broadly expanding the range covered.

You must determine where to put the two bases such that the most pairs of sites can communicate.

Input Format

- The first line contains N, the number of sites being analyzed and M, the number of pairs of sites that could potentially communicate, including through related messages.
- The next *M* lines each have two values identifying all pairs of sites that could directly communicate if a base station is at one of them.

Constraints

- $1 \le N \le 200$
- $1 \le M \le 10,000$
- $0 \le \text{site IDs} < N$

Output Format

Output the count of how many PAIRS of sites would be able to communicate if the communications bases were placed in the optimal position.

Sample Input

5 6

0 1

0 2

0 4

1 4

2 3

2 4

Sample Output

10

Explanation

If communications nodes are placed at sites 0 and 2, they can facilitate communication among all 5 sites, this is allowed because sites 0 and 2 are connected to each other. On top of that, site 0 is connected to 1 and 4, while site 2 is connected to 3 and 4.

With fives sites, there are $(5 \times 4)/2 = 10$ pairs.

Sample Input

- 7 15
- 0 3
- 0 4
- 0 5
- 0 6
- 1 2
- 1 3
- 1 5
- 2 3
- 2 4
- 2 6
- 3 4
- 3 5
- 3 6
- 4 6
- 5 6

Sample Output

Sample Input

- 10 20
- 0 3
- 0 5
- 0 7
- 0 8
- 0 9
- 1 4
- 1 6
- 1 7
- 2 3
- 3 4
- 3 5
- 3 6
- 3 9
- 4 6 4 8
- 4 9
- 5 7
- 6 7
- 7 9 8 9

Sample Output

Sample Input

- 25 30
- 0 20
- 1 3
- 1 21
- 1 24
- 2 9
- 2 10
- 2 11
- 2 17
- 3 5
- 3 9
- 4 6
- 4 16
- 4 17
- 4 19
- 6 21
- 7 13
- 8 17 9 19
- 10 21
- 12 13
- 12 19
- 12 20
- 13 17
- 13 18
- 14 20 15 20
- 15 22
- 19 20
- 21 22
- 23 24

Sample Output

3. Analyzing Communication Clusters

(200 pts)

Description

Over time, exploration teams will build more communications bases and set them up at each site. However, there's still a problem! Some sites simply cannot be connected at all, no matter how many bases are used.

You must analyze a connectivity network and identify how many distinct (unconnected) components a graph is divided into.

Input Format

- The first line contains N, the number of sites being analyzed, and M, the number of connections that will exist when all nodes are in place.
- The next M lines each have two values identifying a pair of connected sites.

Constraints

- $1 \le N \le 200$
- $1 \le M \le 10,000$
- $0 \le \text{site IDs} < N$

Output Format

Print the number of independent site clusters in the input graph.

Sample Input

9 10

0 3

0 4

0 5

1 2

1 6

1 8

2 8

3 7

4 7

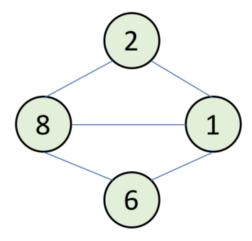
6 8

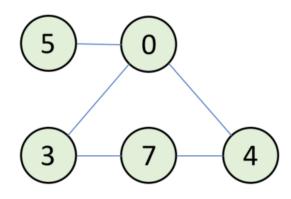
Sample Output

2

Explanation

This graph looks like this:





Sample Input

12 8

0 2

0 9

1 3

1 7

4 6

5 10

6 11

8 10

Sample Output

Sample Input

12 11

0 7

1 10

2 10

3 5

3 7

4 10

6 8

7 8

7 10

9 10

10 11

Sample Output

Sample Input

- 224 318
- 0 75
- 0 113
- 0 215
- 1 65
- 1 95
- 1 138
- 2 15
- 2 168
- 3 106
- 3 118
- 3 212
- 4 5
- 4 101
- 4 169
- 4 198
- 5 101
- 5 113
- 5 169
- 5 198
- 5 170
- 6 16
- 6 18
- 6 72
- 6 167
- 7 23
- 7 80
- 8 43
- 8 145
- 9 53
- 10 38
- 10 82
- 10 158
- 10 183
- 11 50
- 12 166
- 13 15

- 13 66
- 13 160
- 13 191
- 14 48
- 14 88
- 14 142
- 16 105
- 16 127
- 16 164
- 17 128
- 17 182
- 18 72
- 18 129
- 18 164
- 19 54
- 19 122
- 19 159
- 19 196
- 19 199
- 19 207
- 20 42
- 20 138
- 20 218
- 21 133
- 21 161
- 21 189
- 22 84
- 22 140
- 23 200
- 24 45
- 24 195
- 24 214
- 25 28
- 25 41
- 25 130
- 25 173
- 26 58
- 26 79
- 27 63
- 27 165

- 27 216
- 28 96
- 28 112
- 28 219
- 28 222
- 29 163
- 30 32
- 31 47
- 31 137
- 32 160
- 33 53
- 33 135
- 33 182
- 34 43
- 34 55
- 35 67
- 35 180
- 35 204
- 35 221
- 36 80
- 37 56
- 38 71
- 38 183
- 39 218
- 40 48
- 40 88
- 40 115
- 41 152
- 42 218
- 43 194
- 44 63
- 44 103
- 44 157
- 45 163
- 46 118
- 46 212
- 47 210
- 48 147
- 48 220
- 49 93

- 49 144
- 49 212
- 50 210
- 51 71
- 51 94
- 51 121
- 52 122
- 54 102
- 55 117
- 56 154
- 57 170
- 57 182
- 58 86
- 58 98
- 58 110
- 58 176
- 59 221
- 60 179
- 60 181
- 60 203
- 61 97
- 61 136
- 61 163
- 62 86
- 62 111
- 63 103
- 63 162
- 63 165
- 64 153
- 65 95
- 65 141
- 65 218
- 67 68
- 67 116
- 67 171
- 67 204
- 67 221
- 69 74
- 69 100
- 69 114

- 69 133
- 69 134
- 69 201
- 70 170
- 71 158
- 71 217
- 72 127
- 72 167
- 73 203
- 74 133
- 74 134
- 75 101
- 75 198
- 76 116
- 76 171
- 76 180
- 76 204
- 76 221
- 77 96
- 77 112
- 77 219
- 78 83
- 78 89
- 78 181
- 79 86
- 80 153
- 80 200
- 80 206
- 81 195
- 82 94
- 83 125
- 83 181
- 84 184
- 84 193
- 85 131
- 85 135
- 85 170
- 85 182
- 86 98
- 86 110

- 86 111
- 87 214
- 88 151
- 89 125
- 89 203
- 90 117
- 90 194
- 90 197
- 91 163
- 92 205
- 93 106
- 93 185
- 93 212
- 95 141
- 96 173
- 99 183
- 99 223
- 100 114
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- 100 155
- 101 113
- 101 198
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- 109 182
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- 111 176
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- 114 189
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- 115 177
- 118 175
- 119 181
- 120 171
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- 123 193
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- 128 170
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- 133 161 133 189
- 133 201
- 136 163
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- 137 154
- 137 192
- 137 210
- 139 194
- 139 197
- 140 184
- 140 186
- 141 202
- 142 146
- 142 151
- 142 220

- 144 174
- 144 212
- 145 197
- 146 147
- 147 151
- 148 173
- 148 219
- 149 191
- 153 166
- 153 190
- 155 198
- 156 159
- 156 199
- 157 162
- 157 216
- 158 217
- 158 223
- 160 168
- 161 189
- 162 216
- 163 214
- 164 167
- 165 216
- 166 200
- 168 191
- 168 205
- 171 180
- 172 174
- 174 175
- 175 185
- 177 220
- 184 186
- 184 188
- 186 211
- 187 188
- 187 193
- 193 211
- 194 197
- 199 209
- 202 218

204 221 208 214 210 213

Sample Output

4. Pricing Instantaneous Communications

(200 pts)

Description

Great news! Oganesson Dynamics has worked out the technology to provide instantaneous communications between two sites, MUCH faster than the old system. The problem is that it is quite expensive to setup, especially since a sending unit and receiving unit must be built for each site and the price varies depending on how far apart they are and if there are any obstacles in between.

Given a graph where nodes represent sites and edges provide the cost of connecting two sites, determine the cheapest it would be to make sure that ALL of the sites can communicate with one another. (Again, relaying of information is allowed, and much faster this time!)

Input Format

- The first line contains N, the number of sites, and M, the number of possible pairs of sites that could communicate.
- The next *M* lines each have three integer values. The first two identify the two sites involved, while the third provides the cost to setup a communications channel.
- All graphs are guaranteed to be connected.

Constraints

- $1 \le N \le 200$
- $1 \le M \le 10,000$
- $0 \le \cos t$ to setup a communication channel $\le 1,000,000$
- $0 \le \text{site IDs} < N$

Output Format

A single value indicating the minimum total cost to setup communications.

Sample Input

5 6

0 1 5

0 2 10

0 3 2

1 3 3

1 4 8

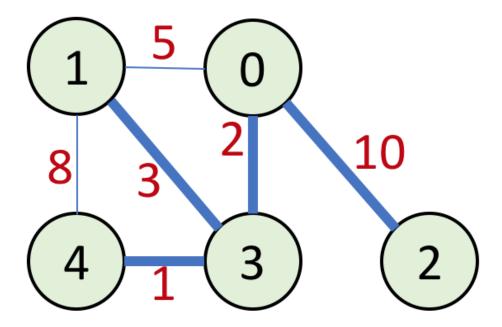
3 4 1

Sample Output

16

Explanation

The graph is as follows, with the minimum-cost communication channels highlighted:



As such, the total cost is 1 + 2 + 3 + 10 = 16.

Sample Input

6 10

0 1 7

0 3 0

0 4 4

1 2 3

1 3 5

1 4 5

2 3 3

2 4 7

3 4 2

3 5 3

Sample Output

Sample Input

- 6 15
- 0 1 17
- 0 2 13
- 0 3 10
- 0 4 11
- 0 5 10
- 1 2 11
- 1 3 15
- 1 4 15
- 1 5 14
- 2 3 12
- 2 4 16
- 2 5 12
- 3 4 14
- 3 5 10
- 4 5 12

Sample Output

(595 pts)

Description

Oganesson Dynamics wants to ensure that all of its communications within a network remain private, but this requires extra equipment for at least one end of each communicating pair of sites.

Given a graph where vertices represent work sites and edges represent pairs of sites that must be able to communicate securely, what are the fewest work sites that can have security-monitoring nodes to guarantee that each pair has at least one?

Hint: this is a brute-force optimization problem in which you will want to use many different optimizations (e.g., backtracking, branch and bound, etc.). As in homework 4's optimization problem, 400 points is considered to be full credit on the problem. However, 595 points are possible.

IMPORTANT NOTE 1: For this problem, your program must be guaranteed to output the correct answer (or timeout in the process of searching). If your program's output is incorrect for a single test case, it will receive 0 credit on this problem (regardless of how many other test cases it passes). Timing out is allowed, however. Test cases on which you time out will result in no points from that test case, but will not zero out points earned on test cases where your program returned the correct output.

IMPORTANT NOTE 2: Only the "Graded testcases" problem (HW5 Q5a on Gradescope) is graded. There is only a single testcase which actually consists of 100 different sub-testcases; each sub-testcase is worth either 2.5 points or 14 points, and the total amount of points possible across all sub-testcases is 592 points.

All other testcases in the "Ungraded testcases" problem (HW5 Q5b on Gradescope) are worth 0 points and are for testing purposes only. The first few test cases for the "Ungraded testcases" problem are visible test cases. This means that if you fail them, Gradescope will tell you the the expected output and the output your code actually gave. Note that these outputs may not be listed in order. The rest of the test cases for the "Ungraded testcases" problem are hidden test cases, which means you will not be shown the input or expected output.

IMPORTANT NOTE 3: Full credit for the problem will be 400 points. Thus, 195 bonus points are available.

Input Format

- The first line contains two values: the number of work sites (N) and the number of pairs of sites that must be able to communicate securely (M).
- The next *M* lines each describe a pair of communicating sites, indicating the unique id of each.

Constraints

• $1 \le N \le 200$ $1 \le M \le 10,000$ $0 \le \text{Work Site IDs} < N$

Output Format

A single number indicating the minimum number of work sites that must have security monitoring nodes to ensure that at least one exists for every communicating pair.

Sample Input

5 6

0 1

0 2

0 4

1 4

2 3

2 4

Sample Output

3

Explanation

Security nodes at sites 0, 1, and 2 will ensure that all communications in the network are safe.