Efficient Feasibility Checking Algorithm of Photovoltaic Array Reconfiguration

Dafang Zhao, Fukohito Ooshita, *Member, IEEE* and Michiko Inoue, *Member, IEEE*

Abstract—Power generation efficiency of photovoltaic (PV) systems are significantly affected by partial shading and PV cell damage. Partial shading or PV cell damage induces mismatched power generation among PV panels. Conducted bypass diodes under mismatch conditions result in loss of efficiency in power generation. Mismatched PV array can be recovered by reconfiguring electrical connections among PV panels in it. In this paper, a feasibility check problem of PV panel reconfiguration is introduced. This problem identifies whether a connection among PV panels can be configured from a given PV module-level solution. The proposed algorithm evaluated by comparison with the exhaustive search through random shading distributed PV array. The experimental results demonstrate that the proposed algorithm can identify feasible configurations more than 49,000 times faster than the exhaustive search with around 0.5% errors.

Index Terms—PV reconfiguration, partial-shading, mismatch, feasibility, heuristic

I. INTRODUCTION

IN recent years, the use of green and renewable energy sources has been increased with the aim to reduce fossil fuel depletion and environment pollution. Photovoltaic (PV) energy is one of the most promising emerging technologies. PV market growth by improvements of converting unlimited solar energy into electrical energy as well as the cost reductions of PV panels.

The use of PV systems for power generation brings many challenges. Due to the nature of PV cell, which is the basic component of PV array. PV system easily suffers from various forms of system faults, which include physical damage, temperature in-homogeneity, or partial shading. Unlike cell damage or other system faults, partial shading sources from cloud, dust or snow are very hard to prevent and predict. Thus, when PV cell could not uniformly generate power when they experience different irradiance or been damaged. This unbalanced working scenario will lead whole system mismatch. Mismatch condition might accelerate heating or aging of PV cells and furthermore hinder operation of maximum power point tracking (MPPT) algorithm, especially when the PV array output P-V curve becomes non-convex [1].

For several series connected PV cells, shaded or damaged cell causing normal cells to produce higher voltages that may reverse bias of "bad" cells. When a large number of series connected cells cause a huge reverse bias across shaded cells, leading to large dissipation of energy in the "bad" cells. This huge energy dissipation occurring in a small area might get overheating or burning of PV cells, or "hot-spots". To protect PV cells from "hot-spots", bypass diode is used to circumvent concentrated energy dissipation. However, the operation of

bypass diode will cause several stop delivering power and generate multiple local maximum power points or the total PV module current are limited by the one of worst "bad" cell. In both situation, available energy is lost. Furthermore, strand bypass diode can not complete eliminate hot-spotting [2].

In order to improve PV system power generation efficiency and protect PV cells from damage, an efficient and effectively PV system manage method is worth to investigate. The key to improve power generation efficiency is to maintain maximum output power. Thus, different maximum power point tracking (MPPT) algorithms have been proposed in this regard. In different MPPT algorithms, P&O [3]–[6]and hill climbing [7]–[9] methods received many attention due to its low complexity and implement cost. Hill climbing and P&O methods are different ways to implement the same fundamental method. Using the approximate linear relationship between V_{OC},I_{SC} and V_{MPP}, I_{MPP}, fractional voltage-based and current-based MPPT are popular due to their linear dependency of PV panel characteristic and irradiance level [10]–[14]. However, under abnormal working condition such as partial shading, it is very hard for conventional MPPT algorithms to find global maximum power point. Meanwhile, without any PV module level improvement, it is impossible to protect PV panel against hot-spotting [15], [16].

Another attractive direction to improve PV system efficiency under different working condition is the reconfiguration of PV arrays. This concept was first proposed by Salameh *et al.* [17] in 1990. Then in 2002, Sherif and Boutros *et al.* proposed a reconfigurable scheme for PV arrays that using transistors as switch network to improve performance [18]. Nguyen *et al.* proposed a method that divide PV array into two parts as the "fixed" part which is static connected PV modules and "adaptive" part that can be attached to "fixed" part with different configurations [19].

In this type of "fixed" - "adaptive" architecture, the mathematical formulation is not clear. Moreover, if the partial shading part is large enough to cover "adaptive" part, this scheme become ineffective. Velasco *et al.* [20]–[22] proposed an principle of operation is referred as "irradiance equalization". Reconfiguration strategy based on this principle aim to relocate the PV panels on the rows so that "irradiance equalization" is achieved. However, optimal configuration requires differences between "row" irradiance level are minimized as shown in [20] therefore it can only be realized on fully reconfigurable array. They proposed algorithm that calculate all possible configurations and stored as look-up table locally. Then, the "best" configuration for real-time shading scenario will be chosen from look-up table. There is no doubt that the number

of possible configurations will increase from the size of the PV array and it will be more difficult to determine the optimal configuration with limited time. Though ~~there are~~ several reconfiguration methods have been proposed, most works consider reconfiguration in PV cell or PV module level. However, those reconfiguration requires a significantly high computation time or a large number of switches to implement reconfiguration. In addition, they require special PV panels with the capability of switching, and cannot be applied to the system constructed with standard PV panels. Orozco-Gutierrez et al. proposed an efficient and effective reconfiguration method [23] where it first selects candidates of configurations using the product of approximated currents and voltages, then finds the best one with precise power simulation. However, these candidates are specified in a PV module level though PV modules could not be fully reconfigured. Actually, we found that some of configuration candidates are not able to be realized. However, the paper [23] does not show any systematic way to identify such a feasibility.

This paper proposes an optimization algorithm to check possibilities of configurations for reconfiguring PV arrays, which uses an approach that limit and sort possible combinations by the index of "Loss-rate". Proposed feasibility checking algorithm can rapidly check feasibility that a given configuration candidate can be actually formed by given PV panels. The experimental results demonstrate the effectiveness of the proposed method where it can identify the feasibility of configurations with a very small false negative rate of less than 1%.

II. PV ARRAY RECONFIGURATION

In this paper, we use following definition for PV array. A PV array is formed by PV panels. Several PV panels connected in series form a $String$. A PV panel is formed by three series-connected PV modules. A PV module consists of several series-connected PV cells with reverse biased bypass diode. Figure 1 represent the definition are introduced above. In a
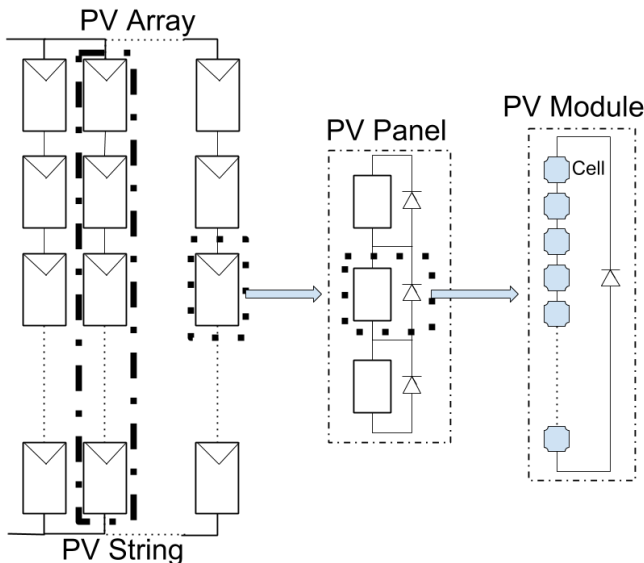

Fig. 1: PV array, string, module and panel

PV array, we have two constraints. PV panels in the same string have the same current level, while all the strings have the same control voltage. When considering reconfiguration of PV panel, we should find out the best configuration while considering these constraints.

As addressed in section I, Orozco-Guierrez et al. proposed an efficient reconfiguration algorithm for mismatched PV arrays. The method utilizes information of MPPs of every PV panel that is extracted using an online monitoring [24] and a power estimation [25].

Firstly, maximum power points of PV panels are extracted by using online monitoring from [24]. Then the current and voltage levels of maximum power point for each PV module s, I_{mpp_s} and V_{mpp_s}, are estimated from maximum power points of PV panels using the method in [25].

Estimated maximum power point current values (I_{mpp_s}) are approximated and close values are grouped into a small number of classes so that the number of possible combinations, in other words, search space is reduced. Estimated maximum power point voltage values (V_{mpp_s}) are averaged and the mean value $\bar{V}_{mpp}$ of all PV modules is used to approximate MPP voltage in the string. MPP voltage is approximated by $m \times \bar{V}_{mpp}$ where m is the number of active modules at the MPP. Then possible combination of configurations are enumerated.

For furthermore power simulation and configuration implementation, the current sequence and number of activating PV modules pre-string (it is called $feasibility$) are examined for each configuration. A simple feasibility checking method provided in [23] are described as follow.

(Feasibility 1)[1] For each n-th highest current level I_n in a current sequence Q for strings, it is required that the number of active modules in a configuration is $N(I_n)/n$ or less, where $N(I)$ is the total number of active modules for a current level I. The maximum number M of active modules for a PV array is given as

$$M = \min \left\{ N(I), \frac{N(I_2)}{2}, \frac{N(I_3)}{3}, \cdots, \frac{N(I_n)}{n} \right\}. \quad (1)$$

Thus, approximated power P can be calculated as follow.

$$P = \sum Q \times M \times \bar{V}_{mpp} \quad (2)$$

Approximated power values with and only with **Feasibility 1** is satisfied are listed into a power matrix like Table I. The configuration with maximum power value and configurations whose powers are more than 77% of maximum power are selected as configuration candidates [23].

III. FEASIBILITY

From Orozco $et\ al.$ method, to complete reconfiguration, we need to assign panels into strings to realize a candidate configuration. For example, consider a candidate configuration for a current sequence (2.5A, 2.5A) with 3 active modules per string as shown in Table I It can be realized by assigning panel P2 to the first string and panel P1 and P4 to the second

---

[1]In [23], only the case of $n = 2$ is considered, and we extended the condition to general n.

TABLE I

APPROXIMATED POWER

# modules per string	currents sequence for strings (A)									
	(3,3)	(3,2.5)	(3,1.5)	(3,0.5)	(2.5,2.5)	(2.5,1.5)	(2.5,0.5)	(1.5,1.5)	(1.5,0.5)	(0.5,0.5)
1	120W	110W	90W	70W	100W	80W	60W	60W	40W	20W
2	-	220W	180W	140W	200W	160W	120W	120W	80W	40W
3	-	**330W**	270W	210W	**300W**	240W	180W	180W	120W	60W
4	-	-	-	-	-	**320W**	240W	240W	160W	80W
5	-	-	-	-	-	-	**300W**	-	200W	100W
6	-	-	-	-	-	-	-	-	240W	120W

string as shown in Figure 2. However, for another current sequence (2.5A, 1.5A), there is no feasible assignment for 4 active modules per string, even though this configuration is listed in Table I by method in [23]. It is same for a configuration candidate with current sequence (3A, 2.5A) for 3 active modules per string. That is, the condition **Feasibility 1** is necessary but not sufficient to give an actual feasibility result.

For a PV array A with s strings, let $M_{A,i}(I)$ denotes the number of modules that can be activated with a current level I in the i-th string. Let $Q = (Q_1, Q_2, \ldots, Q_s)$ be the sequence of current required for strings. We define the feasibility as follows: A sequence of current Q is feasible with m modules if and only if it is possible to form a PV array A such that

$$M_{A,i}(Q_i) \geq m \tag{3}$$
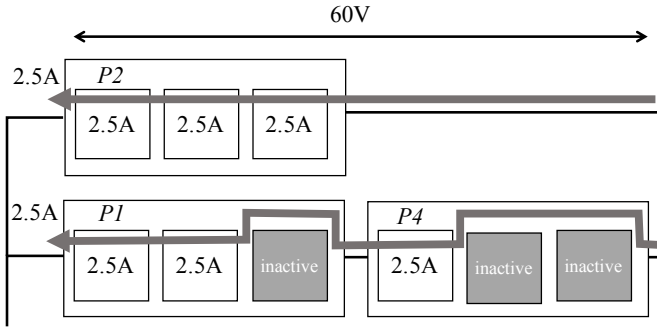
holds for each string i ($1 \leq i \leq s$).

Fig. 2: Feasible configuration for a current sequence (2.5A, 2.5A) with 3 active modules.

IV. Algorithm to Identify Feasibility

A. Feasibility problem

In this section, we introduce an algorithm to identify feasibility for a configuration. First we will formulate the feasibility problem that identifies the feasibility for a given current sequence and the number of active modules per string as follows.

Definition 1 (*Feasibility Problem*):
Input: Information of MPPs (approximated values), a current sequence Q, the number of m of active modules per-string.
Output: Whether Q with m modules is feasible or not.

TABLE II

Maximum number of activate PV modules

Current level	Q_1	Q_2	Q_3
Number of activate PV modules	2	2	3

B. Outline of the algorithm

In a PV array with s PV strings. For a given current sequence $Q = (Q_1, Q_2, \ldots, Q_s)$ ($Q_i \geq Q_j$ if $i \leq j$), we will assign PV panels to strings in the order of $Q_1, Q_2, \ldots, Q_s$. In this step, we will consider how much we lose an opportunity that the panel is more effectively used for the next current level. For example, as the PV panel shown in Table II which contain three PV modules. 2 of 3 modules are able to working at current level of Q_1 and Q_2. All 3 modules are able to activate at current level Q_3. Thus, if this panel been selected working for a PV string at current level $Q_{1 or 2}$, two-third of modules are able to deliver energy while the remaining one module been bypassed. We consider the "*loss*" that how many PV modules we fail to use in other current level and the "*gain*" that how many PV modules we can activate in the target current level together. Since two parameter are considered in the selection process, they can minimize "*loss*" while minimizing the number of selected PV panels. In this algorithm, we choose $Loss-rate$ to give selection priority. The $Loss$-$rate$ is the difference of active modules in panel p for the target current level $Q_{k'}$ in current sequence Q with other current levels per module. Let $M_p(I)$ denote the number of PV modules that can be active at current level I in panel p. In a s strings PV array, $Loss - rate$ for panel p of target current level k' on k-th current level Q_k ($2 \leq k \leq s$) is shown as follow.

$$Loss\text{-}rate(p, k', k) = \begin{cases} \dfrac{M_p(Q_k) - M_p(Q_{k'})}{M_p(Q_{k'})} & \text{if } M_p(Q_{k'}) \neq 0 \\ \infty & \text{if } M_p(Q_{k'}) = 0 \end{cases} \tag{4}$$

For example, as the panel p in Table II $Loss - rate(p, Q_1, Q_2) = 0$, $Loss - rate(p, Q_1, Q_3) = 0.5$, that is, for target current level Q_1 there is no "*losses*" for current level Q_2 and only "*loss*" 0.5 active module for current level Q_3.

To assign panels of target current level k' to k'-th string, the following step are applied. Let m be the required number of active modules per-string.

1) Sort un-selected panels with "*loss*" and "*gain*".
2) Selects sorted panels one by one until reaching the required number of active modules m.

3) Optimize the number of active modules. In step 3, if the number of selected active modules exceeds the required number m, redundant panels will be removed if exist.

As shown in Table III ,there is an example that demonstrate how to select PV panel for a string on current level Q_1 where required active module per string m is 4. Panels are sorted by $Loss - rate$ which seen as "$loss$" and $M_p(Q_1)$ which seen as "$gain$". In this example, panel P1, P4 and P5 have the same $Loss - rate$ is (0, 0.5) which is smaller than other panels. Thus, those three panels are grouped into the **First** "$loss$" level. Similarly, panel P2, P3, and P7 are in the **Second** "$loss$" level and panel P6 is in the **last** "$loss$" level because it does not contain any module that is able to activate on the target current level. Panel selection started from the **First** "$loss$" level, after selected panels active PV modules exceed m ($m = 4$ in this example) selection procedure will terminate. If some redundant panels exist, those panels are removed. In this example, no redundant panel exist and panel P1 and P4 are selected for target current level Q_1 PV string.

TABLE III

EXAMPLE OF PROPOSED ALGORITHM

		Panels						
		P1	P4	P5	P2	P3	P7	P6
	Q1	2	2	2	1	1	1	0
Number of modules	Q2	2	2	2	1	1	1	1
	Q3	3	3	3	2	2	2	2
Loss-rate(p,Q1,Q2)		0	0	0	0	0	0	∞
Loss-rate(p,Q1,Q3)		0.5	0.5	0.5	1	1	1	1
Priority		1	2	3	4	5	6	7

C. Algorithm

Proposed algorithm assigns PV panels of given configuration in order of target current level $Q_1, Q_2, \ldots, Q_s$. Where $Q_i \geq Q_j$ ($i \leq j$) holds for configurations current sequence $Q = (Q_1, Q_2, \ldots, Q_s)$. Pseudo code of proposed algorithm as shown below.

In the step 1, un-selected PV panels are sorted in a lexicographically ascending order of $Loss$-$rate(p, n, n + 1)$, $Loss$-$rate(p, n, n+2), \ldots, Loss$-$rate(p, n, s)$ and lexicographically descending order of $M_p(Q_{k'})$. This sorting procedure gives a basic priority for all un-selected PV panels. A panel with less "$loss$" and more modules for target current level are preferred. The last sorting index $M_p(Q_{k'})$ aim to minimize the number of selected PV modules for target current level $Q_{k'}$, and it will be described in step 3. Moreover, according to equation (4), the $Loss - rate$ is infinity for those PV panels who have no active modules.

In the step 2, sorted PV panels are selected one by one until reaching m, the required number of active modules per-string.

In the step 3, the number of active PV modules and the number of selected panels are optimized. As if selected PV panels contain active PV modules exceed required amount m, redundant panels will be removed if existed. A PV panel is redundant if and only if the number of active PV module still meet requirement or more even if this panel has been removed. This remove procedure starting from the last "$loss$" level, proposed algorithm will check $M_p Q_{k'}$ for each panel p.

If the number of active modules still exceed m after panel p is removed, the same procedure will repeat until all redundant panels are removed or every panel have been checked.

Algorithm 1: Feasibility Problem

Input: $Q = (Q_1, Q_2, \ldots, Q_s)$: current sequence for given configuration;
  m: the number of required active modules per-string;
  M_p: the number of active modules in each panel p;
  s: the number of strings;
Output: Feasibility Result;
1 **for** each target current level Q_k ($1 \leq k \leq s$) and each panel p **do**
2   **if** $\sum M_p(Q_k) < m$ **then**
3   **return :** Feasibility = NO
      /* Not enough active modules to meet requirement */
4   **else**
5     $Loss - rate(p, k, k + 1) = (M_p(Q_{k+1}) - M_p(Q_k))/M_p(Q_k)$
6   **end**
7 **end**
8 **for** each *target current level* Q_k **do**
9   P = a set of un-selected panels with $M_p(Q_k)$;
    // **step 1**
10 Sort PS in a lexicographically ascending order of $Loss - rate$ and descending order of $M_P(Q_j)$
    // **step 2**
11 select panels in the order in PS until $\sum M_p(Q_j)$ for selected panels $\geq m$;
    // **step 3**
12 group selected panels into "$loss$" levels (LV_1, $LV_2, \ldots, LV_h$) based on $Loss - rate$ ascending order.
13 **for** i = h to 1 **do**
14   find and remove redundant panels in "$loss$" level LV_i
15 **end**
16 **end**
17 **return :** Feasibility = YES

V. EVALUATION

A. Algorithm time complexity

In proposed algorithm, first, $Loss - rate(p, k', k)$ is evaluated for all PV panel p and target current level $Q_{k'}$ with other current level Q_k. This step takes $O(sn)$ where is n is the number of panels and s is the number of PV strings. Then the algorithm iterate as follows:

1) Sort panels by $Loss - rate$ and $M_P(Q_{k'})$
2) Select initial panels
3) Remove redundant panels

When selecting panels for target current level PV string, to sort panels, we use a lexicographically order of $Loss - rate$

and $M_P(Q_{k'})$. It can be done by sort panels recursively from the first component, that is, we sort panels by $Loss - rate(p, k', k)$, then, sort panels with same level by $Loss - rate(p, k', k+1)$ and so on. Since values of $Loss - rate$ and $M_P(Q_{k'})$ are related to the number of modules of PV panel and typically there are three PV modules in a commercial PV panel. Thus,we can use a bucket sort with 4 buckets. Therefore, it takes $O(n)$ for each bucket sort and $O(sn)$ for a lexicographical sorting in total. For initial selection, we check the panels in the sorted order, and we check panels again in reverse order to find and remove redundant panels. Therefore, these two steps take $Q(n)$. Since it takes $O(sn)$ for each PV string in addition to $O(sn)$ to obtain "$loss$" levels, the proposed algorithm takes $Q(s^2 n)$ in total.

B. Experimental setup

To evaluate proposed methods, we prepared MPP information of 100 mismatched PV arrays and applied algorithms. Exhaustive searching algorithm is also applied for comparison. MPP information is setup as follows for each PV array. Each PV array has 4 - 20 PV panels and the number of panels is randomly determined. Three current values are randomly assigned to the three PV modules in each panel from 2 - 20 current values, where the number of possible current values are also randomly determined for each PV array. The number of strings is 2 - 8 and it is randomly determined for each PV array.

We first applied the method in [23] to select configuration candidates. Condition **Feasibility 1** is applied to exclude a part of infeasible configurations, and then configurations that generate more than 77% of maximum power are selected from the remaining candidates as mentioned in [23]. In the experiments, 384 candidates are selected from 100 PV arrays as shown in Table IV. That is, for each PV array, 3 - 4 candidates are selected in average.

C. Experimental results

Table IV shows the results of feasibility checking by proposed algorithm and the exhaustive search algorithm. Since the exhaustive search checks all the combinations of panel connections for each configuration candidate. Though the exhaustive search identified total 268 feasible configurations, proposed algorithm identified 266 feasible configurations, only 2 configurations are misidentified. The proposed algorithms had a few false negative cases, but its rate was less than 1%. That is, it can correctly identify feasibility for most cases. Indeed, the proposed algorithm found at least one candidate configuration as feasible among 3 - 4 candidates for each PV arrays.

The calculation was performed using an Intel i7-4770 Quad-Core processor and 16 GB RAM. The time used by proposed algorithm is 1ms, which is significantly lower than 59s needed by the exhaustive search algorithm.

In conclusion, the proposed algorithms provide a solution for checking the feasibility of PV configurations with an error lower than 1% but execution speed are significantly faster than the one required by exhaustive search algorithm(about 49,000X).

TABLE IV

COMPARISON BETWEEN EXHAUSTIVE SEARCH
AND PROPOSED ALGORITHM

	Proposed Algorithm	Exhaustive Search
Number of PV array	100	
Number of MPP candidates	384	
Feasible candidates	266	268
Infeasible candidates	118	116
Error rate (%)	0.5	0
Ave. time pre candidate	0.0012s	59s

VI. CONCLUSION

Power generation of mismatched PV arrays can be recovered by reconfiguring connection of PV panels. In this paper, a feasibility check problem to configure PV panel connections from PV module connections is introduced and an algorithm that efficiently and effectively solve the problem are proposed. The experimental results show that the proposed algorithm can identify feasible configurations more than 49,000X faster than the exhaustive search with around 0.5% errors.

The proposed algorithm is useful to identify or configure actual PV panel connections from PV module level specifications or connections. Since there are many works of PV module level reconfiguration and PV panel level reconfiguration using PV module level optimization as an intermediate solution, the proposed algorithm is useful to actually find PV panel connections from PV module level solutions. PV panel level reconfiguration is a practical and realistic solution for PV arrays composed of standard PV panels, and the proposed algorithms contribute to finding an efficient and effective solution for PV panel level reconfiguration.

REFERENCES

[1] H. Islam, S. Mekhilef, N. B. M. Shah, T. K. Soon, M. Seyedmahmousian, B. Horan, and A. Stojcevski, "Performance evaluation of maximum power point tracking approaches and photovoltaic systems," *Energies*, vol. 11, no. 2, p. 365, 2018.

[2] K. A. Kim and P. T. Krein, "Reexamination of photovoltaic hot spotting to show inadequacy of the bypass diode," *IEEE Journal of Photovoltaics*, vol. 5, no. 5, pp. 1435–1441, 2015.

[3] N. Kasa, T. Iida, and H. Iwamoto, "Maximum power point tracking with capacitor identifier for photovoltaic power system," in *2000 Eighth International Conference on Power Electronics and Variable Speed Drives (IEE Conf. Publ. No. 475)*, pp. 130–135, Sep. 2000.

[4] S. Jain and V. Agarwal, "A new algorithm for rapid tracking of approximate maximum power point in photovoltaic systems," *IEEE power electronics letters*, vol. 2, no. 1, pp. 16–19, 2004.

[5] K. Chomsuwan, P. Prisuwanna, and V. Monyakul, "Photovoltaic grid-connected inverter using two-switch buck-boost converter," in *Conference Record of the Twenty-Ninth IEEE Photovoltaic Specialists Conference, 2002.*, pp. 1527–1530, IEEE, 2002.

[6] O. Wasynezuk, "Dynamic behavior of a class of photovoltaic power systems," *IEEE transactions on power apparatus and systems*, no. 9, pp. 3031–3037, 1983.

[7] E. Koutroulis, K. Kalaitzakis, and N. C. Voulgaris, "Development of a microcontroller-based, photovoltaic maximum power point tracking control system," *IEEE Transactions on power electronics*, vol. 16, no. 1, pp. 46–54, 2001.

[8] W. Teulings, J. Marpinard, A. Capel, and D. O'sullivan, "A new maximum power point tracking system," in *Proceedings of IEEE Power Electronics Specialist Conference-PESC'93*, pp. 833–838, IEEE, 1993.

[9] W. Xiao and W. G. Dunford, "A modified adaptive hill climbing mppt method for photovoltaic power systems," in *2004 IEEE 35th annual power electronics specialists conference (IEEE Cat. No. 04CH37551)*, vol. 3, pp. 1957–1963, Ieee, 2004.

[10] C. Liu, W. Liu, L. Wang, G. Hu, L. Ma, and B. Ren, "A new method of modeling and state of charge estimation of the battery," *Journal of Power Sources*, vol. 320, pp. 1–12, 2016.

[11] K. Kobayashi, H. Matsuo, and Y. Sekine, "A novel optimum operating point tracker of the solar cell power supply system," in *2004 IEEE 35th Annual Power Electronics Specialists Conference (IEEE Cat. No. 04CH37551)*, vol. 3, pp. 2147–2151, IEEE, 2004.

[12] B. Bekker and H. Beukes, "Finding an optimal pv panel maximum power point tracking method," in *2004 IEEE Africon. 7th Africon Conference in Africa (IEEE Cat. No. 04CH37590)*, vol. 2, pp. 1125–1129, IEEE, 2004.

[13] N. Mutoh, T. Matuo, K. Okada, and M. Sakai, "Prediction-data-based maximum-power-point-tracking method for photovoltaic power generation systems," in *2002 IEEE 33rd Annual IEEE Power Electronics Specialists Conference. Proceedings (Cat. No. 02CH37289)*, vol. 3, pp. 1489–1494, IEEE, 2002.

[14] T. Noguchi, S. Togashi, and R. Nakamoto, "Short-current-pulse based adaptive maximum-power-point tracking for photovoltaic power generation system," *IEEJ Transactions on Industry Applications*, vol. 121, no. 1, pp. 78–83, 2001.

[15] C. Olalla, M. N. Hasan, C. Deline, and D. Maksimović, "Mitigation of hot-spots in photovoltaic systems using distributed power electronics," *Energies*, vol. 11, no. 4, pp. 1–16, 2018.

[16] T. Ghanbari, "Permanent partial shading detection for protection of photovoltaic panels against hot spotting," *IET Renewable Power Generation*, vol. 11, no. 1, pp. 123–131, 2016.

[17] Z. Salameh and F. Dagher, "The effect of electrical array reconfiguration on the performance of a PV-powered volumetric water pump," *IEEE Transactions on Energy Conversion*, vol. 5, no. 4, pp. 653–658, 1990.

[18] R. A. Sherif and K. S. Boutros, "Solar module array with reconfigurable tile," Feb. 26 2002. US Patent 6,350,944.

[19] D. Nguyen and B. Lehman, "A reconfigurable solar photovoltaic array under shadow conditions," *Conference Proceedings - IEEE Applied Power Electronics Conference and Exposition - APEC*, pp. 980–986, 2008.

[20] G. Velasco-Quesada, F. Guinjoan-Gispert, R. Piqué-López, M. Román-Lumbreras, and A. Conesa-Roca, "Electrical PV array reconfiguration strategy for energy extraction improvement in grid-connected PV systems," *IEEE Transactions on Industrial Electronics*, vol. 56, no. 11, pp. 4319–4331, 2009.

[21] G. Velasco, F. Guinjoan, and R. Pique, "Grid-connected pv systems energy extraction improvement by means of an electric array reconfiguration (ear) strategy: Operating principle and experimental results," in *2008 IEEE Power Electronics Specialists Conference*, pp. 1983–1988, IEEE, 2008.

[22] G. Velasco, J. Negroni, F. Guinjoan, and R. Piqué, "Energy generation in pv grid-connected systems: A comparative study depending on the pv generator configuration," in *IEEE International Symposium on Industrial Electronics*, pp. 1025–1030, 2005.

[23] M. L. Orozco-Gutierrez, G. Spagnuolo, J. M. Ramirez-Scarpetta, G. Petrone, and C. A. Ramos-Paja, "Optimized Configuration of Mismatched Photovoltaic Arrays," *IEEE Journal of Photovoltaics*, vol. 6, no. 5, pp. 1210–1220, 2016.

[24] P. L. Carotenuto, P. Manganiello, G. Petrone, and G. Spagnuolo, "Online recording a PV module fingerprint," *IEEE Journal of Photovoltaics*, vol. 4, no. 2, pp. 659–668, 2014.

[25] M. L. Orozco-Gutierrez, J. M. Ramirez-Scarpetta, G. Spagnuolo, G. Petrone, and C. A. Ramos-Paja, "Fast estimation of MPPs in mismatched PV arrays based on lossless model," *5th International Conference on Clean Electrical Power: Renewable Energy Resources Impact, ICCEP 2015*, pp. 435–440, 2015.