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Basic Number Theory-2

TUTORIAL PROBLEMS

The concept of prime numbers is a very important concept in math. This article discusses the concept of prime numbers and related properties.

What are prime numbers and composite numbers?

Prime numbers are those numbers that are greater than 1 and have only two factors 1 and itself.

Composite numbers are also numbers that are greater than 1 but they have at least one more divisor other than 1 and itself.

For example, 5 is a prime number because 5 is divisible only by 1 and 5 only. However, 6 is a composite number as 6 is divisible by 1, 2, 3, and 6.

There are various methods to check whether a number is prime.

Naive approach

Traverse all the numbers from 1 to N and count the number of divisors. If the number of divisors is equal to 2 then the given number is prime, else it is not.

```
void checkprime(int N){
    int count = 0;
    for( int i = 1; i <= N; ++i )
        if( N % i == 0 )
            count++;
    if(count == 2)
```

```

        cout << N << " is a prime number." << endl;
    else
        cout << N << " is not a prime number." << endl;
}

```

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Time complexity

The time complexity of this function is $O(N)$ because you traverse from 1 to N .

Better approach

If you have two positive numbers N and D , such that N is divisible by D and D is less than the square root of N .

- (N/D) must be greater than the square root of N .
- N is also divisible by (N/D) . If there is a divisor of N that is less than the square root of N , then there will be a divisor of N that is greater than square root of N . You will have to traverse till the square root of N .

Note: You are generating all the divisors of N and if the count of divisors is greater than 2, then the number is composite.

For example, if $N=50$, $\sqrt{N}=7$ (floor value). You will iterate from 1 to 7 and count the number of divisors of N . The divisors of N are 1, 50; 2, 25; 5, 10. You have 6 divisors of 50, and therefore, it is not prime.

```

void checkprime(int N) {
    int count = 0;
    for( int i = 1; i * i <= N; ++i ) {
        if( N % i == 0 ) {
            if( i * i == N )
                count++;
            else // i < sqrt(N) and (N / i) > sqrt(N)
                count += 2;
        }
    }
    if(count == 2)
        cout << N << " is a prime number." << endl;
    else
        cout << N << " is not a prime number." << endl;
}

```

Time complexity

The time complexity of this function is $O(\sqrt{N})$ because you traverse from 1 to \sqrt{N} .

Sieve of Eratosthenes

You can use the *Sieve of Eratosthenes* to find all the prime numbers that are less than or equal to a given number N or to find out whether a number is a prime number.

The basic idea behind the Sieve of Eratosthenes is that at each iteration one prime number is picked up and all its multiples are eliminated. After the elimination process is complete, all the unmarked numbers that remain are prime.

Pseudo code

1. Mark all the numbers as prime numbers except 1
2. Traverse over each prime numbers smaller than \sqrt{N}
3. For each prime number, mark its multiples as composite numbers
4. Numbers, which are not the multiples of any number, will remain marked as prime number and others will change to composite numbers.

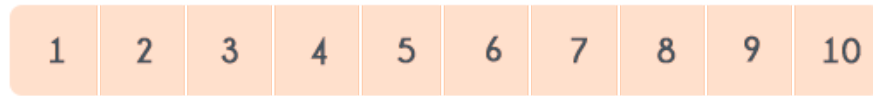
```
void sieve(int N) {
    bool isPrime[N+1];
    for(int i = 0; i <= N; ++i) {
        isPrime[i] = true;
    }
    isPrime[0] = false;
    isPrime[1] = false;
    for(int i = 2; i * i <= N; ++i) {
        if(isPrime[i] == true) {           //Mark all the
multiples of i as composite numbers
            for(int j = i * i; j <= N ;j += i)
                isPrime[j] = false;
        }
    }
}
```

This code will compute all the prime numbers that are smaller than or equal to N .

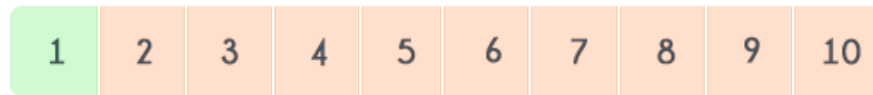
Let us compute prime numbers where $N = 10$.

1. Mark all the numbers as prime
2. Mark 1 as composite

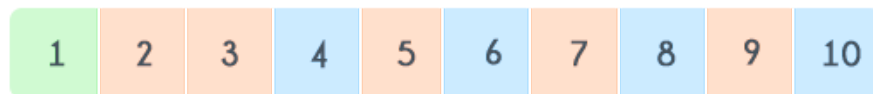
In each iteration, check if a number is prime or not, if it is then mark all of its multiple as composite.



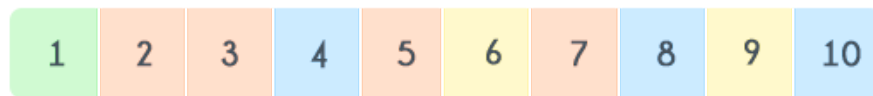
Mark 1 as composite



Mark multiples of 2 as composite



Mark multiples of 3 as composite



Next prime is 5 which is greater than $\sqrt{10}$ so the loop will break

The prime numbers are 2, 3, 5, and 7.

Time complexity

The inner loop that runs for each element is as follows:

1. If $i = 2$, inner loop runs $N / 2$ times
2. If $i = 3$, inner loop runs $N / 3$ times
3. If $i = 5$, inner loop runs $N / 5$ times

Total complexity

$$N * (\frac{1}{2} + \frac{1}{3} + \frac{1}{5} + \dots) = O(N \log \log N)$$

Reference for complexity analysis: [Sieve of Erasthones](#)

Modification of Sieve of Eratosthenes for fast factorization

Factorization in \sqrt{N}

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```
vector<int> factorize(int n) {
    vector<int> res;
    for (int i = 2; i * i <= n; ++i) {
        while (n % i == 0) {
            res.push_back(i);
            n /= i;
        }
    }
    if (n != 1) {
        res.push_back(n);
    }
    return res;
}
```

For example,

$n = 50, i = 2; res = \{2\}$

$n = 25, i = 3; res = \{2\}$

$n = 25, i = 4; res = \{2\}$

$n = 25, i = 5; res = \{2, 5, 5\}$

$n = 1$

When you exit the for loop, *res vector* is the factorization of $N = 50$.

At every step, you must look for the prime number of the least value, which divides the current N . This is the main idea of this modification.

Let us construct an array which will give us this number in $O(1)$ time.

```
int minPrime[n + 1];
for (int i = 2; i * i <= n; ++i) {
    if (minPrime[i] == 0) {          //If i is prime
        for (int j = i * i; j <= n; j += i) {
            if (minPrime[j] == 0) {
                minPrime[j] = i;
            }
        }
    }
}

for (int i = 2; i <= n; ++i) {
    if (minPrime[i] == 0) {
```

```

        minPrime[i] = i;
    }
}

```

Now, use this modification to factorize N in $O(\log(N))$ time.

```

vector<int> factorize(int n) {
    vector<int> res;
    while (n != 1) {
        res.push_back(minPrime[n]);
        n /= minPrime[n];
    }
    return res;
}

```

For example,

$n = 50, \text{minprime}[50] = 2, \text{res} = 2$

$n = 25, \text{minprime}[25] = 5, \text{res} = 2, 5$

$n = 5, \text{minprime}[5] = 5, \text{res} = 2, 5, 5$

$n = 1$

The required factors are in *res*.

Conditions

You can implement this modification only if you are allowed to create an array of integers of size N .

Note: This approach is useful when you need to factorize not-very-large numbers. It is not necessary to build a modified Sieve for every problem in which factorization is required. Moreover, you cannot build it for large numbers N like 10^9 or 10^{12} . Therefore, for such large numbers it is recommended that you factorize in $O(\sqrt{N})$ instead.

Fact

If the factorization of N is $p_1^{q_1} * p_2^{q_2} * \dots * p_k^{q_k}$ where $p_1, p_2 \dots p_k$ are the prime factors of N and $q_1, q_2 \dots q_k$ are the powers of the respective prime factors, then N has $(q_1 + 1) * (q_2 + 1) * \dots * (q_k + 1)$ distinct divisors.

Sieve of Eratosthenes on the segment:

Sometimes you need to find all the primes that are in the range $[L \dots R]$ and not in $[1 \dots N]$, where R is a large number.

Conditions

You are allowed to create an array of integers with size $(R - L + 1)$.

Implementation

```
bool isPrime[r - 1 + 1]; //filled by true
for (long long i = 2; i * i <= r; ++i) {
    for (long long j = max(i * i, (1 + (i - 1)) / i * i); j <= r; j += i) {
        isPrime[j - 1] = false;
    }
}
for (long long i = max(1, 2); i <= r; ++i) {
    if (isPrime[i - 1]) {
        //then i is prime
    }
}
```

The approximate complexity is $O(\sqrt{R})$

Suggestions

It is recommended that you do not build a Sieve to check several numbers for primality. Use the following function instead, which works in $O(\sqrt{N})$ for every number:

```
bool isPrime(int n) {
    for (int i = 2; i * i <= n; ++i) {
        if (n % i == 0) {
            return false;
        }
    }
    return true;
}
```

Contributed by: Shubham Gupta

Did you find this tutorial helpful?



YES



NO

TEST YOUR UNDERSTANDING

Number of Primes

Given a number N , find number of primes in the range $[1, N]$.

Input:

The only line of input consists of a number N .

Output:

Print the number of primes in the range $[1, N]$.

Constraints:

$$1 \leq N \leq 1000000$$

SAMPLE INPUT



10

SAMPLE OUTPUT



4

Enter your code or [Upload your code as file.](#)

Save

C (gcc 4.8.2)



```
1 #include <stdio.h>
2
3 int main()
4 {
5     printf("Hello World!\n");
6     return 0;
7 }
8
```

1:1

☒ Provide custom input

COMPILE & TEST

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Need Help ?

In case you feel you are stuck with the problem, you can view our editorial.
Remember this is just to help you out, in order to learn you should try your best.

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COMMENTS (7) ↻

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Join Discussion...

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**Ashwani Singh** ✍ Edited 4 months ago

I think there is Problem in your Sieve of Eratosthenes code
In sieve() function inner loop should start from $j=i*2$ rather than starting from $j=i*i$
If i am wrong please correct me.

▲ 1 vote ● Reply ● Message ● Permalink

**Shubham Gupta** ⚡ Moderator 4 months ago

The numbers from $2*i$ to $(i-1)*i$ are covered by other primes less than i , so it can be omitted.

▲ 2 votes ● Reply ● Message ● Permalink

**Ashwani Singh** 4 months ago

Thanks

▲ 0 votes ● Reply ● Message ● Permalink

**Shashank Gupta** 2 months ago

In "Modification of Sieve of Eratosthenes for fast factorization" you can remove the below for loop and put `minPrime[i] = i;` in first for loop itself.

```
for (int i = 2; i <= n; ++i) {  
    if (minPrime[i] == 0) {  
        minPrime[i] = i;  
    }  
}
```

▲ 0 votes ● Reply ● Message ● Permalink

**Shashank Shekhar** 2 months ago

The first for loop is going till \sqrt{n} . That is why we are using the second for loop.

▲ 0 votes ● Reply ● Message ● Permalink

**ashish chavan** 12 days ago

```
int checkPrime(int n){
    if(n<2)return 0;
    for(int i=2;i*i<=n;i++){
        if(n%i==0)return 0;
    }
    return 1;
}
```

▲ 0 votes ● Reply ● Message ● Permalink

**Shubham Prasad** 5 days ago

in segmented sieve, how did you get this lower limit for j ?

```
for (long long j = max(i * i, (l + (i - 1)) / i * i); j <= r; j += i) {
    isPrime[j - l] = false;
}
```

▲ 0 votes ● Reply ● Message ● Permalink

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