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### LCM - Editorial

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## SOLUTION TO LCM

The solution is to use Möbius inversion. See for example the Wikipedia for information on the Möbius function  $\mu$  and related concepts.

Assume throughout that  $A \leq B$ . We want to compute  $\sum_{a=1}^A \sum_{b=1}^B abF(\gcd(a, b))$ , where  $F(n) = \frac{1}{n}$  if  $n$  is square-free and 0 otherwise. Adapting a standard approach, we do the following:

Let  $f(n) = \sum_{d|n} \mu(d)F(\frac{n}{d})$ . Then  $F(n) = \sum_{d|n} f(d)$  by the Möbius inversion formula, and

$$\begin{aligned} & \sum_{a=1}^A \sum_{b=1}^B abF(\gcd(a, b)) \\ &= \sum_{a=1}^A \sum_{b=1}^B ab \sum_{d|\gcd(a, b)} f(d) \\ &= \sum_{d=1}^A \sum_{i=1}^{\lfloor A/d \rfloor} \sum_{j=1}^{\lfloor B/d \rfloor} (id)(jd)f(d) \\ &= \sum_{d=1}^A d^2 f(d) \sum_{i=1}^{\lfloor A/d \rfloor} i \sum_{j=1}^{\lfloor B/d \rfloor} j \\ &= \sum_{d=1}^A d^2 f(d) \binom{\lfloor A/d \rfloor + 1}{2} \binom{\lfloor B/d \rfloor + 1}{2} \\ &= \sum_{d=1}^A d^2 f(d) \lfloor A/d \rfloor \lfloor A/d + 1 \rfloor \lfloor B/d \rfloor \lfloor B/d + 1 \rfloor / 4 \end{aligned}$$

For the second equality: Since  $a \leq A \leq B$  and  $d$  divides  $a$ ,  $d$  can range from 1 to  $A$ . Since  $a$  and  $b$  are multiples of  $d$ , we can write them as  $id$  and  $jd$ , respectively, with  $i \leq \lfloor A/d \rfloor$  and  $j \leq \lfloor B/d \rfloor$ .

For any  $n \geq 1$ , we have  $f(n) = \sum_{d|n} \mu(n)F(\frac{n}{d}) = \frac{1}{n} \sum_{d|n} \mu(d)d\mu^2(\frac{n}{d})$ . ( $\mu^2(k) = 1$  iff  $k$  is square-free.) Since  $F$  is multiplicative, so is  $f$ . Using this it is easy to find  $f$ . If  $n$  is divisible by  $p^3$  for some  $p$ , then  $f(n) = 0$ . Otherwise,  $n = p_1^2 p_2^2 \cdots p_s^2 q_1 q_2 \cdots q_t$ . Letting  $k = p_1 p_2 \cdots p_s$  and  $q = q_1 q_2 \cdots q_t$  we have  $f(n) = \frac{1}{n} \prod_{p|k} -p \prod_{p|q} (1-p)$ , where the products are taken over all prime divisors  $p$ .

Let  $g(n) = \prod_{p|k} -p \prod_{p|q} (1-p)$ . That is: Start with 1. For each prime  $p$  such that  $p$  divides  $n$  but  $p^2$  does not divide  $n$ , multiply by  $1-p$ . For each prime  $p$  such that  $p^2$  divides  $n$ , multiply by  $-p$ . If  $p^3$  divides  $n$  for some prime  $p$ ,  $g(n) = 0$ . We can precompute  $g$  as follows: First use a prime sieve to find a prime factor of  $n$  for each  $n \leq 4000000$ . Then compute  $f[n]$  iteratively. The sieve of Eratosthenes is too slow. A wheel sieve should work, but a segmented sieve is simpler and faster in practise. In code:

```
int MAX = 4000000;
int g[4000000];

// Compute prime factors
// .
// .
// Now g[n] equals 0 if n is prime, otherwise g[n] is a prime factor of n.

for(int n=2; n<=MAX; n++) {
    if(g[n]==0) {
        g[n]=1-n; // n is prime
        continue;
    }

    p=g[n];
    int m=n/p;

    if(m%p != 0) g[n]=f[m]*(1-p); // p divides i, but p^2 does not,
                                // so g[n] = g[n/p]*(1-p)
    else { // p^2 divides i
        m/=p;
        if(m%p != 0) g[n] = g[m]*(-p); // p^3 does not divide i
        else g[n]=0; // p^3 divides i
    }
}
```


It is possible to do better, using an all-in-one segmented approach that also avoids divisions, but this is good enough.

To compute  $\sum_{d=1}^A dg(d) \binom{\lfloor A/d \rfloor + 1}{2} \binom{\lfloor B/d \rfloor + 1}{2}$ , note that  $\lfloor A/d \rfloor$  takes on at most  $2\sqrt{A} + 1$  values, and  $\lfloor B/d \rfloor$  takes on at most  $2\sqrt{B} + 1$  values. So, we can precompute  $\sum_{d=1}^k dg(d)$  for all  $k \leq 4000000$  and group terms that have the same values for  $\lfloor A/d \rfloor$  and  $\lfloor B/d \rfloor$ .

Computing the result modulo  $2^{30}$  allows using the “wrap-around” arithmetic of unsigned 32-bit integers, and clearing the two high bits of the final result. For languages that don’t have wrap-around arithmetic,  $x \bmod 2^{30}$  can be computed very quickly as  $x \text{ AND } (2^{30} - 1)$ .

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