

# Delaunay properties of digital straight segments

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# Outline

Definitions: patterns and Delaunay triangulation

Observation: Delaunay triangulation of patterns?

Characterization: proof

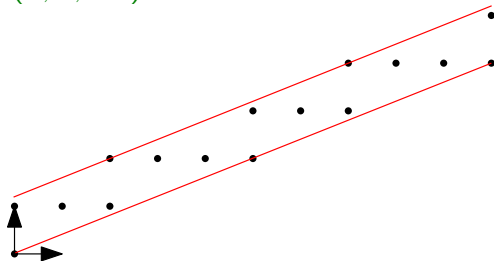
Conclusion and perspectives: new algorithms

# Digital straight line (DSL)

## Standard DSL

The points  $(x, y) \in \mathbb{Z}^2$  verifying  $\mu \leq ax - by < \mu + |a| + |b|$  belong to the standard DSL  $D(a, b, \mu)$  of slope  $\frac{a}{b}$  and intercept  $\mu$  ( $a, b, \mu \in \mathbb{Z}$  and  $\text{pgcd}(a, b) = 1$ ).

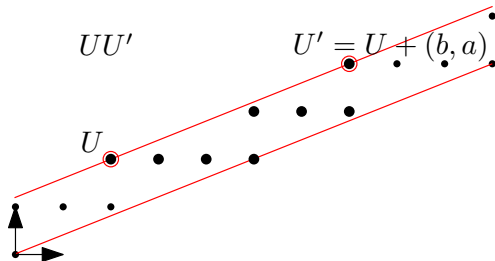
Example:  $D(2, 5, -6)$



# Pattern

- ▶ a pattern is a subsequence of a DSL between two consecutive upper leaning points

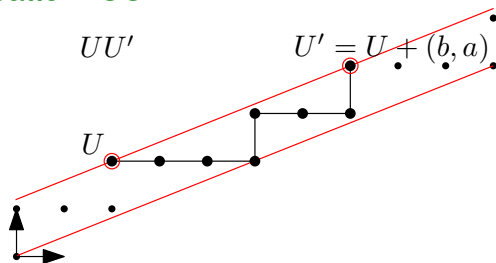
Example: pattern  $UU'$



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- ▶ a pattern is a subsequence of a DSL between two consecutive upper leaning points
- ▶ its staircase representation is the polygonal line linking the points in order

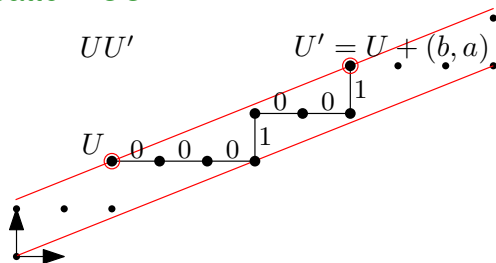
Example: pattern  $UU'$



# Pattern

- ▶ a pattern is a subsequence of a DSL between two consecutive upper leaning points
- ▶ its staircase representation is the polygonal line linking the points in order
- ▶ its chain code is a Christoffel word

Example: pattern  $UU'$



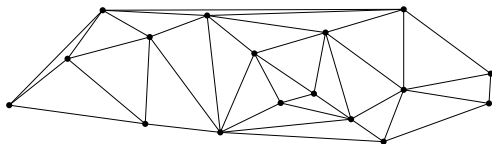
# Delaunay triangulation

## Triangulation of a finite set of points $\mathcal{S}$

Partition of the convex hull of  $\mathcal{S}$  into triangular facets, whose vertices are points of  $\mathcal{S}$ .

## Delaunay condition

The interior of the circumcircle of each triangular facet does not contain any set point.



always exists and is unique (without 4 cocircular points)

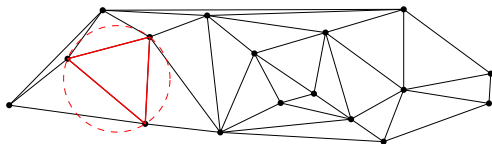
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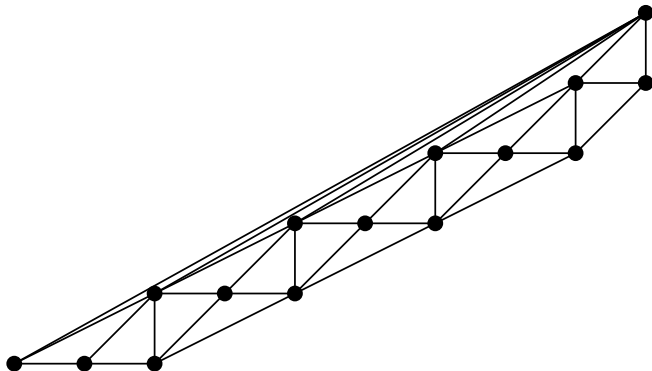
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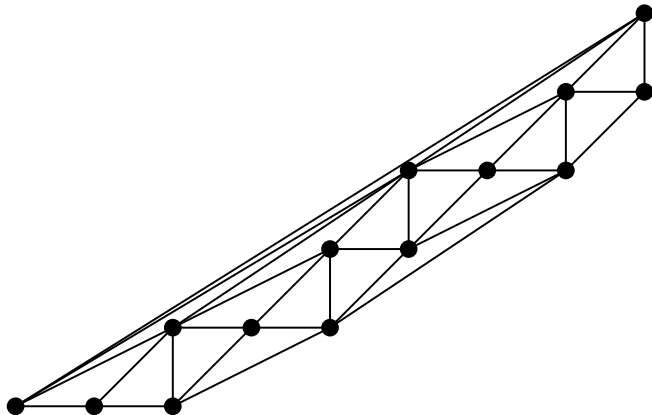
# Delaunay triangulation of patterns

Pattern of slope  $5/9$



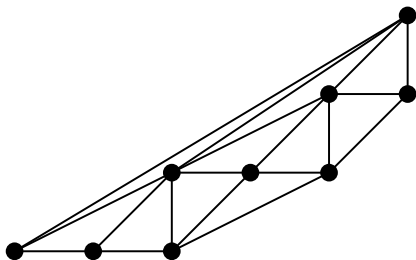
# Delaunay triangulation of patterns

Pattern of slope  $5/8$



# Delaunay triangulation of patterns

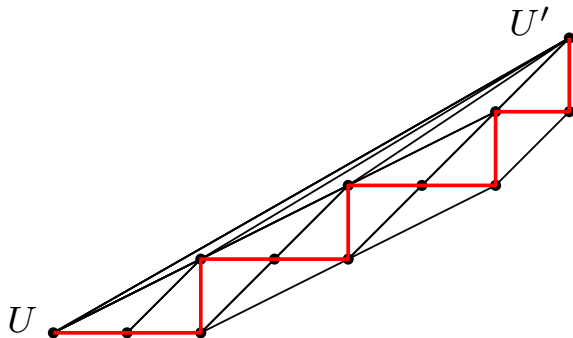
Pattern of slope  $2/5$



## Three remarks

1. the Delaunay triangulation of  $UU'$  contains the staircase representation of  $UU'$ .

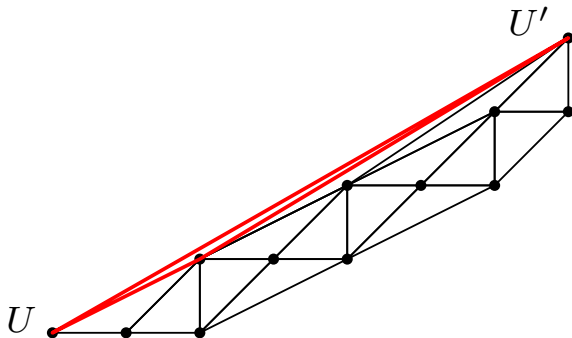
Pattern of slope 4/7



## Three remarks

1. the Delaunay triangulation of  $UU'$  contains the staircase representation of  $UU'$ .
2.  $U$ ,  $U'$  and the closest point of  $UU'$  to  $[UU']$  (Bezout point) define a facet.

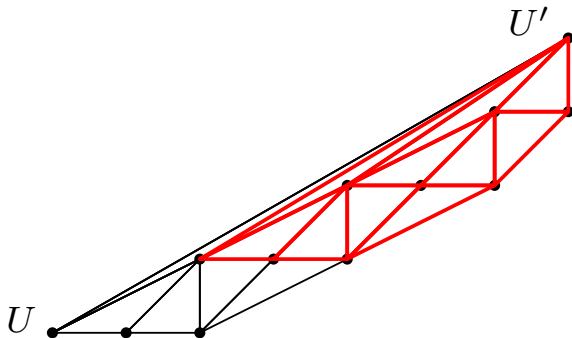
Pattern of slope 4/7



## Three remarks

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2.  $U$ ,  $U'$  and the closest point of  $UU'$  to  $[UU']$  (Bezout point) define a facet.
3. the Delaunay triangulation of some patterns contains the Delaunay triangulation of subpatterns.

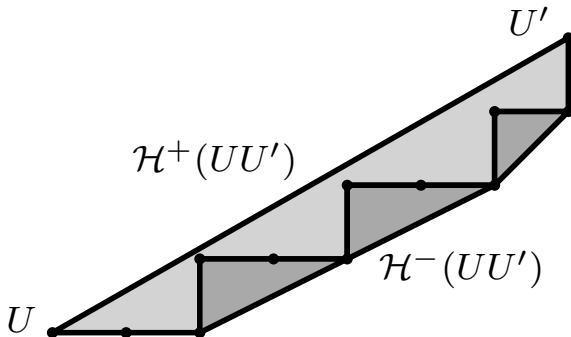
Pattern of slope 4/7



## Dividing the triangulation (remark 1)

- ▶ The convex hull is divided into an upper part  $\mathcal{H}^+(UU')$  and a lower part  $\mathcal{H}^-(UU')$ .

Pattern of slope 4/7

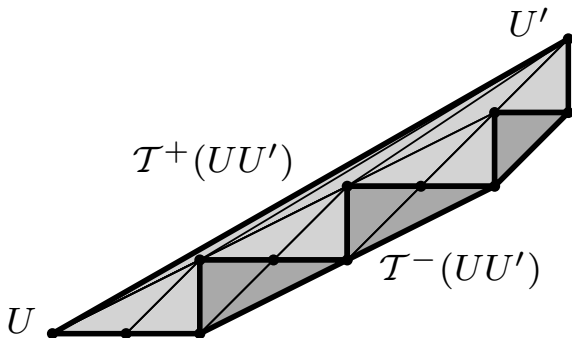




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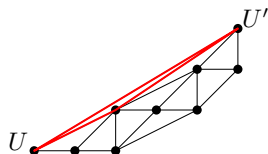
- ▶ The convex hull is divided into an upper part  $\mathcal{H}^+(UU')$  and a lower part  $\mathcal{H}^-(UU')$ .
- ▶ The Delaunay triangulation is divided into an upper part  $\mathcal{T}^+(UU')$  and a lower part  $\mathcal{T}^-(UU')$ .

Pattern of slope 4/7



# Facets of a pattern

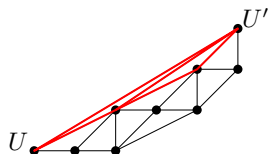
- ▶ main facet (remark 2)
  - ▶ geometrical characterization (Bezout point)
  - ▶ combinatorial characterization (splitting formula)
- ▶ induction (remark 3)



0 0 1 | 0 0 1 0 1

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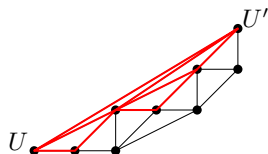
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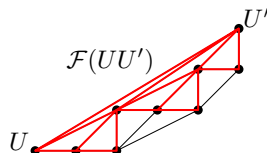
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- ▶ main facet (remark 2)
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0 | 0|1|0 | 0|1|0|1

# Main result

## Theorem

The facets  $\mathcal{F}(UU')$  of the pattern  $UU'$  is a triangulation of  $\mathcal{H}^+(UU')$  such that each facet has points of  $UU'$  as vertices and satisfies the Delaunay property, i.e.  $\mathcal{F}(UU') = \mathcal{T}^+(UU')$ .

the (upper part of the) Delaunay triangulation of a pattern is characterized by the continued fraction expansion of its slope

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# Sketch of the proof

#1

- ▶ no triangular facet of the Delaunay triangulation of a pattern  $UU'$  can cross its staircase representation
- ▶ the set of facets  $\mathcal{F}(UU')$  is the *unique* way of triangulating  $\mathcal{H}^+(UU')$

To be more constructive, we chose:

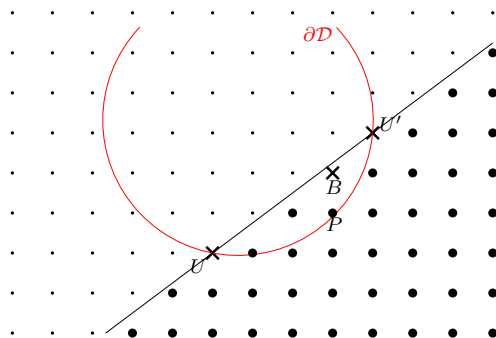
#2

- ▶ the set of facets  $\mathcal{F}(UU')$  is a triangulation of  $\mathcal{H}^+(UU')$  (easy part)
- ▶ the interior of the circumcircle of each facet of  $\mathcal{F}(UU')$  does not contain any point of  $UU'$  (let us focus on that part)



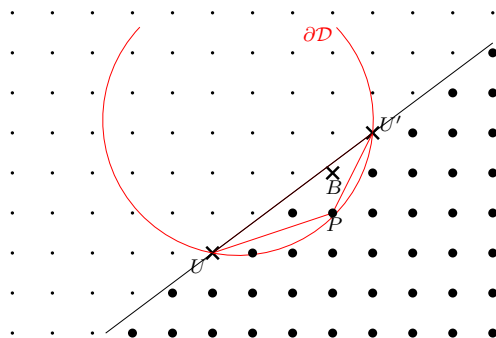
# Lemma 1

Let  $\mathcal{D}$  be a disk whose boundary passes through  $U$  and  $U'$  and whose center is located above  $(UU')$ . Let  $\partial\mathcal{D}$  be its boundary.  $\mathcal{D} \setminus \partial\mathcal{D}$  contains a lattice point below or on  $(UU')$  if and only if it contains (at least)  $B$ , the lower Bezout point of  $[UU']$ .



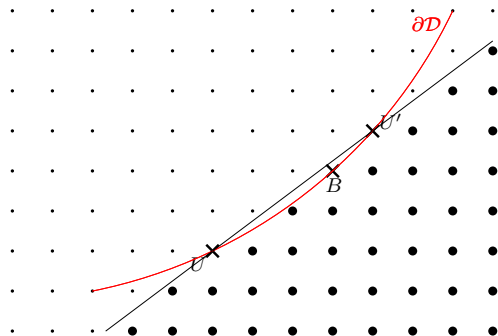
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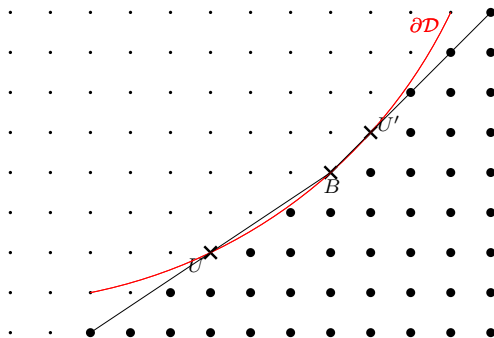
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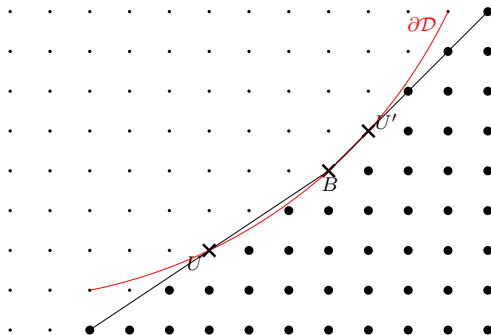
## Lemma 2

Let  $\mathcal{D}$  be a disk whose boundary  $\partial\mathcal{D}$  is the circumcircle of  $UBU'$ .  $\mathcal{D} \setminus \partial\mathcal{D}$  contains none of the *background points* of  $UU'$  (lattice points below  $UB$  or  $BU'$ ).



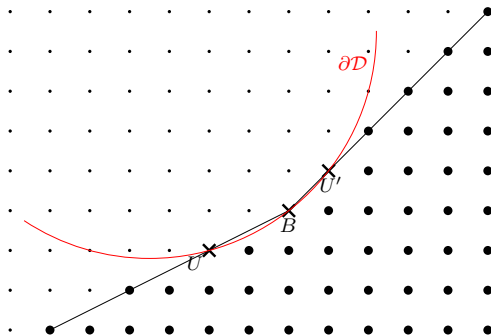
# Induction

- ▶ The circumcircle of the main facet  $UBU'$  contains none of the background points of  $UU'$  in its interior (lemma 2).
- ▶ The background points of  $UB$  and  $BU'$  contain the background points of  $UU'$ , which contains the set points.



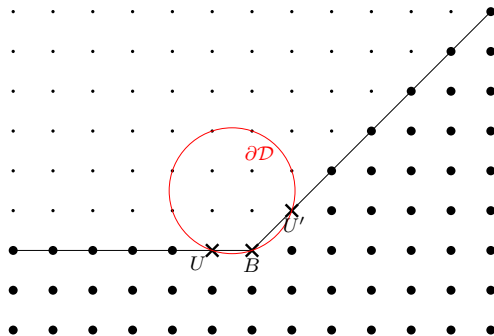
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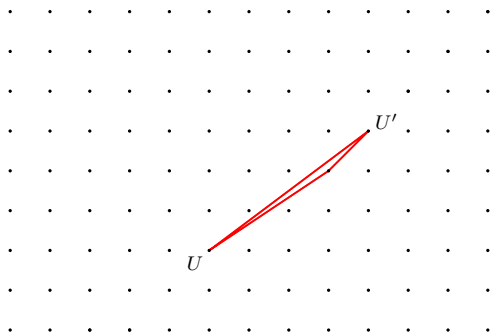
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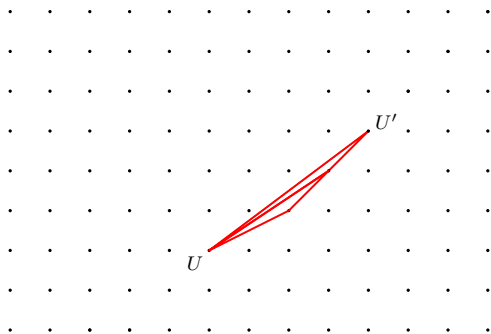
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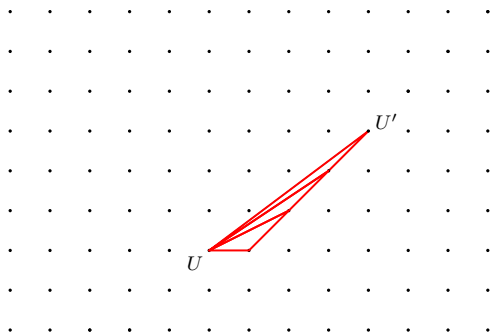
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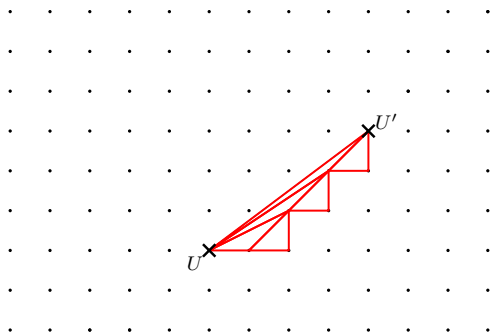
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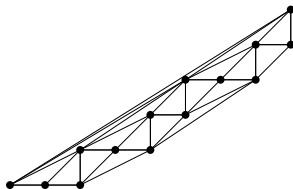
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# Delaunay triangulation computation

## ► Pattern

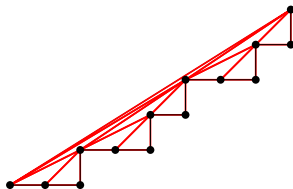
pattern of slope  $8/5$



# Delaunay triangulation computation

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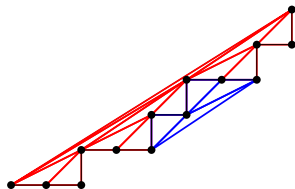
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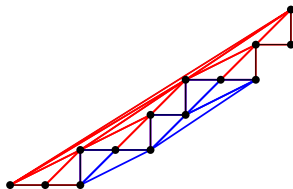
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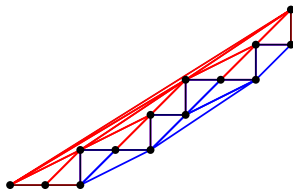




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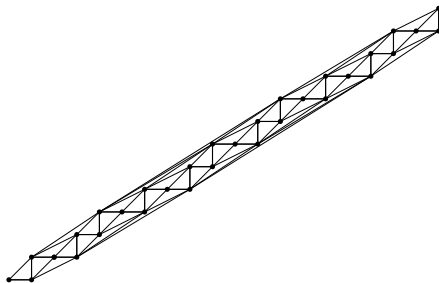
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- ▶ Pattern
- ▶ DSS

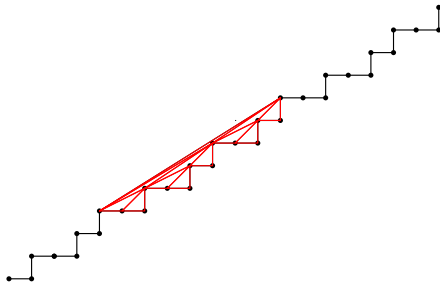
DSS of slope  $8/5$



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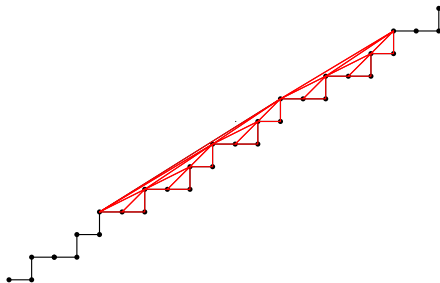
DSS of slope  $8/5$



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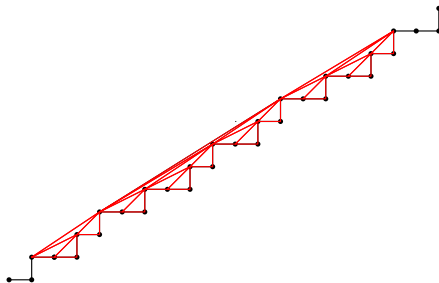
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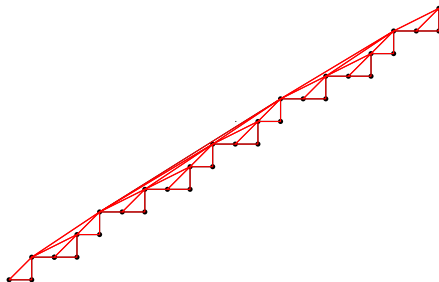
DSS of slope  $8/5$



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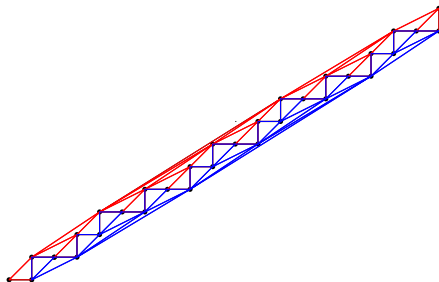
DSS of slope  $8/5$



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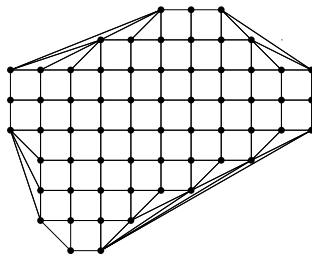
DSS of slope  $8/5$



# Delaunay triangulation computation

- ▶ Pattern
- ▶ DSS
- ▶ Convex digital object

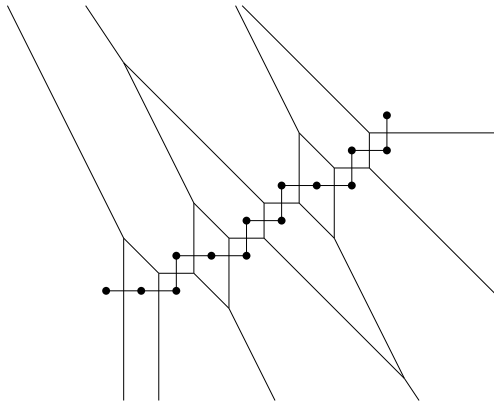
## Convex digital object





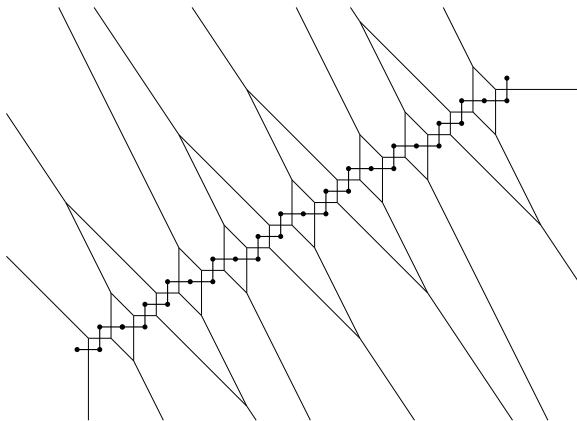
# Voronoi diagram computation

Pattern



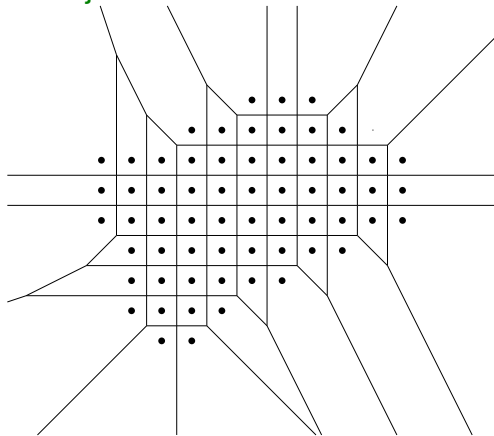
# Voronoi diagram computation

DSS



# Voronoi diagram computation

Convex digital object



# Output-sensitive algorithm for $\alpha$ -hull computation

## Definition

For all  $\alpha \in [0, 1]$ , the  $\alpha$ -hull of a set  $S$  is defined as the intersection of all closed complements of discs of radius  $1/\alpha$  that contain all the points of  $S$ .

figures pour différents  $\alpha$  de 0 à 1.

# Number of vertices of $\alpha$ -hulls: question

Let  $X$  be the Gauss digitization of a convex body of diameter  $\delta$ .  
Let  $\#V(\alpha)$  be the number of the vertices of the  $\alpha$ -hull of  $X$ .

- ▶  $\#V(1) = O(\delta)$  is trivial.
- ▶ it is known that  $\#V(0) = O(\delta^{\frac{2}{3}})$  (convex hull).
- ▶ is there a generic formula for all  $\alpha \in [0, 1]$  such that:  
 $\#V(\alpha) = O(f(\delta, \alpha))$  ?