# SICP

God's Programming Book

Lecture-11 Trees



# Trees

Slides Adapted from cs61a of UC Berkeley

# **Box-and-Pointer Notation**

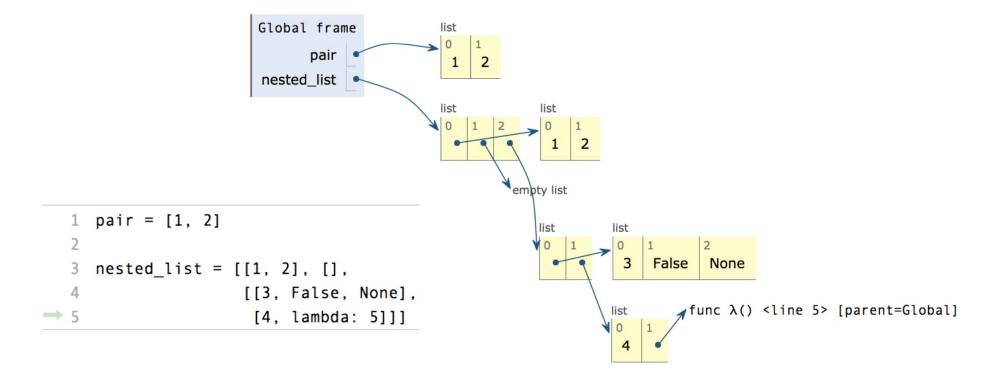
## The Closure Property of Data Types

- A method for combining data values satisfies the closure property if: The result of combination can itself be combined using the same method
- Closure is powerful because it permits us to create hierarchical structures
- Hierarchical structures are made up of parts, which themselves are made up of parts, and so on

Lists can contain lists as elements (in addition to anything else)

# Box-and-Pointer Notation in Environment Diagrams

Lists are represented as a row of index-labeled adjacent boxes, one per element Each box either contains a primitive value or points to a compound value



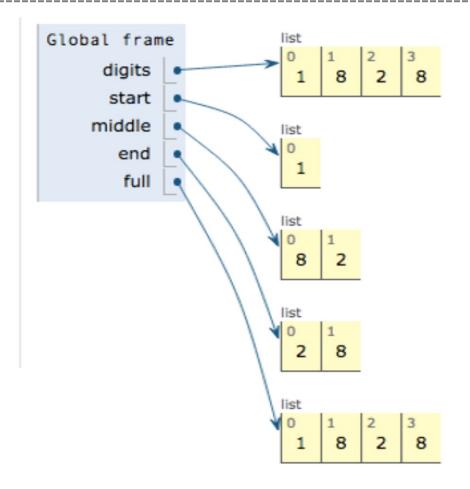
# Slicing

(Demo)

## Slicing Creates New Values

```
1 digits = [1, 8, 2, 8]
2 start = digits[:1]
3 middle = digits[1:3]
4 end = digits[2:]

5 full = digits[:]
```



# Processing Container Values

### Sequence Aggregation

Several built-in functions take iterable arguments and aggregate them into a value

- sum(iterable[, start]) -> value
   Return the sum of a 'start' value (default: o) plus an iterable of numbers.
- max(iterable[, key=func]) -> value

With a single iterable argument, return its largest item.

With two or more arguments, return the largest argument.

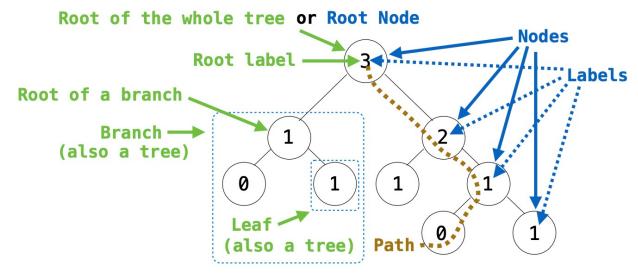
all(iterable) -> bool

Return True if bool(x) is True for all values x in the iterable. If the iterable is empty, return True.

# Trees

### Tree Abstraction

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#### Recursive description (wooden trees):

A tree has a root label and a list of branches
Each branch is a tree
A tree with zero branches is called a leaf
A tree starts at the root

#### Relative description (family trees):

Each location in a tree is called a **node**Each **node** has a **label** that can be any value

One node can be the **parent/child** of another

The top node is the **root node** 

People often refer to labels by their locations: "each parent is the sum of its children"

### Implementing the Tree Abstraction

```
def tree(label, branches=[]):
                                                           · A tree has a root label
                                    Verifies the
    for branch in branches:
                                                             and a list of branches
                                   tree definition
        assert is_tree(branch)

    Each branch is a tree

    return [label] + list(branches)
                       Creates a list
def label(tree):
                       from a sequence
    return tree[0]
                         of branches
def branches(tree):
                      Verifies that
    return tree[1:]
                      tree is bound
                        to a list
def is_tree(tree):
                                                      >>> tree(3, [tree(1),
    if type(tree) != list or len(tree) < 1:</pre>
                                                                    tree(2, [tree(1),
        return False
                                                                             tree(1)1)1)
                                                      [3, [1], [2, [1], [1]]]
    for branch in branches(tree):
        if not is_tree(branch):
                                                 def is leaf(tree):
            return False
                                                     return not branches(tree)
    return True
```

# Tree Processing

### Tree Processing Uses Recursion

Processing a leaf is often the base case of a tree processing function

The recursive case typically makes a recursive call on each branch, then aggregates

```
def count_leaves(t):
    """Count the leaves of a tree."""
    if is_leaf(t):
        return 1
    else:
        branch_counts = [count_leaves(b) for b in branches(t)]
        return sum(branch_counts)
```

### **Discussion Question**

Implement leaves, which returns a list of the leaf labels of a tree Hint: If you sum a list of lists, you get a list containing the elements of those lists

```
>>> sum([ [1], [2, 3], [4] ], []) def leaves(tree):
[1, 2, 3, 4]
                                         """Return a list containing the leaf labels of tree.
>>> sum([ [1] ], [])
                                         >>> leaves(fib tree(5))
                                         [1, 0, 1, 0, 1, 1, 0, 1]
>>> sum([ [[1]], [2] ], [])
[[1], 2]
                                         if is_leaf(tree):
                                             return [label(tree)]
                                         else:
                                             return sum(List of leaf labels for each branch, [])
     branches(tree)
                                                 [b for b in branches(tree)]
     leaves(tree)
                                                 [s for s in leaves(tree)]
      [branches(b) for b in branches(tree)]
                                                 [branches(s) for s in leaves(tree)]
     [leaves(b) for b in branches(tree)]
                                                 [leaves(s) for s in leaves(tree)]
```

### **Creating Trees**

A function that creates a tree from another tree is typically also recursive

```
def increment_leaves(t):
    """Return a tree like t but with leaf labels incremented."""
    if is_leaf(t):
        return tree(label(t) + 1)
    else:
        bs = [increment_leaves(b) for b in branches(t)]
        return tree(label(t), bs)

def increment(t):
    """Return a tree like t but with all labels incremented."""
    return tree(label(t) + 1, [increment(b) for b in branches(t)])
```

# Example: Printing Trees

(Demo)

# Example: Summing Paths

(Demo)

# Thanks for Listening