

# Assignment -1

Answer: 1

$\langle \text{float} \rangle \langle \text{id, limited Square} \rangle \langle [ \rangle \langle \text{id, x} \rangle \langle ] \rangle \langle \{ \rangle$   
 $\langle \text{float} \rangle \langle \text{id, x} \rangle$   
 $\langle \text{return} \rangle \langle \langle \rangle \rangle \langle \text{id, x} \rangle \langle \text{op, "<="} \rangle$   
 $\langle \text{num, -10.0} \rangle \langle \text{op, "11"} \rangle \langle \text{id, x} \rangle \langle \text{num, 10.0} \rangle$   
 $\langle \rangle \rangle \langle \text{op, "?"} \rangle \langle \text{num, 100} \rangle \langle \text{op, ":"} \rangle \langle \text{id, x} \rangle$   
 $\langle \text{op, "*"} \rangle \langle \text{id, x} \rangle \langle \} \rangle$

Answer: 2 (a)

$$S \rightarrow 0S1 \mid 01$$

$$S \rightarrow 0S1 \rightarrow 00S11 \rightarrow 000111$$

language generated is equal number of 0s followed by equal number of 1s

(b)  $S \rightarrow +SS \mid -SS \mid a$

$$S \rightarrow a$$

$$S \rightarrow +SS \rightarrow +aa ; S \rightarrow +SS \rightarrow +(-SS) S \rightarrow +aaa$$

Here the language generated is  $(+,-)^* a^+$

(c)  $S \rightarrow S(S) S \mid \epsilon$

$$S \rightarrow \epsilon$$

$$S \rightarrow S(S)S \rightarrow ()$$

$$S \rightarrow S(S)S \rightarrow S(S(S)S)S$$

$$\rightarrow (())$$

Left Expansion

$$S \rightarrow S(S)S \rightarrow S(S)S(S)S \rightarrow (())$$

Right Expansion

$$S \rightarrow S(S)S \rightarrow S(S)S(S)S \rightarrow (())$$

$\therefore$  The language generated contains same number of left and right brackets (or parenthesis).

$$(d) S \rightarrow aSbS \mid bSaS \mid \epsilon$$

$$S \rightarrow ab,$$

$$S \rightarrow ba;$$

$$S \rightarrow aSbS$$

$$S \rightarrow ~~abba~~ aSbSbS \rightarrow aabbb$$

$$S \rightarrow bSaS \rightarrow bSaSbS \rightarrow baabb$$

language generated contains equal number of a and b and can also generate empty string.

$$(e) S \rightarrow a \mid S+S \mid SS \mid S^* \mid (S)$$

$$S \rightarrow a$$

$$S \rightarrow (a)$$

$$S \rightarrow a+a$$

$$S \rightarrow S+S$$

$$S \rightarrow aa$$

$$S \rightarrow SS+SS \rightarrow aa+aa$$

$$S \rightarrow a^*$$

∴ The set of all strings of +, \*, 'a' and symmetric parenthesis and plus is not beginning or end of partition, multiplication is not the beginning of partition.

Answer: 3

$$S \rightarrow [S|S] [{}|{}] [()|()] [{}|{}] [C|c] [T|t]$$

Answer: 4

$$S \rightarrow (L) | a$$

$$L \rightarrow L, S | S$$

With string (a,a), a, (a)

(a)  $S \rightarrow (L) \rightarrow (L, S) \rightarrow (L, S, S) \rightarrow (S, S, S)$

$(S, S, S) \leftarrow (L, S, S) \leftarrow (L, S, S) \leftarrow (L, S, S)$

$(a, S, S) \rightarrow (a, a, S, S) \rightarrow (a, a, a, (L))$

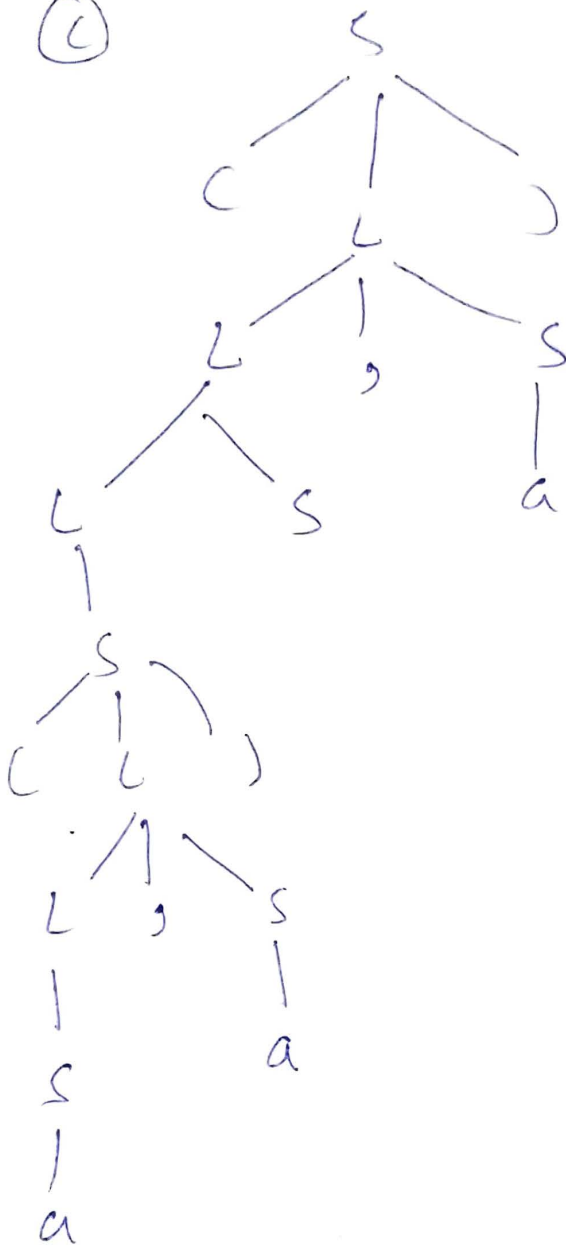
$(a, a, a, (a)) \leftarrow (a, a, a, (S)) \leftarrow (a, a, a, (S))$

(b)  $S \rightarrow (L) \rightarrow (L, S) \rightarrow (L, (L)) \rightarrow (L, (a)) \rightarrow (L, S, (a))$

$(L, a, a) \leftarrow (S, a, a) \leftarrow (L, a, (a))$

$((L, S), a, (a)) \rightarrow ((S, S), a, (a)) \rightarrow ((a, a), a, (a))$

(c)



(d) Unambiguous.

(e) Like a tuple in Python.

s

5 (1)

$$S \rightarrow S + S \mid SS \mid (S) \mid S^* \mid a$$

↓ left factoring

$$S \rightarrow SA \mid (S) \mid a$$

$$A \rightarrow +S \mid \varepsilon \mid *$$

↓ further extraction

$$S \rightarrow SA \mid T$$

$$A \rightarrow +S \mid \varepsilon \mid *$$

$$T \rightarrow (S) \mid a$$

↓ eliminate Left Recursion

$$S \rightarrow TB$$

$$B \rightarrow AB \mid \varepsilon$$

$$A \rightarrow +S \mid TB \mid *$$

$$T \rightarrow (S) \mid a$$

$$\text{First}(S) = \text{First}(T) = \{ (, a \}$$

$$\text{First}(B) = \{ \text{First}(A), \varepsilon \} = \{ +, (, a, *, \varepsilon \}$$

$$\text{First}(A) = \{ +, \text{First}(T), * \} = \{ +, (, a, * \}$$

$$\text{Follow}(S) = \{ \$, ), +, (, a, * \}$$

$$\text{Follow}(B) = \{ \$, ) \}$$

$$\text{Follow}(A) = \{ +, (, a, *, ), \$ \}$$

$$\text{Follow}(T) = \{ +, (, a, *, \$ \}$$

## Parse Table

	(	)	+	*	a	\$
S	$S \rightarrow TB$				$S \rightarrow TB$	
B	$B \rightarrow AB$		$B \rightarrow AB$	$B \rightarrow AB$	$B \rightarrow AB$	$B \rightarrow \epsilon$
A	$A \rightarrow TB$		$A \rightarrow +S$	$A \rightarrow *$	$A \rightarrow TB$	
T	$T \rightarrow (S)$				$T \rightarrow a$	

(2)

$$S \rightarrow (L) | a$$

$$L \rightarrow L, S | S$$

eliminate left recursion

$$S \rightarrow (L) | a$$

$$L \rightarrow SA$$

$$A \rightarrow , SA | \epsilon$$

first  $(S) = \{(, a\}$

first  $(L) = \{(, a\}$

first  $(A) = \{, , \epsilon\}$

$$\text{follow}(S) = \{\$, , )\}$$

$$\text{follow}(L) = \{)\}$$

$$\text{follow}(A) = \{,\}$$

## Parse Table

	(	)	a	,	\$
S	$S \rightarrow (L)$		$S \rightarrow a$		
L	$L \rightarrow SA$		$L \rightarrow SA$		
A	$A \rightarrow \epsilon$			$A \rightarrow , SA$	