	Recurrence Equations
1	Master Theorem Analysis a≥1, b>1, f(n) asymptotically positive
	$T(n) = a T\left(\frac{n}{n}\right) + f(n)$
1-	b = Size of original problem  Size of subproblem    Division   Example: a = 2, b = 2   Work   2 problems of \frac{1}{2} \lequesize
-	f(n)=  Fixtra    Combine   Work    Showing case   Combine   Combin
	special case $a = b \Rightarrow leg_{b}(a) = leg_{b}(b) = 1$
	Case 1: $f(n) = O(n^{1-\epsilon}) \Rightarrow T(n) = O(n)$ "Division Dominates
	Case 2: $f(n) = \theta(n) = f(n) = \theta(n \log n)$ Division + Combine Partnership
	Case 3: $f(n) = \Omega(n^{1+\epsilon}) \Rightarrow T(n) = \Theta(f(n))$ "Combine" Dominates
	t regularly
	General case: $log_b(a) = c$
	Case 1 $f(n) = O(n^{c-\epsilon}) \Rightarrow T(n) = \Theta(n^c)$ Division Dominates
	2: $f(n) = \theta(n^c) \Rightarrow T(n) = \theta(n^c \log n)$ Division + Combine
	3: $f(n) = SL(n^{C+C}) \Rightarrow T(n) = O(f(n))$ + regularly
	Need to know how to compute log, (a)
a	1 · · 56 35 26 6 6 62 63 · 6k ·
log \$(9)	0 1/2 1 2 3 k