

1. $C=1 \Rightarrow 2\sqrt{n} = 2\sqrt{n}$

When $n=5$ $\Rightarrow \sqrt{5} \approx 2.236 = 4.472 < 5$

When $n=4$ $\Rightarrow \sqrt{4} = 4 = 4$

Therefore the smallest $n=5$

2. ^{Proof:} When $\lim_{n \rightarrow \infty} f(n) = L$ and $0 < L < \infty$

also ~~also~~ $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \frac{L}{L_1}$, $0 < L_1 < \infty$.

$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \frac{L}{L_1}$ $0 < \frac{L}{L_1} < \infty$

By the theorem we learned, ~~$f(n) = \theta(g)$~~ $f(n) = \theta(g(n))$

This completes the proof \square

3. ~~if~~ Proof.

If $f(n) = o(g(n))$ then ~~$0 < c_1 g(n) < f(n) < c_2 g(n)$~~

$0 \leq f(n) < Cg(n)$

$\ln f(n) < \ln Cg(n)$

$\ln f(n) \leq \ln C + \ln g(n)$

We can easily find a c' that makes

$c' \ln g(n) > \ln C + \ln g(n)$, e.g. $c' > \frac{\ln C}{\ln g(n)} + 1$

Hence, we get when $f(n) = o(g(n))$, $\ln f(n) = \text{less than } o(\ln g(n))$ \square

4. ~~$\frac{d}{dx}$ Proof:~~

~~$f(n) = (2n)^n$
 $\ln f(n) = \ln (2n)^n$~~

~~$\frac{d}{dn} f(n) = \frac{d}{dn} (2n)^n$
 $= \frac{d}{dn} e^{n \ln(2n)}$
 $= \frac{d}{dn} e^{n \ln(2n)} \cdot \frac{d}{dn} (n \ln(2n))$
 $= e^{n \ln(2n)} \cdot (\ln(2n) + 1)$~~

$f(n) = (2n)^n$

$\ln f(n) = \ln (2n)^n$

$f(n) = e^{n \ln(2n)}$

$f(n)' = e^{n \ln(2n)} \cdot \frac{d}{dn} (n \ln(2n))$

$= e^{n \ln(2n)} \cdot (\ln(2n) + 1)$

$= (2n)^n \cdot (\ln(2n) + 1)$