Wednesday January 20

Proposition Declarative sentence that is true or false (not both).

Propositional variable Variable that represents a proposition.

Compound proposition New propositions formed from existing propositions (potentially) using

logical operators.

Truth table Table with 1 row for each of the possible combinations of truth values

of the input and an additional column that shows the truth value of

the result of the operation corresponding to a particular row.

Note: A propositional variable is one example of a compound proposition.

Logical operators aka propositional connectives

| Conjunction | AND | \wedge | $\ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ $ | 2 inputs | Evaluates to T when both inputs are T |
|--------------|-----|----------|--|----------|---|
| Exclusive or | XOR | \oplus | \oplus | 2 inputs | Evaluates to T when exactly one of inputs is T |
| Disjunction | OR | \vee | \lor | 2 inputs | Evaluates to T when at least one of inputs is T |
| Negation | NOT | \neg | \label{lnot} | 1 input | Evaluates to T when its input is F |

| Inp | out | | Output | | | |
|----------------|-----|--------------|--------------|-------------|----------------|--------------------|
| | | Conjunction | Exclusive or | Disjunction | Input | Output |
| p | q | $p \wedge q$ | $p\oplus q$ | $p \lor q$ | | Output Negation |
| \overline{T} | T | T | F | T | p | $\parallel \neg p$ |
| T | F | F | T | T | \overline{T} | F |
| F | T | F | T | T | F | \parallel T |
| F | F | F | F | F | | 11 |

| J | npu | t | Output |
|----------------|-----|---|---|
| p | q | r | $\mid (p \wedge q) \oplus (\ (p \oplus q) \wedge r\) \mid (p \wedge q) \lor (\ (p \oplus q) \wedge r\)$ |
| \overline{T} | T | T | |
| T | T | F | |
| T | F | T | |
| T | F | F | |
| F | T | T | |
| F | T | F | |
| F | F | T | |
| F | F | F | |

Logical equivalence Two compound propositions are logically equivalent

means that they have the same truth values for all settings of truth values to their propositional variables.

Tautology A compound proposition that evaluates to true for all

settings of truth values to its propositional variables; it

is abbreviated T.

Contradiction A compound proposition that evaluates to false for all

settings of truth values to its propositional variables; it

is abbreviated F.

Contingency A compound proposition that is neither a tautology nor

a contradiction.

of the compound propositions in the table below are logically equivalent?

| Inp | out | | | Output | | |
|----------------|-----|-------------------------|-----------------------------|-------------------|------------------------|---------------|
| p | q | $\neg (p \land \neg q)$ | $\neg (\neg p \lor \neg q)$ | $(\neg p \lor q)$ | $(\neg q \lor \neg p)$ | $(p \land q)$ |
| \overline{T} | T | | | | | |
| T | F | | | | | |
| F | T | | | | | |
| F | F | | | | | |

(Some) logical equivalences) cf. Rosen pp. 26-28

$$\begin{array}{ll} p \vee q \equiv q \vee p & p \wedge q \equiv q \wedge p \\ (p \vee q) \vee r \equiv p \vee (q \vee r) & (p \wedge q) \wedge r \equiv p \wedge (q \wedge r) \\ p \wedge F \equiv F & p \vee T \equiv T & p \wedge T \equiv p & p \vee F \equiv p \\ \neg (p \wedge q) \equiv \neg p \vee \neg q & \neg (p \vee q) \equiv \neg p \wedge \neg q \end{array}$$

Commutativity Ordering of terms
Associativity Grouping of terms
Absorption aka short circuit evaluation
DeMorgan's Laws

Extra Example: Which

Can replace p and q with any compound proposition

Application: design a circuit given a desired input-output relationship.

| Inr | t | Out | rout | |
|----------------|----------------|------------------------------|-----------------------|--|
| Input | | Output $mystery_1 mystery_2$ | | |
| $\frac{P}{T}$ | $\frac{q}{T}$ | T | $\frac{mgseerg_2}{F}$ | |
| T | \overline{F} | T | \overline{F} | |
| \overline{F} | \overline{T} | \overline{F} | \overline{F} | |
| F | F | T | T | |
| | ' | I | | |

| I | npu | Output | |
|----------------|-----|--------|---|
| p | q | r | ? |
| \overline{T} | T | T | T |
| T | T | F | T |
| T | F | T | F |
| T | F | F | T |
| F | T | T | F |
| F | T | F | F |
| F | F | T | T |
| F | F | F | F |

| A compound proposition that gives output $mystery_1$ is: | |
|--|---|
| | |
| | |
| | |
| A compound proposition that gives output $mystery_2$ is: | |
| | |
| | |
| | |
| | |
| Definition A compound proposition is in disjunctive normal form (DNF) means that it is an OR of ANDs of variables and their negations. | f |
| Definition A compound proposition is in conjunctive normal form (CNF) means that it is an AND of ORs of variables and their negations. | f |
| Extra example: A compound proposition that gives output? is: | |
| | |
| | |
| | |
| | |
| | |
| | |

Review

1. (a) Consider the logic circuit



For which of the following settings(s) of input values is the output $y_1 = 0$? (Select all and only those that apply.)

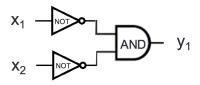
i.
$$x_1 = 0$$
, $x_2 = 0$, $x_3 = 0$, and $x_4 = 0$

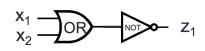
ii.
$$x_1 = 1$$
, $x_2 = 1$, $x_3 = 1$, and $x_4 = 1$

iii.
$$x_1 = 1$$
, $x_2 = 0$, $x_3 = 0$, and $x_4 = 1$

iv.
$$x_1 = 0$$
, $x_2 = 0$, $x_3 = 1$, and $x_4 = 1$

(b) Consider the logic circuits





For which of the following settings(s) of input values do the outputs of these circuits have the same value, i.e. $y_1 = z_1$? (Select all and only those that apply.)

i.
$$x_1 = 1, x_2 = 1$$

ii.
$$x_1 = 1, x_2 = 0$$

iii.
$$x_1 = 0, x_2 = 1$$

iv.
$$x_1 = 0, x_2 = 0$$

2. For each of the following propositions, indicate exactly one of:

- There is no assignment of truth values to its variables that makes it true,
- There is exactly one assignment of truth values to its variables that makes it true, or
- There are exactly two assignments of truth values to its variables that make it true, or
- ullet There are exactly three assignments of truth values to its variables that make it true, or
- $\bullet \ All$ assignments of truth values to its variables make it true.

(a)
$$x \wedge y \wedge (x \vee y)$$

(b)
$$\neg x \land y \land (x \lor y)$$

(c)
$$x \land \neg y \land (x \land y)$$

(d)
$$\neg x \land (y \lor \neg y)$$

(e)
$$x \wedge (y \vee \neg x)$$

Friday January 22

The only way to make the conditional statement $p \to q$ false is to _____

The **hypothesis** of $p \to q$ is ______ The **antecedent** of $p \to q$ is ______

The conclusion of $p \to q$ is _____ The consequent of $p \to q$ is _____

| Input | | Output | | | | | |
|-------|--------------|--------------|-------------|---------------|-----------------------|--|--|
| | Conjunction | Exclusive or | Disjunction | Conditional | Biconditional | | |
| p q | $p \wedge q$ | $p\oplus q$ | $p \lor q$ | $p \to q$ | $p \leftrightarrow q$ | | |
| T T | T = T | F | T | T | T | | |
| T F | F | T | T | F | F | | |
| F T | F | T | T | T | F | | |
| F F | F = F | F | F | $\mid T \mid$ | T | | |

Examples

$$p \to q \equiv \neg p \lor q$$
 because _____

 $p \leftrightarrow q$ is not logically equivalent to $p \land q$ because _____

 $\neg (p \leftrightarrow q) \equiv p \oplus q \text{ because} _$

 $p \to q$ is not logically equivalent to $q \to p$ because _____

 $p \leftrightarrow q \equiv q \leftrightarrow p$ because _____

The **converse** of $p \to q$ is _____

The **inverse** of $p \to q$ is ______ Which of these is logically equivalent to $p \to q$?

The **contrapositive** of $p \to q$ is _____

| Translation : Express each of the following sentences tions. | s as compound propositions, using the given proposi- |
|---|--|
| "A sufficient condition for the warranty to be good is that you bought the computer less than a year ago" | w is "the warranty is good" b is "you bought the computer less than a year ago" |
| "Whenever the message was sent from an unknown system, it is scanned for viruses." | s is "The message is scanned for viruses" u is "The message was sent from an unknown system" |
| "I will complete my to-do list only if I put a reminder in my calendar" | r is "I will complete my to-do list" c is "I put a reminder in my calendar" |

Review

- 1. For each of the following propositions, indicate exactly one of:
 - There is no assignment of truth values to its variables that makes it true,
 - There is exactly one assignment of truth values to its variables that makes it true, or
 - There are exactly two assignments of truth values to its variables that make it true, or
 - There are exactly three assignments of truth values to its variables that make it true, or
 - All assignments of truth values to its variables make it true.
 - (a) $(p \leftrightarrow q) \oplus (p \land q)$
 - (b) $(p \to q) \lor (q \to p)$
 - (c) $(p \to q) \land (q \to p)$
 - (d) $\neg (p \rightarrow q)$
- 2. **Definition**: A collection of compound propositions is called **consistent** if there is an assignment of truth values to the propositional variables that makes each of the compound propositions true.

For each of the following system specifications, identify the compound propositions that give their translations to logic and then determine if the translated collection of compound propositions is consistent.

(a) Specification: If the computer is out of memory, then network connectivity is unreliable. No disk errors can occur when the computer is out of memory. Disk errors only occur when network connectivity is unreliable.

Translation: M = "the computer is out of memory"; N = "network connectivity is unreliable"; D = "disk errors can occur".

i.

$$\neg M \to N$$
$$\neg D \to M$$
$$D \to N$$

ii.

$$M \to \neg N$$
$$\neg D \land M$$
$$N \to D$$

iii.

$$M \to N$$
$$M \to \neg D$$
$$\neg N \to \neg D$$

(b) Specification: Whether you think you can, or you think you can't - you're right. ¹ Translation: T= "you think you can"; C= "you can".

i.

$$T \to C$$
$$\neg T \to \neg C$$

ii.

$$T \wedge C \\ \neg T \wedge \neg C$$

iii.

$$T \to \neg T$$
$$C \to \neg C$$

(c) Specification: A secure password must be private and complicated. If a password is complicated then it will be hard to remember. People write down hard-to-remember passwords. If a password is written down, it's not private. The password is secure.

Translation: S = "the password is secure"; P = "the password is private"; C = "the password is complicated"; H = "the password is hard to remember"; W = "the password is written down".

i.

$$\neg (P \land C) \rightarrow \neg S$$

$$C \rightarrow H$$

$$W \land H$$

$$W \rightarrow \neg P$$

$$S$$

ii.

$$(P \land C) \rightarrow S$$

$$C \rightarrow H$$

$$W \rightarrow H$$

$$W \rightarrow P$$

$$S$$

iii.

$$S \to (P \land C)$$

$$C \to H$$

$$H \to W$$

$$W \to \neg P$$

$$S$$

¹Henry Ford