

Constraint-based Modeling of Cellular Networks

Exercise 3 – Linear optimization

3. 11. 2022

Translate text into LP model

- Two products – pants and jackets
- Material: cotton and polyester

To manufacture one pant or one jacket (i.e. 1 unit), following quantities are required:

- Each unit of pants requires 1m^2 of cotton and 2m^2 of polyester
- Each unit of jackets requires 1.5m^2 of cotton and 1m^2 of polyester

The manufacturer has a total 750m^2 of cotton and 1000m^2 of polyester. The price is fixed at

- 50€ per pant
- 40€ per jacket

What is the number of pants and jackets that the manufacturer must produce to obtain a maximum sale?

Translate text into LP model

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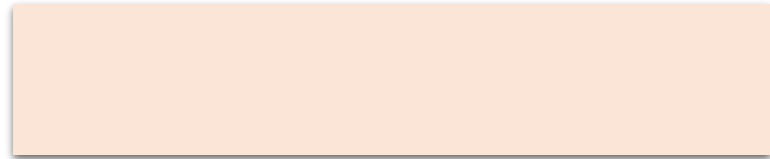
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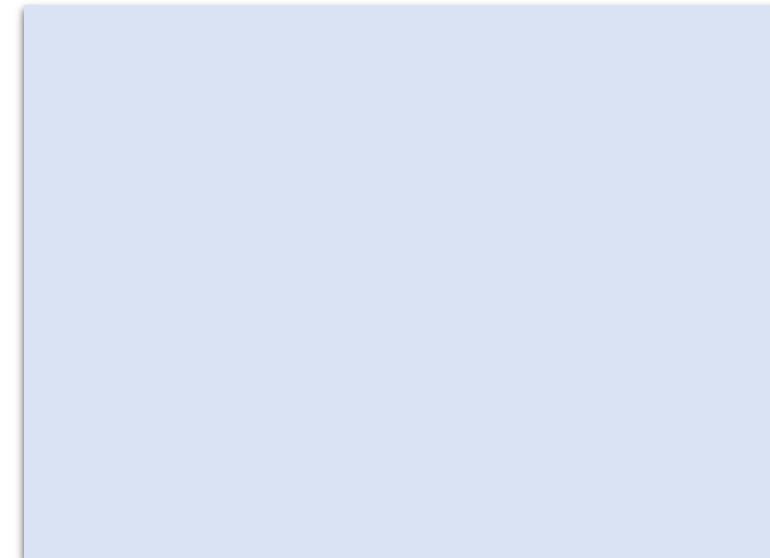
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objective

s.t.



constraints

Q1: What are our variables?

Q2: What constrains the variables?

- Two products – pants and jackets
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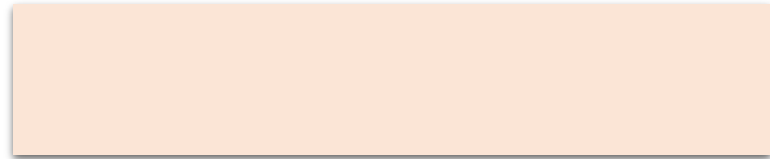
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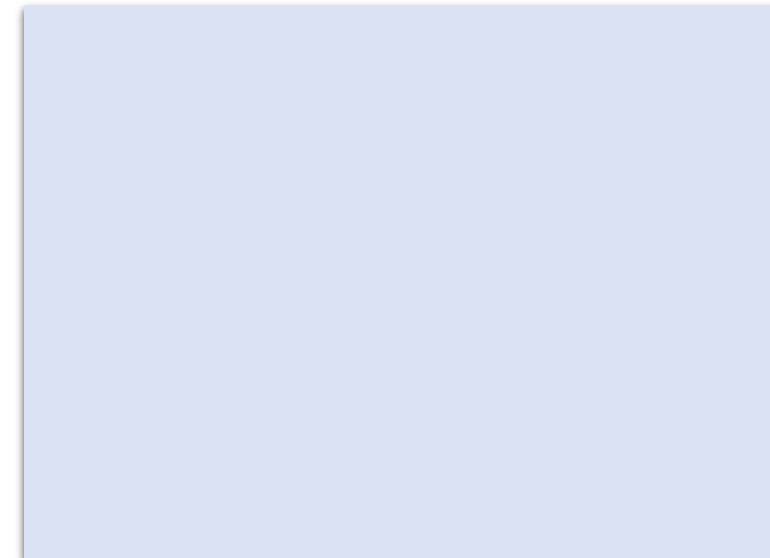
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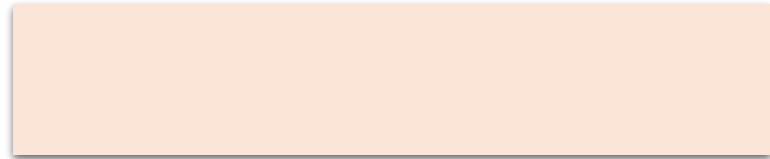
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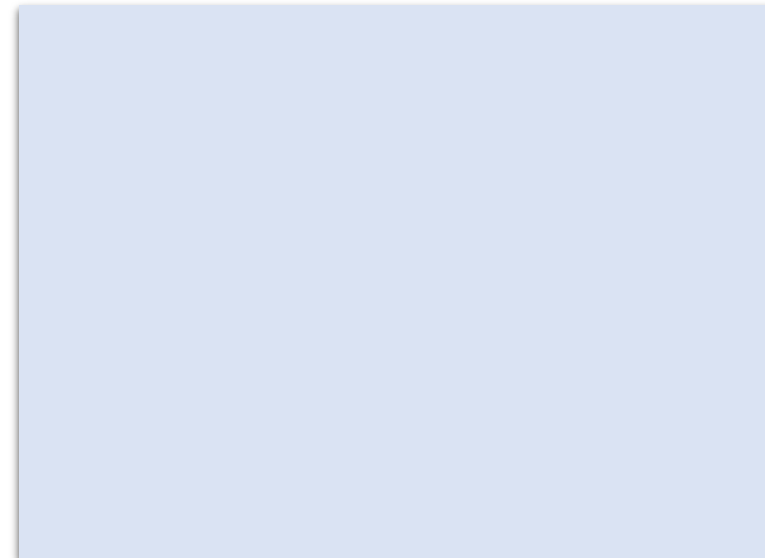
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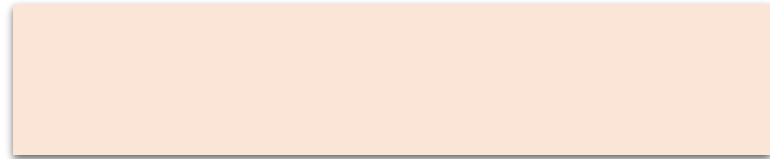
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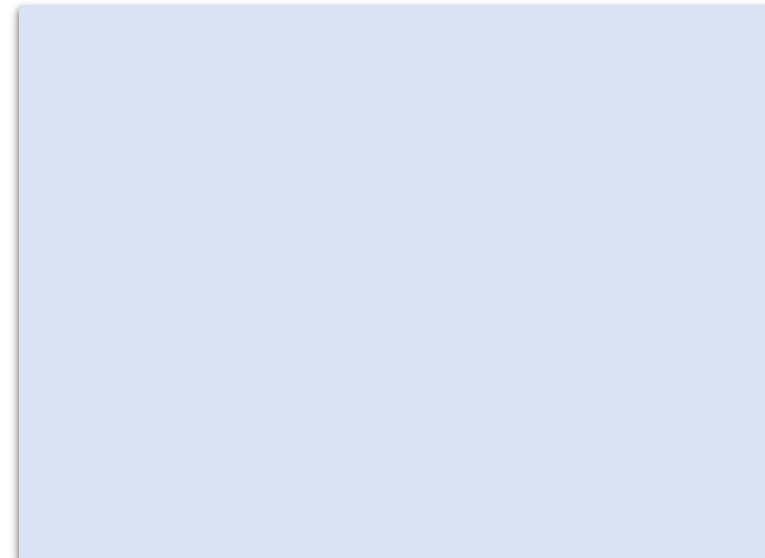
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s.t.



constraints

Availability of material constrains product production

- Two products – pants (P) and jackets (J)
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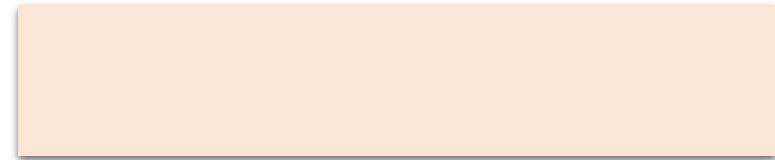
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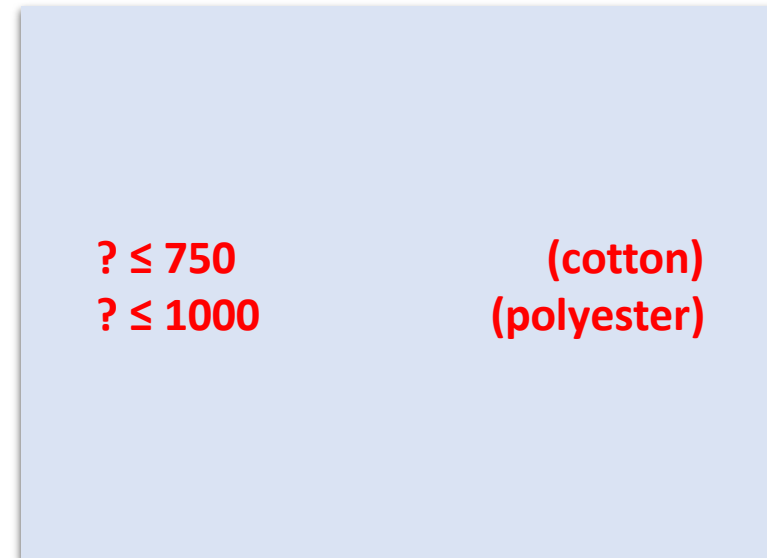
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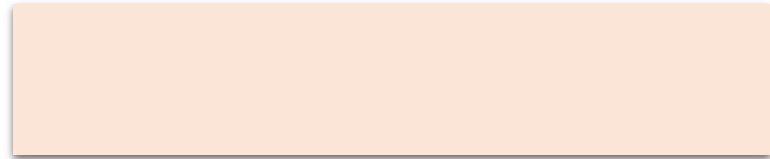
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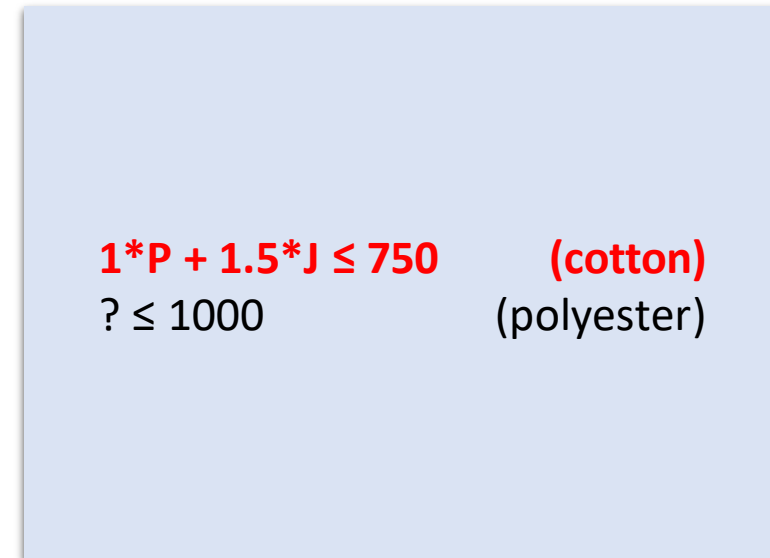
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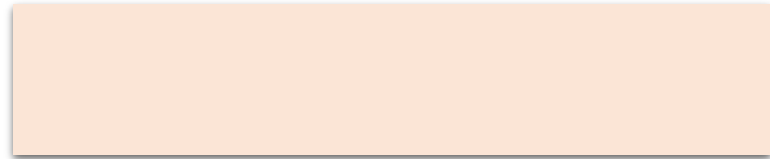
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objective

s.t.

$$\begin{aligned} 1*P + 1.5*J &\leq 750 && \text{(cotton)} \\ \mathbf{2*P + 1*J} &\leq \mathbf{1000} && \mathbf{(polyester)} \end{aligned}$$

constraints

Lower and upper bounds

- Two products – pants (P) and jackets (J)
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$$\max z = 50 P + 40 J$$

objective

s.t.

$$1 \cdot P + 1.5 \cdot J \leq 750 \quad (\text{cotton})$$

$$2 \cdot P + 1 \cdot J \leq 1000 \quad (\text{polyester})$$

$$P \geq 0$$

$$J \geq 0$$

constraints

The objective

- Two products – pants and jackets
- Material: cotton and polyester

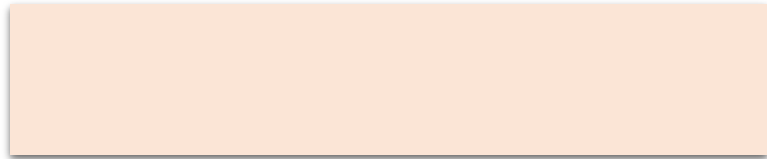
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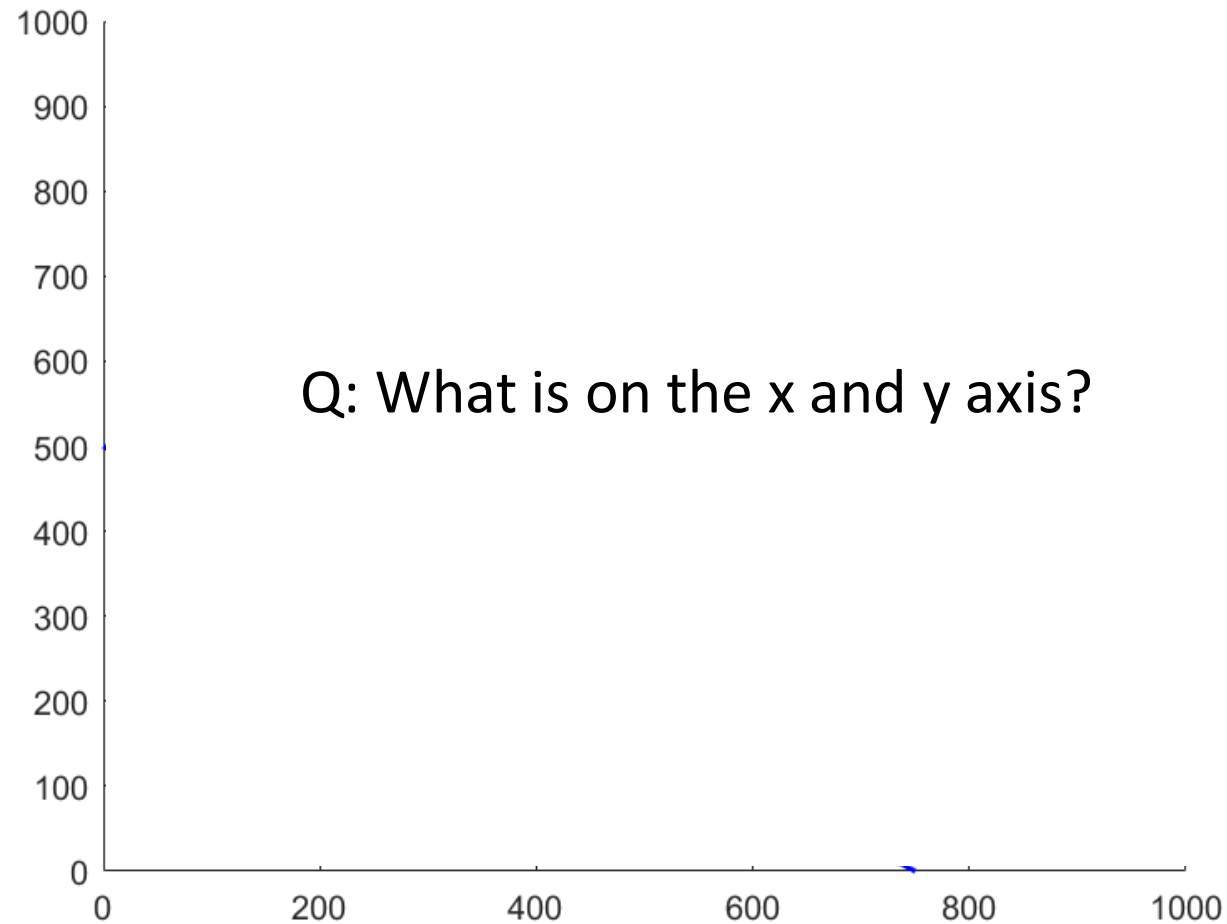
$$2 \cdot P + 1 \cdot J \leq 1000 \quad (\text{polyester})$$

$$P \geq 0$$

$$J \geq 0$$

constraints

Graphical solution



$$\max z = 50 P + 40 J$$

s.t.

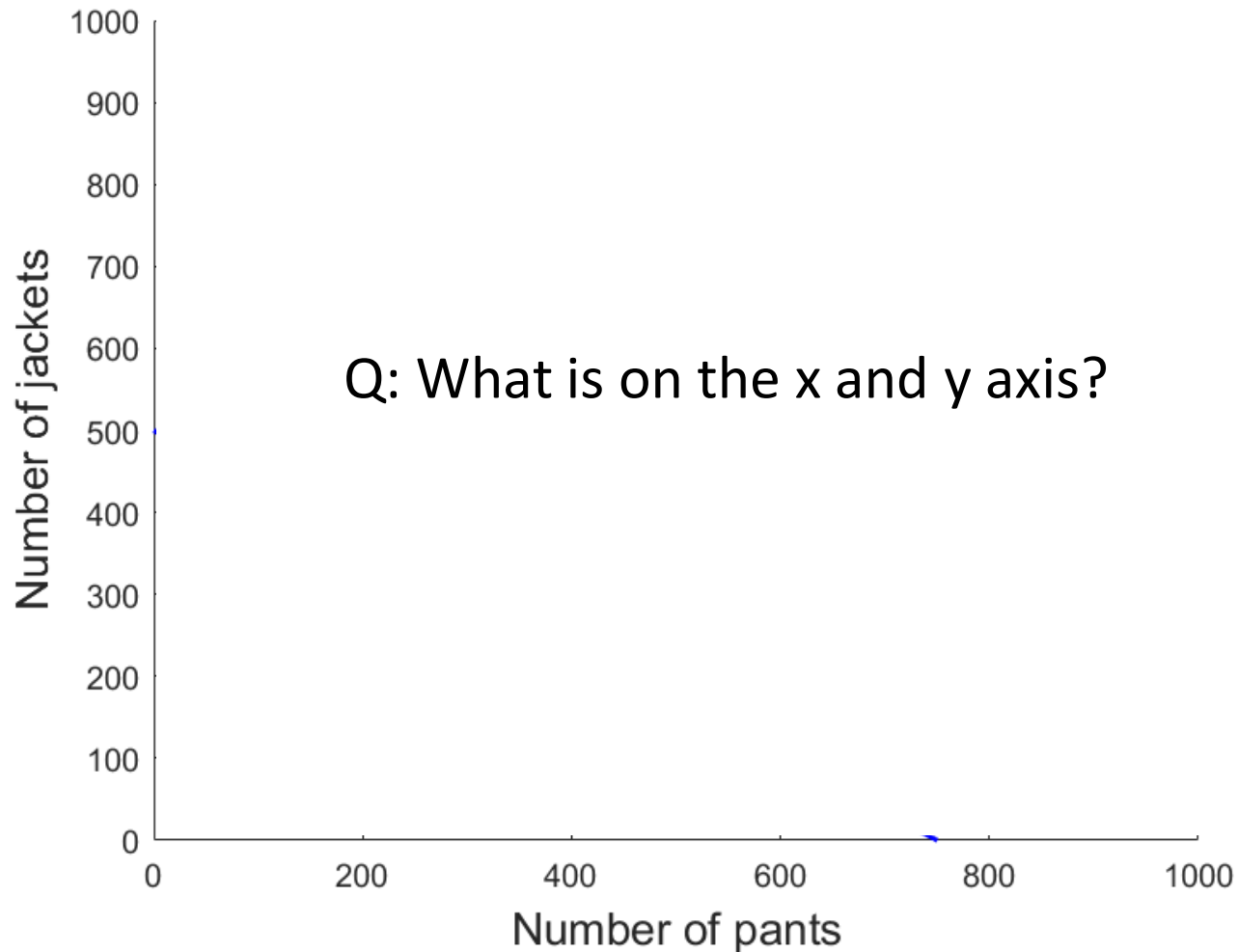
$$1 * P + 1.5 * J \leq 750 \quad (\text{cotton})$$

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$$P \geq 0$$

$$J \geq 0$$

Graphical solution



$$\max z = 50 P + 40 J$$

s.t.

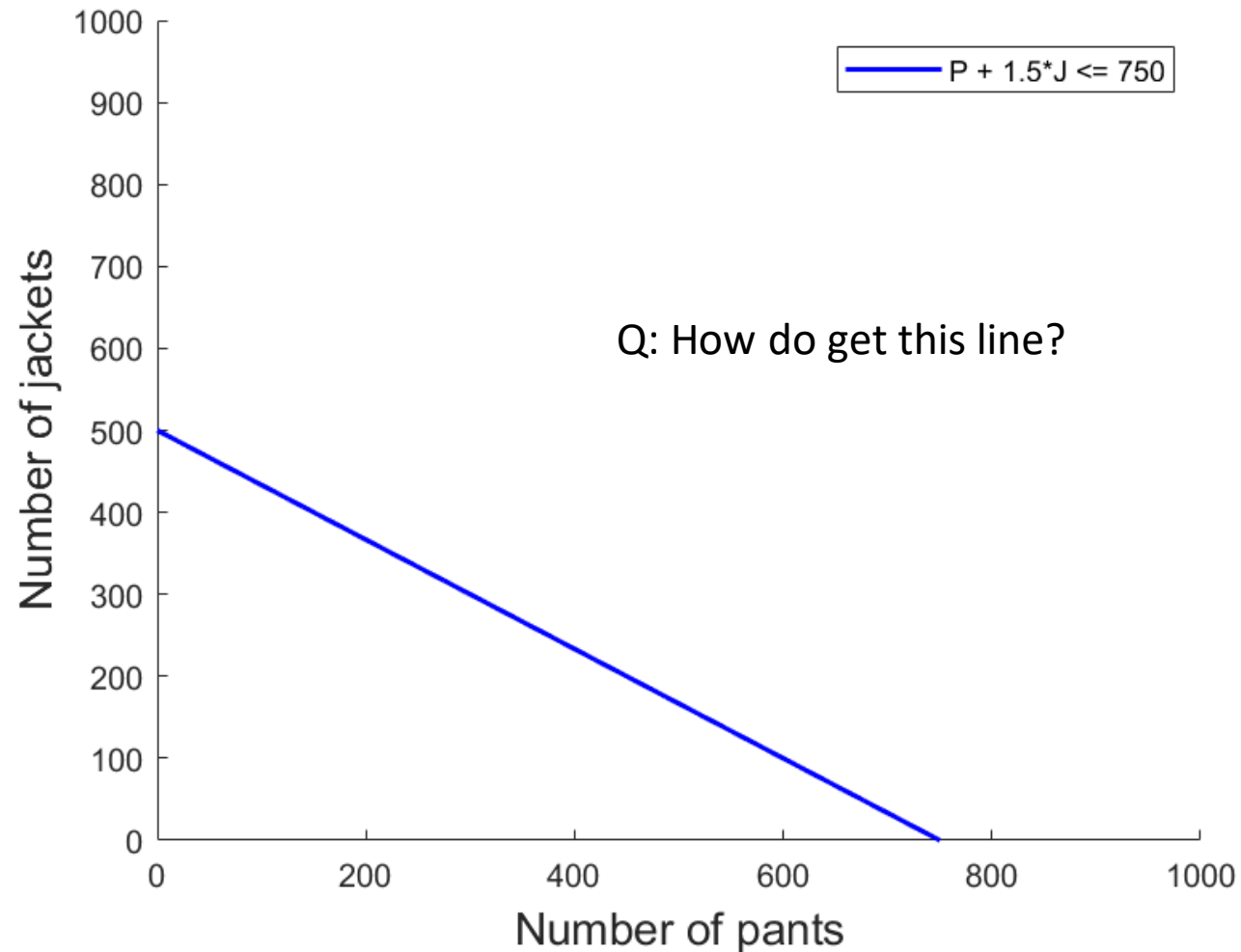
$$1 * P + 1.5 * J \leq 750 \quad (\text{cotton})$$

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$$P \geq 0$$

$$J \geq 0$$

Graphical solution – constraint 1



$$\max z = 50 P + 40 J$$

s.t.

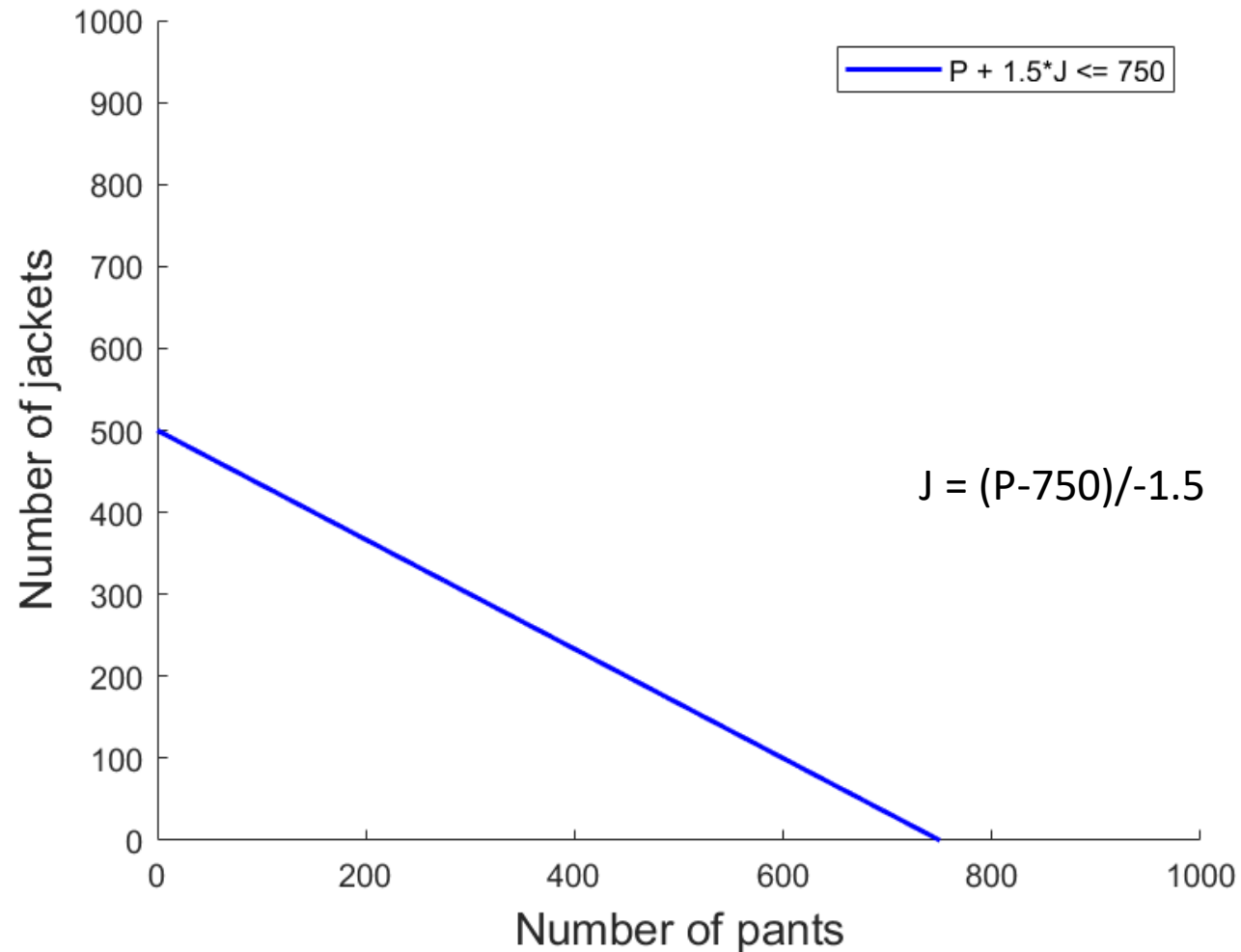
$$1 * P + 1.5 * J \leq 750 \quad (\text{cotton})$$

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$$P \geq 0$$

$$J \geq 0$$

Graphical solution – constraint 1



$$\max z = 50P + 40J$$

s.t.

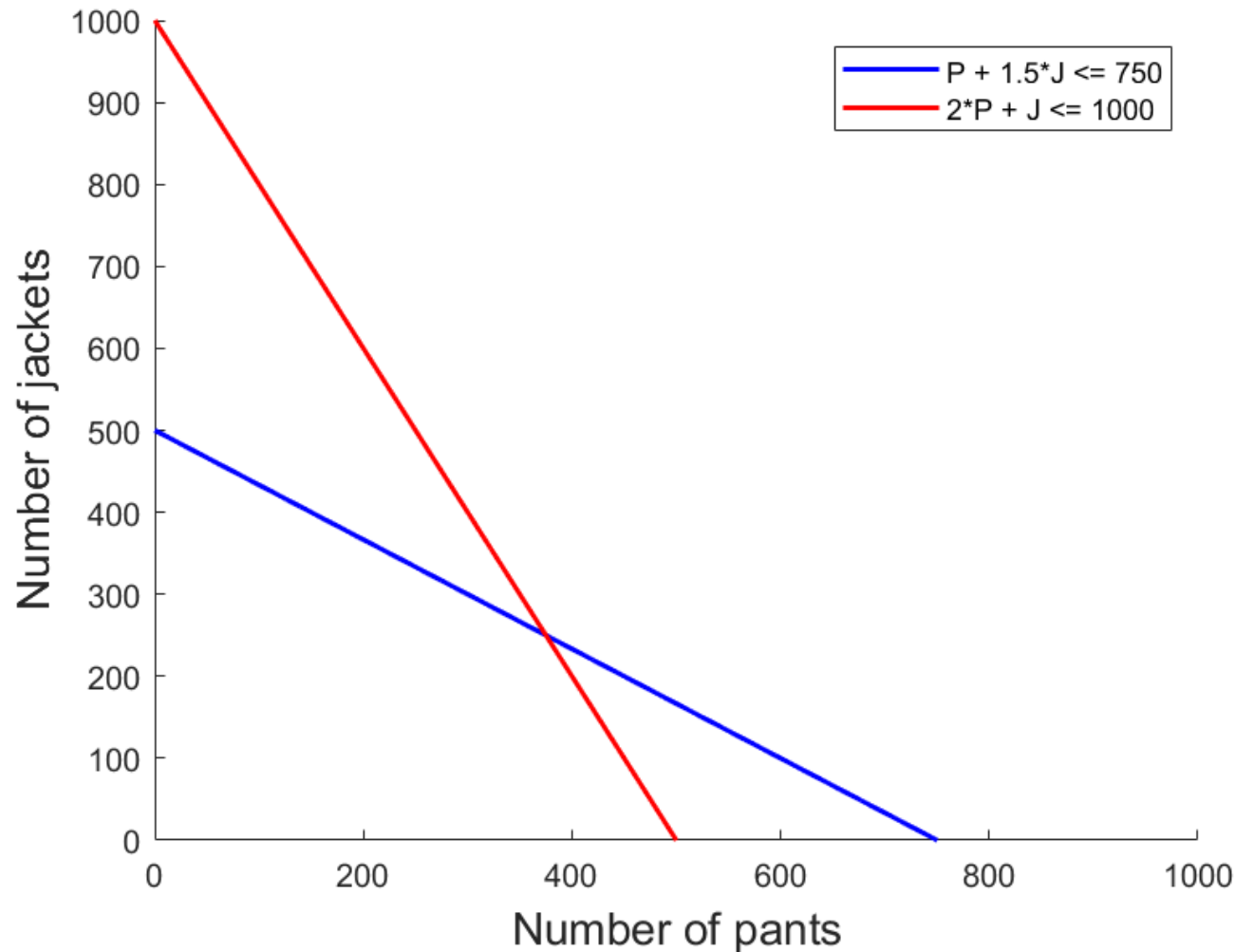
$$1 \cdot P + 1.5 \cdot J \leq 750 \quad (\text{cotton})$$

$$2 \cdot P + 1 \cdot J \leq 1000 \quad (\text{polyester})$$

$$P \geq 0$$

$$J \geq 0$$

Graphical solution – constraint 2



$$\max z = 50P + 40J$$

s.t.

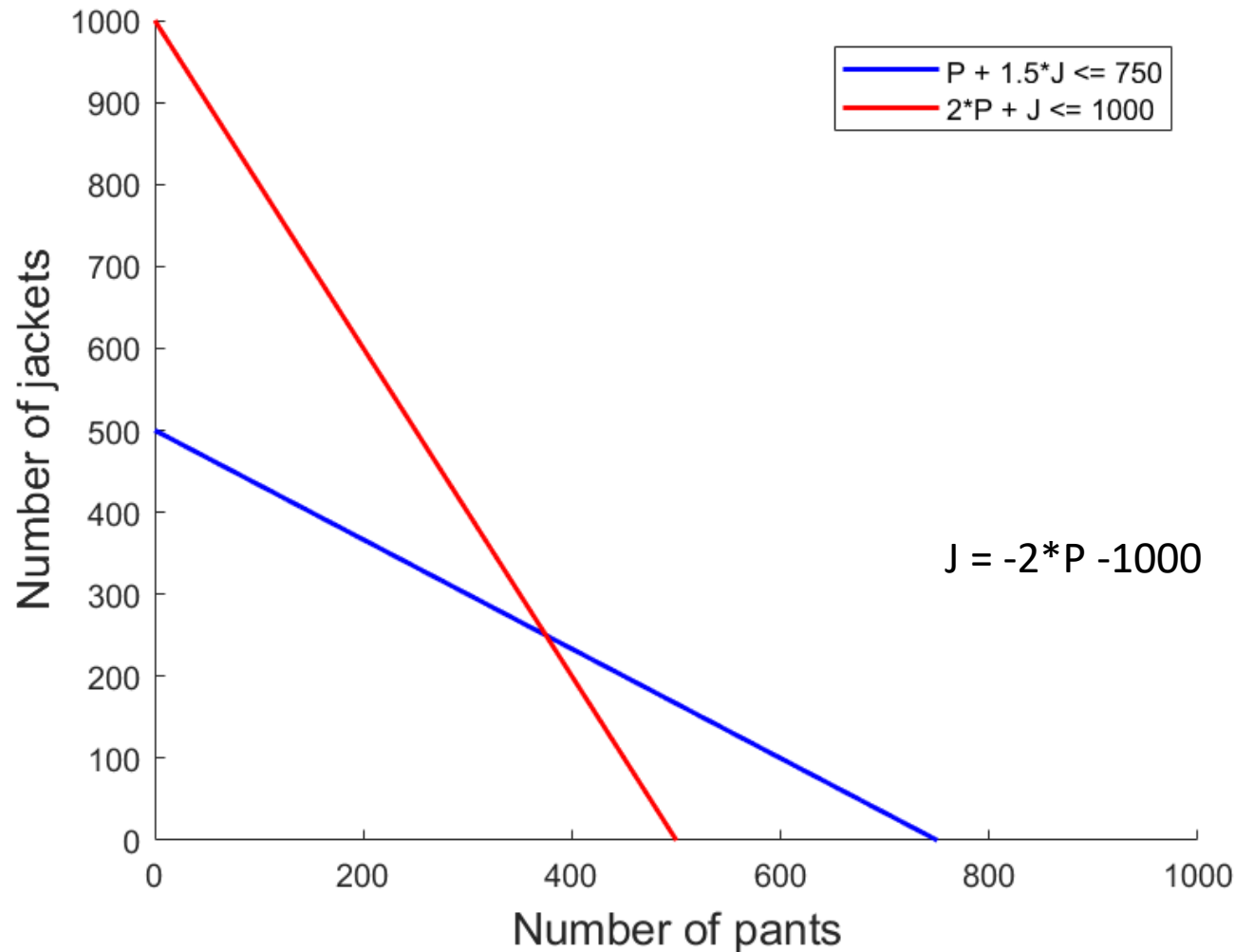
$$1P + 1.5J \leq 750 \quad (\text{cotton})$$

$$2P + 1J \leq 1000 \quad (\text{polyester})$$

$$P \geq 0$$

$$J \geq 0$$

Graphical solution – constraint 2



$$\max z = 50P + 40J$$

s.t.

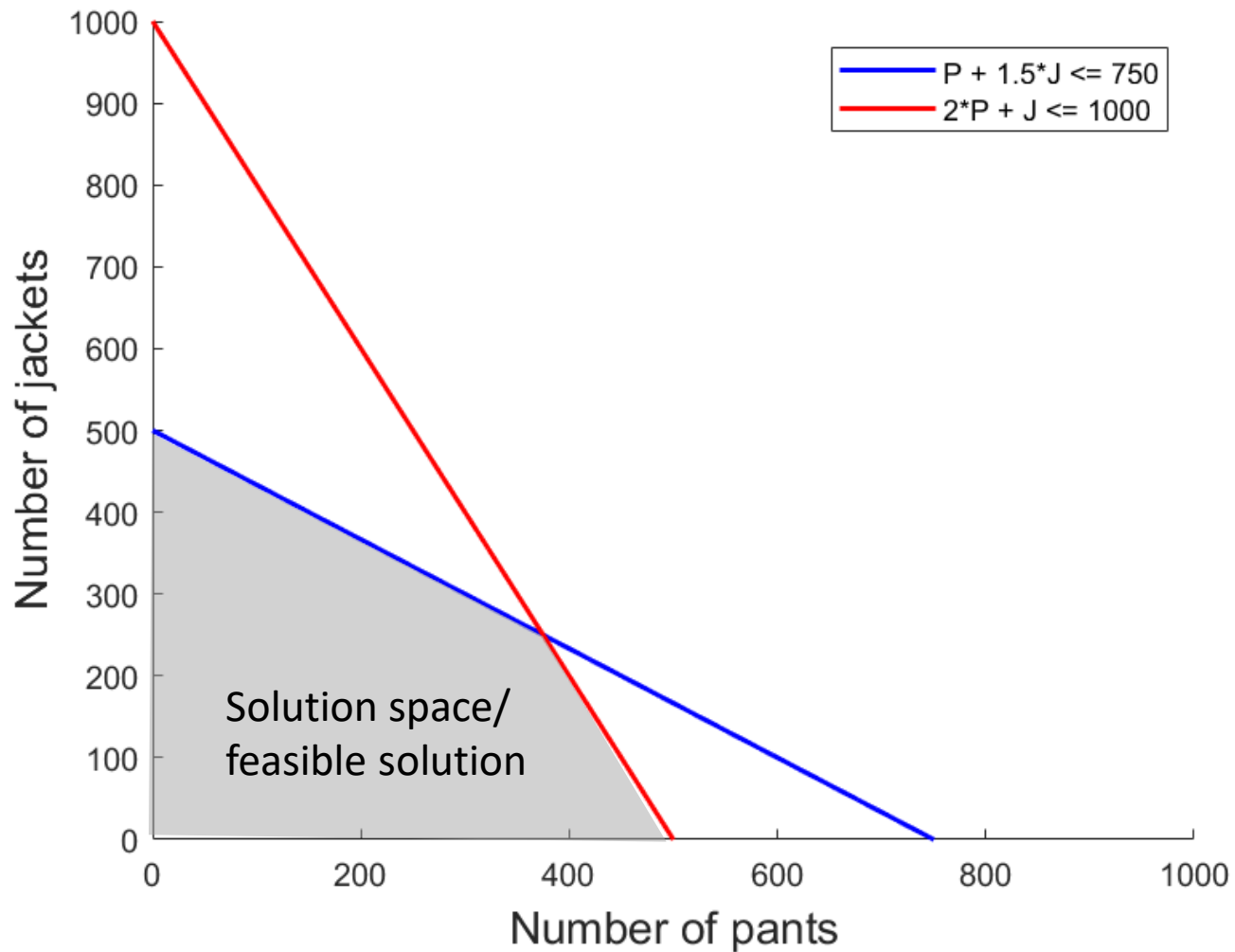
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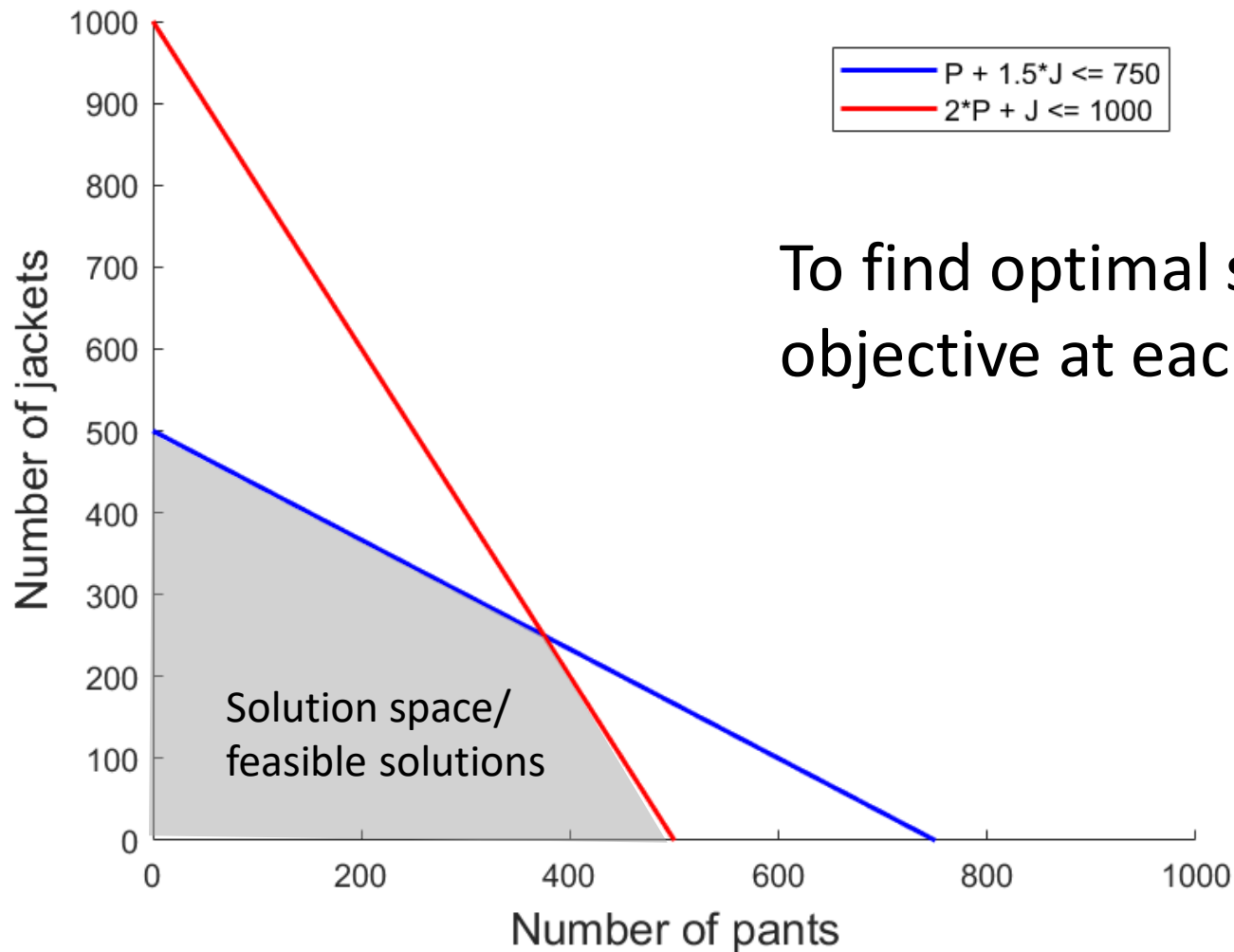
$$P \geq 0$$

$$J \geq 0$$

Graphical solution – solution space

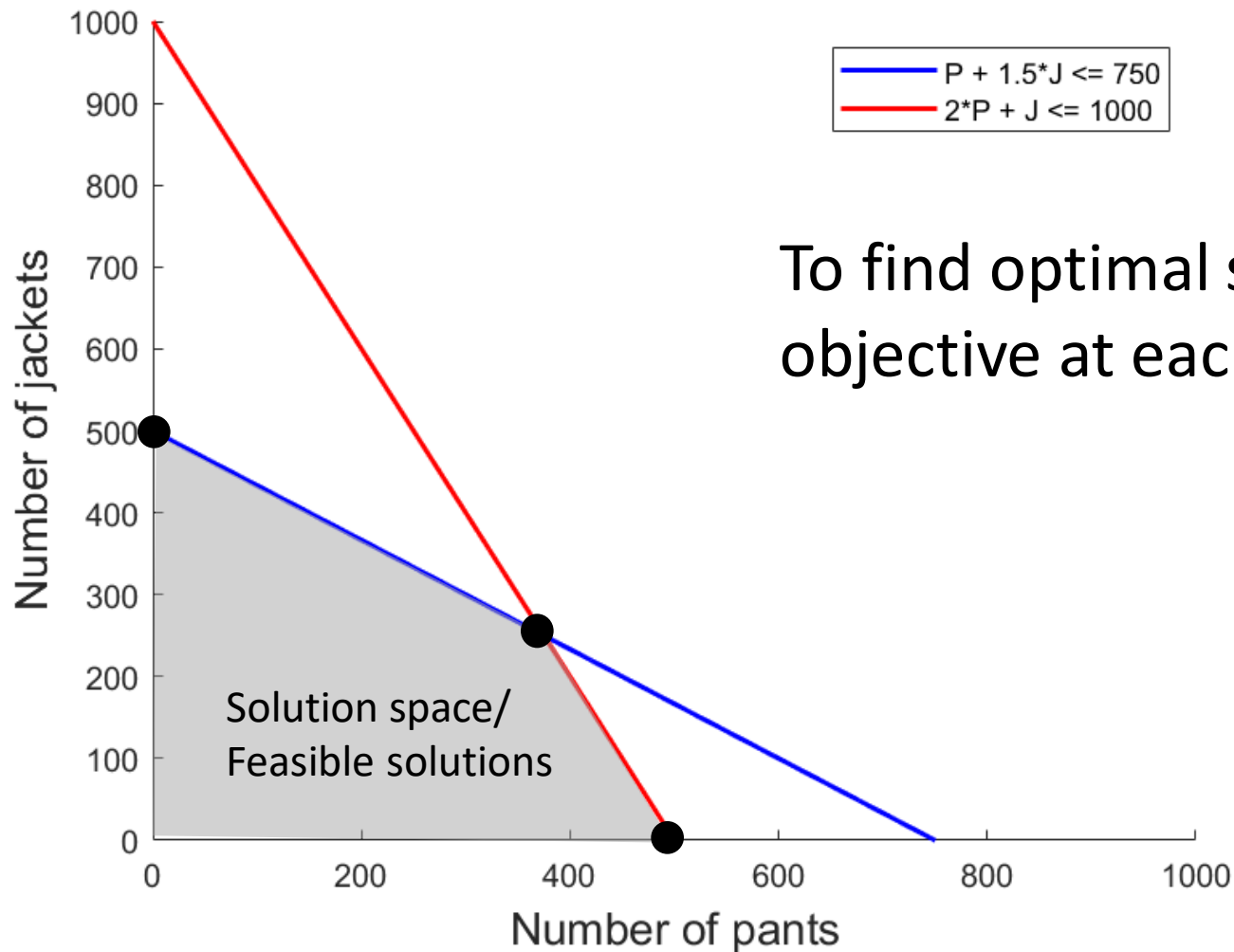


Graphical solution – optimal solution



To find optimal sale evaluate
objective at each **extreme point**

Graphical solution – optimal solution



To find optimal sale evaluate
objective at each **extreme point**

Graphical solution – optimal solution

$$z(P,J) = 50*P + 40*J \quad (\text{objective function})$$

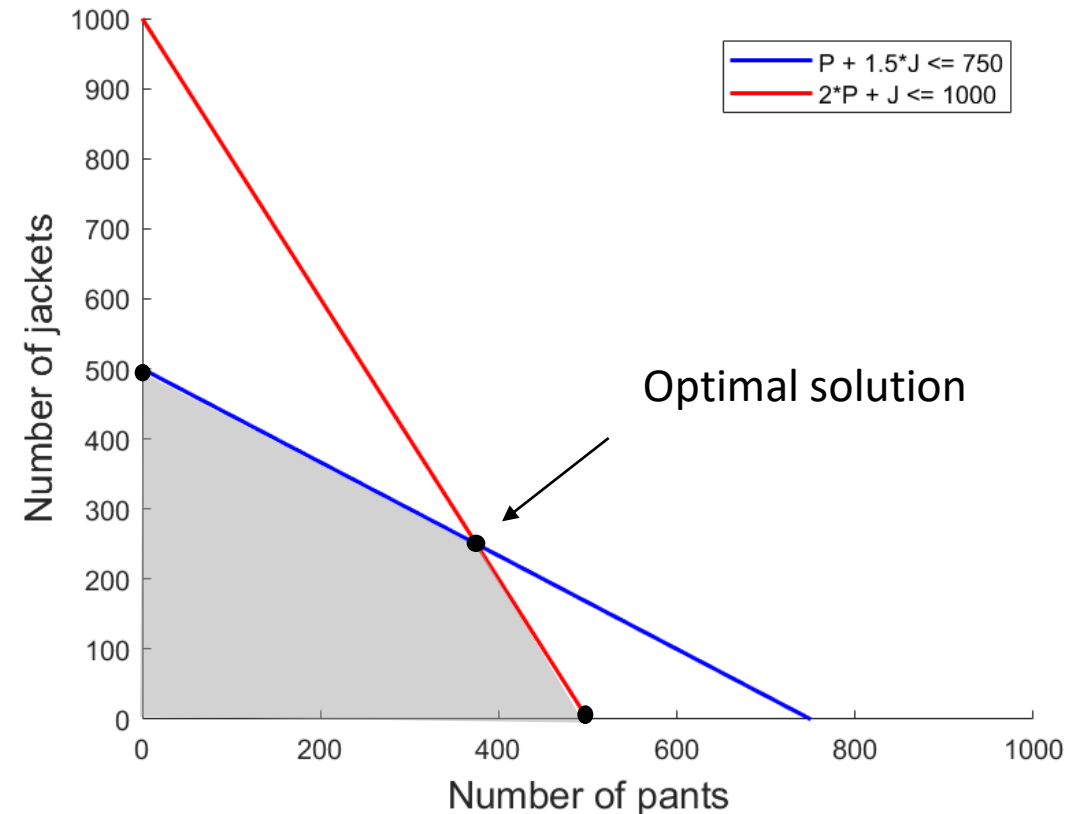
$$z(0,500) = 50*0 + 40*500 = 20\,000\text{€}$$

$$z(500,0) = 50*500 + 40*0 = 25\,000\text{€}$$

$$z(375,250) = 50*375 + 40*250 = \mathbf{28\,750\text{€}}$$

The optimal sale is achieved if 375 pants and 250 jackets are produced.

Q: Is the optimal solution unique?



Graphical solution – optimal solution

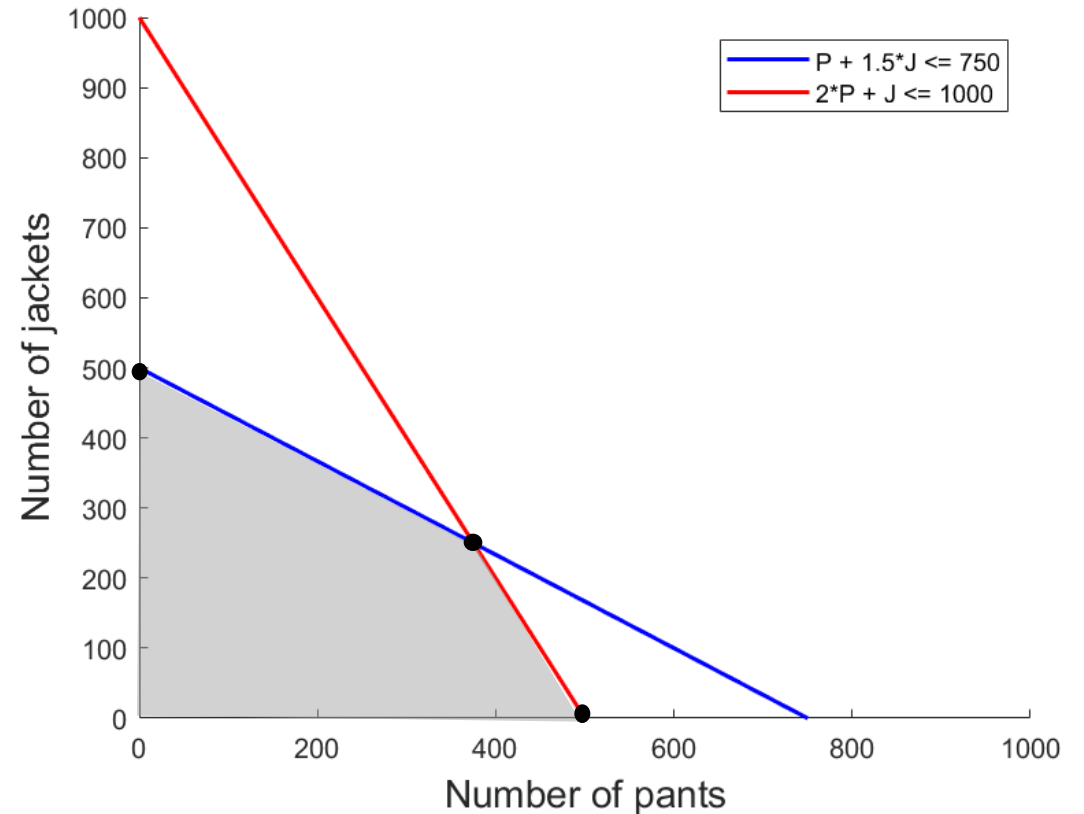
-> price change from 50€ and 40€ for pants and jackets to 20€ and 30€, respectively

$$z(P,J) = 20 \cdot P + 30 \cdot J$$

$$z(0,500) = 20 \cdot 0 + 30 \cdot 500 =$$

$$z(500,0) = 20 \cdot 500 + 30 \cdot 0 =$$

$$z(375,250) = 20 \cdot 375 + 30 \cdot 250 =$$



Graphical solution – constraint 2

-> price change from 50€ and 40€ for pants and jackets to 20€ and 30€, respectively

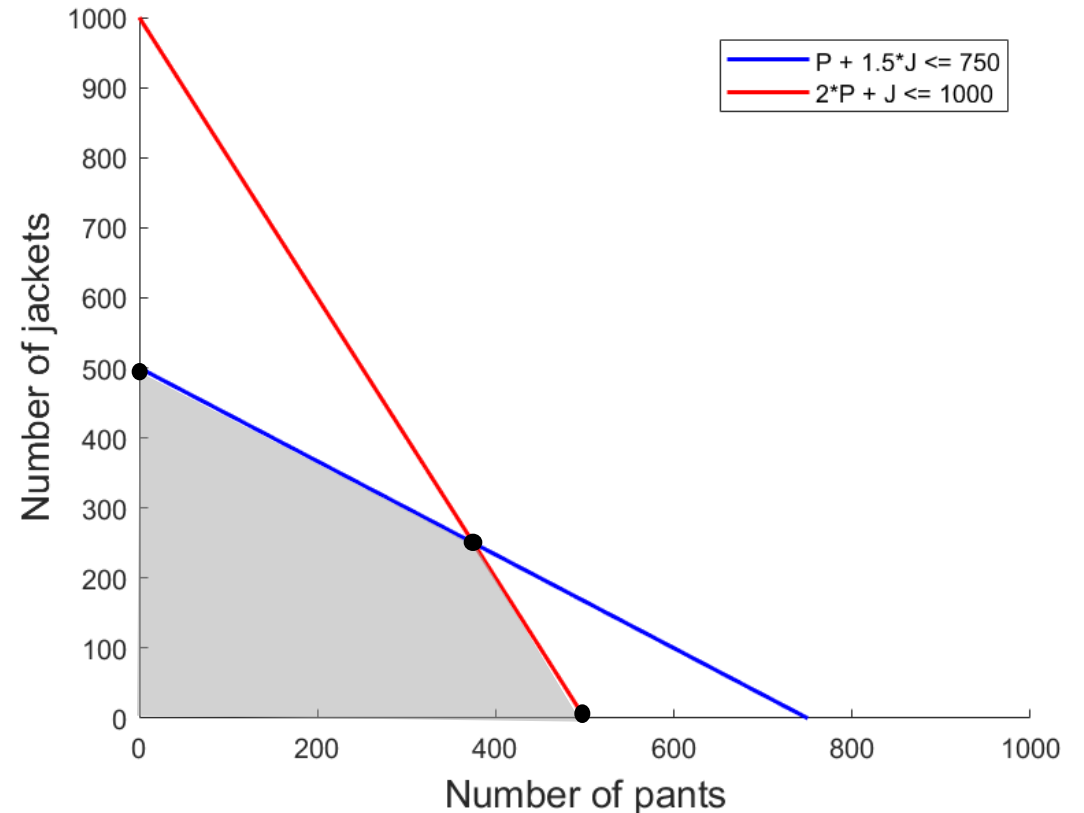
$$z(P,J) = 20 \cdot P + 30 \cdot J$$

$$z(0,500) = 20 \cdot 0 + 30 \cdot 500 = \mathbf{15\ 000\text{€}}$$

$$z(500,0) = 20 \cdot 500 + 30 \cdot 0 = 10\ 000\text{€}$$

$$z(375,250) = 20 \cdot 375 + 30 \cdot 250 = \mathbf{15\ 000\text{€}}$$

Q: What is the optimal solution?



Graphical solution – constraint 2

-> price change from 50€ and 40€ for pants and jackets to 20€ and 30€, respectively

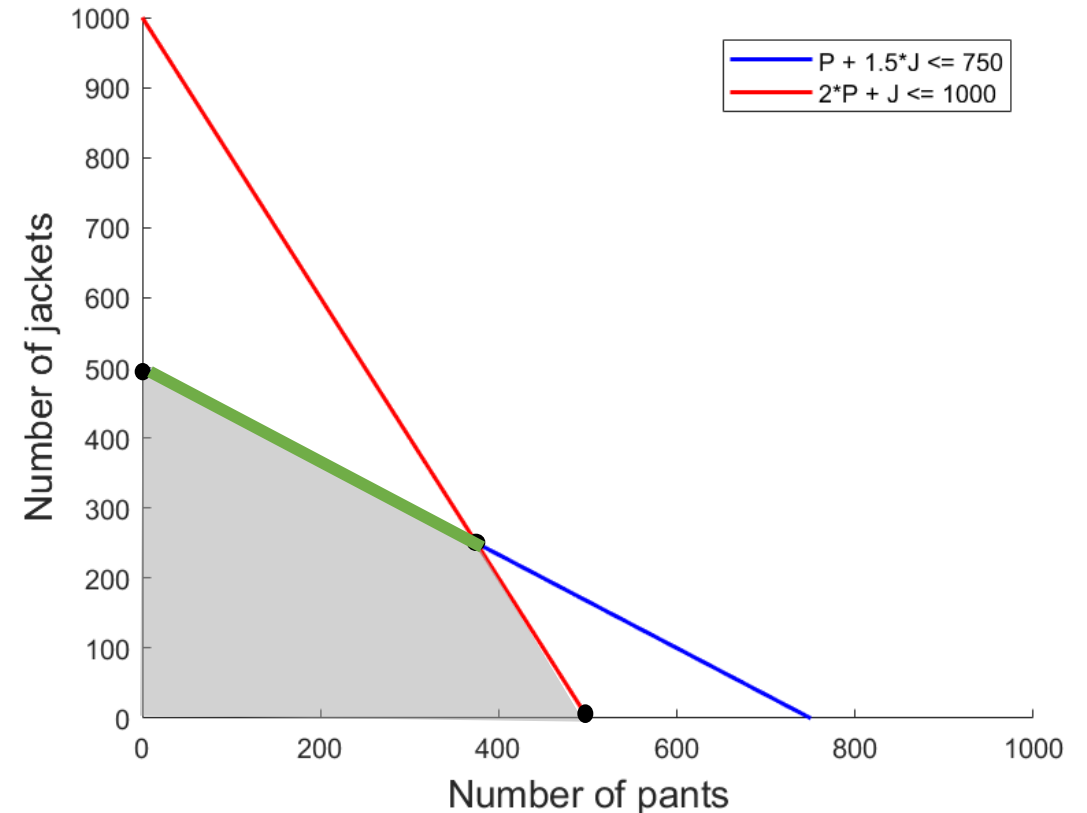
$$z(P,J) = 20 \cdot P + 30 \cdot J$$

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$$z(500,0) = 20 \cdot 500 + 30 \cdot 0 = 10\,000\text{€}$$

$$z(375,250) = 20 \cdot 375 + 30 \cdot 250 = 15\,000\text{€}$$

All pairs of integer values on the segment in green would give the same optimal sale!



Solving LP using Simplex

Reformulation into system of equalities!

$\max z = 50*P + 40*J$ (OBJECTIVE - MAXIMIZE PROFIT)

s.t.

$1*P + 1.5*J \leq 750$ (AVAILABILITY OF COTTON CONSTRAINT)

$2*P + 1*J \leq 1000$ (AVAILABILITY OF POLYESTER CONSTRAINT)

$P \geq 0$

$J \geq 0$

To use simplex we first convert all inequalities to equalities.

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To use simplex we first convert all inequalities to equalities.

Row 1: $z - 50*P - 40*J - 0*s - 0*t = 0$

Row 2: $1*P + 1.5*J + s = 750$ (less than --> + slack, greater than --> - slack)

Row 3: $2*P + 1*J + t = 1000$

$P, J, s, t \geq 0$

Initial tableau

Basic variable: variable isolated in a constraint that can be set to the right-hand side of that constraint

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	-50	-40	0	0	0	
2	s	0	1	1.5	1	0	750	
3	t	0	2	1	0	1	1000	

Basic columns

Optimality conditions:

1. The objective row of the tableau is 0 in the basic columns, except for z
2. There is no negative entry in the objective row ❌

Identify entering variable

Basic variable: variable isolated in a constraint that can be set to the right-hand side of that constraint

Nonbasic variable which is associated the most negative (for maximization) coefficient in the objective

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	-50	-40	0	0	0	
2	s	0	1	1.5	1	0	750	
3	t	0	2	1	0	1	1000	

Identify leaving variable

- Basic variable: variable isolated in a constraint that can be set to the right-hand side of that constraint
- **Entering variable**: nonbasic variable which is associated the most negative (for maximization) coefficient in the objective
- **Leaving variable**: variable which will be changed from a non-zero to zero value in the next solution
-> choose the one that has smallest upper bound on entering variable

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	-50	-40	0	0	0	
2	s	0	1	1.5	1	0	750	$750/1=750$
3	t	0	2	1	0	1	1000	$1000/2=500$

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- Basic variable: variable isolated in a constraint that can be set to the right-hand side of that constraint
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3	t	0	2	1	0	1	1000	$1000/2=500$

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2	s	0	1	1.5	1	0	750	$750/1=750$
3	t	0	2	1	0	1	1000	$1000/2=500$

Second tableau

- Enter the basic variable for the second tableau.
- Update the coefficients of the tableau

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	-50	-40	0	0	0	
2	s	0	1	1.5	1	0	750	750
3	t	0	2	1	0	1	1000	500

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1			0					
2			0					
3	P	0	1	1/2	0	1/2	500	

Becomes basic variable

Second tableau

- Enter the basic variable for the second tableau.
- Update the coefficients of the tableau

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
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2	s	0	1	1.5	1	0	750	750
3	t	0	2	1	0	1	1000	500

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0		0			
2	s	0	0		1			
3	P	0	1	1/2	0	1/2	500	

Becomes basic variable

Second tableau

- Enter the basic variable for the second tableau.
- Update the coefficients of the tableau $T2(i,j) = T1(i,j) - T1(i,entering\ variable) * T2(leaving\ variable,j))$

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0		0			
2	s	0	0		1			
3	P	0	1	1/2	0	1/2	500	

Initial tableau

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Second tableau

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Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	$-15 = -40 - (-50 * 1/2)$	0			
2	s	0	0		1			
3	P	0	1	1/2	0	1/2	500	

Initial tableau

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
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2	s	0	1	1.5	1	0	750	750
3	t	0	2	1	0	1	1000	500

Second tableau

- Enter the basic variable for the second tableau.
- Update the coefficients of the tableau $T2(i,j) = T1(i,j) - T1(i, \text{entering variable}) * T2(\text{leaving variable}, j)$

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	$-15 = -40 - (-50 * 1/2)$	0			
2	s	0	0	$1 = 1.5 - (1 * 1/2)$	1			
3	P	0	1	1/2	0	1/2	500	

Initial tableau

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	-50	-40	0	0	0	
2	s	0	1	1.5	1	0	750	750
3	t	0	2	1	0	1	1000	500

Second tableau

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- Update the coefficients of the tableau $T2(i,j) = T1(i,j) - T1(i,entering\ variable) * T2(leaving\ variable,j))$

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	$-15 = -40 - (-50*1/2)$	0	$25 = 0 - (-50*1/2)$		
2	s	0	0	$1 = 1.5 - (1*1/2)$	1			
3	P	0	1	1/2	0	1/2	500	

Initial tableau

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	-50	-40	0	0	0	
2	s	0	1	1.5	1	0	750	750
3	t	0	2	1	0	1	1000	500

Second tableau

- Enter the basic variable for the second tableau.
- Update the coefficients of the tableau $T2(i,j) = T1(i,j) - T1(i, \text{entering variable}) * T2(\text{leaving variable}, j)$

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	$-15 = -40 - (-50 * 1/2)$	0	$25 = 0 - (-50 * 1/2)$		
2	s	0	0	$1 = 1.5 - (1 * 1/2)$	1	$-1/2 = 0 - (1 * 1/2)$		
3	P	0	1	1/2	0	1/2	500	

Initial tableau

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	-50	-40	0	0	0	
2	s	0	1	1.5	1	0	750	750
3	t	0	2	1	0	1	1000	500

Second tableau

- Enter the basic variable for the second tableau.
- Update the coefficients of the tableau $T2(i,j) = T1(i,j) - T1(i, \text{entering variable}) * T2(\text{leaving variable}, j)$

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	$-15 = -40 - (-50 * 1/2)$	0	$25 = 0 - (-50 * 1/2)$	$25000 = 0 - (-50 * 500)$	
2	s	0	0	$1 = 1.5 - (1 * 1/2)$	1	$-1/2 = 0 - (1 * 1/2)$		
3	P	0	1	1/2	0	1/2	500	

Initial tableau

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	-50	-40	0	0	0	
2	s	0	1	1.5	1	0	750	750
3	t	0	2	1	0	1	1000	500

Second tableau

- Enter the basic variable for the second tableau.
- Update the coefficients of the tableau $T2(i,j) = T1(i,j) - T1(i,entering\ variable) * T2(leaving\ variable,j))$


Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	$-15 = -40 - (-50 * 1/2)$	0	$25 = 0 - (-50 * 1/2)$	$25000 = 0 - (-50 * 500)$	
2	s	0	0	$1 = 1.5 - (1 * 1/2)$	1	$-1/2 = 0 - (1 * 1/2)$	$250 = 750 - (1 * 500)$	
3	P	0	1	1/2	0	1/2	500	

Initial tableau

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	-50	-40	0	0	0	
2	s	0	1	1.5	1	0	750	750
3	t	0	2	1	0	1	1000	500

Second tableau

Optimality conditions:

1. The objective row of the tableau is 0 in the basic columns
2. There is no negative entry in the objective row 

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	-15	0	25	25000	
2	s	0	0	1	1	-1/2	250	
3	P	0	1	1/2	0	1/2	500	

Identify entering and leaving variable

- Basic variable: variable isolated in a constraint that can be set to the right-hand side of that constraint
- **Entering variable**: nonbasic variable which is associated the most negative (for maximization) coefficient in the objective
- **Leaving variable**: variable which will be changed from a non-zero to zero value in the next solution
-> choose the one that has smallest upper bound on entering variable

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	-15	0	25	25000	
2	s	0	0	1	1	-1/2	250	$250/1 = 250$
3	P	0	1	1/2	0	1/2	500	$500/0.5 = 1000$

Third tableau

- Enter the basic variable for the new tableau.

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1				0				
2	J	0		1	1	-1/2	250	
3				0				

previous tableau	Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
	1	z	1	0	-15	0	25	25000	
	2	s	0	0	1	1	-1/2	250	250
	3	P	0	1	1/2	0	1/2	500	1000

Third tableau

- Enter the basic variable for the new tableau.

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	0				
2	J	0	0	1	1	-1/2	250	
3	P	0	1	0				

previous tableau	Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
	1	z	1	0	-15	0	25	25000	
	2	s	0	0	1	1	-1/2	250	250
	3	P	0	1	1/2	0	1/2	500	1000

Third tableau

- Update the coefficients of the tableau ($V(i,j) = O(i,j) - O(i,\text{entering variable}) * V(\text{leaving variable},j)$)

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	0	$15 = 0 - (-15*1)$			
2	J	0	0	1	1	-1/2	250	
3	P	0	1	0				

previous tableau {

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	-15	0	25	25000	
2	s	0	0	1	1	-1/2	250	250
3	P	0	1	1/2	0	1/2	500	1000

Third tableau

- Update the coefficients of the tableau ($V(i,j) = O(i,j) - O(i,\text{entering variable}) * V(\text{leaving variable},j)$)

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	0	$15 = 0 - (-15*1)$			
2	J	0	0	1	1	-1/2	250	
3	P	0	1	0	$-1/2 = 0 - (1/2*1)$			

previous tableau

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	-15	0	25	25000	
2	s	0	0	1	1	-1/2	250	250
3	P	0	1	1/2	0	1/2	500	1000

Third tableau

- Update the coefficients of the tableau ($V(i,j) = O(i,j) - O(i,\text{entering variable}) * V(\text{leaving variable},j)$)

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	0	$15 = 0 - (-15*1)$	$17.5 = 25 - (-15*-1/2)$		
2	J	0	0	1	1	-1/2	250	
3	P	0	1	0	$-1/2 = 0 - (1/2*1)$			

previous tableau

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	-15	0	25	25000	
2	s	0	0	1	1	-1/2	250	250
3	P	0	1	1/2	0	1/2	500	1000

Third tableau

- Update the coefficients of the tableau ($V(i,j) = O(i,j) - O(i,\text{entering variable}) * V(\text{leaving variable},j)$)

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	0	$15 = 0 - (-15*1)$	$17.5 = 25 - (-15*-1/2)$		
2	J	0	0	1	1	-1/2	250	
3	P	0	1	0	$-1/2 = 0 - (1/2*1)$	$1/4 = 0 - (1/2*-1/2)$		

previous tableau {

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	-15	0	25	25000	
2	s	0	0	1	1	-1/2	250	250
3	P	0	1	1/2	0	1/2	500	1000

Third tableau

- Update the coefficients of the tableau ($V(i,j) = O(i,j) - O(i,\text{entering variable}) * V(\text{leaving variable},j)$)

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	0	$15 = 0 - (-15*1)$	$17.5 = 25 - (-15*-1/2)$	$28750 = 25000 - (-15*250)$	
2	J	0	0	1	1	-1/2	250	
3	P	0	1	0	$-1/2 = 0 - (1/2*1)$	$1/4 = 0 - (1/2*-1/2)$		

previous tableau {

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	-15	0	25	25000	
2	s	0	0	1	1	-1/2	250	250
3	P	0	1	1/2	0	1/2	500	1000

Third tableau

- Update the coefficients of the tableau ($V(i,j) = O(i,j) - O(i,\text{entering variable}) * V(\text{leaving variable},j)$)

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	0	$15 = 0 - (-15*1)$	$17.5 = 25 - (-15*-1/2)$	$28750 = 25000 - (-15*250)$	
2	J	0	0	1	1	-1/2	250	
3	P	0	1	0	$-1/2 = 0 - (1/2*1)$	$1/4 = 0 - (1/2*-1/2)$	$375 = 500 - (1/2*250)$	

previous tableau {

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	-15	0	25	25000	
2	s	0	0	1	1	-1/2	250	250
3	P	0	1	1/2	0	1/2	500	1000

Optimality conditions

- ✓ The objective row of the tableau is 0 in the basic columns
- ✓ There is no negative entry in the objective row

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	0	15	17.5	28750	
2	J	0	0	1	1	-1/2	250	
3	P	0	1	0	-1/2	1/4	375	

Solution: $P = ?$, $J = ?$, $z = ?$

Optimality conditions

- ✓ The objective row of the tableau is 0 in the basic columns
- ✓ There is no negative entry in the objective row

Row number	Basic variable	z	P	J	s	t	Right-hand side	Upper bound on entering variable
1	z	1	0	0	15	17.5	28750	
2	J	0	0	1	1	-1/2	250	
3	P	0	1	0	-1/2	1/4	375	

Solution: $P = 375$, $J = 250$, $z = 28750$

Writing the LP using matrix format

$$\max z = c^T x$$

s. t.

$$Ax = b$$

$$lb \leq x \leq ub$$

$$\max z = 50 P + 40 J$$

s.t.

$$1 \cdot P + 1.5 \cdot J \leq 750 \quad (\text{cotton})$$

$$2 \cdot P + 1 \cdot J \leq 1000 \quad (\text{polyester})$$

$$P \geq 0$$

$$J \geq 0$$

Writing the LP using matrix format

$$\max z = c^T x \quad \max z = [50 \quad 40] \begin{bmatrix} P \\ J \end{bmatrix}$$

s. t.

$$Ax = b$$

$$lb \leq x \leq ub$$

$$\max z = 50 P + 40 J$$

s.t.

$$1 \cdot P + 1.5 \cdot J \leq 750 \quad (\text{cotton})$$

$$2 \cdot P + 1 \cdot J \leq 1000 \quad (\text{polyester})$$

$$P \geq 0$$

$$J \geq 0$$

Writing the LP using matrix format

$$\max z = c^T x \qquad \max z = [50 \quad 40] \begin{bmatrix} P \\ J \end{bmatrix}$$

s. t. *s. t.*

$$Ax = b \qquad \begin{bmatrix} \quad \quad \end{bmatrix} \begin{bmatrix} P \\ J \end{bmatrix} = \begin{bmatrix} 750 \\ 1000 \end{bmatrix}$$

$$lb \leq x \leq ub$$

$$\max z = 50 P + 40 J$$

s.t.

$$1 \cdot P + 1.5 \cdot J \leq 750 \qquad (\text{cotton})$$

$$2 \cdot P + 1 \cdot J \leq 1000 \qquad (\text{polyester})$$

$$P \geq 0$$

$$J \geq 0$$

Writing the LP using matrix format

$$\max z = c^T x \qquad \max z = [50 \quad 40] \begin{bmatrix} P \\ J \end{bmatrix}$$

s. t.

$$Ax = b \qquad \begin{bmatrix} 1 & 1.5 \\ 2 & 1 \end{bmatrix} \begin{bmatrix} P \\ J \end{bmatrix} \leq \begin{bmatrix} 750 \\ 1000 \end{bmatrix}$$

$$lb \leq x \leq ub \qquad \begin{bmatrix} 0 \\ 0 \end{bmatrix} \leq \begin{bmatrix} P \\ J \end{bmatrix} \leq \begin{bmatrix} Inf \\ Inf \end{bmatrix}$$

$$\max z = 50 P + 40 J$$

s.t.

$$1 \cdot P + 1.5 \cdot J \leq 750 \quad (\text{cotton})$$

$$2 \cdot P + 1 \cdot J \leq 1000 \quad (\text{polyester})$$

$$P \geq 0$$

$$J \geq 0$$

Writing the LP using matrix format

$$\max z = c^T x \qquad \max z = [50 \quad 40] \begin{bmatrix} P \\ J \end{bmatrix}$$

s. t.

$$Ax = b \qquad \begin{bmatrix} 1 & 1.5 \\ 2 & 1 \end{bmatrix} \begin{bmatrix} P \\ J \end{bmatrix} \leq \begin{bmatrix} 750 \\ 1000 \end{bmatrix}$$

$$x \geq 0 \qquad \begin{bmatrix} 0 \\ 0 \end{bmatrix} \leq \begin{bmatrix} P \\ J \end{bmatrix} \leq \begin{bmatrix} Inf \\ Inf \end{bmatrix}$$

$$\max z = 50 P + 40 J$$

s.t.

$$1 \cdot P + 1.5 \cdot J \leq 750 \quad (\text{cotton})$$

$$2 \cdot P + 1 \cdot J \leq 1000 \quad (\text{polyester})$$

$$P \geq 0$$

$$J \geq 0$$

The *linprog* solver

- part of the MATLAB Optimization Toolbox
- function definition: [\[x,fval,exitflag\] = linprog\(f,A,b,Aeq,beq,lb,ub\)](#)

Input arguments	argument	definition	Size
	f	defines the objective (see c from our example)	n x 1
	A	matrix encoding inequality constraints (\leq)	m x n
	b	Right-hand side vector for inequality constraints	m x 1
	Aeq	matrix encoding equality constraints (=)	m' x n
	beq	Right-hand side vector for equality constraints	b' x 1
	lb	lower limits (bounds) for all variables	n x 1
	ub	upper limits (bounds) for all variables	n x 1
Output	argument	definition	Size
	x	Solution vector	n x 1
	fval	Objective value at solution, fval = f'*x	1 x 1
	exitflag	Integer indicating reason of termination	1 x 1

$$\begin{aligned}
 &\min_x f^T \cdot x \\
 &\quad s.t. \\
 &\quad A \cdot x \leq b \\
 &\quad Aeq \cdot x = beq \\
 &\quad lb \leq x \leq ub
 \end{aligned}$$

Encoding inequality constraints in *linprog*

- all inequalities must be reformulated to the form: $A x \leq b$

examples:

- $5x + 3y \geq 60$

Encoding inequality constraints in *linprog*

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$$\implies -5x - 3y \leq -60$$

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examples:

- $5x + 3y \geq 60$

$$\implies -5x - 3y \leq -60$$

- $x \leq 100y$

Encoding inequality constraints in *linprog*

- all inequalities must be reformulated to the form: $A x \leq b$

examples:

- $5x + 3y \geq 60$

$$\implies -5x - 3y \leq -60$$

- $x \leq 100y$

$$\implies x - 100y \leq 0$$

Encoding inequality constraints in *linprog*

- all inequalities must be reformulated to the form: $A x \leq b$

examples:

- $5x + 3y \geq 60$

$$\implies -5x - 3y \leq -60$$

- $x \leq 100y$

$$\implies x - 100y \leq 0$$

- $x \geq -100 (y-1)$

Encoding inequality constraints in *linprog*

- all inequalities must be reformulated to the form: $A x \leq b$

examples:

- $5x + 3y \geq 60$

$$\implies -5x - 3y \leq -60$$

- $x \leq 100y$

$$\implies x - 100y \leq 0$$

- $x \geq -100(y-1)$

$$\implies -100y - x \leq -100$$

linprog arguments for our problem

function definition: $x = \text{linprog}(f, A, b, Aeq, beq, lb, ub)$

- $f =$
- $A =$
- $b =$
- $Aeq =$
- $beq =$
- $lb =$
- $ub =$

$$\begin{aligned} \min_x & f^T \cdot x \\ \text{s.t.} & \\ & A \cdot x \leq b \\ & Aeq \cdot x = beq \\ & lb \leq x \leq ub \end{aligned}$$

linprog arguments for our problem

function definition: $x = \text{linprog}(f, A, b, Aeq, beq, lb, ub)$

- $f = [-50; -40];$ $\rightarrow f \cdot (-1)$ needed since *linprog* *minimizes* by default, but we want to *maximize*
- $A = \begin{bmatrix} 1 & 1.5; \\ 2 & 1 \end{bmatrix};$
- $b = [750; 1000];$
- $Aeq = [];$
- $beq = [];$
- $lb = [0; 0];$
- $ub = [1000; 1000]$ (arbitrary large number which does not further constrain the solution)

$$\begin{aligned} \min_x & f^T \cdot x \\ \text{s.t.} & \\ & A \cdot x \leq b \\ & Aeq \cdot x = beq \\ & lb \leq x \leq ub \end{aligned}$$