

University of Sheffield

# Lab Report For COM6015- 2022



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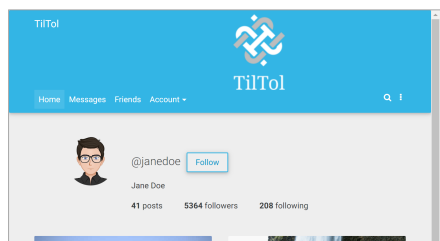
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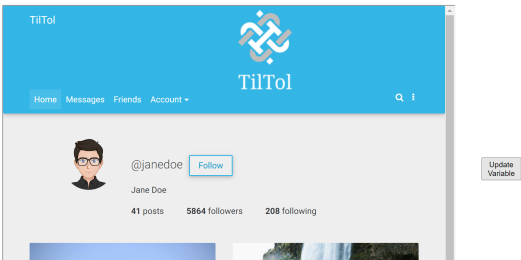
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# Chapter 1

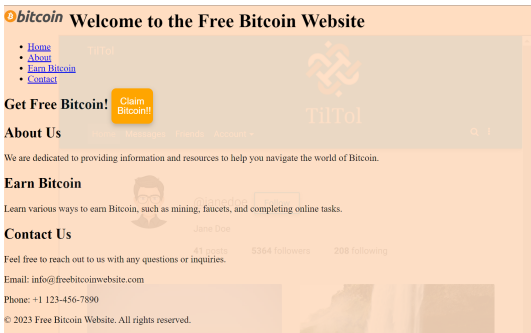
## Question 1:Clickjacking



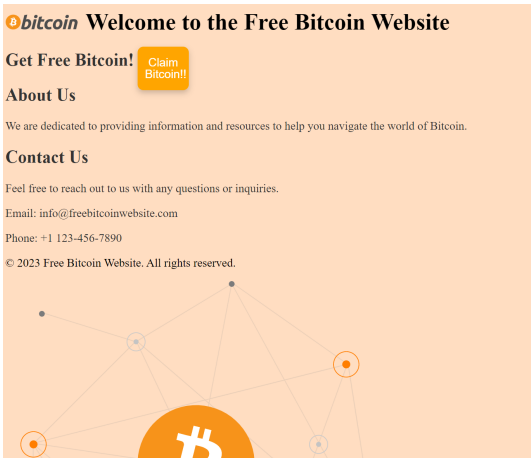
(a) Before Update: followers:5364



(b) After update: followers:5864 (+500 on every press)



(c) Pre-Final



(d) Final Clickbait

Figure 1.1: Scam Chain

The code files are attached as test.html for clickbait webpage, and joan.html for Joan Doe's tiltolpage.

```

1 <!DOCTYPE html>
2 /**
3  * Retrieves and updates followers count of Jane Doe
4  *
5  * The access() function retrieves a reference to the iframe element
6    containing Jane Doe's Webpage.
7  *
8  * The updateVariable() function accesses the document object of the iframe
9    and then retrieves the element with the ID "followers_no", which is the
10   number of followers and
11   * it updates its value by incrementing the current value by 500.
12   * Clicking the button triggers the updateVariable() function.
13  */
14 <script>
15     function access() {
16         var iframe = document.getElementById("victim_website");
17
18     }
19     function updateVariable() {
20         var iframe = document.getElementById("jane");
21         var iframeDocument = iframe.contentDocument || iframe.contentWindow.
22         document;
23
24         // Access and update the variable within the iframe
25         var followers_no = iframeDocument.getElementById("followers_no");
26         if (followers_no) {
27             var currentValue = parseInt(followers_no.textContent);
28             var newValue = currentValue + 500;
29             followers_no.textContent = newValue;
30         }
31     }
32 </script>
33 <iframe id="jane" src="joan.html" onload="access()" class="scam_frame"></
34 iframe>
35 <button onclick="updateVariable()" class="scam">Claim Bitcoin!!</button>

```

Code 1.1: Clickbait Code

The Clickbait puts the Joan Doe's page in an Iframe and then using JavaScript extracts the variable "followers.no" from the joan.html. Everytime the Claim Bitcoin button is clicked the number of followers increase by 500.

# Chapter 2

## Question 2: Wireshark

### 2.1 What is the victim MAC address?

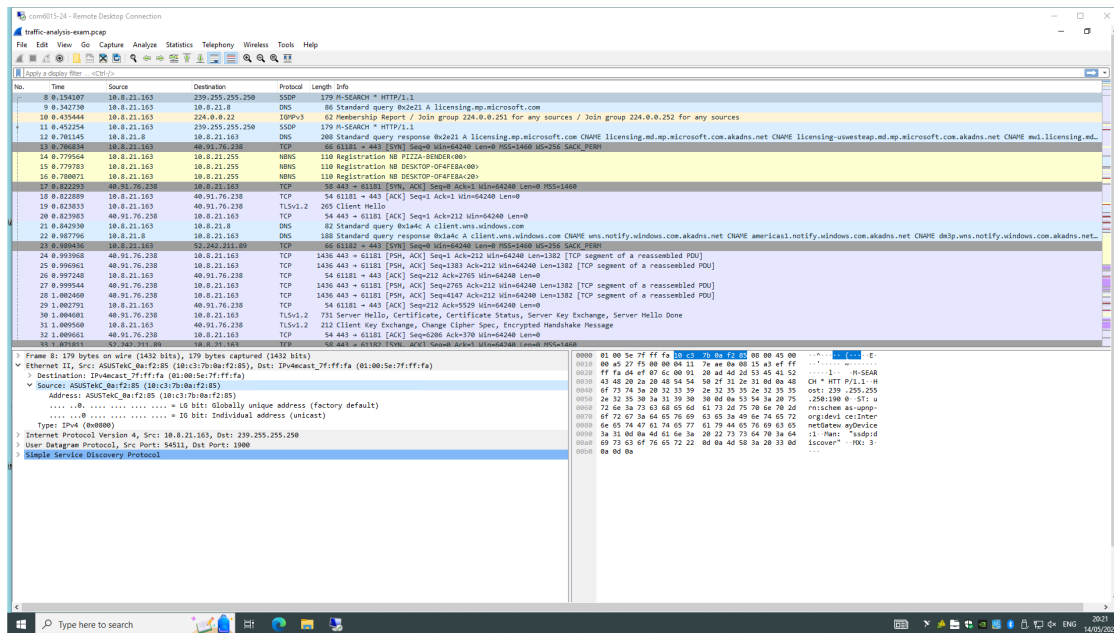


Figure 2.1: Victim MAC Address

As can be seen here the MAC address of the victim with the IP address - 10.8.21.163 is 10:C3:7B:0A:F2:85. This was found under the Source field of Ethernet.



## 2.2 What is the victim Host Name?

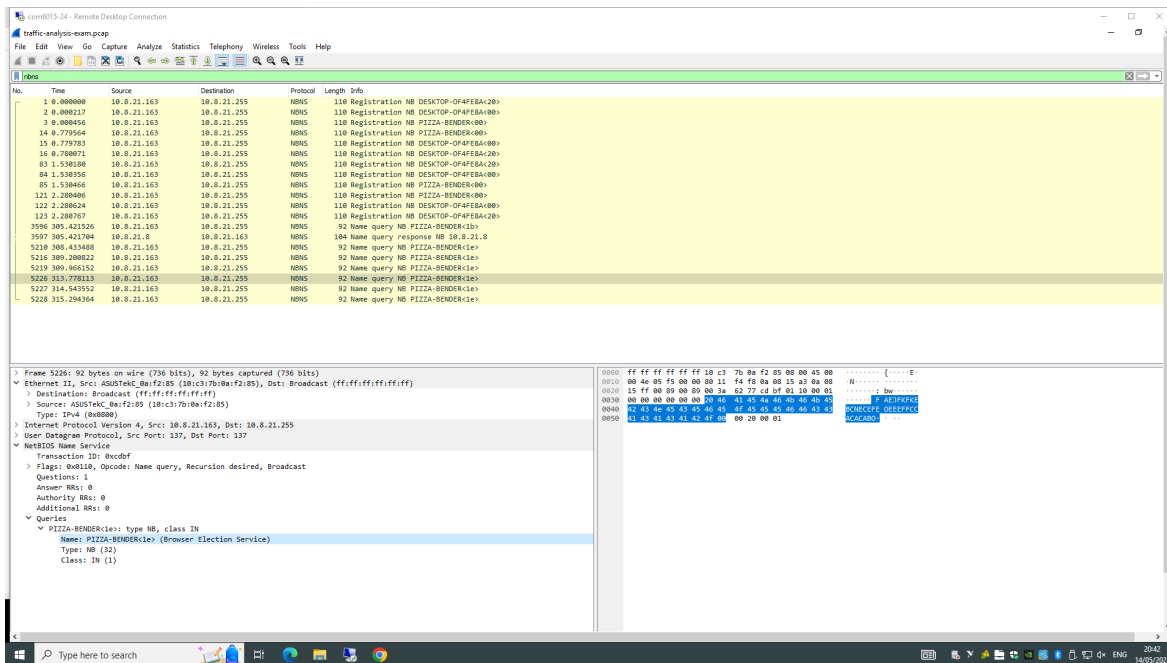
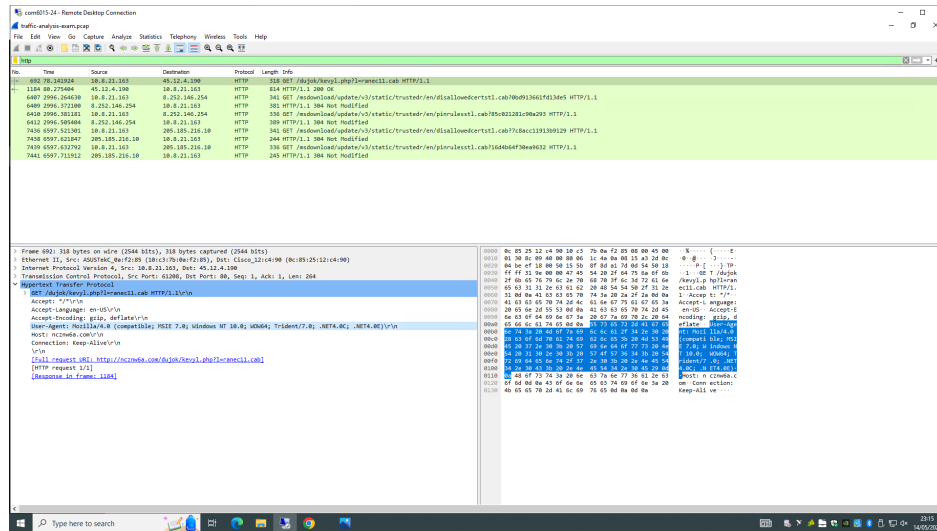


Figure 2.2: Victim Host Name

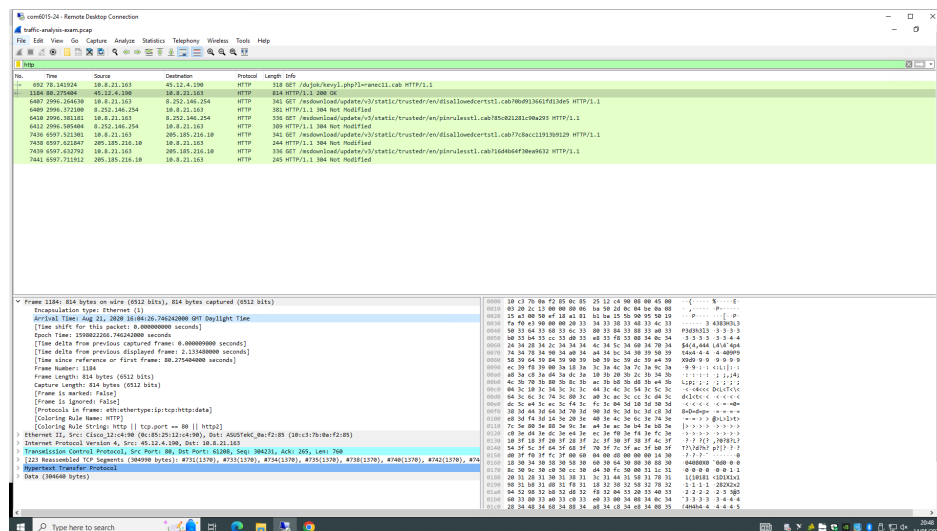
The host name is PIZZA-BENDER. This was found by checking the NetBIOS which is used by Microsoft Windows for its name resolution function.

### 2.3 What is the client operating system, and what day and time the malware was executed?



**Figure 2.3:** *Victim Client OS*

Here HTTP filter was used and as can be seen that there is only one request that gave code 200(success). So the malware might have been downloaded from here. Now as can be seen the HTTP field in the User-Agent that the client operating system is Windows NT 10.0.



**Figure 2.4:** *Date and Time of attack*

Also the day and time on which the attack was executed, which can be seen in the particulars of the frame, is August 21, 2020 at 16:04.

## 2.4 Where was the intruder file location?

As can be seen in the full request URI field the HTTP request is as follows :  
<http://ncznw6a.com/dujok/kevy1.php?1=ranec11.cab> .

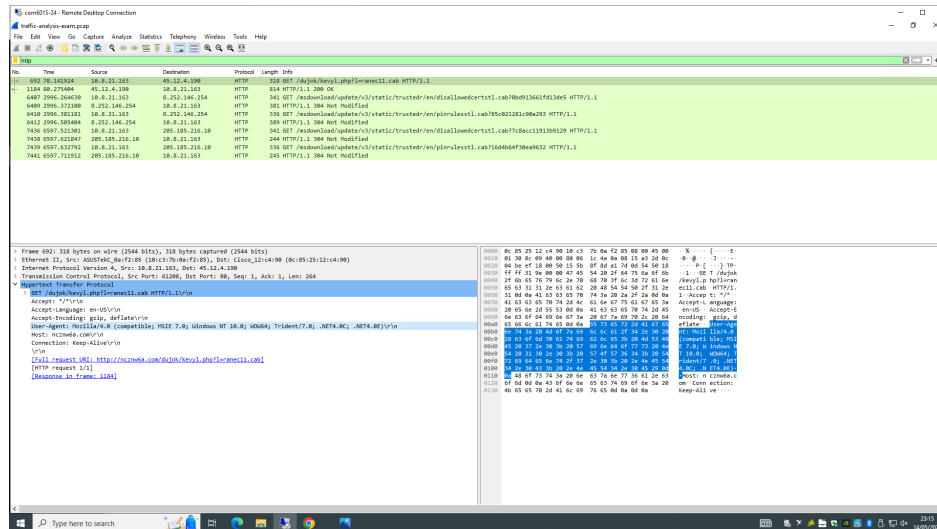


Figure 2.5: File Location

## 2.5 What is the intruder IP address? Show the geographical location of the intruder using the Wireshark endpoint feature.

As we can see in the above images, the IP address is 45.12.4.190. Now using the Wireshark endpoint feature we got the geographical location as:

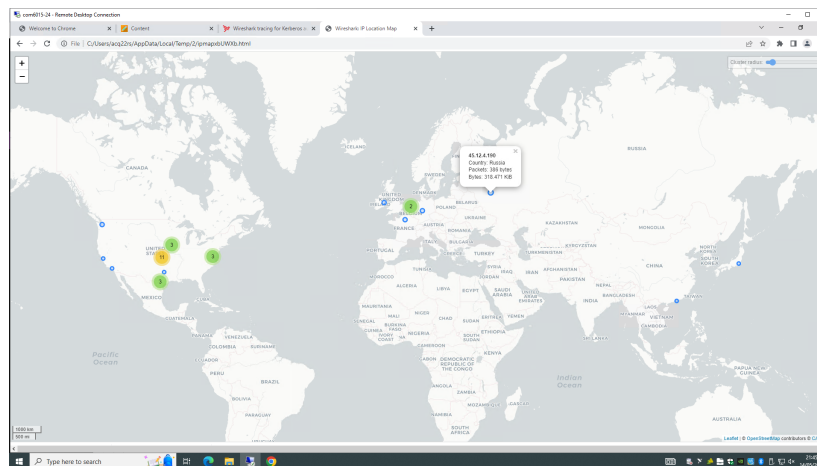


Figure 2.6: IP Address

## 2.6 Extract the malicious file (in your VM) and extract its hash key using Linux terminal/ windows cmd (SHA 256).

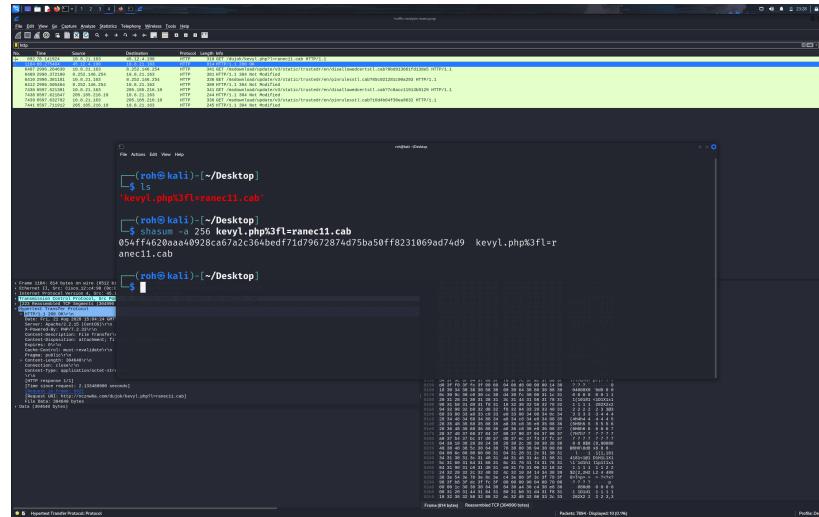


Figure 2.7: *Malicious File*

The pcap was opened in Kali to see the hash as the as the vm system did not have rights to download malicious files. The file was exported from Wireshark and the hash(SHA 256) as seen was extracted.

## 2.7 Use the extracted SHA 256 code, search on the Internet and provide the malware description.

We can see that when the hash was entered on the internet the following details were found. The malware's name is IcedID and it is a trojan that is designed to steal banking credentials of the victims. It steals payment information and can also deliver another viruses. This explains the unauthorised withdrawal of \$5000.

**ANY.RUN** ANALYZE MALWARE

- Huge database of samples and IOCs
- Custom VM setup
- Unlimited submissions
- Interactive approach
- Sign up, it's free

### General Info

File name: ranec11.cab

Full analysis: <https://app.any.run/tasks/3c4a9f63-7c35-496e-a14c-5704fd9cda88>

Verdict: **Suspicious activity**

Threats: **IcedID**

IcedID is a banking trojan-type malware which allows attackers to utilize it to steal banking credentials of the victims. IcedID aka BokBot mainly targets businesses and steals payment information, it also acts as a loader and can deliver another viruses or download additional modules.

Analysis date: August 21, 2020 at 22:25:16

OS: Windows 10 Professional (build: 16299, 64 bit)

Tags: icedid shathak ta551

Indicators:

MIME: application/x-dosexec

File info: PE32 executable (DLL) (GUI) Intel 80386, for MS Windows

MD5: A52A1E151BF4B993EFCFF87B3780D731

SHA1: FOACF0B669505BED5BD7FCB7F32C88C08E7DE768

SHA256: 054FF4620AAA40928CA67A2C364BEDF71D79672874D758A50FF8231069AD74D9

SSDEEP: 6144:ULdSq++cF0Nsj8AbctH97Qv680yDi6S6Zu+GADv9JfJLx+qNs1dbctQveyDou+vGxAF

ANY.RUN is an interactive service which provides full access to the guest system. Information in this report could be distorted by user actions and is provided for user acknowledgement as it is. ANY.RUN does not guarantee maliciousness or safety of the content.

Software environment set and analysis options

### Behavior activities

Figure 2.8: Malicious File

## 2.8 Which user is accountable for downloading the malicious malware from the Internet? provide Account Name by searching inside Kerberos packets (hint: find the answer in the KRBS packet (237) ‘CNameString’ section)

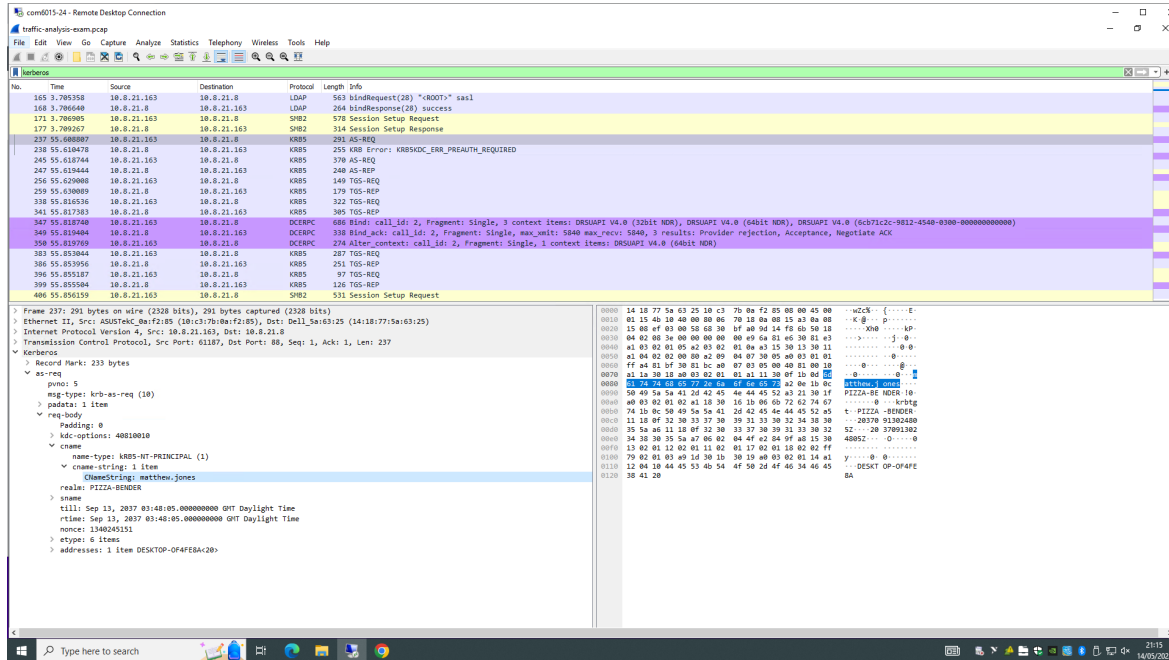


Figure 2.9: User Accountable

As can be seen the kerberos packets were explored. In the cname string the name Matthew.jones is seen. It can be inferred that he is responsible for downloading the malicious malware.

## Chapter 3

### Question 3 : SQL Injection

- 3.1 Extract the column names of the table that contains the user data
- 3.2 Attempt the previous injection in medium security, update your query to bypass the new security measures
- 3.3 Locate the location of the database on the remote system (File path the database is stored)
- 3.4 Read a file (e.g. passwords) from the discovered path
- 3.5 Discover the users with the highest and lowest salaries
- 3.6 Who has the insurance number: 53779132

## Chapter 4

### Question 4 : XSS

- 4.1 Set your DVWA to HIGH security. Analyse the PHP source code and perform an attack that works on Reflected, Stored and DOM XSS pages. Provide a description for each relating to the source code as to why your given attack works.



## Chapter 5

### Question 5 : Defence

5.1 In a language of your choice, write a short function which handles input from a user which would prevent an SQL injection attack and a Stored XSS attack. Provide a short description of where/why your function is safe against these vulnerabilities.

```
1  function sanitizeInput(input) {  
2  // to remove single quotes, double quotes, and backslashes from the input  
   string.  
3  const sanitizedInput = input.replace(/["'\\]/g, '');  
4  
5  // prevents the browser from interpreting the input as HTML and executing  
   potentially harmful scripts.  
6  const safeInput = document.createElement('div');  
7  safeInput.appendChild(document.createTextNode(sanitizedInput));  
8  const escapedInput = safeInput.innerHTML;  
9  
10 return escapedInput;  
11 }
```

Code 5.1: Code Sanitizer

## Chapter 6

# Question 6: Buffer Overflow

### 6.1 Choose one of the known buffer overflow vulnerabilities and write a half-page description of it. Aspects that you should cover are:

The FORCEDENTRY buffer overflow vulnerability, also known as CVE-2021-30860, was an integer overflow vulnerability when collating referenced segments<sup>[1]</sup> resulting in triggering a heap buffer overflow in the ImageIO JBIG2 decoder, used by ImageIO library in CoreGraphicsAPI used in Apple. The entry point for attack was a fake gif containing JBIG2(decoder codec) for PDF which was sent as iMessage, the JBIG2 codec contained code which was Turing Complete i.e. operators like AND, OR, XOR, and XNOR, enabled the attacker to perform bit-level operations on memory regions at arbitrary offsets. Tallowing for arbitrary memory access and computation of any computable function.

#### 6.1.1 Which systems were affected by the vulnerability?

The vulnerability affected all versions of iOS and macOS prior to iOS 14.8 and macOS Big Sur 11.3.1. It is estimated that over 1 billion devices were vulnerable to the attack at the time it was discovered.

#### 6.1.2 When was it discovered, reported and fixed?

This vulnerability was discovered in March 2021 by Citizen Lab, The vulnerability was patched by Apple in iOS 14.8 and macOS Big Sur 11.3.1 updates released for its devices.

#### 6.1.3 What were the known attacks exploiting it, and what were their consequences?

This vulnerability was used by NSO Group for its zero-click iMessage exploit to spy on Saudi activist and Al-jazeera journalists<sup>[2]</sup>.The affected phone was analysed by Citizen Lab and further analysis was done by Ian Beer of Project Zero<sup>[1]</sup> from Google.NSO was added the "Entity List" by the US government, severely restricting the ability of US companies to do business with NSO.

## 6.2 Propose a modification of the C programming language that would mitigate buffer overflow vulnerabilities, and discuss the implications of this modification.

Buffer Overflow vulnerabilities have plagued C language since the 1980s, as C allows for direct access to memory and lacks strong object typing.

One Modification that would help in mitigating buffer overflows would be the use of canaries or stack cookies<sup>[3]</sup> during runtime, They are random values which are placed on stack after each buffer to check for buffer overflows. They may contain other bytes, such as newline characters, that frequently terminate the copying responsible for string-based buffer overflows<sup>[4]</sup> or some bits unknown to the attacker to prevent return-address clobbering with an integer overflow.

However, they cannot protect against a direct overwrite (by exploiting an indexing error) as they only check for corruption at function exit.

### Implications:

- Overhead: Adding stack canaries may add some overhead to the program and therefore slow execution.
- Compatibility issues may arise due to so many legacy code not using stack canaries.
- Using Stack Canaries leads to Reduced Vulnerability Surface and Enhanced Security

Other modifications that are more robust and can guard against more sophisticated attacks are :

1. Address Space Layout Randomisation(ASLR), it provides most comprehensive protection against attacks and is only vulnerable to information leakage attacks.
2. Using Data execution Prevention(DEP) for non-executable stacks, However this approach is vulnerable to Return-Oriented Programming Attacks.
3. Rearrangement of local variables, so that scalar variables are above array variables, so that in case of overflow, static variables are not affected.
4. Bound Checking to check for variables are in appropriate buffers
5. Null termination for handling string buffers
6. Using Static Analysis tools and using secure alternatives to vulnerable functions such as strcpy(), and writing secure and high quality code.

# Bibliography

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- [3] Laszló Szekeres, Mathias Payer, Tao Wei, and Dawn Song. SoK: Eternal war in memory. In *2013 IEEE Symposium on Security and Privacy*, pages 48–62. IEEE, 2013. URL <https://people.eecs.berkeley.edu/~dawnsong/papers/Oakland13-SoK-CR.pdf>.
- [4] Steve Anderson. Low-level security by example. *University of Pennsylvania*, 2019. URL <https://www.cis.upenn.edu/~sga001/classes/cis331f19/resources/low-level-security-by-example.pdf>.

# Appendices