Computer System Architecture Programming Project Part II

Project Description

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Project Description

In the first part of the project we implemented the Load/Store instructions and a simple memory, for this part of the project at the beginning, we have started to code for the rest of the instructions, except for CHK, floating point/vector operations, and trap. Then we expanded the User Interface to demonstrate the 1st program.

Starting with the instructions:

Transfer Instructions:

All of the instructions are implemented in the **Simulator.java and their** opcodes in Opcodes.java.

The transfer instructions move data between memory and the general-purpose and processor registers, processor registers, and I/O devices ,from one processor register to another ,and perform operations such as conditional moves. The conditional transfer test the values in the registers.

The illustration below shows the binary instruction code format of transfer instructions:

	Opcode		R		lx		l		Address	
0		5	6	7	8	9	10	11		15

Field	Bits	Note
Opcode	6	Specifies one of 64 possible instructions; Not all may be defined in this project
R	2	Specifies the General-Purpose Register GPR (RO-R3)
IX	2	Index Register (IXRO - IXR3)
I	1	If I =0, It is indirect addressing O.W no addressing
Addr	5	Specifies one of 32 locations

Opcode	Instruction	Description
10	JZ r, x, address[,I]	Jump If Zero: If c(r) = 0, then PC <- EA Else PC <- PC+1

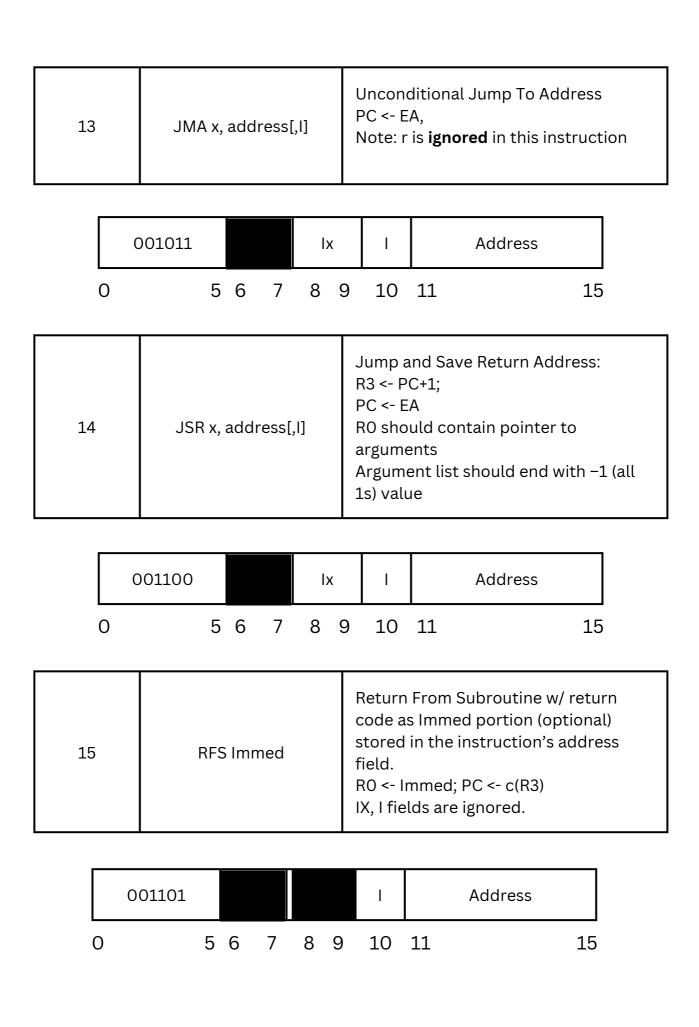
001000	R		lx		I		Address	
0	5 6	7	8	9	10	11		 15

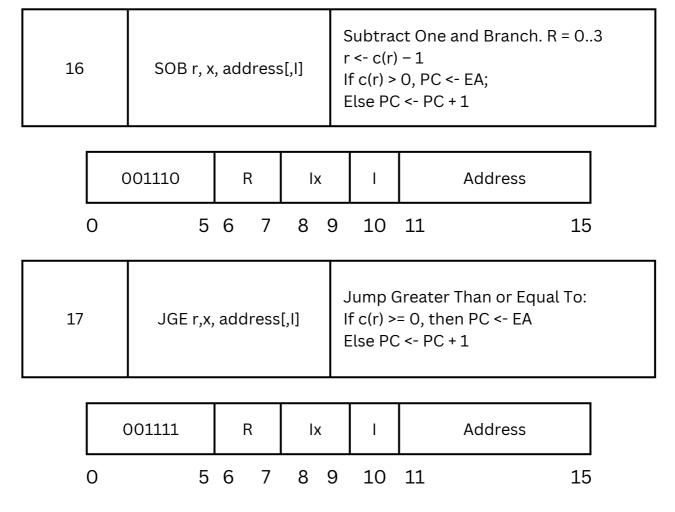
11	JNE r, x, address[,l]	Jump If Not Equal: If c(r) != 0, then PC < EA Else PC <- PC + 1
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001001		R		lx		_		Address	
0	5	6	7	8	9	10	11		_ 15

12	JCC cc, x, address[,I]	Jump If Condition Code cc replaces r for this instruction cc takes values 0, 1, 2, 3 0 Overflow, 1 Underflow, 2 Divzero, 3 Equalornot Register is to check; If cc bit = 1, PC <- EA Else PC <- PC + 1
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	001010		F	3	l	X	l		Address	
0		5	6	7	8	9	10	11		 15





c(r) shows the content of register r.

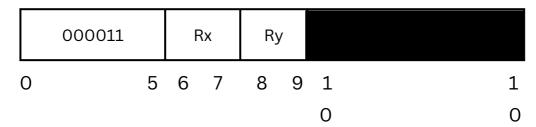
All these transfer codes are specified in Simulator.java after the Load and Store instructions.

Arithmetic and Logical Instructions:

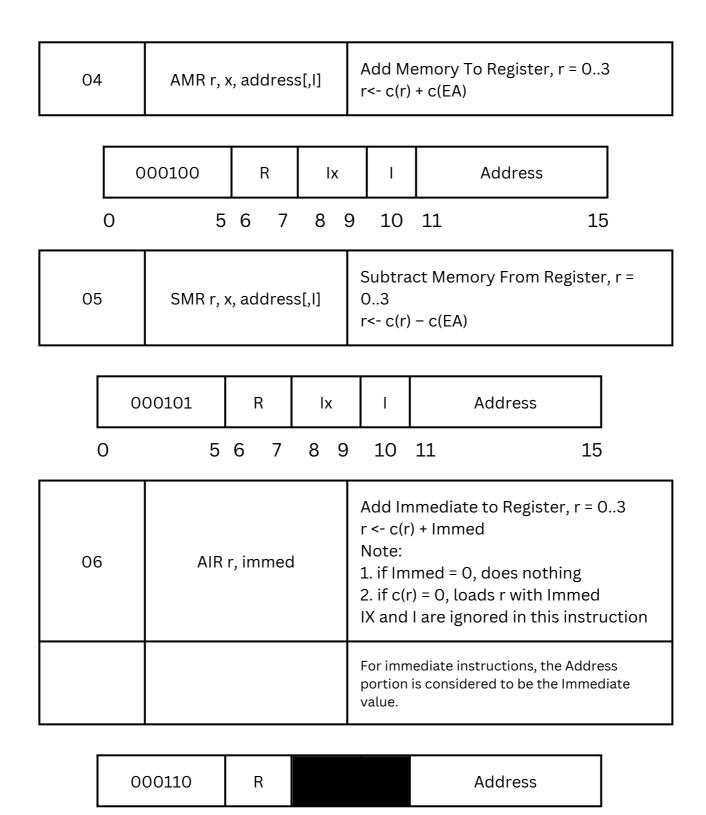
Logical instructions support and, or operations. Arithmetic instructions support addition, subtraction, multiplication, and division operations.

Arithmetic:

Certain arithmetic and logical instructions are register to register operations. The format of these instructions is:

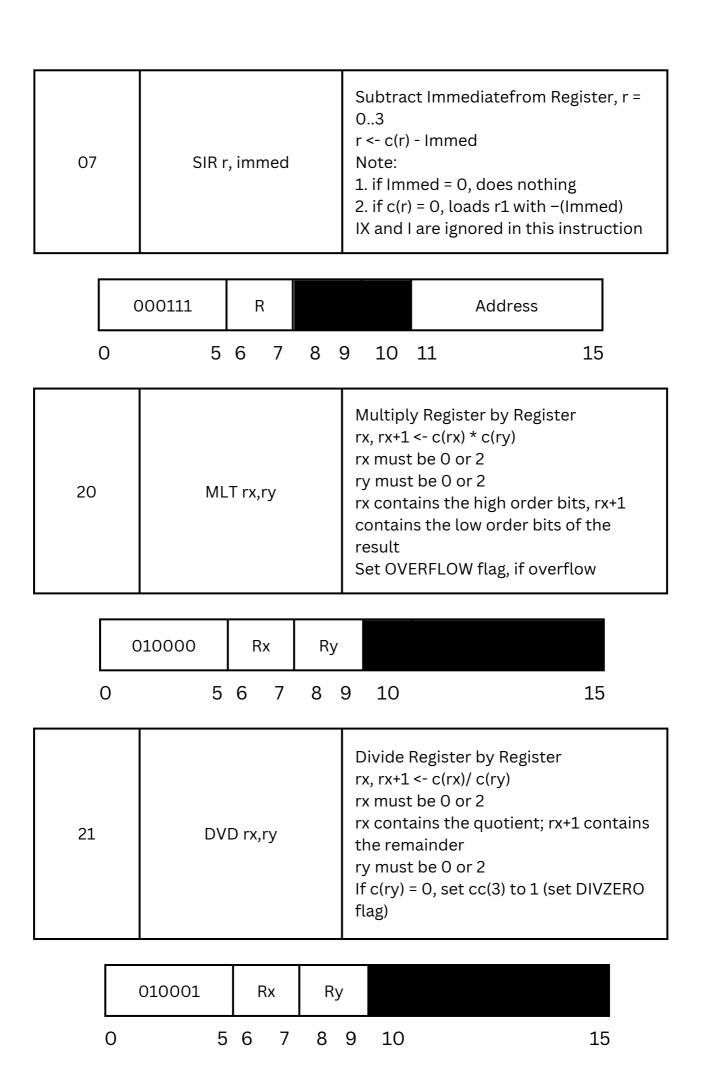


The blacked out portion means that portion of the instruction is ignored. Rx and Ry refer to one of RO-R3.



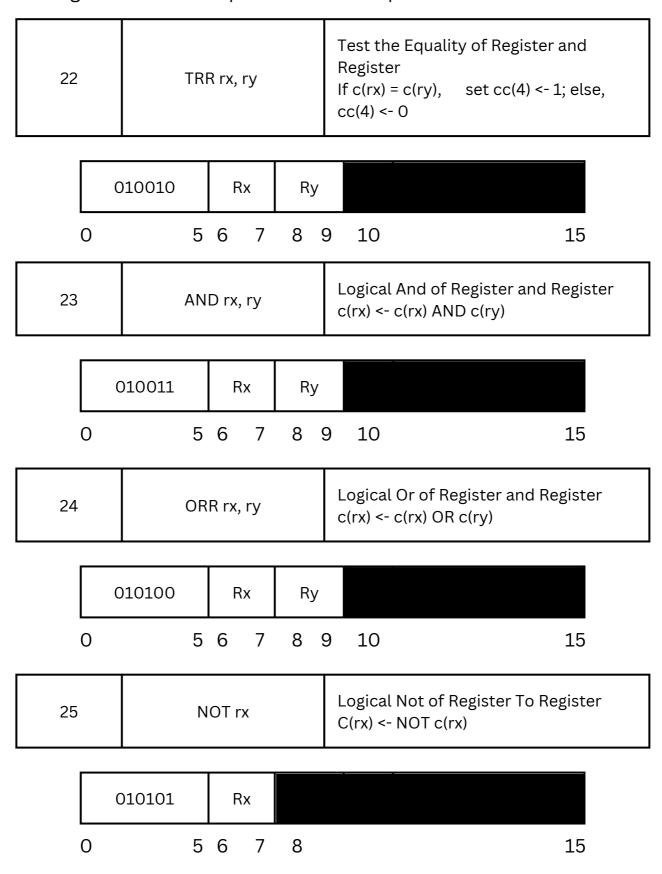
5 6

8 9



Logical:

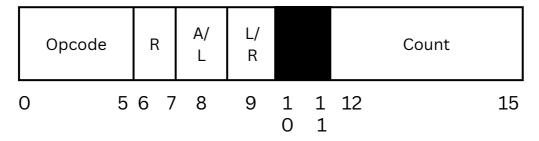
The logical instructions perform bitwise operations.



Shift/Rotate Operations:

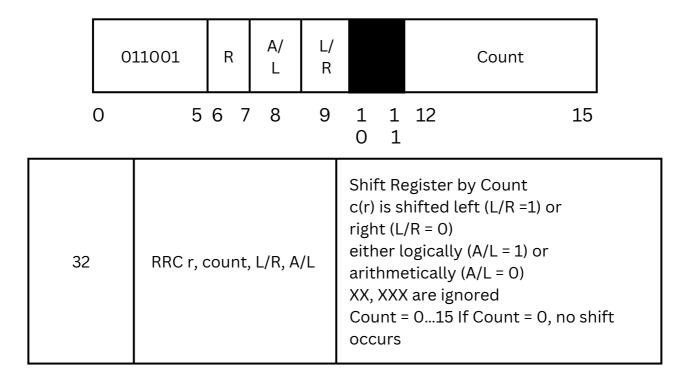
Shift and Rotate instructions control a datum in a sign in. they're used to shift the bits in the destination operand via one or more positions either to the left or proper.

The binary instruction code format of Shift and Rotate commands is as follows:



Field	Bits	Note
Opcode	6	Specifies one of 64 possible instructions; Not all may be defined in this project
R	2	Specifies the General-Purpose Register GPR (RO-R3)
A/L	2	Arithmetic Shift (A/L = 0); Logical Shift (A/L = 1)
L/R	2	Logical Rotate (L/R = 1)
Count	4	Specifies the Count for Operation

31	SRC r, count, L/R, A/L	Shift Register by Count c(r) is shifted left (L/R =1) or right (L/R = 0) either logically (A/L = 1) or arithmetically (A/L = 0) XX, XXX are ignored Count = 015 If Count = 0, no shift occurs
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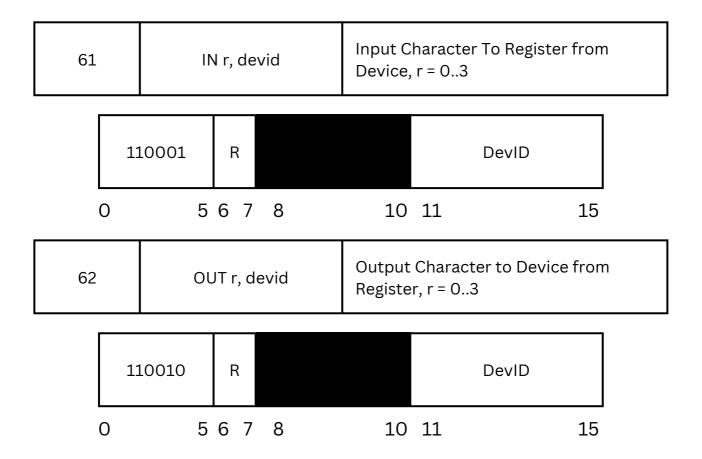


	011010	R		A/ L	L/ R				Count	
0	5	6	7	8	9	1	1	12		15

I/O Operations

I/O operations permits the imperative gadget of the computer to talk with the outdoor environment i.e., peripheral gadgets. For man or woman I/O, the instruction layout is:

Орсос	le	R				DevID	
0	5	6 7	8	10	11		15
DEVID	Dev	vice					
0	Co	nsule	: Keyboard	b			
1	Co	nsole	Printer				
2	Card Reader						
3-31	Co	nsole	Registers	s, switc	hes, et	C	



UI Display



- The **GPRO-3** are **4**, 16-bit General Purpose Registers. These can store both memory address and data.
- The IXR1-0 are 3, 16- bit Index Registers, they hold the current offset of a memory location. These are used for pointing to operand addresses when running the program.
- **PC: PC** is a 12-bit program counter which has the address of the next instruction to be executed from memory. It is a digital counter needed for faster execution of tasks as well as for tracking the current execution point.
- MAR: Memory Address Register is a 12-bit register which is used to access data and instructions from memory during the execution phase of instruction. MAR holds the memory location of data that needs to be accessed. When reading from memory, data addressed by MAR is fed into the MBR (memory buffer register).
- MFR: It is a 4-bit Machine Fault Register.

UI Display



- MBR: Memory Buffer Register is a 16-bit register which is used to store the data being transferred to and from the immediate access store. It contains the copy of designated memory locations specified by MAR. It acts as a buffer allowing the processor and memory units to act independently without being affected by minor differences in operation.
- **IR: IR** is a 16- bit register that holds the instruction currently being executed or decoded. Each instruction to be executed is loaded into the instruction register, which holds it while it is decoded, prepared and ultimately executed, which can take several steps.
- CC: The overflow and divided by zero flag will be shown in the CC
- **SS: The SS** is the Single Step button, it is used to execute the one step at a time in order to determine functioning.
- Run: This button is used to execute all of the instructions specified in the input and produces the final output.
- INIT: Initializing PC, Instruction, and memory
- **Program1**: This button is used to run the program 1