Project 2 Eyouel Kibret (G29247050) October 15, 2025

# Project 2

*Author Name Eyouel Kibret*

## 1 Problem Statement

We are required to analyze the following program/code sample.

ALGORITHM FindMedianOfMedians(nums\_list, k):

N = length(nums\_list)

IF N ≤ 5 THEN

SORT nums\_list using InsertionSort

RETURN nums\_list[k - 1]

END IF

// Divide array into groups of 5

partitions = []

FOR i FROM 0 TO N-1 INCREMENTING 5 FROM I EACH TIME

partition = elements of nums\_list from index i to i+4

ADD partition TO partitions

END FOR

// Find the median of each group

medians = []

FOR each partition IN partitions DO

SORT partition using InsertionSort

median = element at position floor(length(group) / 2)

ADD median TO medians

END FOR

// Recursively find the median of all medians

medianOfMedians = FindMedianOfMedians(medians, (length(medians) + 1) / 2)

// Partition nums\_list around the medianOfMedians

left = all elements in nums\_list less than medianOfMedians

equal = all elements in nums\_list equal to medianOfMedians

right = all elements in nums\_list greater than medianOfMedians

// Determine which partition contains the k-th smallest element

IF k ≤ length(left) THEN

RETURN FindMedianOfMedians(left, k)

ELSE IF k ≤ length(left) + length(equal) THEN

RETURN medianOfMedians

ELSE

RETURN FindMedianOfMedians(right, k - length(left) - length(equal))

END IF

END FUNCTION

## 2 Theoretical Analysis

Explain your theoretical estimate in 3-4 sentences.

The deterministic version of the Quick Select (median of medians method) uses divide and conquer strategy to select the median from an unsorted list of numbers. Until the first calling of “findMedianOfMedians” it takes because sorting takes because it is sorting a maximum of 5 elements. Finding the median of medians “findMedianOfMedians” method calling, takes Partitioning also takes O(n) because we are working with at most n elements in the list. After pivot selection, it is guaranteed that the pivot is between 30% of the elements in both directions. Therefore, we are working with 70% of the partition at a time which makes it T In conclusion it takes:

to find median of medians, to recursively search the larger portion if median is not found, and O(n) time for sorting, partitioning and grouping.

After solving using the method of substitution the time complexity becomes (Arora, 2022, Section 4.6.2)

## 3 Experimental Analysis

### 3.1 Program Listing

(Feel free to include only selected portions if you like. For example, I would like to know which values of “n” you ran the program for.)

def findMedianOfMedians(nums\_list, k):

# actual function to find the k-th smallest element using the median of medians algorithm

N = len(nums\_list)

if N <= 5:

# if the length of the numbers is less than or equal to 5,

# just sort it using the insertionSort helper function and return the median

nums\_list = insertionSort(nums\_list)

return nums\_list[k-1]

# Divide the array into groups of 5. There are n/5 groups.

partitions = [nums\_list[i:i + 5] for i in range(0, N, 5)]

# Sort the small groups using insertion sort. Since each group is

# of 5 elements, it takes a constant amount of time to sort those

# groups.

# Collect all the n/5 medians from the n/5 groups.

medians = [insertionSort(partition)[len(partition) // 2] for partition in partitions]

# recursively find the median of the medians

medianOfMedians = findMedianOfMedians(medians, (len(medians) + 1) // 2)

# Partition the array on the median of medians.

left, equal, right = partitionArray(nums\_list, medianOfMedians)

# recursively call the function based on the value of k

# if k is less than or equal to the length of the less than list

# it implies that the k-th smallest element is in the less than list

if k <= len(left):

return findMedianOfMedians(left, k)

# if k is greater than the length of the less than list

# and less than or equal to the length of the less than list + length of the equal to list

# it implies that the k-th smallest element is the median of medians

elif k <= len(left) + len(equal):

return medianOfMedians

# else it implies that the k-th smallest element is in the greater than list

# so recursively call the function on the greater than list

else:

return findMedianOfMedians(right, k - len(left) - len(equal))

if \_\_name\_\_ == "\_\_main\_\_":

power = 1

n = 10 \*\* power

k = n // 2 # Find the median value

nums\_list = [random.randint(1, n \* 10) for \_ in range(n)] # Random list of n values

start\_time = time.time()

result = findMedianOfMedians(nums\_list, k)

end\_time = time.time()

execution\_time = end\_time - start\_time

print(f"Execution time: ", execution\_time)

print(f"The {k}th smallest number is :", result)

Github repo:

I ran the program for values of n by changing the value of power from 1,2,3,4,5,6,7,8. “n” then becomes 10,100,1000,, *.*

### 3.2 Data Normalization Notes

Do you normalize the values by some constant? How did you derive that constant?

Yes, I normalized the values by a constant. Since the experimental values are so low usually in the microsecond range and theore5cal values much higher which doesn’t have a specific unit of measurement like seconds. I used the average ra5o to normalize the values of theore5cal values. I calculated the scaling factor by taking the sum of the experimental values and dividing by the sum of the theore5cal values.

Scaling factor(c) = ∑ 𝐸𝑥𝑝𝑜𝑛𝑒𝑛𝑡𝑖𝑎𝑙 𝑉𝑎𝑙𝑢𝑒𝑠(𝑖) / ∑ 𝑇ℎ𝑒𝑜𝑟𝑒𝑡𝑖𝑐𝑎𝑙 𝑉𝑎𝑙𝑢𝑒𝑠(𝑖)

I chose this approach because unlike one-point scaling factor, the scaling factor I chose balances the whole dataset. So, c = 2.04E+02/ 111111110 = 1.84E-06. Therefore, multiplying each theoretical value by this constant will result in values that are near each other so comparison will be easy.

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### 3.3 Output Numerical Data

|  |  |  |
| --- | --- | --- |
| n | Experimental | Theoretical |
| 10 | 1.31E-05 | 10 -> 1.84E-05 |
| *Remember, you have to choose the RIGHT values of n for the charts to be meaningful.* | | |
| 100 | 0.000114918 | 100 -> 1.84E-04 |
| 1000 | 0.000946522 | 1000 -> 1.84E-03 |
| 10000 | 0.009229422 | 10000 -> 1.84E-02 |
| 100000 | 0.095644951 | 100000 -> 1.84E-01 |
| 1000000 | 1.2084589 | 1000000 -> 1.84E-00 |
| 1.00E+07 | 15.29438496 | 1.00E+07 -> 1.84E-01 |
| 1.00E+08 | 187.0289783 | 1.00E+08 -> 1.84E-02 |

**3.4 Graph**

**3.5 Graph Observations**

From the above graph I can specify that:

* The theoretical values align well with the experimental values. There is no tremendous change(error) between them.
* Without the scaling factor the experimental data points would have been a flat line and the theoretical values would have been an increasing one. Therefore, a scaling constant of 1.84E-06 modified the magnitude of the theoretical values so that they can be compared.
* When the value of n increases, the experimental values grow in sync with the values of the theoretical showing minimal difference. It implies that it follows the expected linear trend
* For much higher value of n both lines grow in strong concordance, this implies that the theoretical implementatin closely aligns with and supports the experimental results.
* There is a minor difference between the theoretical and experimental values for n and but the growth trend is similar.

## 4 Conclusions

After careful analysis the theoretical analysis of code became . But to test it, a basic python program was built using built in methods and functions. After writing the python program, when working with smaller values of “n” there was no memory allocation issue. However, when working with higher values of “n” especially from n starting at , memory allocation was a huge problem so the program couldn’t be tested for this value and above. However, for values under the program worked well. The python program returned the median and the time taken for each n (by changing the value of power in the code), in seconds. The usage of a scaling factor was necessary because the values of the theoretical values with diﬀerent values of “n” where large when compared to the experimental results. A scaling factor of 1.84E-06 was taken for easy comparison. It can be concluded that this algorithm always ensures a good partition and running at linear time makes it preferable rather than sorting a list and finding the median which takes

**Reference**

Arora, A. (2022). *Analysis and Design of Algorithms* (3rd ed.). Cognella, Inc.