

Lecture 5: Function Approximation + The Curse of Dimensionality

Dynamic Programming

Thomas Jørgensen



Recap: Linear interpolation

- **Recap (# + 1 known nodes!):**

- Let $\hat{x} = [\hat{x}_1, \hat{x}_2, \dots, \hat{x}_{\#+1}]$ be *grid points*
- Let $\hat{f} = [\hat{f}_1, \hat{f}_2, \dots, \hat{f}_{\#+1}] = [f(\hat{x}_1), f(\hat{x}_2), \dots, f(\hat{x}_{\#+1})]$ be the *function values* at these points
- **Linear interpolation** then is

$$f(x) \approx \check{f}(x; \hat{f}) = \hat{f}_n + \frac{\hat{f}_{n+1} - \hat{f}_n}{\hat{x}_{n+1} - \hat{x}_n} (x - \hat{x}_n)$$

where $\hat{x}_n \leq x < \hat{x}_{n+1}$

- **Requirements for good interpolants:**

- ① Fast to set up and evaluate
- ② Shape preserving (monotonicity and concavity)
- ③ High flexibility per grid point
- ④ Continuously differentiable



Beyond linear interpolation

- ① **Finite element methods** (local)
 - ① Piecewise splines
 - ② Local basis functions (e.g. B-splines)
- ② **Spectral methods** (global basis functions)



1a. Piecewise splines

- **Linear interpolation** can also be written as

$$\check{f}(x; \hat{f}) = \sum_{n=1}^{\#} \mathbf{1}_{x \in [\hat{x}_n; \hat{x}_{n+1})} \phi_n(x; \hat{f})$$

$$\phi_n(x; \hat{f}) = [a_n + b_n(x - \hat{x}_n)]$$

where the (a_n, b_n) 's are chosen such that

$$\begin{aligned} \phi_n(\hat{x}_n) &= \hat{f}_n, \forall n \in \{1, 2, \dots, \#\} \\ \phi_n(\hat{x}_{n+1}) &= \hat{f}_{n+1}, \forall n \in \{1, 2, \dots, \#\} \end{aligned}$$

- **Higher order piecewise splines:**

$$\phi_n(x; \hat{f}) = a_n + b_n(x - x_n) + c_n(x - x_n)^2 + \dots$$



1a. Piecewise cubic splines I

- Piecewise cubic splines:**

$$\phi_n(x; \hat{f}) = a_n + b_n(x - x_n) + c_n(x - x_n)^2 + d_n(x - x_n)^3$$

where the (a_n, b_n, c_n, d_n) 's are chosen such that

$$\phi_n(\hat{x}_n) = \hat{f}_n, \forall n \in \{1, 2, \dots, \#\}$$

$$\phi_n(\hat{x}_{n+1}) = \hat{f}_{n+1}, \forall n \in \{1, 2, \dots, \#\}$$

$$\phi'_n(\hat{x}_{n+1}) = \phi'_{n+1}(\hat{x}_{n+1}), \forall n \in \{1, 2, \dots, \# - 1\}$$

$$\phi''_n(\hat{x}_{n+1}) = \phi''_{n+1}(\hat{x}_{n+1}), \forall n \in \{1, 2, \dots, \# - 1\}$$

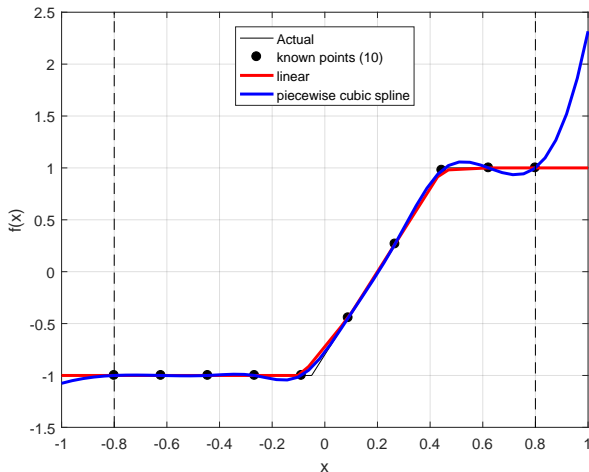
$$\phi''_1(\hat{x}_1) = 0 \text{ (could be something else)}$$

$$\phi''_n(\hat{x}_{\# + 1}) = 0 \text{ (could be something else)}$$

(number of equations: $2\# + 2(\# - 1) + 2 = 4\#$)



Illustration: Linear and cubic spline

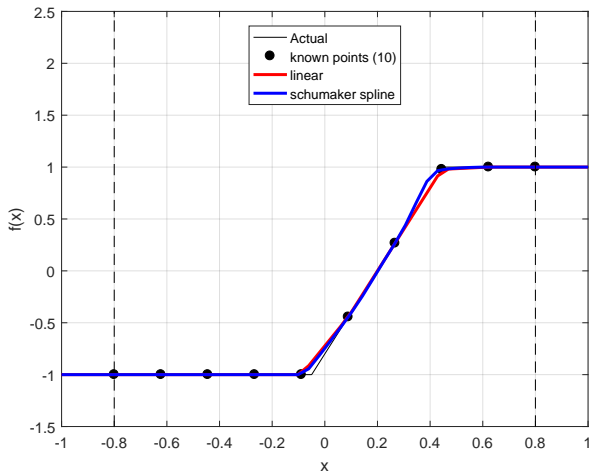


1a. Piecewise cubic splines II

- **Pro piecewise cubic splines:**
 - ① No error at known points
 - ② More flexible than linear interpolation
 - ③ Twice continuously differentiable
- **Con piecewise cubic splines:**
 - ① Slower than linear interpolation
 - ② Not shape-preserving
 - ③ Poor extrapolation
- **Shape-preserving:** Schumaker (1983)



Illustration: Schumaker (1983) spline



- Discrete choice example on white-board



1b. B-splines

- Until now **piecewise interpolation**:

$$\check{f}(x) = \sum_{n=1}^{\#} \mathbf{1}_{x \in [x_n; x_{n+1})} \phi_n(x; \hat{f})$$

- **Alternative finite element method:** *Local basis functions*

$$\check{f}(x) = \sum_{n=1}^{\#} \omega_n(\hat{f}) \phi_n(x)$$

where we note that the basis functions, $\phi_n(x)$, are independent of \hat{f}



1b. Linear B-spline

① Choose

$$\omega_n(\hat{f}) = f_n$$

$$\phi_n(x) = \begin{cases} 0 & \text{if } x \notin [x_{n-1}, x_{n+1}] \\ \frac{x - x_{n-1}}{x_n - x_{n-1}} & \text{if } n > 1 \text{ and } x \in [x_{n-1}, x_n] \\ \frac{x_{n+1} - x}{x_{n+1} - x_n} & \text{if } n < \# + 1 \text{ and } x \in [x_n, x_{n+1}] \end{cases}$$

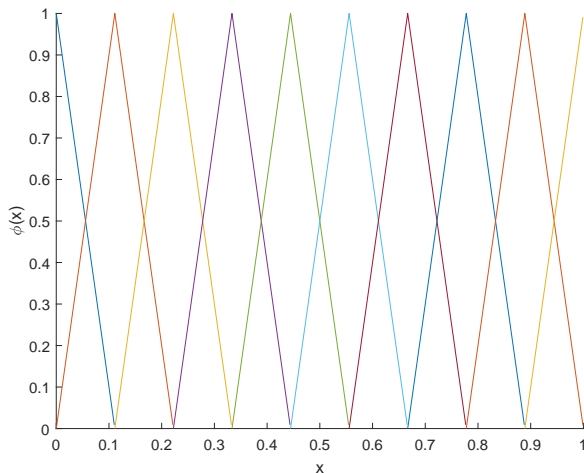
② then this is yet another way to do **linear interpolation**

$$\begin{aligned} \check{f}(x) &= 0 + \cdots + 0 + \hat{f}_n \phi_n(x) + \hat{f}_{n+1} \phi_{n+1}(x) + 0 + \cdots + 0 \\ &= \hat{f}_n \frac{x_{n+1} - x}{x_{n+1} - x_n} + \hat{f}_{n+1} \frac{x - x_n}{x_{n+1} - x_n} \\ &= \hat{f}_n + \frac{\hat{f}_{n+1} - \hat{f}_n}{x_{n+1} - x_n} (x - x_n) \end{aligned}$$

③ More general $\phi_n(x) \rightarrow$ higher-order **B-splines**



1b. The $\phi_n(x)$'s for a linear B-spline (10 nodes)



2. Regression with polynomials

- Interpolant with P *global* basis functions

$$\check{f}(x) = [\phi_1(x) \phi_2(x) \cdots \phi_P(x)] \omega(\hat{f}) = \sum_{i=1}^P \omega_i(\hat{f}) \phi_i(x)$$

- $\omega(\hat{f})$ can be found by **OLS regression**:

$$X = \begin{pmatrix} \phi_1(\hat{x}_1) & \phi_2(\hat{x}_1) & \cdots & \phi_P(\hat{x}_1) \\ \phi_1(\hat{x}_2) & \phi_2(\hat{x}_2) & \cdots & \phi_P(\hat{x}_2) \\ \vdots & \vdots & \ddots & \vdots \\ \phi_1(\hat{x}_{\# + 1}) & \phi_2(\hat{x}_{\# + 1}) & \cdots & \phi_P(\hat{x}_{\# + 1}) \end{pmatrix}$$

$$Y = [\hat{f}_1 \hat{f}_2 \cdots \hat{f}_{\# + 1}]'$$

$$\omega(\hat{f}) = [\omega_1 \omega_2 \cdots \omega_P]' = (X'X)^{-1}X'Y$$

- **Convention:** $\phi_i(x) = T_i(g(x))$
- **Ordinary polynomials:** $T_i(z) = z^{i-1}$ and $g(x) = x$



2. Regression with Chebyshev polynomials

- **Interval** $x \in [a; b]$: Use $g(x) = -1 + 2\frac{x-a}{b-a} \in [0, 1]$
- **Chebyshev polynomials**

$$T_i(z) = \cos(i \cos^{-1}(z))$$

- **Orthogonal**

$$\sum_{n=1}^{\#} T_i(z_n) T_j(z_n) = 0 \text{ for } i \neq j$$

if nodes are chosen as

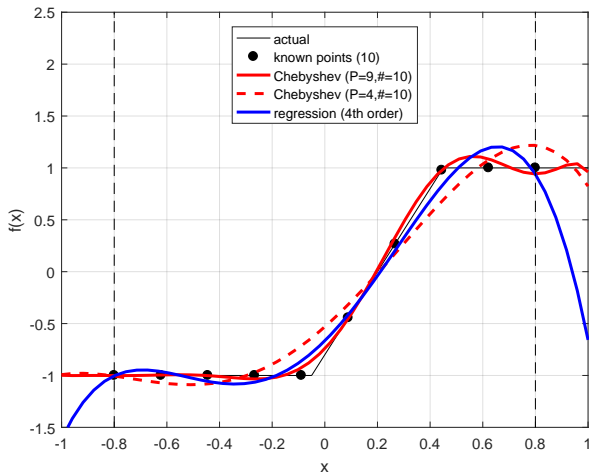
$$z_n = -\cos\left(\frac{2n-1}{2(\#+1)}\pi\right) \in [-1, 1] \text{ for } n = 1, \dots, \# + 1$$

$$x_n = g^{-1}(z_n) = a + \frac{z_n + 1}{2}(b - a)$$

- **Drawback:** Not shape-preserving; can be added at a cost



Illustration: Chebyshev and regression



Multi-dimensional interpolation

- **Linear interpolation:** Simple (next slide)
- **Higher order piecewise splines:** Complicated to find the parameters and no simple shape preserving splines available
- **Basis functions:** Include cross products of one dimensional basis functions
 - ① Global polynomial regression
 - ② Local B-splines
- **Frontier topics:**
 - ① *Global sparse grids* (not all cross products) (Judd et. al. 2014)
 - ② *Locally adaptive sparse grids* (+ hierarchy of basis functions) (Brumm and Scheidegger, 2014)
 - ③ *Scattered* data (triangulation and barycentric interpolation)



Bi-linear interpolation

- For $f : \mathbb{R}^2 \rightarrow \mathbb{R}$ known at the grid points

$$(x_1, y_1), (x_2, y_2), \dots, (x_{\#+1}, y_{\#+1})$$

- 1 Locate **neighboring points**

$$x_n \leq x < x_{n+1}$$

$$y_m \leq y < y_{m+1}$$

- 2 Interpolate in x -dimension for **constant** y

$$f_{n,m} \equiv f(x_n, y_m) + \frac{f(x_{n+1}, y_m) - f(x_n, y_m)}{x_{n+1} - x_n} (x - x_n)$$

$$f_{n,m+1} \equiv f(x_n, y_{m+1}) + \frac{f(x_{n+1}, y_{m+1}) - f(x_n, y_{m+1})}{x_{n+1} - x_n} (x - x_n)$$

- 3 Interpolate **across** y

$$\check{f}(x, y) = f_{n,m} + \frac{(f_{n,m+1} - f_{n,m})}{y_{n+1} - y_n} (y - y_n)$$

- TASK:** illustrate this in 2 dimensions
- Similar in higher dimensions



The Three Curses of Dimensionality

- ❶ **Multiple states:** Exponential growth in total number of grid points for tensor product grids
- ❷ **Multiple choices:** Harder to solve the optimization problem given the states
- ❸ **Multiple shocks:** Exponential growth in the number of quadrature points needed to approximate the continuation value

→ lots of tips and tricks to alleviate the curse in practice



Interpolation

The Curse of
Dimensionality

Programming

Until next

1. Think!

- ① Only put in stuff you need
- ② Can you use a discrete state or choice instead of a continuous one?



2. Use analytical structure

- ① The problem might contain **intra-temporal sub-problems**, which can be solved fast (e.g. in closed form)

$$V_t(M_t) = \max_{C_t, D_t} u(C_t^\alpha D_t^{1-\alpha}) + \beta \mathbb{E}_t [V_{t+1}(M_{t+1})]$$

s.t.

$$A_t = M_t - (p_C C_t - p_D D_t)$$

$$M_{t+1} = RA_t + Y_{t+1}$$

- ② There might be freedom of choice wrt. to **state variables**

$$V_t(A_{t-1}, Y_t) = \max_{C_t, D_t} u(C_t^\alpha D_t^{1-\alpha}) + \beta \mathbb{E}_t [V_{t+1}(A_t, Y_{t+1})]$$

s.t.

$$A_t = \underbrace{RA_{t-1} + Y_t}_{=M_t} - p_C C_t - p_D D_t$$

- ③ The problem might be **scaleable** in a state \rightarrow the problem can be solved in ratio form with one fewer state variable



3. Taste shocks I

- ① **i.i.d. taste shocks** for working or not

$$\begin{aligned}
 V_t(M_t, \varepsilon_t^0, \varepsilon_t^1) &= \max_{L_t \in \{0,1\}} \left\{ v_t(M_t | L_t) + \sigma_\varepsilon \varepsilon_t^{L_t} \right\} \\
 v_t(M_t | L_t) &= \max_{C_t} u(C_t, L_t) + \beta \mathbb{E}_t [V_{t+1}(\bullet_{t+1})] \\
 &\text{s.t.} \\
 M_{t+1} &= R(M_t - C_t) + W \cdot L_t
 \end{aligned}$$

- ② Assume that ε_t^0 and ε_t^1 are **Extreme Value Type I** then

$$\begin{aligned}
 \mathbb{E}[V_t(M_t, \varepsilon_t^0, \varepsilon_t^1) | M_t] &= \sigma_\varepsilon \log \left(\sum_{j \in \{0,1\}} \exp \left(\frac{v_t(M_t | L_t)}{\sigma_\varepsilon} \right) \right) \\
 &\equiv \mathcal{W}_t(M_t)
 \end{aligned}$$



3. Taste shocks II

- 1 We only need to **find choice-specific value functions**

$$v_t(M_t|L_t) = \max_{C_t} u(C_t, L_t) + \beta \mathbb{E}_t [\mathcal{W}_t(M_{t+1})]$$

- 2 The **choice probabilities** for the discrete choices are

$$\begin{aligned} \Pr(L_t = 1|M_t) &= \Pr(v_t(M_t|1) - v_t(M_t|0) \geq \sigma_\varepsilon(\varepsilon_t^0 - \varepsilon_t^1)) \\ &= \frac{\exp(v_t(M_t|1)/\sigma_\varepsilon)}{\sum_{j \in \{0,1\}} \exp(v_t(M_t|j)/\sigma_\varepsilon)} \end{aligned}$$

- **Question:** What happens as $\sigma_\varepsilon \rightarrow 0$ or $\sigma_\varepsilon \rightarrow \infty$?



4. Pre-computations

- For the problem

$$V_t(M_t) = \max_{C_t} u(C_t) + \beta \mathbb{E}_t [V_{t+1}(M_{t+1})]$$

s.t.

$$A_t = M_t - C_t$$

$$M_{t+1} = RA_t + Y_{t+1}$$

- 1 Construct a grid of the **post-decision state** A_t and pre-compute the **post-decision value function**

$$W_t(A_t) \equiv \mathbb{E}_t [V_{t+1}(RA_t + Y_{t+1})]$$

- 2 Solve the **simpler problem**

$$V_t(M_t) = \max_{C_t} u(C_t) + \beta W_t(A_t)$$

- In **infinite horizon**: Can be a good idea to iterate on W_t instead of V_t (see Hull 2015)



5. Grids and inverse interpolation

- **Grid points should be spend wisely**
→ most where there is most curvature

① e.g. *Rust-spaced*

$$\begin{aligned} i \geq 2 : x_i &= x_{i-1} + \frac{\bar{x} - x_{i-1}}{(n - i + 1)^\phi} \\ x_1 &= \underline{x} \end{aligned}$$

② or *polynomially spaced*

$$x_i = \underline{x} + (\bar{x} - \underline{x}) \left(\frac{i - 1}{N} \right)^\phi$$

- **Value function inherits curvature of utility function:**
Interpolate $u^{-1}(V_t)$, instead of V_t , and convert back.
(require that the transformation $u^{-1}(\bullet)$ is monotone)



6. Time-iterations

- We know that optimal interior choices satisfy the **Euler-equation**

$$u'(C_t) = \beta R \mathbb{E}_t [u'(C_{t+1}^*(\Gamma(M_t, C_t)))]$$

- **Idea:** Find $C_t^*(M_t)$ by **solving the Euler-equation in** $C_t < M_t$, else $C_t = M_t$
 - This can be easier to solve than the optimization problem itself
 - We do not need the value function at all
 - More accurate - why?
- Basis for the *endogenous grid point method* we will discuss in a number of lectures



Programming principles

- ① **Correct code** beats fast/smart incorrect code
- ② **Understandable code** is alfa and omega - others and your future-self need to be able to understand it
 - **Names.** Variables and functions should have precise names
 - **Structure:** Each part of the code should have its own special purpose; preferably numbered. Repeated complex calculations should be in functions.
 - **Comments:** Short and add new information
 - **Testing:** Use `assert(x==1, 'error msg.')`
 - **Replicability:** Everything should be called from a single file
- ③ Code should only be **optimized** when the bottleneck has been located
 - **Time:** `tic; fun(x) toc;`
 - **Profile:**
`profile on; fun(x); profile off; profile viewer;`

More: `appendix_good_programming.mlx`
in the MATLAB online course



Optimization principles

- ❶ **Optimize the algorithm you use.** Only calculate what you need and especially avoid repeating expensive operations. Below are some “computational costs”
 - addition, subtraction, comparison = 1
 - multiplication = 4
 - division = 10
 - exp/log = 50
 - power = 100
- ❷ **Tell the computer what you know in advance**
 - pre-allocate memory
- ❸ **Work on consecutive chunks of memory**
 - correct loop order (MATLAB is “column-major”: first index as the inner-most loop)
 - vectorize all you can
- ❹ **(Parallize) (*parfor* in MATLAB)**

Do not optimize your code before you are sure it is working and you know where the bottleneck is



Beyond MATLAB

- Numerical precision (floating point arithmetic)
- Especially **loops** are slow in MATLAB
- **Alternatives:**
 - ① **Python:** Similar speed and complexity (but free)
 - ② **Fortran:** Very fast, but more complex (easy parallization)
 - ③ **C++:** Very fast, but more complex (easy parallization)
 - ④ **Julia:** Almost as fast as C++/Fortran, but still simple (still not fully developed)
- MATLAB can call e.g. C++ code using the **mex** interface
→ what I am doing all the time:
 - ① MATLAB for setup and figures
 - ② C++ with OpenMP for parallization of the central stuff
 - ③ NLOpt as an optimizer in C++



Until next

- **Ensure that you understand:**
 - ① Linear interpolation
 - ② Interpolation by regression
 - ③ The curse of dimensionality
 - ④ The use of taste shocks
 - ⑤ The method of time iteration
- Go to **PadLet** and ask or answer a question
(<https://padlet.com/tjo2/dp>)
- **Next time:** Recap! Send me an email (tjo@econ.ku.dk) with stuff you want me to recap.
Remaining time will be devoted to exercises.

