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A Comparison of Approaches to  
Combinatorial Optimisation for  
Touristic Route Planning

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# Abstract

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## 1 Introduction

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## 2 Problem Formulation

### 2.1 Problem Description

Given a positive integer  $d$  and a graph  $G = (V, E)$ , where  $V$  is set of locations including a designated starting point  $s$  and  $E$  is a set of weighted edges linking every location to every other location, find a route that:

1. Visits all nodes  $V \setminus c$  once.
2. Starts and finishes at  $s$ , having visited it  $d$  times, without ever visiting consecutively.
3. Minimizes both the cumulative edge weights in the route and the variance in cumulative weight between each visit to  $s$ .

### 2.2 Inputs and Outputs

Inputs:

- $d$ : The number of times  $s$  should be visited in a route. Contextually,  $d$  represents the number of days a tourist will spend on their trip.  $d \in \mathbb{Z}, d > 0$
- $G = (V, E)$ : A pair comprising of:
  - $V$ : A set of nodes representing locations the tourist would like to visit.  $v \in V, v = (x, y, t)$ , a triple comprising of:
    - $x$ : Longitude, indicating the location's geographic east-west position on the earth<sup>1</sup>.  
 $x \in \mathbb{Q}, -180 \leq x \leq 180$ .
    - $y$ : Latitude, indicating the location's geographic north-south position on the earth<sup>2</sup>.  
 $y \in \mathbb{Q}, -90 \leq y \leq 90$ .
    - $t$ : Duration, in minutes, indicating how much time to spend at this location.  
 $t \in \mathbb{Z}, t > 0$ .
  - $E$ : A set of edges  $e \in E$  that connects every node to every other node, bidirectionally.  $e = (v_1, v_2, w)$ , a triple comprising of:

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<sup>1</sup>While the coordinates of our locations are included in  $V$ , they are not directly tied to the weight of our edges  $E$ , which are based on time and not distance.

<sup>2</sup>See footnote 1

- $v_1$ : A location representing the origin of the edge.  
 $v_1 \in V$ .
- $v_2$ : A location representing the destination of the edge.  
 $v_2 \in V$ .
- $w$ : A weight indicating the sum of the time it takes to travel from  $v_1$  to  $v_2$  and the time the tourist wishes to spent at  $v_2$ .  
 $w \in \mathbb{Z}, w > 0$ .
- $s$ : Starting point that should be visited  $d$  times. Contextually,  $s$  represents where the tourist is staying and will return to at the end of each day.  
 $s \in V$ .

Outputs:

- $R$ : A valid route satisfying all constraints, represented as an ordered sequence of locations.  
 $R = [r_1, r_2, \dots, r_n], r_i \in V$ .

## 2.3 Optimisation

As previously mentioned in the Problem Description, our goal is to find a route that minimises the cumulative weight and the variance in route weight between each visit to  $s$ . To accomplish this the following cost function is applied to each route:

$$Cost(R) = W \times (1 + \sigma) \quad (1)$$

Where  $W$  is the sum of the weights of all edges traversed in the route and  $\sigma$  is the standard deviation of the sum of weights between each visit to  $s$ :

$$\sigma = \sqrt{\frac{\sum_{i=0}^d (x_i - \mu)^2}{d}}, x_i \in X \quad (2)$$

Where  $R$  is divided into sections between each visit to  $s$  and  $X$  is a list of the sum of weights within these sections.  $\mu$  is the mean cumulative weight of each  $x_i$ .

### 3 Literature Review

Write literature review

## 4 Algorithms Investigated

Paragraph describing different types of algorithm used (Routing, Cluster and Routing, Evolutionary, etc.)

### 4.1 Brute Force

Write brute force explanation

### 4.2 Clustering

Write clustering explanation

Once divided into clusters, each cluster can be solved using our brute force method (with  $d = 1$ ) or with traditional TSP algorithms. TSP algorithms implemented in this project include:

### 4.3 Greedy

Write greedy explanation

## 5 Evaluation and Comparison

Write paragraph about experiment process. Comparison based on computation time and route evaluation. Describe how route is evaluated. Describe data being tested on.



## 6 Conclusion

Write conclusion

## 7 Bibliography

Fill in bibliography