

## Announcement

- Postponed HW 6! HW5 due today midnight.
- See Canvas announcements re midterm (Fri Feb 16, in-class)
- Solving B&V book problems is good practice.
- TA review session this Fri (recorded, but with live Zoom office hours afternoon; please see Natalia's earlier announcement on Canvas.

## 6. Convex optimization problems: GP, SDP, and multi-objective optimization

- • geometric programming (*GP*)
  - generalized inequality constraints
- • semidefinite programming
  - vector (multi-objective) optimization

# Geometric programming

application:

- mechanical design/structures
- circuit design (amplifiers, chemical reactions PLLs)
- power allocation in wireless networks

## monomial function

$$f(x) = cx_1^{a_1}x_2^{a_2}\cdots x_n^{a_n}, \quad \text{dom } f = \underline{\mathbf{R}}_{++}^n$$

with  $c > 0$ ; exponent  $a_i$  can be any real number

posynomial function: sum of monomials

$$f(x) = \sum_{k=1}^K c_k x_1^{a_{1k}} x_2^{a_{2k}} \cdots x_n^{a_{nk}}, \quad \text{dom } f = \mathbf{R}_{++}^n$$

## geometric program (GP)

$$\left[ \begin{array}{ll} \text{minimize} & f_0(x) \\ \text{subject to} & \frac{f_i(x)}{h_i(x)} \leq 1, \quad i = 1, \dots, m \\ & \underline{\frac{h_i(x)}{h_i(x)} = 1}, \quad i = 1, \dots, p \end{array} \right.$$

with  $f_i$  posynomial,  $h_i$  monomial

## Geometric program in convex form

change variables to  $y_i = \log x_i$ , and take logarithm of cost, constraints

- monomial  $f(x) = cx_1^{a_1} \cdots x_n^{a_n}$  transforms to  $\log c + \sum_{i=1}^n a_i \log x_i$   
 $\longrightarrow \log f(e^{y_1}, \dots, e^{y_n}) = a^T y + b \quad (b = \log c) \quad \text{affine in } y$

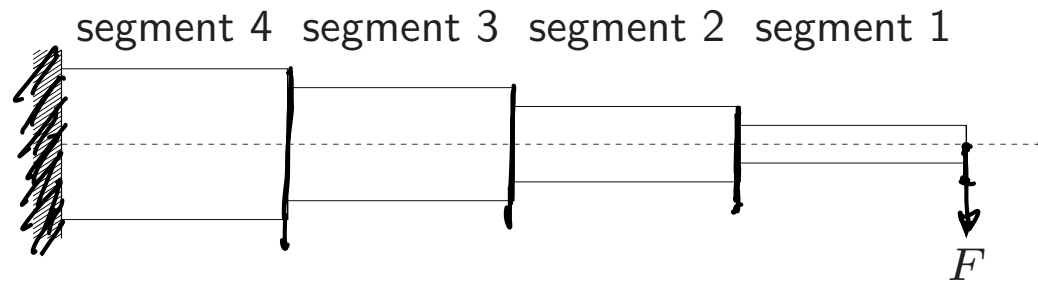
- posynomial  $f(x) = \sum_{k=1}^K c_k x_1^{a_{1k}} x_2^{a_{2k}} \cdots x_n^{a_{nk}}$  transforms to

$$\log f(e^{y_1}, \dots, e^{y_n}) = \log \left( \sum_{k=1}^K e^{a_k^T y + b_k} \right) \quad (b_k = \log c_k)$$

- geometric program transforms to convex problem

$$\left[ \begin{array}{ll} \underset{y \in \mathbb{R}^n}{\text{minimize}} & \log \left( \sum_{k=1}^K \exp(\underline{a_{0k}^T} y + \underline{b_{0k}}) \right) \\ \text{subject to} & \log \left( \sum_{k=1}^K \exp(\underline{a_{ik}^T} y + \underline{b_{ik}}) \right) \leq 0, \quad i = 1, \dots, m \\ & Gy + d = 0 \end{array} \right.$$

## Design of cantilever beam



- $N$  segments with unit lengths, rectangular cross-sections of size  $w_i \times h_i$
- given vertical force  $F$  applied at the right end

### design problem

minimize total weight  
 subject to upper & lower bounds on  $w_i, h_i$   
upper bound & lower bounds on aspect ratios  $h_i/w_i$   
upper bound on stress in each segment  
upper bound on vertical deflection at the end of the beam

variables:  $w_i$ ,  $h_i$  for  $i = 1, \dots, N$

## objective and constraint functions

- total weight  $w_1 h_1 + \dots + w_N h_N$  is posynomial

$$\alpha h_i w_i^{-1} \leq 1$$

$$\beta h_i^{-1} w_i \leq 1$$

- aspect ratio  $h_i/w_i$  and inverse aspect ratio  $w_i/h_i$  are monomials

$$w_i^{-1} h_i^{-2} \leq 1 \quad \forall i$$

- maximum stress in segment  $i$  is given by  $6iF/(w_i h_i^2)$ , a monomial

- the vertical deflection  $y_i$  and slope  $v_i$  of central axis at the right end of segment  $i$  are defined recursively as

$$\underline{v_i} = 12(i-1/2) \frac{F}{E \underline{w_i h_i^3}} + \underline{v_{i+1}} \rightarrow$$

$$\underline{y_i} = 6(i-1/3) \frac{F}{E \underline{w_i h_i^3}} + \underline{v_{i+1}} + \underline{y_{i+1}}$$

for  $i = N, N-1, \dots, 1$ , with  $v_{N+1} = y_{N+1} = 0$  ( $E$  is Young's modulus)

$v_i$  and  $y_i$  are posynomial functions of  $w$ ,  $h$

## formulation as a GP

$$\left[ \begin{array}{ll} \underset{w, h}{\text{minimize}} & w_1 h_1 + \cdots + w_N h_N \\ \text{subject to} & w_{\max}^{-1} w_i \leq 1, \quad w_{\min} w_i^{-1} \leq 1, \quad i = 1, \dots, N \\ & h_{\max}^{-1} h_i \leq 1, \quad h_{\min} h_i^{-1} \leq 1, \quad i = 1, \dots, N \\ & S_{\max}^{-1} w_i^{-1} h_i \leq 1, \quad S_{\min} w_i h_i^{-1} \leq 1, \quad i = 1, \dots, N \\ & 6iF\sigma_{\max}^{-1} w_i^{-1} h_i^{-2} \leq 1, \quad i = 1, \dots, N \\ & y_{\max}^{-1} y_1 \leq 1 \end{array} \right.$$

note

- we write  $w_{\min} \leq w_i \leq w_{\max}$  and  $h_{\min} \leq h_i \leq h_{\max}$

$$\underbrace{w_{\min}/w_i \leq 1, \quad w_i/w_{\max} \leq 1, \quad h_{\min}/h_i \leq 1, \quad h_i/h_{\max} \leq 1}$$

- we write  $S_{\min} \leq h_i/w_i \leq S_{\max}$  as

$$S_{\min} w_i / h_i \leq 1, \quad h_i / (w_i S_{\max}) \leq 1$$

# Minimizing spectral radius of nonnegative matrix (skip)

## Perron-Frobenius eigenvalue $\lambda_{\text{pf}}(A)$

- exists for (elementwise) positive  $A \in \mathbf{R}^{n \times n}$
- a real, positive eigenvalue of  $A$ , equal to spectral radius  $\max_i |\lambda_i(A)|$
- determines asymptotic growth (decay) rate of  $A^k$ :  $A^k \sim \lambda_{\text{pf}}^k$  as  $k \rightarrow \infty$
- alternative characterization:  $\lambda_{\text{pf}}(A) = \inf\{\lambda \mid Av \preceq \lambda v \text{ for some } v \succ 0\}$

## minimizing spectral radius of matrix of posynomials

- minimize  $\lambda_{\text{pf}}(A(x))$ , where the elements  $A(x)_{ij}$  are posynomials of  $x$
- equivalent geometric program:

$$\begin{array}{ll} \text{minimize} & \lambda \\ \text{subject to} & \sum_{j=1}^n A(x)_{ij} v_j / (\lambda v_i) \leq 1, \quad i = 1, \dots, n \end{array}$$

variables  $\lambda, v, x$



$$x \in K \quad x - y \in K$$

$$\underline{x \succeq_K 0} \quad \underline{x \succeq_K y}$$

## Generalized inequality constraints

convex problem with generalized inequality constraints

$$\left[ \begin{array}{ll} \text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \preceq_{K_i} 0, \quad i = 1, \dots, m \\ & Ax = b \end{array} \right. \quad \begin{array}{l} f_1(x) \preceq_{K_1} 0 \\ \vdots \\ f_m(x) \preceq_{K_m} 0 \end{array}$$

- $f_0 : \mathbf{R}^n \rightarrow \mathbf{R}$  convex;  $\underline{f_i : \mathbf{R}^n \rightarrow \mathbf{R}^{k_i}}$   $K_i$ -convex w.r.t. proper cone  $K_i$
- same properties as standard convex problem (convex feasible set, local optimum is global, etc.)

conic form problem: special case with affine objective and constraints

*conic-linear*

$$\left[ \begin{array}{ll} \text{minimize} & \underline{c^T x} \\ \text{subject to} & \underline{\frac{Fx + g \preceq_K 0}{Ax = b}} \end{array} \right. \quad \left\{ \begin{array}{l} K = \mathbf{R}_+^m \\ K = S_+^m \end{array} \right.$$

extends linear programming ( $K = \mathbf{R}_+^m$ ) to nonpolyhedral cones

## Semidefinite program (SDP)

$$\rightarrow \begin{cases} \text{minimize} & \underline{c}^T \underline{x} \\ \text{subject to} & \underline{x}_1 \underline{F}_1 + \underline{x}_2 \underline{F}_2 + \cdots + \underline{x}_n \underline{F}_n + \underline{G} \preceq 0 \\ & \underline{Ax} = \underline{b} \end{cases}$$

negative semidefinite

with  $\underline{F}_i, \underline{G} \in \underline{\mathbf{S}}^k$

- inequality constraint is called linear matrix inequality (LMI)
- includes problems with multiple LMI constraints: for example,

$$\rightarrow x_1 \hat{F}_1 + \cdots + x_n \hat{F}_n + \hat{G} \preceq 0, \quad x_1 \tilde{F}_1 + \cdots + x_n \tilde{F}_n + \tilde{G} \preceq 0$$

is equivalent to single LMI

$$\begin{bmatrix} A & 0 \\ 0 & B \end{bmatrix} \succeq 0 \Leftrightarrow A \succeq 0 \ \& \ B \succeq 0$$

$$\rightarrow x_1 \begin{bmatrix} \hat{F}_1 & 0 \\ 0 & \underline{\tilde{F}_1} \end{bmatrix} + x_2 \begin{bmatrix} \hat{F}_2 & 0 \\ 0 & \underline{\tilde{F}_2} \end{bmatrix} + \cdots + x_n \begin{bmatrix} \hat{F}_n & 0 \\ 0 & \underline{\tilde{F}_n} \end{bmatrix} + \begin{bmatrix} \hat{G} & 0 \\ 0 & \underline{\tilde{G}} \end{bmatrix} \preceq 0$$


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$$A(x) = A_0 + \underbrace{x_1}_{\in \mathbb{R}} A_1 + \underbrace{x_2}_{\in \mathbb{R}} A_2 + \dots + \underbrace{x_n}_{\in \mathbb{R}} A_n$$

$$x \in \mathbb{R}^n$$

$$A(x): \mathbb{R}^n \rightarrow S^n$$

$$A(x) \preceq 0$$

$$S^n_+$$

$$\begin{bmatrix} 1-x_1 & x_2 & 0 \\ x_2 & 5 & -x_3+1 \\ 0 & -x_3+1 & 0 \end{bmatrix} \preceq 0 \Leftrightarrow \underbrace{\begin{bmatrix} 1 & 0 & 0 \\ 0 & 5 & 1 \\ 0 & 1 & 0 \end{bmatrix}}_{A_0} + \underbrace{\begin{bmatrix} -1 & 0 \\ 0 & 0 \end{bmatrix}}_{A_1} x_1 + \underbrace{\begin{bmatrix} & & \\ & & \\ & & \end{bmatrix}}_{A_2} x_2 + \begin{bmatrix} & & \\ & & \\ & & \end{bmatrix} x_3$$

$$x \equiv \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix}$$

# LP and SOCP as SDP

## LP and equivalent SDP

$$\text{LP: } \begin{cases} \text{minimize} & c^T x \\ \text{subject to} & Ax \preceq_{\mathbb{R}_+^n} b \end{cases}$$

$$\text{SDP: } \begin{cases} \text{minimize} & c^T x \\ \text{subject to} & \text{diag}(Ax - b) \preceq_{\mathbb{S}_+^n} 0 \end{cases}$$

(note different interpretation of generalized inequality  $\preceq$ )

$$\begin{bmatrix} a_1^T x - b_1 & & 0 \\ & \ddots & \\ 0 & & a_m^T x - b_m \end{bmatrix} \preceq_{\mathbb{S}_+^n} 0$$

## SOCP and equivalent SDP

$$\text{SOCP: } \begin{cases} \text{minimize} & f^T x \\ \text{subject to} & \|A_i x + b_i\|_2 \leq c_i^T x + d_i, \quad i = 1, \dots, m \end{cases}$$

$$\text{SDP: } \begin{cases} \text{minimize} & f^T x \\ \text{subject to} & \begin{bmatrix} (c_i^T x + d_i)I & A_i x + b_i \\ (A_i x + b_i)^T & c_i^T x + d_i \end{bmatrix} \succeq 0, \quad i = 1, \dots, m \end{cases}$$

$$\rightarrow \|A_i x + b_i\|_2^2 \leq (c_i^T x + d_i)^2$$

assume:  $c_i^T x + d_i > 0, \forall i$

$$\rightarrow \frac{(A_i x + b_i)^T (A_i x + b_i)}{(c_i^T x + d_i)^2} \leq \frac{(c_i^T x + d_i)^2}{(c_i^T x + d_i)^2}$$

Schur complements lemma,

$$\begin{bmatrix} A & B \\ B^T & C \end{bmatrix} \succeq 0, A \succ 0, C \succeq 0$$

$$\Leftrightarrow \underbrace{C - B^T A^{-1} B}_{\text{Schur complement}} \succeq 0$$

$$- \underbrace{(A_i x + b_i)^T}_{\text{green}} \underbrace{\left( \frac{1}{(c_i^T x + d_i)^2} \right)}_{\text{blue}} \underbrace{(A_i x + b_i)}_{\text{green}} + \underbrace{1}_{\text{orange}} \succeq 0$$

$$\Leftrightarrow \begin{matrix} 1 \uparrow \\ n \uparrow \end{matrix} \begin{bmatrix} \underbrace{(c_i^T x + d_i)^2}_{\text{orange}} & \overset{n}{\underbrace{(A_i x + b_i)^T}_{\text{green}}} \\ \underbrace{(A_i x + b_i)}_{\text{green}} & \underset{n \times n}{I} \end{bmatrix} \succeq 0 \rightarrow \text{not an LMI}$$

$$- (A_i x + b_i)^T \frac{1}{c_i^T x + d_i} (A_i x + b_i) + (c_i^T x + d_i) \succeq 0$$

$$\begin{matrix} 1 \uparrow \\ n \uparrow \end{matrix} \begin{bmatrix} (c_i^T x + d_i) & \underbrace{(A_i x + b_i)^T}_{\text{green}} \\ \underbrace{(A_i x + b_i)}_{\text{green}} & (c_i^T x + d_i) \underset{n \times n}{I} \end{bmatrix} \succeq 0 \Rightarrow \text{LMI} \checkmark$$

## Eigenvalue minimization

$$\underset{x}{\text{minimize}} \quad \lambda_{\max}(A(x))$$

$$\left[ \begin{array}{l} \min. \quad t \\ x, t \\ \lambda_{\max}(A(x)) \leq t \end{array} \right.$$

where  $A(x) = \underline{A_0} + \underline{x_1}A_1 + \cdots + \underline{x_n}A_n$  (with given  $A_i \in \mathbf{S}^k$ )

equivalent SDP

$$\left[ \begin{array}{l} \underset{x, t}{\text{minimize}} \quad t \\ \text{subject to} \quad \underline{A(x)} \preceq \underline{tI} \end{array} \right.$$

- variables  $x \in \mathbf{R}^n, t \in \mathbf{R}$
- follows from

$$\underline{\lambda_{\max}(A) \leq t} \iff \underline{A \preceq tI}$$

$$A - tI \leq 0 \iff$$

$$\lambda_i(A - tI) \leq 0, \forall i \iff$$

$$\lambda_i(A) - t \leq 0, \forall i \iff$$

$$\lambda_{\max}(A) \leq t$$

## Matrix norm minimization

$$\rightarrow \text{minimize } \underbrace{\|A(x)\|_2}_{\sigma_{\max}(A(x))} = \left( \lambda_{\max}(\underbrace{A(x)^T A(x)}) \right)^{1/2}$$

where  $\underline{A(x)} = A_0 + x_1 A_1 + \cdots + x_n A_n$  (with given  $\underline{A_i} \in \mathbf{R}^{p \times q}$ )  
equivalent SDP

$$\left[ \begin{array}{l} \text{minimize } \underline{t} \\ \text{subject to } \left[ \begin{array}{cc} \underline{tI} & \underline{A(x)} \\ \underline{A(x)^T} & \underline{tI} \end{array} \right] \succeq 0 \end{array} \right.$$

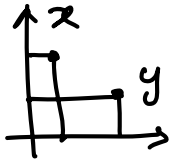
- variables  $x \in \mathbf{R}^n$ ,  $t \in \mathbf{R}$
- constraint follows from

$$\begin{aligned} \lambda_{\max}(A^T A) &\leq t^2 \\ \underline{\|A\|_2 \leq t} &\iff \underline{A^T A \leq t^2 I}, \quad \underline{t \geq 0} \\ &\iff \left[ \begin{array}{cc} \underline{tI} & \underline{A} \\ \underline{A^T} & \underline{tI} \end{array} \right] \succeq 0 \end{aligned}$$

- LP — QP — QCQP — SOCP — SDP
- GP

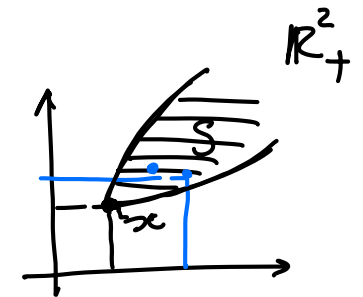
# Minimum and minimal elements of a set

$\preceq_K$  is not in general a linear ordering: we can have  $x \not\preceq_K y$  and  $y \not\preceq_K x$



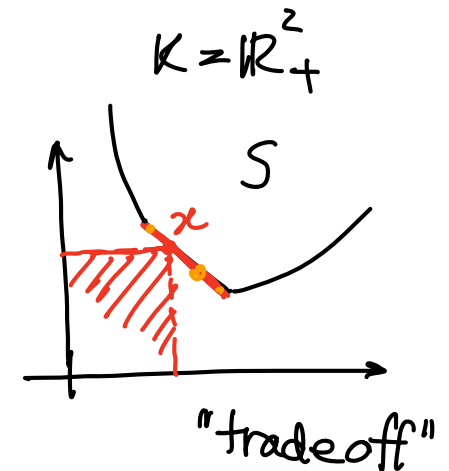
- $x \in S$  is the minimum element of  $S$  with respect to  $\preceq_K$  if

$$\underline{y \in S \implies x \preceq_K y}$$



- $x \in S$  is a minimal element of  $S$  with respect to  $\preceq_K$  if

$$\underline{y \in S, \quad y \preceq_K x \implies y = x}$$





# Multiobjective (vector) optimization

general vector optimization problem

$$\begin{array}{ll} \text{minimize (w.r.t. } K) & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_i(x) \leq 0, \quad i = 1, \dots, p \end{array}$$

- $K = \mathbb{R}_+^q$  in all our multi-obj problems in this course!

vector objective  $f_0 : \mathbf{R}^n \rightarrow \mathbf{R}^q$ , minimized w.r.t. proper cone  $K \in \mathbf{R}^q$

convex vector optimization problem

$$\left[ \begin{array}{ll} \text{minimize (w.r.t. } K) & \underline{f_0(x)} \\ \text{subject to} & \underline{f_i(x)} \leq 0, \quad i = 1, \dots, m \\ & \underline{Ax = b} \end{array} \right.$$

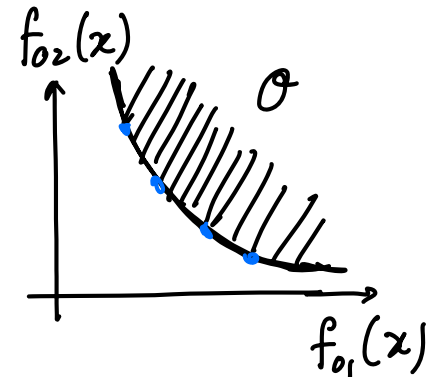
$$x \longrightarrow \begin{bmatrix} f_0(x) \\ \vdots \\ f_q(x) \end{bmatrix} \in \mathbb{R}^q$$

with  $f_0$   $K$ -convex,  $f_1, \dots, f_m$  convex

# Optimal and Pareto optimal points

set of achievable objective values

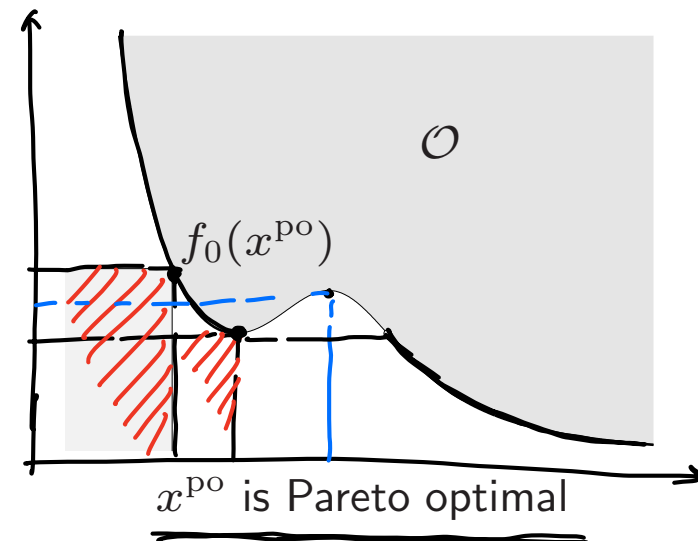
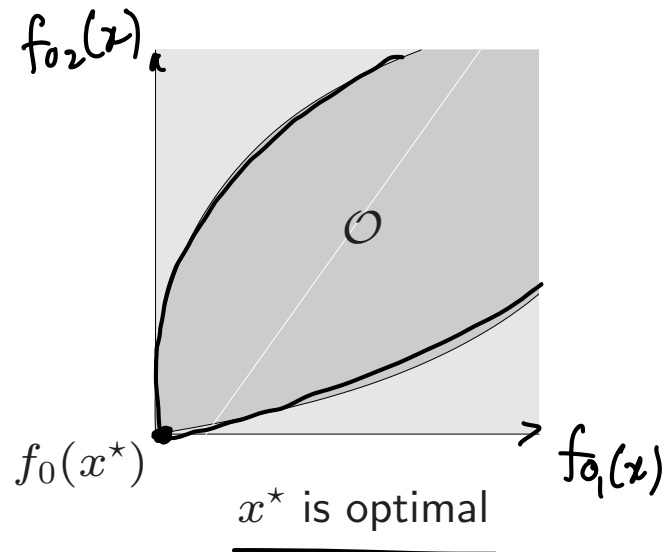
$$f_0(x) = \begin{bmatrix} f_{01}(x) \\ f_{02}(x) \end{bmatrix}$$



$$\underline{\underline{O}} = \{ \underline{f_0(x)} \mid \underline{x \text{ feasible}} \}$$

•  $O$  is a convex set for a *conv* opt prob.

- feasible  $x$  is **optimal** if  $f_0(x)$  is a minimum value of  $O$
- feasible  $x$  is **Pareto optimal** if  $f_0(x)$  is a minimal value of  $O$



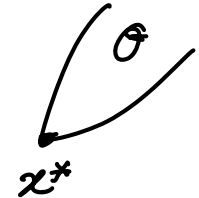
# Multiobjective optimization

vector optimization problem with  $K = \mathbf{R}_+^q$

$$f_0(x) = (\underbrace{F_1(x)}, \dots, \underbrace{F_q(x)})$$

- $q$  different objectives  $F_i$ ; roughly speaking we want all  $F_i$ 's to be small
- feasible  $x^*$  is optimal if

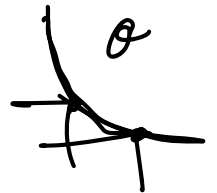
$$y \text{ feasible} \implies \underline{f_0(x^*) \preceq f_0(y)}$$



if there exists an optimal point, the objectives are noncompeting

- feasible  $x^{p^o}$  is Pareto optimal if

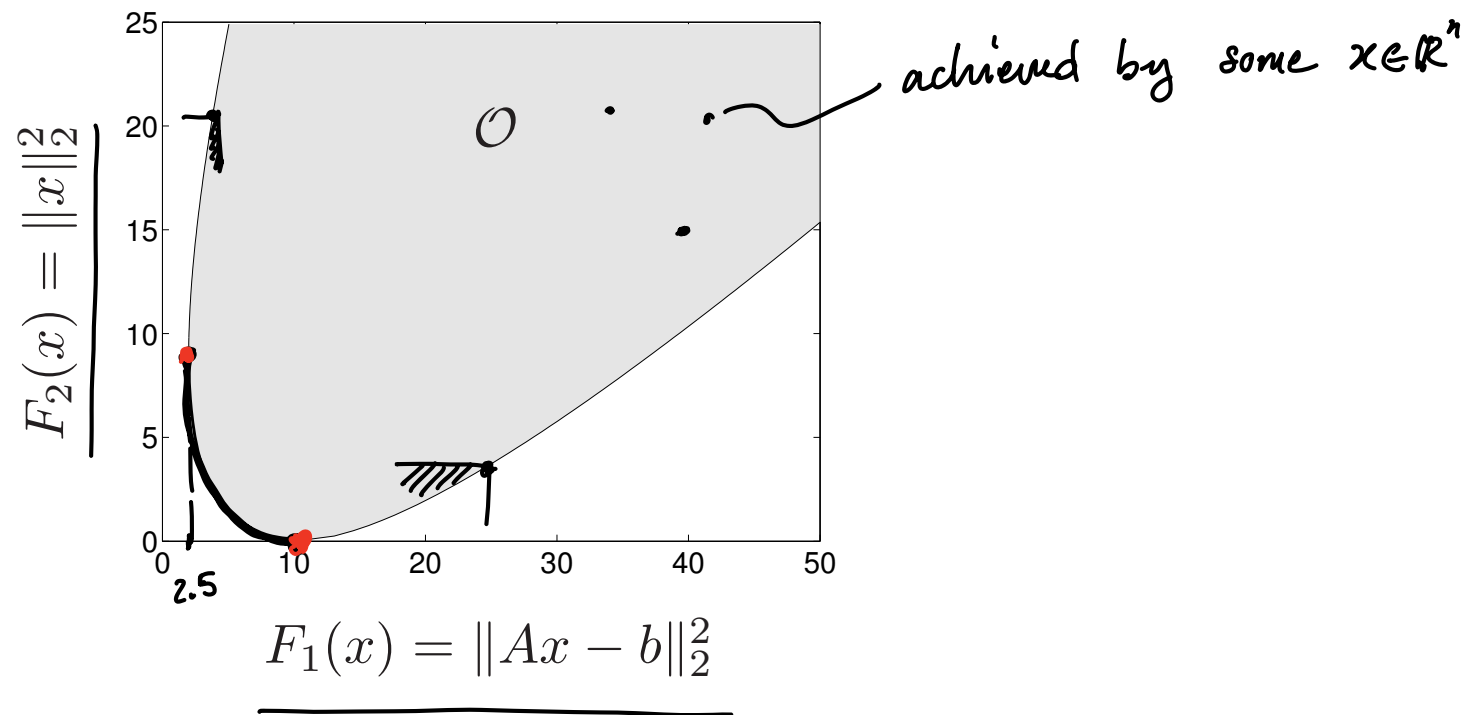
$$y \text{ feasible, } f_0(y) \preceq f_0(x^{p^o}) \implies \underline{f_0(x^{p^o}) = f_0(y)}$$



if there are multiple Pareto optimal values, there is a trade-off between the objectives

# Regularized least-squares

$$\underset{x \in \mathbb{R}^n}{\text{minimize (w.r.t. } \mathbf{R}_+^2)} \quad (\underbrace{\|Ax - b\|_2^2}_{F_1(x)}, \underbrace{\|x\|_2^2}_{F_2(x)})$$



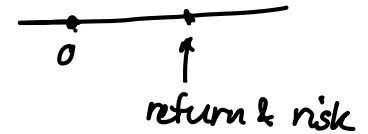
example for  $A \in \mathbf{R}^{100 \times 10}$ ; heavy line is formed by Pareto optimal points

$$A_{m \times n}, n=10, b \in \mathbb{R}^m, m=100 \quad \|b\|_2^2=10 \quad \min_x \|Ax - b\|_2^2 \quad ? \sim 2.5$$

1950's Markowitz portfolio opt.

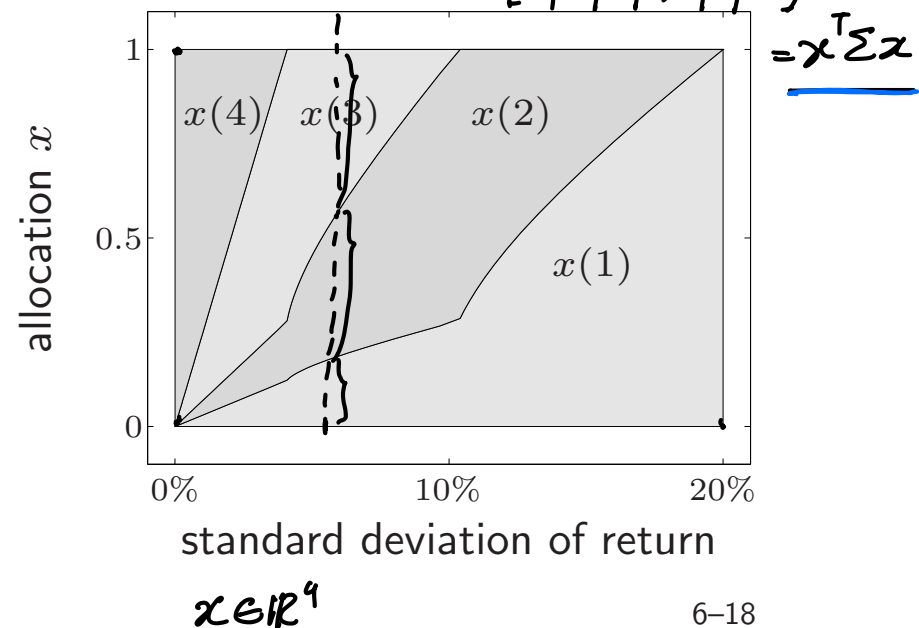
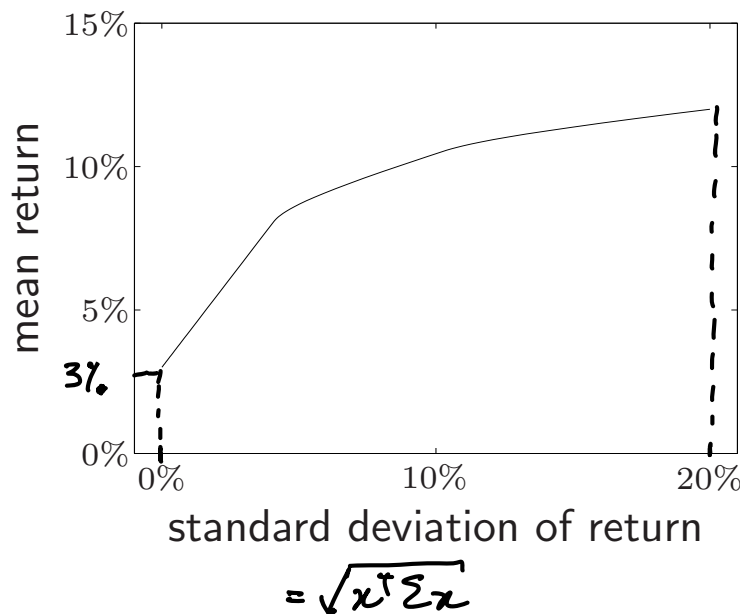
## Risk return trade-off in portfolio optimization

$$\begin{cases} \text{minimize (w.r.t. } \mathbf{R}_+^2) & (-\bar{p}^T x, x^T \Sigma x) \\ \text{subject to} & \mathbf{1}^T x = 1, \quad x \succeq 0 \end{cases}$$



- $x \in \mathbf{R}^n$  is investment portfolio;  $x_i$  is fraction invested in asset  $i$
- $p \in \mathbf{R}^n$  is vector of relative asset price changes; modeled as a random variable with mean  $\bar{p}$ , covariance  $\Sigma$   $p \sim \mathcal{D}(\bar{p}, \Sigma)$
- $\bar{p}^T x = \mathbf{E} r$  is expected return;  $x^T \Sigma x = \text{var } r$  is return variance  $x^T (p - \bar{p})$   
 $\text{return} = \sum p_i \pi_i = p^T x$   $\mathbb{E}_p(p^T x) = \bar{p}^T x$   $\mathbb{E}_p(p^T x) = \mathbb{E}_p(p^T x - \bar{p}^T x + \bar{p}^T x)$   
 $= x^T [\mathbb{E}_p(p - \bar{p})(p - \bar{p})^T] x$   
 $= x^T \Sigma x$

example



# Scalarization for multicriterion problems

to find Pareto optimal points, minimize positive weighted sum

$$\underline{\lambda^T f_0(x)} = \underline{\lambda_1 F_1(x)} + \cdots + \underline{\lambda_q F_q(x)}$$

## examples

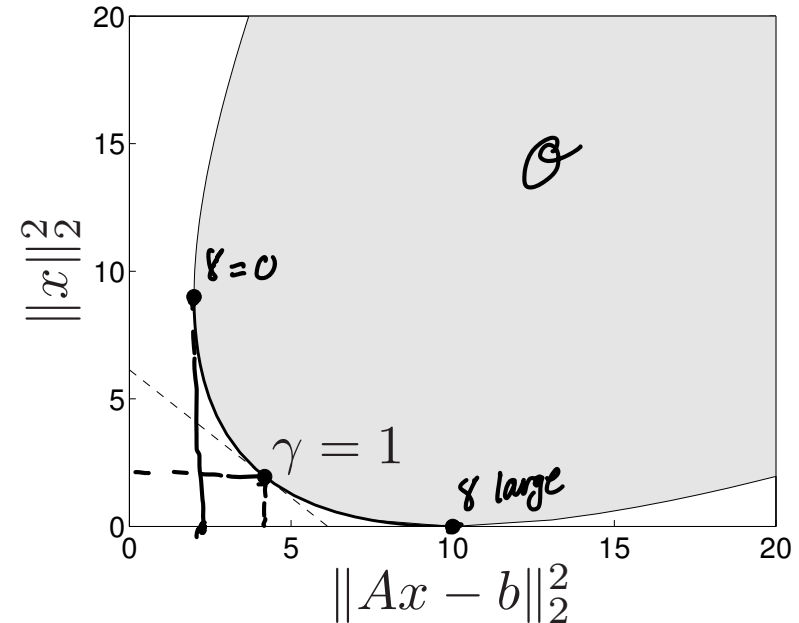
- regularized least-squares problem of page 6–17

take  $\lambda = (1, \gamma)$  with  $\gamma > 0$

$$\left[ \begin{array}{l} \text{minimize} \\ x \end{array} \quad \|Ax - b\|_2^2 + \gamma \|x\|_2^2 \right]$$

for fixed  $\gamma$ , a LS problem

$$\left\| \underbrace{\begin{bmatrix} A \\ \sqrt{\gamma} I \end{bmatrix}}_{\tilde{A}} x - \underbrace{\begin{bmatrix} b \\ 0 \end{bmatrix}}_{\tilde{b}} \right\|_2^2 \quad x^* = \tilde{A}^+ \tilde{b} = (A^T A + \gamma I)^{-1} A^T b$$



- risk-return trade-off of page 6–18

$$\begin{cases} \underset{x}{\text{minimize}} & -\bar{p}^T x + \gamma x^T \Sigma x \\ \text{subject to} & \mathbf{1}^T x = 1, \quad x \succeq 0 \end{cases}$$

for fixed  $\gamma > 0$ , a quadratic program

*in practice, discretize  $\gamma$ , solve one QP for each  $\gamma$*