

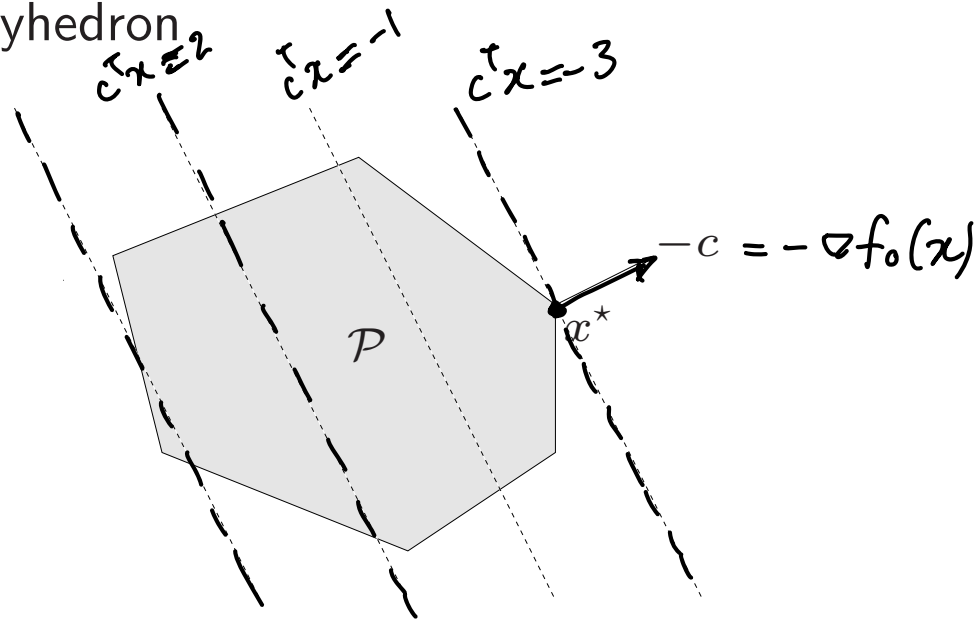
5. Linear and quadratic programming

- linear programming
- quadratic programming
- second-order cone programming

Linear program (LP)

$$\left\{ \begin{array}{ll} \text{minimize} & \widetilde{c^T x + d} \\ \text{subject to} & Gx \preceq h \\ & Ax = b \end{array} \right.$$

- convex problem with affine objective and constraint functions
- feasible set is a polyhedron



Examples

diet problem: choose quantities $\underline{x_1}, \dots, \underline{x_n}$ of \underline{n} foods

- one unit of food \underline{j} costs $\underline{c_j}$, contains amount $\underline{a_{ij}}$ of nutrient i
- healthy diet requires nutrient i in quantity at least b_i

to find cheapest healthy diet,

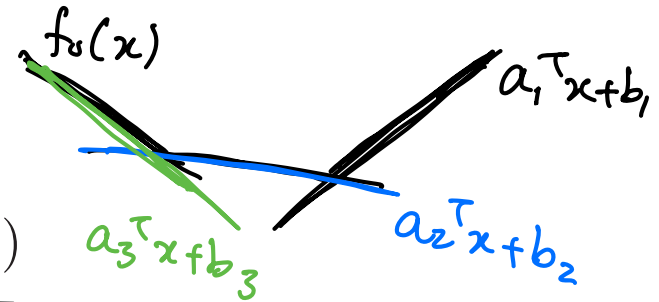
$$\left\{ \begin{array}{ll} \text{minimize} & c^T x = \sum c_j x_j \\ \text{subject to} & \underbrace{Ax \succeq b}_{a_i^T x \geq b_i}, \quad \underbrace{x \succeq 0} \end{array} \right.$$

one of the earliest applications of LP during and after World War II:
optimal diet for troops

{ (see G. Dantzig's interview for the whole history:
<https://www.informs.org/Explore/History-of-O.R.-Excellence/Oral-Histories/George-Dantzig>)

(convex) piecewise-linear minimization

$$\begin{cases} \text{minimize} & \max_{i=1, \dots, m} (a_i^T x + b_i) \end{cases}$$



equivalent to an LP (epigraph form):

$$\begin{cases} \text{minimize} & t \\ \text{subject to} & a_i^T x + b_i \leq t, \quad i = 1, \dots, m \end{cases}$$

$$f_0(x) = \max_i (a_i^T x + b_i)$$

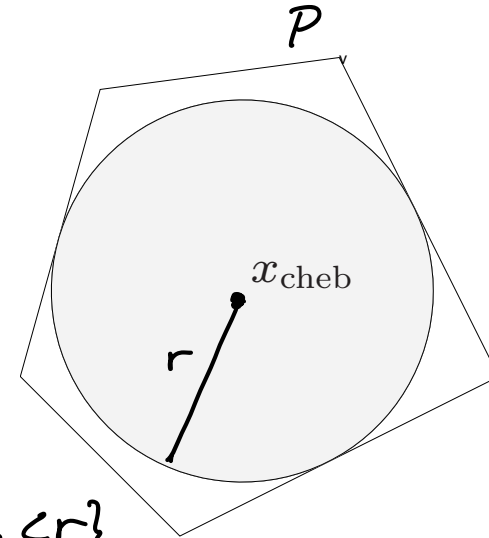
$$\begin{cases} \min & t \\ x \in \mathbb{R}^n, t \in \mathbb{R} \\ \text{s.t.} & \max_{i=1, \dots, m} (a_i^T x + b_i) \leq t \end{cases} \Leftrightarrow a_i^T x + b_i \leq t, \quad i=1, \dots, m$$

Chebyshev center of a polyhedron

Chebyshev center of

$$\mathcal{P} = \{x \mid \underline{a_i^T x \leq b_i}, \quad i = 1, \dots, \underline{m}\}$$

is center of largest inscribed ball



$$\mathcal{B} = \{\underline{x_c} + u \mid \underline{\|u\|_2 \leq r}\} = \{x \mid \|x - x_c\|_2 \leq r\}$$

- $a_i^T x \leq b_i$ for all $x \in \mathcal{B}$ if and only if

$$\rightarrow \sup\{a_i^T(x_c + u) \mid \|u\|_2 \leq r\} = a_i^T x_c + r\|a_i\|_2 \leq b_i$$

- hence, x_c, r can be determined by solving the LP

$$\begin{cases} \text{maximize} & r \\ \text{subject to} & \underline{a_i^T x_c} + \underline{r\|a_i\|_2} \leq b_i, \quad i = 1, \dots, m \end{cases}$$

$$a_i^T x \leq b_i \quad \text{for all } x \in \{ \underbrace{x_c + u} \mid \|u\|_2 \leq r \}$$

$$\sup_{\|u\|_2 \leq r} a_i^T(x_c + u) \leq b_i$$

$$\sup_{\|u\|_2 \leq r} (\overbrace{a_i^T x_c + a_i^T u}) \leq b_i$$

$$a_i^T x_c + \underbrace{\sup_{\|u\|_2 \leq r} a_i^T u}_{r \|a_i\|_2} \leq b_i \quad i=1, \dots, m$$

$$a_i^T \underline{x_c} + \|a_i\|_2 \underline{r} \leq b_i \quad i=1, \dots, m$$

$$|a_i^T u| \leq \|a_i\|_2 \|u\|_2$$

↑
eg. if u in direction
of a_i

Announcements

- Check Natalia's announcements about her TA review & OH on Zoom (due to travel out of country) — back on Feb 14.
- HW5 due this Wed — note this HW is a bit harder than previous, ask Q's if you need help!
- Short HW6 will be assigned Wed & due next Monday (midnight)
- Midterm Fri Feb 16, in-class, closed book/notes, 2 letter-sized sheets double-sided

More on LP

LP:

$$\left[\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Gx \preceq h \\ & \underline{Ax = b} \end{array} \right. \leftarrow$$

'standard' form LP:

$$\left[\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax = b \\ & x \succeq 0 \end{array} \right.$$

converting to 'standard' form:

- inequality constraints: write $\underline{a_i^T x \leq b_i}$ as $Ax + s = b$
 $\underline{a_i^T x + s_i = b_i}, \quad \underline{s_i \geq 0}$

$$A(x^+ - x^-) + s = b$$

$$[A \quad -A \quad I] \begin{bmatrix} x^+ \\ x^- \\ s \end{bmatrix} = b$$

$$\underbrace{\begin{bmatrix} x^+ \\ x^- \\ s \end{bmatrix}}_{\geq 0}$$

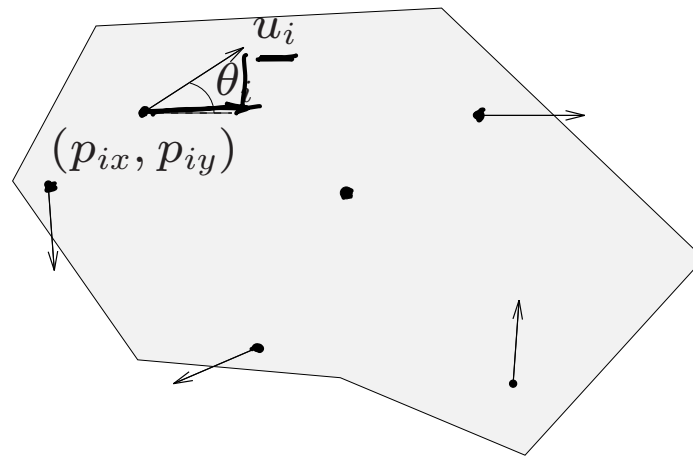
s_i is *slack variable* associated with $a_i^T x \leq b_i$
- unconstrained variables: write $x_i \in \mathbf{R}$ as

$$\underline{x_i} = \underline{x_i^+} - \underline{x_i^-}, \quad \underline{x_i^+}, \underline{x_i^-} \geq 0$$

Example

force/moment generation with thrusters

- rigid body with center of mass at origin $p = 0 \in \mathbf{R}^2$
- n forces with magnitude $\underline{u_i}$, acting at $p_i = (\underline{p_{ix}}, \underline{p_{iy}})$, in direction θ_i



- resulting horizontal force: $\underline{F_x} = \sum_{i=1}^n \underline{u_i} \cos \theta_i$
- resulting vertical force: $\underline{F_y} = \sum_{i=1}^n \underline{u_i} \sin \theta_i$
- resulting torque: $\underline{T} = \sum_{i=1}^n (\underline{p_{iy}} \underline{u_i} \cos \theta_i - \underline{p_{ix}} \underline{u_i} \sin \theta_i)$

- force limits: $0 \leq u_i \leq 1$ (thrusters)
- fuel usage: $\underline{u_1 + \cdots + u_n}$

problem: find thruster forces $\underline{u_i}$ that yield given desired forces and torques and minimize fuel usage (if feasible)

can be expressed as LP:

$$\left[\begin{array}{ll} \text{minimize} & \mathbf{1}^T u \\ \text{subject to} & Fu = f^{\text{des}} \\ & 0 \leq \underline{u_i} \leq 1, \ i = 1, \dots, n \end{array} \right.$$

where

$$F = \begin{bmatrix} \cos \theta_1 & \cdots & \cos \theta_n \\ \sin \theta_1 & \cdots & \sin \theta_n \\ p_{1y} \cos \theta_1 - p_{1x} \sin \theta_1 & \cdots & p_{ny} \cos \theta_n - p_{nx} \sin \theta_n \end{bmatrix}_{3 \times n},$$

$$f^{\text{des}} = (F_x^{\text{des}}, F_y^{\text{des}}, T^{\text{des}}), \quad \mathbf{1} = (1, 1, \cdots 1)$$

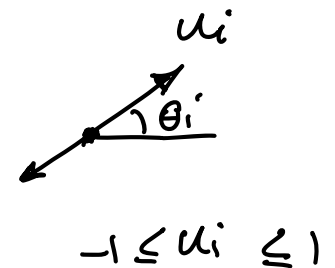
thruster problem in 'standard' form

$$\begin{aligned} & \text{minimize} && \begin{bmatrix} \mathbf{1}^T & 0 \end{bmatrix} \begin{bmatrix} u \\ s \end{bmatrix} \\ & \text{subject to} && \begin{bmatrix} u \\ s \end{bmatrix} \succeq 0 \\ & && \begin{bmatrix} F & 0 \\ I & I \end{bmatrix} \begin{bmatrix} u \\ s \end{bmatrix} = \begin{bmatrix} f^{\text{des}} \\ \mathbf{1} \end{bmatrix} \end{aligned}$$

some extensions (can express these as LP):

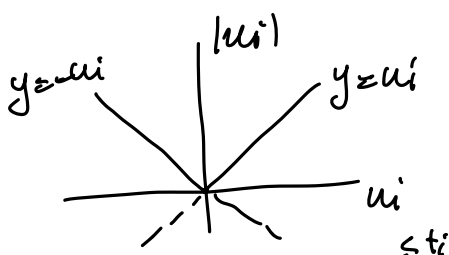
- opposing thruster pairs

$$\left\{ \begin{array}{ll} \text{minimize} & \sum_i |u_i| \\ \text{subject to} & Fu = f^{\text{des}} \\ & |u_i| \leq 1, \quad i = 1, \dots, n \end{array} \right.$$



- given f^{des} ,

$$\begin{aligned} & \text{minimize} && \|Fu - f^{\text{des}}\|_{\infty} \quad \nearrow \text{Chebyshev error} \\ & \text{subject to} && 0 \leq u_i \leq 1, \quad i = 1, \dots, n \end{aligned}$$



$$|u_i| = \max \{u_i, -u_i\}$$

$$\begin{cases} \min. & \sum_{i=1}^n |u_i| \\ u \in \mathbb{R}^n & \\ \text{s.t.} & Fu = f_{des} \\ & |u_i| \leq 1 \end{cases}$$

epigraph form:

epigraph variables $t_i \quad i=1, \dots, n$

$$\begin{cases} \min. & \sum t_i \\ \underline{u} \in \mathbb{R}^n & \\ \underline{t} \in \mathbb{R}^n & \\ \text{s.t.} & \underline{F}\underline{u} = f_{des} \\ & |u_i| \leq t_i \rightarrow -\underline{t}_i \leq \underline{u}_i \leq \underline{t}_i \\ & |u_i| \leq 1 \rightarrow -1 \leq \underline{u}_i \leq 1 \end{cases}$$

$$F = \begin{bmatrix} f_1^T \\ f_2^T \\ f_3^T \end{bmatrix}$$

• Chebyshev error:

$$\begin{cases} \min. & \|Fu - f_{des}\|_\infty \rightarrow \max_i |f_i^T u - f_{des,i}| \\ u & \\ \text{s.t.} & |u_i| \leq 1 \end{cases}$$

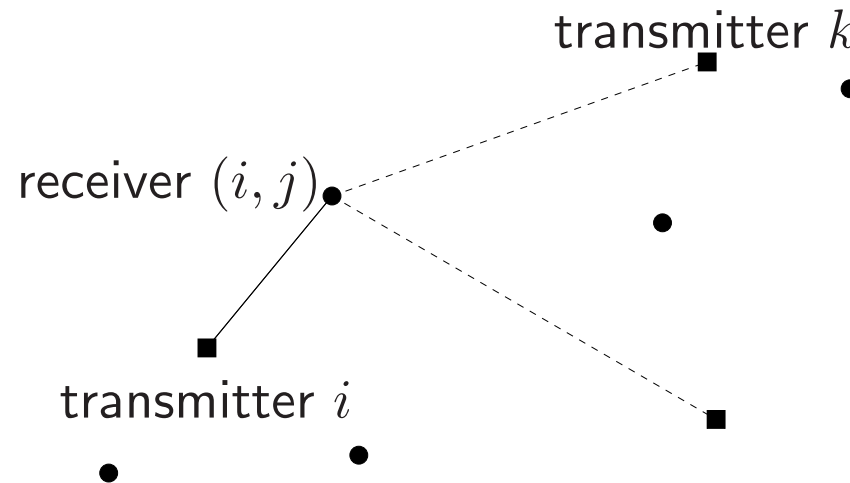
epigraph form:

$$\begin{cases} \min. & t \\ u \in \mathbb{R}^n & \\ t \in \mathbb{R} & \\ -t \leq f_i^T u - f_{des,i} \leq t & i=1,2,3 \\ & -1 \leq u_i \leq 1 \end{cases}$$

Optimal transmitter power allocation

(generalized) linear fractional program (refer to section 4.2.3[?])

- m transmitters, mn receivers all at same frequency
- transmitter i wants to transmit to n receivers labeled (i, j) , $j = 1, \dots, n$



- A_{ijk} is path gain from transmitter k to receiver (i, j)
- N_{ij} is (self) noise power of receiver (i, j)
- variables: transmitter powers p_k , $k = 1, \dots, m$

at receiver (i, j) :

- signal power: $S_{ij} = A_{iji}p_i$
- noise plus interference power: $I_{ij} = \sum_{k \neq i} A_{ijk}p_k + N_{ij}$
- signal to interference/noise ratio (SINR): S_{ij}/I_{ij}

problem: choose p_i to maximize smallest SINR:

$$\begin{array}{ll} \text{maximize} & \min_{i,j} \frac{A_{iji}p_i}{\sum_{k \neq i} A_{ijk}p_k + N_{ij}} \\ \text{subject to} & 0 \leq p_i \leq p_{\max} \end{array}$$

... a (generalized) linear fractional program

Minimum-time optimal control

$$x(t) \in \mathbb{R}^n, u(t) \in \mathbb{R}^m$$

$$\underline{x(t+1) = Ax(t) + Bu(t)}, \quad \text{with } \underline{x(0) = x_0} \quad \text{and} \quad \underline{\|u(t)\|_\infty \leq 1, t = 0, 1, \dots, K}$$

variables: $\underline{u(0)}, \dots, \underline{u(K)}$

settling time $f(u(0), \dots, u(K))$ is

$$\rightarrow \inf \{ \underline{T} \mid \underline{x(t) = 0 \text{ for } T \leq t \leq K+1} \}$$

→ set of $K+1-T$ linear equation in $u(t)$'s

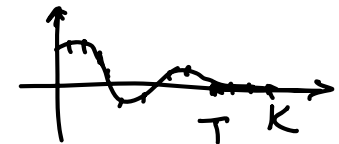
$$x(1) = Ax(0) + Bu(0)$$

$$x(2) = Ax(1) + Bu(1) = A(Ax(0) + Bu(0)) + Bu(1)$$

$$= A^2x(0) + AB\underline{u(0)} + B\underline{u(1)}$$

⋮

$x(t)$ depends linearly on all $u(t)$'s



Minimum-time optimal control

$$x(t+1) = Ax(t) + Bu(t), \quad \text{with } x(0) = x_0 \quad \text{and} \quad \underline{\|u(t)\|_\infty \leq 1, \quad t = 0, 1, \dots, K}$$

variables: $u(0), \dots, u(K)$

settling time $f(u(0), \dots, u(K))$ is

$$\inf \{T \mid x(t) = 0 \text{ for } T \leq t \leq K+1\}$$

f is quasiconvex function of $(u(0), \dots, u(K))$: $S_T = \{u(0), \dots, u(K) \mid f(u(0), \dots, u(K)) \leq T\}$

$f(u(0), u(1), \dots, u(K)) \leq T$ if and only if for all $t = \underline{T}, \dots, \underline{K+1}$ *intersection of*

$$x(t) = A^t x_0 + A^{t-1} \underbrace{Bu(0)} + \dots + \underbrace{Bu(t-1)} = 0 \quad \begin{array}{l} \text{hyperplanes} \Rightarrow \\ \text{convex set} \end{array}$$

min-time optimal control problem:

$\Rightarrow f$ is Q-cvx.

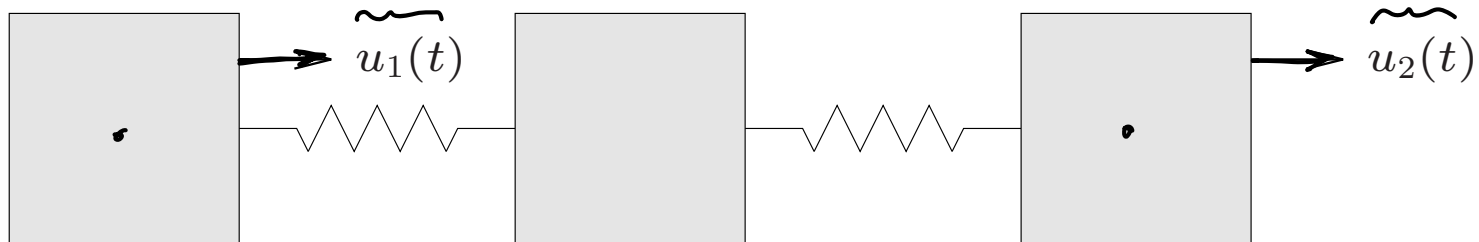
$$\left[\begin{array}{ll} \text{minimize} & f(u(0), u(1), \dots, u(K)) \\ \text{subject to} & \|u(t)\|_\infty \leq 1, \quad t = 0, \dots, K \end{array} \right.$$



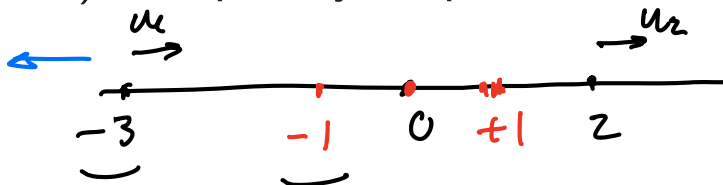
Min-time control example

three unit masses, connected by two unit springs with equilibrium length one

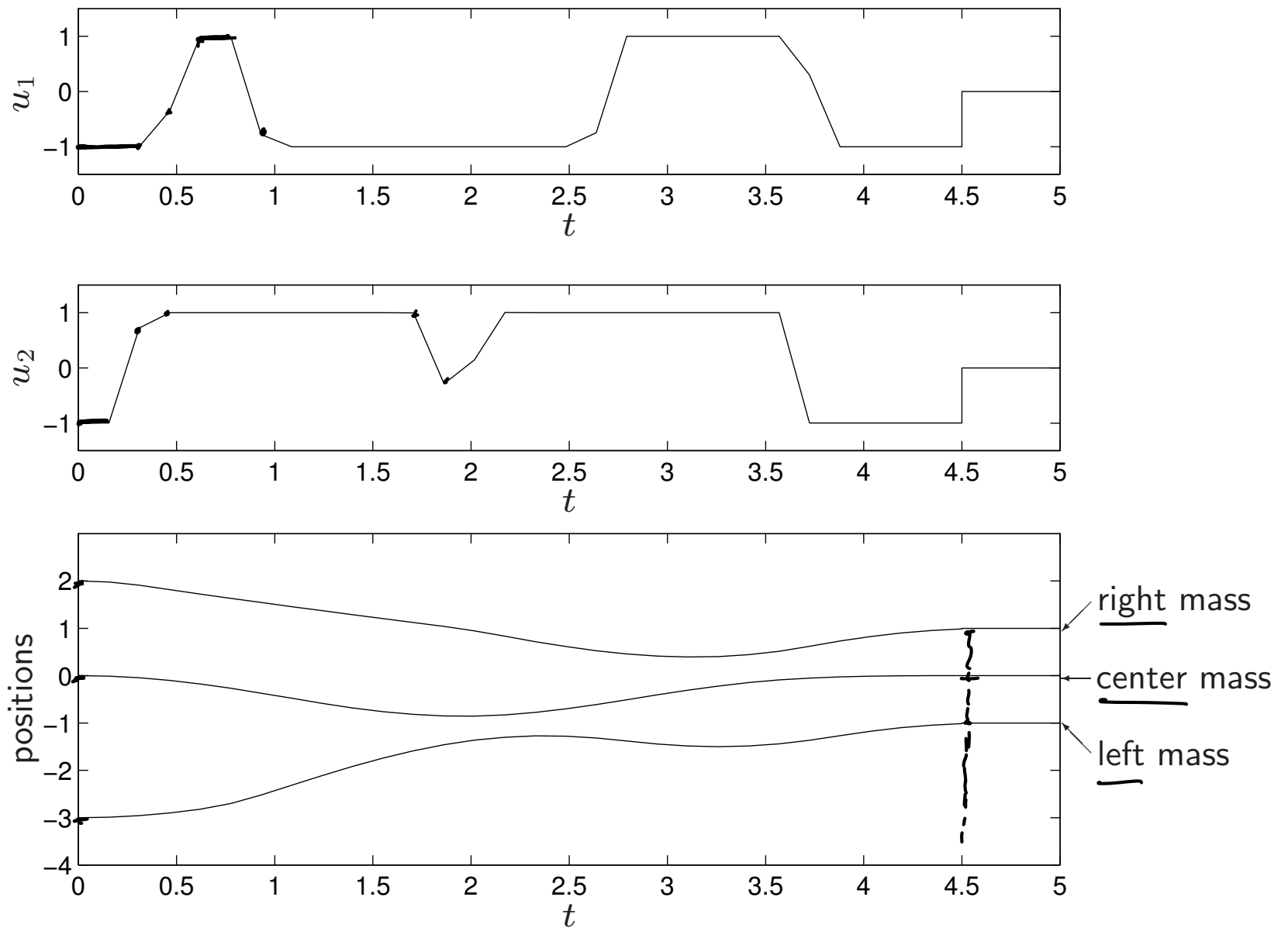
$u(t) \in \mathbf{R}^2$ is force on left and right mass over time interval $[0.15t, 0.15(t+1)]$



problem: pick $u(0), \dots, u(K)$ to bring masses to positions $(-1, 0, 1)$ (at rest), as quickly as possible, from initial condition $(-3, 0, 2)$ (at rest)



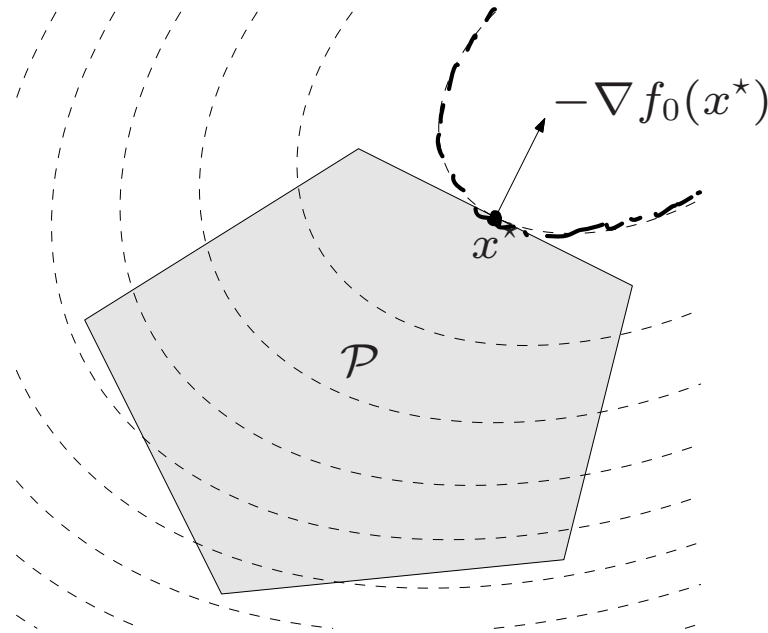
optimal solution:



Quadratic program (QP)

$$\left[\begin{array}{ll} \text{minimize} & \frac{(1/2)x^T P x + q^T x + r}{\text{subject to} \quad \frac{Gx \preceq h}{Ax = b}} \end{array} \right.$$

- $P \in \mathbf{S}_{+}^n$, so objective is convex quadratic
- minimize a convex quadratic function over a polyhedron



Examples

least-squares

$$\text{minimize} \quad \|Ax - b\|_2^2$$

- analytical solution $x^* = A^\dagger b$ (A^\dagger is pseudo-inverse)
- can add linear constraints, $e.g., l \preceq x \preceq u$

linear program with random cost

$$\begin{cases} \text{minimize} & \bar{c}^T x + \gamma x^T \Sigma x = \mathbf{E} \underbrace{c^T x}_{\text{risk}} + \gamma \underbrace{\text{var}(c^T x)} \\ \text{subject to} & Gx \preceq h, \quad Ax = \underline{b} \end{cases}$$

- c is random vector with mean \bar{c} and covariance Σ
- hence, $c^T x$ is random variable with mean $\bar{c}^T x$ and variance $x^T \Sigma x$
- $\gamma > 0$ is risk aversion parameter; controls the trade-off between expected cost and variance (risk)

$$\begin{cases} \min_x & \underline{c^T x} \\ \text{s.t.} & Gx \leq h \end{cases} \quad \text{problem data: } c, G, h$$

• in practice, uncertainty in data, e.g., uncertainty in $c \in \mathbb{R}^n$

$c \in \mathbb{R}^n$ is a random vector $c \sim D(\bar{c}, \underline{\Sigma}) \rightsquigarrow$ given

$c^T x$ is also random variable:

$$\bullet \mathbb{E}_c(c^T x) = \bar{c}^T x$$

$$\bullet \mathbb{E}_c[(c^T x - \bar{c}^T x)(c^T x - \bar{c}^T x)^T] = \mathbb{E}_c[x^T (c - \bar{c})(c - \bar{c})^T x] = x^T \underbrace{\mathbb{E}_c[(c - \bar{c})(c - \bar{c})^T]}_{\Sigma} x = x^T \Sigma x$$

$$\begin{cases} \min_x & \bar{c}^T x + \gamma \overbrace{x^T \Sigma x} \\ \text{s.t.} & Gx \leq h \end{cases} \rightarrow \text{a QP}$$

Quadratically constrained quadratic program (QCQP)

$$\left\{ \begin{array}{ll} \text{minimize} & (1/2)x^T P_0 x + q_0^T x + r_0 \\ \text{subject to} & \underbrace{(1/2)x^T P_i x + q_i^T x + r_i \leq 0, \quad i = 1, \dots, m}_{Ax = b} \end{array} \right.$$

- $P_i \in \mathbf{S}_{+}^n$; objective and constraints are convex quadratic
- if $P_1, \dots, P_m \in \mathbf{S}_{++}^n$, feasible region is intersection of m ellipsoids and an affine set

Second-order cone programming

$$x \in \mathbb{R}^n \quad \begin{cases} \text{minimize} & f^T x \\ \text{subject to} & \|A_i x + b_i\|_2 \leq \underline{c_i^T x + d_i}, \quad i = 1, \dots, m \\ & Fx = g \end{cases}$$

$\{(x, t) \mid \|x\|_2 \leq t, t > 0\}$
 $x \rightarrow \begin{bmatrix} A_i x + b_i \\ c_i^T x + d_i \end{bmatrix} \updownarrow n_i$

$$(A_i \in \mathbb{R}^{n_i \times n}, F \in \mathbb{R}^{p \times n})$$

- inequalities are called second-order cone (SOC) constraints:

$$(\underline{A_i x + b_i}, \underline{c_i^T x + d_i}) \in \text{second-order cone in } \mathbb{R}^{n_i+1}$$

- for $n_i = 0$, reduces to an LP; if $c_i = 0$, reduces to a QCQP
- more general than QCQP and LP

$$\begin{array}{c} \text{LP} \\ \hline \text{QP} \\ \hline \text{QCQP} \\ \hline \text{SOCP} \\ \hline \text{SDP} \end{array}$$

Robust linear programming

the parameters in optimization problems are often uncertain, *e.g.*, in an LP

$$\begin{cases} \text{minimize} & c^T x \\ \text{subject to} & \underbrace{a_i^T x}_{\leq b_i}, \quad i = 1, \dots, m, \end{cases}$$

there can be uncertainty in \underbrace{c} , $\underbrace{a_i}$, $\underbrace{b_i}$

two common approaches to handling uncertainty (in a_i , for simplicity)

- deterministic model: constraints must hold for all $\underbrace{a_i \in \mathcal{E}_i}$

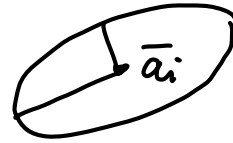
$$\begin{cases} \text{minimize} & c^T x \\ \text{subject to} & \underbrace{a_i^T x}_{\leq b_i} \text{ for all } \underbrace{a_i \in \mathcal{E}_i}, \quad i = 1, \dots, m, \end{cases}$$

- stochastic model: $\underbrace{a_i}_{\text{is random variable}}$ is random variable; constraints must hold with probability η

$$\begin{cases} \text{minimize} & c^T x \\ \text{subject to} & \underbrace{\text{prob}(a_i^T x \leq b_i) \geq \eta}, \quad i = 1, \dots, m \end{cases}$$

deterministic approach via SOCP

- choose an ellipsoid as \mathcal{E}_i :



$$\mathcal{E}_i = \{\bar{a}_i + P_i u \mid \|u\|_2 \leq 1\} \quad (\bar{a}_i \in \mathbf{R}^n, \quad P_i \in \mathbf{R}^{n \times n})$$

center is \bar{a}_i , semi-axes determined by singular values/vectors of P_i

- robust LP

$$\begin{cases} \text{minimize} & c^T x \\ \text{subject to} & \underline{a_i^T x \leq b_i} \quad \forall a_i \in \mathcal{E}_i, \quad i = 1, \dots, m \end{cases}$$

is equivalent to the SOCP

$$\begin{cases} \text{minimize} & c^T x \\ \text{subject to} & \bar{a}_i^T x + \|P_i^T x\|_2 \leq b_i, \quad i = 1, \dots, m \end{cases}$$

(follows from $\sup_{\|u\|_2 \leq 1} (\bar{a}_i + P_i u)^T x = \bar{a}_i^T x + \|P_i^T x\|_2$)

$$a_i^T x \leq b_i \quad \text{for all } a_i \in \{ \bar{a}_i + P_i u \mid \|u\|_2 \leq 1 \}$$

\Updownarrow

$$\sup_{\|u\|_2 \leq 1} (\bar{a}_i + P_i u)^T x \leq b_i$$

$$\|u\|_2 \leq 1$$

$$\bar{a}_i^T x + \sup_{\|u\|_2 \leq 1} u^T (P_i^T x) \leq b_i$$

$$\underbrace{\sup_{\|u\|_2 \leq 1} u^T (P_i^T x)}_{= \|u\|_2 \|P_i^T x\|_2 \leq 1} \leq b_i - \bar{a}_i^T x$$

$$= \|P_i^T x\|_2$$

Cauchy-Schwartz

$$|x^T y| \leq \|x\|_2 \|y\|_2$$

$$\bar{a}_i^T x + \|P_i^T x\|_2 \leq b_i \quad i = 1, \dots, m$$

$$\|P_i^T x\|_2 \leq b_i - \bar{a}_i^T x$$

stochastic approach via SOCP

- assume a_i is Gaussian with mean \bar{a}_i , covariance Σ_i ($a_i \sim \mathcal{N}(\bar{a}_i, \Sigma_i)$)

- $a_i^T x$ is Gaussian r.v. with mean $\bar{a}_i^T x$, variance $x^T \Sigma_i x$; hence $\mathcal{N}(0, 1)$

$$\text{prob}(a_i^T x \leq b_i) = \Phi \left(\frac{b_i - \bar{a}_i^T x}{\|\Sigma_i^{1/2} x\|_2} \right)$$

where $\Phi(x) = (1/\sqrt{2\pi}) \int_{-\infty}^x e^{-t^2/2} dt$ is CDF of $\mathcal{N}(0, 1)$

- robust LP

$$\begin{cases} \text{minimize} & c^T x \\ \text{subject to} & \text{prob}(a_i^T x \leq b_i) \geq \eta, \quad i = 1, \dots, m, \end{cases}$$

with $\eta \geq 1/2$, is equivalent to the SOCP

$$\begin{cases} \text{minimize} & c^T x \\ \text{subject to} & \bar{a}_i^T x + \Phi^{-1}(\eta) \|\Sigma_i^{1/2} x\|_2 \leq b_i, \quad i = 1, \dots, m \end{cases}$$

$$\Pr(a_i^T x \leq b_i) = \Pr\left(\frac{a_i^T x - \bar{a}_i^T x}{\underbrace{\sqrt{x^T \Sigma_i x}}_{= \|\Sigma_i^{1/2} x\|_2}} \leq \underbrace{\frac{b_i - \bar{a}_i^T x}{\sqrt{x^T \Sigma_i x}}}_{\sim N(0,1)}\right)$$

$$= \Phi\left(\frac{b_i - \bar{a}_i^T x}{\|\Sigma_i^{1/2} x\|_2}\right)$$

$$\Phi\left(\frac{b_i - \bar{a}_i^T x}{\|\Sigma_i^{1/2} x\|_2}\right) \geq \eta$$

$$\frac{b_i - \bar{a}_i^T x}{\|\Sigma_i^{1/2} x\|_2} \geq \Phi^{-1}(\eta) \Rightarrow$$

$$b_i \geq \underbrace{\Phi^{-1}(\eta)}_{\text{SOC constraint}} \underbrace{\|\Sigma_i^{1/2} x\|_2}_{\text{SOC constraint}} + \bar{a}_i^T x$$

Φ : CDF of $N(0,1)$

