

Expectation of Discrete random variables

Definition : The expectation of a discrete rv X is given by

$$\begin{aligned} E(X) &= \sum_{x: p(x) > 0} x p(x) \\ &= \sum_{x: p(x) > 0} x \cdot P(X=x) \end{aligned}$$

When $\sum_{x: p(x) > 0} |x| \cdot p(x) < \infty$.

Ex : A fair die is rolled twice
 $X = \#$ 1s observed

$$\begin{aligned} E(X) &= 0 \cdot P(X=0) + 1 \cdot \underbrace{P(X=1)}_{10/36} + 2 \cdot \underbrace{P(X=2)}_{1/36} \\ &= \frac{1}{3} \end{aligned}$$

Ex : Suppose X has a pmf

$$p(n) = \frac{1}{n(n+1)}, \quad n=1, 2, 3, \dots$$

$$\sum_{x=1}^{\infty} |x| \cdot p(x) = \sum_{x=1}^{\infty} x \cdot \frac{1}{x(x+1)} = \sum_{x=1}^{\infty} \frac{1}{x+1} \text{ diverges.}$$

$\Rightarrow E(X)$ is undefined.

Theorem: Let X be a discrete
rv & $g: \mathbb{R} \rightarrow \mathbb{R}$ such that

$$\sum_x |g(x)| p(x) < \infty.$$

Then $E[g(X)] = \sum g(x) p(x)$

Proof: Let $\{y_j\}$ denote all the possible
values $g(X)$ can take.

Define $A_j := \{x : g(x) = y_j\}$

$\Rightarrow (g(A_j) = \{y_j\})$

Let $Y = g(X)$.

$$E[g(X)] = E[Y] = \sum_j y_j P(Y = y_j)$$

$$= \sum_j y_j \cdot P(X \in A_j)$$

$$= \sum_j y_j \sum_{x \in A_j} P(X = x)$$

$$= \sum_j \sum_{x \in A_j} g(x) P(X = x) \quad \text{since } g(x) = y_j \quad \forall x \in A_j.$$

$$= \sum_x g(x) P(X=x) \quad \text{since } A_i \cap A_j = \emptyset, \forall i \neq j.$$

Prop :

$$(i) E[X+Y] = E[X] + E[Y]$$

$$(ii) E[kX] = kE[X], \quad \forall k \in \mathbb{R}.$$

$$(iii) X \leq Y \Rightarrow E(X) \leq E(Y)$$

$$(iv) X = k \text{ (constant)}, \quad E(X) = k.$$

$$(v) E(1_A) = P(A)$$

Proof : (i) Let $\{x_i\}$ be the values X takes
 $\{y_j\}$ denote the values Y takes &
 $\{z_k\}$ be the values $Z = X + Y$ takes.

$$A_k = \{(i, j) : x_i + y_j = z_k\}$$

$$\text{Notice that } P(X+Y = z_k) = \sum_{(i,j) \in A_k} P(X = x_i, Y = y_j)$$

$$E(Z) = \sum_k z_k P(X+Y = z_k)$$

$$= \sum_k z_k \sum_{(i,j) \in A_k} P(X = x_i, Y = y_j)$$

$$= \sum_k \sum_{A_k} (x_i + y_j) P(X = x_i, Y = y_j)$$

$$= \sum_i \sum_j (x_i + y_j) P(X = x_i, Y = y_j) = \sum_i \sum_j x_i P(X = x_i, Y = y_j)$$

$$+ \sum_i \sum_j y_j P(X = x_i, Y = y_j) = \sum_i x_i P(X = x_i) + \sum_j y_j P(Y = y_j) = E(X) + E(Y) \quad \square$$

Moments

Definition: Let X be a rv with $E(X) < \infty$. The variance of X is defined by

$$V(X) = E(X - E(X))^2$$

Prop: $V(X) = E(X^2) - E(X)^2$

Prop: $V(aX + b) = a^2 V(X)$.

Ex: Prediction:

Let X be a rv. What t best predicts X (i.e. minimizes $E[(X - t)^2]$)?

$$R(t) = E[(X - t)^2]$$

$$= E[X^2 - 2tX + t^2]$$

$$R(t) = E[X^2] - 2t E(X) + t^2$$

$$\frac{d}{dt} R(t) = -2E(X) + 2t = 0 \Rightarrow t = E(X).$$

The value that best predicts X is $E(X)$.

Special Discrete Distributions

① Bernoulli Distribution.

We say X is Bernoulli if

$$P(X=1) = p, \quad P(X=0) = 1-p =: q$$

Ex: If A is an event, then 1_A is a Bernoulli RV.

$$E(X) = p$$

$$V(X) = p(1-p)$$

② Binomial RV

Let X_1, X_2, \dots, X_n be independent Bernoulli RVs, Then

$X = \sum_{i=1}^n X_i$ is called a binomial

RV with parameters n, p

$X = \#$ successes of n independent Bernoulli RVs.

$$P(X=k) = \binom{n}{k} p^k (1-p)^{n-k}, \quad k=0, 1, 2, \dots, n$$

$$\sum_{k=0}^n \binom{n}{k} p^k (1-p)^{n-k} = (p + (1-p))^n = 1 \quad \checkmark$$

Notation $X \sim B(n, p)$

Proposition : $E(X) = np$
 $V(X) = np(1-p)$

Ex : $X \sim B(n, p)$

$$\begin{aligned} E(a^X) &= \sum_{k=0}^n a^k \binom{n}{k} p^k (1-p)^{n-k} \\ &= \sum_{k=0}^n \binom{n}{k} (ap)^k (1-p)^{n-k} \end{aligned}$$

By the binomial thm

$$= \boxed{(ap + (1-p))^n}$$