Performance and Simulation of Social Networks

Jakob Wyatt
19477143

October 23, 2019

1 Abstract

2 Background

This report focuses on simulating simple social media networks, and evaluates the effectiveness of a network given parameters of that network.
The social network consists of a set of users that may follow eachother, which has been represented in code as a directed graph. Users may not follow themselves, or follow eachother more than once.
There exists only one post at a time, with a new post being loaded when the current post has not had any activity in the last timestep. The original poster always likes their own post. A user can only like a post once.
The simulation consists of timesteps, with a function `update()` to move between timesteps. The update algorithm works as follows:

1. Check that there exists some users that have liked the post in the previous timestep. If there are none, the update ends and the next post is loaded.
2. Iterate through all users who liked the post in the previous timestep. Each of their followers is 'exposed' to the post, and have a chance of liking the post. This chance is sampled from a Bernoulli distribution with probability $\mathit{clamp}\left(\mathit{prob\_like} \times \mathit{clickbait\_factor}, 0, 1\right)$.
3. If a user likes a post in the current timestep, they have a chance of following the original poster. This is sampled using the same technique as above, with global probability $\mathit{prob\_foll}$.

Note that in the above algorithm, if a user does not like a post, they may potentially be exposed to it later via a different friend. This behaviour is intentional, as it incentivises a highly connected network.

Some parameters that will be varied in the creation of the social network include:

- Probability of liking a post
- Probability of following a user
- Number of users

Some parameters that will be tracked throughout the simulation include:

- Clustering Coefficient
- Average and s.d. of followers per user
- Likes per user per post

3 Methodology

Terminology in this section includes the word 'predecessor' and 'successor' in the context of graphs. If there is a directed edge between verticies $A \to B$, then B is a successor of A and A is a predecessor of B.

3.1 Parameter Gridsearch

To profile multiple runs of the simulation a gridsearch algorithm was used to vary the below parameters:

- Average and s.d. of followers per user
- Like/Follow probabilities
- Number of users in the network

First, arrays that contain the parameter values to be tested were created. Next, for all possible combinations of these parameters, a network matching these parameters was created. A simulation was then run on all networks, with simulation statistics being logged to a file as CSV data.

Parameters that are varied in this function are the like and follow probabilities, network size, clickbait standard deviation, post/timestep number, and follower average (density) / follower standard deviation (relative popularity).

An advantage of this algorithm is that it is very easy to implement compared to other search space algorithms, such as gradient descent and evolution. However, a disadvantage is that it searches a very large volume when the parameter space has many dimensions (curse of dimensionality), causing a large amount of time to generate the data. The algorithmic complexity of the gridsearch algorithm is $O(x^n)$, where n is the number of parameters that are varied and x is the number of datapoints per parameter.

3.2 Statistics

Throughout the simulation, various statistics about the state of the network were computed. These included the average/s.d. followers per user and the average number of likes per user. One particularly interesting statistic that was calculated was the Clustering Coefficient. To calculate the clustering coefficient of a vertex V, the neighbourhood of that vertex is first found.
$N = V_{successors} \bigcup V_{predecessors}$
Next, the number of connections between nodes in this neighbourhood is found.
$C = |e_{jk} : v_j, v_j \in N, e_{jk} \in E|$
Finally, this is divided by the total number of potential connections within the neighbourhood.
$ClusteringCoefficient = \frac{C}{|N| \times (|N|-1)}$
To find the network average clustering coefficient, the clustering coefficient is calculated for all verticies in the network and averaged. As the number of followers will always increase over time in this simulation, to enable comparison between clustering coefficients at different timesteps, the network average clustering coefficient was then divided by the number of edges in the graph.

3.3 Graph Generation

To generate the simulated networks, a unique algorithm was used that was created by the author, and to their knowledge, does not exist anywhere else.

This algorithm allows efficient generation of structured (non-random) networks, with a consistent runtime complexity of $O(V \times E)$.

First, a number of verticies are created. Next, these verticies are iterated over, and the number of successors of each vertex is calculated by sampling from a distribution $X \sim N\left(follower\_av, follwer\_sd^2\right)$.

To select the edges of the network, an approach similar in nature to a variogram (https://en.wikipedia.org/wiki/Variogram) was used (taken from the field of geostatistics). A variogram measures the amount of correlation between two points that are some distance apart. In this algorithm, this distance is discrete rather than continuous, and is measured by finding some path between two nodes. By increasing the spacial correlation at short distances, a graph can be created with a higher clustering coefficient. Reducing the spacial correlation to 0 causes the graph to become random.

A property of a graph that is analogous to the variogram is the clustering coefficient, which determines the correlation between verticies at a distance of 1. This clustering coefficient can be generalized to include all verticies within some distance h of a central vertex (Xiao et al. 2007). It is expected that this clustering coefficient will decrease as the distance increases in a real life network. An example of this is that your friends are likely to be friends of eachother, but your friends of friends are less likely to be friends of eachother. By assuming the correlation between verticies, realistic social networks can be generated for testing.

To generate an edge in the graph, a random walk is performed starting at a given vertex, V_0. This random walk occurs along any nodes that are predecessors of the current node, and does not loop back onto previously visited nodes. At each step on the random walk, the probability of creating an edge with V_0 is calculated by evaluating the variogram using the distance from the original node (length of the random walk). If an edge is created, the process starts again and repeats until all necessary edges have been created. If there are no available verticies to walk to, or the random walk has reached some distance threshold from the original node, an edge is created with any random vertex in the graph. This allows initial 'bootstrapping' of the network when no edges have yet been created.

3.4 Time and Space Complexity

The post propogation algorithm used in this program has a time complexity of $O(V \times E)$. However, many algorithms that calculate statistics about the network have a much worse time complexity. One example is the calculation of the clustering coefficient, which has a time complexity of $O\left(V^3\right)$ when implemented using a hash table and $O\left(V^4\right)$ when implemented using a linked list.

3.5 Execution

The generated data for this simulation can be found in report/data.csv. This data can be created by running the program found in src/SocialNetworkSimRunner.py.

This program takes approximately 30 minutes to run on an i7-7700k processor. Gridsearch parameters can be changed within the gridsearch function.

4 Results

4.1 Time

It was found that clustering coefficient did not change significantly over the course of the simulation, and either increases or decreases by around 1% over the course of the simulation. This suggests that over time, friend groups do not become more or less clustered, and remain approximately the same. This makes sense, as the proportion of posts a friend group will be exposed to within the group is proportional to the clustering of the group itself.

Another result that was found was that the standard deviation of the follower count decreased over time in almost all cases. The only exception to this was where the standard deviation of the follower count was 0. This shows that, over time, the social network will become more equitable, with famous people becoming less popular over time. However, this is based off the assumption that there is no difference in content between x and y. This can be observed in real life.

As expected, follower count and number of likes consistently increased over time. However, what is interesting is how they increased over time. (Insert graphs). This is applicable to social networks, where a large factor in a persons online popularity is related to how many posts they have.

4.2 Size

Even when scaled for network size, the number of likes consistently increased when the size of the network was increased.

5 Conclusion and Future Work

Future work that could be done on this simulation involve both better analysis of the current data, generation of networks using different algorithms, and implementing more graph statistics to measure.

Due to time constraints, only a limited amount of basic analysis could be done on generated data. This analysis was mostly constrained to setting constant values for certain parameters, and varying a single parameter at a time to see the outcome. This assumes that each parameter is independent, however this may not be the case.

A more in-depth analysis could involve principal component analysis to rigorously determine the most important factors in the performance of a social network.

The current algorithm used to generate random graphs with structure is not mathematically rigorous. One method of verifying the results of this algorithm

would be to implement a different graph generation algorithm such as ClustRNet (Bansal, Khandelwal, and Meyers, 2009), and compare the results with the current method of generating random graphs. This was not done due to time constraints.

Only a small amount of statistics related to graphs were analysed. To improve this analysis in future, a wider range of statistics about the graph could be calculated at each time step, including but not limited to:

- Mutuality (The probability that someone is followed back)
- Distance (The minimum number of edges required to connect any two verticies in the graph)
- Centrality of highly followed users (How much influence does a node have)

This would allow a more in depth analysis into the factors behind social network performance, and a better comparison with real life networks.

More work that could be done to improve this simulation would be to increase its complexity. One possible way of doing this could be to introduce individual like probabilities per user. However, this added complexity would not be worth the implementation time required, with the previous three suggestions offering greater benefit for less development time.

6 References

Xiao, Wenjun, Wenhong Wei, Weidong Chen, Yong Qin, Behrooz Parhami. 2007. *Extended Clustering Coefficients: Generalization of Clustering Coefficients in Small-World Networks.* Journal of Systems Science and Systems Engineering. https://www.ece.ucsb.edu/ parhami/pubs_folder/parh07-jssse-ext-clustering-coeff.pdf

Bansal, Shweta, Shashank Khandelwal, and Lauren Ancel Meyers. 2009. *Exploring Biological Network Structure with Clustered Random Networks.* BMC Bioinformatics. https://bmcbioinformatics.biomedcentral.com/articles/10.1186/1471-2105-10-405