Jalen Powell

1. a.) baab

S- AaBb

- → ba Bb (A=>b)
- -> baab (8=7a)
- · Valid Sentence

b.) bbbab

S-> AaBb

- -> AbaBb (A=> Ab)
- > AbbaBb (A=>Ab)
- > bbbaBb (A=>b)
- To get the desired outcome, this next step can't be executed. Invalid

c.) bbaaaaa

5-AaBb

- -) AbaBb (A=7Ab)
- → bbaBb (A=>b)
- > bbaqBb (B=>aB)
- The last step is a b, so it wont reach desired outputs
- · Involid sentence

d.) bbaab

S>AaBb

- →AbaBb (A=>Ab)
- → bbaBb (A=>b)
- > bboab (B=>a)
- · Valid Sentence
- 2. Identify all of the tokens, and which lexemes they categorize. Put them in a toke.

Tokens	Levemes
Assignmen+ Operator	=
Arithmetic Operators	+,*
Iden+17iers	Ay B, C
Crouping Symbols	(1)

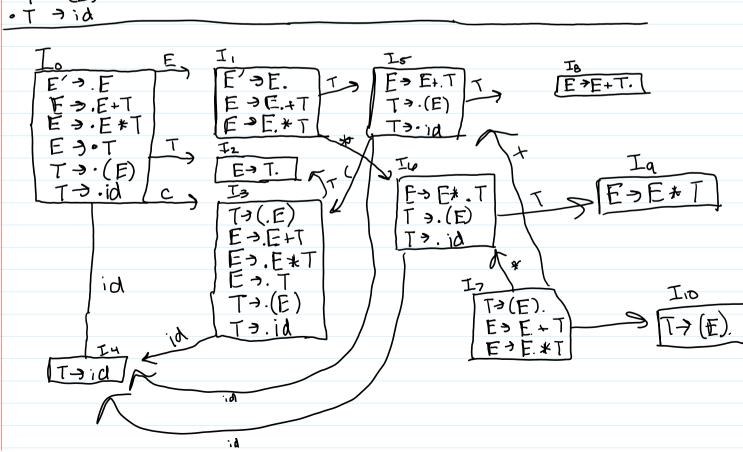
3.
$$B = B + (c + (A + A))$$
 $\langle assign \rangle \rightarrow \langle id \rangle = \langle expr \rangle$
 $\Rightarrow B = \langle expr \rangle$
 $\Rightarrow B = \langle expr \rangle$
 $\Rightarrow B = B + \langle expr \rangle$

>B= B+ (C+(A*A)

assign
$$\begin{array}{rcl}
\langle id \rangle &=& \langle e \times pr \rangle \\
0 & \langle id \rangle &+& \langle e \times pr \rangle \\
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- · A > Aal AbC
- · A > (
- · A > AA' => A' → albC
- · S -> Aalbb

- · A > AA
- · A' Jal bC
- · B > C
- · B > 5166
- · C > C
- 5. A > abclacla B> blab
 - · A > a K
 - ·A' → Bc/c/E ·B > blaB
- G. E → E+ TIE* TIT T > (E) lid
 - ・ビーラビ
 - · E>E+T
 - · E > E * T
 - · T > (E)



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	State	Action				C-c	AC	1		
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	9	r ²	r2	r2	12	rz	r2	- /		
1	10	r4	r4	r4	l r4	r4	1r4/			

- 1) E > E+T 2) E > E*T
- 3) E> T 4) T>(E) 5) T>id

7 (id+id) + ia

Stack	1 Buffer	Action	
\$	(ia+ ia) + ia 5	_	
\$	(ia+ia)*ia\$	shif+	
\$ (id+ia)*ia\$	Shif+	
\$ 1d	/ +ia) * id\$	reduce T=id	
\$(T	+ia)* id\$	teduce E>T	
\$ (E	+id)# id\$	Shift	
D (E+	1d)* id\$	Shift	
\$(E+id	? id\$	reduce T > id	
\$ (E+T)* id\$	reduce E>E+T	
3 (E		shift	
\$ (E)	* id\$	reduce $T \Rightarrow (E)$	
\$ T	* 10\$	reduce FTT	
\$ E	* id \$	5Nift	
\$ E*	id \$	Shif+	
\$ E* 19	\$	reduce T→id	
\$ E * T	\$	reduce E>E*T	
\$E	\$	Accept	
	•		

Ψ ·		• - • • • •
\$ E	* id \$	5Nift
\$ E*	id \$	8hi++
\$ E* 19	\$	reduce T→id
SE*T	\$	reduce E>E*T
\$E	\$	Accept
	•	•

8. (id+ia)+id

$$E \rightarrow E + T$$

$$\rightarrow E + id$$

$$\rightarrow T + id$$

$$\rightarrow E \rightarrow T$$

$$\rightarrow (E) + id$$

$$\rightarrow (E + T) + id$$

$$T \rightarrow (E)$$

$$\rightarrow (E + ia) + id$$

$$T \rightarrow id$$

$$T \rightarrow id$$

$$T \rightarrow id$$

$$T \rightarrow id$$