## SC223 - Linear Algebra

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Lecture 8



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### LU Decomposition

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$$PA = LU$$
,

where  $P \in \mathbb{R}^{m \times m}$ ,  $L \in \mathbb{R}^{m \times m}$ ,  $U \in \mathbb{R}^{m \times n}$ .

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 $Ax = b \rightarrow LUx = b$ .

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- ▶ In the example, where is y?

$$\begin{bmatrix} \mathbf{1} & -1 & 2 & 3 & -1 & 1 \\ 0 & \mathbf{2} & 5 & 4 & 2 & 2 \\ 0 & 0 & -\mathbf{23} & -16 & 2 & -12 \\ 0 & 0 & 0 & -\mathbf{137} & -95 & 64 \end{bmatrix}$$

#### Example

$$\bullet \ U = \begin{bmatrix} \mathbf{1} & -1 & 2 & 3 & -1 \\ 0 & \mathbf{2} & 5 & 4 & 2 \\ 0 & 0 & -\mathbf{23} & -16 & 2 \\ 0 & 0 & 0 & -\mathbf{137} & -95 \end{bmatrix}, y = \begin{bmatrix} 1 \\ 2 \\ -12 \\ 64 \end{bmatrix}.$$

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Solve for x in

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## LU Decomposition Algorithm

• Let 
$$A = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & \dots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nn} \end{bmatrix}$$

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► Thus, 
$$L = \begin{bmatrix} 1 & 0 & \dots & 0 \\ \frac{a_{21}}{a_{11}} & 1 & \dots & 0 \\ \frac{a_{31}}{a_{11}} & \frac{a_{32}}{a_{22}} & \dots & 0 \\ \vdots & \vdots & \vdots & \vdots \\ \frac{a_{n1}}{a_{11}} & \frac{a_{n2}^1}{a_{22}} & \dots & 1 \end{bmatrix}$$

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- Why should one use *LU* decomposition?