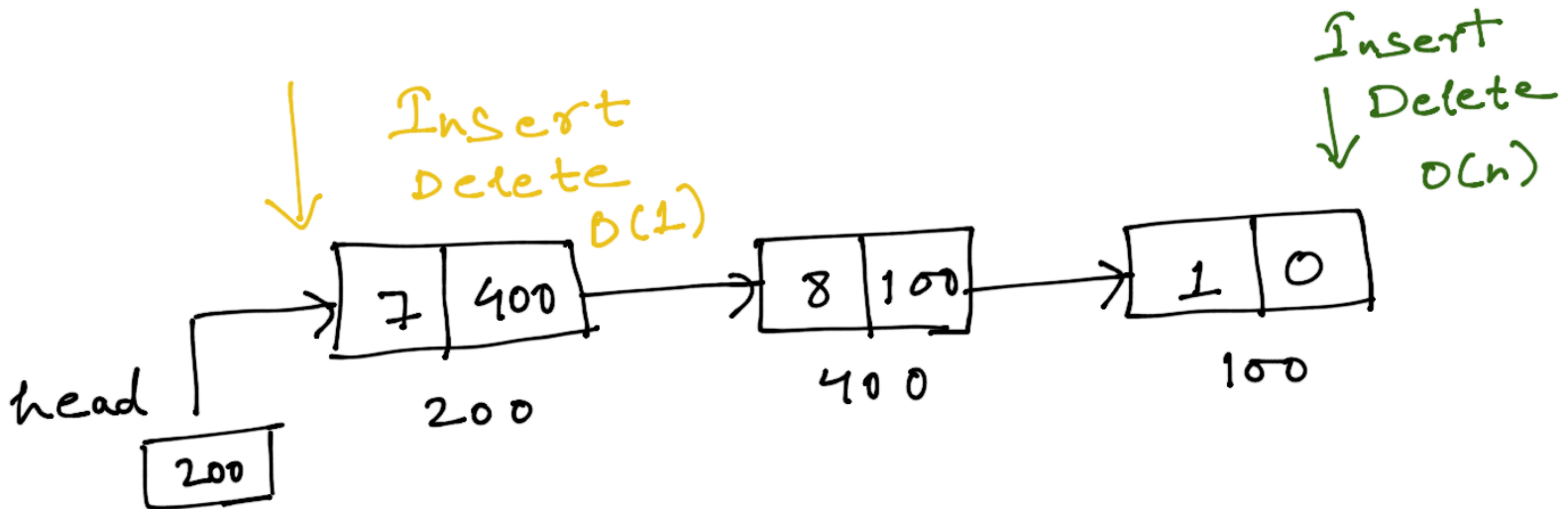


Implementation of Stack using Linked List

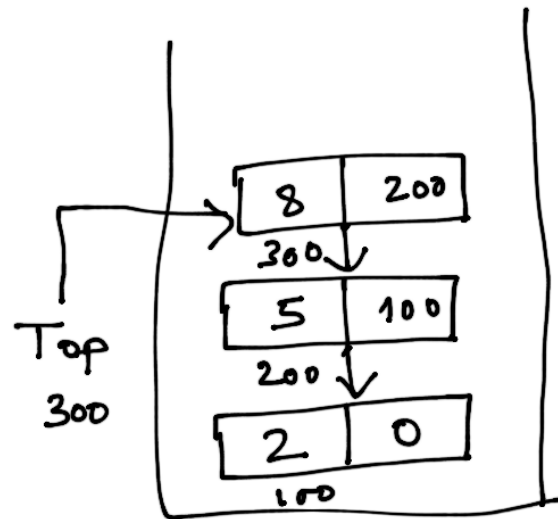


Stack

LIFO

push: $O(1)$

pop: $O(1)$

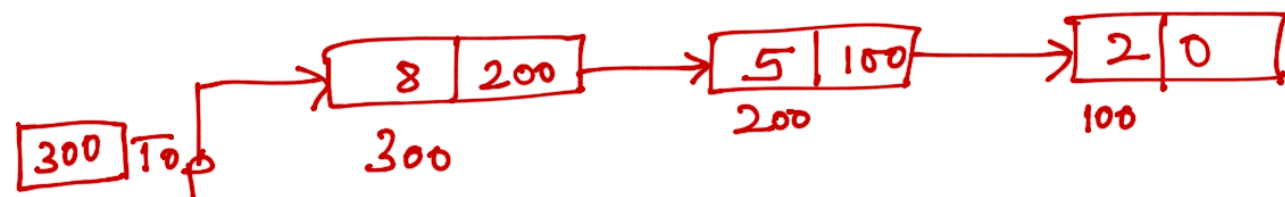


push(2)

push(5)

push(8)

push(10)



push (value)

Steps

1. Create a new Node and set data part = value
2. check whether stack is Empty ($top == NULL$)
3. If it is Empty, then set $newNode \rightarrow next = NULL$
4. If it is NOT Empty, then set $newNode \rightarrow next = top$
5. Then, set $top = newNode$

pop()

Steps

1. Check whether stack is Empty ($top == NULL$)
2. If it is Empty, then display "Stack underflow / stack is Empty"
3. If it is NOT Empty, then define a Node pointer 'temp' and set it to 'top'.
4. Then, set $top = top \rightarrow next$
5. Finally, delete 'temp'.

Infin / Prefin / Postfin

Infin:

$$5+1$$

which is correct

$$5+1*6$$

$$= 36$$

$$= 11$$

Precedence

1) $() , [] , \{ \}$

2) $\wedge \rightarrow R-L$

3) $* / \rightarrow L-R$

4) $+ - \rightarrow L-R$

Eg: $1 + \underline{2*5} + 30/5$

$$= 1 + 10 + 30/5$$

$$= 1 + 10 + 6$$

$$= 11 + 6$$

$$= 17$$

Eg: $2^{\wedge} 2^{\wedge} 3$

$$= 2^{\wedge} 8$$

$$= 256$$

NOT

$$2^{\wedge} 2^{\wedge} 3$$

$$\rightarrow 4^{\wedge} 3$$

$$= 64 \quad X$$

Prefix / Polish Notation

$\langle \text{operator} \rangle \langle \text{operand} \rangle \langle \text{operand} \rangle$

$$5 + 1 \Rightarrow + 5 1$$

$$a * b + c \Rightarrow * a b + c$$

$$\Rightarrow + * a b c$$

Postfix / Reverse Polish Notation

$\langle \text{operand} \rangle \langle \text{operand} \rangle \langle \text{operator} \rangle$

$$5 + 1 \Rightarrow 5 1 +$$

$$a * b + c \Rightarrow a b * + c$$

$$\Rightarrow a b * c +$$

- Memory consumption is less for prefix/postfix
- Easy for computers to evaluate

Infix to Postfix Using Stack

Precedence Rule

()

^ — R-L

* / — L-R

+ - — L-R



$$\begin{aligned} \text{Eg: } 1) \quad & A + B * C \\ &= A + (B * C) \quad (\text{Scan 1}) \\ &= A + \underbrace{BC *}_{\text{Scan 2}} \end{aligned}$$

$$= A \underbrace{BC *}_{\text{postfix Exp}} +$$

$$\begin{aligned} 2) \quad & K + L - M * N + (O \wedge P) * W \quad (\text{Scan 1}) \\ &= K + L - \underbrace{M * N}_{\text{Scan 2}} + O \wedge P * W \quad (\text{Scan 2}) \\ &= K + L - \underbrace{MN *}_{\text{Scan 3}} + O \wedge P * W \quad (\text{Scan 3}) \\ &= K + L - \underbrace{MN *}_{\text{Scan 4}} + O \wedge P * W \quad (\text{Scan 4}) \\ &= \underbrace{KL +}_{\text{Scan 5}} - \underbrace{MN *}_{\text{Scan 6}} + O \wedge P * W \quad (\text{Scan 5}) \\ &= \underbrace{KL + MN *}_{\text{Scan 6}} - O \wedge P * W \quad (\text{Scan 6}) \end{aligned}$$

→ Using Stacks, we can evaluate this expression in just one scan

→ Once you reach the right end of the expression, we will get the equivalent postfix expression.

Eg: $K + L - M$

Input Exp

K

+

L

-

Stack

+

+

-

Postfix Exp

K

K

K L

K L +

('+' pops out of the stack bcoz of LR asso. though same precedence)

M

-

K L + M

Postfix Exp → $K L + M -$

Eg: $K + L - M * N$

Input Exp

Stack

Postfix Exp

K

K

+

+

K

L

+

KL

-

-

KL +

M

-

KL + M

*

$\frac{*}{-}$

KL + M

N

- *

KL + MN

Postfix Exp \rightarrow $KL + MN * -$

Eg: $K + L - M * N + (O \wedge P) * W$

Input Exp

Stack

Postfix Exp

K

K

+

+

K

L

+

KL

-

-

KL+

M

-

KL+M

*

-*

KL+M

N

-*

KL+MN

+

+

KL+MN*-

(

+(

KL+MN*-

O

+(

KL+MN*-O

^

+(^

KL+MN*-O

P

+(^

KL+MN*-OP

)

+

KL+MN*-OP^

*
W

+

KL+MN*-OP^

KL+MN*-OP^W*+