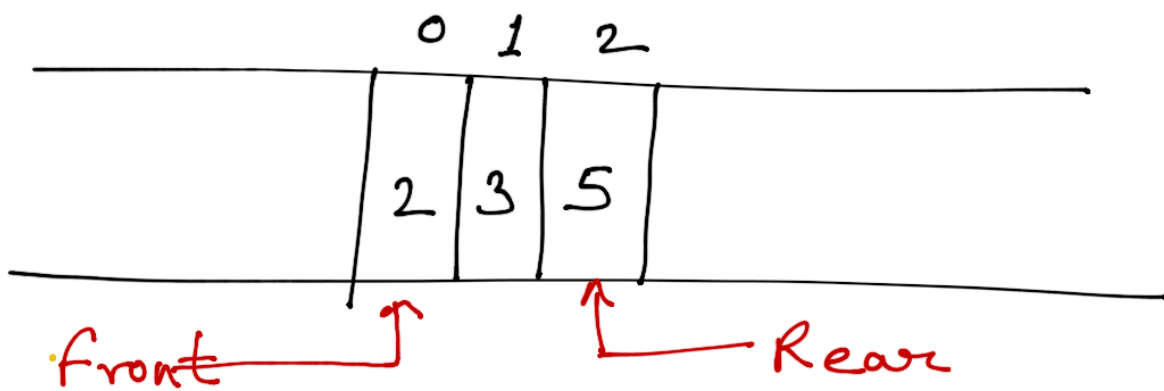


Queue

- linear data structure
- as ADT
- Based on FIFO

Enqueue \rightarrow Insertion — Rear/Tail

Dequeue \rightarrow Deletion — Head/front



Operations

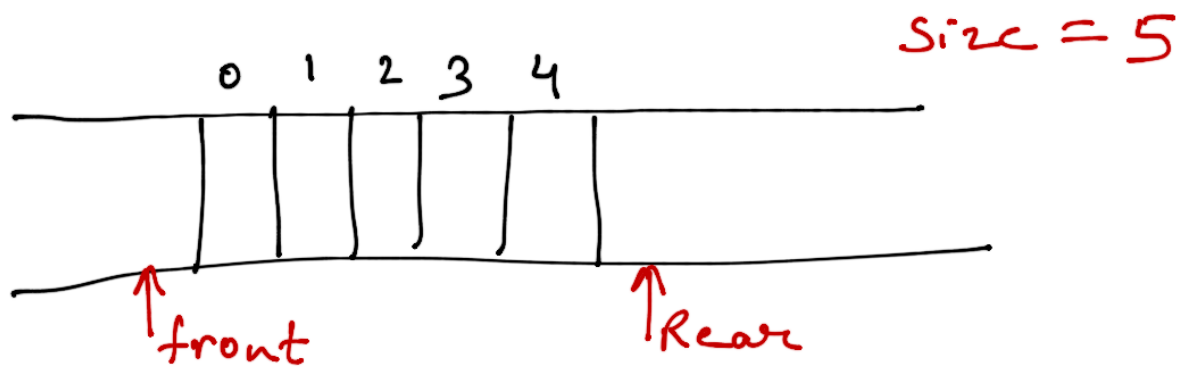
- Enqueue(2)

- Dequeue()

- Peek() / front()

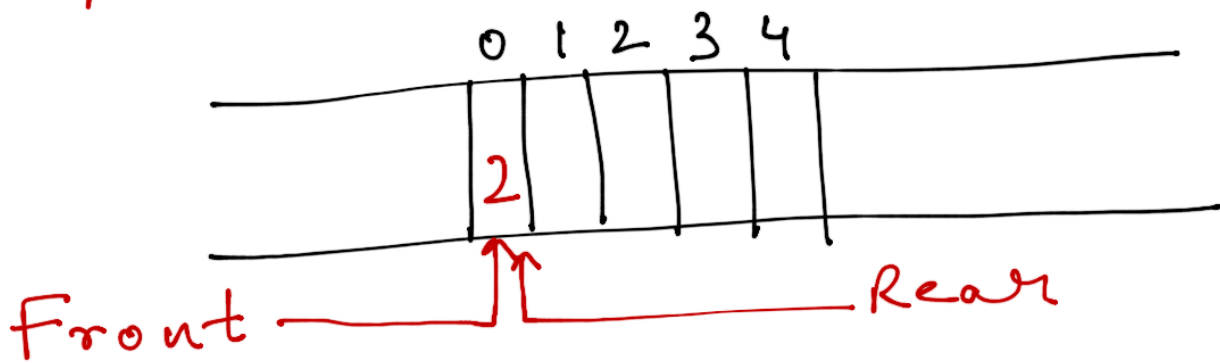
- isFull()

- isEmpty()

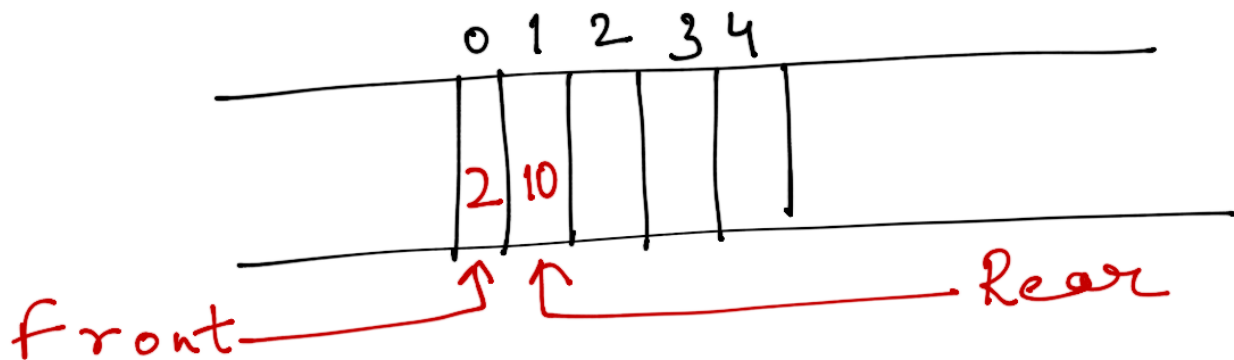


Initial front = Rear = -1 Empty
(No Element)

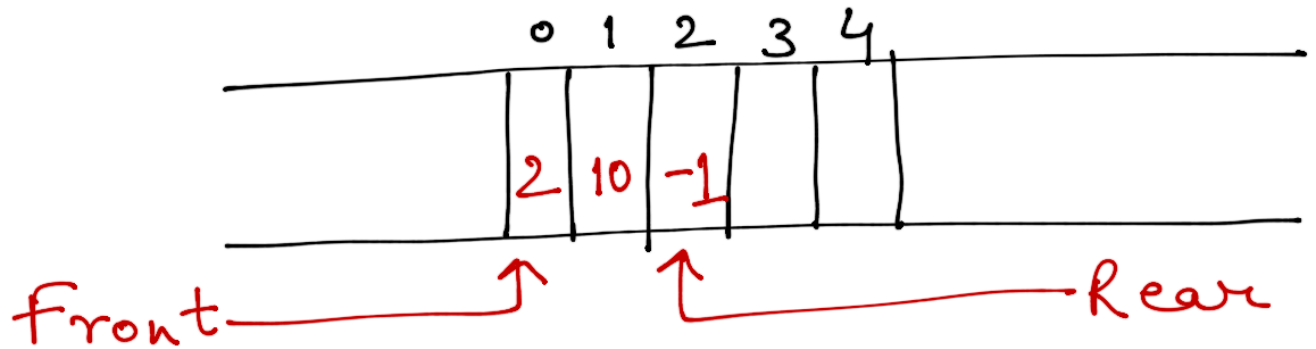
Enqueue(2)



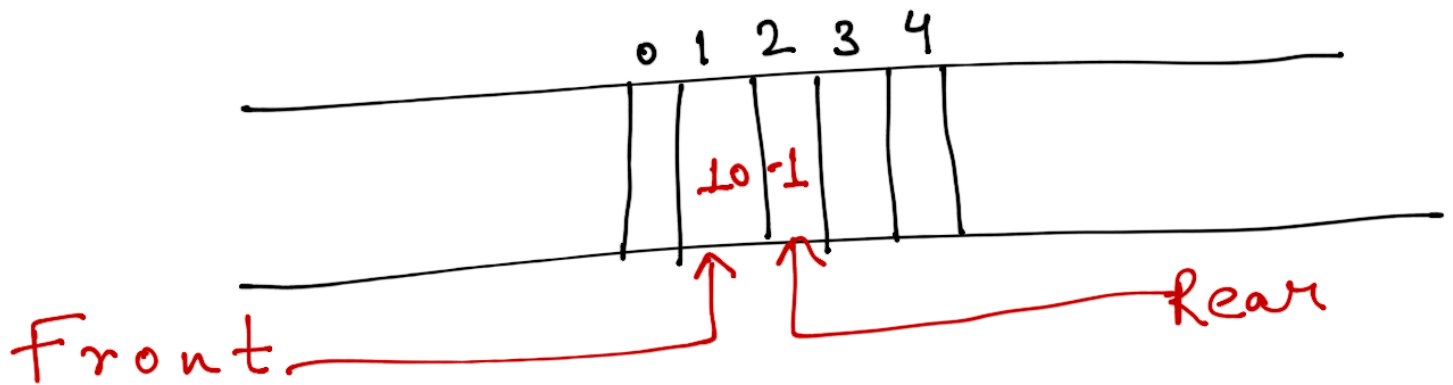
Enqueue(10)



Enqueue(-1)

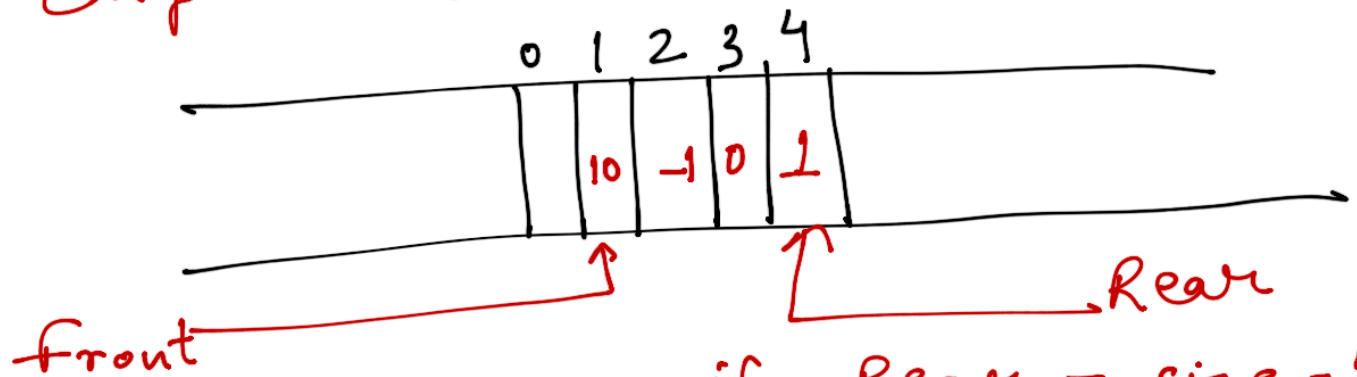


Dequeue()



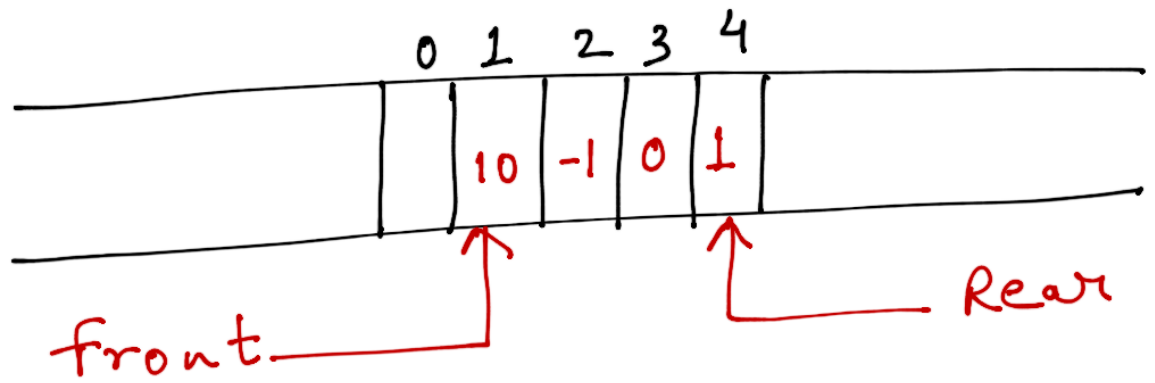
(Increment the front)

Enqueue(0), Enqueue(1)



Enqueue(5)

if $\text{Rear} = \text{size} - 1$
"overflow"

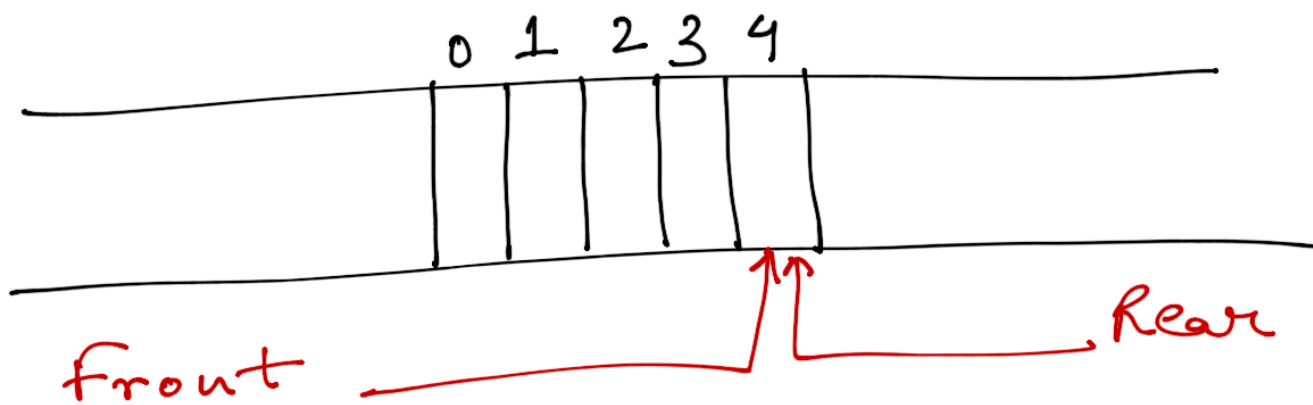


Peek() \rightarrow 10 (Element present at front position)

Dequeue()

Dequeue()

Dequeue()

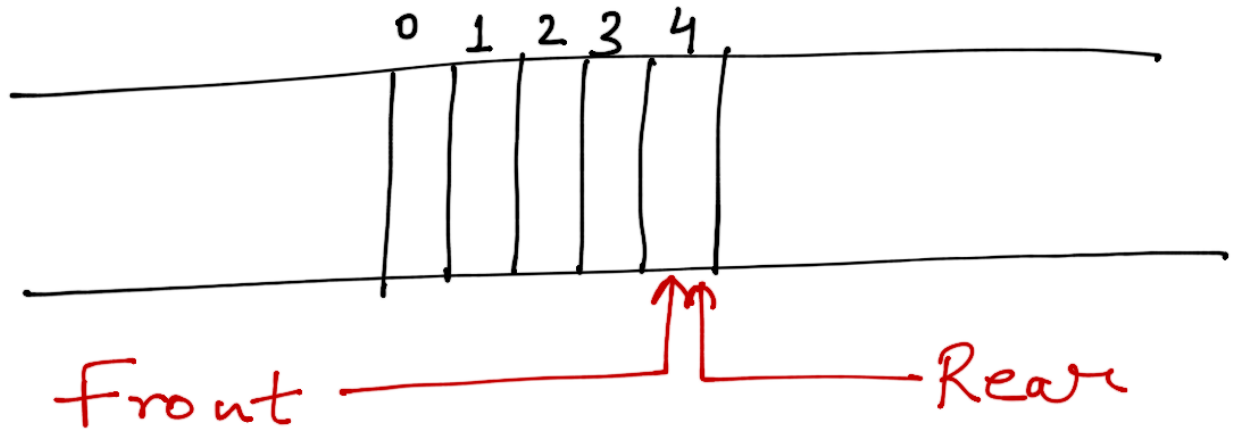


front == Rear, only 1 element

After removing, front = Rear = -1
OR, front > Rear \Rightarrow Empty

Enqueue (-15)

size = 5

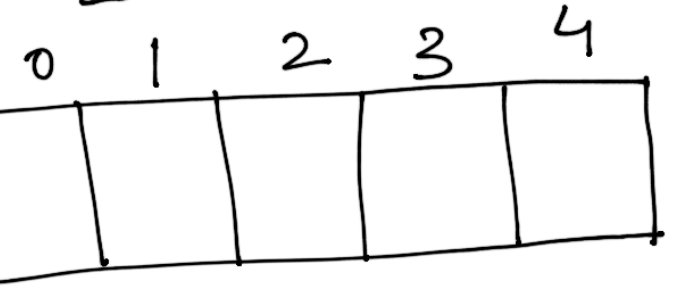


Can we enqueue here ??

Implementation Using Arrays

int queue [size]

Enqueue(DataElement)



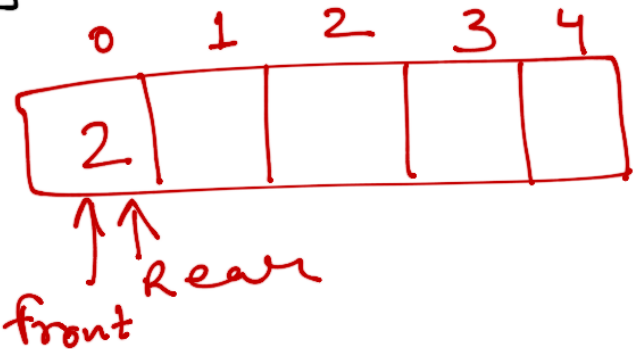
if (Rear == Size - 1)

{ "overflow" }

elseif (Front == -1 && Rear == -1)

{ front = Rear = 0

queue[Rear] = DataElement



else {

Rear++

queue[Rear] = DataElement

}

Dequeue()

```
{  
    if (Front == -1 && Rear == -1)  
        { "Underflow" }  
    elseif (Front == Rear)  
        { Front = Rear = -1 }  
    else  
        { Front++ }  
}
```

Display ()

```
{  
  if (front == -1 && Rear == -1)  
  { Empty }  
  else  
  {  
    for (i = front; i < Rear + 1; i++)  
    {  
      PRINT queue[i]  
    }  
  }  
}
```

