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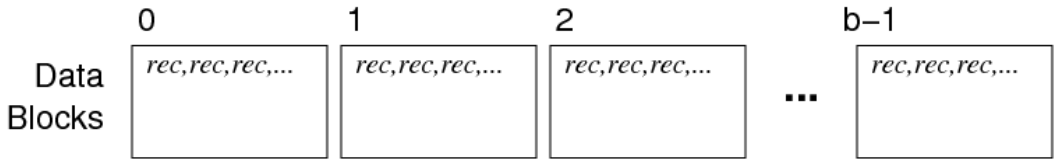
Scanning

❖ Scanning

Consider executing the query:

```
select * from Rel;
```

where the relation has a file structure like:



This would done by a simple scan of all records/tuples.

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❖ Scanning (cont)

Abstract view of how the scan might be implemented:

```
for each tuple T in relation Rel {
    add tuple T to result set
}
```

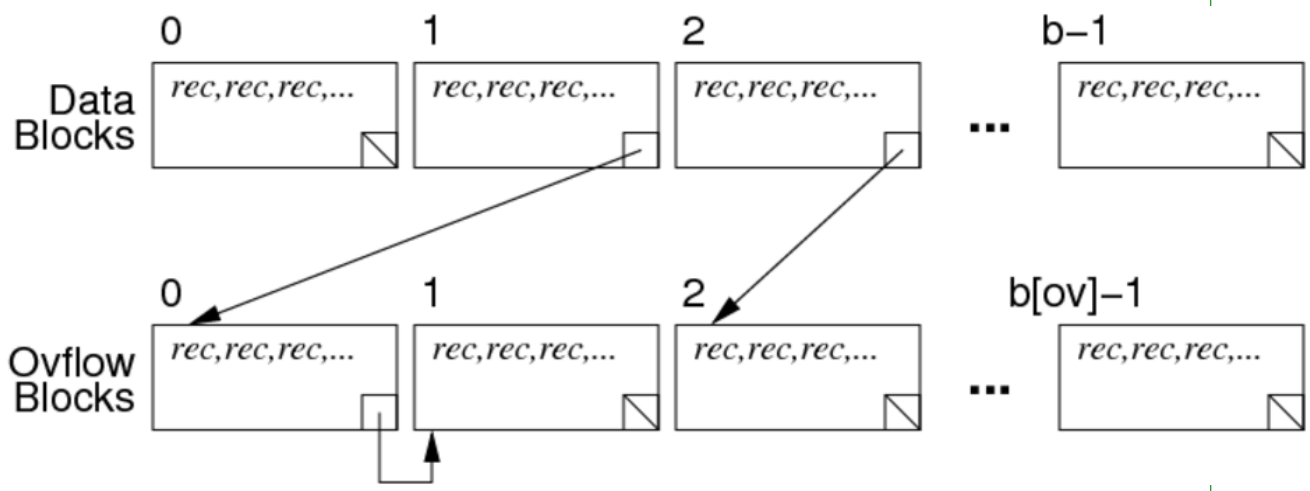
Operational view:

```
for each page P in file of relation Rel {
   for each tuple T in page P {
      add tuple T to result set
   }
}
```

Cost = read every data page once = *b*

❖ Scanning (cont)

Consider a file with overflow pages, e.g.



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❖ Scanning (cont)

In this case, the implementation changes to:

```
for each page P in data file of relation Rel {
    for each tuple t in page P {
        add tuple t to result set
    }
    for each overflow page V of page P {
        for each tuple t in page V {
            add tuple t to result set
}   }   }
```

Cost: read each data page and each overflow page once

$Cost = b + b_{Ov}$

where b_{Ov} = total number of overflow pages

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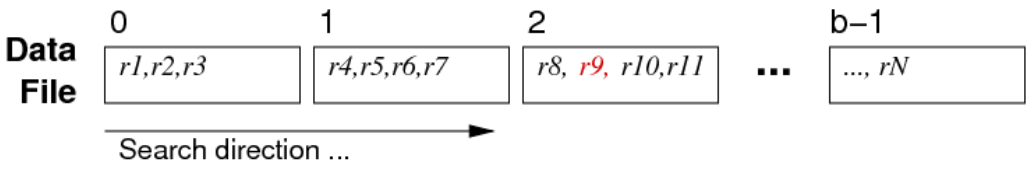
❖ Selection via Scanning

Consider a *one* query like:

```
select * from Employee where id = 762288;
```

In an unordered file, search for matching tuple requires:



Guaranteed at most one answer; but could be in any page.

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❖ Selection via Scanning (cont)

Overview of scan process:

```
for each page P in relation Employee {
    for each tuple t in page P {
        if (t.id == 762288) return t
}   }
```

Cost analysis for *one* searching in unordered file

- best case: read one page, find tuple
- worst case: read all *b* pages, find in last (or don't find)
- average case: read half of the pages (*b/2*)

Page Costs: $Cost_{avg} = b/2$ $Cost_{min} = 1$ $Cost_{max} = b$

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❖ Iterators

Access methods typically involve iterators, e.g.

```
Scan s = start_scan(Relation r, ...)
```

- commence a scan of relation **`r`**
- **`Scan`** may include condition to implement **`WHERE`**-clause
- **`Scan`** holds data on progress through file (e.g. current page)

```
Tuple next_tuple(Scan s)
```

- return **`Tuple`** immediately following last accessed one
- returns **`NULL`** if no more **`Tuples`** left in the relation

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❖ Example Query

Example: simple scan of a table ...

```
select name from Employee
```

implemented as:

```
DB db = openDatabase("myDB");
Relation r = openRelation(db,"Employee",READ);
Scan s = start_scan(r);
Tuple t;  // current tuple
while ((t = next_tuple(s)) != NULL) {
   char *name = getStrField(t,2);
   printf("%s\n", name);
}
```

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❖ `next_tuple()` Function

Consider the following possible **`Scan`** data structure

```
typedef ScanData *Scan;

typedef struct {
   Relation rel;
   Page     *page;    // Page buffer
   int      curPID;   // current pid
   int      curTID;   // current tid
} ScanData;
```

Assume tuples are indexed 0..**`nTuples(p)-1`**

Assume pages are indexed 0..**`nPages(rel)-1`**

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❖ `next_tuple()` Function (cont)

Implementation of **`Tuple next_tuple(Scan)`** function

```
Tuple next_tuple(Scan s)
{
   if (s->curTID >= nTuples(s->page)-1) {
      // get a new page; exhausted current page
      s->curPID++;
      if (s->curPID >= nPages(s->rel))
         return NULL;
      else {
         s->page = get_page(s->rel, s->curPID);
         s->curTID = -1;
      }
   }
   s->curTID++;
   return get_tuple(s->rel, s->page, s->curTID);
}
```

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❖ Relation Copying

Consider an SQL statement like:

```
create table T as (select * from S);
```

Effectively, copies data from one table to a new table.

Process:

```
make empty relation T
s = start scan of S
while (t = next_tuple(s)) {
    insert tuple t into relation T
}
```

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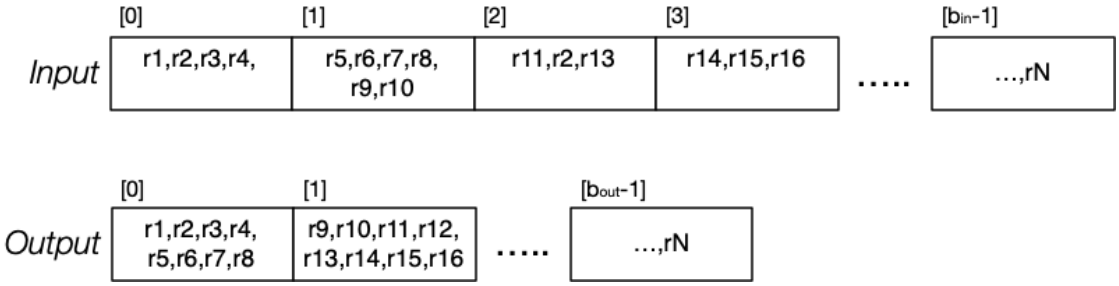
❖ Relation Copying (cont)

It is possible that **T** is smaller than **S**

- may be unused free space in **S** where tuples were removed
- if **T** is built by simple append, will be compact



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❖ Relation Copying (cont)

In terms of existing relation/page/tuple operations:

```
Relation in;       // relation handle (incl. files)
Relation out;      // relation handle (incl. files)
int ipid,opid,tid; // page and record indexes
Record rec;        // current record (tuple)
Page ibuf,obuf;    // input/output file buffers

in = openRelation("S", READ);
out = openRelation("T", NEW|WRITE);
clear(obuf);  opid = 0;
for (ipid = 0; ipid < nPages(in); ipid++) {
    ibuf = get_page(in, ipid);
    for (tid = 0; tid < nTuples(ibuf); tid++) {
        rec = get_record(ibuf, tid);
        if (!hasSpace(obuf,rec)) {
            put_page(out, opid++, obuf);
            clear(obuf);
        }
        insert_record(obuf,rec);
}   }
if (nTuples(obuf) > 0) put_page(out, opid, obuf);
```

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❖ Scanning in PostgreSQL

Scanning defined in: backend/access/heap/heapam.c

Implements iterator data/operations:

- **`HeapScanDesc`** ... struct containing iteration state
- **`scan = heap_beginscan(rel,...,nkeys,keys)`**
- **`tup = heap_getnext(scan, direction)`**
- **`heap_endscan(scan)`** ... frees up **`scan`** struct
- **`res = HeapKeyTest(tuple,...,nkeys,keys)`**
  ... performs **`ScanKeys`** tests on tuple ... is it a result tuple?

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❖ Scanning in PostgreSQL (cont)

```
typedef HeapScanDescData *HeapScanDesc;

typedef struct HeapScanDescData
{
  // scan parameters
  Relation      rs_rd;         // heap relation descriptor
  Snapshot      rs_snapshot;   // snapshot ... tuple visibility
  int           rs_nkeys;      // number of scan keys
  ScanKey       rs_key;        // array of scan key descriptors
  ...
  // state set up at initscan time
  PageNumber    rs_npages;     // number of pages to scan
  PageNumber    rs_startpage;  // page # to start at
  ...
  // scan current state, initally set to invalid
  HeapTupleData rs_ctup;       // current tuple in scan
  PageNumber    rs_cpage;      // current page # in scan
  Buffer        rs_cbuf;       // current buffer in scan
   ...
} HeapScanDescData;
```

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❖ Scanning in other File Structures

Above examples are for heap files

- simple, unordered, maybe indexed, no hashing

Other access file structures in PostgreSQL:

- **`btree`**, **`hash`**, **`gist`**, **`gin`**
- each implements:
  - startscan, getnext, endscan
  - insert, delete (update=delete+insert)
  - other file-specific operators

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