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# COMP9315 14s2 ... Q8

(a) not conclict serializable
    - conflict between T1:R(X) -> T2:W(X)
    - conflict between T2:W(X) -> T1:W(X)
    - precedence graph thus has a cycle

(b) not view serializable
    - supplied schedule
      - first reader is T1
      - T reads initial version of X
      - last writer is T1
    - is it view equivalent to T1 then T2?
      - first reader is T1
      - T1 reads initial version of X
      - last writer is T2
      - not view equivalent
    - is it view equiavalent to T2 then T1?
      - first reader is T1
      - T1 reads version of X written by T2
      - last writer is T1

(c) t0: X(T1) = 1,   X(T2) = 1,   RTS(X) = 0,   WTS(X) = 0
    t1: X(T1) = 1,   X(T2) = 1,   RTS(X) = t1,  WTS(X) = ?
    t2: X(T1) = 10,  X(T2) = 5,   RTS(X) = t1,  WTS(X) = t2
    t3: X(T1) = 10,  X(T2) = 5,   RTS(X) = t1,  WTS(X) = t3

(d) X = 5
    T2 writes X at t2 giving value 5
    T1 writes X at t3 giving value 10
    T2 commits succesfully
    T1 attempts to commit, but sees that X has been modified and
aborts
    (the answer was given away in part (e) of the question)

(e) There will be two versions of X with different timestamps
    Which version should be used is ambiguous for later Txs
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