ISM6218-901 Final Project James Long USF Spring 2020

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Executive Summary

This project is the culmination of effort performed over the span of a 12-week semester. The project incorporates several earlier assignments along with additional work that builds upon and extends those assignments. The work is not organized chronologically due to the required format as specified in Table 1. To minimize confusion, I labeled previous assignments and added introductory comments where appropriate to aid the reader in understanding the sequence of events.

Table 1
Organization and Point Weightings

Topic Area	Description	Points
Database Design	Logical database design, normalization, integrity constraints, design considerations for loading bulk data	30
Query Writing	Queries, a stored procedure	20
Performance Tuning	Btree+ and bitmap indexing, table and column statistics, optimizer modes, table partitioning	25
Other Topics	Loading bulk data via Python	25

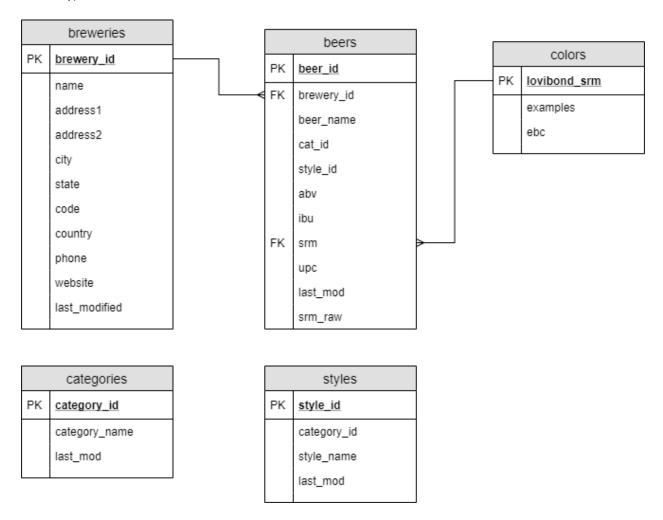
To maximize my learning outcomes, I chose to work alone this semester. Doing so forced me to complete 100% of the work as opposed to 25% in a group of four. While this benefited me, some of my solutions might be less polished than they would otherwise be had I the benefit of a peer group. Of course, this would not be the case if I were an experienced programmer and DBA. However, I am new to both, which is why I wanted to complete all the work for myself. The scope of work that I could undertake alone prevented me from fully exploring every obstacle I encountered. So, I left some exercises to future students of this course (search for "future students"). Hopefully, these will be of interest and provide some benefit to other students.

Of the choices given, I opted to remediate the beersdb schema rather than venture into the unknown. I began by copying the existing schema into my personal schema (DB870). Then I made several corrections before extending the design to accommodate competitions, sponsors, awards, and reviews (including numerical ratings and textual comments). The original data served as the basis for the first round of query and indexing experiments. I then loaded bulk data from the RateBeers dataset using a Python script. This new data served as the basis for the second round of query and performance experiments.

Database Design

Correct the Existing BeerDB Design (Assignment 2)

The following ERD depicts the beerdb schema as it currently exists on CDB9. This is slightly different than the ERD depicted on the assignment page in Canvas (which shows some FK constraints that are missing in reality).



The following SQL transactions were issued to replicate the beerdb tables and recreate the pre-existing key constraints within my personal schema (db870). All other constraints were "NOT NULL" and were preserved during the CTAS transactions.

```
CREATE TABLE
beers
AS
SELECT
*
FROM
```

```
beerdb.beers;
CREATE TABLE
  breweries
AS
  SELECT
  FROM
    beerdb.breweries;
CREATE TABLE
  categories
AS
  SELECT
  FROM
    beerdb.categories;
CREATE TABLE
  colors
AS
  SELECT
 FROM
    beerdb.colors;
CREATE TABLE
  styles
AS
  SELECT
 FROM
    beerdb.styles;
ALTER TABLE
  beers
ADD CONSTRAINT
  beers_pk
PRIMARY KEY(beer_id);
ALTER TABLE
  breweries
ADD CONSTRAINT
  breweries_pk
PRIMARY KEY(brewery_id);
ALTER TABLE
  categories
```

```
ADD CONSTRAINT
 categories_pk
PRIMARY KEY(category_id);
ALTER TABLE
 colors
ADD CONSTRAINT
 colors_pk
PRIMARY KEY(lovibond_srm);
ALTER TABLE
  styles
ADD CONSTRAINT
 styles pk
PRIMARY KEY(style_id);
ALTER TABLE
 beers
ADD CONSTRAINT
 breweries_fk
FOREIGN KEY
 (brewery_id)
REFERENCES
 breweries(brewery_id);
ALTER TABLE
  beers
ADD CONSTRAINT
  colors fk
FOREIGN KEY
 (srm)
REFERENCES
 colors(lovibond_srm);
```

According to Craft Beer & Brewing (https://beerandbrewing.com/tools/color-calculator/), Lovibond is not the same thing as SRM. Using that website, I was able to determine the "lovibond_srm" values in the colors table are SRM values, not Lovibond values. So, the following change was made to correct the attribute name, enable natural joins between the beers and colors tables, and make the key relationship more obvious. The change did not break the existing key relationship.

ALTER TABLE colors
RENAME COLUMN lovibond_srm
TO srm;

The following changes were made to eliminate NULL values in the srm attribute of the beers table and prevent new records with a NULL srm value.

```
INSERT INTO
colors
VALUES
(-1, 'unknown', -1);

UPDATE
beers
SET
srm = -1
WHERE
srm IS NULL;

ALTER TABLE
beers
ADD CONSTRAINT
srm_not_null
CHECK ("SRM" IS NOT NULL);
```

Some records in the beers table had a value of -1 for style_id, but no such value existed in the styles table. Likewise, some records in the beers table had a value of -1 for category_id, but no such value existed in the categories table. I examined the relationship of style_id to category_id in the beers table as follows.

```
SELECT
 *
FROM
 beerdb.beers
WHERE
 cat_id = -1
 AND style id <> -1;
```

No records were returned, which indicated all beers with a category_id of -1 also had a style_id of -1. Therefore, the following changes were made.

```
INSERT INTO
categories
VALUES
(-1, 'unknown', CURRENT_DATE);
INSERT INTO
styles
VALUES
```

```
(-1, -1, 'unknown', CURRENT_DATE);
```

The following foreign key constraints were then implemented to enable referential integrity. All attributes used as foreign keys in any table either already had a CHECK constraint to enforce NOT NULL or had a new such constraint defined as shown above (e.g., beers(srm)).

```
ALTER TABLE
  beers
ADD CONSTRAINT
  styles fk
FOREIGN KEY
  (style_id)
REFERENCES
  styles(style_id);
ALTER TABLE
  styles
ADD CONSTRAINT
  categories_fk
FOREIGN KEY
  (category_id)
REFERENCES
  categories(category_id);
```

At this point, I realized the category_id attribute was stored redundantly in the beers and styles tables. This gives rise to potential data integrity problems:

- The user knows the category_id of a beer but not the style_id. The record is inserted into the beers table with a category_id of something other than -1 and a style_id of -1. Meanwhile, the styles table has only one record with style_id = -1, and the associated category_id is -1. We cannot force the user to also insert a new record into the styles table to define the new category_id / style_id pairing.
- The user knows the style_id of a beer but not the category_id. The user can run a query on the styles table to lookup the category_id associated with the known style_id. However, the user may instead choose to insert the record into the beers table with a style_id of something other than -1 and a category_id of -1. Meanwhile, the styles table has only one record with category_id = -1, and the associated style_id is -1. We cannot force the user to also insert a new record into the styles table to define the new category_id / style_id pairing.

This example illustrates why we normalize tables. The category_id of a given beer can be looked up via a join of the beers table with the styles table and then another join with the categories table. While the performance impact of two joins is obviously not ideal, I think the performance penalty of one or more CHECK constraints to ensure data integrity would be higher (assuming such constraints are even possible). So, I dropped the category_id attribute from the beers table.

```
-- cat_id NOT NULL
```

```
ALTER TABLE
beers
DROP CONSTRAINT
sys_c00108212;

ALTER TABLE
beers
DROP COLUMN
cat_id;
```

Without knowing which attribute was modified, the last_mod attribute is of little use. Furthermore, the attribute's name and format are inconsistent from table to table. Finally, this attribute tells us a fact about another non-key attribute, since the key value cannot change. By contrast, a date_created attribute would describe the key. So, I dropped the last_mod/last_modified attribute from all tables that contained it.

```
ALTER TABLE
beers
DROP COLUMN
last_mod;
```

ALTER TABLE categories DROP COLUMN last_mod;

ALTER TABLE styles DROP COLUMN last_mod;

ALTER TABLE breweries DROP COLUMN last_modified;

The srm_raw attribute in the beers table is redundant with the srm attribute. The values do not all align, but most do. Those that do not align can be proven to have incorrect values in the srm_raw attribute according to the following query cross-referenced with the Craft Beer & Brewing website.

```
select
c.ebc,
b.srm,
b.srm_raw
FROM
beers b
```

```
INNER JOIN colors c
ON b.srm = c.srm

WHERE
b.beer_id = 5885
OR b.beer_id = 5874;

So, I dropped the srm_raw attribute from the beers table.

ALTER TABLE
beers
DROP COLUMN
srm_raw;
```

The upc attribute in the beers table stores a value of 0 for all but six beers. Furthermore, only two values are shared by those six distinct beers. Thus, the upc attribute provides no useful information. So, I dropped the upc column.

```
ALTER TABLE
beers
DROP COLUMN
upc;
```

All remaining attributes in all tables are atomic except for examples in the colors table. So, I created a new table and migrated the examples attribute data. I used the example attribute as the primary key since an example can only belong to a single SRM value, thus requiring example to be unique. The Weissbier value appeared twice, so I did not migrate the instance associated with an SRM of 4, which I determined to be incorrect based on the results of an Internet search.

```
CREATE TABLE
  color_examples (
    example VARCHAR2(50 BYTE),
    srm NUMBER(2,0) NOT NULL,
    CONSTRAINT color examples pk PRIMARY KEY (example),
    CONSTRAINT color_fk FOREIGN KEY (srm) REFERENCES colors(srm)
    );
INSERT ALL
  INTO color examples VALUES ('unknown', -1)
  INTO color_examples VALUES ('Pale Lager', 2)
  INTO color_examples VALUES ('Witbier', 2)
  INTO color_examples VALUES ('Pilsener', 2)
  INTO color examples VALUES ('Berliner', 2)
  INTO color examples VALUES ('Weisse', 2)
  INTO color examples VALUES ('Maibock', 3)
  INTO color examples VALUES ('Blonde Ale', 3)
  INTO color examples VALUES ('American Pale Ale', 6)
```

```
INTO color examples VALUES ('India Pale Ale', 6)
  INTO color_examples VALUES ('Weissbier', 8)
  INTO color examples VALUES ('Saison', 8)
  INTO color_examples VALUES ('English Bitter', 10)
  INTO color examples VALUES ('ESB', 10)
  INTO color examples VALUES ('Biere de Garde', 13)
  INTO color examples VALUES ('Double IPA', 13)
  INTO color_examples VALUES ('Dark Lager', 17)
  INTO color_examples VALUES ('Vienna Lager', 17)
  INTO color examples VALUES ('Marzen', 17)
  INTO color examples VALUES ('Amber Ale', 17)
  INTO color examples VALUES ('Brown Ale', 20)
  INTO color_examples VALUES ('Bock', 20)
  INTO color examples VALUES ('Dunkel', 20)
  INTO color examples VALUES ('Dunkelweizen', 20)
  INTO color_examples VALUES ('Irish Dry Stout', 24)
  INTO color examples VALUES ('Doppelbock', 24)
  INTO color examples VALUES ('Porter', 24)
  INTO color_examples VALUES ('Stout', 29)
  INTO color_examples VALUES ('Foreign Stout', 35)
  INTO color examples VALUES ('Baltic Porter', 35)
  INTO color examples VALUES ('Imperial Stout', 40)
SELECT * FROM dual;
ALTER TABLE
  colors
DROP COLUMN
  examples;
```

Brewers can own and operate multiple breweries. For example, the Molson Coors Beverage Company has seven breweries that offer tours (and possibly others that do not offer tours). Furthermore, a given beer can be produced at multiple breweries. Therefore, it is clear the intent of the breweries table is to track brewers, not breweries. So, I renamed the table brewers and renamed the brewery_id attribute brewer id in the brewers and beers tables.

```
ALTER TABLE
breweries
RENAME TO
brewers;

ALTER TABLE
brewers
RENAME COLUMN
brewery_id
TO
brewer id;
```

ALTER TABLE
beers
RENAME COLUMN
brewery_id
TO
brewer_id;

Given the high number of attributes named "name" that will be included in the extended design, I decided to rename all such attributes to be more descriptive. Only one attribute was affected in the current database schema.

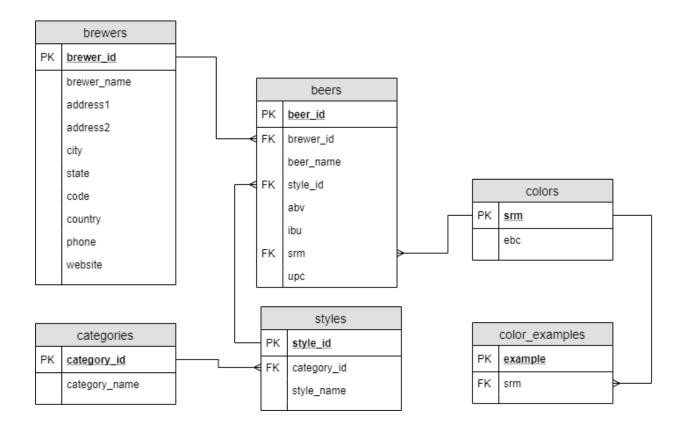
ALTER TABLE
brewers
RENAME COLUMN
name
TO
brewer_name;

A beer's name might not provide a unique identifier. For national or international beers, the name might suffice depending on state, national, and international trademark registrations. But this cannot be guaranteed. And for locally brewed beers, uniqueness is almost certainly not guaranteed. Thus, a composite natural key or synthetic key is required. As part of the extended design, the beers key must be used in the associative table beer_reviews, in the comments table, and in the beer_comp table. A composite key would increase storage requirements and could impact join performance. So, the short, numeric (synthetic) key that was already in place was kept.

Similarly, a brewer's name might not provide a unique identifier for the same reasons that a beer's name might not. This is less likely since the number of brewers is significantly lower than the number of beers. However, brewers periodically merge, reorganize, rebrand, and divest. Any of those events could precipitate a name change. Thus, the short, numeric (synthetic) key that was already in place was kept.

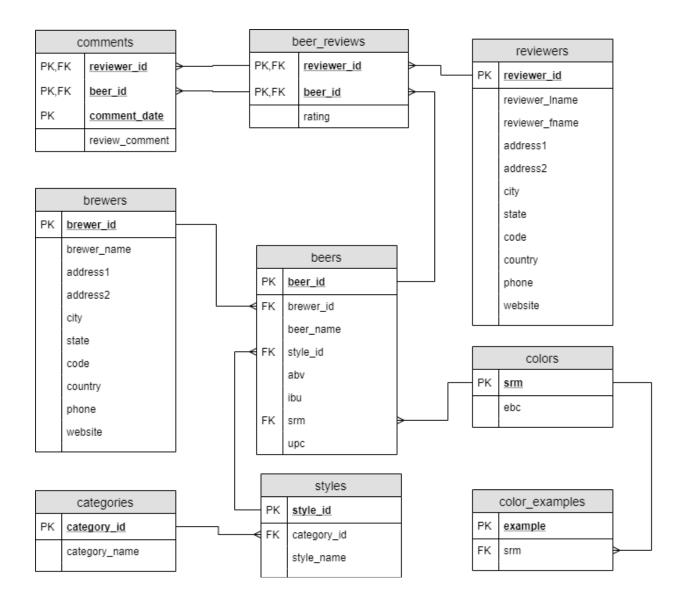
The style and category names are unique, but they are text strings up to 50 bytes long. Moreover, it is possible that additional names added in the future could be even longer. Therefore, using either name as a primary key could impact performance. Also, it is theoretically possible (but not likely) for a style or category name to change. So, the short, numeric (synthetic) keys that were already in place were kept.

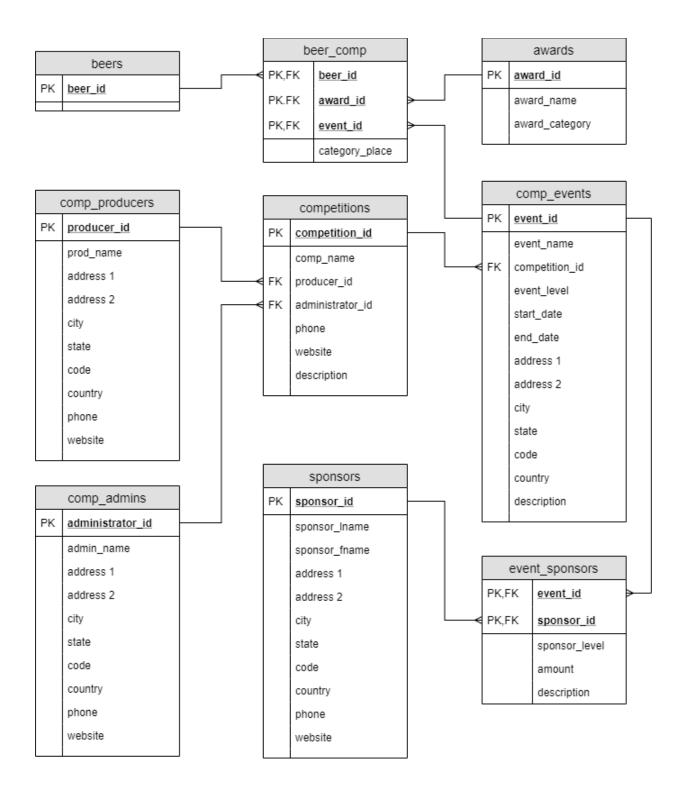
The resulting schema is in 3rd normal form. The updated ERD follows.



Update the Design to Fulfill the New Requirements (Assignment 2)

Next, I updated the design to incorporate the new requirements. The extended ERD follows. Other extension ideas for future students include tracking event judges and their credentials, tracking breweries and associating each with a brewer (especially those breweries that offer public tours), and tracking event revenues from sources other than sponsors (e.g., ticket sales).





The beer_reviews table is an associative table, so the choice for primary key was made using conventional wisdom. The chosen rating is based on the Beer Connoisseur website (https://beerconnoisseur.com/how-we-score).

Beer reviewers are individuals, so the beer reviewer's name will not provide a unique identifier. Thus, a composite natural key or synthetic key is required. The reviewers key must be used in the associative

table beer_reviews and in the comments table, which would increase storage requirements and could impact join performance. So, a short, numeric (synthetic) key was chosen. Some reviewers will be professional beer critics, so a website attribute (nullable) is included in the reviewers table.

Comments are tied to a specific review, so the obvious choice for primary key includes the composite primary key of the beer_reviews table plus the date attribute of the comment (which includes both date and time, so multiple comments can be submitted per day). Comments could be stored in the beer_reviews table, but the rating attribute (a small integer) would be redundantly stored. The alternative is to store comments in a separate table referencing the reviewer_id and beer_id in the beer_reviews table. This results in even more redundant data, but the only redundant data are key attributes. While this proper 3rd normal form, it might make sense (when implementing this design) to denormalize the comments table into the beer_reviews table (with date as part of the primary key) to reduce the total amount of redundant data and preclude the need for an extra join operation. However, this would introduce the possibility of inconsistent rating values across related records.

A competition's name might not be unique. Small, local competitions might be named after the city or county in which they take place (e.g., Springfield Beer Competition). Many cities and counties share the same name. Therefore, name and producer_id would be needed to form a natural key. However, that would require a composite foreign key in the comp_events table, which would increase storage requirements and could impact join performance. Also, a competition's name could change over time. So, a short, numeric (synthetic) key was chosen.

If two or more companies jointly produce a competition, I assume they form a Limited Liability Partnership (LLP) or other similar legal entity to facilitate the joint venture. Thus, each competition has a single producer. Of course, a producer may produce many different competitions.

A competition may be administered by a professional organization or by the producer. Either way, I assume a single administrator is associated with each competition. For producer-administered competitions, the administrator_id attribute is left empty (NULL).

Competition producers and administrators are assumed to be organizations in most (if not all) cases. As stated previously about brewers, organizations can be expected to change their names over time. Therefore, a short, numeric (synthetic) key was chosen for the comp_producers and comp_admins tables.

Events need to be stored separately from competitions because event names, dates, locations, sponsors, etc. can vary from event to event for a given competition. An event's name might be reused from year to year (e.g., Central Valley Annual Local Brew-off), so event names might not be unique. Thus, a natural key would need to be a composite of name and something like start_date. But the many-to-many relationship between comp_events and sponsors would require the composite key to be used in the associative table event_sponsors, which would increase storage requirements and could impact join performance. So, a short, numeric (synthetic) key was chosen. A description attribute (nullable) is included in case additional details are desirable (e.g., "25th anniversary of the competition"). I assume events do not have a website. Instead, they rely on the competition's website for all marketing and registrations. The event_level attribute should store a limited set of values such as local, state, regional, national, and international. To ensure only these values are used, a separate lookup table named event_levels could be created with attributes event_level_id and event_level_name (similar to the categories, styles, color examples, and colors tables). Additionally, read-only access restrictions could

be placed on all such tables to prevent proliferation of lookup values; only admins would have readwrite access. These options are left to future classes to implement.

The event_sponsors table is an associative table, so the choice for primary key was made using conventional wisdom. Event sponsors may be grouped into levels (e.g., Platinum, Gold, etc.), so a sponsor_level attribute (nullable) was created to capture this information. An amount attribute (nullable) was also created to capture the dollar amount (or dollar value equivalent) of the sponsorship. A description attribute (nullable) was also created to capture any unique information that may be of value (e.g., "advertiser" or "loaned tables and chairs free of charge").

A sponsor's name might not provide a unique identifier. Sponsors might be individuals for small, local competition events. Thus, a composite natural key or synthetic key is required. The sponsors key must be used in the associative table event_sponsors, and corporate sponsors may change names over time. For the reasons stated previously, a short, numeric (synthetic) key was chosen. For non-individual sponsors, I assume the organization's name will be stored in the sponsor_Iname attribute, and the sponsor_fname attribute will be empty.

The beer comp table is an associative table, so the choice for primary key was made using conventional wisdom. A beer may win zero, one, or many awards across categories and/or within a single category at a single competition event. My design accommodates all these scenarios. The beer_comp table includes award id as part of the primary key. If a beer does not win an award, an award id value of -1 is used, which has an award_name of "no award" and an award_category of "no award category". For such a beer, the category_place attribute in the beer_comp table represents the beer's place within the primary category in which the beer competed. For winning beers, the category_place attribute represents the beer's placement in the associated award_category. If a beer wins multiple awards in a single category, the award id and award name will differ but the award category and category place will be the same. If a beer wins multiple awards across categories, the award_id, award_name, award category, and category place will all differ. The award grade (e.g., gold, silver, 1st place, 2nd place, etc.) is assumed to be reflected in the award's name. The award level (e.g., local, state, regional, etc.) is captured in the comp events table instead of the awards table because the level applies to all awards at the event (i.e., level tell us a fact about the event). The level is not captured in the competitions table because a competition may hold numerous events at different levels (local, state, regional, etc.) leading up to a national or international final event.

The award_name attribute is likely to hold many instances of "gold medal", "first place", and similar such names. Likewise, award_category is likely to contain many common values since different competitions are likely to categorize beers in a similar manner. Therefore, a short, numeric (synthetic) key was chosen. This also improves join performance with the beer_comp table. The award_category is set to NOT NULL. For an award that does not pertain to a category, this attribute can be set to "overall".

All name attributes are set to NOT NULL.

In newly created tables, all primary keys that can be auto-generated are auto-generated. In pre-existing tables, no changes were made to the way primary keys are generated because such a change would require dropping existing primary key columns and re-creating them as ID columns.

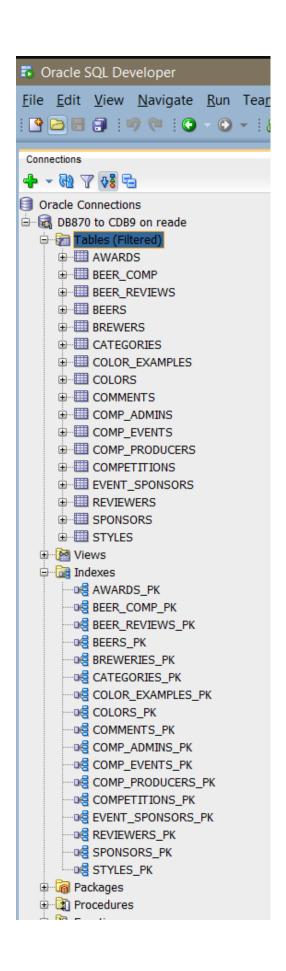
The design was implemented with the following code.

```
CREATE TABLE
  reviewers (
    reviewer id NUMBER GENERATED by default on null as IDENTITY,
    reviewer_Iname VARCHAR2(50 BYTE) NOT NULL,
    reviewer fname VARCHAR2(50 BYTE) NOT NULL,
    address1 VARCHAR2(255 BYTE),
    address2 VARCHAR2(255 BYTE),
    city VARCHAR2(255 BYTE),
    state VARCHAR2(255 BYTE),
    code VARCHAR2(25 BYTE),
    country VARCHAR2(255 BYTE),
    phone VARCHAR2(50 BYTE),
    website VARCHAR2(255 BYTE),
    CONSTRAINT reviewers pk PRIMARY KEY (reviewer id)
    );
CREATE TABLE
  beer reviews (
    reviewer id NUMBER,
    beer_id NUMBER,
    rating NUMBER(3,0) NOT NULL,
    CONSTRAINT beer reviews pk PRIMARY KEY (reviewer id, beer id),
    CONSTRAINT reviewers_fk FOREIGN KEY (reviewer_id) REFERENCES reviewers(reviewer_id),
    CONSTRAINT beers_fk FOREIGN KEY (beer_id) REFERENCES beers(beer_id)
   );
CREATE TABLE
  comments (
    reviewer id NUMBER,
    beer id NUMBER,
    comment date DATE,
    review_comment VARCHAR2(255 BYTE) NOT NULL,
    CONSTRAINT comments pk PRIMARY KEY (reviewer id, beer id, comment date),
    CONSTRAINT beer_reviews_fk FOREIGN KEY (reviewer_id, beer_id) REFERENCES
       beer reviews(reviewer id, beer id)
   );
CREATE TABLE
  comp producers (
    producer id NUMBER GENERATED by default on null as IDENTITY,
    prod_name VARCHAR2(255 BYTE) NOT NULL,
    address1 VARCHAR2(255 BYTE),
    address2 VARCHAR2(255 BYTE),
    city VARCHAR2(255 BYTE),
    state VARCHAR2(255 BYTE),
    code VARCHAR2(25 BYTE),
    country VARCHAR2(255 BYTE),
    phone VARCHAR2(50 BYTE),
```

```
website VARCHAR2(255 BYTE),
    CONSTRAINT comp_producers_pk PRIMARY KEY (producer_id)
   );
CREATE TABLE
  comp admins (
    administrator id NUMBER GENERATED by default on null as IDENTITY,
    admin_name VARCHAR2(255 BYTE) NOT NULL,
    address1 VARCHAR2(255 BYTE),
    address2 VARCHAR2(255 BYTE),
    city VARCHAR2(255 BYTE),
    state VARCHAR2(255 BYTE),
    code VARCHAR2(25 BYTE),
    country VARCHAR2(255 BYTE),
    phone VARCHAR2(50 BYTE),
    website VARCHAR2(255 BYTE),
    CONSTRAINT comp admins pk PRIMARY KEY (administrator id)
    );
CREATE TABLE
  competitions (
    competition id NUMBER GENERATED by default on null as IDENTITY,
    comp_name VARCHAR2(255 BYTE) NOT NULL,
    producer id NUMBER,
    administrator id NUMBER,
    phone VARCHAR2(50 BYTE),
    website VARCHAR2(255 BYTE),
    description VARCHAR2(255 BYTE),
    CONSTRAINT competitions pk PRIMARY KEY (competition id),
    CONSTRAINT comp producers fk FOREIGN KEY (producer id) REFERENCES
      comp producers(producer id),
    CONSTRAINT comp_admins_fk FOREIGN KEY (administrator_id) REFERENCES
     comp admins(administrator id)
   );
CREATE TABLE
  comp events (
    event id NUMBER GENERATED by default on null as IDENTITY,
    event name VARCHAR2(255 BYTE) NOT NULL,
    competition id NUMBER,
    event_level VARCHAR2(255 BYTE) NOT NULL,
    start date DATE NOT NULL,
    end_date DATE NOT NULL,
    address1 VARCHAR2(255 BYTE),
    address2 VARCHAR2(255 BYTE),
    city VARCHAR2(255 BYTE),
    state VARCHAR2(255 BYTE),
    code VARCHAR2(25 BYTE),
```

```
country VARCHAR2(255 BYTE),
    description VARCHAR2(255 BYTE),
    CONSTRAINT comp events pk PRIMARY KEY (event id),
    CONSTRAINT competitions fk FOREIGN KEY (competition id) REFERENCES
      competitions(competition id)
   );
CREATE TABLE
  sponsors (
    sponsor id NUMBER GENERATED by default on null as IDENTITY,
    sponsor Iname VARCHAR2(255 BYTE) NOT NULL,
    sponsor fname VARCHAR2(255 BYTE),
    address1 VARCHAR2(255 BYTE),
    address2 VARCHAR2(255 BYTE),
    city VARCHAR2(255 BYTE),
    state VARCHAR2(255 BYTE),
    code VARCHAR2(25 BYTE),
    country VARCHAR2(255 BYTE),
    phone VARCHAR2(50 BYTE),
    website VARCHAR2(255 BYTE),
    CONSTRAINT sponsors pk PRIMARY KEY (sponsor id)
    );
CREATE TABLE
  event_sponsors (
    event id NUMBER,
    sponsor id NUMBER,
    sponsor level VARCHAR2(255 BYTE),
    amount NUMBER,
    description VARCHAR2(255 BYTE),
    CONSTRAINT event_sponsors_pk PRIMARY KEY (event_id, sponsor_id),
    CONSTRAINT comp_events_fk FOREIGN KEY (event_id) REFERENCES comp_events(event_id),
    CONSTRAINT sponsors fk FOREIGN KEY (sponsor id) REFERENCES sponsors(sponsor id)
    );
CREATE TABLE
  awards (
    award id NUMBER GENERATED by default on null as IDENTITY,
    award name VARCHAR2(255 BYTE) NOT NULL,
    award category VARCHAR2(255 BYTE) NOT NULL,
    CONSTRAINT awards_pk PRIMARY KEY (award_id)
    );
ALTER TABLE
  awards
DISABLE CONSTRAINT
  awards_pk;
```

```
INSERT INTO
  awards
VALUES
 (-1, 'no award', 'no award category');
ALTER TABLE
  awards
ENABLE CONSTRAINT
  awards_pk;
CREATE TABLE
  beer_comp (
    beer_id NUMBER,
    award_id NUMBER,
    event_id NUMBER,
    category_place NUMBER,
    CONSTRAINT beer_comp_pk PRIMARY KEY (beer_id, award_id, event_id),
    CONSTRAINT beer_fk FOREIGN KEY (beer_id) REFERENCES beers(beer_id),
    CONSTRAINT awards_fk FOREIGN KEY (award_id) REFERENCES awards(award_id),
    CONSTRAINT comp_event_fk FOREIGN KEY (event_id) REFERENCES comp_events(event_id)
    );
```



Design Considerations for Loading Bulk Data

Initially, I used the data pre-loaded in the beerdb database. In various parts of this project, I also loaded data manually with SQL INSERT statements to facilitate experiments. However, that solution proved inadequate in some respects. So, between assignments three and four, I loaded data from the Stanford SNAP / RateBeer dataset (available at https://www.ratebeer.com/api.asp) into the beers, brewers, styles, reviewers, beer_reviews, and comments tables. Basic statistics about the RateBeer dataset are available at https://snap.stanford.edu/data/web-ratebeer.html and shown below.

Number of reviews	2,924,127
Number of users	40,213
Number of beers	110,419
Users with > 50 reviews	4,798
Median # of words per review	54
Timespan	Apr 2000 - Nov 2011

A sample of the data provided for each beer follows.

beer/name: John Harvards Simcoe IPA

beer/beerld: 63836 beer/brewerld: 8481 beer/ABV: 5.4

beer/style: India Pale Ale (IPA)

review/appearance: 4/5
review/aroma: 6/10
review/palate: 3/5
review/taste: 6/10
review/overall: 13/20
review/time: 1157587200
review/profileName: hopdog

review/text: On tap at the Springfield, PA location. Poured a deep and cloudy orange (almost a copper) color with a small sized off white head. Aromas or oranges and all around citric. Tastes of oranges, light caramel and a very light grapefruit finish. I too would not believe the 80+ IBUs - I found this one to have a very light bitterness with a medium sweetness to it. Light lacing left on the glass.

Since beer names are not guaranteed to be unique, some of the provided beer names might have been duplicates of beer names pre-populated in the beers table. Also, the beer_name attribute in the beers table did not have a pre-defined UNIQUE constraint, which indicated duplicate beer names might have already existed in the beers table. However, each (brewer_id, beer_name) tuple should be unique. So, I attempted to implement that constraint on the beers table. Eight duplicates were found, so I identified and removed them as follows. The duplicate records with the least amount of valid data were deleted.

```
brewer_id,
  beer_name,
  COUNT(*)
FROM
  beers
GROUP BY
  brewer_id,
  beer_name
  HAVING COUNT(*) > 1
ORDER BY
  brewer_id;
SELECT
  b.beer_id,
  b.brewer_id,
  br.brewer_name,
  br.country,
  b.beer_name,
  b.style_id,
  b.abv,
  b.ibu,
  b.srm,
  b.upc
FROM
  beers b
  INNER JOIN brewers br
    ON b.brewer_id=br.brewer_id
WHERE
  b.brewer_id IN (49,236,1391,1417)
  AND b.beer_name IN ('Guinness Draught','Avalanche Amber','Dos Perros','Hefeweizen',
                      'Onward Stout', 'Pale Ale', 'Sly Rye Porter', 'Watershed IPA')
ORDER BY
  b.beer_name;
DELETE FROM
  beers
WHERE
  beer_id IN (3721,4457,4473,4462,4460,4459,4458,5896);
Then I implemented the UNIQUE constraint on beers(brewer_id, beer_name) as follows.
ALTER TABLE
  beers
ADD CONSTRAINT
  brewer beer unq
UNIQUE (brewer_id, beer_name);
```

Some of the provided beer names contained "(" and ")" representing "(" and ")" respectively. Other encodings may have been present as well. I did not convert these strings. Cleaning up the beer_names values was left as an exercise for future students.

The RateBeer dataset provided a brewer_id field but no other information about the brewers. Nonetheless, the provided brewer_id field served to prevent the conflation of distinct beers sharing a duplicate name. Therefore, I removed the NOT NULL constraint on brewers(brewer_name). In the future, RateBeer.com might make available more information about brewers, at which time the brewer table can be populated by matching records on brewer_id. Afterward, the NOT NULL constraint can be re-implemented on brewer_name. My load script was written such that any records in the RateBeer dataset that did not have a brewerld value did not get inserted into the database. However, my full scan of the file indicated no such records existed.

```
-- brewer_name NOT NULL
ALTER TABLE
  brewers
DROP CONSTRAINT
  sys_c00108215;
```

Some of the provided beerId values could have been duplicates of pre-populated beer_id values in the beers table. The RateBeer dataset claimed to have less than 111,000 unique beerIds. However, there was no way to know if the beerIds were numbered sequentially and contiguously starting at 0. To resolve this potential issue, I temporarily disabled the beer_id foreign key constraint on the beer_comp and beer_reviews tables. Then I updated all beer_id values in the beers table by adding 500,000 to each record. Then I repeated the update on the beer_comp table. Finally, I re-enabled the beer_id foreign key constraint on the beer_comp and beer_reviews tables. My solution avoided collisions as new records were added and made it very easy to distinguish new from pre-existing records using the beer_id range. In the event a RateBeer record had no beerId, the record was not inserted into the database. However, my full scan of the file indicated no such records existed.

```
ALTER TABLE
beer_comp
DISABLE CONSTRAINT
beer_fk;

ALTER TABLE
beer_reviews
DISABLE CONSTRAINT
beers_fk;

UPDATE
beers
SET
beer_id = (beer_id + 500000);

UPDATE
```

```
beer_comp
SET
  beer_id = (beer_id + 500000);

ALTER TABLE
  beer_reviews
ENABLE CONSTRAINT
  beers_fk;

ALTER TABLE
  beer_comp
ENABLE CONSTRAINT
  beer_fk;
```

Similarly, some of the provided brewerld values could have been duplicates of pre-populated brewer_id values in the brewers table. I implemented the same solution as for beer_ids.

```
ALTER TABLE
  beers
RENAME CONSTRAINT
  breweries_fk
TO
  brewers_fk;
ALTER TABLE
  beers
DISABLE CONSTRAINT
  brewers_fk;
UPDATE
  brewers
SET
  brewer_id = (brewer_id + 500000);
UPDATE
  beers
SET
  brewer_id = (brewer_id + 500000);
ALTER TABLE
  beers
ENABLE CONSTRAINT
  brewers_fk;
```

Some of the provided abv values were "-" indicating an unknown abv. Since unknown abv values were already represented in the beers table as 0, I converted all "-" values to 0 upon insertion into the beers table.

The RateBeer dataset provided a style field containing a style name but no category information. Furthermore, many (perhaps all) of the provided style names differed from the style names already in the beerdb styles table. And for any matching values, there was no way to know if the category mapping pre-established in beerdb was valid for the new records (since style names can be legitimately duplicated across categories). So, I inserted all RateBeer style names as new records in the styles table with new (generated) style_ids mapped to a category_id of -1. The styles_id attribute proved to be too small, so I had to increase the size.

```
ALTER TABLE
styles
MODIFY
style_id NUMBER;
```

The RateBeer dataset provided review data for each beer including separate ratings for appearance, aroma, palate, taste, and overall. To complicate matters, each rating used a different scale. So, I normalized the rating scales to 100 and modified the beer_reviews table as follows. None of the new attributes was given a NOT NULL constraint, but the previously defined NOT NULL constraint on rating (now overall_rating) was retained. In the event a RateBeer record had no overall_rating, the review and comment portions of the record were not inserted into the database. However, my full scan of the file indicated no such records existed.

```
ALTER TABLE
beer_reviews
RENAME COLUMN rating
TO overall_rating;

ALTER TABLE
beer_reviews
ADD (
appearance_rating NUMBER(3,0),
aroma_rating NUMBER(3,0),
palate_rating NUMBER(3,0),
taste_rating NUMBER(3,0)
);
```

The time stamp of each review was a large integer representing a date between April 2000 and November 2011. Decoding and transforming those values into a more user-friendly format was left as an exercise for future students. I loaded this field unmodified into comments(comment_date). To do so, I first had to modify the data type of the comment_date attribute as follows.

```
ALTER TABLE comments
```

```
MODIFY comment_date NUMBER;
```

In the event a RateBeer record had no comment_date, the record was not inserted into the comments table, but other portions of the record were inserted into appropriate tables. However, my full scan of the file indicated no such records existed. By contrast, many RateBeer records existed with no data in the review/text field. Obviously, such records were not inserted into the comments table.

Most of the review comments exceeded 255 bytes, and some exceeded 4000 bytes. So, I increased the size of comments(review_comment) by changing the data type to LONG.

```
ALTER TABLE
comments
MODIFY
review_comment LONG;

ALTER TABLE
comments
MODIFY
(review_comment NOT NULL);
```

The reviewers were identified by a profileName field instead of first and last name. I mapped the profileName field into reviewers(reviewer_Iname) and dropped the NOT NULL constraint on reviewers(reviewer_fname). In the event a RateBeer record had no profileName, the review and comment portions of the record were not inserted into the database. However, my full scan of the file indicated no such records existed. Since real names can be duplicated (even first/last pairings), I did not implement any UNIQUE constraints on any combination of the name attributes.

-- fname NOT NULL ALTER TABLE reviewersDROP CONSTRAINT sys c00108390;

All records inserted into the beers table were given a value of -1 (unknown) for srm. Since a value of 0 was already used in the beers table to represent an unknown ibu value, I assigned an ibu value of 0 to all inserted records.

I did not log SQL errors during the load process. With so many records, losing a few due to errors is not a problem. Programatically identifying any missing records and inserting them into the database might be a good exercise for future students. Similarly, some records in the RateBeer dataset contained invalid characters. I replaced those with "?" prior to insertion into the database. Programatically finding and replacing those characters with the correct characters might be an interesting exercise for future students. However, deciphering what the correct characters are would likely entail a manual review of each affected record.

Prior to loading the new dataset, I made all indexes unusable on the affected tables to expedite the load process. This action resulted in errors related to primary key indexes. I did not have high confidence in the RateBeer dataset or my code, so I did not want to disable integrity constraints. Thus, I re-enabled all indexes associated with primary or foreign key constraints on the affected tables. The following indexes remained unusable until after the load completed.

```
ALTER INDEX beer_name_btree UNUSABLE;
ALTER INDEX brewer_name_btree UNUSABLE;
ALTER INDEX style_name_btree UNUSABLE;
```

I also renamed the breweries_pk index to match the new name of the table.

```
ALTER INDEX breweries_pk RENAME TO brewers_pk;
```

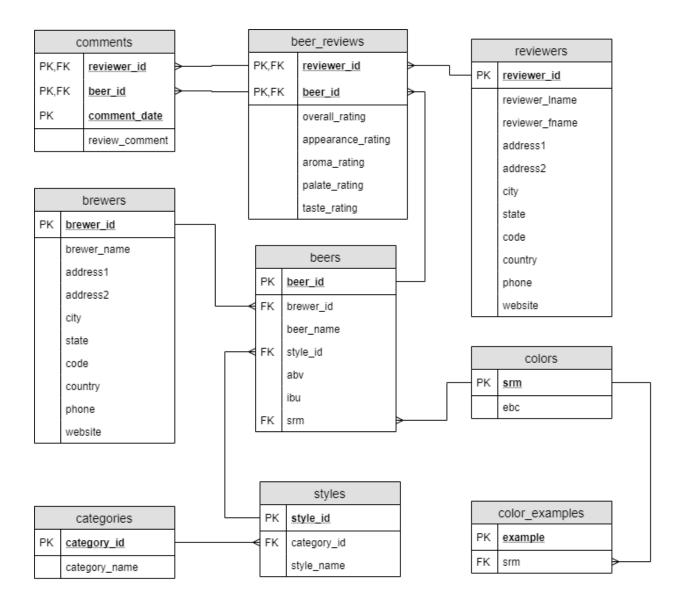
After the load process, I refreshed the statistics for the impacted tables.

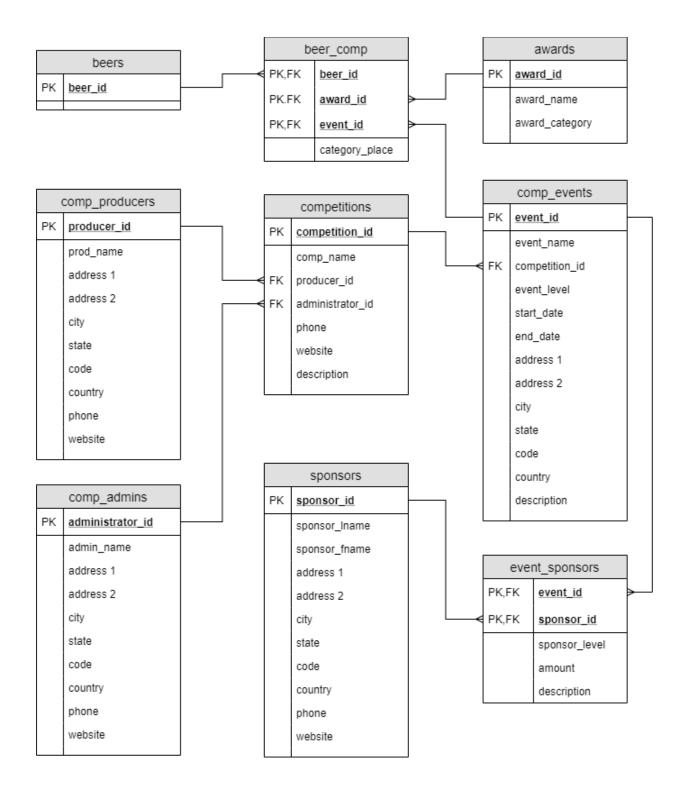
```
EXECUTE DBMS_STATS.GATHER_TABLE_STATS('db870', 'beer_reviews'); EXECUTE DBMS_STATS.GATHER_TABLE_STATS('db870', 'comments'); EXECUTE DBMS_STATS.GATHER_TABLE_STATS('db870', 'reviewers'); EXECUTE DBMS_STATS.GATHER_TABLE_STATS('db870', 'brewers'); EXECUTE DBMS_STATS.GATHER_TABLE_STATS('db870', 'styles'); EXECUTE DBMS_STATS.GATHER_TABLE_STATS('db870', 'beers');
```

Then I rebuilt the unusable indexes.

```
ALTER INDEX beer_name_btree REBUILD;
ALTER INDEX brewer_name_btree REBUILD;
ALTER INDEX style_name_btree REBUILD;
```

The updated ERD follows.





Query Writing

Write Two Interesting Queries (Assignment 2)

1. For each level of competition (local, state, etc.), list the top five competition producers in terms of sponsorship revenue. Sort results by event level alphabetically, and from highest to lowest revenue within each level.

```
-- insert dummy data to validate query results
INSERT INTO
       comp_producers (prod_name)
VALUES
       ('producer1');
INSERT INTO
       comp producers (prod name)
VALUES
       ('producer2');
INSERT INTO
       comp_producers (prod_name)
VALUES
       ('producer3');
INSERT INTO
       comp producers (prod name)
VALUES
       ('producer4');
INSERT INTO
       comp_producers (prod_name)
VALUES
       ('producer5');
INSERT INTO
       comp_producers (prod_name)
VALUES
       ('producer6');
INSERT INTO
       competitions (comp_name, producer_id)
VALUES
       ('comp1', 1);
INSERT INTO
       competitions (comp name, producer id)
VALUES
       ('comp2', 4);
INSERT INTO
       competitions (comp_name, producer_id)
VALUES
       ('comp3', 5);
INSERT INTO
       competitions (comp_name, producer_id)
```

```
VALUES
       ('comp4', 6);
INSERT INTO
       competitions (comp_name, producer_id)
VALUES
       ('comp5', 7);
INSERT INTO
       competitions (comp_name, producer_id)
VALUES
       ('comp6', 8);
INSERT INTO
       comp_events (event_name, competition_id, event_level, start_date, end_date)
VALUES
       ('event1', 1, 'local', CURRENT_DATE, CURRENT_DATE);
INSERT INTO
       comp_events (event_name, competition_id, event_level, start_date, end_date)
VALUES
       ('event2', 1, 'local', CURRENT_DATE, CURRENT_DATE);
INSERT INTO
       comp_events (event_name, competition_id, event_level, start_date, end_date)
VALUES
       ('event3', 4, 'local', CURRENT_DATE, CURRENT_DATE);
INSERT INTO
       comp_events (event_name, competition_id, event_level, start_date, end_date)
VALUES
       ('event4', 4, 'local', CURRENT_DATE, CURRENT_DATE);
INSERT INTO
       comp_events (event_name, competition_id, event_level, start_date, end_date)
VALUES
       ('event5', 5, 'state', CURRENT_DATE, CURRENT_DATE);
INSERT INTO
       comp_events (event_name, competition_id, event_level, start_date, end_date)
VALUES
       ('event6', 5, 'state', CURRENT DATE, CURRENT DATE);
INSERT INTO
       comp events (event name, competition id, event level, start date, end date)
VALUES
       ('event7', 6, 'state', CURRENT DATE, CURRENT DATE);
INSERT INTO
       comp_events (event_name, competition_id, event_level, start_date, end_date)
VALUES
       ('event8', 6, 'state', CURRENT_DATE, CURRENT_DATE);
INSERT INTO
       comp_events (event_name, competition_id, event_level, start_date, end_date)
VALUES
       ('event9', 2, 'regional', CURRENT_DATE, CURRENT_DATE);
INSERT INTO
```

```
comp_events (event_name, competition_id, event_level, start_date, end_date)
VALUES
       ('event10', 2, 'regional', CURRENT_DATE, CURRENT_DATE);
INSERT INTO
       comp_events (event_name, competition_id, event_level, start_date, end_date)
VALUES
       ('event11', 3, 'regional', CURRENT_DATE, CURRENT_DATE);
INSERT INTO
       comp_events (event_name, competition_id, event_level, start_date, end_date)
VALUES
       ('event12', 3, 'regional', CURRENT_DATE, CURRENT_DATE);
INSERT INTO
       sponsors (sponsor_Iname)
VALUES
       ('corp_sponsor1');
INSERT INTO
       event_sponsors (event_id, sponsor_id, amount)
VALUES
       (2, 1, 100);
INSERT INTO
       event_sponsors (event_id, sponsor_id, amount)
VALUES
       (3, 1, 150);
INSERT INTO
       event_sponsors (event_id, sponsor_id, amount)
VALUES
       (4, 1, 200);
INSERT INTO
       event_sponsors (event_id, sponsor_id, amount)
VALUES
       (5, 1, 250);
INSERT INTO
       event sponsors (event id, sponsor id, amount)
VALUES
       (6, 1, 300);
INSERT INTO
       event_sponsors (event_id, sponsor_id, amount)
VALUES
       (7, 1, 350);
INSERT INTO
       event_sponsors (event_id, sponsor_id, amount)
VALUES
       (8, 1, 400);
INSERT INTO
       event_sponsors (event_id, sponsor_id, amount)
VALUES
```

```
(9, 1, 450);
INSERT INTO
       event_sponsors (event_id, sponsor_id, amount)
VALUES
       (14, 1, 500);
INSERT INTO
       event_sponsors (event_id, sponsor_id, amount)
VALUES
       (15, 1, 550);
INSERT INTO
       event_sponsors (event_id, sponsor_id, amount)
VALUES
       (16, 1, 600);
INSERT INTO
       event_sponsors (event_id, sponsor_id, amount)
VALUES
       (17, 1, 650);
-- run query
(SELECT
  ce.event level,
  cp.prod name as producer,
  SUM(es.amount) as revenue
FROM
  comp_producers cp
  INNER JOIN competitions c
    ON cp.producer_id = c.producer_id
  INNER JOIN comp events ce
    ON c.competition_id = ce.competition_id
  INNER JOIN event_sponsors es
    ON ce.event_id = es.event_id
WHERE
  ce.event_level LIKE 'local'
GROUP BY
  ce.event level,
  cp.prod name
FETCH FIRST 5 ROWS ONLY)
UNION
(SELECT
  ce.event level,
  cp.prod_name as producer,
  SUM(es.amount) as revenue
FROM
  comp producers cp
  INNER JOIN competitions c
    ON cp.producer id = c.producer id
  INNER JOIN comp events ce
    ON c.competition id = ce.competition id
```

```
INNER JOIN event_sponsors es
    ON ce.event_id = es.event_id
WHERE
  ce.event_level LIKE 'state'
GROUP BY
  ce.event_level,
  cp.prod_name
FETCH FIRST 5 ROWS ONLY)
UNION
(SELECT
  ce.event_level,
  cp.prod name as producer,
  SUM(es.amount) as revenue
FROM
  comp_producers cp
  INNER JOIN competitions c
    ON cp.producer_id = c.producer_id
  INNER JOIN comp_events ce
    ON c.competition_id = ce.competition_id
  INNER JOIN event_sponsors es
    ON ce.event id = es.event id
WHERE
  ce.event_level LIKE 'regional'
GROUP BY
  ce.event_level,
  cp.prod_name
FETCH FIRST 5 ROWS ONLY)
UNION
(SELECT
  ce.event level,
  cp.prod_name as producer,
  SUM(es.amount) as revenue
FROM
  comp_producers cp
  INNER JOIN competitions c
    ON cp.producer_id = c.producer_id
  INNER JOIN comp_events ce
    ON c.competition_id = ce.competition_id
  INNER JOIN event sponsors es
    ON ce.event_id = es.event_id
WHERE
  ce.event_level LIKE 'national'
GROUP BY
  ce.event level,
  cp.prod name
FETCH FIRST 5 ROWS ONLY)
UNION
(SELECT
```

```
ce.event level,
  cp.prod_name as producer,
  SUM(es.amount) as revenue
FROM
  comp producers cp
  INNER JOIN competitions c
    ON cp.producer_id = c.producer_id
  INNER JOIN comp_events ce
    ON c.competition_id = ce.competition_id
  INNER JOIN event sponsors es
    ON ce.event_id = es.event_id
WHERE
  ce.event_level LIKE 'international'
GROUP BY
  ce.event level,
  cp.prod_name
FETCH FIRST 5 ROWS ONLY)
UNION
SELECT
  null,
  null,
  null
FROM
  dual
WHERE
  null IS NOT NULL
ORDER BY
  1 ASC,
  3 DESC;
```

I had difficulty with the sort operation related to syntax (parentheses required here but not there), the fetch limit requirement per select clause, and the sort requirement across all select clauses. I ended up doing some kludgy stuff with dual to make the sort work properly. There might be a more elegant way to accomplish this.

2. For each level of competition (local, state, etc.), list the top five beers in terms of the number of awards won. Sort results by event level alphabetically, and from highest to lowest award count within each level.

```
-- insert dummy data to validate query results
INSERT INTO
awards (award_name, award_category)
VALUES
('gold', 'pilsner');
INSERT INTO
awards (award_name, award_category)
VALUES
```

```
('silver', 'pilsner');
INSERT INTO
  awards (award_name, award_category)
VALUES
  ('bronze', 'pilsner');
INSERT INTO
  awards (award_name, award_category)
VALUES
  ('gold', 'lager');
INSERT INTO
  awards (award_name, award_category)
VALUES
  ('silver', 'lager');
INSERT INTO
  awards (award_name, award_category)
VALUES
  ('bronze', 'lager');
INSERT INTO
  awards (award_name, award_category)
VALUES
  ('gold', 'ale');
INSERT INTO
  awards (award_name, award_category)
VALUES
  ('silver', 'ale');
INSERT INTO
  awards (award_name, award_category)
VALUES
  ('bronze', 'ale');
INSERT INTO
  beer_comp (beer_id, award_id, event_id)
VALUES
  (10, 1, 2);
INSERT INTO
  beer_comp (beer_id, award_id, event_id)
VALUES
  (10, 2, 3);
INSERT INTO
  beer_comp (beer_id, award_id, event_id)
VALUES
  (20, 2, 2);
INSERT INTO
  beer_comp (beer_id, award_id, event_id)
VALUES
  (20, 1, 3);
INSERT INTO
  beer_comp (beer_id, award_id, event_id)
VALUES
```

```
(20, 1, 4);
INSERT INTO
  beer_comp (beer_id, award_id, event_id)
VALUES
  (30, 4, 6);
INSERT INTO
  beer_comp (beer_id, award_id, event_id)
VALUES
  (30, 5, 7);
INSERT INTO
  beer_comp (beer_id, award_id, event_id)
VALUES
  (40, 5, 6);
INSERT INTO
  beer_comp (beer_id, award_id, event_id)
VALUES
  (40, 4, 7);
INSERT INTO
  beer_comp (beer_id, award_id, event_id)
VALUES
  (40, 4, 8);
INSERT INTO
  beer_comp (beer_id, award_id, event_id)
VALUES
  (50, 7, 14);
INSERT INTO
  beer_comp (beer_id, award_id, event_id)
VALUES
  (50, 8, 15);
INSERT INTO
  beer_comp (beer_id, award_id, event_id)
VALUES
  (60, 8, 14);
INSERT INTO
  beer_comp (beer_id, award_id, event_id)
VALUES
  (60, 7, 15);
INSERT INTO
  beer_comp (beer_id, award_id, event_id)
VALUES
  (60, 7, 16);
-- run query
(SELECT
  ce.event_level,
  b.beer name,
  COUNT(bc.award id) as awards
FROM
```

```
beers b
  INNER JOIN beer_comp bc
    ON b.beer_id = bc.beer_id
  INNER JOIN comp_events ce
    ON bc.event id = ce.event id
WHERE
  ce.event_level LIKE 'local'
  AND bc.award_id <> -1
GROUP BY
  ce.event_level,
  b.beer_name
FETCH FIRST 5 ROWS ONLY)
UNION
(SELECT
  ce.event_level,
  b.beer_name,
  COUNT(bc.award_id) as awards
FROM
  beers b
  INNER JOIN beer_comp bc
    ON b.beer_id = bc.beer_id
  INNER JOIN comp_events ce
    ON bc.event_id = ce.event_id
WHERE
  ce.event_level LIKE 'state'
  AND bc.award_id <> -1
GROUP BY
  ce.event level,
  b.beer_name
FETCH FIRST 5 ROWS ONLY)
UNION
(SELECT
  ce.event_level,
  b.beer_name,
  COUNT(bc.award_id) as awards
FROM
  beers b
  INNER JOIN beer_comp bc
    ON b.beer_id = bc.beer_id
  INNER JOIN comp_events ce
    ON bc.event_id = ce.event_id
WHERE
  ce.event_level LIKE 'regional'
  AND bc.award_id <> -1
GROUP BY
  ce.event level,
  b.beer_name
FETCH FIRST 5 ROWS ONLY)
```

```
UNION
(SELECT
  ce.event_level,
  b.beer_name,
  COUNT(bc.award_id) as awards
FROM
  beers b
  INNER JOIN beer_comp bc
    ON b.beer_id = bc.beer_id
  INNER JOIN comp_events ce
    ON bc.event_id = ce.event_id
WHERE
  ce.event_level LIKE 'national'
  AND bc.award_id <> -1
GROUP BY
  ce.event_level,
  b.beer_name
FETCH FIRST 5 ROWS ONLY)
UNION
(SELECT
  ce.event_level,
  b.beer_name,
  COUNT(bc.award_id) as awards
FROM
  beers b
  INNER JOIN beer_comp bc
    ON b.beer_id = bc.beer_id
  INNER JOIN comp_events ce
    ON bc.event_id = ce.event_id
WHERE
  ce.event_level LIKE 'international'
  AND bc.award_id <> -1
GROUP BY
  ce.event_level,
  b.beer name
FETCH FIRST 5 ROWS ONLY)
UNION
SELECT
  null,
  null,
  null
FROM
  dual
WHERE
  null IS NOT NULL
ORDER BY
  1 ASC,
  3 DESC;
```

Additional Queries

The following queries were run after loading bulk data into the schema (see the Other Topics section).

1. Display the count of beer reviews that do not have an associated comment.

```
-- count 1710949

SELECT
COUNT(*)

FROM
beer_reviews rev
LEFT JOIN comments c
ON rev.reviewer_id=c.reviewer_id
AND rev.beer_id=c.beer_id

WHERE
c.reviewer_id IS NULL
AND c.beer_id IS NULL;
```

2. Display the count of beer reviews that have more than one associated comment.

```
-- count 0
SELECT
    COUNT(*)
FROM
    (SELECT
          DISTINCT c1.reviewer_id,
          c1.beer_id
FROM
          comments c1
          INNER JOIN comments c2
                ON c1.reviewer_id=c2.reviewer_id
                AND c1.beer_id=c2.beer_id
WHERE
                c1.comment_date<>c2.comment_date);
```

3. Display the count of beer reviews that have exactly one associated comment.

```
-- count 1094588
SELECT
COUNT(*)
FROM
beer_reviews rev
INNER JOIN comments c
```

```
ON rev.reviewer_id=c.reviewer_id
AND rev.beer_id=c.beer_id
WHERE
(rev.reviewer_id, rev.beer_id) NOT IN
(SELECT
DISTINCT c1.reviewer_id,
c1.beer_id
FROM
comments c1
INNER JOIN comments c2
ON c1.reviewer_id=c2.reviewer_id
AND c1.beer_id=c2.beer_id
WHERE
c1.comment_date<>c2.comment_date);
```

Note the total of the counts returned from the previous three queries is 2805537, which is the exact number of rows in the beer_reviews table.

4. Display the brewer ID and (if available) brewer name of every brewer that has an average overall rating of 80 or higher and an average ABV of 4% or higher. For the purpose of calculating average ABV, exclude beers with an ABV of 0. Sort the results by average rating highest to lowest and then by average ABV highest to lowest.

```
-- 108 rows returned
SELECT
 br.brewer id,
  br.brewer name,
  ROUND(AVG(rev.overall rating),1) AS Avg Rating,
  ROUND(AVG(b.abv),2) AS Avg_ABV
FROM
  brewers br
 INNER JOIN beers b
    ON br.brewer id=b.brewer id
 INNER JOIN beer_reviews rev
    ON b.beer id=rev.beer id
WHERE
 b.abv > 0
GROUP BY
 br.brewer_id,
  br.brewer name
 HAVING ROUND(AVG(rev.overall_rating),1) >= 80
  AND ROUND(AVG(b.abv),2) >= 4
ORDER BY
 Avg Rating DESC,
 Avg_ABV DESC;
```

Why are none of the brewer names displayed? Because the brewer name is not available for beers imported from the RateBeer dataset, and ratings are available only for beers imported from the RateBeer dataset.

5. Determine if any beers imported from the RateBeer dataset had a style name that was already in the styles table. If so, display the category and style names along with the count of beers for each style name. As a point of reference, include the count of all other imported beers in a single row. Sort by beer count highest to lowest. Hints: recall that all pre-existing beers were reassigned beer ID values over 500000, and all style names in the RateBeer dataset were mapped to category_id = -1 if the style name did not already exist in the styles table.

```
-- 6 rows returned
SELECT
  DISTINCT ca.category_name,
  s.style name,
  COUNT(*) AS Beer_Count
FROM
  categories ca
  INNER JOIN styles s
    ON ca.category_id=s.category_id
  INNER JOIN beers b
    ON s.style_id=b.style_id
WHERE
  b.beer id < 500000
  AND ca.category_id <> -1
GROUP BY
  ca.category_name,
  s.style name
UNION ALL
SELECT
  'Unknown' AS category_name,
  'Aggregated' AS style_name,
  COUNT(*) AS Beer Count
FROM
  categories ca
  INNER JOIN styles s
    ON ca.category id=s.category id
  INNER JOIN beers b
    ON s.style_id=b.style_id
WHERE
  b.beer_id < 500000
  AND ca.category_id = -1
ORDER BY
  Beer Count DESC;
```

The total number of beers counted in the previous query was 108194. The total number of pre-existing beers in the beers table after cleaning the original data was 5890. Those numbers sum to 114084, which is the total number of beers in the beers table.

6. Display all beers with an average overall rating that is equal to the highest overall rating given by any reviewer to any beer within the same style. Sort results by style name alphabetically and then by beer name alphabetically.

While working on this query, I discovered a data anomaly when I noticed some numbers did not add up properly. The updated styles table had 226 total rows. The original styles table had 141. Thus, 85 new styles were added during the RateBeer import. Imported beers were also associated to five of the pre-existing styles, for a total of 90 styles used by imported beers. However, the following query returned only 89 rows.

```
SELECT
DISTINCT style_id
FROM
beers
WHERE
beer_id < 500000;
```

During the import, several unrecognized characters were encountered in the RateBeer dataset. My Python script replaced them with "?". However, one style name (Kölsch) occurred twice with a different number of unrecognized characters. Consequently, two versions of this style name were imported, K?lsch and K???lsch. The strange thing is that no beers were associated to K???lsch. I have no explanation for this.

```
SELECT
  DISTINCT style_id,
  style_name
FROM
  styles
WHERE
  category id = -1
  AND style id <> -1
MINUS
SELECT
  DISTINCT s.style_id,
  style_name
FROM
  styles s
  INNER JOIN beers b
    ON s.style id=b.style id
WHERE
  b.beer id < 500000
  AND s.category id = -1;
```

```
SELECT
FROM
  styles
WHERE
  style_name LIKE 'K%';
SELECT
  beer_id,
  beer_name,
  b.style_id,
  s.style_name
FROM
  beers b
  INNER JOIN styles s
    ON b.style_id=s.style_id
WHERE
  b.style_id = 226
  OR b.style_id = 144;
To resolve this issue, I deleted the K???lsch style name from the styles table.
DELETE FROM
  styles
WHERE
  style_id = 226;
EXECUTE dbms_stats.gather_table_stats('DB870', 'styles');
Now back to the query at hand.
-- 53 rows returned
WITH
  style_ratings
AS
  (SELECT
    b.style_id,
    MAX(rev.overall_rating) AS max_rating
  FROM
    beers b
    INNER JOIN beer_reviews rev
      ON b.beer_id=rev.beer_id
  GROUP BY
    b.style id
  ORDER BY
```

```
style_id ASC)
SELECT
  s.style_name,
  b.beer_name,
  sr.max rating AS rating
FROM
  styles s
  INNER JOIN beers b
    ON s.style_id=b.style_id
  INNER JOIN style ratings sr
    ON b.style id=sr.style id
  INNER JOIN (SELECT
                b.beer id,
                ROUND(AVG(rev.overall_rating),0) AS avg_rating
             FROM
                beers b
             INNER JOIN beer reviews rev
                ON b.beer_id=rev.beer_id
             GROUP BY
                b.beer id
             ORDER BY
                beer id ASC) rat
    ON b.beer id=rat.beer id
WHERE
  rat.avg_rating=sr.max_rating
ORDER BY
  style_name ASC,
  beer name ASC;
```

A Stored Procedure

Goal: Identify the top five beers within each style based on average overall rating. Include ties even if doing so returns more than five beers per style. Display the style id, style name, average overall rating of the beer, beer id, beer name, and brewer id. (Brewer name would be nice but is not available.) Sort the results by style name alphabetically and then by rating highest to lowest.

An inelegant solution would use the following query as the base building block for a very long UNION series (currently 89 styles, so 89 base blocks). Sorting could be done with a "dummy" block at the end as was done with the "two interesting queries" from Assignment 2. However, each style_id would need to be coded manually into the WHERE clause of each block, and prior knowledge of every style_id in the beer_reviews table would be needed. Obviously, such code would be onerous to maintain and brittle in the presence of data updates.

```
WITH beer_ratings AS
```

```
(SELECT
    b.style_id,
    b.beer id,
    ROUND(AVG(rev.overall_rating),0) AS avg_rating
  FROM
    beers b
    INNER JOIN beer_reviews rev
      ON b.beer_id=rev.beer_id
  GROUP BY
    b.style id,
    b.beer_id
  ORDER BY
    style_id ASC,
    avg_rating DESC)
SELECT
  br.style_id,
  s.style_name,
  br.avg_rating,
  br.beer_id,
  b.beer_name,
  b.brewer_id
FROM
  beer_ratings br
  INNER JOIN beers b
    ON br.beer_id=b.beer_id
  INNER JOIN styles s
    ON b.style_id=s.style_id
WHERE
  br.style_id=13
ORDER BY
  avg_rating DESC
FETCH FIRST 5 ROWS WITH TIES;
```

A more elegant solution is the following stored procedure. It would be easier to maintain and more robust in the presence of data updates. However, I could not figure out how to sort the results as required and display column headers. Perhaps future students can tackle these challenges.

```
create or replace PROCEDURE
top_beers
IS
beer_lst DBMS_OUTPUT.CHARARR;
CURSOR
style_cur
IS
SELECT
DISTINCT b.style_id
FROM
```

```
beers b
      INNER JOIN beer_reviews rev
        ON b.beer_id=rev.beer_id
    ORDER BY
      style_id ASC;
BEGIN
  DBMS_OUTPUT.ENABLE (100000);
  FOR rec IN style_cur
    LOOP
      DECLARE
        CURSOR
          loop_cur
        IS
          SELECT
            br.style_id,
            s.style_name,
            br.avg_rating,
            br.beer_id,
            b.beer_name,
            b.brewer_id
          FROM
            (SELECT
              b.style_id,
              b.beer_id,
              ROUND(AVG(rev.overall_rating),0) AS avg_rating
            FROM
              beers b
            INNER JOIN beer_reviews rev
              ON b.beer_id=rev.beer_id
            GROUP BY
              b.style_id,
              b.beer_id
            ORDER BY
              style_id ASC,
              avg rating DESC) br
            INNER JOIN beers b
              ON br.beer_id=b.beer_id
            INNER JOIN styles s
              ON b.style_id=s.style_id
          WHERE
            br.style_id=rec.style_id
          ORDER BY
            avg_rating DESC
          FETCH FIRST 5 ROWS WITH TIES;
      BEGIN
      FOR style_rec IN loop_cur
        LOOP
```

```
DBMS_OUTPUT.PUT_LINE(style_rec.style_id || ':' || style_rec.style_name || ':' || style_rec.avg_rating || ':' || style_rec.beer_id || ':' || style_rec.beer_name || ':' || style_rec.brewer_id);

END LOOP;

END;

END LOOP;

END;

-- 665 rows returned

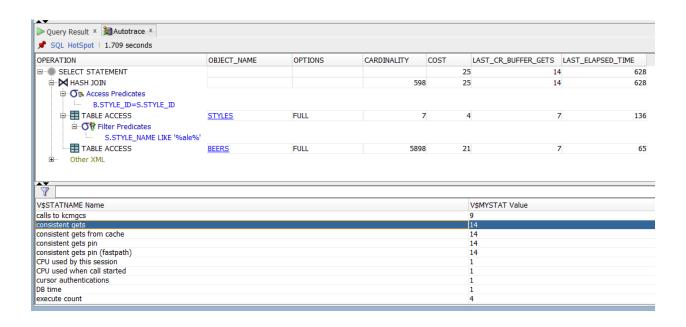
EXECUTE top_beers;
```

Performance Tuning

Basic Indexing & Query Performance (Assignment 3)

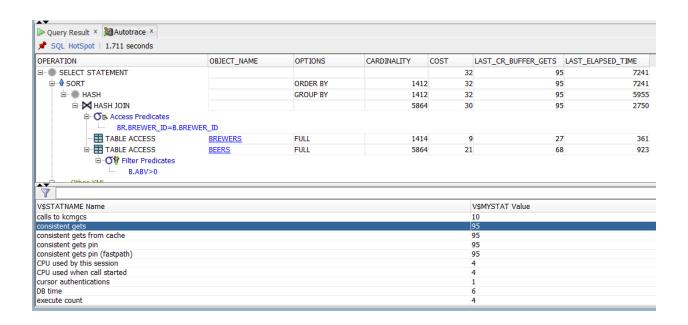
Using my implemented design from Assignment 2, I ran some basic queries with only the default indexes. The first query did a full scan on each joined table and returned 1025 rows.

```
SELECT
b.beer_name,
s.style_name
FROM
beers b
INNER JOIN styles s
ON b.style_id=s.style_id
WHERE
style_name LIKE '%ale%';
```

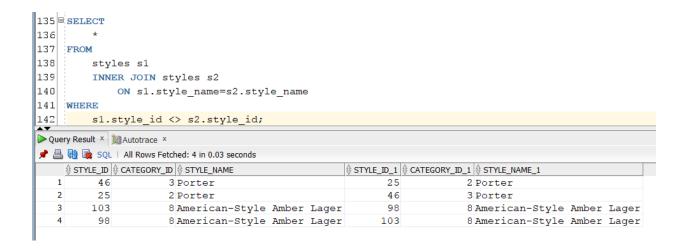


The second query also did a full scan on each joined table. It returned 754 rows.

```
SELECT
br.brewer_name,
MIN(b.abv) AS Min_ABV,
MAX(b.abv) AS Max_ABV,
ROUND(AVG(b.abv),2) AS Avg_ABV
FROM
brewers br
INNER JOIN beers b
ON br.brewer_id=b.brewer_id
WHERE
b.abv > 0
GROUP BY
br.brewer_name
ORDER BY
Avg_ABV DESC;
```



Since searches often include name attributes, it might be advantageous to create some indexes on frequently used name attributes. Beer drinkers tend to be interested in specific styles, so I started with the style_name attribute. Logically, all (category_name, style_name) tuples should be unique. However, I discovered a duplicate style_name record in the North American Lager category.



No beers were associated with this style, so I was able to delete the duplicate record without impacting other tables.

```
DELETE FROM
styles
WHERE
style_id = 103;
```

Then I created a unique constraint to prevent this incident from recurring. Oracle automatically created an index (also named cat_style_unq) on the constraint.

```
ALTER TABLE styles
ADD CONSTRAINT cat_style_unq
UNIQUE (category_id, style_name);
```

Next, I created a B+-tree index on style_name with default settings.

```
CREATE INDEX
style_name_btree
ON
styles (style_name);

EXECUTE DBMS_STATS.GATHER_TABLE_STATS('db870', 'styles');
```

It makes sense that beer and brewer names would be used in many searches, so I also created B+-tree indexes on each of them using default settings.

CREATE INDEX

```
beer_name_btree
ON
   beers (beer_name);

CREATE INDEX
   brewer_name_btree
ON
   brewers (brewer_name);

EXECUTE DBMS_STATS.GATHER_TABLE_STATS('db870', 'beers');

EXECUTE DBMS_STATS.GATHER_TABLE_STATS('db870', 'brewers');
```

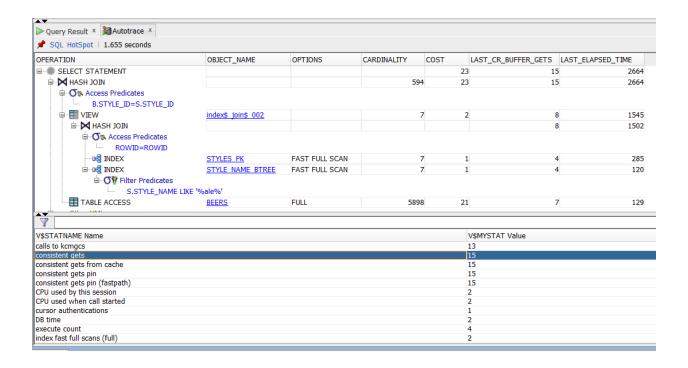
Before doing any more experimentation, I created the usual indexes on foreign keys. However, I did not create indexes for beers(srm), color_examples(srm), or styles(category_id) because the number of foreign key values is very low and likely to stay very low. Thus, a full scan of the referenced tables would not be costly, which means the optimizer might not use the indexes.

```
CREATE INDEX
 brewer_fk_btree
ON
 beers (brewer_id);
CREATE INDEX
  style_fk_btree
ON
 beers (style_id);
CREATE INDEX
  comp_fk_btree
ON
 comp_events (competition_id);
CREATE INDEX
  prod_fk_btree
ON
  competitions (producer_id);
CREATE INDEX
  admin_fk_btree
ON
  competitions (administrator_id);
EXECUTE DBMS_STATS.GATHER_TABLE_STATS('db870', 'beers');
EXECUTE DBMS_STATS.GATHER_TABLE_STATS('db870', 'comp_events');
```

EXECUTE DBMS_STATS.GATHER_TABLE_STATS('db870', 'competitions');

Next, I re-ran the simple queries from above. The first query used the new style_name index because I filtered on style_name. The new styles foreign key index was not used, but the styles primary key index was used. The consistent_gets increased by one, perhaps related to the extra object access (two indexes and one table versus two tables).

```
SELECT
b.beer_name,
s.style_name
FROM
beers b
INNER JOIN styles s
ON b.style_id=s.style_id
WHERE
style_name LIKE '%ale%';
```

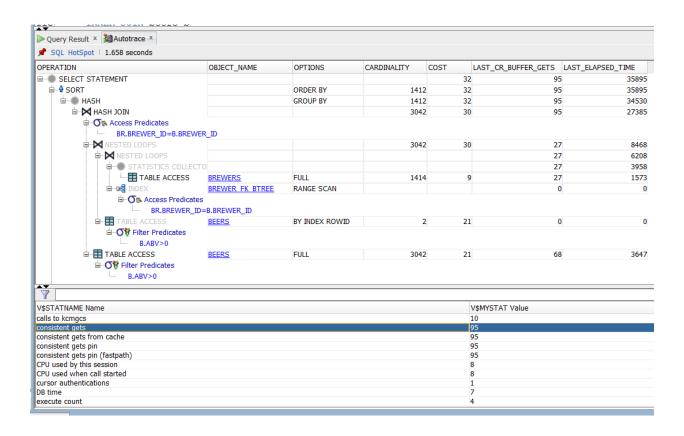


The second query did not use the new brewer_name index because the query required all brewers to be included in the results. The new brewer foreign key index was used. The consistent_gets remained the same.

SELECT

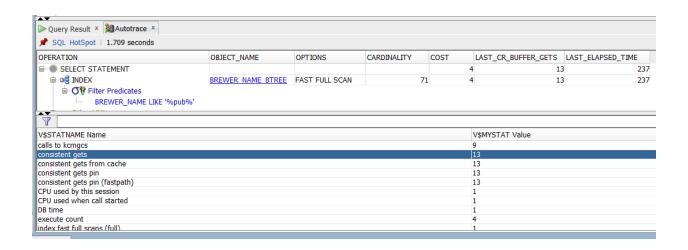
```
br.brewer_name,
MIN(b.abv) AS Min_ABV,
MAX(b.abv) AS Max ABV,
```

```
ROUND(AVG(b.abv),2) AS Avg_ABV
FROM
brewers br
INNER JOIN beers b
ON br.brewer_id=b.brewer_id
WHERE
b.abv > 0
GROUP BY
br.brewer_name
ORDER BY
Avg_ABV DESC;
```



Filtering on brewer_name to find pubs with inhouse microbreweries demonstrates the use of the brewer_name index. Only 25 rows are returned.

```
SELECT
br.brewer_name
FROM
brewers br
WHERE
brewer_name LIKE '%pub%';
```



Here's what I have so far. All the new indexes end with BTREE to make it easy to spot them.

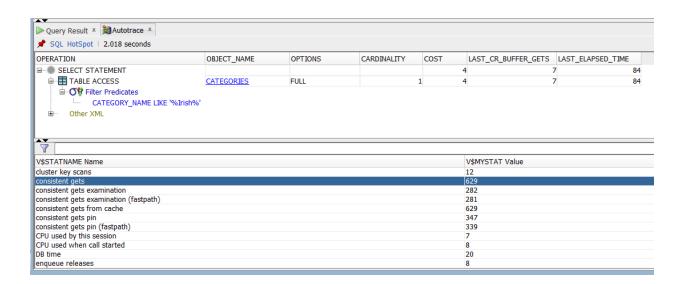


Intermediate Indexing & Query Performance (Assignment 3)

Bitmap indexes could be useful for attributes such as comp_events(event_level) with five levels and event_sponsors(sponsor_level) with a currently unknown but expected to be small number of values (such as titanium, platinum, gold, silver, and bronze). In time, those tables will each hold many instances. Another possibility is categories(category_name) with 12 values, but the number of instances likely will not grow beyond the current 12. Thus, we have high cardinality (all unique values) combined with a low number of values. That should make an interesting experiment. Since the categories tables is pre-populated and (indirectly) referenced from the pre-populated beers table, I have data with which to work. I started with the categories table and worked backward to the beers table.

```
SELECT

*
FROM
categories
WHERE
category_name LIKE '%Irish%';
```



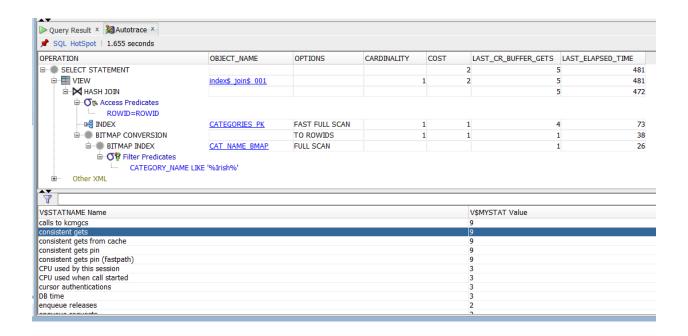
The query did a full table scan, and the consistent_gets came in at 629. I do not understand how such a small table could require that many I/Os. The Last_CR_Buffer_Gets column indicates much lower I/O. Anyway, I created the index as follows.

```
CREATE BITMAP INDEX
    cat_name_bmap
ON
    categories(category_name);

EXECUTE DBMS_STATS.GATHER_TABLE_STATS('db870', 'categories');

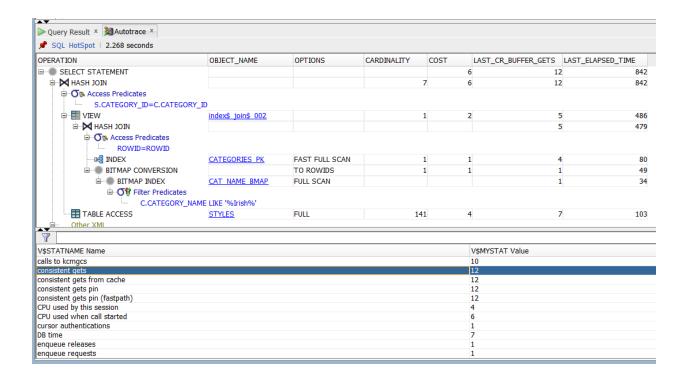
SELECT
```

```
FROM categories
WHERE category_name LIKE '%Irish%';
```



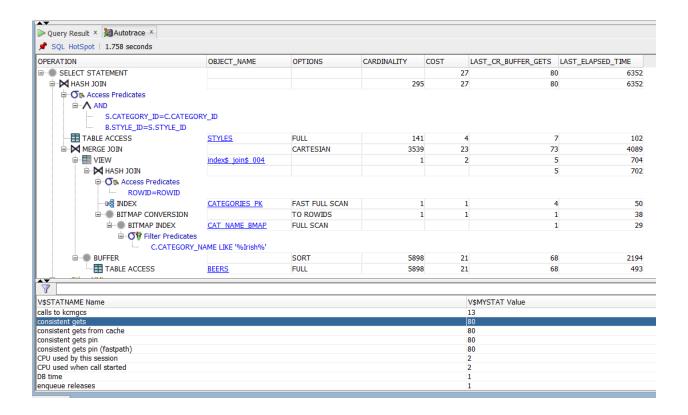
The query used the new index, and the consistent_gets dropped to just nine. The Last_CR_Buffer_Gets dropped from seven to five. Let's see what happens when I query the styles table and then the beers table.

```
SELECT
style_name
FROM
styles s
INNER JOIN categories c
ON s.category_id = c.category_id
WHERE
category_name LIKE '%Irish%';
```



The index is used again with a low consistent_gets count of 12. The index-related Last_CR_Buffer_Gets remained at five.

```
SELECT
beer_name
FROM
beers b
INNER JOIN styles s
ON b.style_id=s.style_id
INNER JOIN categories c
ON s.category_id = c.category_id
WHERE
category_name LIKE '%Irish%';
```



The index is used again, but the consistent_gets count rose to 80. However, the index-related Last CR Buffer Gets remained at five. Now let's try it without the bitmap index.

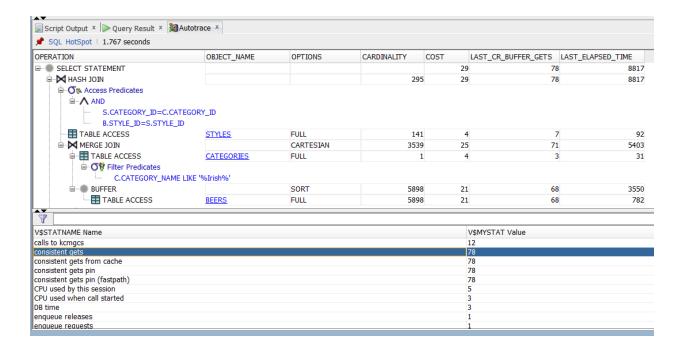
```
DROP INDEX
cat_name_bmap;

EXECUTE DBMS_STATS.GATHER_TABLE_STATS('db870', 'categories');

SELECT
beer_name

FROM
beers b
INNER JOIN styles s
ON b.style_id=s.style_id
INNER JOIN categories c
ON s.category_id = c.category_id

WHERE
category_name LIKE '%Irish%';
```



The consistent_gets dropped by two. The five index-related Last_CR_Buffer_Gets were replaced by three for a full table scan of categories. In a more realistic scenario (with a large table), the bitmap index would undoubtedly be more efficient than a full table scan. Of course, that scenario would imply a low cardinality attribute. So, we have confirmed a low number of values is, by itself, insufficient to guarantee a bitmap index efficiency advantage. However, I cannot state low cardinality is required for bitmap indexes to be advantageous because (1) my experiment did not prove that and (2) the OTN article (https://www.oracle.com/technical-resources/articles/sharma-indexes.html) provided in module five provides evidence to the contrary.

Moving on, I wanted to experiment with a composite index versus a pair of single indexes. Perhaps a beer lover wants to attend a beer competition in his/her area. A search based on city and state would be useful. For this experiment, I needed more data, so I inserted records as follows.

```
UPDATE
comp_events

SET
city = 'Miami',
state = 'FL'

WHERE
event_id IN (2,3);

UPDATE
comp_events

SET
city = 'Key West',
state = 'FL'

WHERE
event_id IN (4,5);

UPDATE
```

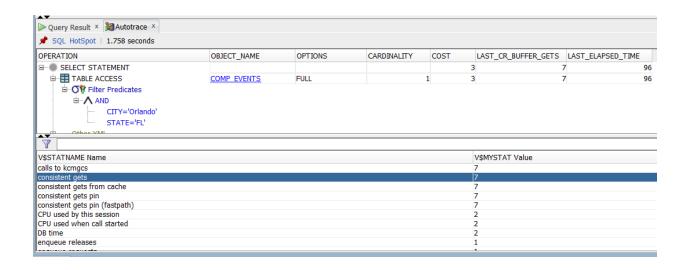
```
comp_events
SET
  city = 'Orlando',
  state = 'FL'
WHERE
  event_id IN (6,7);
UPDATE
  comp_events
SET
  city = 'Fort Myers',
  state = 'FL'
WHERE
  event_id IN (8,9);
UPDATE
  comp_events
SET
  city = 'Sarasota',
  state = 'FL'
WHERE
  event_id IN (14,15);
UPDATE
  comp_events
SET
  city = 'Tampa',
  state = 'FL'
WHERE
  event_id IN (16,17);
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event13', 1, 'regional', CURRENT_DATE, CURRENT_DATE, 'Gainesville', 'FL');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event14', 1, 'regional', CURRENT_DATE, CURRENT_DATE, 'Gainesville', 'FL');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event15', 1, 'state', CURRENT_DATE, CURRENT_DATE, 'Jacksonville', 'FL');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event16', 1, 'state', CURRENT DATE, CURRENT DATE, 'Jacksonville', 'FL');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event17', 2, 'state', CURRENT DATE, CURRENT DATE, 'Boca Raton', 'FL');
```

```
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event18', 2, 'state', CURRENT DATE, CURRENT DATE, 'Boca Raton', 'FL');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event19', 2, 'national', CURRENT_DATE, CURRENT_DATE, 'Orlando', 'FL');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event20', 2, 'national', CURRENT DATE, CURRENT DATE, 'Orlando', 'FL');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event21', 3, 'national', CURRENT_DATE, CURRENT_DATE, 'Tallahassee', 'FL');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event22', 3, 'national', CURRENT_DATE, CURRENT_DATE, 'Tallahassee', 'FL');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event23', 3, 'national', CURRENT DATE, CURRENT DATE, 'Lakeland', 'FL');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event24', 3, 'national', CURRENT DATE, CURRENT DATE, 'Lakeland', 'FL');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event25', 4, 'state', CURRENT_DATE, CURRENT_DATE, 'Lake City', 'FL');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event26', 4, 'state', CURRENT_DATE, CURRENT_DATE, 'Lake City', 'FL');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event27', 4, 'national', CURRENT DATE, CURRENT DATE, 'Hollywood', 'FL');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event28', 4, 'national', CURRENT DATE, CURRENT DATE, 'Hollywood', 'FL');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event29', 4, 'local', CURRENT DATE, CURRENT DATE, 'Melbourne', 'FL');
```

```
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event30', 4, 'local', CURRENT DATE, CURRENT DATE, 'Melbourne', 'FL');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event31', 5, 'national', CURRENT DATE, CURRENT DATE, 'Tampa', 'FL');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event32', 5, 'national', CURRENT DATE, CURRENT DATE, 'Tampa', 'FL');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event33', 5, 'international', CURRENT_DATE, CURRENT_DATE, 'Orlando', 'FL');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event34', 5, 'international', CURRENT_DATE, CURRENT_DATE, 'Orlando', 'FL');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event35', 5, 'national', CURRENT DATE, CURRENT DATE, 'Jacksonville', 'FL');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event36', 5, 'national', CURRENT DATE, CURRENT DATE, 'Jacksonville', 'FL');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event37', 6, 'international', CURRENT_DATE, CURRENT_DATE, 'Ocala', 'FL');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event38', 6, 'international', CURRENT DATE, CURRENT DATE, 'Ocala', 'FL');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event39', 6, 'international', CURRENT DATE, CURRENT DATE, 'West Palm Beach', 'FL');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event40', 6, 'international', CURRENT DATE, CURRENT DATE, 'West Palm Beach', 'FL');
```

Then I ran the following targeted query.

```
FROM comp_events
WHERE city = 'Orlando'
AND state = 'FL';
```

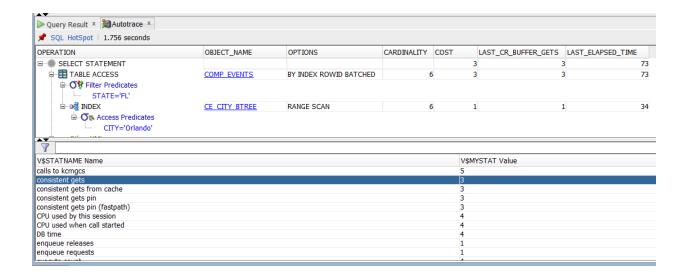


Of course, a full table scan was required because neither city nor state were indexed attributes. Then I created the following B+-tree indexes and ran the query again.

```
CREATE INDEX
    ce_city_btree
ON
    comp_events(city);
CREATE INDEX
    ce_state_btree
ON
    comp_events(state);

EXECUTE DBMS_STATS.GATHER_TABLE_STATS('db870', 'comp_events');

SELECT
    *
FROM
    comp_events
WHERE
    city = 'Orlando'
    AND state = 'FL';
```



The city index was used, but the state index was not. So, I added some events in other states to see if the state index would be used.

```
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event41', 1, 'national', CURRENT_DATE, CURRENT_DATE, 'Atlanta', 'GA');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event42', 1, 'national', CURRENT DATE, CURRENT DATE, 'Dallas', 'TX');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event43', 2, 'regional', CURRENT_DATE, CURRENT_DATE, 'Chicago', 'IL');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event44', 2, 'regional', CURRENT DATE, CURRENT DATE, 'San Diego', 'CA');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event45', 3, 'local', CURRENT DATE, CURRENT DATE, 'Huntsville', 'AL');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event46', 3, 'local', CURRENT_DATE, CURRENT_DATE, 'Jackson', 'MS');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event47', 4, 'state', CURRENT_DATE, CURRENT_DATE, 'Cheyenne', 'WY');
INSERT INTO
```

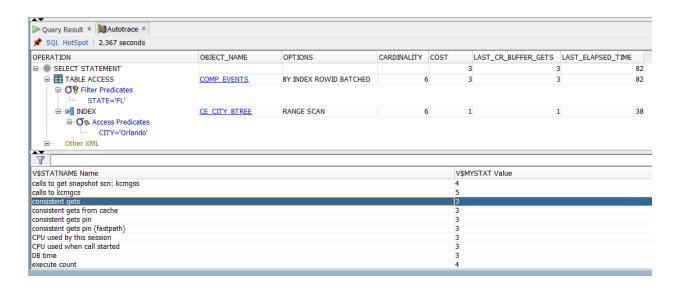
```
comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event48', 4, 'state', CURRENT DATE, CURRENT DATE, 'Tulsa', 'OK');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event49', 5, 'international', CURRENT DATE, CURRENT DATE, 'Anchorage', 'AK');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event50', 5, 'international', CURRENT_DATE, CURRENT_DATE, 'Roanoke', 'VA');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event51', 1, 'national', CURRENT DATE, CURRENT DATE, 'Denver', 'CO');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event52', 1, 'national', CURRENT DATE, CURRENT DATE, 'Boise', 'ID');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event53', 2, 'regional', CURRENT DATE, CURRENT DATE, 'Reno', 'NV');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event54', 2, 'regional', CURRENT_DATE, CURRENT_DATE, 'Tucson', 'AZ');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event55', 3, 'local', CURRENT_DATE, CURRENT_DATE, 'Albuquerque', 'NM');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event56', 3, 'local', CURRENT DATE, CURRENT DATE, 'Seattle', 'WA');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event57', 4, 'state', CURRENT DATE, CURRENT DATE, 'Portland', 'OR');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event58', 4, 'state', CURRENT_DATE, CURRENT_DATE, 'Billings', 'MT');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event59', 5, 'international', CURRENT DATE, CURRENT DATE, 'Bismarck', 'ND');
INSERT INTO
```

```
comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event60', 5, 'international', CURRENT DATE, CURRENT DATE, 'Rapid City', 'SD');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event61', 6, 'national', CURRENT DATE, CURRENT DATE, 'Lincoln', 'NE');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event62', 6, 'national', CURRENT_DATE, CURRENT_DATE, 'Wichita', 'KS');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event63', 6, 'regional', CURRENT DATE, CURRENT DATE, 'Provo', 'UT');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event64', 6, 'regional', CURRENT DATE, CURRENT DATE, 'Little Rock', 'AR');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event65', 6, 'local', CURRENT DATE, CURRENT DATE, 'Shreveport', 'LA');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event66', 6, 'local', CURRENT DATE, CURRENT DATE, 'St. Louis', 'MO');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event67', 6, 'state', CURRENT_DATE, CURRENT_DATE, 'Des Moines', 'IA');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event68', 6, 'state', CURRENT DATE, CURRENT DATE, 'Rochester', 'MN');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event69', 6, 'international', CURRENT DATE, CURRENT DATE, 'Green Bay', 'WI');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event70', 6, 'international', CURRENT_DATE, CURRENT_DATE, 'Fort Wayne', 'IN');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event71', 1, 'national', CURRENT DATE, CURRENT DATE, 'Atlanta', 'GA');
INSERT INTO
```

```
comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event72', 2, 'national', CURRENT_DATE, CURRENT_DATE, 'Atlanta', 'GA');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event73', 3, 'regional', CURRENT_DATE, CURRENT_DATE, 'Atlanta', 'GA');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event74', 4, 'regional', CURRENT_DATE, CURRENT_DATE, 'Atlanta', 'GA');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event75', 5, 'local', CURRENT DATE, CURRENT DATE, 'Atlanta', 'GA');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event76', 6, 'local', CURRENT DATE, CURRENT DATE, 'Atlanta', 'GA');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event77', 6, 'state', CURRENT_DATE, CURRENT_DATE, 'Atlanta', 'GA');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event78', 6, 'state', CURRENT_DATE, CURRENT_DATE, 'Atlanta', 'GA');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event79', 6, 'international', CURRENT_DATE, CURRENT_DATE, 'Atlanta', 'GA');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event80', 6, 'international', CURRENT DATE, CURRENT DATE, 'New York', 'NY');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event81', 1, 'national', CURRENT DATE, CURRENT DATE, 'Atlanta', 'GA');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event82', 2, 'national', CURRENT_DATE, CURRENT_DATE, 'Atlanta', 'GA');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event83', 3, 'regional', CURRENT DATE, CURRENT DATE, 'Atlanta', 'GA');
INSERT INTO
```

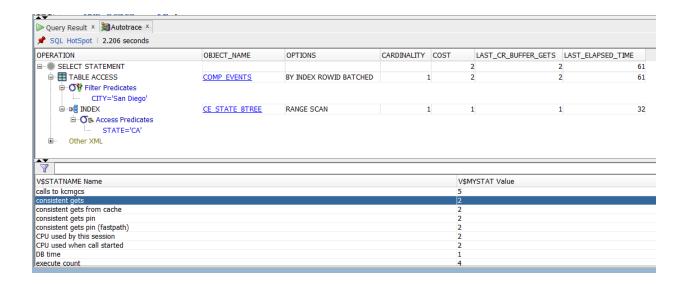
```
comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event84', 4, 'regional', CURRENT DATE, CURRENT DATE, 'Atlanta', 'GA');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event85', 5, 'local', CURRENT DATE, CURRENT DATE, 'Atlanta', 'GA');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event86', 1, 'local', CURRENT_DATE, CURRENT_DATE, 'Atlanta', 'GA');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event87', 2, 'state', CURRENT DATE, CURRENT DATE, 'Atlanta', 'GA');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event88', 3, 'state', CURRENT DATE, CURRENT DATE, 'Atlanta', 'GA');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event89', 4, 'international', CURRENT DATE, CURRENT DATE, 'Atlanta', 'GA');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event90', 5, 'international', CURRENT_DATE, CURRENT_DATE, 'Atlanta', 'GA');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event91', 1, 'national', CURRENT_DATE, CURRENT_DATE, 'Pittsburgh', 'PA');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event92', 2, 'national', CURRENT DATE, CURRENT DATE, 'Pittsburgh', 'PA');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event93', 3, 'regional', CURRENT DATE, CURRENT DATE, 'Pittsburgh', 'PA');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event94', 4, 'regional', CURRENT_DATE, CURRENT_DATE, 'Pittsburgh', 'PA');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event95', 5, 'local', CURRENT DATE, CURRENT DATE, 'Pittsburgh', 'PA');
INSERT INTO
```

```
comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event96', 1, 'local', CURRENT_DATE, CURRENT_DATE, 'Pittsburgh', 'PA');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event97', 2, 'state', CURRENT_DATE, CURRENT_DATE, 'Pittsburgh', 'PA');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event98', 3, 'state', CURRENT_DATE, CURRENT_DATE, 'Pittsburgh', 'PA');
INSERT INTO
  comp_events (event_name, competition_id, event_level, start_date, end_date, city, state)
VALUES
  ('event99', 4, 'international', CURRENT DATE, CURRENT DATE, 'Pittsburgh', 'PA');
INSERT INTO
  comp events (event name, competition id, event level, start date, end date, city, state)
VALUES
  ('event100', 5, 'international', CURRENT DATE, CURRENT DATE, 'Pittsburgh', 'PA');
SELECT
FROM
  comp_events
WHERE
  city = 'Orlando'
  AND state = 'FL';
```



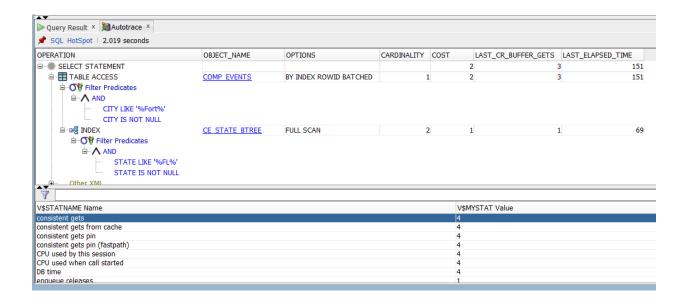
Same result, so I changed the query parameters.

```
FROM comp_events
WHERE city = 'San Diego'
AND state = 'CA';
```



The state index was used, but not the city index. Next, I tried a wildcard query.

```
FROM
comp_events
WHERE
city LIKE '%Fort%'
AND state LIKE '%FL%';
```



The results were the same except for the type of scan done on the state index. Next, I dropped the city and state indexes and created a composite index. Then I ran the same three queries.

```
DROP INDEX
    ce_city_btree;
DROP INDEX
    ce_state_btree;

CREATE INDEX
    ce_city_state_btree
ON
    comp_events(city, state);

EXECUTE DBMS_STATS.GATHER_TABLE_STATS('db870', 'comp_events');

SELECT
    *

FROM
    comp_events
WHERE
    city = 'Orlando'
    AND state = 'FL';
```

SELECT STATEMENT TABLE ACCESS COMP EVENTS BY INDEX ROWID BATCHED CE CITY STATE BTREE RANGE SCAN 2 1 CITY='Orlando' STATE='FL' V\$STATNAME Name bytes sent via SQL*Net to client calls to get snapshot scn: kcmgss calls to get snapshot scn: kcmgss calls to kcmgcs CCursor + sql area evicted consistent gets prin consist	CGETS I	S LAST_ELAPSED_TI	ME
CE CITY STATE BTREE RANGE SCAN 2 1 AND CITY=Orlando' STATE='FL' V\$STATNAME Name V\$MYSTAT Value bytes sent via SQL*Net to client calls to get snapshot scn: kcmgss calls to kcmgcs 7 CCursor + sql area evicted 1 consistent gets from cache consistent gets from cache consistent gets from cache CPU used by this session CPU used when call started 3 CPU used when call started	3	3	
### Access Predicates #### CITY='Orlando'	3	3	
### Access Predicates #### CITY='Orlando'	1	1	
CITY='Orlando' STATE='FL' V\$MYSTAT Value			
CITY='Orlando' STATE='FL'			
STATE='FL' Other YM V\$MYSTAT Value V\$STATNAME Name V\$MYSTAT Value V\$MYST			
V\$STATNAME Name V\$MYSTAT Value bytes sent via SQL**Net to client 52554 calls to get snapshot scn: kcmgss 8 calls to kcmgcs 7 CCursor + sql area evicted 1 consistent gets 4 consistent gets from cache 4 consistent gets pin 4 consistent gets pin (fastpath) 4 CPU used by this session 3 CPU used when call started 3			
V\$TATNAME Name V\$MYSTAT Value bytes sent via SQL*Net to client 52554 calls to get snapshot scn: kcmgss 8 calls to kcmgcs 7 CCursor + sql area evicted 1 consistent gets 4 consistent gets from cache 4 consistent gets pin 4 CPU used by this session 3 CPU used when call started 3			
V\$STATNAME Name V\$MYSTAT Value bytes sent via SQL*Net to client 52554 calls to get snapshot scn: kcmgss 8 calls to kcmgcs 7 CCursor + sql area evicted 1 consistent gets 4 consistent gets from cache 4 consistent gets pin 4 consistent gets pin (fastpath) 4 CPU used by this session 3 CPU used when call started 3			
bytes sent via SQL*Net to client 5254 calls to get snapshot scn: kcmgss 8 calls to kcmgcs 7 CCursor + sql area evicted 1 consistent gets 4 consistent gets from cache 4 consistent gets pin 4 consistent gets pin (fastpath) 4 CPU used by this session 3 CPU used when call started 3			
calls to get snapshot scn: kcngss 8 calls to kcngcs 7 CCursor + sql area evicted 1 consistent gets 4 consistent gets from cache 4 consistent gets pin 4 consistent gets pin (fastpath) 4 CPU used by this session 3 CPU used when call started 3			_
calls to kcmgcs 7 CCursor + sql area evicted 1 consistent gets 4 consistent gets from cache 4 consistent gets pin 4 consistent gets pin (fastpath) 4 CPU used by this session 3 CPU used when call started 3			
CCursor + sql area evicted 1 consistent gets 4 consistent gets from cache 4 consistent gets pin 4 consistent gets pin (fastpath) 4 CPU used by this session 3 CPU used when call started 3			
consistent gets 4 consistent gets from cache 4 consistent gets pin 4 consistent gets pin (fastpath) 4 CPU used by this session 3 CPU used when call started 3			
consistent gets from cache 4 consistent gets pin 4 consistent gets pin (fastpath) 4 CPU used by this session 3 CPU used when call started 3			
consistent gets pin 4 consistent gets pin (fastpath) 4 CPU used by this session 3 CPU used when call started 3			
consistent gets pin (fastpath) 4 CPU used by this session 3 CPU used when call started 3			
CPU used by this session 3 CPU used when call started 3			
CPU used when call started 3			
DB time 3			
us urie 3 3 enquer e la serie 1 1 1 enquer e la serie 2 1 1 enquer e la serie 2			

SELECT

*

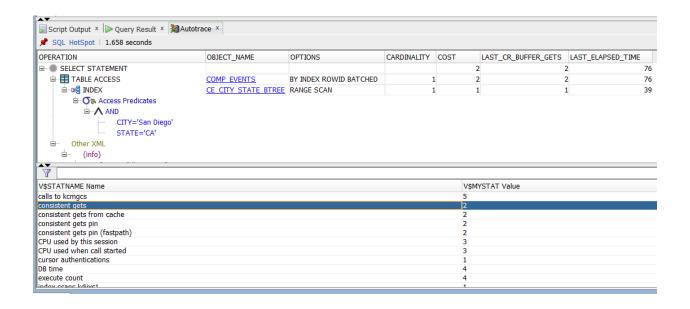
FROM

comp_events

WHERE

city = 'San Diego'

AND state = 'CA';



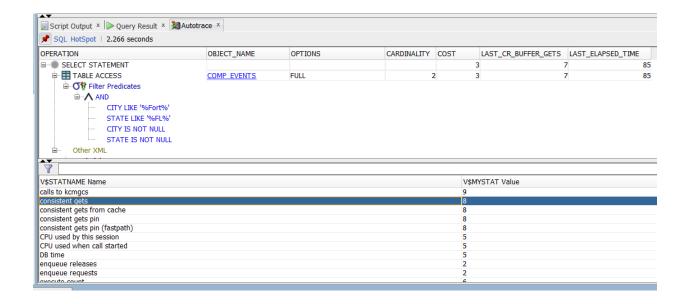
SELECT

*

FROM

comp_events

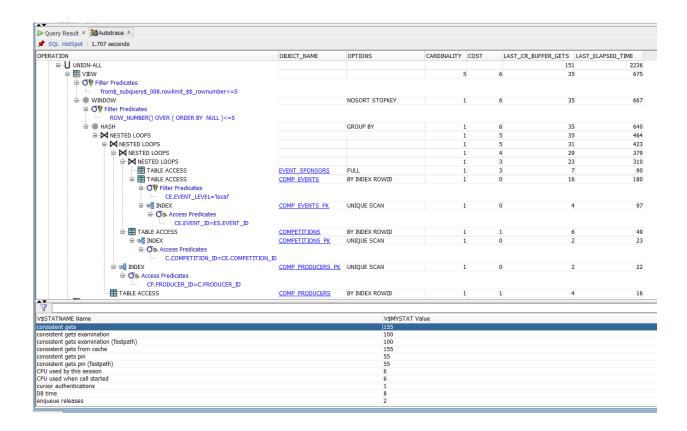
WHERE city LIKE '%Fort%' AND state LIKE '%FL%';



The first two queries used the composite index. The third query did not, which I suppose was due to the less targeted nature of the wildcard query. This seems conceptually consistent with the execution of the wildcard query using separate indexes in the sense that the type of index scan went from range to full. My take-away is that composite indexes are best used when you know a certain type of targeted query will be run repeatedly. Otherwise, stick to separate indexes so you have a better chance of the indexes being used.

With the composite index, the consistent_gets jumped from two for the second targeted query to eight for the wildcard query. In absolute terms, that is not a significant difference. However, in relative terms, that is a very significant increase, which speaks to the potential benefits of indexes on large tables. In all the queries on the comp_events table, the consistent_gets and Last_CR_Buffer_Gets were very small, most likely due to the small size of the table. I suppose I would get more interesting results from a much larger table.

For my final experiment, I wanted to optimize my complex union query (the first one) from the last assignment. I thought a bitmap index on comp_events(event_level) would help. Likewise, as the database grows, a B+-tree on comp_producers(prod_name) might help too. For brevity's sake, the query is not repeated here. Click here to review it.



The consistent_gets are 155 with all the indexes I implemented so far. Looking at the autotrace results, it appears the primary key index of each table except event_sponsors was scanned to facilitate the joins. Then the rowids from the index scans were used to facilitate table access. A full scan was done on event_sponsors presumably due to the aggregation function in the SELECT clause. The new indexes were implemented as follows.

```
CREATE BITMAP INDEX
    event_level_bmap

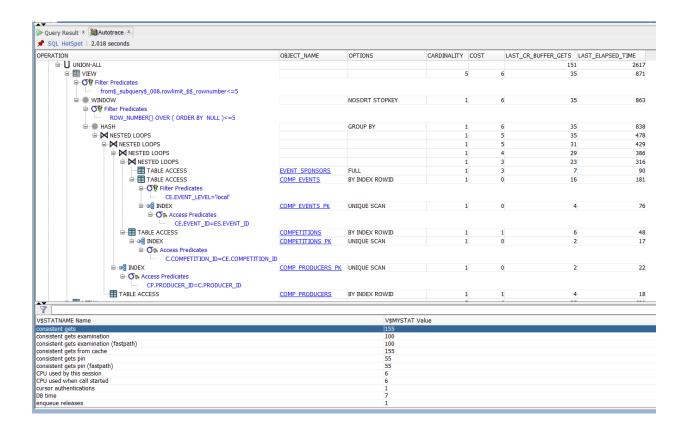
ON
    comp_events(event_level);

CREATE INDEX
    prod_name_btree

ON
    comp_producers(prod_name);

EXECUTE DBMS_STATS.GATHER_TABLE_STATS('db870', 'comp_events');

EXECUTE DBMS_STATS.GATHER_TABLE_STATS('db870', 'comp_producers');
```

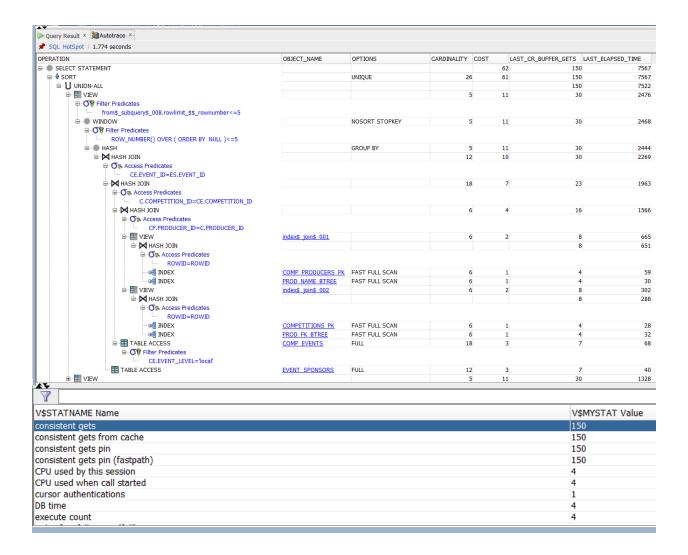


The consistent_gets remained 155 with the new indexes. I did not expect the prod_name index to be used since the table only has six records, but I thought the event_level index would be used (based on the results of the category_name experiment). Making comp_events_pk a covering index for event_level would eliminate the table access, but that would potentially preclude use of the index for many other queries. That seems like a bad idea for a primary key index. Maybe as the table grows, the bitmap index would become advantageous.

Just for kicks, I re-ran the gather_table_stats statements with sample size at 100%. This time I also included the event_sponsors table, which had not been analyzed since its creation.

EXECUTE DBMS_STATS.GATHER_TABLE_STATS ('db870', 'comp_events', estimate_percent => 100); EXECUTE DBMS_STATS.GATHER_TABLE_STATS ('db870', 'competitions', estimate_percent => 100); EXECUTE DBMS_STATS.GATHER_TABLE_STATS ('db870', 'comp_producers', estimate_percent => 100); EXECUTE DBMS_STATS.GATHER_TABLE_STATS ('db870', 'event_sponsors', estimate_percent => 100);

When I re-ran the query, I got interesting autotrace results.



Unlike last time, the prod_name_btree index was used (surprisingly) as was the prod_fk_btree index, but the comp_events_pk index was not. Like last time, the competitions_pk and comp_producers_pk indexes were used. The consistent_gets dropped by five, while the Last_CR_Buffer_Gets dropped by one. Even though the performance did not improve much, this experiment illustrates the potential variance among execution plans. I am not sure if the execution plan changed as a result of increasing the sample size or gathering stats on the event_sponsors table or both.

Next, I changed all LIKE operators to =. The execution plan was identical, and the I/O cost was nearly identical. This makes sense considering the equality returned multiple matches just as the LIKE did. Had the event_level attribute been subject to a UNIQUE constraint, the equality would have returned just one result per SELECT clause, and the execution plan probably would have been different.

I wondered about the relationship of indexes to set operations, union in particular. In chapter 15 of the textbook (7th edition), query processing is discussed. Little guidance is offered about optimizing unions (disjunctives). If I understand correctly, the authors indicate optimization is performed within each SELECT statement to retrieve pointers optimally (same as non-set operations). The only difference is

that the pointers are aggregated prior to record retrieval from the file(s). Likewise, the lecture slides did not provide any further insight. So, there does not seem to be any union-specific indexing technique.

I decided to drop the bitmap index on event_level since it was never used. I kept the B+-tree index on prod_name since that attribute will likely be used as a predicate in future queries as the database grows, and my experiment results indicate it will be used at times.

```
DROP INDEX
  event_level_bmap;
```

This left the database with no active bitmap indexes.



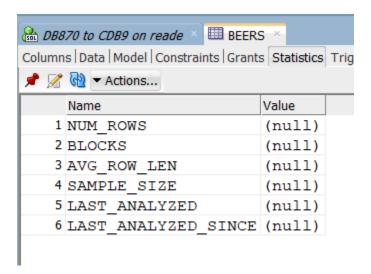
Column Statistics (Assignment 4)

Between assignments three and four, I loaded bulk data into my schema. The "Other Topics" section of this document contains the details of that effort. The following table summarizes the results. The comments table has a LONG attribute that contains textual comments that exceed the maximum size of VARCHAR2.

Table	Total Records	Pre-existing Records	New Records
beers	114,084	5,890	108,194
brewers	8,789	1,414	7,375
styles	226	141	85
reviewers	29,097	0	29,097
beer_reviews	2,805,537	0	2,805,537
comments	2,775,682	0	2,775,682

For my first experiment, I cleared all statistics in my schema and ran some queries to establish a performance baseline.

EXECUTE dbms_stats.delete_schema_stats('DB870');



```
SELECT
b.beer_name,
b.beer_id,
avg(rev.overall_rating) as avg_rating
FROM
beers b
INNER JOIN beer_reviews rev
ON b.beer_id = rev.beer_id
GROUP BY
```

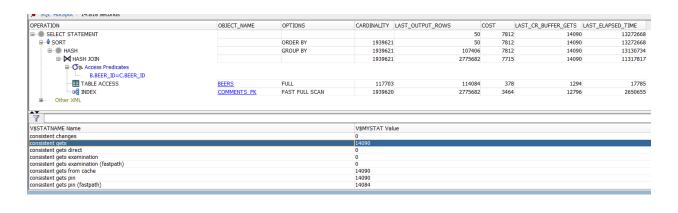
```
b.beer_name,
 b.beer_id
ORDER BY
  avg_rating DESC,
 beer_name ASC;
SELECT
 b.beer_name,
 b.beer_id,
 count(c.review_comment) as count
FROM
 beers b
 INNER JOIN comments c
    ON b.beer_id = c.beer_id
GROUP BY
  b.beer_name,
  b.beer_id
ORDER BY
  count DESC,
 beer_name ASC;
```

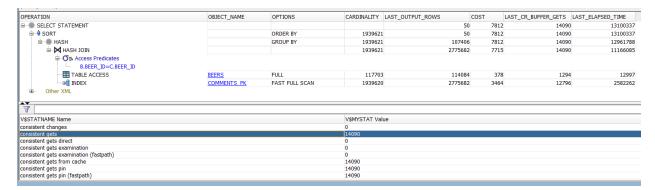
OPERATION	OBJECT_NAME	OPTIONS	CARDINALITY	LAST_OUTPUT_ROWS	COST	LAST_CR_BUFFER_GETS	LAST_ELAPSED_TIME
■ SELECT STATEMENT				50	7812	14090	14790543
SORT		ORDER BY	1939621	50	7812	14090	14790543
⊟ MASH		GROUP BY	1939621	107406	7812	14090	1463817
			1939621	2775682	7715	14090	1275956
Access Predicates							
C.BEER_ID=B.BEER_ID							
TABLE ACCESS	BEERS	FULL	117703	114084	378	1294	9792
od INDEX	COMMENTS PK	FAST FULL SCAN	1939620	2775682	3464	12796	4273223
■ Other XML							
AY							
7							
V\$STATNAME Name			V\$MYSTAT Val	ue			
consistent changes			0				
consistent gets			78644				
			0				
consistent gets direct			58768				
consistent gets direct consistent gets examination							
consistent gets direct consistent gets examination consistent gets examination (fastpath)			58766				
consistent gets direct consistent gets examination			58766 78644 19876				

OPERATION	OBJECT_NAME	OPTIONS	CARDINALITY	LAST_OUTPUT_ROWS	COST	LAST_CR_BUFFER_GETS	LAST_ELAPSED_TIME		
■- SELECT STATEMENT				50	7812	14090	14074578		
⇒ • SORT		ORDER BY	1939621	50	7812	14090	14074578		
⊞ HASH		GROUP BY	1939621	10740	7812	14090	13932128		
⊞ M HASH JOIN			1939621	277568	7715	14090	12007238		
Access Predicates B.BEER_ID=C.BEER_ID									
TABLE ACCESS	BEERS	FULL	117703	114084	378	1294	11485		
	COMMENTS PK	FAST FULL SCAN	1939620	277568	3464	12796	3120219		
7			VANDOTAT VAL						
V\$STATNAME Name			V\$MYSTAT Val	ue					
consistent changes			0 111416						
consistent gets direct			0						
consistent gets examination			61422						
consistent gets examination (fastpath)			61422						
consistent gets from cache			111416						
consistent gets pin			49994						
consistent gets pin (fastpath)			49994						

Curiously, the primary key index of the comments table was used by both queries even though the first query did not join the comments table. Next, I gathered statistics on all tables in the schema and re-ran the same queries in the same order to determine how much performance improved.

EXECUTE dbms stats.gather schema stats('DB870');





As we can see, the last_elapsed_time decreased slightly for both queries, while the consistent_gets dropped significantly for both queries. The execution plans did not change.

Next, I experimented with column group statistics using the ratings columns in the beer_reviews table. First, I cleared table statistics and ran a query to get a baseline.

EXECUTE dbms_stats.delete_table_stats('DB870', 'beer_reviews');

```
SELECT
b.beer_name,
b.beer_id,
rev.taste_rating,
rev.aroma_rating
FROM
beers b
INNER JOIN beer_reviews rev
ON b.beer id = rev.beer id
```

```
WHERE

rev.taste_rating >= 80

AND rev.aroma_rating >= 80

AND rev.palate_rating >= 80

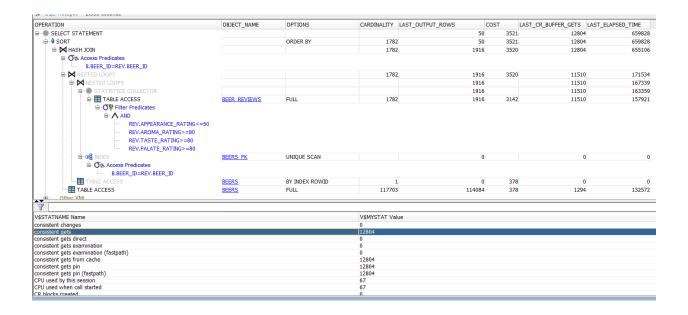
AND rev.appearance_rating <= 50

ORDER BY

taste_rating DESC,

aroma_rating DESC,

beer_name ASC;
```



Then, I gathered column group statistics and re-ran the query to determine how much performance improves.

SELECT

```
dbms_stats.create_extended_stats('DB870', 'beer_reviews', '(overall_rating, appearance_rating, aroma_rating, palate_rating, taste_rating)')
FROM
dual;
```

EXECUTE dbms_stats.gather_table_stats('DB870', 'beer_reviews');

DPERATION	OBJECT_NAME	OPTIONS	CARDINALITY	LAST_OUTPUT_ROWS	COST	LAST_CR_BUFFER_GETS	LAST_ELAPSED_TIME
SELECT STATEMENT				50	3521	12804	515748
SORT		ORDER BY	1782	50	3521	12804	515748
⊞-M HASH JOIN			1782	1916	3520	12804	512221
Access Predicates							
B.BEER_ID=REV.BEER_ID							
■ M NESTED LOOPS			1782	1916	3520	11510	131604
B NESTED LOOPS				1916		11510	128519
■ STATISTICS COLLECTOR				1916		11510	12544
☐ ■ TABLE ACCESS	BEER REVIEWS	FULL	1782		3142		
⊕ O ♥ Filter Predicates	DELIC REFIELD	1 022	1702	1510	52.12		12271
REV.APPEARANCE_RATING<=50							
⊕-∧ AND ———————————————————————————————————							
REV.AROMA_RATING>=80							
REV.TASTE_RATING>=80							
REV.PALATE RATING>=80							
□ □ □ INDEX	BEERS PK	UNIQUE SCAN		0		0	(
□ On Access Predicates							
B.BEER ID=REV.BEER ID							
TABLE ACCESS	BEERS	BY INDEX ROWID	1	0	378	0	(
TABLE ACCESS	BEERS	FULL	117703		378		
	DEEKS	POLL	117703	114004	3/0	1294	101005
The Other XMI							
/\$STATNAME Name			V\$MYSTAT Val	ue			
consistent changes			0				
consistent gets			12804				
consistent gets direct			0				
consistent gets examination			0				
consistent gets examination (fastpath)			0				
consistent gets from cache			12804				
consistent gets pin			12804 12804				
consistent gets pin (fastpath) CPU used by this session			53				

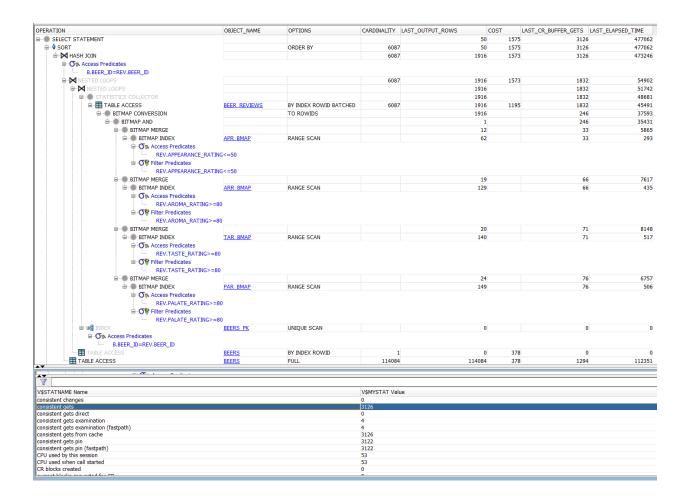
As we can see, last_elapsed_time dropped moderately, while consistent_gets and the execution plan remained unchanged. That was not what I was expecting to see. Several online articles indicated the optimizer can benefit from column group statistics when several columns are used together in the predicate. I suppose the specific nature of the relationships between the columns in the predicate could influence the results, but I was not sure what to do next. So, I moved on.

Bitmap Indexes (Assignment 4)

Since the ratings columns in beer_reviews all have low cardinality with a small number of values, I thought they would be good candidates for another bitmap index experiment.

```
CREATE BITMAP INDEX
ovr_bmap
ON
beer_reviews(overall_rating);
CREATE BITMAP INDEX
apr_bmap
ON
beer_reviews(appearance_rating);
CREATE BITMAP INDEX
arr_bmap
ON
beer_reviews(aroma_rating);
CREATE BITMAP INDEX
arr_bmap
ON
beer_reviews(aroma_rating);
CREATE BITMAP INDEX
par_bmap
ON
```

```
beer_reviews(palate_rating);
CREATE BITMAP INDEX
  tar_bmap
ON
  beer_reviews(taste_rating);
EXECUTE dbms_stats.gather_table_stats('DB870', 'beer_reviews');
I tested with the same query from the previous experiment.
SELECT
  b.beer_name,
  b.beer_id,
  rev.taste_rating,
  rev.aroma_rating
FROM
  beers b
  INNER JOIN beer_reviews rev
    ON b.beer_id = rev.beer_id
WHERE
  rev.taste_rating >= 80
  AND rev.aroma_rating >= 80
  AND rev.palate_rating >= 80
  AND rev.appearance_rating <= 50
ORDER BY
  taste_rating DESC,
  aroma_rating DESC,
  beer_name ASC;
```



As we can see, last_elapsed_time dropped moderately, while consistent_gets dropped significantly. The execution plan changed quite a bit and used four of the new bitmap indexes.

Optimizer Modes (Assignment 4)

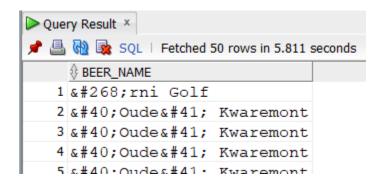
Next, I experimented with session-level optimizer modes. I first tried to query substrings in the review_comment attribute of the comments table and discovered that is not possible with the LONG data type. So, I tried to convert review_comment from LONG to CLOB. Unfortunately, a tablespace error occurred.

ORA-01652: unable to extend temp segment by 1024 in tablespace STUDENTS 01652. 00000 - "unable to extend temp segment by %s in tablespace %s" *Cause: Failed to allocate an extent of the required number of blocks for a temporary segment in the tablespace indicated.

So, I used the comment_date attribute instead.

ALTER SESSION SET OPTIMIZER_MODE = ALL_ROWS;

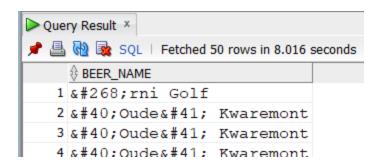
```
SELECT
b.beer_name
FROM
beers b
INNER JOIN comments c
ON b.beer_id = c.beer_id
WHERE
c.comment_date <= 1200000000
ORDER BY
b.beer_name ASC;
```

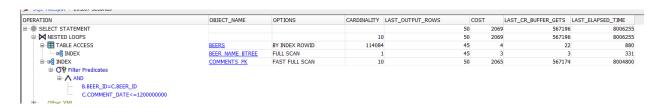




ALTER SESSION SET OPTIMIZER MODE = FIRST ROWS 10;

```
SELECT
b.beer_name
FROM
beers b
INNER JOIN comments c
ON b.beer_id = c.beer_id
WHERE
c.comment_date <= 1200000000
ORDER BY
b.beer_name ASC;
```





I was confused by these results at first. In FIRST_ROWS_10 mode, the optimizer used the comments_pk and beer_name_btree indexes, accessed the beers table BY INDEX ROWID, and avoided doing a HASH JOIN. In ALL_ROWS mode, only the comments_pk index was used, and a FULL SCAN was done on the beers table, and a HASH JOIN was employed. The FIRST_ROWS_10 execution plan should be faster since more indexes were used, access BY INDEX ROWID generally takes less time than a FULL SCAN, and a HASH JOIN entails startup cost. Yet, the FIRST_ROWS_10 query took longer. I hypothesize that, due to the excessively large size of the inner result set from the comments_pk index, the time required for the NESTED LOOP JOIN more than offset the time required for the FULL SCAN of the relatively small beers table and the HASH JOIN startup cost. Since I was one for three in Assignment 4, I decided to move on to partitioning where I felt confident that I could get "friendly" results.

Table Partitioning (Assignment 4)

I found a way to partition an unpartitioned table while keeping the table online. I tried it on the comments table using the following code:

```
ALTER TABLE comments

MODIFY

PARTITION BY RANGE (comment_date) INTERVAL (100000000)

(PARTITION p1 VALUES LESS THAN (100000000),

PARTITION p2 VALUES LESS THAN (1100000000),

PARTITION p3 VALUES LESS THAN (1200000000),

PARTITION p4 VALUES LESS THAN (1300000000),

PARTITION p5 VALUES LESS THAN (1400000000))

ONLINE;
```

Unfortunately, I got an error. Further research revealed this method only works in Oracle 12c R2, and the course server appears to be running Oracle 12c R1. An alternative was to create a partitioned table, copy the data from the comments table into the new table, enable any lost constraints on the new table, disable downstream constraints, drop the original comments table, rename the new table "comments", re-enable downstream constraints, redefine any lost indexes on the new comments table, and refresh the new table's stats. Since I had already twice received tablespace errors indicating limited available storage space, this method would require deletion of a significant amount of data from the existing comments table prior to moving any data to the new table. This solution was not ideal, but I chose to move forward because I wanted to see the benefits of partitioning in action on a very large table. Since I did not know the distribution of rows across dates, I proceeded with caution.

```
DELETE FROM
comments
WHERE
comment_date < 1100000000;

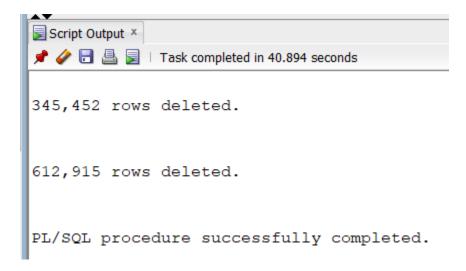
DELETE FROM
comments
WHERE
comment_date > 1300000000;
```

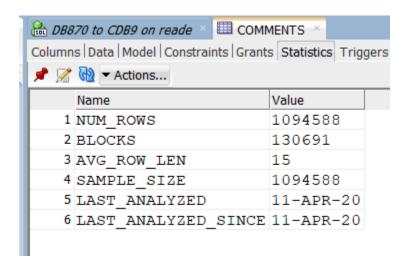


```
DELETE FROM
comments
WHERE
comment_date < 1150000000;

DELETE FROM
comments
WHERE
comment_date > 1250000000;

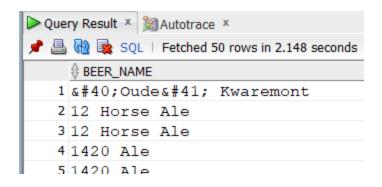
EXECUTE dbms stats.gather table stats('DB870', 'beer reviews');
```

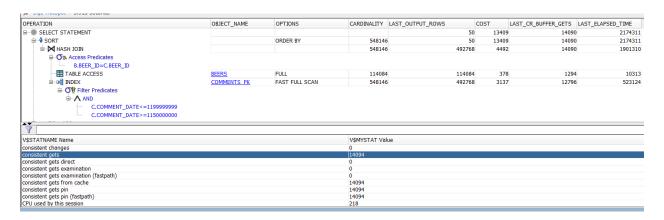




That left nearly 1.1 million rows between date values 1150000000 and 1250000000. I ran a baseline query prior to partitioning.

```
SELECT
b.beer_name
FROM
beers b
INNER JOIN comments c
ON b.beer_id = c.beer_id
WHERE
c.comment_date BETWEEN 1150000000 AND 119999999
ORDER BY
b.beer_name ASC;
```





Then I tried to create a new table from the comments table.

```
CREATE TABLE
pcomments
PARTITION BY RANGE (comment_date) INTERVAL (25000000)
(PARTITION p1 VALUES LESS THAN (1175000000),
PARTITION p2 VALUES LESS THAN (1200000000),
PARTITION p3 VALUES LESS THAN (1225000000),
PARTITION p4 VALUES LESS THAN (1250000000))
AS
SELECT
*
FROM
comments;
```

Sadly, I got the following error.

ORA-00997: illegal use of LONG datatype 00997. 00000 - "illegal use of LONG datatype"

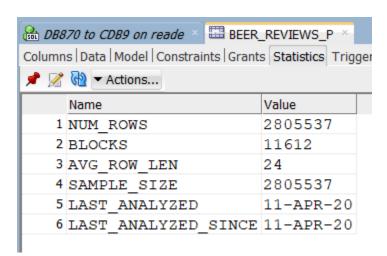
At this point, I tried very hard to forget I ever created the comments table. Two valuable lessons were learned. First, plan accordingly and create your partitions up front. Second, never use the LONG data type; use xLOB instead.

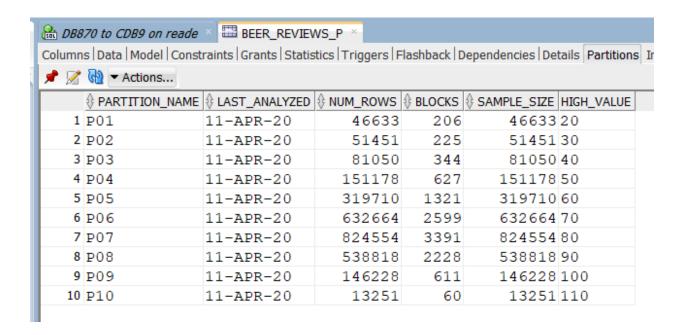
Moving on, I decided to set my sights on the beer_reviews table. It too was very large in terms of row count, but the block count was much smaller.

```
CREATE TABLE
  beer_reviews_p
  PARTITION BY RANGE (overall rating) INTERVAL (10)
  (PARTITION p01 VALUES LESS THAN (20),
  PARTITION p02 VALUES LESS THAN (30),
  PARTITION p03 VALUES LESS THAN (40),
  PARTITION p04 VALUES LESS THAN (50),
  PARTITION p05 VALUES LESS THAN (60),
  PARTITION p06 VALUES LESS THAN (70),
  PARTITION p07 VALUES LESS THAN (80),
  PARTITION p08 VALUES LESS THAN (90),
  PARTITION p09 VALUES LESS THAN (100),
  PARTITION p10 VALUES LESS THAN (110))
AS
  SELECT
  FROM
```

beer_reviews;

EXECUTE dbms_stats.gather_table_stats('DB870', 'beer_reviews_p');





ALTER TABLE

beer reviews p

ADD CONSTRAINT beer reviews p pk PRIMARY KEY (reviewer id, beer id);

ALTER TABLE

beer_reviews_p

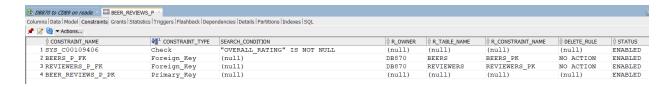
ADD CONSTRAINT reviewers p fk FOREIGN KEY (reviewer id) REFERENCES reviewers(reviewer id);

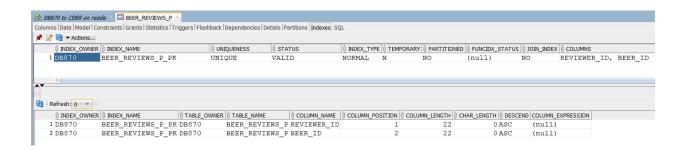
ALTER TABLE

beer_reviews_p

ADD CONSTRAINT beers_p_fk FOREIGN KEY (beer_id) REFERENCES beers(beer_id);

The NOT NULL constraint on the overall rating attribute carried over to the new table automatically.





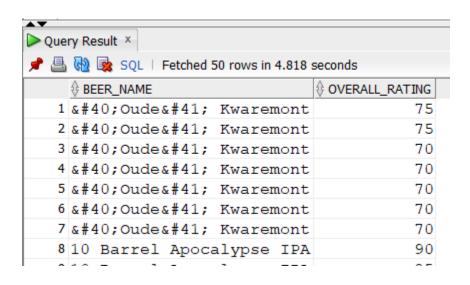
The only remaining step was to configure the bitmap indexes for the new table.

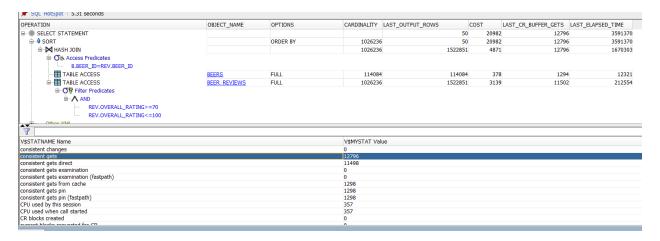
```
CREATE BITMAP INDEX
 povr_bmap
ON
 beer_reviews_p(overall_rating) LOCAL;
CREATE BITMAP INDEX
 papr_bmap
ON
 beer_reviews_p(appearance_rating) LOCAL;
CREATE BITMAP INDEX
 parr_bmap
ON
 beer_reviews_p(aroma_rating) LOCAL;
CREATE BITMAP INDEX
 ppar bmap
ON
 beer_reviews_p(palate_rating) LOCAL;
CREATE BITMAP INDEX
 ptar bmap
ON
 beer_reviews_p(taste_rating) LOCAL;
EXECUTE dbms_stats.gather_table_stats('DB870', 'beer_reviews_p');
```

olumns Data Model Co	nstraints Grants Statistics Triggers	Flashback Dependencies	Details Partitions Index	ces SQL					
Actions									
									COLUMNS
1 DB870	PAPR_BMAP	NONUNIQUE	N/A	BITMAP	N	YES	(null)	NO	APPEARANCE_RATING
2 DB870	PARR_BMAP	NONUNIQUE	N/A	BITMAP	N	YES	(null)	NO	AROMA_RATING
3 DB870	POVR_BMAP	NONUNIQUE	N/A	BITMAP	N	YES	(null)	NO	OVERALL_RATING
4 DB870	PPAR BMAP	NONUNIQUE	N/A	BITMAP	N	YES	(null)	NO	PALATE RATING
5 DB870	PTAR BMAP	NONUNIQUE	N/A	BITMAP	N	YES	(null)	NO	TASTE RATING
6 DB870	BEER REVIEWS P PK	UNIQUE	VALID	NORMAL	N	NO	(null)	NO	REVIEWER ID, BEER I

At this point, I gave some thought to previously planned next steps. In a production environment, it would make sense to backup and then drop (or simply rename) the old table and then rename the new table to assume the old table's name. However, this is a lab environment, and we have plenty of space since I deleted over half the rows in the comments table. So, I decided to keep the old table, and I left the comments table unchanged (i.e., still pointing to the beer_reviews table for its FK constraints). With this schema, we can run queries against the old and new review tables for comparison. So, that is what I did next.

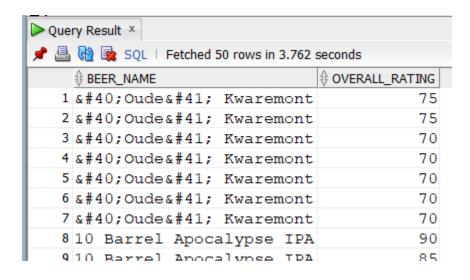
```
SELECT
b.beer_name,
rev.overall_rating
FROM
beers b
INNER JOIN beer_reviews rev
ON b.beer_id = rev.beer_id
WHERE
rev.overall_rating BETWEEN 70 AND 100
ORDER BY
b.beer_name ASC,
rev.overall_rating DESC;
```

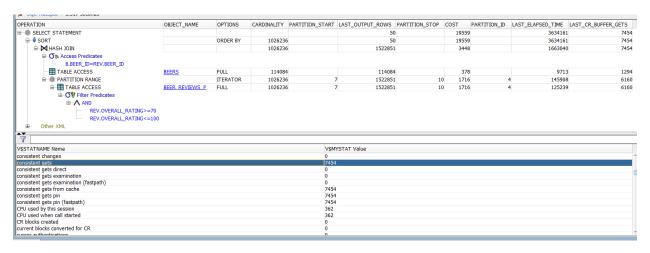




SELECT
b.beer_name,
rev.overall_rating
FROM

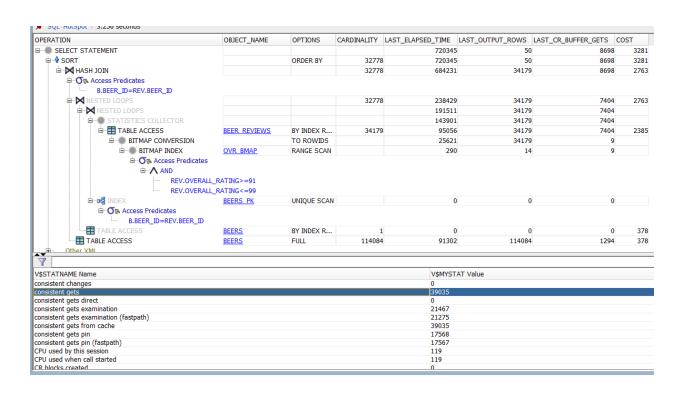
beers b
INNER JOIN beer_reviews_p rev
ON b.beer_id = rev.beer_id
WHERE
rev.overall_rating BETWEEN 70 AND 100
ORDER BY
b.beer_name ASC,
rev.overall_rating DESC;





As we can see, the execution plan for the partitioned table made use of the partitions. The partitioned table performed better with significantly fewer consistent_gets and a much faster time to return results (3.762 seconds versus 4.818 seconds). However, the last_elapsed_time was slightly higher for the partitioned table. According to Oracle documentation, that column represents the "Elapsed time (in microseconds) corresponding to this operation, during the last execution." So, I am confused how it can be higher for the partitioned table. I decided to try a more selective query.

```
SELECT
  b.beer_name,
  rev.overall_rating
FROM
  beers b
  INNER JOIN beer_reviews rev
    ON b.beer_id = rev.beer_id
WHERE
  rev.overall_rating BETWEEN 91 AND 99
ORDER BY
  b.beer_name ASC,
  rev.overall_rating DESC;
SELECT
  b.beer_name,
  rev.overall_rating
FROM
  beers b
  INNER JOIN beer_reviews_p rev
    ON b.beer_id = rev.beer_id
WHERE
  rev.overall_rating BETWEEN 91 AND 99
ORDER BY
  b.beer_name ASC,
  rev.overall_rating DESC;
```

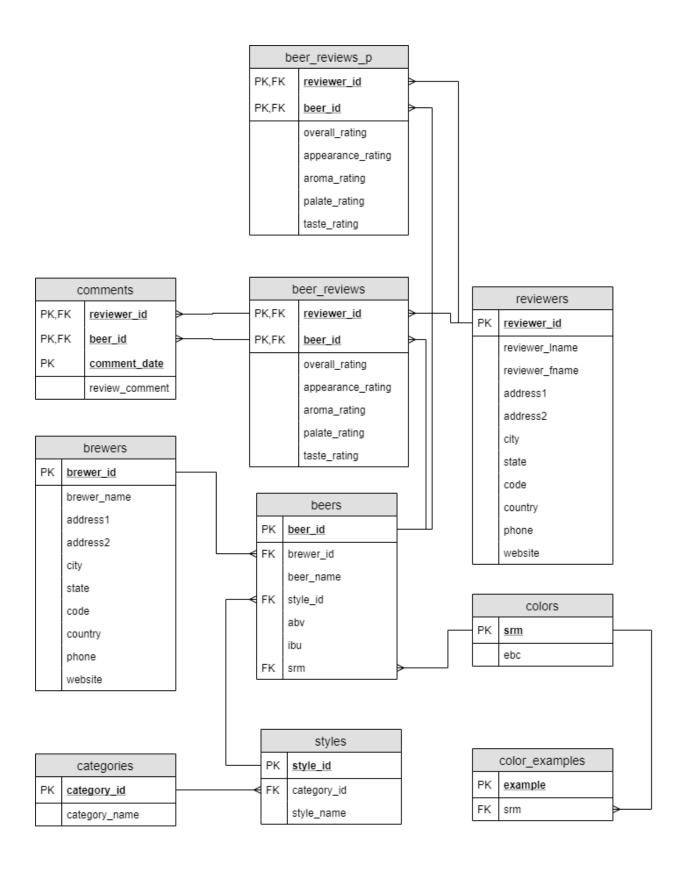


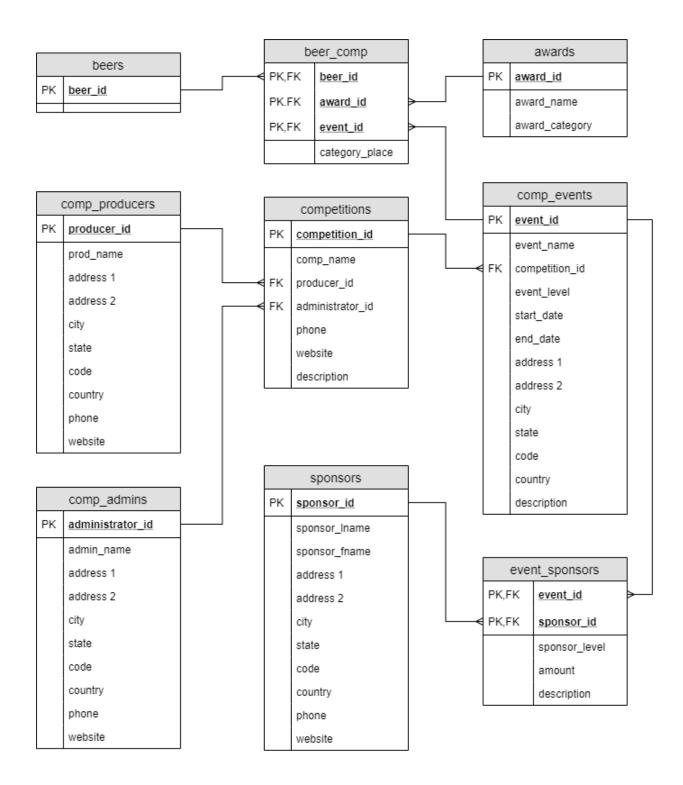
PERATION	OBJECT_NAME	OPTIONS	CARDINALITY	PARTIT	LAST_OUTPUT_ROWS	PARTITION_STOP	LAST_CR_BUFFER_GETS	COST	PARTITION_ID	LAST_ELAPSED_TIME
SELECT STATEMENT					50)	1889	1490		74704
⊕ • SORT		ORDER BY	59839)	50)	1889	1490		74704
			59839		34179		1889	546		70881
B.BEER_ID=REV.BEER_ID			50000		0.4476		505			20000
			59839		34179		595			23008
					34179		595			18121
					34179		595			13210
		SINGLE	34179				595			6 8218
	BEER REVIEWS P	FULL	34179	9	34179	!	595	168		7 3309
□ ○ □ ○ □ □ □ □ □ □ □ □ □ □ □ □ □ □ □ □ □ □ □ □										
□ · ∧ AND										
REV.OVERALI	L_RATING>=91									
REV.OVERALI	L_RATING<=99									
⊞- u ∰ INDEX	BEERS PK	UNIQUE SCAN			()	()		
⊕ Ot Access Predicates										
B.BEER ID=REV.BEER II	D									
TABLE ACCESS	BEERS	BY INDEX R	2		()	(378		
TABLE ACCESS	BEERS	FULL	114084		114084		1294	378		9989
		,								
7										
\$STATNAME Name					V\$MYSTAT Value					
onsistent changes					0					
onsistent gets					1902					
onsistent gets direct onsistent gets examination					0					
onsistent gets examination (fastpath)					0					
onsistent gets from cache					1902					
onsistent gets pin					1902					
onsistent gets pin (fastpath)					1902					
PU used by this session					83					
					83					
PU used when call started R blocks created					0					

I ran each query a dozen times. The first query averaged 0.088 seconds to return results. The second query (partitioned) averaged 0.07 seconds. That equates to a 20% performance improvement over the non-partitioned query. Again, we see the last_elapsed_time is higher for the partitioned table on the last run of the query. I did not average that statistic across all runs, so it might have been lower on average for the partitioned table. However, it seems strange that it would ever be higher for the partitioned table given the large performance advantage of the partitioned table in terms of time to return results to the client. In any event, I counted this as a win, which brought my score to two for four in Assignment 4.

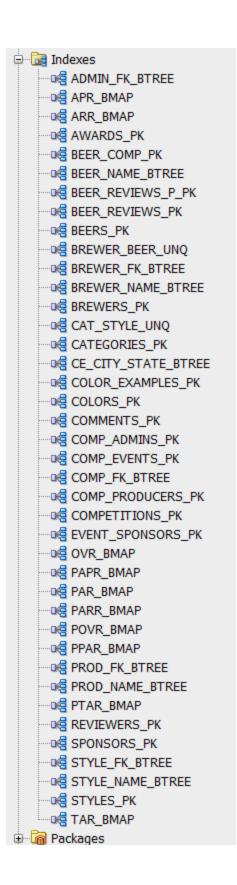
Note the access path to the beer_reviews table was the ovr_bmap index, which significantly increased performance versus a FULL SCAN, thereby offsetting some of the partitioning advantage. The access path to the beer_reviews_p table was a FULL SCAN of partition p9 (owing to the relatively small size of p9 I suppose).

The updated ERD follows and shows the addition of the partitioned table.





The updated list of indexes follows.



Other Topics

Loading Bulk Data via Python

For my other topic, I chose to experiment with Python, cx_Oracle, and Oracle Instant Client to load the RateBeer dataset into the database. I used Spyder 3.3.6 and Python 3.7.4 (both installed by Anaconda3) with cx_Oracle 7.3.0 and Oracle Instant Client (Basic Light) 19.5.0.0.0. I had other options for loading the RateBeer dataset, but I saw this as an opportunity to improve my limited Python skills. Since I am new to coding, experienced coders will likely see many potential improvements in my code. Nonetheless, here it is.

```
# -*- coding: utf-8 -*-
Created on Tue Mar 17 22:40:00 2020
@author: James Long
import cx_Oracle
import re
# Connect as user "DB870" with password <redacted>
# running on <redacted>
conn = cx_Oracle.connect("DB870", <redacted>, <redacted>)
cur = conn.cursor()
#rbfh = open('ratebeer_sample.txt') # test dataset, first 5 records
rbfh = open('ratebeer db.txt') # full dataset
#rbfh = open('ratebeer_db_trimmed.txt') # partial dataset to speed disconnect recovery
# if a restart is needed due to disconnect before the script completes, you can
# comment the following lines down to (but not including) the error logging block
cur.execute(""
      SELECT
        beer id
      FROM
        beers
      111
      )
beer_id_lst = list(cur.fetchall())
cur.execute(""
      SELECT
        brewer_id,
        beer_name
      FROM
        beers
```

```
beer_unq_lst = list(cur.fetchall())
cur.execute(""
      SELECT
         style_name,
         style_id
      FROM
        styles
      )
style_lst = list(cur.fetchall())
style_dict = {}
for tup in style_lst:
  style_dict[tup[0]] = tup[1]
cur.execute(""
      SELECT
         style_id
      FROM
         styles
      )
style_lst = list(cur.fetchall())
new_style_id = int(max(style_lst)[0])
del style_lst
cur.execute(""
      SELECT
         reviewer_Iname,
         reviewer_id
      FROM
         reviewers
reviewer_lst = list(cur.fetchall())
reviewer_dict = {}
for tup in reviewer_lst:
  reviewer_dict[tup[0]] = tup[1]
del reviewer_lst
cur.execute(""
      SELECT
         brewer_id
       FROM
         brewers
       111
brewer_lst = list(cur.fetchall())
```

```
cur.execute(""
      SELECT
         reviewer_id,
         beer_id,
         comment_date
       FROM
         comments
comment_lst = list(cur.fetchall())
cur.execute(""
      SELECT
         reviewer_id,
         beer_id
       FROM
         beer_reviews
beer_review_lst = list(cur.fetchall())
# error logging
reviewer_err_lst = list()
reviewer_err_cnt = 0
brewer_err_lst = list()
brewer_err_cnt = 0
style_err_lst = list()
style_err_cnt = 0
beer_err_lst = list()
beer_err_cnt = 0
beer_review_err_lst = list()
beer_review_err_cnt = 0
comment_err_lst = list()
comment_err_cnt = 0
print()
print('****** BEGIN ******')
print()
for line in rbfh:
  if len(line) == 1:
    continue
  line = line.encode('ascii','replace').decode('ascii')
  if line.startswith('?'):
    line = line.translate(line.maketrans(", ", '?'))
  if line.startswith('beer/name'):
    try:
```

```
beer_name = str(re.findall('^beer/name: (.+)', line)[0])
    continue
  except:
    beer_name = 0
    continue
elif line.startswith('beer/beerId'):
  try:
    beer_id = int(re.findall('^beer/beerId: (.+)', line)[0])
    continue
  except:
    beer id = 0
    continue
elif line.startswith('beer/brewerId'):
  try:
    brewer_id = int(re.findall('^beer/brewerld: (.+)', line)[0])
    continue
  except:
    brewer id = -1
    continue
elif line.startswith('beer/ABV'):
  try:
    abv = float(re.findall('^beer/ABV: (.+)', line)[0])
    continue
  except:
    abv = 0
    continue
elif line.startswith('beer/style'):
    style_name = str(re.findall('^beer/style: (.+)', line)[0])
    continue
  except:
    style_name = 'unknown'
    continue
elif line.startswith('review/appearance'):
    appearance_rating = str(re.findall('^review/appearance: (.+)', line)[0])
    appearance rating = str(round((float(appearance rating[0])/float(appearance rating[2]))*100,0))
    continue
  except:
    appearance_rating = None
    continue
elif line.startswith('review/aroma') and len(line) == 19:
    aroma_rating = str(re.findall('^review/aroma: (.+)', line)[0])
    aroma_rating = str(round((float(aroma_rating[0])/float(aroma_rating[2:]))*100,0))
    continue
  except:
    aroma_rating = None
```

```
continue
elif line.startswith('review/aroma') and len(line) == 20:
  try:
    aroma rating = str(re.findall('^review/aroma: (.+)', line)[0])
    aroma rating = str(round((float(aroma rating[:2])/float(aroma rating[3:]))*100,0))
    continue
  except:
    aroma_rating = None
    continue
elif line.startswith('review/palate'):
  try:
    palate_rating = str(re.findall('^review/palate: (.+)', line)[0])
    palate_rating = str(round((float(palate_rating[0])/float(palate_rating[2]))*100,0))
    continue
  except:
    palate_rating = None
    continue
elif line.startswith('review/taste') and len(line) == 19:
  try:
    taste_rating = str(re.findall('^review/taste: (.+)', line)[0])
    taste rating = str(round((float(taste rating[0])/float(taste rating[2:]))*100,0))
    continue
  except:
    taste_rating = None
    continue
elif line.startswith('review/taste') and len(line) == 20:
  try:
    taste rating = str(re.findall('^review/taste: (.+)', line)[0])
    taste_rating = str(round((float(taste_rating[:2])/float(taste_rating[3:]))*100,0))
    continue
  except:
    taste_rating = None
    continue
elif line.startswith('review/overall') and len(line) == 21:
    overall_rating = str(re.findall('^review/overall: (.+)', line)[0])
    overall rating = str(round((float(overall rating[0])/float(overall rating[2:]))*100,0))
    continue
  except:
    overall rating = -1
    continue
elif line.startswith('review/overall') and len(line) == 22:
    overall rating = str(re.findall('^review/overall: (.+)', line)[0])
    overall rating = str(round((float(overall rating[:2])/float(overall rating[3:]))*100,0))
    continue
  except:
    overall rating = -1
```

```
continue
elif line.startswith('review/time'):
    comment_date = int(re.findall('^review/time: (.+)', line)[0])
    continue
  except:
    comment_date = -1
    continue
elif line.startswith('review/profileName'):
    reviewer_Iname = str(re.findall('^review/profileName: (.+)', line)[0])
    continue
  except:
    reviewer_Iname = 0
    continue
elif line.startswith('review/text'):
    review_comment = str(re.findall('^review/text: (.+)', line)[0])
  except:
    review_comment = None
if reviewer_lname != 0:
  reviewer_id = reviewer_dict.get(reviewer_lname,0)
  if reviewer_id == 0:
    try:
      cur.execute(""
            INSERT INTO
               reviewers(reviewer_Iname)
             VALUES
               (:reviewer_Iname)
             [reviewer_lname]
      conn.commit();
      cur.execute(""
             SELECT
               reviewer_id
             FROM
               reviewers
             WHERE
               reviewer_Iname = :reviewer_Iname
             [reviewer_Iname]
      reviewer_dict[reviewer_lname] = cur.fetchone()[0]
      reviewer id = reviewer dict.get(reviewer lname)
    except Exception as nam_err:
      reviewer_err_lst.append(reviewer_lname)
```

```
reviewer_err_cnt = reviewer_err_cnt + 1
      print('reviewer error',reviewer_err_cnt)
      print(nam_err)
if (brewer id,) not in brewer lst:
  try:
    cur.execute(""
           INSERT INTO
             brewers(brewer_id)
           VALUES
             (:brewer_id)
           ш,
           [brewer_id]
    conn.commit();
    brewer_lst.append((brewer_id,))
  except Exception as brew err:
    brewer_err_lst.append(brewer_id)
    brewer_err_cnt = brewer_err_cnt + 1
    print('brewer error',brewer_err_cnt)
    print(brew err)
style_id = style_dict.get(style_name,0)
if style_id == 0:
  new_style_id = new_style_id + 1
  style_id = new_style_id
  try:
    cur.execute("
           INSERT INTO
             styles(style_id,category_id,style_name)
           VALUES
             (:style_id,-1,:style_name)
           [style_id, style_name]
    conn.commit();
    style_dict[style_name] = style_id
  except Exception as sty_err:
    style err lst.append(style id)
    style_err_cnt = style_err_cnt + 1
    print('style error',style_err_cnt)
    print(sty_err)
if beer name == 0:
  continue
if beer id == 0:
  continue
if (brewer_id, beer_name) in beer_unq_lst:
```

```
continue
  if (beer_id,) not in beer_id_lst:
    try:
      cur.execute(""
             INSERT INTO
               beers(beer_id, brewer_id, beer_name, style_id, abv, ibu, srm)
             VALUES
               (:beer_id, :brewer_id, :beer_name, :style_id, :abv, 0, -1)
             [beer_id, brewer_id, beer_name, style_id, abv]
      conn.commit();
      beer_id_lst.append((beer_id,))
    except Exception as beer err:
      beer_err_lst.append(beer_id)
      beer_err_cnt = beer_err_cnt + 1
      print('beer error',beer err cnt)
      print(beer_err)
  if reviewer_Iname == 0:
    continue
  if overall_rating == -1:
    continue
  if (reviewer_id, beer_id) not in beer_review_lst:
    try:
      cur.execute(""
             INSERT INTO
               beer reviews(reviewer id, beer id, overall rating, appearance rating, aroma rating,
palate_rating, taste_rating)
             VALUES
               (:reviewer_id, :beer_id, :overall_rating, :appearance_rating, :aroma_rating,
:palate_rating, :taste_rating)
             [reviewer_id, beer_id, overall_rating, appearance_rating, aroma_rating, palate_rating,
taste rating]
      conn.commit();
      beer_review_lst.append((reviewer_id, beer_id))
    except Exception as rev err:
      beer review err lst.append((reviewer id, beer id))
      beer_review_err_cnt = beer_review_err_cnt + 1
      print('beer review error',beer_review_err_cnt)
      print(rev_err)
  if comment date == -1:
    continue
  if (reviewer id, beer id, comment date) not in comment lst:
    try:
```

```
len(review_comment)
      try:
        cur.execute(""
               INSERT INTO
                 comments(reviewer id, beer id, comment date, review comment)
               VALUES
                 (:reviewer_id, :beer_id, :comment_date, :review_comment)
               [reviewer_id, beer_id, comment_date, review_comment]
        conn.commit();
        comment_lst.append((reviewer_id, beer_id, comment_date))
      except Exception as com_err:
        comment_err_lst.append((reviewer_id, beer_id, comment_date))
        comment_err_cnt = comment_err_cnt + 1
        print('comment error',comment_err_cnt)
        print(com_err)
    except:
      continue
conn.commit()
cur.close()
Before starting, I performed a full scan of the file using the following code.
# -*- coding: utf-8 -*-
Created on Tue Mar 23 22:40:00 2020
@author: James Long
import re
review_lst = list()
rec cnt = 0
rev_err = 0
beer err = 0
rbfh = open('RateBeer_db.txt')
for line in rbfh:
  if len(line) == 1:
    continue
  line = line.encode('ascii','replace').decode('ascii')
  if line.startswith('?'):
    line = line.translate(line.maketrans(", ", '?'))
```

```
if line.startswith('beer/beerId'):
    try:
      beer id = str(re.findall('^beer/beerld: (.+)', line)[0])
      continue
    except:
       beer id = None
       beer_err = beer_err + 1
      continue
  elif line.startswith('review/profileName'):
    rec cnt = rec cnt + 1
    try:
       reviewer Iname = str(re.findall('^review/profileName: (.+)', line)[0])
      if (reviewer_Iname, beer_id) not in review_lst:
         review_lst.append((reviewer_lname, beer_id))
      continue
    except:
       reviewer Iname = None
      rev_err = rev_err + 1
      continue
print()
print('unique reviews =',len(review lst))
print('total records =',rec cnt)
```

The scan counted 2,924,163 total records. That is 36 more than the 2,924,127 reviews reported by Stanford SNAP. The scan also counted 2,855,232 unique reviews. The lower unique review count indicates some reviewers reviewed the same beer more than once. Since the requirements were to store a single review per beer per reviewer, I did not insert duplicate reviews. (Doing so was not possible based on the composite key for the beer_reviews table.) In case of duplicate reviews, the first one encountered in the dataset was inserted. Without writing additional code to log these duplicate records and compare their timestamps, I cannot say whether the first, intermediate, or last review was inserted. This might be an interesting exercise for future students. The beer_err and rev_err counts were both zero, which indicated no records were missing beerld or profileName values, and no such values were corrupt.

I cached certain data from the database upon startup to reduce total run time by eliminating millions of roundtrips to the server. Caching also made resumption of new record insertion quicker whenever the script had to be restarted. As it happens, my Internet connection dropped several times while the script was running, and I had to restart the script. So, these code changes really saved a lot of time. However, once my script reached around 1.7 million inserted records, restarting took so long that my database connection timed out while my script was reading through the locally cached records to locate the first uninserted record. To get around this, I wrote the following code to trim the first x million records from the dataset file (where x was set for each restart as close as possible to the actual number of inserted rows in the beer_reviews table). I then used the trimmed file to resume record insertion.

```
# -*- coding: utf-8 -*-
```

```
@author: james
"""

oldfh = open('ratebeer_db.txt')

newfh = open('ratebeer_db_trimmed.txt', mode='w')

linecnt = 0

for line in oldfh:
    linecnt = linecnt + 1
    if linecnt <= 28000000:
        continue
    else:
        newfh.write(line)

newfh.close()</pre>
```

With my early code, the BREWER_BEER_UNQ constraint was violated 48 times during insertion of new records. Since that constraint is based on the (brewer_id, beer_name) tuple in the beers table, two possible explanations exist. First, if the records generating these errors were indeed duplicates in the RateBeer dataset, then each would have a beerld matching the beer_id of a previously inserted record. Thus, these errors could not have occurred due to the beer_id check condition already in my code at that time. However, it is possible that a data entry error occurred when the RateBeer dataset was created, and a new beerld was assigned to certain records that already existed, thereby creating duplicate records. If this were the case, the choice would come down to inserting such records into the database to capture as much data as possible versus dropping such records to preserve the data integrity of the database. Given the total number of records, I decided it was better to drop some data than to introduce unreliable or misleading data into the database.

The other possibility is that each record represents a distinct beer (as indicated by the unique beerId) that shares a common name with another distinct beer produced by the same brewer. Such beers could be distinguishable (for the benefit of consumers) by some other characteristic, such as the color of the can/bottle or a sub-name that was excluded from the RateBeer dataset. For example, Summertime Ale - Lemon and Summertime Ale - Lime might both have been entered as Summertime Ale. If this were the case, the choice would come down to modifying such records with distinguishing data prior to inserting them into the database versus dropping such records versus dropping the BREWER_BEER_UNQ constraint. Dropping the BREWER_BEER_UNQ constraint would permit other illegitimate cases of non-unique (brewer_id, beer_name) records to be inserted, so I decided against that option. And since I did not know the actual feature that distinguished these beers, I had no way to reliably modify the records (e.g., by appending the missing name data to the beer name). Thus, I decided to drop such records.

The following debug data were collected for the 48 errors in case future students want to investigate this further. The fact that each offending beerld occurs more than once suggests that the root cause was

not entry of an erroneous beerld. Rather, the associated brewerld and/or name values are likely the cause.

	1
Error Count	beerld
1	19336
2	19336
3	19336
4	19336
5	19336
6	19336
7	19336
8	19336
9	19336
10	19336
11	19336
12	19336
13	19336
14	19336
15	102105
16	102105
17	102105
18	102105
19	102105
20	102105
21	102105
22	102105
23	102105
24	102105
25	102105
26	102105
27	102105
28	102105
29	102105
30	102105
31	102105
32	102105
33	102105
34	102105
35	102105
36	102105
37	102105
38	102105
39	102105
40	102105
41	102105
42	102105

43	102105
44	102105
45	102105
46	102105
47	102105
48	102105

Once my code was able to insert several tens of thousands of records with no errors, I dropped all newly inserted records from the database and re-ran my script to ensure only clean data was inserted. The script ran for 8+ days excluding downtime due to disconnects. In total, the effort took three weeks. Unfortunately, I was unable to complete the load process because the STUDENTS tablespace filled up near the very end of my script.

ORA-01653: unable to extend table DB870.COMMENTS by 8192 in tablespace STUDENTS

The following table shows the number of records in each of the affected tables upon completion of the load process.

Table	Total Records	Pre-existing Records	New Records
beers	114,084	5,890	108,194
brewers	8,789	1,414	7,375
styles	226	141	85
reviewers	29,097	0	29,097
beer_reviews	2,805,537	0	2,805,537
comments	2,775,682	0	2,775,682

The total comment count is lower than the total review count because some reviews did not have a comment. While it was possible that some reviews had more than one comment, that turned out not to be the case as indicated by the following query, which returned zero rows.

```
SELECT
c1.beer_id,
c1.reviewer_id,
c1.comment_date
FROM
comments c1
INNER JOIN comments c2
ON c1.beer_id=c2.beer_id
AND c1.reviewer_id=c2.reviewer_id
WHERE
c1.comment_date<>c2.comment_date;
```

If the review/time values in the RateBeer dataset represented distinct days without a time stamp (i.e., date only), then my code was only able to capture multiple comments by a single reviewer about a single beer if the comments were made on separate days. Multiple comments submitted on the same day by the same reviewer about the same beer would have the same review/time value. Thus, my code would have inserted only the first of such comments encountered in the dataset. There is no way to know whether the inserted comment was the first, intermediate, and last comment made on that day about that beer by that reviewer. However, if the time stamps were also embedded in the review/time values, then same-day comments would have unique review/time values, and my code captured all comments by the same reviewer about the same beer.

I used the following SQL statement in Oracle SQL Developer to monitor the Python session.

```
SELECT

program,
event,
p2text,
wait_class,
state,
wait_time_micro,
time_since_last_wait_micro
FROM
v$session
WHERE
username = 'DB870';
```