

http://www.rlvision.com/blog/method-for-determining-candidate-keys-and-highest-normal-form-of-a-relation-based-on-functional-dependencies/

02 Database normalization

Database normalization

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02 Database normalization

First normal form (1NF)

Relational schema is in 1NF if all occurrences of rows in the respective relational table contain the same number of fields and include the atomic values only, (there is no repeating fields and groups)

e#	name	car used
950001	Peter	Toyota, PKR234, Ford, WER545
932345	Paul	Honda, RTQ456
960020	Joan	Holden, KLR197, Holden, KLR567

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Keys

Superkey
Superkey is a nonempty subset X of relational schema $R = (A_1, \dots, A_n)$ such that for any two rows t_1, t_2 in a relational table defined over R $t_1[X] \neq t_2[X]$
If X is a superkey in R then $X \rightarrow (A_1, \dots, A_n)$

Minimal key
Minimal key is a superkey K with an additional property such that removal of any attribute from K will cause K not to be a superkey

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Keys

Primary key
Primary key is an arbitrarily selected minimal key

Candidate key
Candidate key is any other minimal key which is not primary

important key def here

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Functional dependencies and keys

Let $R = (A_1, \dots, A_n)$ be a relational schema (a header of relational table) and let X, Y be nonempty subsets of R such that $X \cup Y = R$

- (1) If $X \rightarrow Y$ is valid in R then X is a key (!!!)
- (2) If X is a key then $X \rightarrow Y$ is valid in R

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Attributes

Prime attribute
Prime attribute is an attribute from relational schema R which is a member of at least one candidate key in R

Nonprime attribute
Nonprime attribute is an attribute which is a not prime

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Full and partial functional dependencies

Full functional dependency

Full functional dependency is a functional dependency $X \rightarrow Y$ such that removal of any attribute A from X causes that $(X-A) \not\rightarrow Y$

Partial functional dependency

Partial functional dependency is a functional dependency which is not full

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Second normal form (2NF)

Relational schema R is in 2NF if every nonprime attribute A in R is fully functionally dependent on a primary key of schema R

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Second normal form (2NF)

Inventory

part	quantity	warehouse	warehouse-address
------	----------	-----------	-------------------

$\text{warehouse} \rightarrow \text{warehouse-address}$

$\text{part, warehouse} \rightarrow \text{quantity}$

$\text{part, warehouse} \rightarrow \text{warehouse-address}$

$\text{part, warehouse} \rightarrow \text{warehouse-address, quantity}$

Minimal key = (part, warehouse)

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Second normal form (2NF)

Inventory

part	quantity	warehouse	warehouse-address
------	----------	-----------	-------------------

$\text{warehouse} \rightarrow \text{warehouse-address}$

$\text{part, warehouse} \rightarrow \text{quantity}$

Minimal key = (part, warehouse)

Schema **Inventory** is **not in 2NF** because nonprime attribute **warehouse-address** depends on a part (warehouse) of a key (part, warehouse)

Functional dependency which "violates 2NF" is **warehouse \rightarrow warehouse-address**

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Second normal form (2NF)

Inventory

part	quantity	warehouse	warehouse-address
------	----------	-----------	-------------------

$\text{warehouse} \rightarrow \text{warehouse-address}$

$\text{part, warehouse} \rightarrow \text{quantity}$

Minimal key = (part, warehouse)

If all minimal keys in a schema consist of only one attribute (single attribute keys) then such schema is always in 2NF

This is because any nonprime attribute in the schema does not depend on a part of a key (because each key consists of one attribute only !)

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Second normal form (2NF)

Inventory

part	quantity	warehouse	warehouse-address
------	----------	-----------	-------------------

Schema **Inventory** should be decomposed into the following schemas

Store

part	quantity	warehouse
------	----------	-----------

Location

warehouse	warehouse-address
-----------	-------------------

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Second normal form (2NF)

Store

part quantity warehouse

 $\text{part, warehouse} \rightarrow \text{quantity}$

Minimal key = (part, warehouse)

Schema **Store** is in 2NF because no nonprime attributes (quantity) depends on a part of a key (part, warehouse)

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Second normal form (2NF)

Location

warehouse warehouse-address

 $\text{warehouse} \rightarrow \text{warehouse-address}$

Minimal key = (warehouse)

Schema **Location** is in 2NF because no nonprime attributes (warehouse-address) depends on a part of a key (warehouse)

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Second normal form (2NF)

Location

warehouse warehouse-address

 $\text{warehouse} \rightarrow \text{warehouse-address}$

Minimal key = (warehouse)

Every relational schema, which consists of at most 2 attributes is always in 2NF

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Transitive functional dependency

Functional dependency $X \rightarrow Y$ valid in a relational schema R is a transitive functional dependency if there exists a nonempty subset Z of R , that is not a subset of any key in R and such that the functional dependencies $X \rightarrow Z$ and $Z \rightarrow Y$ are valid in R

We say that Y is transitively dependent on X in schema R if $X \rightarrow Y$ is valid in R and $X \rightarrow Y$ is a transitive functional dependency

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Third normal form (3NF)

Relational schema R is in 3NF if it is in 2NF and no nonprime attribute of R is transitively dependent on the primary key

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Third normal form (3NF)

Supplier

s# sname company-name city

 $s\# \rightarrow sname$ $s\# \rightarrow company-name$ $s\# \rightarrow city$ $company-name \rightarrow city$

Minimal key = (s#)

Schema **Supplier** is in 2NF because there are no nonprime attributes which are functionally dependent on a part of minimal key $s\#$

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02 Database normalization

Third normal form (3NF)

Supplier			
s#	sname	company-name	city

$s\# \rightarrow sname$
 $s\# \rightarrow company-name$
 $s\# \rightarrow city$
 $company-name \rightarrow city$
Minimal key = (s#)

Schema **Supplier** is not in 3NF because attribute **city** is transitively dependent on **s#**
 $s\# \rightarrow company-name \ \& \ company-name \rightarrow city$

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Third normal form (3NF)

Supplier			
s#	sname	company-name	city

Schema **Supplier** should be decomposed into the following schemas

Supplier-new		
s#	sname	company-name

Company	
company-name	city

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02 Database normalization

Third normal form (3NF)

Supplier-new		
s#	sname	company-name

$s\# \rightarrow sname$
 $s\# \rightarrow company-name$
Minimal key = (s#)

Schema **Supplier-new** is in 3NF because no attribute is transitively dependent on **s#**

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Third normal form (3NF)

Company	
company-name	city

$company-name \rightarrow city$
Minimal key = (company-name)

Schema **Company** is in 3NF because no attribute is transitively dependent on **company-name**

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Third normal form (3NF)

Company	
company-name	city

$company-name \rightarrow city$

Every relational schema, which consists of at most 2 attributes is always in 3NF

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Third normal form (3NF)

Alternative definition
A relational schema **R** is in 3NF if whenever a functional dependency $X \rightarrow A$ is valid in **R** then either:

- (1) **X** is a superkey in **R**, or
- (2) **A** is a prime attribute in **R**

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02 Database normalization

Third normal form (3NF)

Supplier

s#	sname	company-name	city
----	-------	--------------	------

 $s\# \rightarrow sname$ $s\# \rightarrow company-name$ $s\# \rightarrow city$ $company-name \rightarrow city$

Minimal key = (s#)

Schema Supplier is not in 3NF because

- (1) attribute company_name is not a superkey and
 (2) attribute city is not prime

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Third normal form (3NF)

Location

city	street	zip-code
------	--------	----------

 $city, street \rightarrow zip-code$ $zip-code \rightarrow city$ Minimal key₁ = (city, street)Minimal key₂ = (zip-code, street)

Schema Location is in 3NF because attribute city used
 on the right hand side of functional dependency
 $zip-code \rightarrow city$ is prime

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Third normal form (3NF)

Location

city	street	zip-code
NY	55	484
NY	56	484
LA	55	473
LA	56	473
LA	57	474

Repetition of [LA ... 473]
 and [NY ... 484] is forced
 by functional dependency
 $zip-code \rightarrow city$

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Boyce-Codd normal form (BCNF)

A relational schema R is in BCNF if whenever functional
 dependency $X \rightarrow A$ holds in R then:

- (1) X is a superkey in R

- ~~(2) or A is a prime attribute in R~~ (3NF only !!!)

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Boyce-Codd normal form (BCNF)

Location

city	street	zip-code
------	--------	----------

 $city, street \rightarrow zip-code$ $zip-code \rightarrow city$ Minimal key₁ = (city, street)Minimal key₂ = (zip-code, street)

Schema Location is not in BCNF because zip-code is
 not a superkey

 $zipcode \rightarrow city \ \& \ zip_code \not\rightarrow street$

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Boyce-Codd normal form (BCNF)

Location

city	street	zip-code
------	--------	----------

Schema Location should be decomposed into the
 following schemas

CS

street	zip-code
--------	----------

CZ

city	zip-code
------	----------

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Boyce-Codd normal form (BCNF)

CS

street zip-code

Minimal key₂ = (zip-code, street)

Schema **CS** is in BCNF because does not exist a functional dependency whose left hand side is not a superkey

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Boyce-Codd normal form (BCNF)

CZ

city zip-code

zip-code → city

Minimal key = (zip-code)

Schema **CZ** is in BCNF because zip-code is a superkey

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Boyce-Codd normal form (BCNF)

CS

street zip-code

Minimal key = (zip-code, street)

CZ

city zip-code

zip-code → city

Minimal key = (zip-code)

Every relational schema, which consists of at most 2 attributes is always in BCNF

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Boyce-Codd normal form (BCNF)

CS

street zip-code

Minimal key = (zip-code, street)

CZ

city zip-code

zip-code → city

Minimal key = (zip-code)

Normalization to BCNF costs a functional dependency
city, street → zip-code

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Normalization of relational schemas

Given a relational schema $R = \{ A_1, \dots, A_n \}$

Identify all functional dependencies

Use functional dependencies to derive all minimal keys

Use functional dependencies and minimal keys to identify the highest normal form satisfied by R

Decompose R into the schemas in BCNF (3NF)

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Normalization of relational schemas

Shipment

s# city status p# quantity

s# → city

s# → status

city → status

s#, p# → quantity

s#, p# → city

s#, p# → status

Minimal key = (s#, p#)

Schema **Shipment** is in 1NF

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02 Database normalization

Normalization of relational schemas

Shipment

s# city status p# quantity

 $s\# \rightarrow city$ $s\# \rightarrow status$ $city \rightarrow status$ $s\#, p\# \rightarrow quantity$ $s\#, p\# \rightarrow city$ $s\#, p\# \rightarrow status$ Schema **Shipment** is not in 2NFNonprime attribute **city** depends on a part of a key (**s#, p#**)Minimal key = (**s#, p#**)

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Normalization of relational schemas

Shipment

s# city status p# quantity

Schema **Shipment** should be decomposed into the following schemas

SP

s# p# quantity

Supplier

s# city status

Schema **SP** is in BCNF $s\#, p\# \rightarrow quantity$ Schema **Supplier** is not in 3NF(& BCNF) $s\# \rightarrow city$ $city \rightarrow status$ $s\# \rightarrow status$

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Normalization of relational schemas

Supplier

s# city status

Schema **Supplier** should be decomposed into the following schemas

Supplier'

s# city

Schema **Supplier'** is in BCNF $s\# \rightarrow city$

Location

city status

Schema **Location** is in BCNF $city \rightarrow status$

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Normalization of relational schemas

Supplier

s# city status

Schema **Supplier** may be alternatively decomposed into the following schemas

Supplier'

s# city

Schema **Supplier'** is in BCNF $s\# \rightarrow city$

Supplier''

s# status

Schema **Supplier''** is in BCNF $s\# \rightarrow status$

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References

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