

Primal/Dual Decomposition Methods

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Outline of Lecture

- Subgradients
- Subgradient methods
- Primal decomposition
- Dual decomposition
- Summary

(Acknowledgement to Stephen Boyd for material for this lecture.)

Gradient and First-Order Approximation

- Recall basic inequality for convex differentiable f :

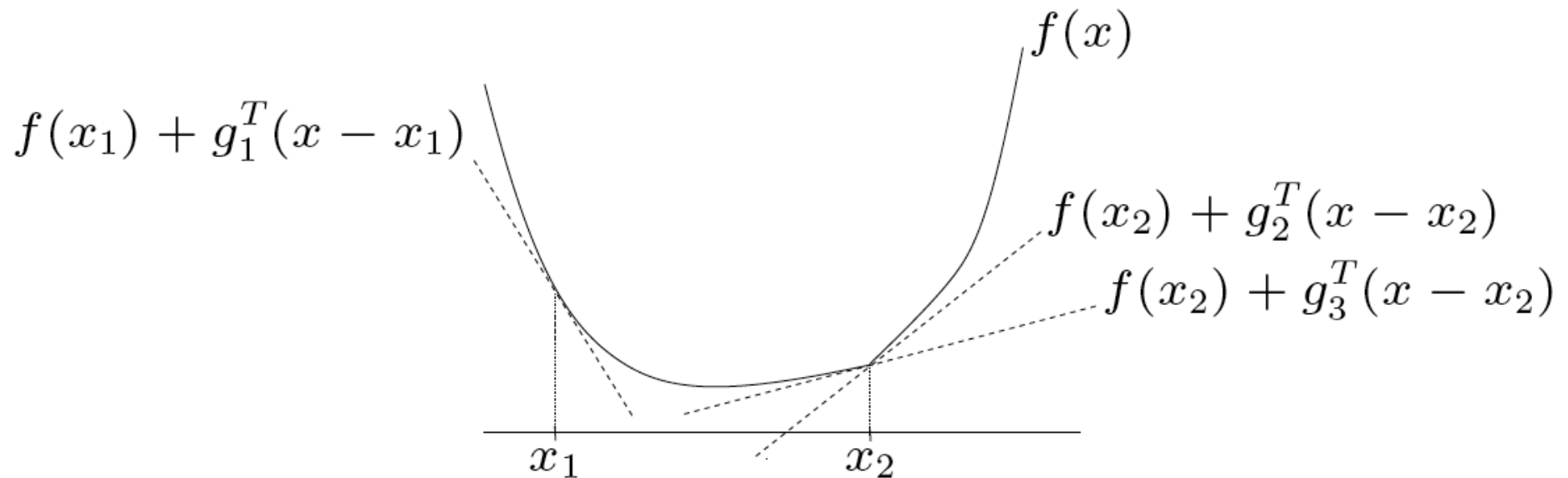
$$f(\mathbf{y}) \geq f(\mathbf{x}) + \nabla f(\mathbf{x})^T (\mathbf{y} - \mathbf{x}) \quad \forall \mathbf{y}$$

- First-order approximation of f at \mathbf{x} is a global underestimator, a supporting hyperplane.
- What if f is not differentiable?
- The answer is given by the concept of *subgradient*.

Subgradient of a Function

- g is a *subgradient* of f (not necessarily convex) at x if

$$f(y) \geq f(x) + g^T(y - x) \quad \forall y$$

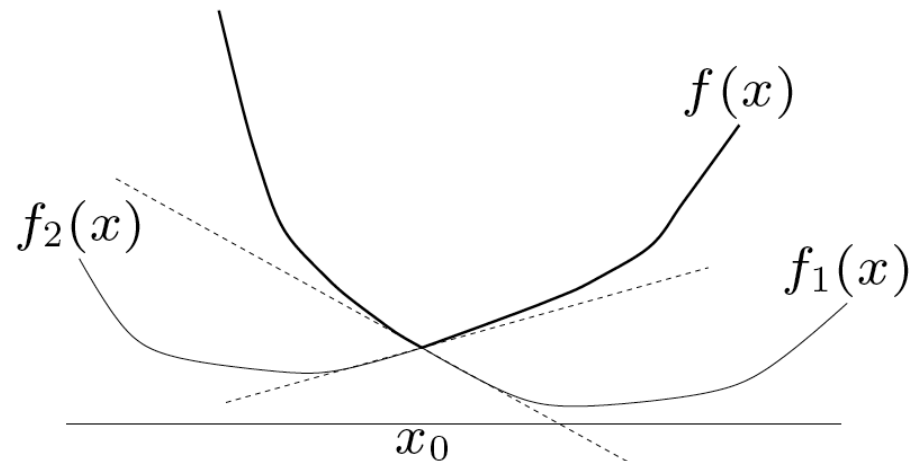


Subgradient of a Function

- \mathbf{g} is a subgradient if and only if $f(\mathbf{x}) + \mathbf{g}^T(\mathbf{y} - \mathbf{x})$ is a global affine underestimator of f .
- If f is convex and differentiable, then $\nabla f(\mathbf{x})$ is a subgradient of f at \mathbf{x} .
- Subgradients come up in several contexts:
 - algorithms for nondifferentiable convex optimization
 - convex analysis, e.g., optimality conditions and duality for nondifferentiable problems.

Example of Subgradient

- Consider $f = \max \{f_1, f_2\}$ with f_1 and f_2 convex and differentiable



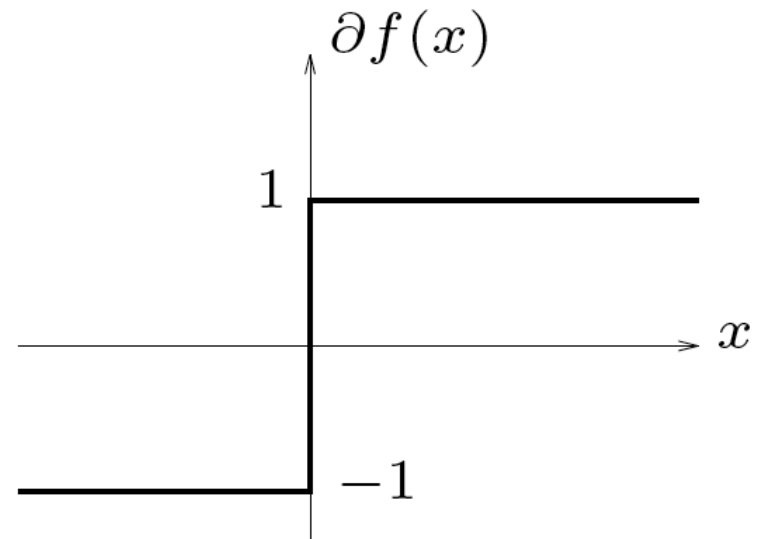
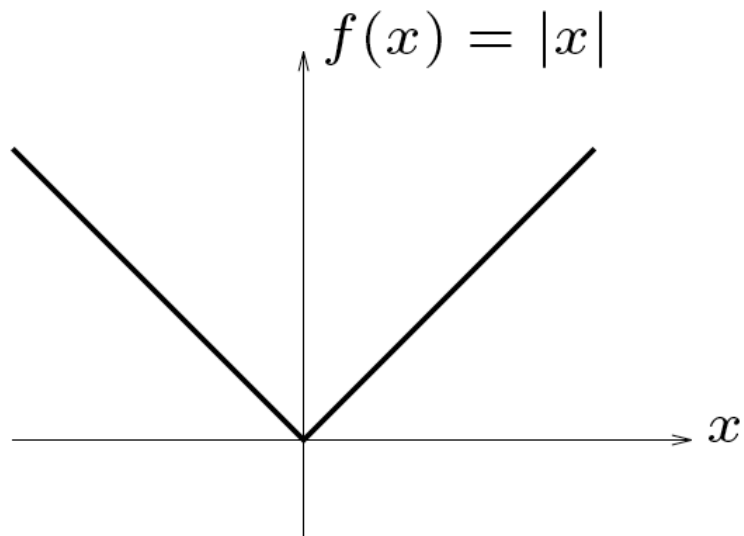
- Subgradient at point \mathbf{x} :
 - If $f_1(\mathbf{x}) > f_2(\mathbf{x})$, the subgradient is unique $\mathbf{g} = \nabla f_1(\mathbf{x})$.
 - If $f_2(\mathbf{x}) > f_1(\mathbf{x})$, the subgradient is unique $\mathbf{g} = \nabla f_2(\mathbf{x})$.
 - If $f_1(\mathbf{x}) = f_2(\mathbf{x})$, the subgradients form an interval $[\nabla f_1(\mathbf{x}), \nabla f_2(\mathbf{x})]$.

Subdifferential

- The set of all subgradients of f at \mathbf{x} is called the *subdifferential* of f at \mathbf{x} and is denoted $\partial f(\mathbf{x})$.
- $\partial f(\mathbf{x})$ is a closed convex set (can be empty).
- If f is convex:
 - $\partial f(\mathbf{x})$ is nonempty for $\mathbf{x} \in \text{relint dom } f$
 - if f is differentiable at \mathbf{x} , then $\partial f(\mathbf{x}) = \{\nabla f(\mathbf{x})\}$ (i.e., a singleton).
 - If $\partial f(\mathbf{x}) = \{\mathbf{g}\}$, then f is differentiable at \mathbf{x} and $\mathbf{g} = \nabla f(\mathbf{x})$.

Example of Subdifferential

- Consider $f(x) = |x|$:



Subgradient Calculus

- *Weak subgradient calculus*: formulas for finding one subgradient $\mathbf{g} \in \partial f(\mathbf{x})$.
- *Strong subgradient calculus*: formulas for finding the whole subdifferential $\partial f(\mathbf{x})$, i.e., all subgradients of f at \mathbf{x} .
- Many algorithms for nondifferentiable convex optimization require only one subgradient, so weak calculus suffices.
- Some algorithms and optimality conditions need the whole subdifferential.
- In practice, if you can compute $f(\mathbf{x})$, you can usually compute a subgradient $\mathbf{g} \in \partial f(\mathbf{x})$.

Some Basic Rules

(From now on, we will assume that f is convex and $\mathbf{x} \in \text{relint dom } f$.)

- $\partial f(\mathbf{x}) = \{\nabla f(\mathbf{x})\}$ if f is differentiable at \mathbf{x}
- Scaling: $\partial(\alpha f) = \alpha \partial f$ (assuming $\alpha > 0$)
- Addition: $\partial(f_1 + f_2) = \partial f_1 + \partial f_2$ (addition of sets)
- Affine transformation of variables: if $g(\mathbf{x}) = f(\mathbf{A}\mathbf{x} + \mathbf{b})$, then $\partial g(\mathbf{x}) = \mathbf{A}^T \partial f(\mathbf{A}\mathbf{x} + \mathbf{b})$
- Finite pointwise maximum: if $f = \max_i f_i$, then

$$\partial f(\mathbf{x}) = \text{Co} \bigcup \{ \partial f_i \mid f_i(\mathbf{x}) = f(\mathbf{x}) \},$$

i.e., convex hull of union of subdifferentials of active functions at \mathbf{x} .

Optimality Conditions: Unconstrained Case

- Recall that for f convex and differentiable, \mathbf{x}^* minimizes $f(\mathbf{x})$ if and only if:

$$\mathbf{0} = \nabla f(\mathbf{x}^*).$$

- The generalization to nondifferentiable convex f is

$$\mathbf{0} \in \partial f(\mathbf{x}^*).$$

Proof. By definition:

$$f(\mathbf{y}) \geq f(\mathbf{x}^*) + \mathbf{0}^T (\mathbf{y} - \mathbf{x}^*) \quad \forall \mathbf{y}.$$

Example: Piecewise Linear Maximization

- We want to minimize $f(\mathbf{x}) = \max_i \{ \mathbf{a}_i^T \mathbf{x} + b_i \}$.
- \mathbf{x}^* minimizes $f(\mathbf{x}) \iff \mathbf{0} \in \partial f(\mathbf{x}^*) = \text{Co} \{ \mathbf{a}_i \mid \mathbf{a}_i^T \mathbf{x}^* + b_i = f(\mathbf{x}^*) \} \iff$
there is a $\boldsymbol{\lambda}$ with

$$\boldsymbol{\lambda} \geq \mathbf{0}, \quad \mathbf{1}^T \boldsymbol{\lambda} = 1, \quad \sum_i \lambda_i \mathbf{a}_i = \mathbf{0}$$

where $\lambda_i = 0$ if $\mathbf{a}_i^T \mathbf{x}^* + b_i < f(\mathbf{x}^*)$.

- Interestingly, these are exactly the KKT conditions for the problem in epigraph form:

$$\begin{array}{ll} \underset{t, \mathbf{x}}{\text{minimize}} & t \\ \text{subject to} & \mathbf{a}_i^T \mathbf{x} + b_i \leq t, \quad i = 1, \dots, m. \end{array}$$

Optimality Conditions: Constrained Case

$$\begin{array}{ll}\underset{\mathbf{x}}{\text{minimize}} & f_0(\mathbf{x}) \\ \text{subject to} & f_i(\mathbf{x}) \leq 0, \quad i = 1, \dots, m\end{array}$$

where each f_i is convex, defined on \mathbb{R}^n (hence subdifferentiable) and strict feasibility holds (Slater's condition).

- The KKT necessary and sufficient conditions are

$$f_i(\mathbf{x}^*) \leq 0, \quad \lambda_i^* \geq 0$$

$$\mathbf{0} \in \partial f_0(\mathbf{x}^*) + \sum_{i=1}^m \lambda_i^* \partial f_i(\mathbf{x}^*)$$

$$\lambda_i^* f_i(\mathbf{x}^*) = 0 \quad (\text{complementary slackness}).$$

Numerical Methods for Nondifferentiable Problems

- In \mathbb{R} , we can always use the bisection method for nondifferentiable functions to reduce the interval (uncertainty) by half at each step.
- Can we generalize this to \mathbb{R}^n ? The problem is that \mathbb{R}^n is not ordered, as opposed to \mathbb{R} .
- The answer is: yes. These methods are called *localization methods*:
 - *cutting-plane methods* (date back to the 1960s in the Russian literature): the uncertainty set is a polyhedron
 - *ellipsoid method* (goes back to the 1970s in the Russian literature): the uncertainty set is an ellipsoid.
- Another convenient method is the *subgradient method*.

Subgradient Method

- *Subgradient method* is a simple algorithm (looks similar to a gradient method) to minimize a nondifferentiable convex function f :

$$\mathbf{x}^{(k+1)} = \mathbf{x}^{(k)} - \alpha_k \mathbf{g}^{(k)}$$

where

- $\mathbf{x}^{(k)}$ is the k th iterate
 - $\mathbf{g}^{(k)}$ is any subgradient of f at $\mathbf{x}^{(k)}$
 - $\alpha_k > 0$ is the k th stepsize
- Note that it is not a descent method (unlike a gradient method), so we need to keep track of the best point so far:

$$f_{\text{best}}^{(k)} = \min_{i=1, \dots, k} f(\mathbf{x}^{(i)}).$$

Stepsize Rules

- Different stepsize methods:
 - constant stepsize: $\alpha_k = \alpha$
 - constant step length: $\alpha_k = \gamma / \|\mathbf{g}^{(k)}\|$ (so $\|\mathbf{x}^{(k+1)} - \mathbf{x}^{(k)}\|_2 = \gamma$)
 - square summable but not summable: stepsizes satisfying

$$\sum_{k=1}^{\infty} \alpha_k^2 < \infty, \quad \sum_{k=1}^{\infty} \alpha_k = \infty$$

- nonsummable diminishing: stepsizes satisfying

$$\lim_{k \rightarrow \infty} \alpha_k = 0, \quad \sum_{k=1}^{\infty} \alpha_k = \infty.$$

Convergence Results

Under some technical conditions (boundedness of optimum value, boundedness of subgradients by G), the limiting value of the subgradient method $\bar{f} = \lim_{k \rightarrow \infty} f_{\text{best}}^{(k)}$ satisfies:

- constant stepsize: $\bar{f} - f^* \leq G^2\alpha/2$, i.e., converges to $G^2\alpha/2$ -suboptimal (converges to f^* if f differentiable and α small enough)
- constant step length: $\bar{f} - f^* \leq G\gamma/2$, i.e., converges to $G\gamma/2$ -suboptimal
- diminishing stepsize rule: $\bar{f} = f^*$, i.e., converges.

Example: Piecewise Linear Minimization

- Consider the following nondifferentiable optimization problem:

$$\text{minimize } f(\mathbf{x}) = \max_i \{ \mathbf{a}_i^T \mathbf{x} + b_i \}.$$

- To find a subgradient, simply choose an index j for which

$$\mathbf{a}_j^T \mathbf{x} + b_j = \max_i \{ \mathbf{a}_i^T \mathbf{x} + b_i \}$$

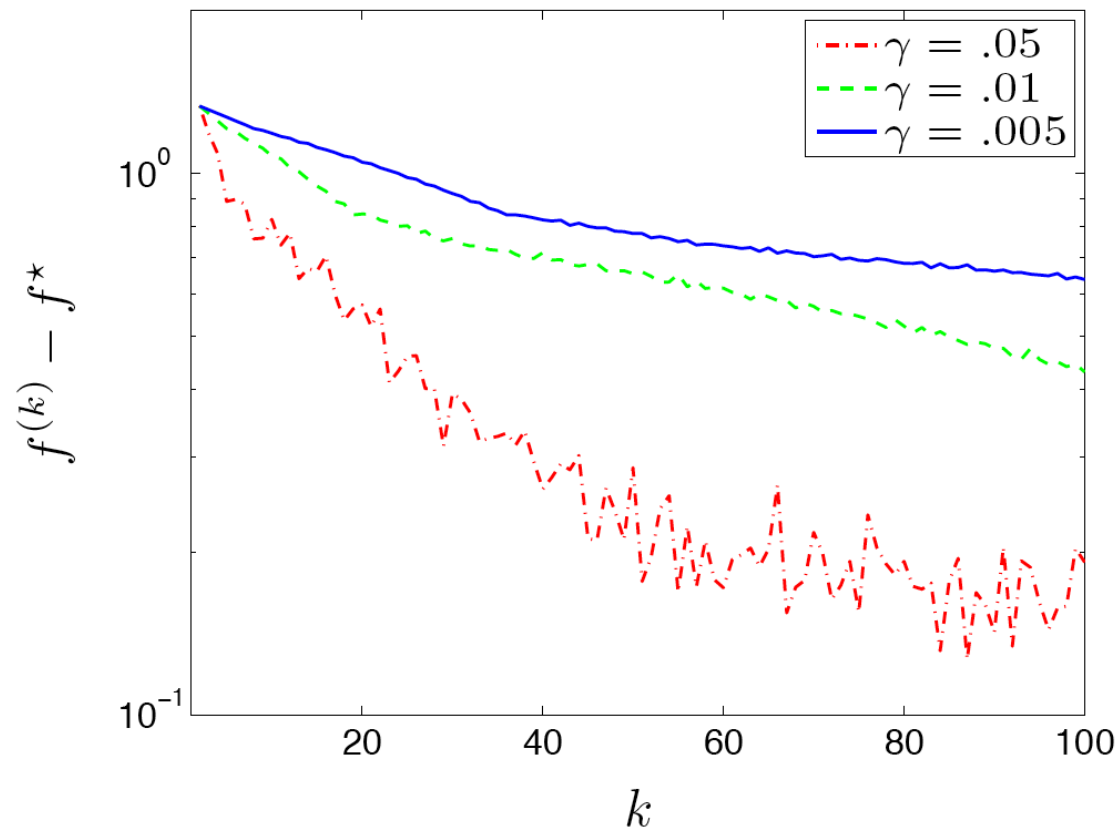
and take $\mathbf{g} = \mathbf{a}_j$.

- The subgradient update is

$$\mathbf{x}^{(k+1)} = \mathbf{x}^{(k)} - \alpha_k \mathbf{a}_j^{(k)}.$$

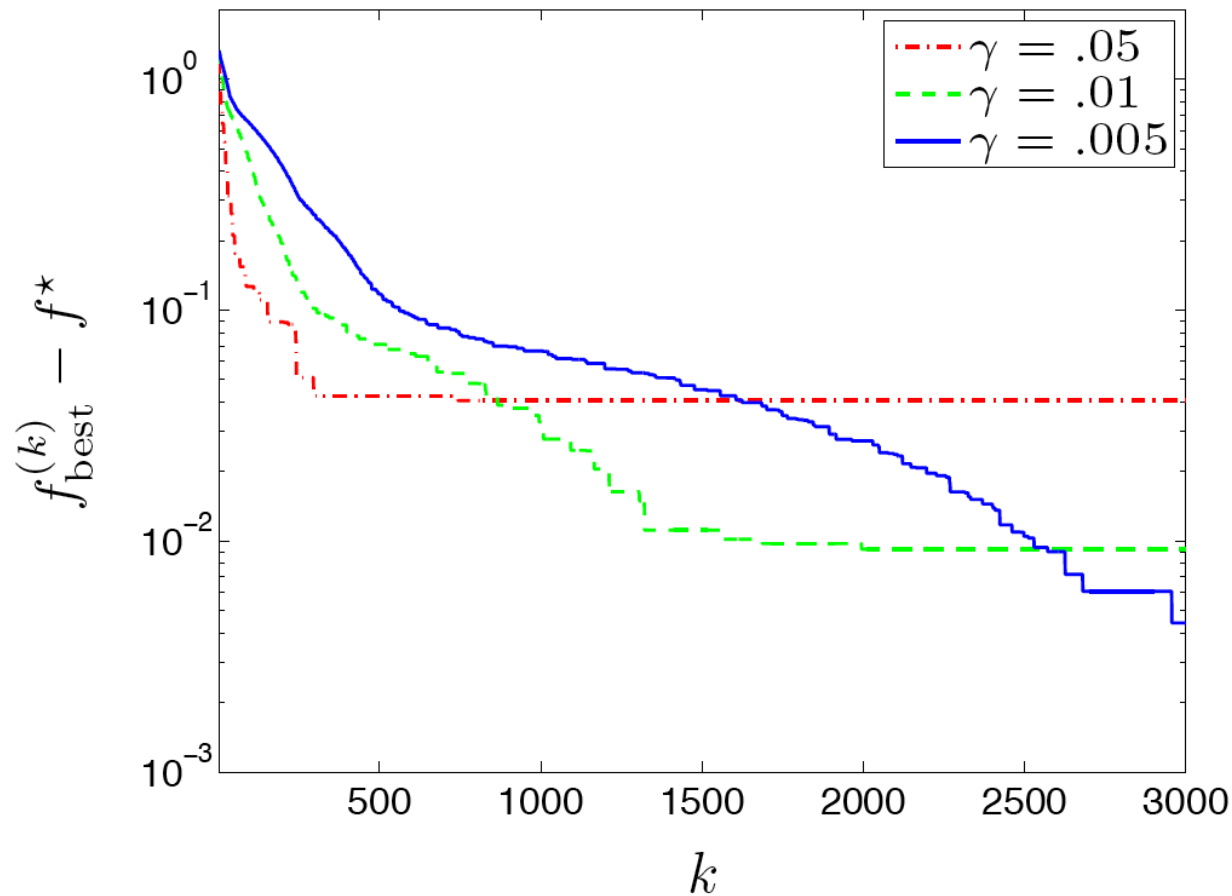
Example: Piecewise Linear Minimization (II)

- Problem instance with $n = 20$ variables, $m = 100$ terms, $f^* \approx 1.1$
- Constant step length, first 100 iterations, $f^{(k)} - f^*$:



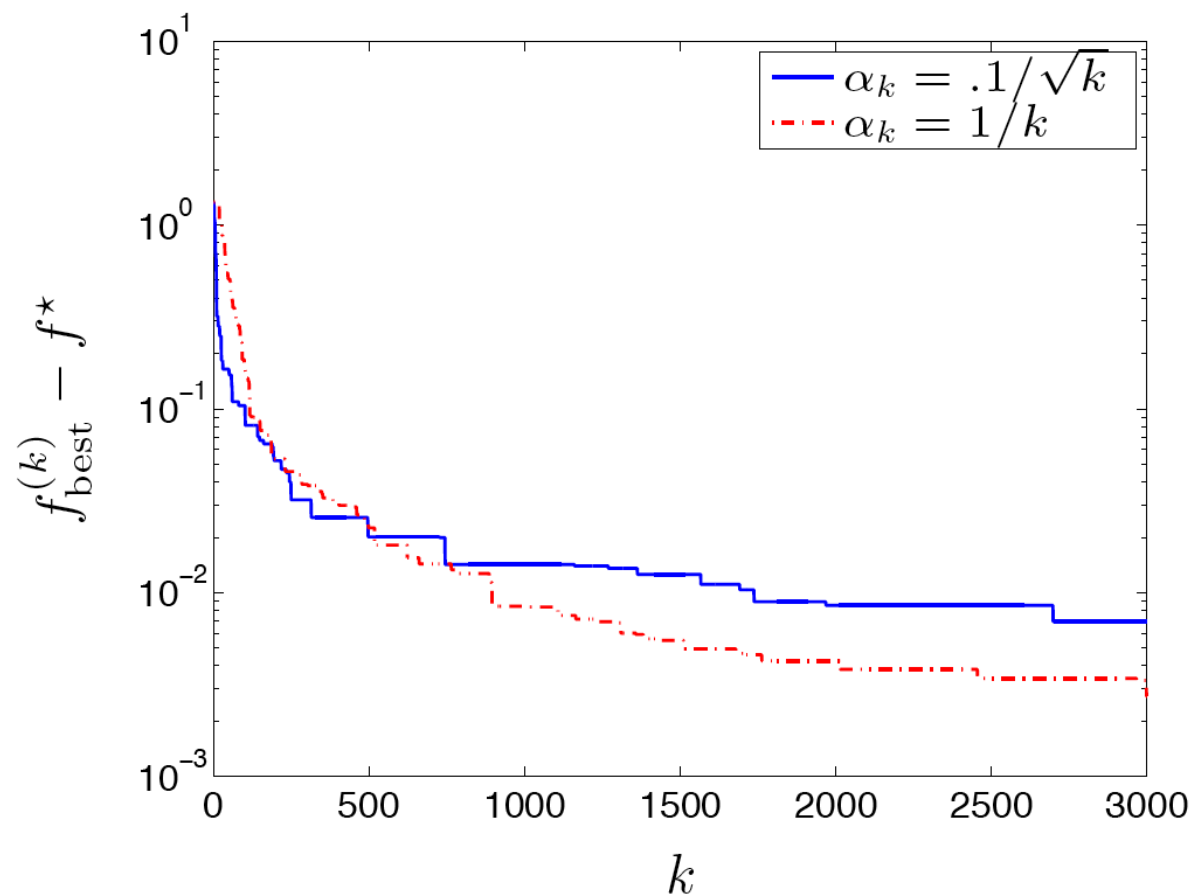
Example: Piecewise Linear Minimization (III)

- Constant step length, $f_{\text{best}}^{(k)} - f^*$:



Example: Piecewise Linear Minimization (IV)

- Diminishing stepsize rule ($\alpha_k = 0.1/\sqrt{k}$) and square summable stepsize rule ($\alpha_k = 1/k$):



Projected Subgradient Method for Constrained Optimization

- Consider the following constrained optimization problem:

$$\begin{array}{ll}\underset{\mathbf{x}}{\text{minimize}} & f(\mathbf{x}) \\ \text{subject to} & \mathbf{x} \in \mathcal{X}.\end{array}$$

- The *projected subgradient method* guarantees that each update produces a feasible point via a projection:

$$\mathbf{x}^{(k+1)} = \left[\mathbf{x}^{(k)} - \alpha_k \mathbf{g}^{(k)} \right]_{\mathcal{X}}$$

where $[\cdot]_{\mathcal{X}}$ denotes projection onto the convex set \mathcal{X} .

Decomposition Methods

- The idea of a decomposition method is to solve a problem by solving smaller subproblems coordinated by a master problem.
- Different reasons to use a decomposition method:
 - to allow the resolution of a problem otherwise unsolvable for memory reasons (useful in areas such as biology or image processing)
 - to speed up the resolution of the problem via parallel computation
 - to solve the problem in a distributed way (desirable for some wireless networks)
 - to derive nice, insightful, and efficient numerical algorithms as alternative to the use of general-purpose interior-point methods.

Decomposition Methods (II)

- In some (few) fortunate cases, problems decouple naturally:

$$\begin{array}{ll} \underset{\mathbf{x}}{\text{minimize}} & f_1(\mathbf{x}_1) + f_2(\mathbf{x}_2) \\ \text{subject to} & \mathbf{x}_1 \in \mathcal{X}_1, \quad \mathbf{x}_2 \in \mathcal{X}_2. \end{array}$$

- We can solve for \mathbf{x}_1 and \mathbf{x}_2 separately (in parallel).
- Even if they are solved sequentially, there is still the advantage of memory and of speed (if the computational effort is superlinear in problem size).

Decomposition Methods (III)

- In general, however, problems do not decouple so easily and that's when we can resort to decomposition methods.
- There are two main types of decomposition methods:
 - *primal decomposition*:
 - * deals with complicating variables that couple the subproblems
 - * the primal master problem controls directly the resources
 - *dual decomposition*:
 - * deals with complicating constraints that couple the subproblems
 - * the dual master problem controls the prices of the resources.

Primal Decomposition

- Consider the following problem with a **coupling variable**:

$$\begin{array}{ll} \underset{\mathbf{x}, \mathbf{y}}{\text{minimize}} & f_1(\mathbf{x}_1, \mathbf{y}) + f_2(\mathbf{x}_2, \mathbf{y}) \\ \text{subject to} & \mathbf{x}_1 \in \mathcal{X}_1, \quad \mathbf{x}_2 \in \mathcal{X}_2 \\ & \mathbf{y} \in \mathcal{Y}. \end{array}$$

- \mathbf{y} is the complicating variable or coupling variable.
- When \mathbf{y} is fixed the problem is separable in \mathbf{x}_1 and \mathbf{x}_2 and then decouples into two subproblems that can be solved independently.
- \mathbf{x}_1 and \mathbf{x}_2 are local or private variables; and \mathbf{y} can be interpreted as a global or public variable that serves as an interface or boundary variable between the two subproblems.

Primal Decomposition (II)

- For a given fixed \mathbf{y} we define the *subproblems*:

subproblem 1 : minimize $_{\mathbf{x}_1 \in \mathcal{X}_1}$ $f_1(\mathbf{x}_1, \mathbf{y})$

subproblem 2 : minimize $_{\mathbf{x}_2 \in \mathcal{X}_2}$ $f_2(\mathbf{x}_2, \mathbf{y})$

with optimal values $f_1^*(\mathbf{y})$ and $f_2^*(\mathbf{y})$.

- Since $\min_{\mathbf{x}, \mathbf{y}} f \equiv \min_{\mathbf{y}} \min_{\mathbf{x}} f$, it follows that the original problem is equivalent to the *master primal problem*:

$$\text{minimize}_{\mathbf{y} \in \mathcal{Y}} \quad f_1^*(\mathbf{y}) + f_2^*(\mathbf{y}).$$

Primal Decomposition (III)

Observations:

- subproblems can be solved independently for a given \mathbf{y}
- we don't have a closed-form expression for each function $f_i^*(\mathbf{y})$ and its corresponding gradient or subgradient (differentiable?)
- instead, to evaluate each function $f_i^*(\mathbf{y})$ at some point \mathbf{y} we need to solve an optimization problem
- interestingly, we can easily obtain a subgradient of $f_i^*(\mathbf{y})$ “for free” when we evaluate the function.

Primal Decomposition: Solving the Master Problem

- If the original problem is convex, so is master problem.
- To solve the master problem, we can use different methods such as
 - bisection (if y is scalar)
 - gradient or Newton method (if f_i^* differentiable)
 - subgradient, cutting-plane, or ellipsoid method.
- The subgradient method is very simple and allows itself to distributed implementation; however, its converge is slow in practice.
- Projected subgradient method:

$$\mathbf{y}(k+1) = [\mathbf{y}(k) - \alpha_k (\mathbf{s}_1(k) + \mathbf{s}_2(k))]_{\mathcal{Y}}$$

where $\mathbf{s}_i(k)$ is a subgradient of $f_i^*(\mathbf{y}(k))$.

Primal Decomposition Algorithm

repeat 1. Solve the subproblems:

Find $\mathbf{x}_1(k) \in \mathcal{X}_1$ that minimizes $f_1(\mathbf{x}_1(k), \mathbf{y}(k))$, and
a subgradient $\mathbf{s}_1(k) \in \partial f_1^*(\mathbf{y}(k))$.

Find $\mathbf{x}_2(k) \in \mathcal{X}_2$ that minimizes $f_2(\mathbf{x}_2(k), \mathbf{y}(k))$, and
a subgradient $\mathbf{s}_2(k) \in \partial f_2^*(\mathbf{y}(k))$.

2. Update the complicating variable to minimize the primal master subproblem:

$$\mathbf{y}(k+1) = [\mathbf{y}(k) - \alpha_k (\mathbf{s}_1(k) + \mathbf{s}_2(k))]_{\mathbf{y}}.$$

3. $k = k + 1$

Primal Decomposition: Subgradients

Lemma: Let $f^*(\mathbf{y})$ be the optimal value of the convex problem

$$\begin{array}{ll}\underset{\mathbf{x}}{\text{minimize}} & f_0(\mathbf{x}) \\ \text{subject to} & h_i(\mathbf{x}) \leq y_i, \quad i = 1, \dots, m.\end{array}$$

A subgradient of $f^*(\mathbf{y})$ is $-\boldsymbol{\lambda}^*(\mathbf{y})$.

Proof. The Lagrangian is $L(\mathbf{x}, \boldsymbol{\lambda}) = f_0(\mathbf{x}) + \boldsymbol{\lambda}^T(\mathbf{h}(\mathbf{x}) - \mathbf{y})$. Then,

$$\begin{aligned} f^*(\mathbf{y}_0) &= f_0(\mathbf{x}^*(\mathbf{y}_0)) \\ &= g(\boldsymbol{\lambda}^*(\mathbf{y}_0)) \\ &\leq L(\mathbf{x}, \boldsymbol{\lambda}^*(\mathbf{y}_0)) = f_0(\mathbf{x}) + \boldsymbol{\lambda}^{*T}(\mathbf{h}(\mathbf{x}) - \mathbf{y}_0) \\ &= f_0(\mathbf{x}) + \boldsymbol{\lambda}^{*T}(\mathbf{h}(\mathbf{x}) - \mathbf{y}) + \boldsymbol{\lambda}^{*T}(\mathbf{y} - \mathbf{y}_0) \\ &\leq f_0(\mathbf{x}) + \boldsymbol{\lambda}^{*T}(\mathbf{y} - \mathbf{y}_0) \quad \forall \mathbf{y} \end{aligned}$$

where the last inequality holds for any \mathbf{x} such that $\mathbf{h}(\mathbf{x}) \leq \mathbf{y}$. In particular,

$$\begin{aligned} f^*(\mathbf{y}_0) &\leq \min_{\mathbf{h}(\mathbf{x}) \leq \mathbf{y}} f_0(\mathbf{x}) + \boldsymbol{\lambda}^{*T}(\mathbf{y} - \mathbf{y}_0) \\ &= f^*(\mathbf{y}) + \boldsymbol{\lambda}^{*T}(\mathbf{y} - \mathbf{y}_0) \end{aligned}$$

or

$$f^*(\mathbf{y}) \geq f^*(\mathbf{y}_0) - \boldsymbol{\lambda}^{*T}(\mathbf{y} - \mathbf{y}_0).$$

□

Lemma: Let $f^*(\mathbf{y})$ be the optimal value of the convex problem

$$\begin{aligned} &\underset{\mathbf{x}}{\text{minimize}} && f_0(\mathbf{x}, \mathbf{y}) \\ &\text{subject to} && h_i(\mathbf{x}) \leq y_i, \quad i = 1, \dots, m. \end{aligned}$$

A subgradient of $f^*(\mathbf{y})$ is $(\mathbf{s}_0(\mathbf{x}^*(\mathbf{y}), \mathbf{y}) - \boldsymbol{\lambda}^*(\mathbf{y}))$.

Dual Decomposition

- Consider the following problem with a **coupling constraint**:

$$\begin{array}{ll}\underset{\mathbf{x}}{\text{minimize}} & f_1(\mathbf{x}_1) + f_2(\mathbf{x}_2) \\ \text{subject to} & \mathbf{x}_1 \in \mathcal{X}_1, \quad \mathbf{x}_2 \in \mathcal{X}_2 \\ & \mathbf{h}_1(\mathbf{x}_1) + \mathbf{h}_2(\mathbf{x}_2) \leq \mathbf{h}_0\end{array}$$

- $\mathbf{h}_1(\mathbf{x}_1) + \mathbf{h}_2(\mathbf{x}_2) \leq \mathbf{h}_0$ is the complicating or coupling constraint.
- If the coupling constraint is relaxed with a Lagrange multiplier λ , then the problem decouples into two subproblems that can be solved independently.
- \mathbf{x}_1 and \mathbf{x}_2 are local or private variables; and λ can be interpreted as a global or public price variable that serves as an interface or boundary variable between the two subproblems.

Dual Decomposition (II)

- The partial Lagrangian after relaxing the coupling constraint is:

$$L(\mathbf{x}, \boldsymbol{\lambda}) = f_1(\mathbf{x}_1) + f_2(\mathbf{x}_2) + \boldsymbol{\lambda}^T (\mathbf{h}_1(\mathbf{x}_1) + \mathbf{h}_2(\mathbf{x}_2) - \mathbf{h}_0).$$

- The dual function is

$$\begin{aligned} g(\boldsymbol{\lambda}) &= \inf_{\mathbf{x} \in \mathcal{X}} L(\mathbf{x}, \boldsymbol{\lambda}) \\ &= \inf_{\mathbf{x}_1 \in \mathcal{X}_1} \left\{ f_1(\mathbf{x}_1) + \boldsymbol{\lambda}^T \mathbf{h}_1(\mathbf{x}_1) \right\} \\ &\quad + \inf_{\mathbf{x}_2 \in \mathcal{X}_2} \left\{ f_2(\mathbf{x}_2) + \boldsymbol{\lambda}^T \mathbf{h}_2(\mathbf{x}_2) \right\} - \boldsymbol{\lambda}^T \mathbf{h}_0 \end{aligned}$$

which clearly decouples.

Dual Decomposition (III)

- For a given fixed λ we define the *subproblems*:

$$\text{subproblem 1 :} \quad \text{minimize}_{\mathbf{x}_1 \in \mathcal{X}_1} \quad f_1(\mathbf{x}_1) + \lambda^T \mathbf{h}_1(\mathbf{x}_1)$$

$$\text{subproblem 2 :} \quad \text{minimize}_{\mathbf{x}_2 \in \mathcal{X}_2} \quad f_2(\mathbf{x}_2) + \lambda^T \mathbf{h}_2(\mathbf{x}_2)$$

with optimal values $g_1(\lambda)$ and $g_2(\lambda)$.

- From strong duality, the original problem is equivalent to the *master dual problem*:

$$\text{maximize}_{\lambda \geq 0} \quad g_1(\lambda) + g_2(\lambda) - \lambda^T \mathbf{h}_0.$$

Dual Decomposition (IV)

Observations:

- subproblems can be solved independently for a given λ
- we don't have a closed-form expression for each function $g_i(\lambda)$ and its corresponding gradient or subgradient (differentiable?)
- instead, to evaluate each function $g_i(\lambda)$ at some point λ we need to solve an optimization problem
- as in the primal case, we can easily obtain a subgradient of $g_i(\lambda)$ “for free” when we evaluate the function.

Dual Decomposition: Solving the Master Problem

- The dual master problem is always convex regardless of the original problem. However, we still need convexity to have strong duality (under some constraint qualifications like Slater's condition).
- To solve the master problem, we can use different methods such as
 - bisection (if λ is scalar)
 - gradient or Newton method (if g_i differentiable)
 - subgradient, cutting-plane, or ellipsoid method.
- Projected subgradient method:

$$\boldsymbol{\lambda}(k+1) = [\boldsymbol{\lambda}(k) + \alpha_k (\mathbf{s}_1(k) + \mathbf{s}_2(k) - \mathbf{h}_0)]^+$$

where $\mathbf{s}_i(k)$ is a subgradient of $g_i(\boldsymbol{\lambda}(k))$.

Dual Decomposition Algorithm

repeat 1. Solve the subproblems:

Find $\mathbf{x}_1(k) \in \mathcal{X}_1$ that minimizes $f_1(\mathbf{x}_1(k)) + \boldsymbol{\lambda}(k)^T \mathbf{h}_1(\mathbf{x}_1(k))$,
and a subgradient $\mathbf{s}_1(k) \in \partial g_1(\boldsymbol{\lambda}(k))$.

Find $\mathbf{x}_2(k) \in \mathcal{X}_2$ that minimizes $f_2(\mathbf{x}_2(k)) + \boldsymbol{\lambda}(k)^T \mathbf{h}_2(\mathbf{x}_2(k))$,
and a subgradient $\mathbf{s}_2(k) \in \partial g_2(\boldsymbol{\lambda}(k))$.

2. Update the complicating variable to minimize the primal master subproblem:

$$\boldsymbol{\lambda}(k+1) = [\boldsymbol{\lambda}(k) + \alpha_k (\mathbf{s}_1(k) + \mathbf{s}_2(k) - \mathbf{h}_0)]^+.$$

3. $k = k + 1$

Dual Decomposition: Subgradients

Lemma: Let $g(\boldsymbol{\lambda})$ be the dual function corresponding to the problem

$$\begin{array}{ll}\underset{\mathbf{x}}{\text{minimize}} & f_0(\mathbf{x}) \\ \text{subject to} & h_i(\mathbf{x}) \leq 0, \quad i = 1, \dots, m.\end{array}$$

A subgradient of $g(\boldsymbol{\lambda})$ is $\mathbf{h}(\mathbf{x}^*(\boldsymbol{\lambda}))$.

Proof. The Lagrangian is $L(\mathbf{x}, \boldsymbol{\lambda}) = f_0(\mathbf{x}) + \boldsymbol{\lambda}^T \mathbf{h}(\mathbf{x})$ and the dual function $g(\boldsymbol{\lambda}) = \inf_{\mathbf{x}} L(\mathbf{x}, \boldsymbol{\lambda})$. Then,

$$\begin{aligned}g(\boldsymbol{\lambda}) &= \inf_{\mathbf{x}} f_0(\mathbf{x}) + \boldsymbol{\lambda}^T \mathbf{h}(\mathbf{x}) \\ &\leq f_0(\mathbf{x}^*(\boldsymbol{\lambda}_0)) + \boldsymbol{\lambda}^T \mathbf{h}(\mathbf{x}^*(\boldsymbol{\lambda}_0)) \\ &= f_0(\mathbf{x}^*(\boldsymbol{\lambda}_0)) + \boldsymbol{\lambda}_0^T \mathbf{h}(\mathbf{x}^*(\boldsymbol{\lambda}_0)) + (\boldsymbol{\lambda} - \boldsymbol{\lambda}_0)^T \mathbf{h}(\mathbf{x}^*(\boldsymbol{\lambda}_0)) \\ &= g(\boldsymbol{\lambda}_0) + (\boldsymbol{\lambda} - \boldsymbol{\lambda}_0)^T \mathbf{h}(\mathbf{x}^*(\boldsymbol{\lambda}_0)).\end{aligned}$$

□

Summary

- We have described the concept of subgradient as a generalization of gradient.
- We have considered subgradient methods as formally similar to gradient methods.
- Finally, we have derived the two basic decomposition techniques:
 - primal decomposition
 - dual decomposition.