

Shri Ramdababu College of
Engineering and Management.

Programme Name
:- Computer Science
and Engineering

Examination
:- Winter 2020
ESE

Course Code :- CST 252

Roll no :- CSU20W-3006

Course Name :-
Data Structures and
Algorithm

Date :- 01/12/2020

Time :- 10 am to
(2 pm

Q.1) a) Consider Structure array.

Struct Array

```
d int *a;  
int length;  
int size;
```

3 Write a C program to delete all
duplicate entries.

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Q. b) Consider an array stores 30 students in 4 Subj. Row represents subjects and column represents students. base address = 2500. Lower bound = 0.

(a) Compute address of marker [3][5] in row major.

(b) Compute address of marker [2][20] in column major.

P. 2) Q. 2) (a) Consider a Stack ADT is already created. write a function to check whether parenthesis are balanced or not.

(b) Design a C program for storing two queues stored in one array. Queue 1 grows from left to right and Queue 2 from right to left.

Q.3) (a) Write a C program to delete all negative numbers from linked list.

(b) Write an algo. to interchange the k^{th} and $(k+1)^{th}$ node of a doubly linked list. Ensure that nodes are swapped and not just values.

Q.4) (a) A company maintains an array of employees which stores employee name, id and salary. Array is always kept in sorted order according to employee id. Design a C function which will search and display the details of a particular employee when ID is specified.

(b) Consider a hash table of size = 11. Use double hashing, insert keys 27, 72, 63, 42, 70, 36, 38 and 10 into table. $h_1 = k \bmod 10$ and $h_2 = k \bmod 8$.

Q.5) (a) Construct a B+ tree of order 5.

80, 40, 15, 25, 30, 90, 35, 50
60, 70

(b)

Write a function to check whether root node is balanced or not.

Q.6) (a)

Consider four cities - (1) New Delhi
(2) Mumbai (3) Chennai
(4) Bangalore. Use dikstra's algo to find shortest path to each node from New Delhi.

Flight no.

101

102

103

104

105

106

107

108

Origin

4

1

2

3

4

Destination

5

4

2

6

4

6

2

3

4

3

1

1

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- (b) Explain following terms:-
 (a) Complete graph.
 (b) Strongly connected graph.

Answers

(a) Word size = 2 \rightarrow taken
 (b) address of [3][5] in row major.

$$= \alpha + [i \times \cancel{N} + j \cancel{\times}] \times \frac{\text{bpw}}{\text{bpw}}$$

$$= 2500 + [3 \times 30 + 5] \times 2$$

$$= 2500 + [90 + 5] \times 2$$

$$= 2500 + 190$$

$$= 2690$$

(b) address of [2][20] in column major

$$\approx \alpha + [i + j \times M] \times \text{bpw}$$

$$\begin{aligned}
 &= 2500 + [2 + 20 \times 4] \times 2 \\
 &= 2500 + (2 + 80) \times 2 \\
 &= 2664
 \end{aligned}$$

Q.2) ①

~~int checkBalance (Stack sk, char a[])~~

②

```

int i = 0;
while (i < len(a)) {
    if (a[i] == ')') {
        if (sk isEmpty(sk))
            return 0;
    }
}

```

} else if (peek(sk) == '(')

~~pop (sk);~~

}

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else if (peek (sk) == ')')
 return 0;

}

} else {

 push (sk);

 i++;

}

if () isEmpty (sk)) {

 return 0;

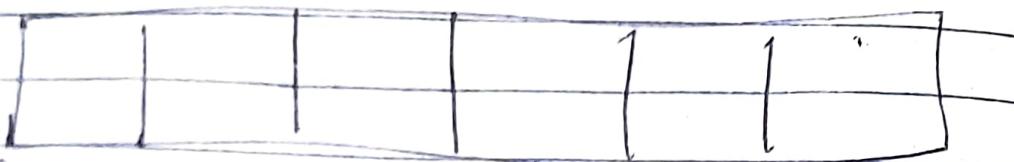
}

return 1;

// isEmpty is a function which checks if
// stack is empty or not
// peek checks the top element
// pop and push removes and add respectively.

Q. 2 b)

Queue 1 \Rightarrow



front 1 \rightarrow

rear 1 \rightarrow

~~tops~~

\leftarrow front 2
 \leftarrow rear 2

~~if~~ To insert in Queue 2

\Rightarrow rear 1 ++

To remove in Queue 1

\Rightarrow front 1 ++

To insert in Queue 2

\Rightarrow ~~rear~~ --

To remove in Queue 1

\Rightarrow front --

Underflow condition is

Queue 1 or if $\text{front} == -1 \& \&$
 $\text{rear} == -1$

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* int ~~return~~

void Insert (int opt, int x, int q[], int front)

int front2, int rear1, int rear2);

if (opt == -1) {

Case of Queue 1

if (front1 == 0 & rear1 == len(q))
~~& front1 > rear1)~~

{ cout ("Overflow\n");

} else {

q[rear1] = x;

3

3

else if ($opt == 2$) {

// Case of Queue 2

if ($front2 == n - 1 \& rear2 == 0$) {

printf ("Overflow\n");

} else {

$q[-rear2] = x;$

}

}

void

remove (

int opt, int q[], int front1,

int front2, int rear1, int rear2);

if ($opt == 1$) {

if ($front1 > rear1$) {

printf ("Underflow\n");

} else

int ans = q[front1];

}

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3. else if ($opt == 2$) {
 if ($front[2] < rear[2]$) {
 printf ("Underflow\n");
 } else {
 int ans = q[-front[2]];
 }
}

3
~~return~~ ans;
 printf ("%d", ans);

Q. 3) (a)

→ // Assuming linked list created.

Void remove (struct node** head) {

Struct ~~Node~~ * ttemp = *head;
temp != NULL && (ttemp -> data < 0) {

ttemp = ttemp -> next;
} // checking till the first positive integer.

~~struct~~ Node *tp = *temp;

while (!tp) = NULL) {

* if (tp -> data > 0) {

Create newnode (tp -> data);

temp -> next = tp -> newnode;
temp = temp -> next;

3
tp = tp -> next;

3

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return temp;

Q. 3) (b)

void delete_kth (struct Node** head, int k)

~~while (- - k > 0)~~

~~while (k != 1)~~

if (head == NULL || head->next
= = NULL)

return head;

3
~~else {~~

struct Node* node = *head;

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while ($-k \neq 1$) {

 node = node \rightarrow next;

}

 node \rightarrow next = node \rightarrow next \rightarrow next \rightarrow
 next;

 return head;

}

Q. 4) (b)

\rightarrow size = m = 11

$$h_1(k) = k \bmod 10$$

$$h_2(k) = k \bmod 8$$

$$h(k, i) = [h_1(k) + i * h_2(k)] \bmod m$$

0 1 2 3 4 5 6 7 8 9 10

--	--	--	--	--	--	--	--	--	--	--

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$$\textcircled{1} \quad h(27, 0) = [27 \% 10 + 0 \times 27 \% 8]$$

$$= 7 \% 11 \\ = 7$$

$$\textcircled{2} \quad h(72, 0)$$

$$= [72 \% 10 + 0 \times 72 \% 8] \% 11$$

$$= [2 + 0] \% 11 \\ = 2$$

Hash table is

0	1	2	3	4	5	6	7	8	9	10
-1	-1	72	-1	-1	-1	-1	27	-1	-1	7

$$\textcircled{3} \quad h(63, 0) = [63 \% 10 + 0 \times 63 \% 8] \% 11$$

0	1	2	3	4	5	6	7	8	9	10
-1	72	63	-1	-1	-1	27	-1	-1	-1	7

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④ $n(42, 0)$

$$= [42 \% \cdot 10 + 0 \times 42 \% \cdot 8] \% \cdot 11$$

= 2

0	1	2	3	4	5	6	7	8	9	10
-1	-1	72	63	-1	-1	-1	27	-1	-1	-1

As we can see at index $\times 2$, 72 is present

$n(42, 1)$

$$= [42 \% \cdot 10 + 1 \times 42 \% \cdot 8] \% \cdot 11$$

$$= 8 [2 + 2] \% \cdot 11$$

$$= 9$$

0	1	2	3	4	5	6	7	8	9	10
-1	-1	72	63	42	-1	-1	27	-1	-1	-1

⑤ $n(44, 0)$

$$= [44 \% \cdot 10 + 0 \times 44 \% \cdot 8] \% \cdot 11$$

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$$= 4$$

We can see at index 4, 42 is present.

$$\therefore h(44, 1)$$

$$= [44\%10 + 1 \times 44 \% 8] \% 11$$

$$= [4 + 4] \% 11$$

$$= 8$$

0	1	2	3	4	5	6	7	8	9	10
-1	1	72	63	42	-1	-1	27	45	-1	-1

⑥ 36

$$\Rightarrow h(36, 0)$$

$$= (36 \% 10 + 0 \times 36 \% 8) \% 11$$

$$= 6$$

0	1	2	3	4	5	6	7	8	9	10
-1	-1	72	63	42	-1	36	27	44	-1	-1

$$\text{① } h(38, 0)$$

$$= (38 \text{ } \% 10 + 0 \times 38 \text{ } \% 8) \% 11$$

$$= 8$$

We can see at index 8, 44 is present.

$$h(38, 1)$$

$$= (38 \text{ } \% 10 + 1 \times 38 \text{ } \% 8) \% 11$$

$$= (8 + 4) \% 11$$

$$= 1$$

0	1	2	3	4	5	6	7	8	9	10
-1	38	72	63	42	-1	36	27	44	-1	-1

~~At~~ index 3, 63 is present.

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$$h(\cancel{38}, 4)$$

$$= (38\% \times 10 + 2 \times 38\% \cdot 8) \% \\ = \cancel{8}$$

⑧ $h(101, 0)$

$$= (101\% \times 10 + 0 \times 101\% \cdot 8) \% \\ = (1) \% \\ = 1$$

A + 1, 38 is present.

$$h(10, 1)$$

$$= (10\% \times 10 + 1 \times 10\% \cdot 8) \% \\ = (1 + 5) \% \\ = 6\%$$

A + 6, 36 is present.

$$h(101, 2)$$

$$= (101 \cdot 1/10 + 2 \times 101^0 / 08) \% 11$$

$$= (1 + 10) \% 11$$

$$= 0$$

0	1	2	3	4	5	6	7	8	9	10
101	38	72	63	42	-1	36	27	43	-1	-1

Hash table ↴

Q.5) (a)

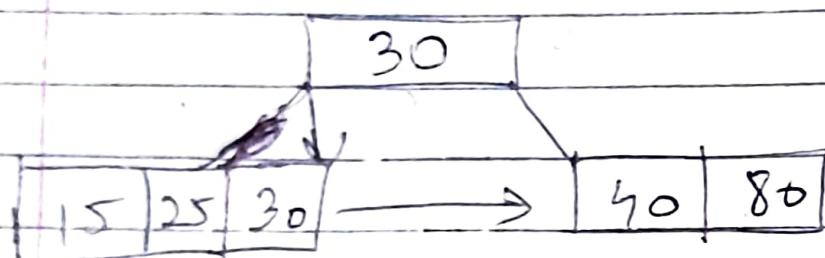
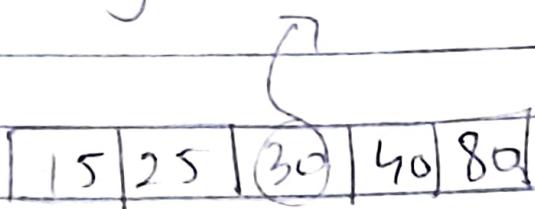
→ Order = m = 5

∴ Max no of keys = m-1 = 4

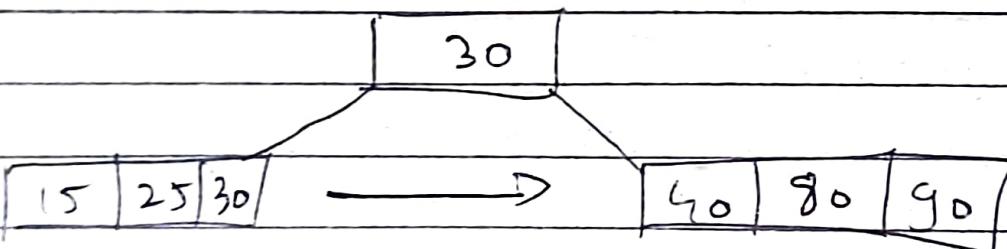
(i) Inserting 80, 40, 15, 25

15	25	40	80
----	----	----	----

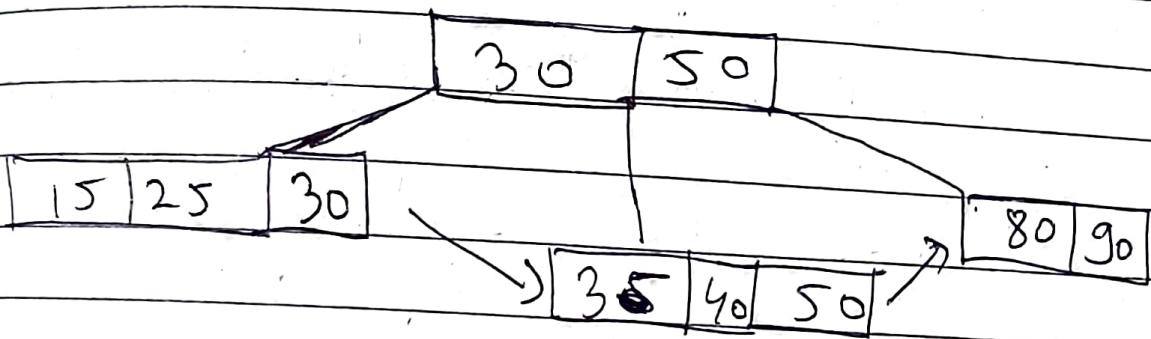
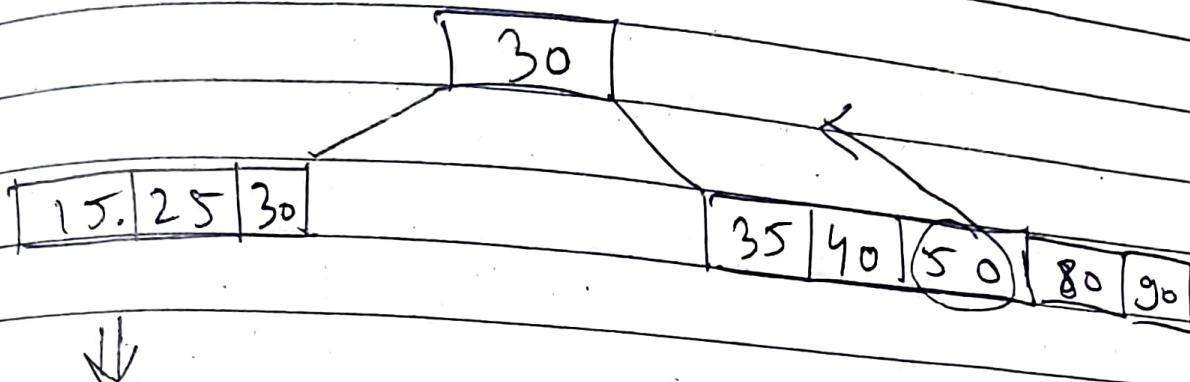
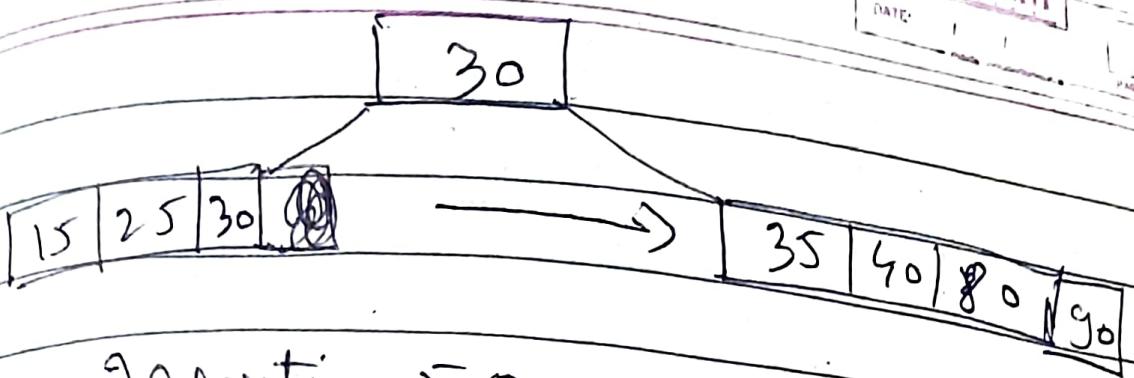
② Inserting 30



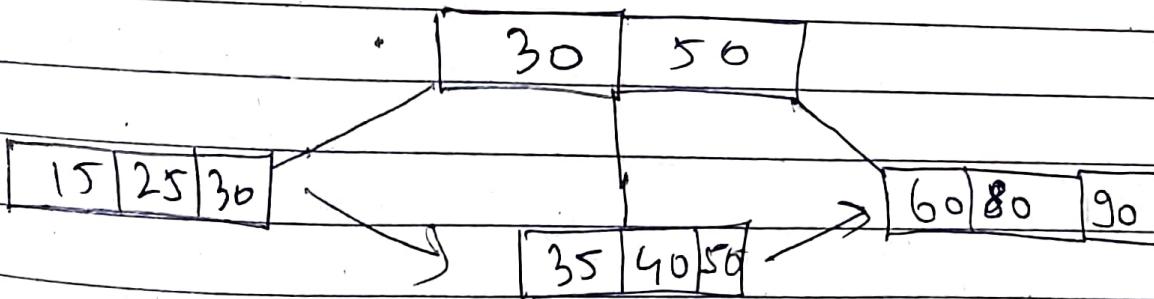
③ Inserting 90,



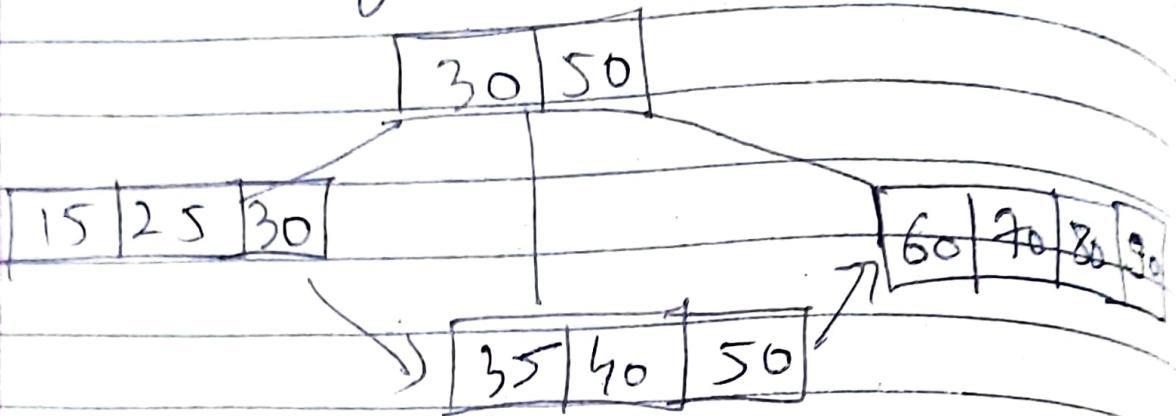
④ Inserting 35,



6 Inserting 60,



Q.6) Inserting 70,



Q.5) (b)



int height (treenode *root) {

if (root == NULL) {

return 0;

}

int l = height (root->left);

int r = height (root->right);

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Kelihongz

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return $1 + \max(l, r)$;

int checkbalance (treeNode *root) {

if (root == NULL) {

return 1;

}

int l = height (~~as~~ root \rightarrow left);

int r = height (root \rightarrow right);

int ans = l - r;

if (ans == 1 || ans == -1 || ans == 0) {

return 1;

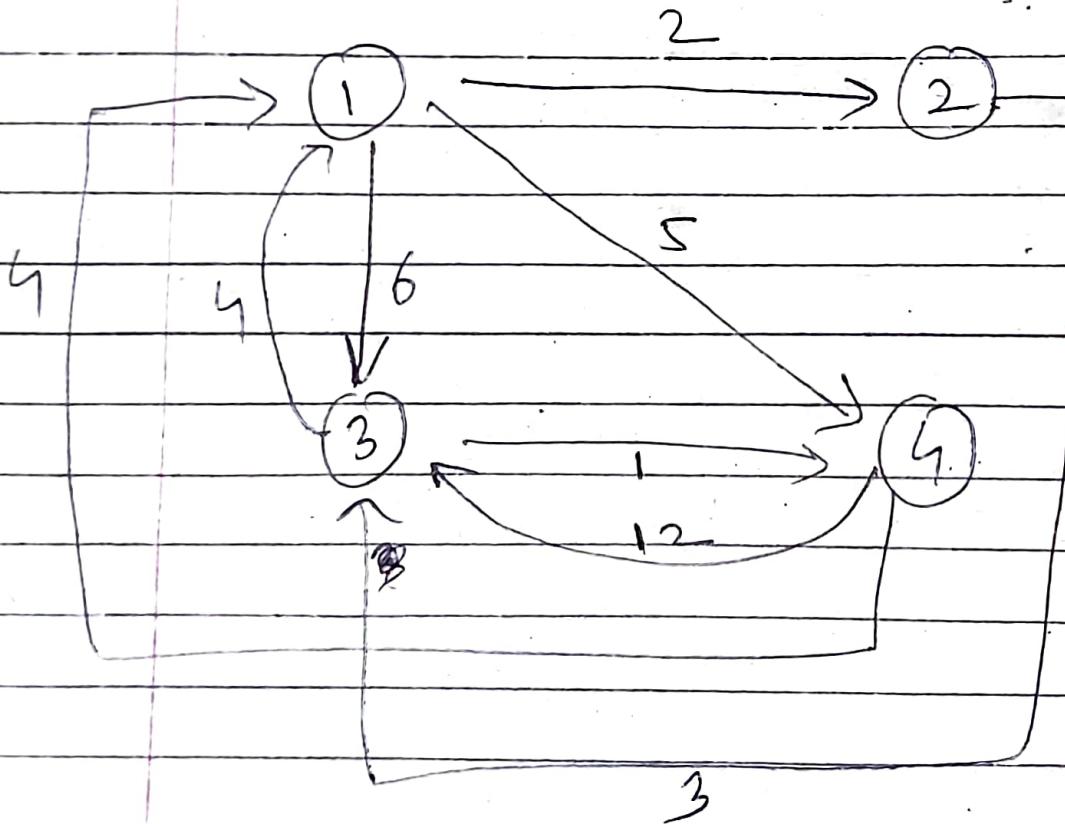
}

return 0;

Q. 6) (a)

Graph using the table will be

- (1) for New Delhi
- (2) for Mumbai
- (3) for Chennai
- (4) for Bangalore



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Dijkstra's Algo from New Delhi to each node:

	(1)	(2)	(3)	(4)
Vertices	-	-	-	-
Prev. vertex	-	-	-	-
weight	0	-	-	-

Neighboring nodes of (1) are.

(2), (3), and (4).

	(1)	(2)	(3)	(4)
Vertices	-	-	-	-
Prev. vertex	-	(1)	(3)	(3)
weight	0	2	5	5

	(1)	(2)	(3)	(4)
Vertices	-	-	-	-
weight	0	2	6	5
Prev. vertex	-	(1)	(1)	(1)

But from (1) to (3). weight is 6 while min (6, path 123)
 $= \min(6, 5) = 5$

Vertices	(1)	(2)	(3)	(4)
weight	0	2	5	
prev. vertex	-	(1)	(2)	

Minimum from New Delhi to

Mumbai is 2

Chennai is 5

Bangalore is 5

Q. (6) b

① Complete graph: A graph is said to be completely connected if every node is connected to every node.

A ~~graph~~ complete graph with 'v' vertices has $\frac{v * (v-1)}{2}$ edges.

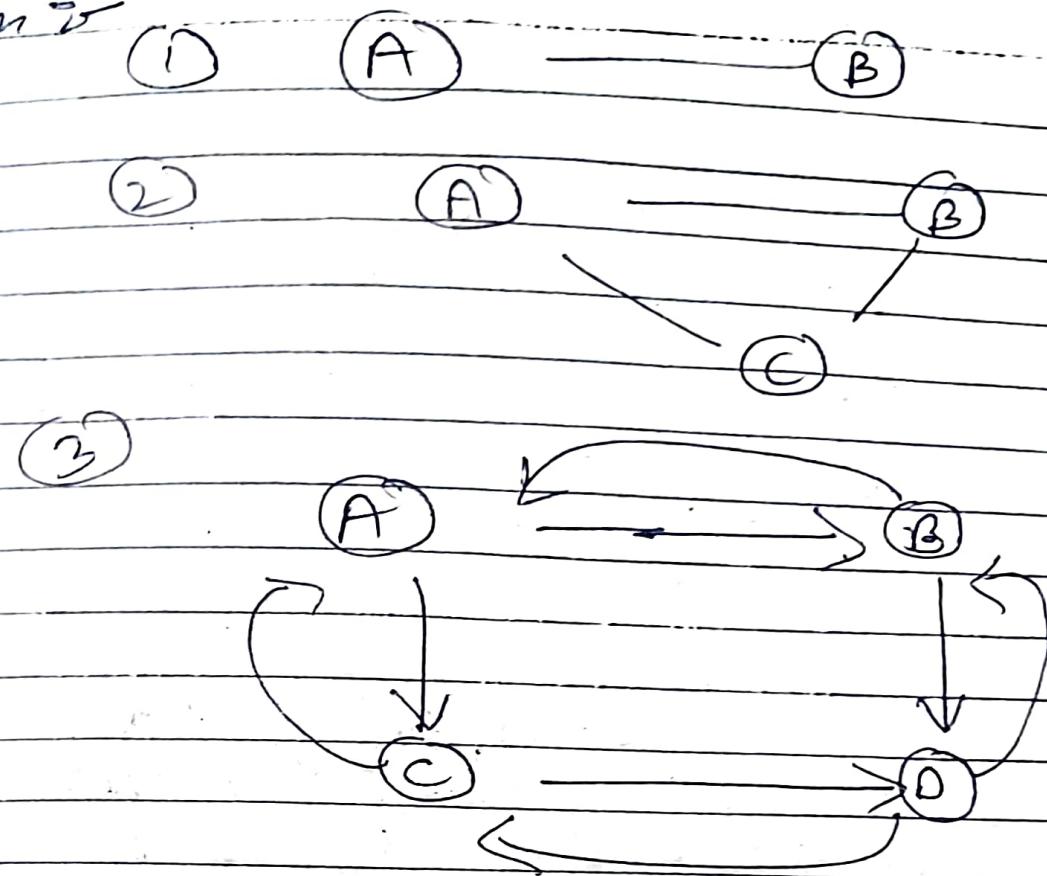
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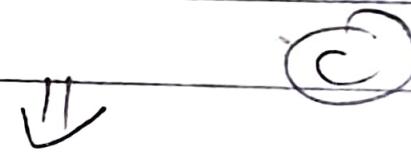
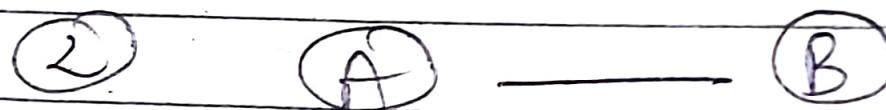
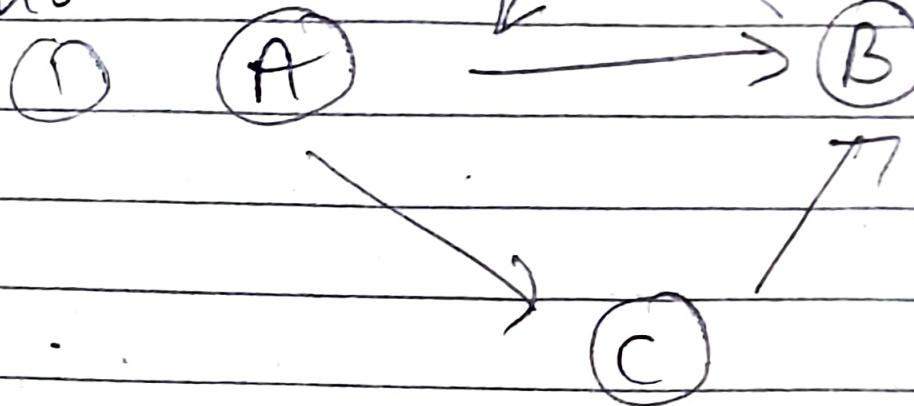
unit



2 Strongly Connected graph:

A graph in which every node is reachable to every node. If there is a path from u to v then surely there will be a path from v to u .

Ex:



It is not a strongly
connected graph.

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Balkharni

P.D.Q

void delete (struct Array & arr, int n) {

for (int i = 0; i < n; i++) {

for (int j = i + 1; j < n; j++) {

if (~~arr~~ . arr[i] ==

arr . arr[j]) :

del (~~arr~~ . arr, ~~i~~, n) ;

del (arr . arr, j, n) ;

}

}

}

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void del (~~struct~~ Array~~*~~ arr , int j, int n) {

~~defn~~ if (j == n - 1) {

$$n = n - 1 ;$$

} else {

while (j < n) {

~~Array . arr [j] =~~

~~Array . arr [j]~~

= ~~Array . arr [j + 1]~~

j ++;

}

}

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(1)

→ Struct emp {

int id;

int salary;

char name[100];

};

Void print (struct emp e[], int id) {

for (int i = 0; i < 10; i++) {

~~if (emp.e[i].id == id)~~

if (e[i].id == id) {

printf ("Name: %s",

e[i].name);

printf ("Salary: %d",

e[i].salary);

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Relchongazi

```
int main ( ) {  
    struct emp e[10];  
    // 10 employees  
    int id=5;  
    print ( emp , id );  
    return 0;  
}
```

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