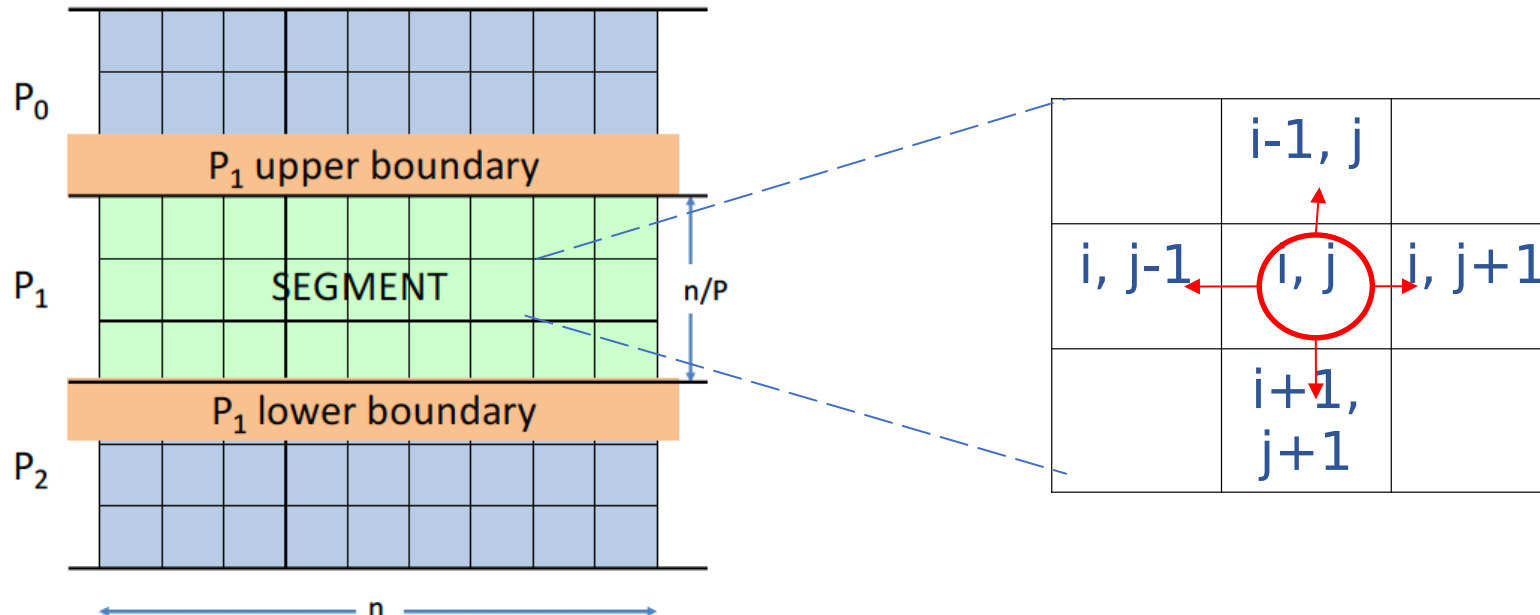


Stencil computation using Jacobi: Row-wise Distribution of the Matrices

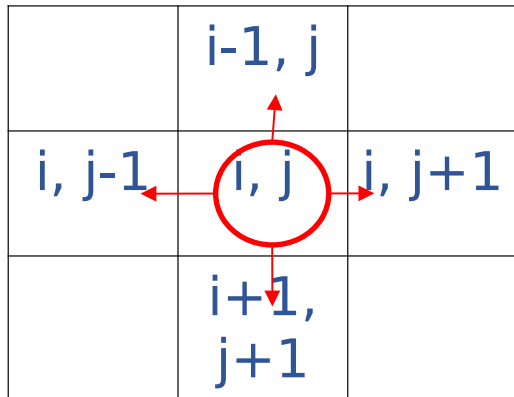
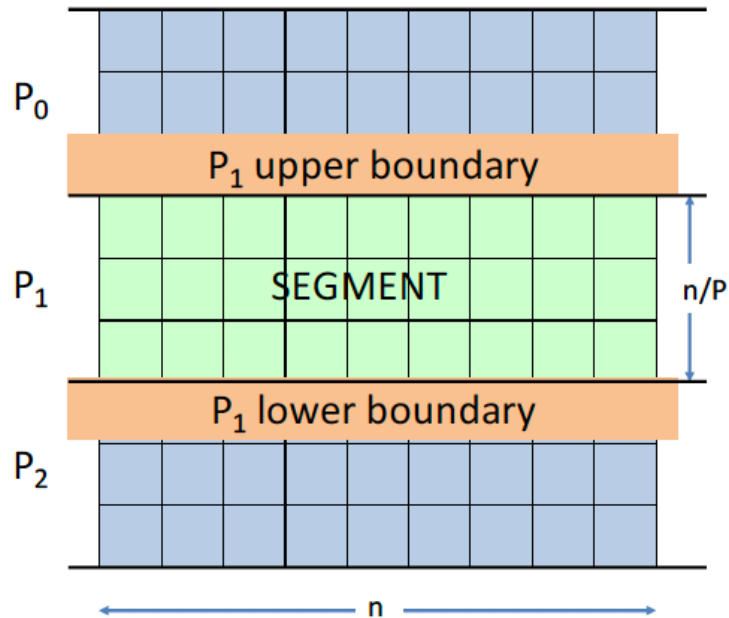
```
void compute(int n, double *u, double *utmp) {  
    int i, j;  
    double tmp;  
  
    for (i = 1; i < n-1; i++) {  
        for (j = 1; j < n-1; j++) {  
            tmp = u[n*(i+1) + j] + u[n*(i-1) + j] + // elements u[i+1][j] and u[i-1][j]  
                  u[n*i + (j+1)] + u[n*i + (j-1)] - // elements u[i][j+1] and u[i][j-1]  
                  4 * u[n*i + j]; // element u[i][j]  
            utmp[n*i + j] = tmp/4; // element utmp[i][j]  
        }  
    }  
}
```



Computation of each element requires access for reading four neighbors (upper, lower, left, right) of matrix u

Stencil computation using Jacobi: Row-wise Distribution of the Matrices

Different memory access situations can arise depending on where the element appears:

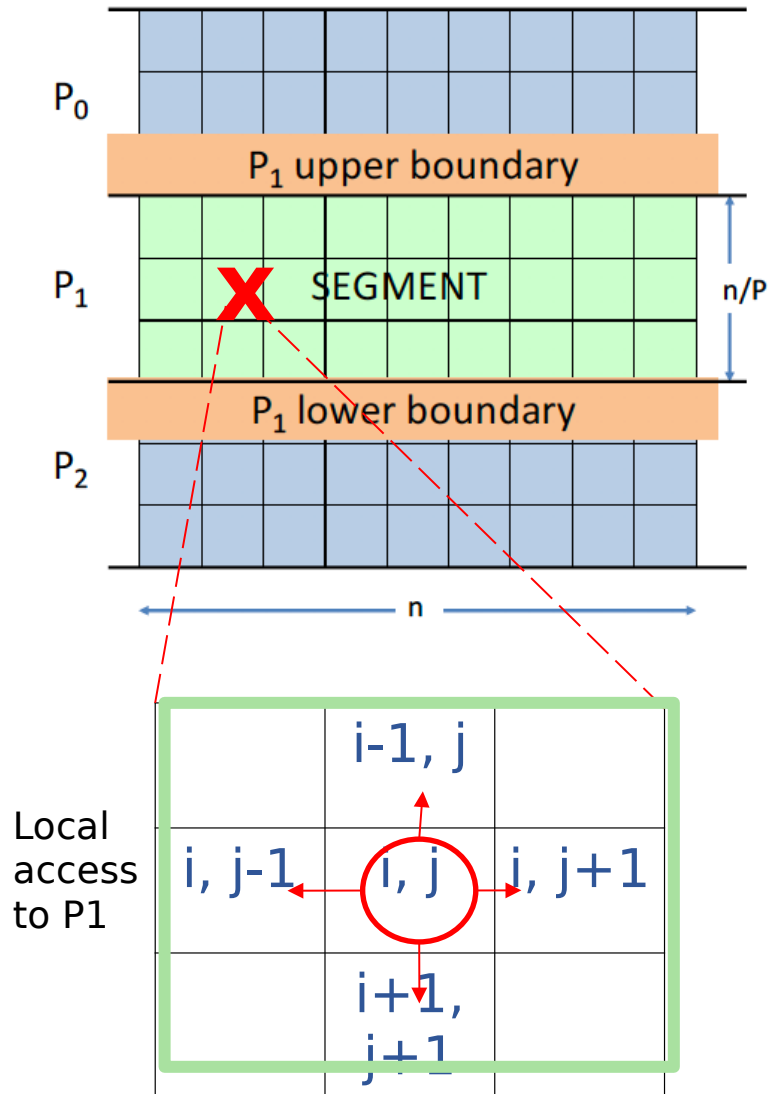


Stencil computation using Jacobi: Row-wise Distribution of the Matrices

Different memory access situations can arise depending on where the element appears:

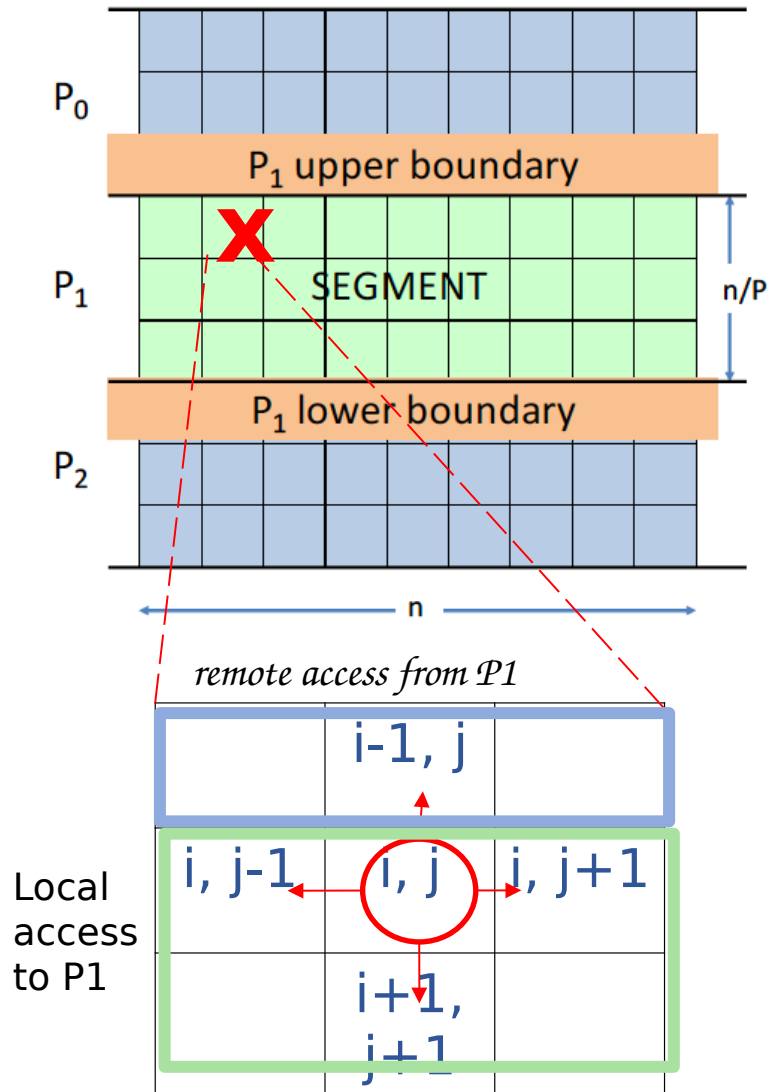
1. *The element is in a row that is not a border:*

All 4 neighbors are local to the current processor (P_1 in the example): local access to data, we assume zero overhead in our data sharing model



Stencil computation using Jacobi: Row-wise Distribution of the Matrices

Different memory access situations can arise depending on where the element appears:



1. The element is in a row that is not a border:

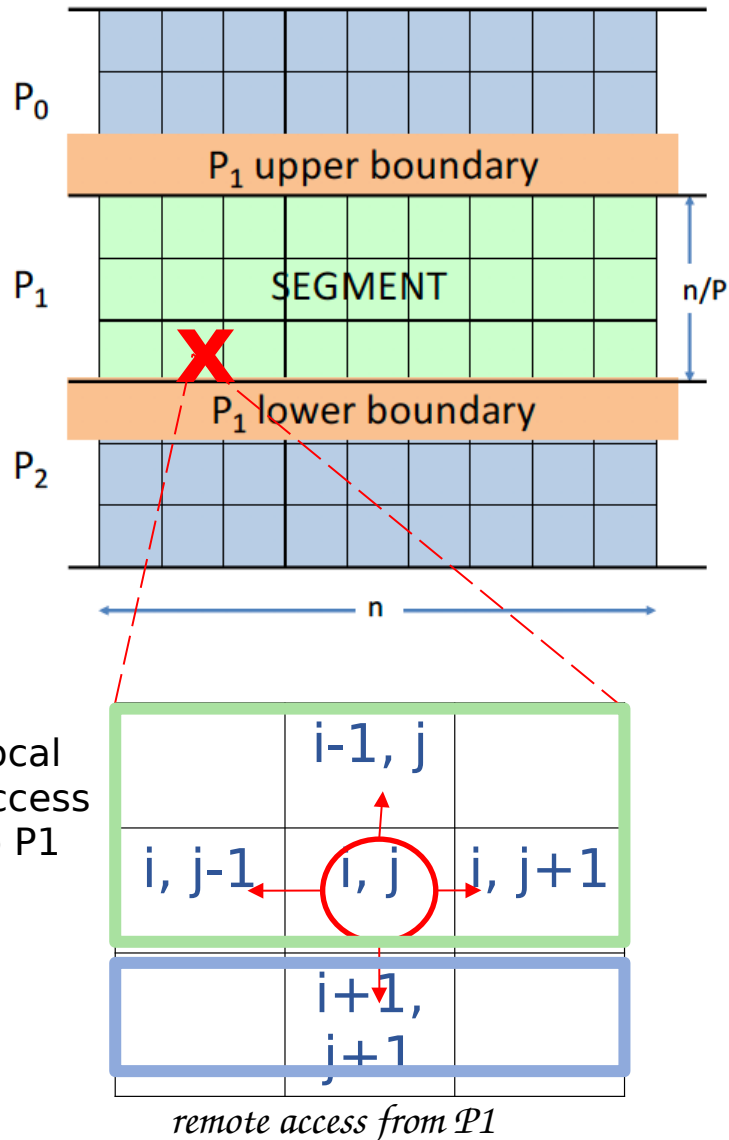
All 4 neighbors are local to the current processor (P_1 in the example): local access to data, we assume zero overhead in our data sharing model

2. The element is situated in a row in the upper border:

3 neighbors (left, right, lower) local to the current processor (P_1 in the example) and 1 neighbor (upper) with remote access (to P_0 in the example) with data sharing overhead

Stencil computation using Jacobi: Row-wise Distribution of the Matrices

Different memory access situations can arise depending on where the element appears:



1. The element is in a row that is not a border:

All 4 neighbors are local to the current processor (P_1 in the example): local access to data, we assume zero overhead in our data sharing model

2. The element is situated in a row in the upper border:

3 neighbors (left, right, lower) local to the current processor (P_1 in the example) and 1 neighbor (upper) with remote access (to P_0 in the example) with data sharing overhead

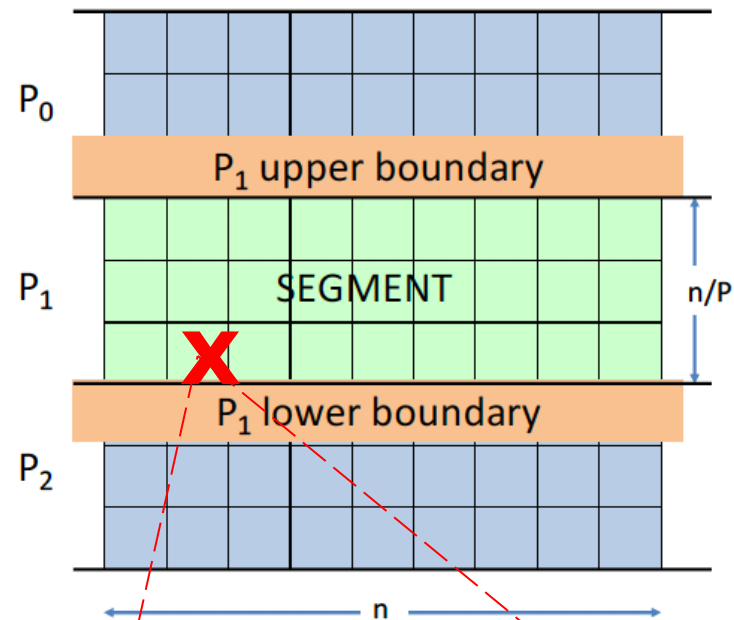
3. The element is situated in a row in the lower border:

3 neighbors (left, right, upper) local to the current processor (P_1 in the example) and 1 neighbor (lower) with remote access (to P_2 in the example) with data sharing overhead

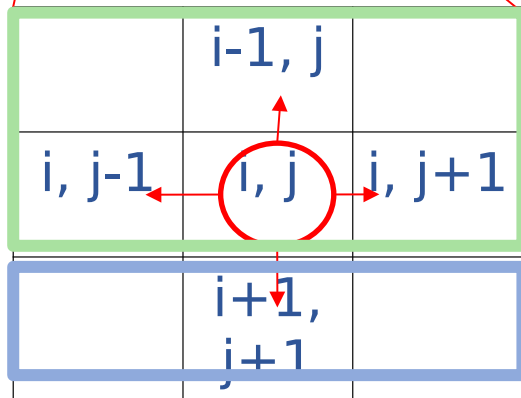
Stencil computation using Jacobi: Row-wise Distribution of the Matrices

Summarizing, there is remote memory access when processing elements in the rows:

- in the upper border of a segment
- in the lower border of a segment



Local
access
to P_1



remote access from P_1