THEORY OF INFORMATION, ARCHITECTURE OF COMPUTERS AND OPERATING SYSTEMS (TIACOS)

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Unit 5 Input/Output. Exercises

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1.	The first parameter of os.write(1,b'hello')
	\triangle represents a virtual device \checkmark
	\Diamond represents a logical device \times
	\bigcirc represents a physical device \times
	\Box is wrong \times
2.	The device driver
	\odot can execute privilege instructions \times
	\triangle implements access/control operations on physical devices \times
	\diamondsuit implements a different interface for each OS \times
	\Box all of the above \checkmark
3.	The first parameter of os.open("nom", os.O_WRONLY)
	\bigcirc represents a phsyical device \times
	\triangle represents a virtual device \times
	\Diamond represents a logical device \checkmark
	\square is wrong \times
4.	After executing the sequence: ret=os.close(0);fd=os.open(fn,os.O_WRONLY)
	\triangle ret==-1 and fd == -1 \times
	\odot ret==0 and we cannot know the value of fd $ imes$
	\Diamond ret==0 and fd==0 \checkmark
	\Box ret==0 and fd==-1 \times
5.	Mark the wrong answer.
	A process is using the file /etc/passwd. Executing a new open("/etc/passwd",O_WRONLY)
	\triangle If there is no error, creates a new entry in the File Descriptor Table \times
	\bigcirc If there is no error, creates a new entry in the File Table \times
	\diamondsuit If there is no error, creates a new entry in the Inode table \checkmark
	\square In the File Table we will have at least 2 entries related to /etc/passwd \times

6.	Mark the wrong answer.
	fd is a valid virtual device. Then os.close(fd)
	\triangle Always modifies the File Descriptor Table \times
	\bigcirc Always modifies the Inode Table \checkmark
	\diamond Always modifies the File Table \times
	\Box The inode in disk may be modified \times
7.	The contents of the file named nom are "123". After executing fd=os.open("nom",os.O_RDONLY); buf=os.read(fd,10)
	\triangle len(buf)==-1 and buf empty \times
	\odot len(buf)==3 and buf=="123" \checkmark
	\Diamond len(buf)==-1 and buf==123 \times
	\Box len(buf)==0 and buf=="123" \times
8.	The file "nom" contains "123" and buf==b'4'. After executing fd=os.open("nom",os.O_WRONLY) os.write(fd,buff)
	\triangle len(buf)==1 and nom=="423" \checkmark
	\odot len(buf)==4 and nom=="1234" \times
	\diamond len(buf)==-1 and nom=="123" \times
	\Box len(buf)==0 and buf=="123" \times
9.	The file nom contains "hola". After executing fd=os.open("nom",os.O_RDWR); os.read(fd,1); write(fd,b'2'), nom contains
	\triangle "20la" $ imes$
	\bigcirc hola $2 \times$
	$\diamond~ \mathrm{h2la}~ \checkmark$
	□ 2hola ×
10.	The following system call never blocks a process
	$ riangle$ pipe $ ilde{\checkmark}$
	\bigcirc read $ imes$
	\diamondsuit open $ imes$
	\square write $ imes$
11.	Would the following system call block the process? os.open("nom",os.O_RDWR)
	\triangle Never \checkmark
	\bigcirc Always \times
	\Diamond Only if nom is a named pipe \times
	\square Only if nom is a file \times

12. Would the following system call block the process? os.open("nom",os.O_WRONLY)
 △ Never × ○ Always × ◇ Only if nom is a named pipe and there aren't readers on it ✓ □ For all type of files, will block if there aren't readers ×
13. Could a process A blocked on a read receive a signal SIGPIPE?
 △ Yes, only if A is reading from a pipe without writers × ○ No, because reading do not cause SIGPIPEs × ◇ Yes, only if other process sends this signal to A ✓ □ None of the above ×
14. A process execute the following code: fd=os.pipe();c=os.read(fd[0],1)
 △ The process will block forever in the read ✓ ○ read will end returning -1 because there are no writers × ◇ read will end returning 0 because there are no writers × □ None of the above ×
15. We create A, a soft link to B. Mark the wrong answer
 △ File A will have a new inode × ○ if we delete B, the contents are deleted from disk × ◇ result of cat A is the same as cat B × □ This operation modifies the inode of B ✓
16. We create A, a hard link to B. Mark the wrong answer
 △ File A will have a new inode ✓ ○ If we delete B, the contents are not deleted from disk × ◇ The result of cat A is the same as cat B × □ This operation modifies the inode of B ×
17. How many blocks accesses an open system call?
 △ Depends on the path of the file ✓ ○ This system call does not accesses blocks × ◇ Only the first block of the file × □ None of the above ×
18. How many blocks accesses a read system call?
 △ Depends on the path of the file × ○ This system call does not accesses blocks × ◇ Only the inode of the file ✓ □ None of the above ×