TCP/IP in hardware using SME

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KU

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TCP/IP in hardware using SME

Mark siger introduktion og 2-3 saetninger "abstrakt"

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2. Implementation
3. Evaluation
4. Discussion
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Background and Motivation

FPGAs are making their way into data centers to boost the computing power and the overall power efficiency.

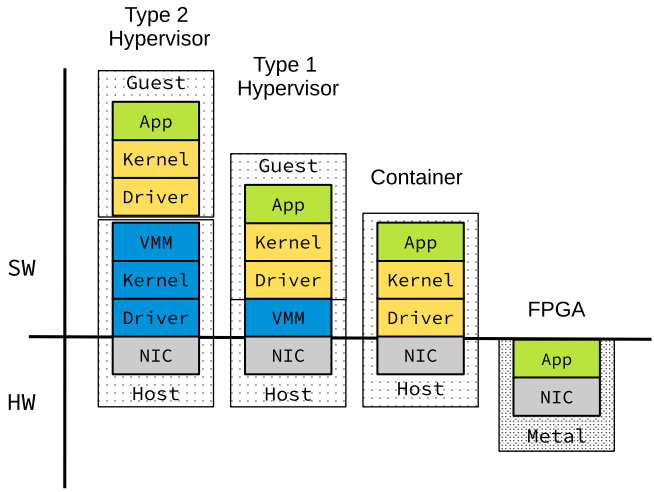
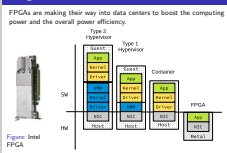


Figure: Intel FPGA



Applikationer og Big-Data udregninger flytter til Cloud, drevet af store data centre.

data-centre kraever meget plads, store maengder af stroem og er svaere at vedligeholde og udvide.

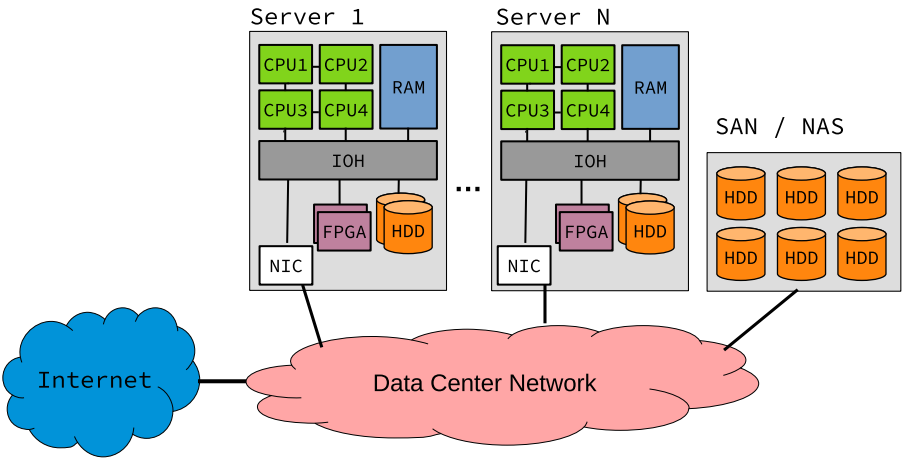
DC optimerer servere for at få mest værdi muligt

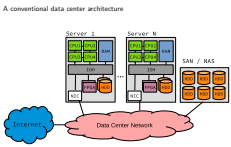
Tendens til aflaste beregninger til FPGAer, fjerne overhead

FPGA er hardware kan udføre beregninger hurtigt pga. dens parallele programmerbare natur. Den er hurtigt fordi instruktioner skrives direkte ned i hardwaren

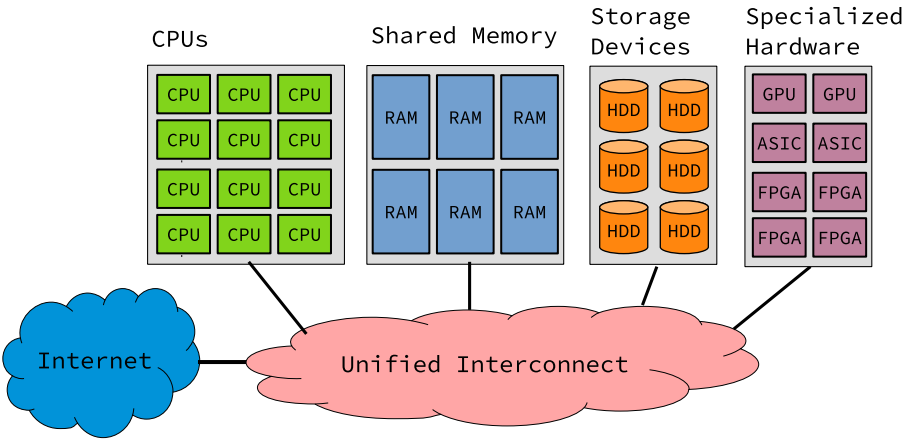
— GRAF HER —
PROBLEMET er at der kun kan vaere en begreanset antal af FPGAer i

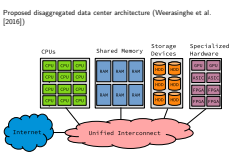
A conventional data center architecture





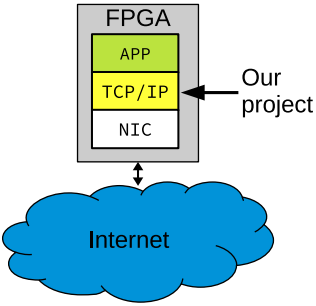
Proposed disaggregated data center architecture (Weerasinghe et al. [2016])





Hvis man splitter resourcerne op, kan man takket været FPGA få bedre ydeevne på det samme areal, samt nemmere håndtering af servere og deres komponenter.

FPGA usage



FPGA usage

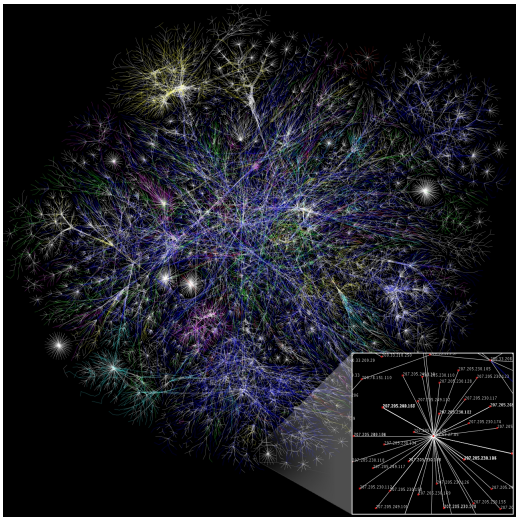
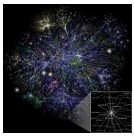


Figure: Map of about 30% of the accessible the endpoints on the Internet

The Internet



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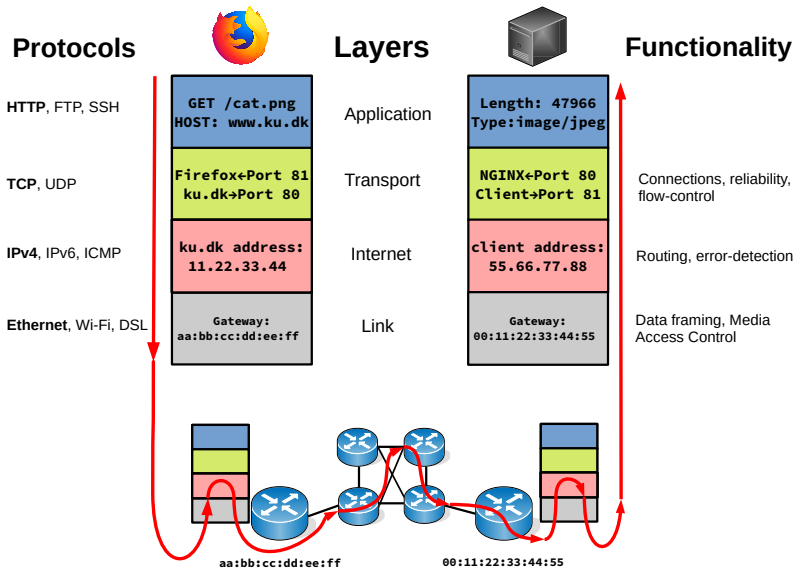
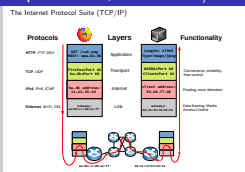
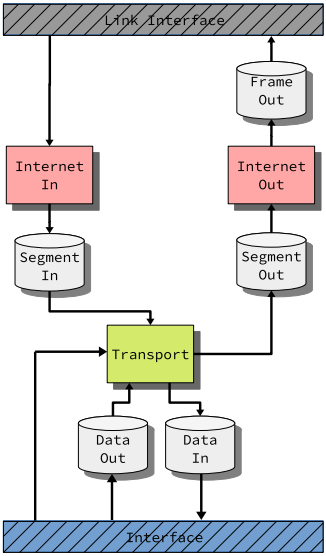




TCP/IP er samling af standarder og protokoller

Link: Overførsel på det fysiske medium
Internet: bestemmer data-veje, addressering, fejl-kontrol
Transport: pålidelighed, forbindelser, kontrol flow
Application: Defineret af selve applikationen

Design with the 4 layers in mind



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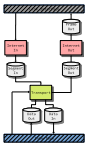

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2019-09-19

TCP/IP in hardware using SME
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2 Implementation
- SME introduction
- Processes
- Buffers
- Interface signal protocol

3 Evaluation
4 Discussion
5 Conclusion
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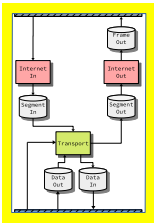
8 Demonstration



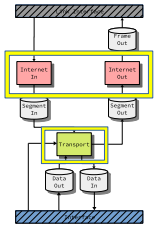
SME(Synchronous Message Exchange) introduction

- Processes and Busses
- Higher abstraction
- Handling of clocks
- Easy testing
- Not fully feature complete with C#(No threads, no allocation)

Implementation
SME introduction
SME(Synchronous Message Exchange) introduction
- Processes and Busses
- Higher abstraction
- Handling of clocks
- Easy testing
- Not fully feature complete with C#(No threads, no allocation)

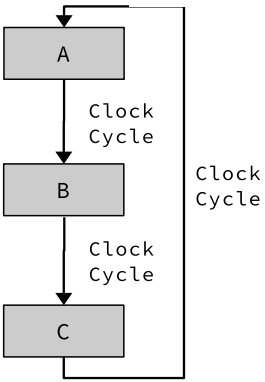
- What is a bus and a process
- No VHDL code
- Clocks abstracted away behind the management of processes and busses
- Testing straight in the simulator, but also in afterwards in the GHDL compiler, via an clock lookup table
- Since not feature complete, only simple structures can be used. We choose state diagrams since they are possible to make, and easy to understand

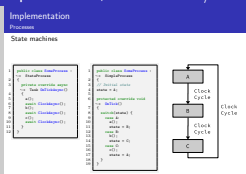
State machines




```
1  public class SomeProcess :
   ↪ StateProcess
2  {
3    private override async
     ↪ Task OnTickAsync()
4    {
5      a();
6      await ClockAsync();
7      b();
8      await ClockAsync();
9      c();
10     await ClockAsync();
11   }
12 }
```

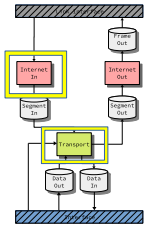
```
1  public class SomeProcess :
   ↪ SimpleProcess
2  {
3  // Initial state
4  state = A;
5
6  protected override void
   ↪ OnTick()
7  {
8    switch(state) {
9      case A:
10       a();
11       state = B;
12     case B:
13       b();
14       state = C;
15     case C:
16       c();
17       state = A;
18   }
19 }
```



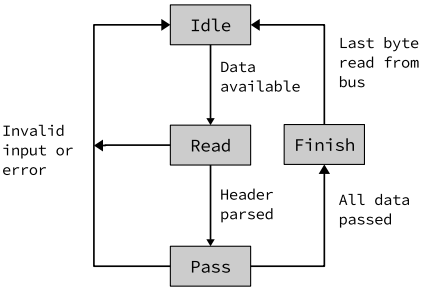


State machines

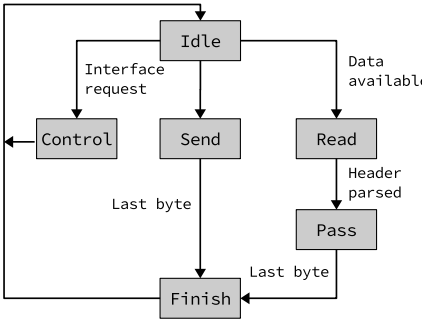
- `StateProcess`
  Eksekvering kan stoppes når som helst(i bidder)
- `SimpleProcess`
  Run er en clock altid, state machine håndteres med en switchcase.
  Algoritme kan splittes op i flere bidder, men kræver en state per bid



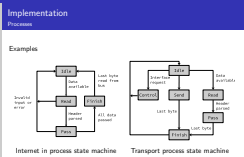
Examples



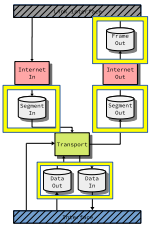
Internet in process state machine



Transport process state machine



- Gå igennem state diagrammer
- Snak om grundlaget for de forskellige typer brug

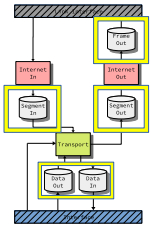


Why buffers?

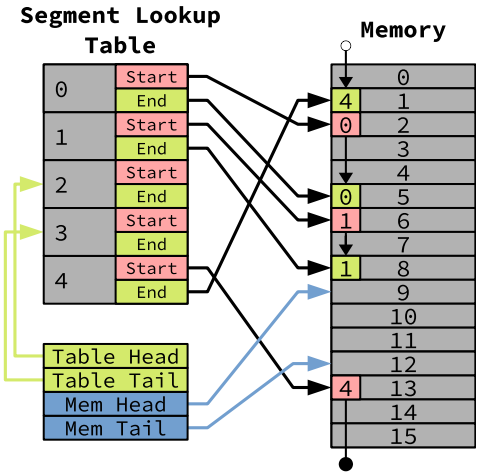
- Fixes segmentation
- Processes can get data at their leisure

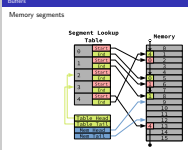
Implementation
Buffers
Why buffers?
- Fixes segmentation
- Processes can get data at their leisure

Hvorfor bruer vi buffers?

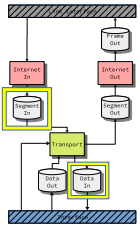


Memory segments

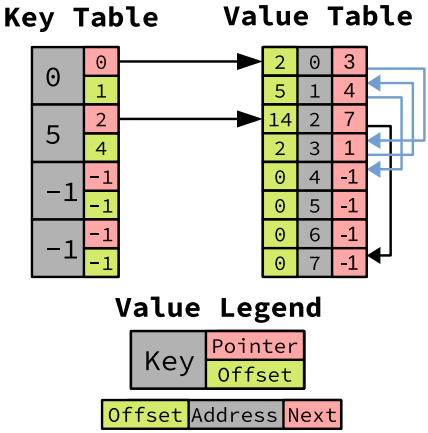




- Reason behind?
  - Segment handling
  - References to other segment to concatting of segments later



Memory dictionary

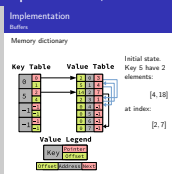


Initial state.
Key 5 have 2 elements:

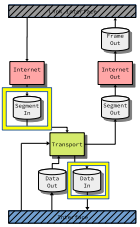
$$[4, 18]$$

at index:

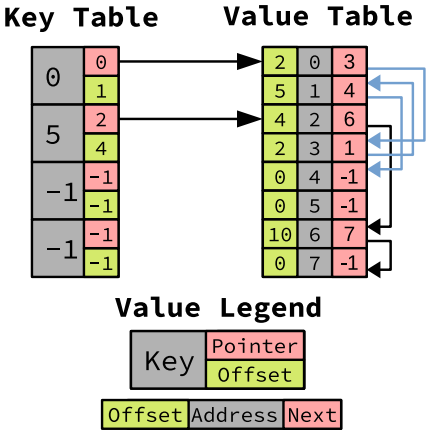
$$[2, 7]$$



Snak om input



Memory dictionary

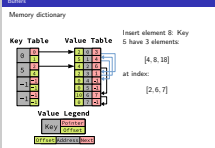


Insert element 8: Key 5 have 3 elements:

$$[4, 8, 18]$$

at index:

$$[2, 6, 7]$$



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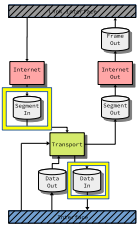
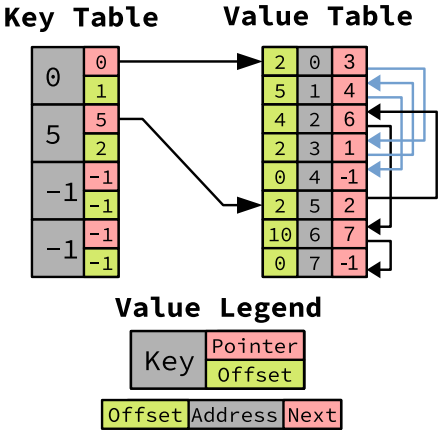


Memory dictionary

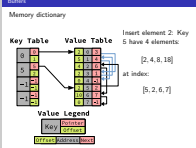


Insert element 2: Key 5 have 4 elements:

$$[2, 4, 8, 18]$$

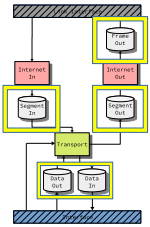
at index:

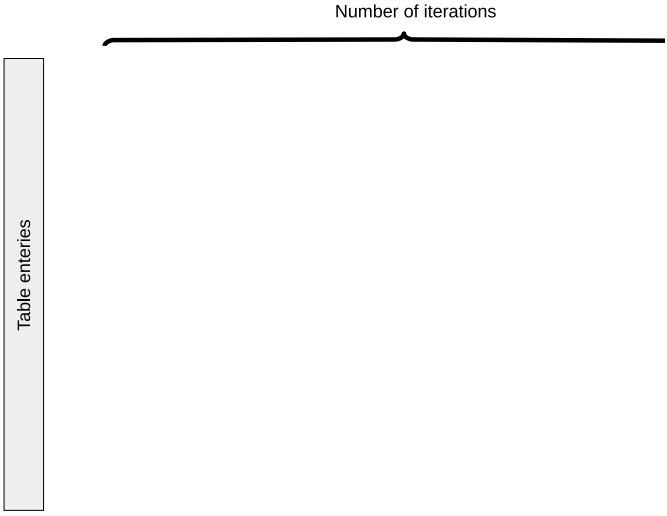
$$[5, 2, 6, 7]$$

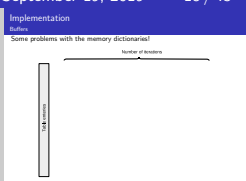


Snak om input

Some problems with the memory dictionaries!



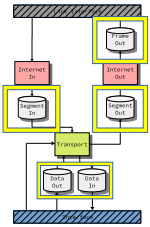


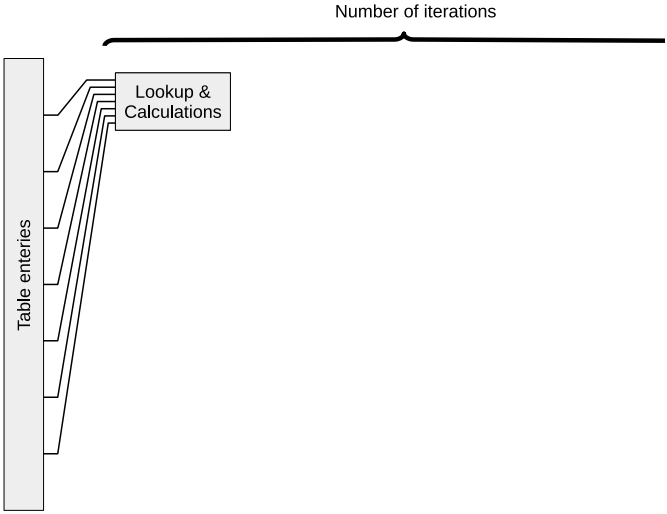


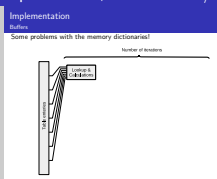
Overflow!
kan læses ved:

- Kør løkken en gang per clock
- Brug en anden model end en linked list, måske et fast offset?

Some problems with the memory dictionaries!







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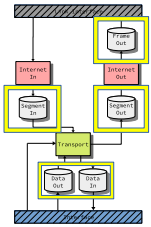

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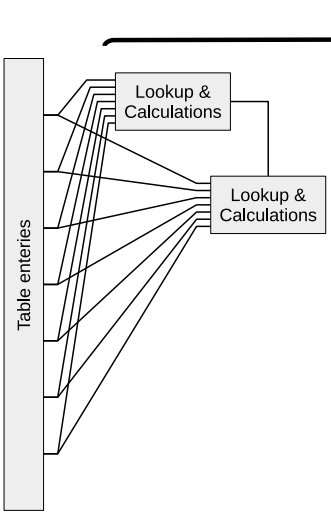

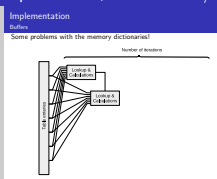

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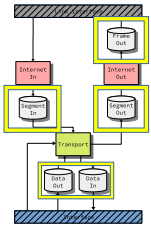

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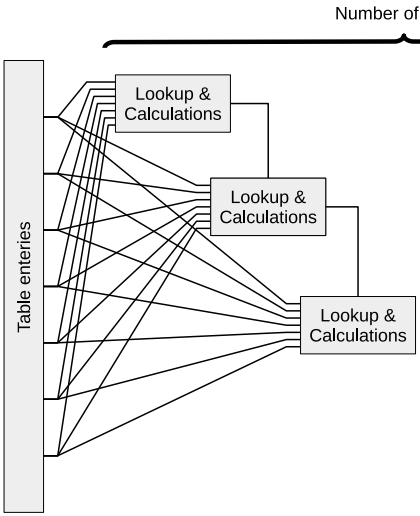

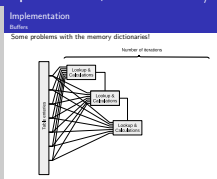

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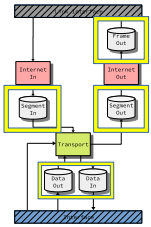

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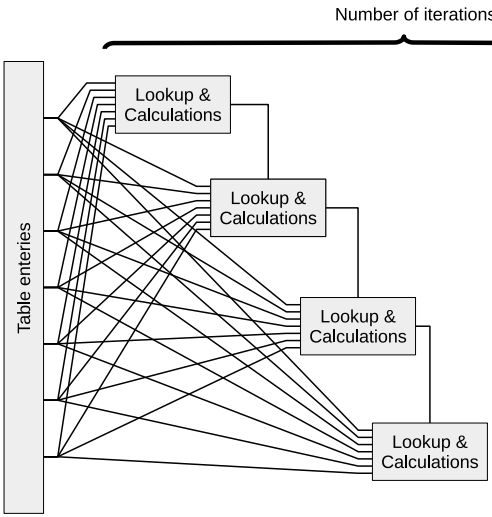

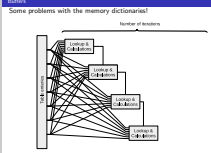

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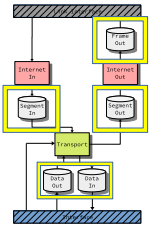

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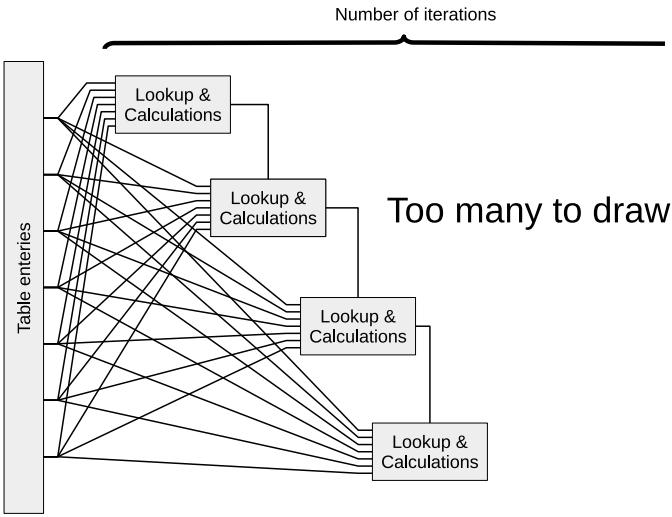

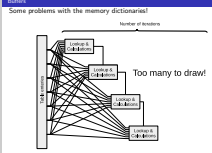

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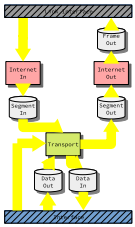
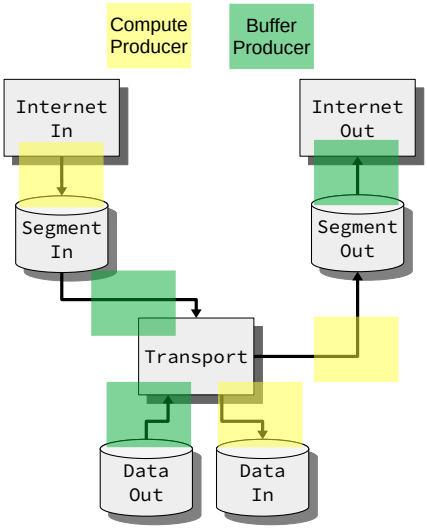
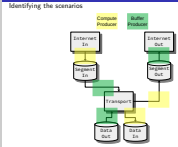


Identifying the scenarios

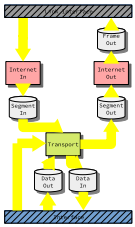


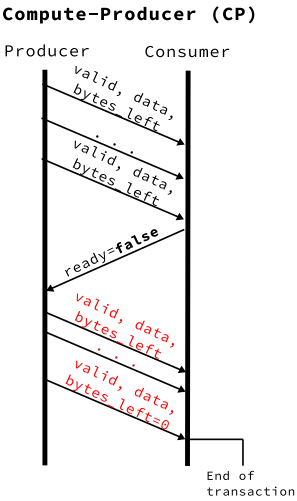


Data skal overføres hurtigst muligt, og det må ikke gå tabt

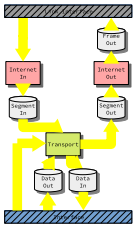
2 scenarier: fra "compute" til buffer, og omvendt

- CP kan ikke vente
- BP har stor buffer, og consumer starter transaktion

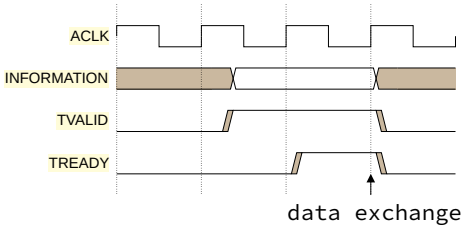




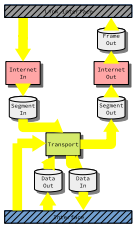
Implementation
Interface signal protocol



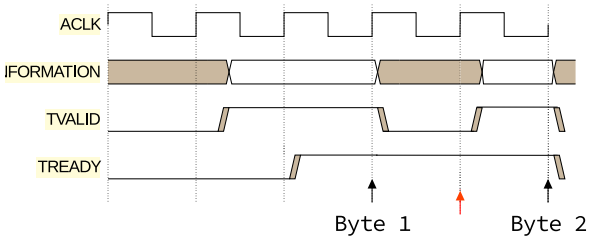
Buffer-Producer: Inspired by AXI4

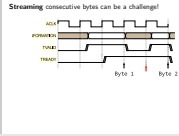


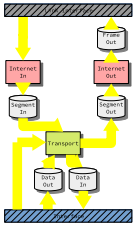
Implementation
Interface signal protocol
Buffer-Producer: Inspired by AXI4
ACLK
INFORMATION
TVALID
TREADY
data exchange



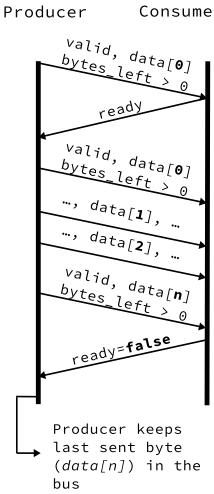
Streaming consecutive bytes can be a challenge!







Buffer-Producer (BP)



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Graph file simulation

- Full input - output
- Does not take latency between packets into account
- Simplifies test cases

Evaluation
Setup
Graph file simulation
- Full input - output
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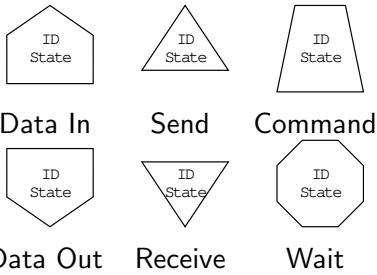
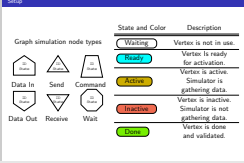
Definer send og receive bedre

Graph simulation node types



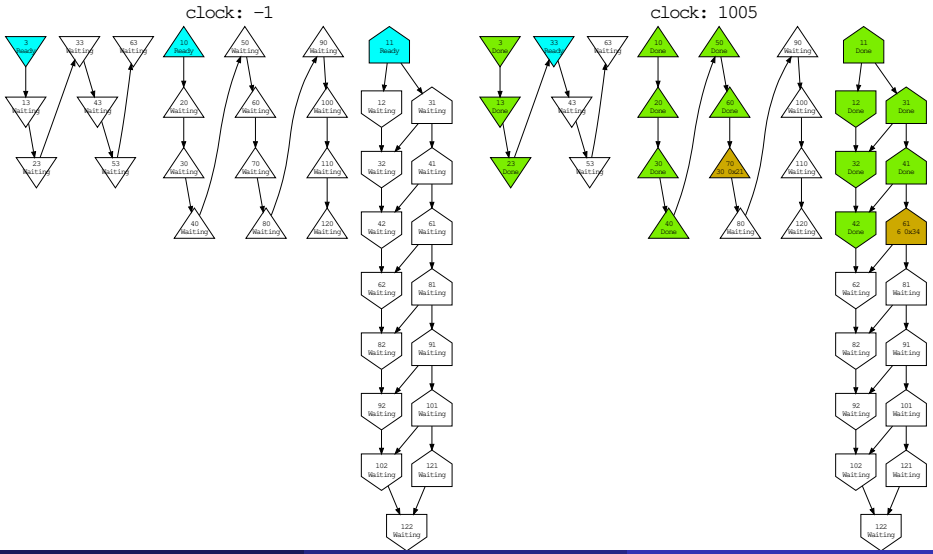
State and Color	Description
Waiting	Vertex is not in use.
Ready	Vertex Is ready for activation.
Active	Vertex is active. Simulator is gathering data.
Inactive	Vertex is inactive. Simulator is not gathering data.
Done	Vertex is done and validated.

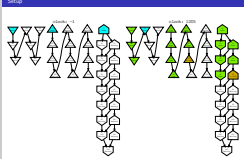


Hop til illustrationen på næste slide nå du snakker om det!

Evaluation

Setup





Evaluation
Test

Senario

- Real life scenario
- Test at high workloads
- Remove garbage
- Respond to packet
- Differ between concurrent connections

Evaluation
Test
Senario
Real life scenario
Test at high workloads
Remove garbage
Respond to packet
Differ between concurrent connections

Fortæl kun om hvad vi vil have, ikke hvad vi har lavet af test

The test

- 17283 packets in total
- Two "sessions"
- 640*2 UDP packets that needs a response
- 640 well formed UDP packets with no session (discard)
- Rest of data is "background noise" (TCP packets with state, data, etc)
- Total data sent through: 1832958 bytes
- 1.83 Million clocks used

Evaluation
Test
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Total data sent through: 1832958 bytes
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Latency calculations:

n_D : The number of bytes in the data part of the protocol. This excludes both headers from transport and internet.

n_I : The internet header size.

n_T : The transport header size.

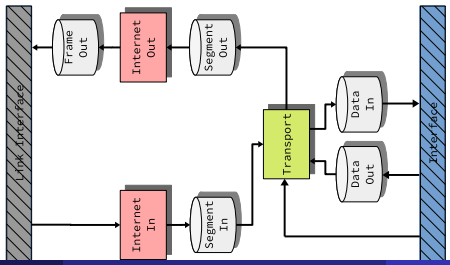
n : The total packet size.

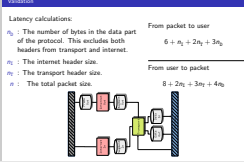
From packet to user

$$6 + n_I + 2n_T + 3n_D$$

From user to packet

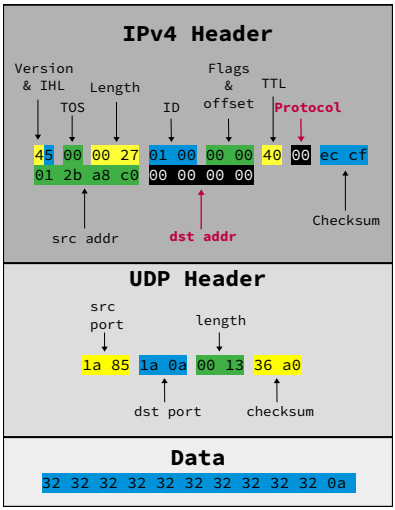
$$8 + 2n_I + 3n_T + 4n_D$$



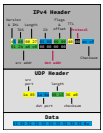


Bufferen kan ikke videresende data dirrekte, da den skal gemme segmentet først

Outgoing packet validation:



Evaluation
Validation
Outgoing packet validation:



Protocol ikke sat korrekt, destination ip ikke sat korrekt

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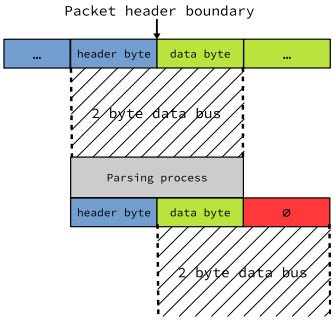
Discussion

Estimated performance:

$$1\ \textbf{Byte} * 10\ \textbf{MHz} = 80\ \textbf{Mbps}$$

Improving the performance:



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Improving the performance:

Discussion

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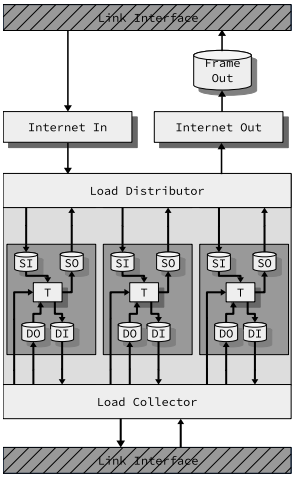

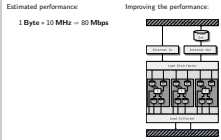

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Conclusion

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- All 17283 ingoing packets parsed correctly
- Few easily fixable errors in outgoing packets
- SME was great for implementation, albeit with a few small errors and bugs

Conclusion
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Design: Distribueret hukommelse er effektiv og giver funktionalitet

SME: C# nemt og hurtigt
Simulationen gjorde udviklingen hurtigere
Manglende features (structs), manglende library

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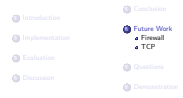
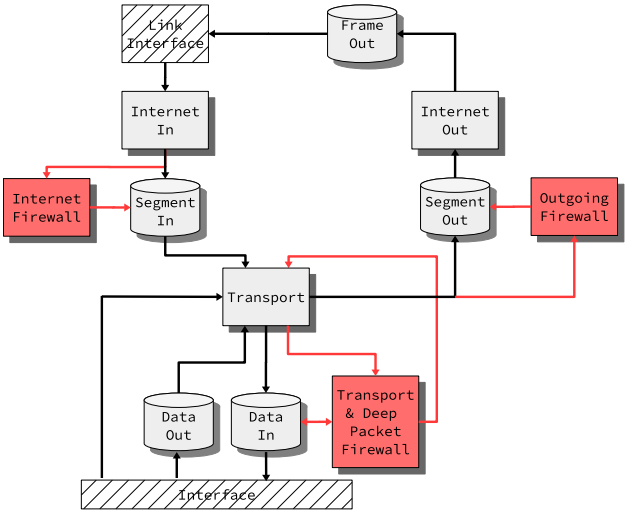
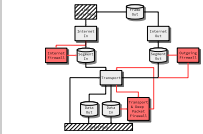
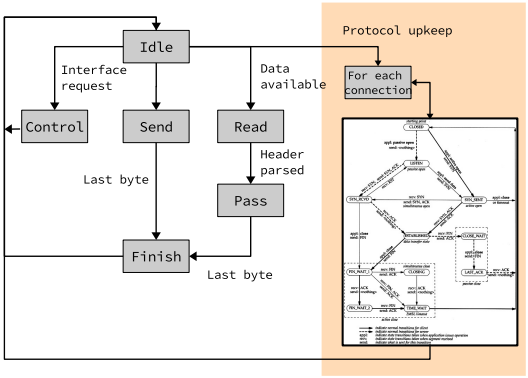






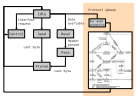
Integration med buffere. Hvad ville det indebære

Implementing TCP



TCP/IP in hardware using SME
└Future Work
└Future Work
2019-09-19

Future Work
TCP
Implementing TCP



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Bibliography

[1] J. Weerasinghe, F. Abel, C. Hagleitner, and A. Herkersdorf. Disaggregated fpgas: Network performance comparison against bare-metal servers, virtual machines and linux containers. In *2016 IEEE International Conference on Cloud Computing Technology and Science (CloudCom)*, pages 9–17, Dec 2016. doi: 10.1109/CloudCom.2016.0018.

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Demonstration

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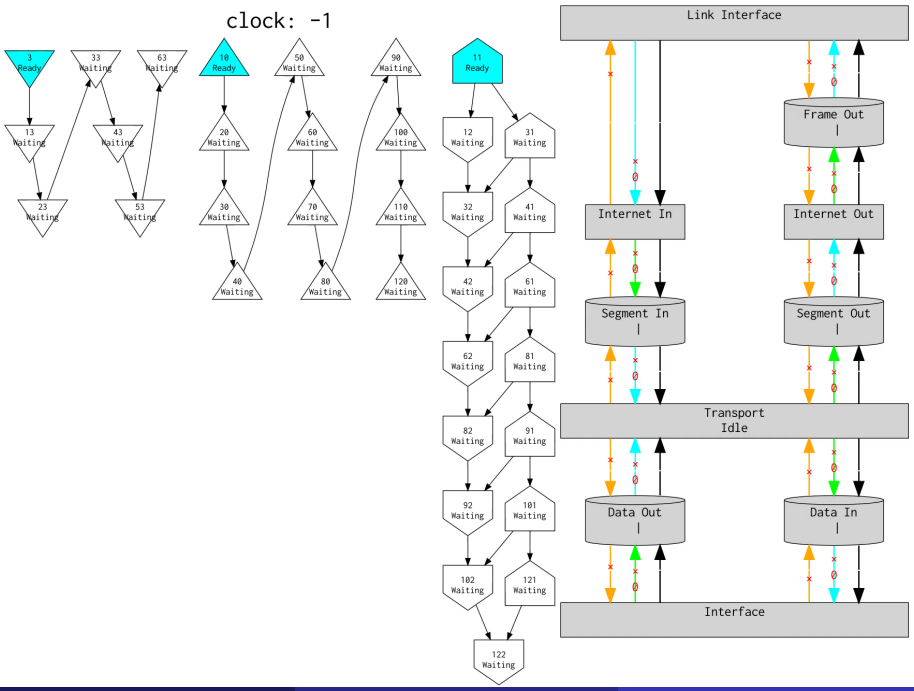

2019-09-19

TCP/IP in hardware using SME

└Demonstration