TCP/IP in hardware using SME

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KU

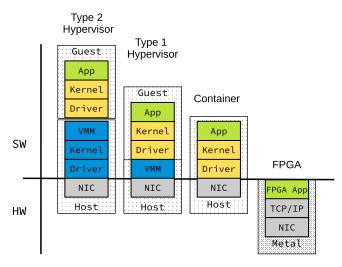
September 16, 2019

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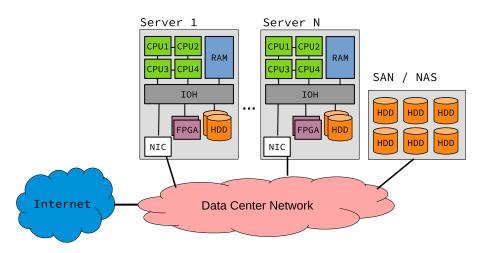
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Background and Motivation

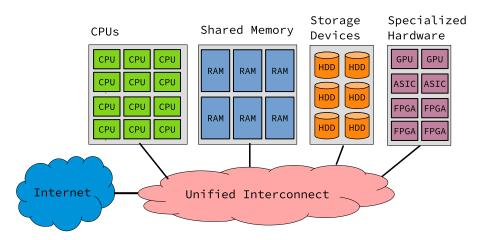
FPGAs are making their way into data centers to boost the computing power and the overall power efficiency.



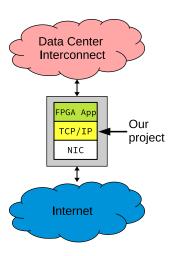
A conventional data center architecture



Proposed, disaggregated data center architecture



FPGA usage



The Internet

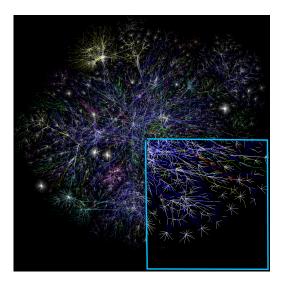
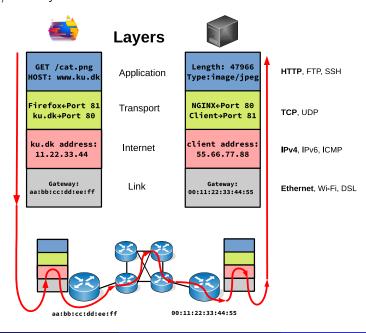


Figure: Map of the 30% of accessible the endpoints on the Internet

TCP/IP 4 layers



Design with the 4 layers in mind

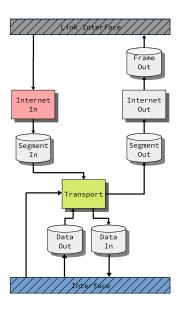


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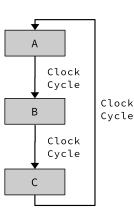
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- Processes
 - State machines
- Buffers
 - Memory segments
 - Dictionary
- Interface signal control
 - Buffer-Producer
 - Compute-Producer
- Interface control
 - Usage
 - Limitations

Processes

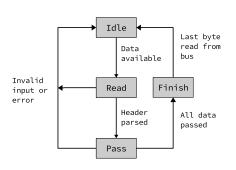
State machines

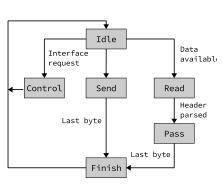
```
public class SomeProcess :
    \hookrightarrow SimpleProcess
     // Initial state
     state = A;
 5
     protected override void
    → OnTick()
 7
       switch(state) {
          case A:
10
            a();
11
            state = B;
12
          case B:
13
            b();
14
            state = C:
15
          case C:
16
            c();
17
            state = A;
18
19
```



Processes

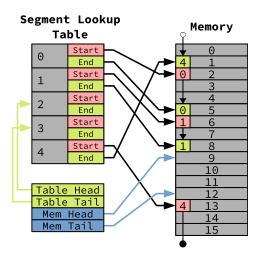
Examples





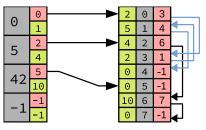
Buffers

Memory segments

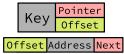


Buffers

Memory dictionary

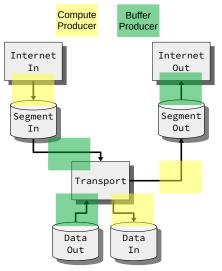


Value Legend



Interface signal protocol

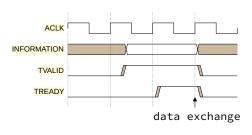
Identifying the scenarios



Interface signal protocol

Inspired by AXI4

- Single clock offset when sending data.
- Indicate end of stream with bytes_left.



Interface signal protocol

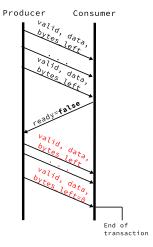
Buffer-Producer (BP) Consumer Producer valid, data[0] ready valid, data[**0**] ..., data[**1**],, data[2], ... valid, data[**n**] ready=false

Producer keeps

last sent byte

(data[n]) in the

Compute-Producer (CP)



bus

Interface protocol

The interface structures

```
enum InterfaceFunction : byte
 2
         INVALID = 0,
         // BIND = 1,
         LISTEN = 2.
         CONNECT = 3.
         ACCEPT = 4,
         CLOSE = 7.
10
         OPEN = 255,
11
     7
12
13
     struct InterfaceData
14
15
         public int socket:
16
         public uint ip;
17
         public byte protocol;
         public ushort port;
18
19
```

```
interface InterfaceBus : IBus
 2
         bool valid;
         byte interface_function;
 5
         InterfaceData request;
     }
     interface InterfaceControlBus : IBus
9
     {
10
         bool valid;
11
12
         byte exit_status;
13
         byte interface_function;
14
         InterfaceData request;
15
         InterfaceData response:
16
    }
```

Interface protocol

Limitations

- One request at a time.
- Arbitrary delay between request and response.

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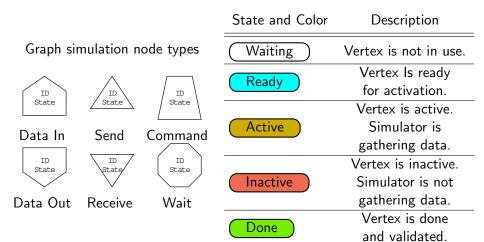
- Setup
 - Graph file simulator
- Test
- Validation
 - Latency
 - Outgoing packet validation
 - Internet Protocol Suite compliancy as per RFC 1122

Setup

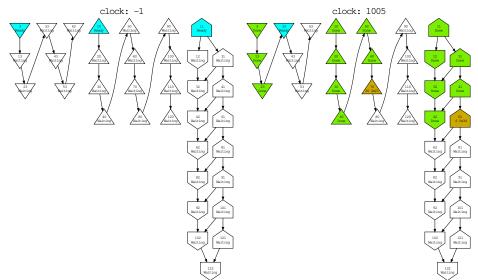
Graph file simulation

- Full input output
- Does not take latency between packets into account
- Simplifies test cases

Setup



Setup



Test

Senario

- Real life scenario
- Test at high workloads
- Remove garbage
- Respond to packet
- Differ between concurrent connections

Test

The test

- 17283 packets in total
- Two "sessions"
- 640*2 UDP packets that needs a response
- 640 well formed UDP packets with no session (discard)
- Rest of data is "background noise" (TCP packets with state, data, etc)
- Total data sent through: 1832958 bytes
- 1.83 Million clocks used

Validation

Latency calculations:

n_D: The number of bytes in the data part of the protocol. This excludes both headers from transport and internet.

 $n_{\rm T}$: The internet header size.

 $n_{\rm T}$: The transport header size.

n: The total packet size.

From packet to user

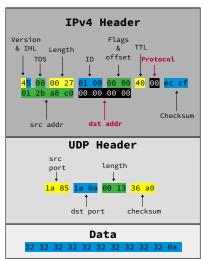
$$6+n_{\mathtt{I}}+2n_{\mathtt{T}}+3n_{\mathtt{D}}$$

From user to packet

$$8+2n_{\mathtt{I}}+3n_{\mathtt{T}}+4n_{\mathtt{D}}$$

Validation

Outgoing packet validation:



Validation

Internet Protocol Suite compliancy as per RFC 1122

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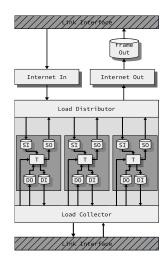
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Discussion

Improving the performance:

Estimated performance:

$$1 \text{ Byte} * 10 \text{ MHz} = 80 \text{ Mbps}$$



Usability SOMETHING

Using C#

State modelling Simulation Concurrency

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Conclusion

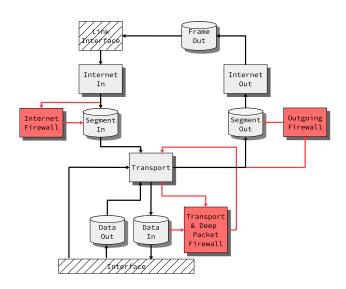
- Lot of trial and error to find the optimal design in the beginning
- In 10 mio. simulated clock cycles, 17283 packets were handled, 1280 of which were correctly received by the Application layer, and then sent out again
- Even with a few flaws, SME is a great framework for hardware modelling

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Future Work

Firewall



Future Work

How to implement TCP

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Questions

? Some random citation so it does not complain[Andrew S Tanenbaum(2013)]

Bibliography



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