



CRICOS PROVIDER 00123M

School of Computer Science

COMP SCI 1103/2103 Algorithm Design & Data Structure

MSSP + Sorting

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seek LIGHT

Previously on ADDS

- Binary search
- Benefits:
 - halve the search space every time
 - don't have to search every element
- Complexity – $O(\log n)$
 - Oh by the way, log with base 2 is denoted by lg. I should correct my previous word: the default base of log is 10
 - But we usually mean base 2 in computer science, and it does not make a difference in terms of Big O notation.
- Sorted data can be searched faster
- Sort once, search a lot

Overview

- See one more problem with different solutions (algorithms)
- Start the topic of Sorting

Example

- Maximum Subsequence Sum Problem
- Given (possibly negative) integers $A_1, A_2 \dots, A_n$, the target of the problem is to find the maximum value of $\sum_{k=i}^j A_k$, where $i, j \in [1, n]$.
- Example: -1, 2, 3, 6, -12, 13
- -1, 2, 3, 6, -8, 13
- There are many different algorithms to solve it and the performance of these algorithms varies drastically.

Algorithm1

//input: arr

Int maxsum=0

for(l=0 to arr.size)

 for(j=l to arr.size)

 {

 int sum=0;

 for(k=l to j)

 sum+=arr[k]

 if(sum> maxsum)

 maxsum=sum

 }

return maxsum

$O(n^3)$

Example -MSSP

- Algorithm 2
- $\sum_{k=i}^j A_k = A_j + \sum_{k=i}^{j-1} A_k$

$O(n^2)$

```
int maxSubSum2(int a[], int size){  
    int maxSum = 0;  
    for(int i=0; i<size; i++){  
        int sum = 0;  
        for(int j=i; j<size; j++){  
            sum+= a[j];  
            if(sum>maxSum)  
                maxSum = sum;  
        }  
    }  
    return maxSum;  
}
```

Divide and conquer strategy

- Divide – split the problem into two roughly equal subproblems which are then solved recursively
- Conquer – patch together the two solutions and possibly do a small amount of additional work to arrive at a solution to the whole problem.
- Algorithm 3 for MSSP
 - Can we use divide and conquer?
 - Yes. With complexity $O(n \log n)$

Example - MSSP

- Algorithm 4 with complexity in $O(n)$

```
int maxSubSum4(int a[], int size){  
    int maxSum, sum = 0;  
  
    for(int j=0; j<size; j++){  
        sum += a[j];  
        if(sum>maxSum)  
            maxSum = sum;  
        else if(sum<0)  
            sum = 0;  
    }  
    return maxSum;  
}
```


Sorting Algorithms

- Insertion Sort
- Selection Sort
- Bubble Sort
- Quicksort
- Merge Sort
- Heapsort (later; if we have enough time)
- Bucket sort

Insertion Sort

```
function insertionsort(list){  
  for i = 1 to length(list)  
    x = list[i]  
    j = i - 1  
    for j = i-1 to 0  
      if A[j] > x  
        A[j+1] ← A[j]  
      else  
        break for loop          // end inner for loop  
    A[j+1] = x  
  }
```

- Insertion sort is a simple sorting algorithm
- Complexity
 - worst-case $O(n^2)$
 - average-case $O(n^2)$
 - best-case $O(n)$

Selection Sort

- In selection sort, the list is divided into two part:
 - Sorted part (considering all elements)
 - Unsorted part

Input: list

For i=0 to n

{

 j= find the index with min value among list[i] to list[n]

 swap list[i] and list[j]

}

- Complexity
 - worst-case $O(n^2)$
 - average-case $O(n^2)$
 - best-case $O(n^2)$



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