

CRICOS PROVIDER 00123M

School of Computer Science

# COMP SCI 1103/2103 Algorithm Design & Data Structure Recursion 4

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#### Previously on ADDS

#### Recursion

- Checklist
- Recursive helper function
- Tail recursion
- Memorization

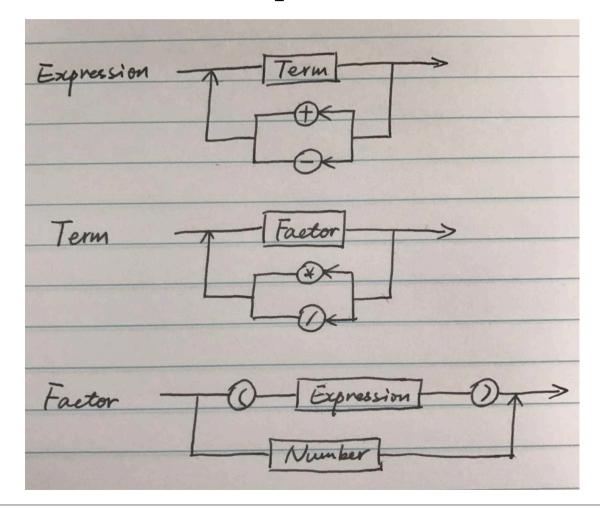
#### Indirect Recursion

- Harder to track and control
- Example: processing arithmetic expressions

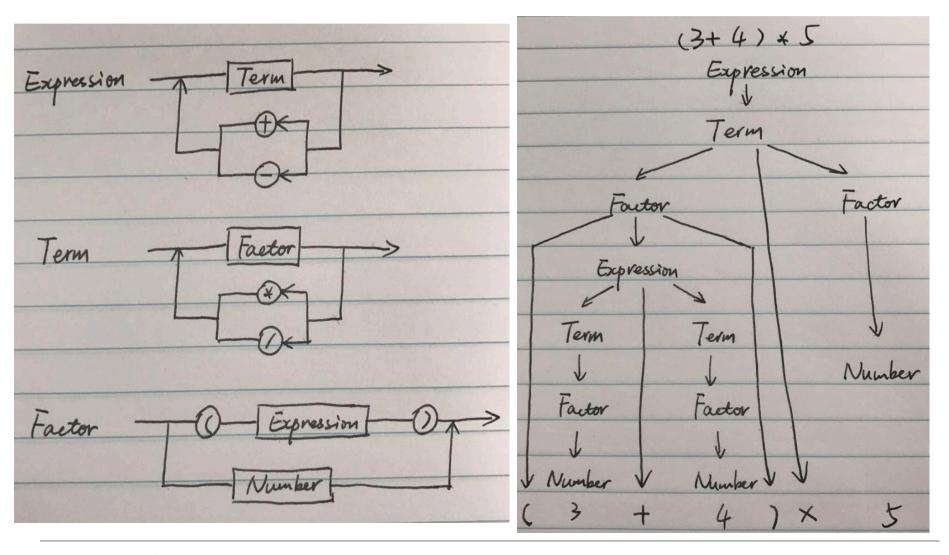
#### Overview

- In this lecture we will:
  - See a sample code for a function that can process arithmetic expressions (indirect recursion)
  - Discuss Dynamic Programming
    - Fibonacci
    - Counting Coin Problem

- Compute the values of arithmetic expressions
- Example
  - 3+4\*5
  - -(3+4)\*5



#### (3+4\*5)



### Sample code for getExptession

- Assume we have
  - getChar()looks ahead in input string
  - removeChar()removes a char from input string
  - getNumber()
     reads a number from the input string and removes the
     corresponding characters from that

### A sample code

```
int getExp(){
  int value = getTerm();
  char next = getChar();
  while ((next == '+') || (next == '-')) \{
        removeChar();
        int value2 = getTerm();
        if(next == '+')
           value += value2;
        else
           value -= value2;
  return value;
```

```
int getTerm(){
 int value = getFactor();
  char next = getChar();
 while ((next == '*') | | (next == '/')) 
         removeChar();
         int value2 = getFactor();
         if(next == "")
           value *= value2;
         else
           value /= value2;
 return value:
int getFactor(){
  int value;
  char next = getChar();
  if(next=='('){
       removeChar();
       value = getExp();
       removeChar(); // discard ')'
  }else
     value = getNumber();
  return value;
```

- For calculating (3+4)\*5
- getExp()
  - getTerm()
    - getFactor() ->consume '('
      - getExp() ->3+4, return 7
    - ->consume ')'
    - getFactor() -> return 5
  - 7\*5 -> return 35
- return 35

#### **Dynamic Programming**

- What do we do in dynamic programming?
  - Again, break the problem down to some smaller subproblems
  - solving each of them once
  - storing the solutions into some data structure (usually a table).
- Similar to Memorization, but DP does the computation from problems with smaller values to larger values.
- We need base cases and recursion relationship as well.

- Let's look back to the Fibonacci number again.
- What about calculating in a bottom-up order?

```
int fib(int n){
  int * fibTable= new int[n];

fibTable[0] = fibTable[1] = 1;

for(int i=2; i<n; i++){
  fibTable[i] = fibTable[i-1]+fibTable[i-2];
 }

return fibTable[n-1];
}</pre>
```

- Counting Coins
- Given a value n, if we need n cents, how many ways can we make the change?
  - Assume infinite supply of each kind of coins
- int coinValue[] = {5, 10, 20, 50, 100, 200};
- How many configurations can you find for 20 cents?
- What is the result for value n when you include i (0<=i<=maxPossible) coins of value x?
  - The result for n-i\*x, without considering coins of value x for that.

### Counting Coins Recursive Algorithm

```
int facevalue[] = {5, 10, 20, 50, 100, 200};
int count(int n, int coinIndex){
  if (n == 0)
           return 1:
   if (n < 0)
          return 0;
   if (coinIndex >= 6)
          return 0;
                                                 How can we also do this loop recursively?
  int counter = 0;
  for (int i = 0; i<= n/facevalue[coinIndex]; i++)</pre>
          counter += count(n - i*facevalue[coinIndex], coinIndex+1);
 return counter;
```

### Example 2 version 2

- Counting Coins
- Given a value n, if we need n cents, how many ways can we make the change?
  - Assume infinite supply of each kind of coins
- int coinValue[] = {5, 10, 20, 50, 100, 200};
- What is the result for value n when you include 0 or 1 coins of value x?
  - The results for n and next type of coins + the results for n-x and coins of value x again.

### Counting Coins Recursive Algorithm v2

```
int facevalue[] = {5, 10, 20, 50, 100, 200};
int countV2(int n, int coinIndex){

  if (n == 0)
      return 1;
  if (n < 0)
      return 0;
  if (coinIndex >= 6)
      return 0;

  return countV2(n, coinIndex+1) + countV2(n-facevalue[coinIndex], coinIndex);
}
```

	Coin types used					
						5
					10	10
				20	20	20
			50	50	50	50
		100	100	100	100	100
Total Amt	200	200	200	200	200	200
0	1	1	1	1	1	1
5	0 —	0	0	0	0	1
10	0	0	0	0	1	2
15	0	0	0	0	0	2
20	0	0	0 ?	1	2	4
25	0	0	0	0	0	4
30	0	0	0	0	2	6
35	0	0	0	0	0	6

#### Counting Coins with DP

```
int dptable[1000][6];
int countDP (int n, int coinIndex){
  // initialize the first row
  for(int j=0; j<6; j++){
   dptable[0][j] = 1;
  //fill in the table downwards
  for(int i=1; i<=n; i++){</pre>
    for(int j=5; j>=0; j--){
      if(i>=facevalue[j])
          dptable[i][j] = dptable[i][j+1] + dptable[i-facevalue[j]][j];
      else
          dptable[i][j] = dptable[i][j+1];
  return dptable[n][coinIndex];
```

### **Dynamic Programming**

• DP is just about filling table.

- Memory usage!
  - In some situations we can save the memory usage by reusing the table, e.g. using a circular data structure.

#### Summary

- Indirect recursion, we saw an example
- Dynamic programming
  - Fill out some tables and use the values recursively

