We acknowledge and pay our respects to the Kaurna people, the traditional custodians whose ancestral lands we gather on.

We acknowledge the deep feelings of attachment and relationship of the Kaurna people to country and we respect and value their past, present and ongoing connection to the land and cultural beliefs.



Computer Systems

Lecture 10: Assembler Review and Exercise



Exercise

When this program has finished assembling, what is the resulting machine code?

```
// Computes 1+...+RAM[0]
// And stored the sum in RAM[1]
    @i
        // i = 1
   M=1
   @sum
         // sum = 0
(LOOP)
         // if i>RAM[0] goto WRITE
    @i
   D=M
    @R0
   D=D-M
   @WRITE
   D; JGT
         // sum += i
   D=M
    @sum
   M=D+M
    @i
         // i++
   M=M+1
   @LOOP // goto LOOP
    0;JMP
(WRITE)
    @sum
   D=M
   @R1
   M=D // RAM[1] = the sum
(END)
    @END
    0;JMP
```



Symbol Table: Pre-defined Symbols

```
// Computes 1+...+RAM[0]
// And stored the sum in RAM[1]
    @i
   M=1
        // i = 1
    @sum
   M=0
         // sum = 0
(LOOP)
         // if i>RAM[0] goto WRITE
   D=M
    @R0
   D=D-M
   @WRITE
   D; JGT
         // sum += i
   D=M
    @sum
    M=D+M
    @i
        // i++
   M=M+1
   @LOOP // goto LOOP
    0;JMP
(WRITE)
    @sum
   D=M
   @R1
   M=D // RAM[1] = the sum
(END)
    @END
    0;JMP
```



Symbol Table: Pre-defined Symbols

Symbol	Address
R0	0
R1	1
	•••
R15	15
SP	0
LCL	1
ARG	2
THIS	3
THAT	4
SCREEN	16384
KBD	24576

```
// Computes 1+...+RAM[0]
// And stored the sum in RAM[1]
    @i
   M=1
        // i = 1
   @sum
   M=0
         // sum = 0
(LOOP)
         // if i>RAM[0] goto WRITE
   D=M
    @R0
   D=D-M
   @WRITE
   D; JGT
         // sum += i
   D=M
    @sum
    M=D+M
    @i
       // i++
   M=M+1
   @LOOP // goto LOOP
    0;JMP
(WRITE)
    @sum
   D=M
   @R1
   M=D // RAM[1] = the sum
(END)
    @END
    0;JMP
```



Code Clean-up

```
// Computes 1+...+RAM[0]
// And stored the sum in RAM[1]
   @i
        // i = 1
   M=1
   @sum
         // sum = 0
   M=0
(LOOP)
         // if i>RAM[0] goto WRITE
   D=M
   @R0
   D=D-M
   @WRITE
   D; JGT
         // sum += i
   D=M
   @sum
   M=D+M
       // i++
   @i
   M=M+1
   @LOOP // goto LOOP
   0;JMP
(WRITE)
   @sum
   D=M
   @R1
   M=D // RAM[1] = the sum
(END)
   @END
   0;JMP
```



Code Clean-up

```
@i
M=1
@sum
M=0
(LOOP)
@i
D=M
@R0
D=D-M
@WRITE
D;JGT
@i
D=M
@sum
M=D+M
@i
M=M+1
@LOOP
0;JMP
(WRITE)
@sum
D=M
@R1
M=D
(END)
@END
0;JMP
```



Symbol Table: After the 1st Pass

```
@i
M=1
@sum
M=0
(LOOP)
@i
D=M
@R0
D=D-M
@WRITE
D; JGT
@i
D=M
@sum
M=D+M
@i
M=M+1
@LOOP
0;JMP
(WRITE)
@sum
D=M
@R1
M=D
(END)
@END
0;JMP
```



Symbol Table: After the 1st Pass

Symbol	Address
Predefined symbols	015
LOOP	4
WRITE	18
END	22

```
@i
1 M=1
  @sum
3 M=0
  (LOOP)
  D=M
  @R0
7 D=D-M
8 @WRITE
9 D; JGT
10 @i
11 D=M
12 @sum
13 M=D+M
14 @i
15 M=M+1
16 @LOOP
17 0; JMP
↓ (WRITE)
18 @sum
19 D=M
20 @R1
21 M=D
↓ (END)
22 @END
23 0; JMP
```



Symbol Table: During the 2nd Pass

Symbol	Address
Predefined symbols	015
LOOP	4
WRITE	18
END	22

```
@i
1 M=1
  @sum
  M=0
   @i
  D=M
7 D=D-M
  @18
9 D; JGT
10 @i
11 D=M
12 @sum
13 M=D+M
14 @i
15 M=M+1
16 @<mark>4</mark>
17 0; JMP
18 @sum
19 D=M
20 @1
21 M=D
22 @22
23 0; JMP
```



Symbol Table: During the 2nd Pass

Symbol	Address
Predefined symbols	015
LOOP	4
WRITE	18
END	22
i	16
sum	17

```
@<u>i</u>
1 M=1
  @sum
3 M=0
   @<u>i</u>
  D=M
7 D=D-M
  @18
9 D; JGT
10 @i
11 D=M
12 @sum
13 M=D+M
14 @i
15 M=M+1
16 @<mark>4</mark>
17 0; JMP
18 @sum
19 D=M
20 @1
21 M=D
22 @22
23 0; JMP
```



During Code Generation

Symbol	Address				
Predefined symbols	015				
LOOP	4				
WRITE	18				
END	22				
i	16				
sum	17				

```
0 @16
1 M=1
   @17
  M=0
  @16
  D=M
7 D=D-M
  @18
9 D;JGT
10 @16
11 D=M
12 @17
13 M=D+M
14 @16
15 M=M+1
16 @<del>4</del>
17 0;JMP
18 @17
19 D=M
20 @1
21 M=D
22 @22
23 0; JMP
```



Translating / assembling A-instructions

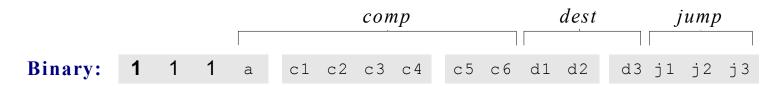
Translation to binary:

- □ If *value* is a non-negative decimal number, simple
- □ If *value* is a symbol...



Translating / assembling C-instructions

```
Symbolic: dest=comp; jump // Either the dest or jump fields may be empty. // If dest is empty, the "=" is ommitted; // If jump is empty, the ";" is omitted.
```



(when a=0)		_			_		(when a=1)	d1	d2	d3	Mnemonic	Destination	Destination (where to store the computed value)		
comp	c1	c2	c 3	c4	c5	c6	сотр	0	0	0	null	The value is not stored anywhere			
0	1	0	1	0	1	0		0	0	1	М	Memory[A	Memory[A] (memory register addressed by A)		
1	1	1	1	1	1	1		0	1	0	D	D register			
-1	1	1	1	0	1	0		0					er		
D	0	0	1	1	0	0		1	0	0	A	A register	,		
A	1	1	0	0	0	0	М					-	13.6	1	
!D	0	0	1	1	0	1		1 0 1 AM A register and Memory[A]					L]		
! A	1	1	0	0	0	1	! M	1	1	0	AD	A register and D register			
-D	0	0	1	1	1	1		1	1	1	AMD	A register,	A register, Memory[A], and D register		
-A	1	1	0	0	1	1	-M		j1		" j2	' јз	I		
D+1	0	1	1	1	1	1		(0	ut <	0)	(out = 0)	(out > 0)	Mnemonic	Effect	
A+1	1	1	0	1	1	1	M+1		0		0	0	null	No jump	
D-1	0	0	1	1	1	0			0		0	1	JGT	If $out > 0$ jump	
A-1	1	1	0	0	1	0	M-1		0		1	0	JEQ	If $out = 0$ jump	
D+A	0	0	0	0	1	0	D+M		0		1	1	JGE	If out ≥0 jump	
D-A	0	1	0	0	1	1	D-M		1		0	0	JLT	If out <0 jump	
A-D	0	0	0	1	1	1	M-D		1		0	1	JNE	If out ≠ 0 jump	
D&A	0	0	0	0	0	0	D∉M		1		1	0	JLE	If out ≤0 jump	
DIA	0	1	0	1	0	1	D M		1		1	1	JMP	Jump	



Question 1 1 pts

Look at the following (incomplete) diagram of the Hack CPU. Look at the wire pointed to by the large red arrow.

Where does the signal on this wire come from and what action does this signal trigger?

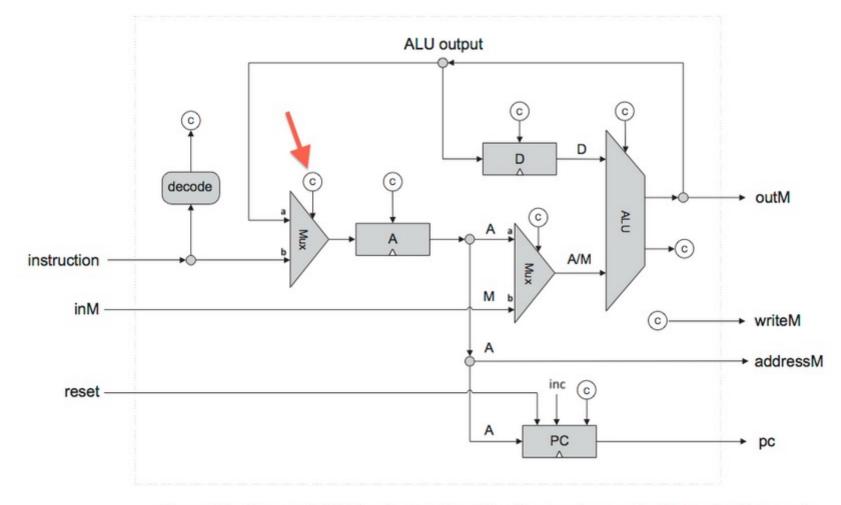


Figure 5.9 Proposed CPU implementation. The diagram shows only *data* and *address paths*, namely, wires that carry data and addresses from one place to another. The diagram does not show the CPU's *control logic*, except for inputs and outputs of control bits, labeled with a circled "c". Thus it should be viewed as an incomplete chip diagram.



Question 2

Look at the following (incomplete) diagram of the Hack CPU. Look at the wire (and it is a single wire) pointed to by the large red arrow.

Where does the signal on this wire come from and what action does this signal trigger?

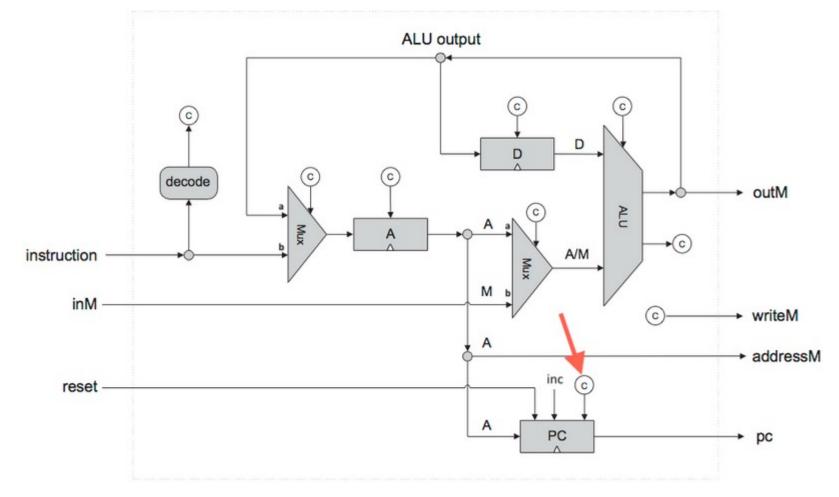


Figure 5.9 Proposed CPU implementation. The diagram shows only data and address paths, namely, wires that carry data and addresses from one place to another. The diagram does not show the CPU's control logic, except for inputs and outputs of control bits, labeled with a circled "c". Thus it should be viewed as an incomplete chip diagram.



Jump Unit (Assignment 4)

The Jump Unit provides a method for us to tell the CPU when to Jump to a different part of the running program.

This will allow us to have loops and conditional statements. In the Hack architecture, this is achieved using the Jump bits of the C-instruction, and comparing these to the ALU's output using its status bits.

- The j1, j2 and j3 bits of the C-instruction tell the CPU whether to perform a Jump if a condition is met.
 - If the j1 bit is set true, a jump should occur if the ALU's output is less than 0
 - If the j2 bit is set true, a jump should occur if the ALU's output is equal to 0
 - If the j3 bit is set true, a jump should occur if the ALU's output is greater than 0
- We can combine these bits:
 - If j1 and j2 are both true, a jump occurs if either the ALU's output is less than 0, or if the ALU's output is equal to 0
 - If all 3 bits are true, a jump always occurs (because the ALU's output is either < 0 or > 0 or 0)
 - If none of 3 bits are true, a jump never occurs
- We can determine whether the ALU's output is < 0 or > 0 or 0 by checking the ALU's status bits:
 - the zr bit will be true if the ALU's output is 0.
 - the ng bit will be true if the ALU's output is < 0.



Question 2 1 pts

Look at the following (incomplete) diagram of the Hack CPU. Look at the wire (and it is a single wire) pointed to by the large red arrow.

Where does the signal on this wire come from and what action does this signal trigger?

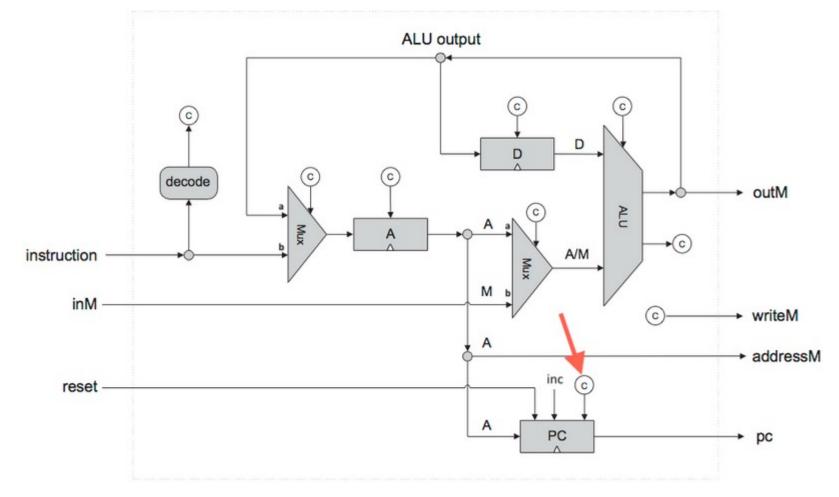
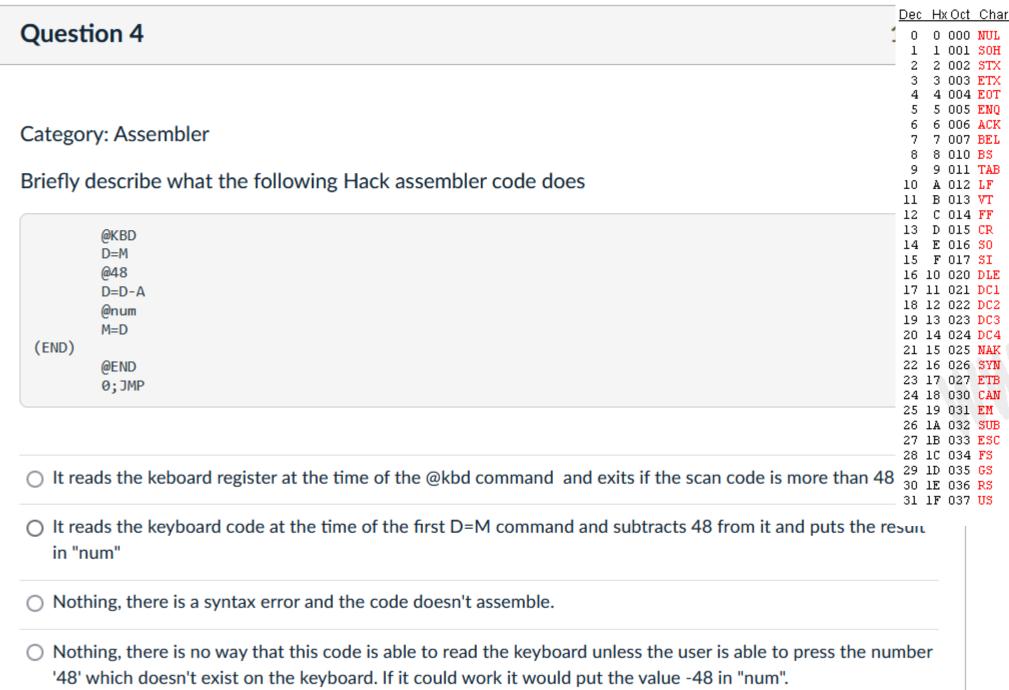


Figure 5.9 Proposed CPU implementation. The diagram shows only data and address paths, namely, wires that carry data and addresses from one place to another. The diagram does not show the CPU's control logic, except for inputs and outputs of control bits, labeled with a circled "c". Thus it should be viewed as an incomplete chip diagram.



Question 3 1 pts What does the following Hack assembler code always do to the current value in register D? D=!DD=D+1 O Sets D to be 0 O Sets D to be 1 - D O Sets D to be 1 O Sets D to be -D





Dec Hx Oct Html Chr Dec Hx Oct Html Chr Dec Hx Oct Html Chr 0 0 000 NUL (null) 32 20 040 Spac 64 40 100 @ 33 21 041 6#33; ! 65 41 101 A A 97 61 141 @#97; 8 1 1 001 SOH (start of heading) 34 22 042 4#34; " 66 42 102 B B 98 62 142 @#98; b 2 002 STX (start of text) 35 23 043 4#35; # 67 43 103 C C 3 3 003 ETX (end of text) 99 63 143 4#99; 36 24 044 \$ \$ 68 44 104 D D 100 64 144 @#100; d 4 4 004 EOT (end of transmission) 37 25 045 4#37; % 69 45 105 @#69; E 101 65 145 e e 5 005 ENQ (enquiry) 6 006 ACK (acknowledge) 38 26 046 4#38; 4 70 46 106 @#70; F 102 66 146 @#102; f 39 27 047 @#39; ' 71 47 107 4#71; 🚱 |103 67 147 @#103; g 7 7 007 BEL (bell) 40 28 050 6#40; (72 48 110 @#72; H 104 68 150 @#104; h 8 010 BS (backspace) 41 29 051 6#41;) 73 49 111 @#73; I 105 69 151 @#105; (horizontal tab) 9 9 011 **TAB** (NL line feed, new line) 42 2A 052 @#42; * 74 4A 112 @#74; J 106 6A 152 @#106; A 012 LF 43 2B 053 + + 75 4B 113 6#75; K 107 6B 153 k k 11 B 013 VT (vertical tab) 44 2C 054 , , 76 4C 114 L L 108 6C 154 l 1 (NP form feed, new page) 12 C 014 FF (carriage return) 45 2D 055 @#45; -77 4D 115 @#77; M 109 6D 155 @#109; 10 13 D 015 CR 14 E 016 SO (shift out) 46 2E 056 . . 78 4E 116 N N 110 6E 156 n n 47 2F 057 / / 79 4F 117 O 0 15 F 017 SI (shift in) |111 6F 157 o 0 48 30 060 4#48; 0 80 50 120 P P 112 70 160 @#112; p 16 10 020 DLE (data link escape) 49 31 061 4#49; 1 81 51 121 6#81; 0 113 71 161 @#113; <mark>q</mark> 17 11 021 DC1 (device control 1) 50 32 062 4#50; 2 82 52 122 6#82; R 18 12 022 DC2 (device control 2) 115 73 163 @#115; S 51 33 063 3 3 83 53 123 4#83; 5 19 13 023 DC3 (device control 3) 52 34 064 6#52; 4 84 54 124 @#84; T 116 74 164 t t 20 14 024 DC4 (device control 4) 53 35 065 4#53; 5 85 55 125 @#85; U |117 75 165 @#117; <mark>u</mark> 21 15 025 NAK (negative acknowledge) 54 36 066 6#54; 6 86 56 126 V V 118 76 166 v ♥ 22 16 026 SYN (synchronous idle) 23 17 027 ETB (end of trans. block) 55 37 067 4#55; 7 87 57 127 @#87; W |119 77 167 w ₩ 56 38 070 8 8 88 58 130 X X |120 78 170 x <mark>×</mark> 24 18 030 CAN (cancel) 57 39 071 4#57; 9 89 59 131 Y Y 121 79 171 @#121; Y 25 19 031 EM (end of medium) 58 3A 072 @#58;: 90 5A 132 6#90; Z 122 7A 172 @#122; Z 26 1A 032 SUB (substitute) 27 1B 033 ESC (escape) 59 3B 073 ;; 91 5B 133 [[123 7B 173 { 60 3C 074 < < 92 5C 134 @#92; 124 70 174 @#124; (file separator) 28 1C 034 FS 125 7D 175 @#125; (group separator) 61 3D 075 = = 93 5D 135] 29 1D 035 **GS** 30 1E 036 RS (record separator) 62 3E 076 >> 94 5E 136 @#94; 126 7E 176 @#126; ~ 127 7F 177 DEL 63 3F 077 ? ? 31 1F 037 US (unit separator) 95 5F 137 _ _ Source: www.LookupTables.com



Question 5 1 pts Category: Assembler What does the following Hack assembler code do? @pix M=1D=M M=M+D M=M+1D=M M=M+D M=M+1D=M@SCREEN M=D (END) @END 0;JMP It draws one pixel black. O It draws no pixels black because of overflow. O It draws the following pixels black, pixel 0, pixel 16 and pixel 32. It draws the topmost leftmost three pixels on the screen black.



Question 5

1 pts

Category: Assembler

What does the following Hack assembler code do?

		Α	D	M	@pix	@SCREEN
	@pix	16	?	?	?	?
	M=1					
	D=M					
	M=M+D					
	M=M+1					
	D=M					
	M=M+D					
	M=M+1					
	D=M					
	@SCREEN	16384				
	M=D					
(END)						
	@END					
	0;JMP					



This Week

- Review Chapters 5 & 6 of the textbook (if you haven't already)
- Assignment 3 extension due this Sunday.
- Assignment 4 due this Sunday.
- Review Chapter 7 of the textbook before next week.