

Introduction to Embedded Systems

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Exercise 4 – Real Time

Overview

- ▶ **Basics**
- ▶ **OSEK**
- ▶ **Real Time and Resources**
- ▶ **Periodic Scheduling**

Basics

Task 1:



Task 1:

a) Name the **two requirements** needed for real time.

- ① computation finishes before Deadline
- ② Computation is correct

for embedded systems

- ① reaction/response time must be \leq maximal value

real time = in time

Task 1:

DL = Deadline

b) Explain the terms:

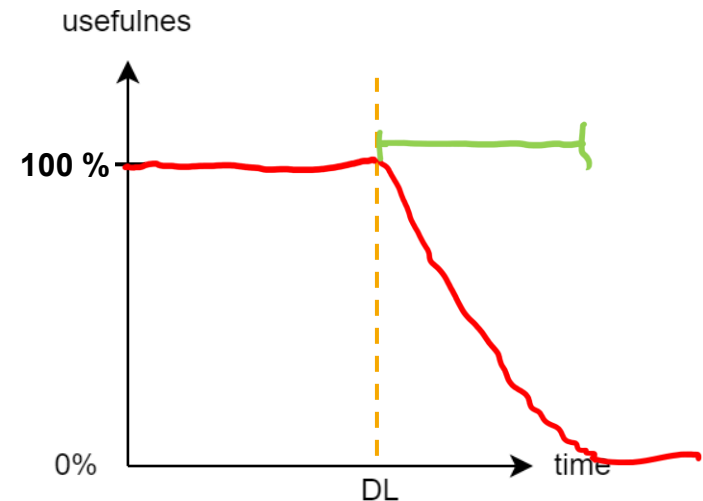
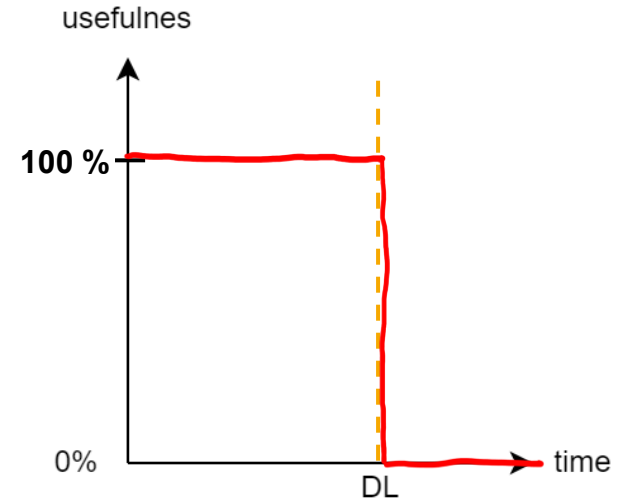
Hard real time

- useless after DL
- airbag

- more important to adhere DL

Soft real time

- still useful for certain amount of time
- marker on navi
- less important to adhere DL





OSEK

Task 2

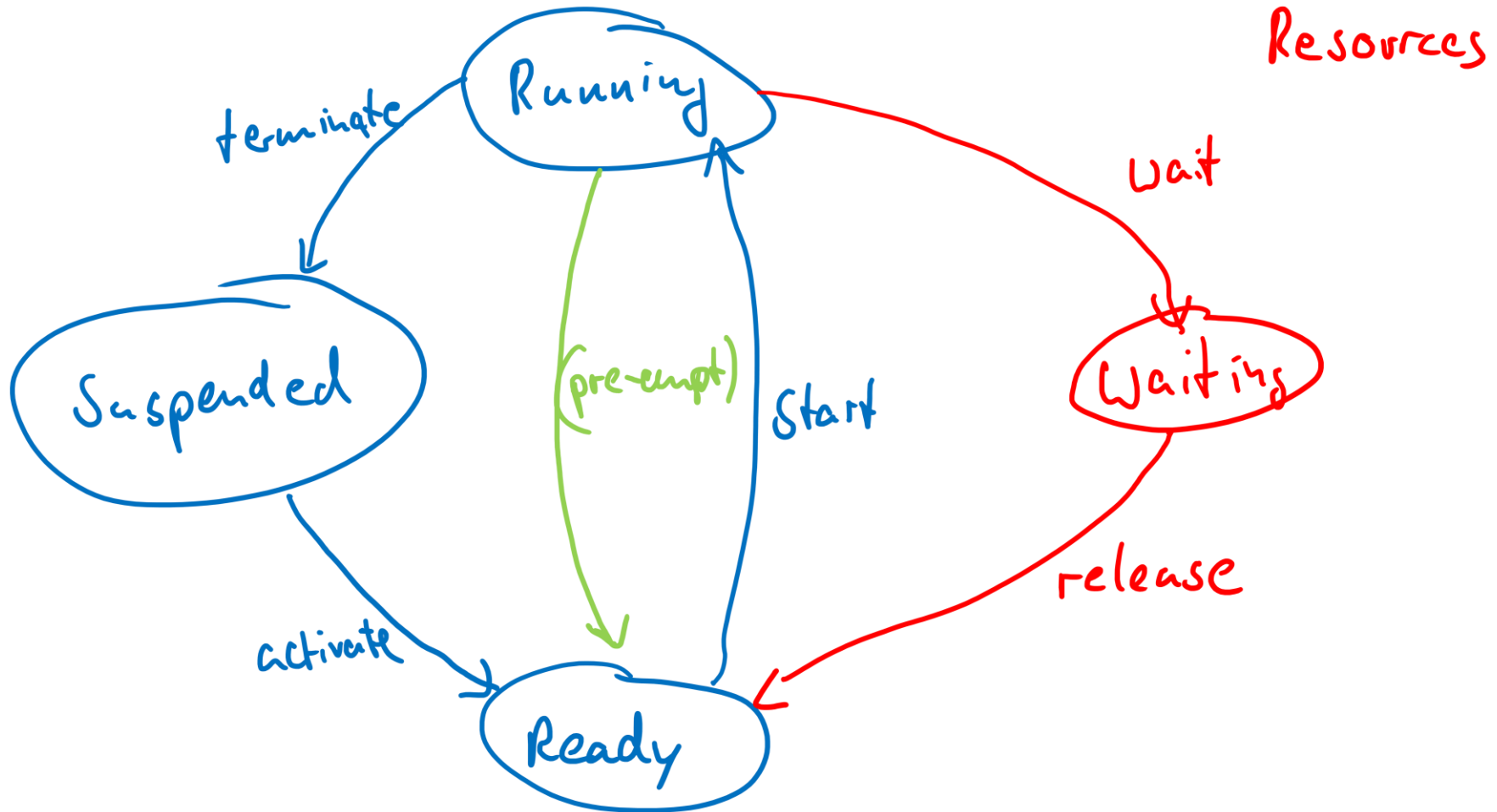


Task 2

basic

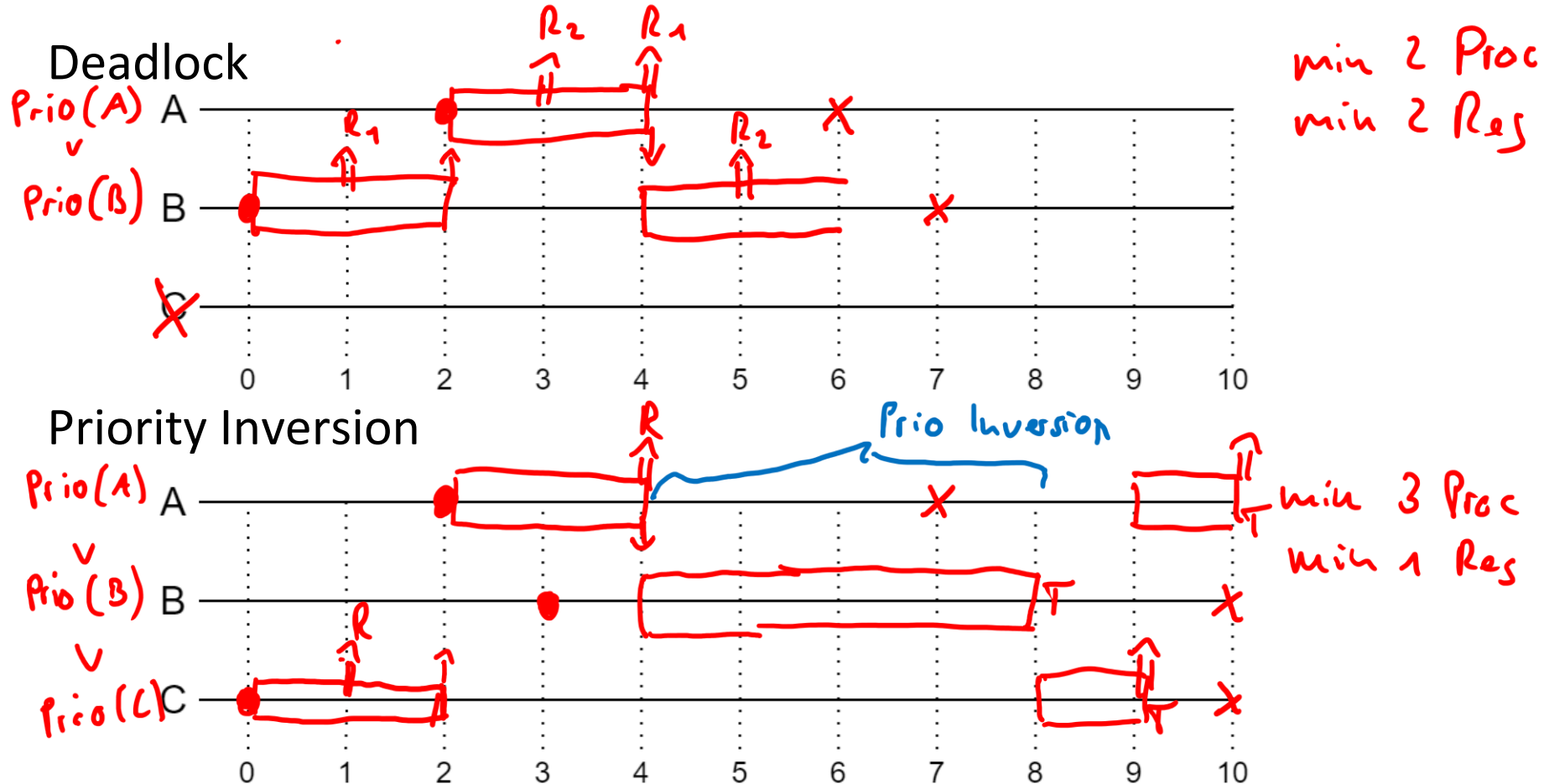
(with pre-emption)

a) Sketch the extended OSEK task model.



Task 2

b) How many **processes** and **resources** are needed for:



Real Time and Resources

Task 3



Task 3

Schedule these tasks (sorted by priority; first task has highest priority)

Task	Execution	Pattern	Deadline (absolute)
Prio(1) Task A	Start @ 5	runs 1 Req. runs 1 T	DL @ 10
⋮ Task B	Start @ 3	runs 1 T	DL @ 5
✓ Task C	Start @ 5	runs 3 T	DL @ 13
Prio(0) Task D	Start @ 1	runs 3 Req. runs 3 T	DL @ 13

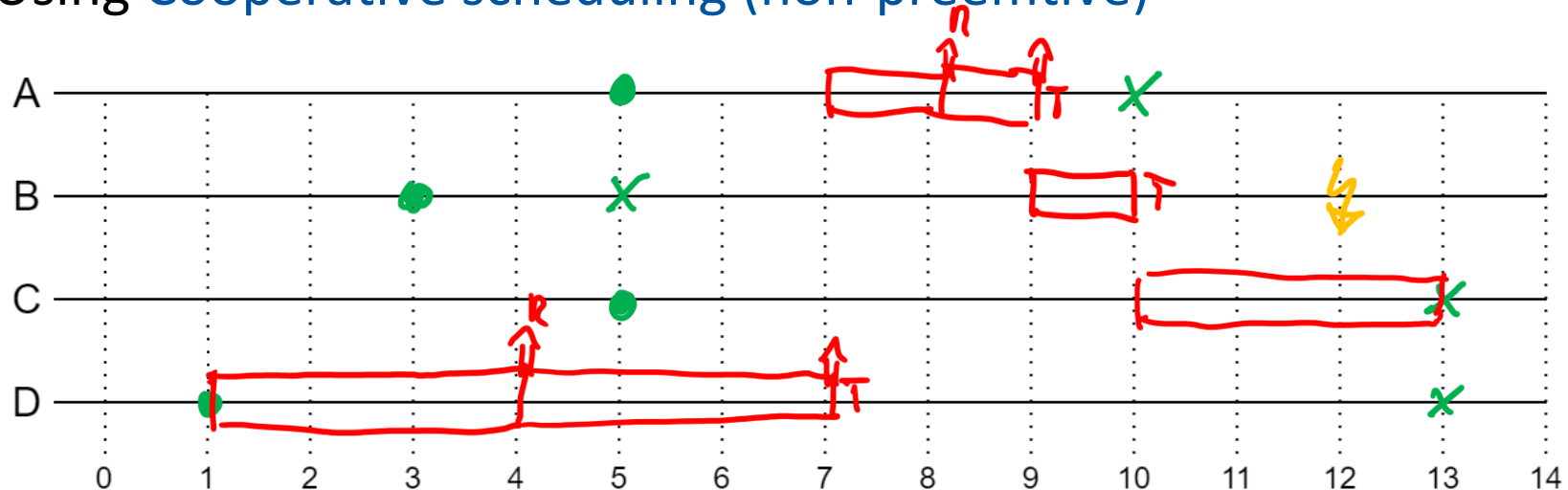
Activate

Task 3

Task	Execution	Deadline (absolute)
Task A	Start @ 5 runs 1 Req. runs 1 T	DL @ 10
Task B	Start @ 3 runs 1 T	DL @ 5
Task C	Start @ 5 runs 3 T	DL @ 13
Task D	Start @ 1 runs 3 Req. runs 3 T	DL @ 13

Activate

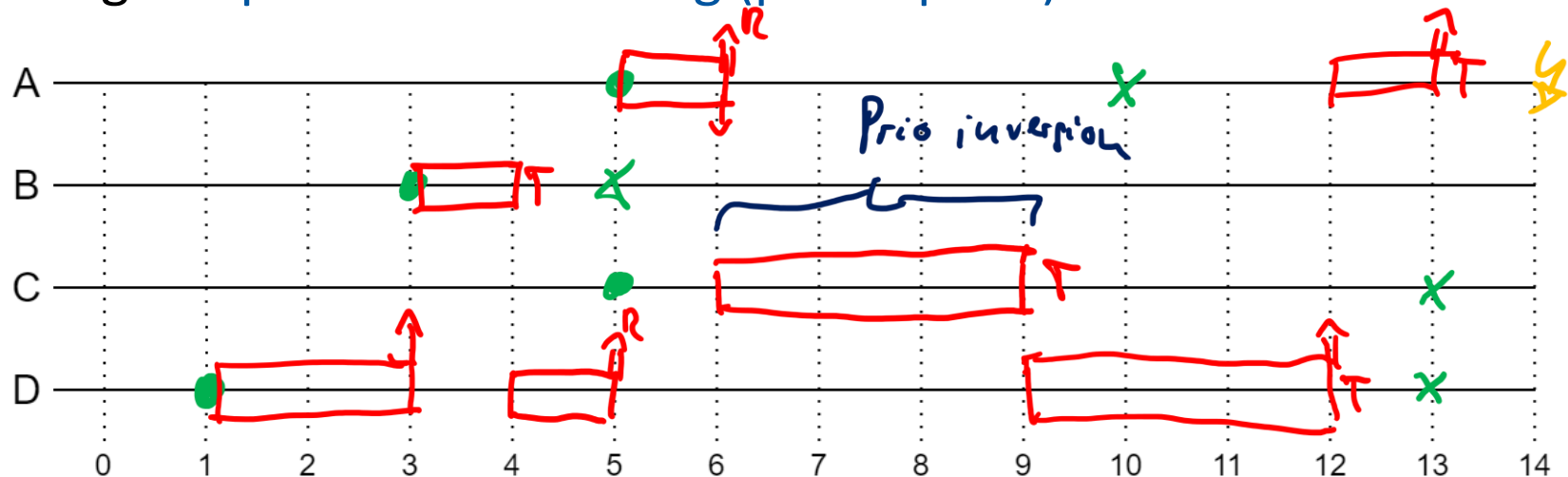
Using Cooperative scheduling (non-preemptive)



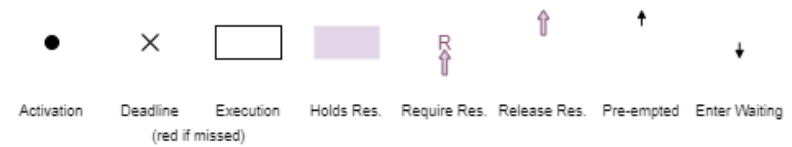
Task 3

Task	Execution	Deadline (absolute)
Task A	Start @ 5 runs 1 Req. runs 1 T	DL @ 10
Task B	Start @ 3 runs 1 T	DL @ 5
Task C	Start @ 5 runs 3 T	DL @ 13
Task D	Start @ 1 runs 3 Req. runs 3 T	DL @ 13

Using ~~Cooperative~~ ^{preemptive} scheduling (preemptive)



Task 3

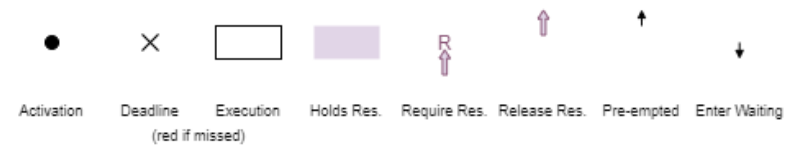


Task	Execution	Deadline (absolute)
Task A	Start @ 5 runs 1 Req. runs 1 T	DL @ 10
Task B	Start @ 3 runs 1 T	DL @ 5
Task C	Start @ 5 runs 3 T	DL @ 13
Task D	Start @ 1 runs 3 Req. runs 3 T	DL @ 13

Using Priority inheritance protocol

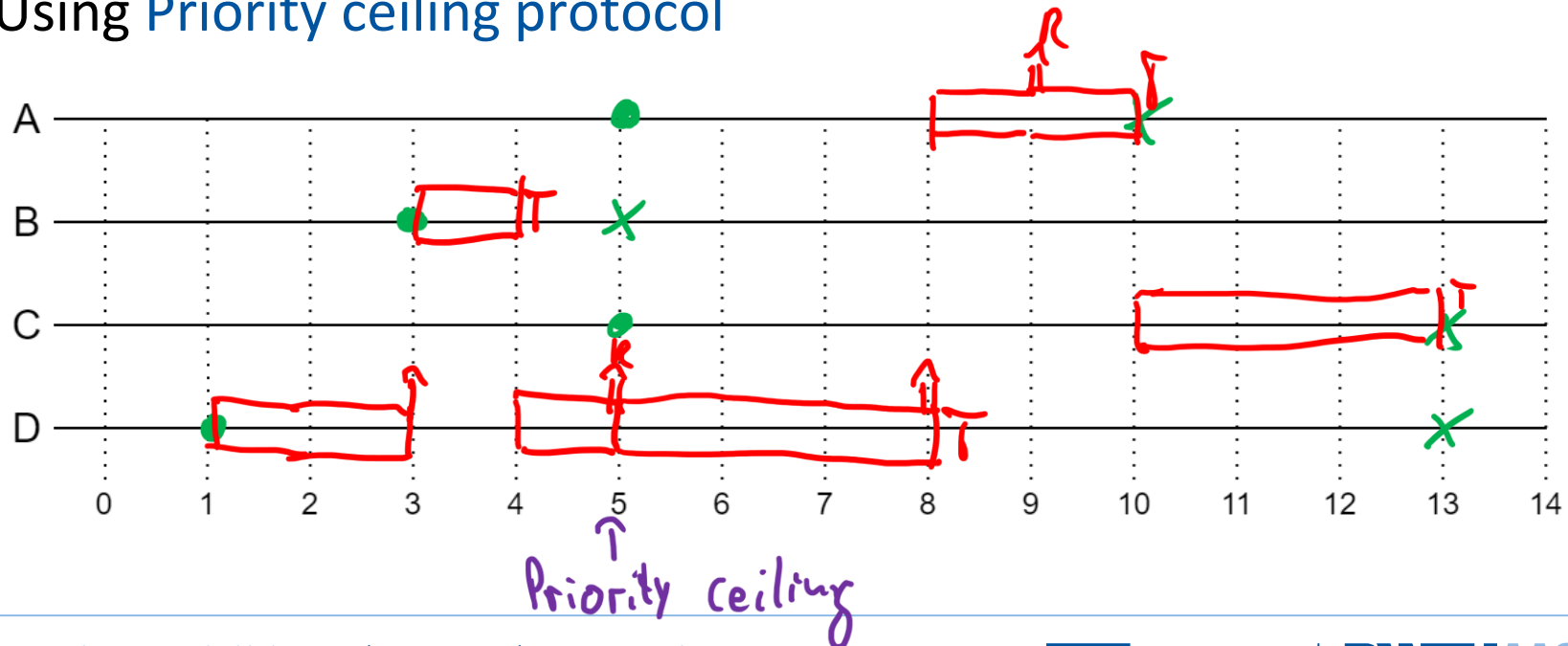


Task 3



Task	Execution	Deadline (absolute)
Task A	Start @ 5 runs 1 Req. runs 1 T	DL @ 10
Task B	Start @ 3 runs 1 T	DL @ 5
Task C	Start @ 5 runs 3 T	DL @ 13
Task D	Start @ 1 runs 3 Req. runs 3 T	DL @ 13

Using Priority ceiling protocol



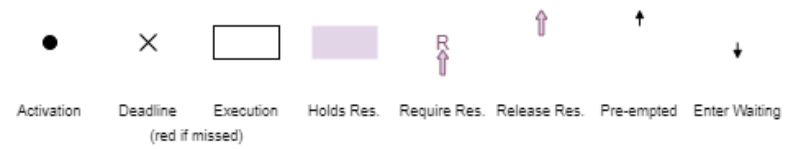
Periodic Scheduling

Task 4



Task 4

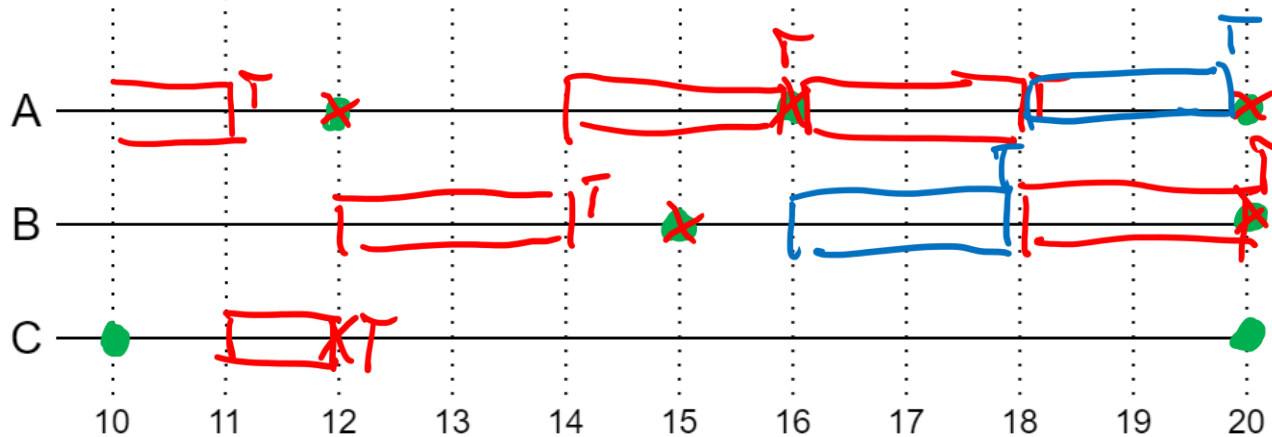
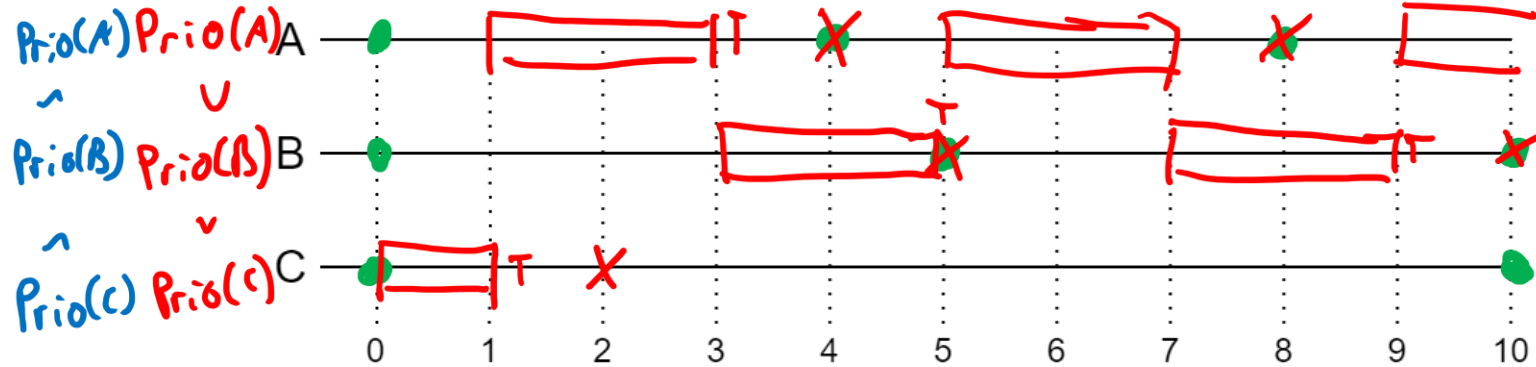
$$P_i := (\tilde{\tau}_i, T_i, DL_i) \quad \text{EDF}$$



a) Use earliest deadline first to schedule this task system.

A := (4, 2, 4), B := (5, 2, 5), C := (10, 1, 2)

all tasks activated @ 0



RMS
Rate
Monotonic
Scheduling
Rate = $\frac{1}{T_i}$

Task 4

b) Why is the following task system not schedulable?

A := (3, 2, 2), B := (6, 2, 7), C := (10, 3, 10)

$$\text{Utilization } (u) = \sum_{i=1}^n \frac{T_i}{\tilde{T}_i}$$

$$= \frac{2}{3} + \frac{2}{6} + \frac{3}{10} = 1,3$$

Handwritten annotations: A bracket under $\frac{2}{3} + \frac{2}{6}$ is labeled '1'. Below $\frac{3}{10}$ is '0,3'. To the right, a lightning bolt symbol is drawn over the inequality ≤ 1 .

not schedulable



necessary condition
for schedulability

\Rightarrow not sufficient!