

# CS 354 - Machine Organization & Programming

## Tuesday Sept 12th and Thursday Sept 14, 2023

**Project p1: DUE on or before Friday 9/22 (submit this week if possible)**

**Project p2A:** Released Friday and due on or before Friday 9/29

**Homework hw1:** Assigned soon

**Exam Conflicts (check entire semester):** Report by 9/29 to: <http://tiny.cc/cs354-conflicts>

**TA Lab Consulting & PM Activities** are scheduled. See links on course front page.

### Week 2 Learning Objectives (at a minimum be able to)

- ◆ state and show in memory diagrams the name, value, type, address, size of variable
- ◆ understand and show binary representation and byte ordering for int, char, address, values
- ◆ declare, assign, and dereference pointer “address” variables
- ◆ code, describe, and diagram 1D arrays on stack and on heap
- ◆ understand and show byte representation of character array vs “C string” variables
- ◆ understand and use `<string.h>` library functions with string literals and “C string” variables

### This Week

Tuesday	Thursday
Finish COMPILE, RUN, DEBUG Recall Variables and Meet Pointers Practice Pointers Recall 1D Arrays 1D Arrays and Pointers	Passing Addresses 1D Arrays on the Heap Pointer Caveats Meet C Strings Meet <code>string.h</code>
Read before Thursday K&R Ch. 7.8.5: Storage Management (malloc and calloc) K&R Ch. 5.5: Character Pointers and Functions K&R Ch. 5.6: Pointer Arrays; Pointers to Pointers	

### Next Week

**Topic:** 2D Arrays and Pointers

**Read:**

- K&R Ch. 5.7: Multi-dimensional Arrays
- K&R Ch. 5.8: Initialization of Pointer Arrays
- K&R Ch. 5.9: Pointers vs. Multi-dimensional Arrays
- K&R Ch. 5.10: Command-line Arguments

**Do:** Finish project p1 and start p2A

## Recall Variables

**What?** A scalar variable is primitive a unit of storage whose contents can change

→ Draw a basic memory diagram for the variable in the following code:

```
void someFunction() {  
    int i = 44;  
}
```

*int i* | 44

### Aspects of a Variable

identifier: name

value:

data stored,

type:

representation of bit pattern.

address:

starting location

size:

num of bytes

\* A scalar variable used as a source operand

e.g., `printf("%i\n", i);`

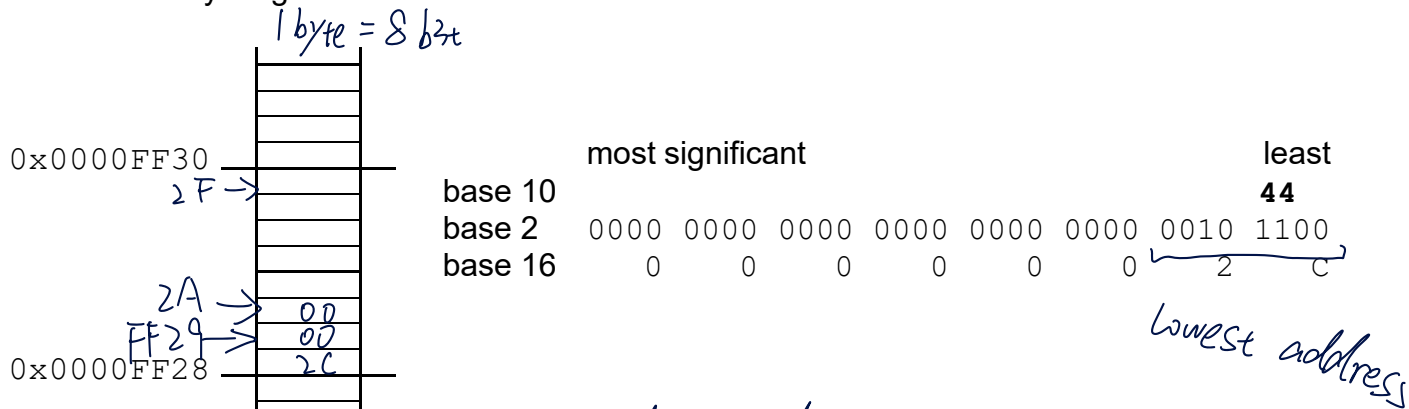
*read value*

\* A scalar variable used as a destination operand

e.g., `i = 11;`

### Linear Memory Diagram

A linear memory diagram is



byte addressability: each address identified 1 byte

endianess: byte ordering for variables with more 1 byte

little endian: IA-32: least significant byte in lowest address

big endian: most byte in lowest

# Meet Pointers

**What?** A pointer variable is

- ♦ a scalar variable whose value is address
- ♦

**Why?**

- ♦ for indirect access to memory.
- ♦  $\downarrow \quad \downarrow \quad \downarrow$  to function.
- ♦ Common in library.
- ♦ to access memory map hardware.

**How?**

→ Consider the following code:

```
void someFunction() {  
    int i = 44;
```

```
    int *ptr = NULL;  
    point to what type
```

$ptr = \&i;$

Basic Diag.

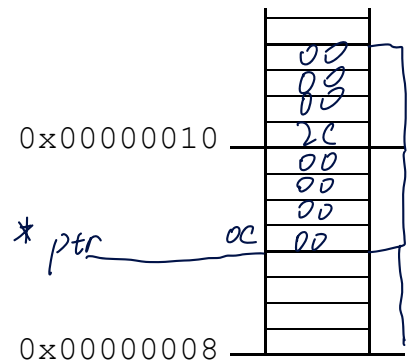
44

i

\*

ptr

Linear Diag.



→ What is ptr's initial value?  $0 \times 0$  / null address?  $0 \times \dots DC$  type?  $\text{int}^*$  size? 4 bytes

pointer: contain address to point

pointee: what is pointed to.

& address of operator:  $\&i$   $\&ptr \rightarrow$  get address of

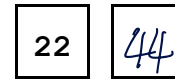
\* dereferencing operator:  $*ptr$ .  
↑  
only for pointer

## Practice Pointers

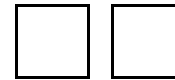
→ Complete the following diagrams and code so that they all correspond to each other:

```
void someFunction() {
    int i = 22;
    int j = 44;
    int *p1 = &j;
    int *p2; //at addr 0xFC0100EC
```

Basic Diag:



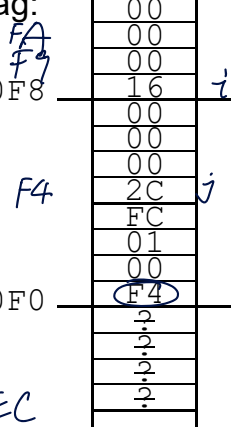
i j



p1 p2

Linear Diag:

0xFC0100F8



0xFC0100F0

→ What is p1's value?

→ Write the code to display p1's pointee's value.

→ Write the code to display p1's value.

→ Is it useful to know a pointer's exact value?

→ What is p2's value?

→ Write the code to initialize p2 so that it points to nothing.

→ What happens if the code below executes when p2 is NULL?

```
printf("%i\n", *p2);
```

→ What happens if the code below executes when p2 is uninitialized?

```
printf("%i\n", *p2);
```

→ Write the code to make p2 point to i.

→ How many pointer variables are declared in the code below?

```
void someFunction() {
    int* p1, p2;
```

→ What does the code below do?

```
int **q = &p1;
```

## Recall 1D Arrays

**What?** An array is

- a compound unit of storage, elem of same type.
- access via identifier and index
- allocate as contiguous fixed size block of memory

**Why?**

- store collection of data of same type fast.
- easier to deal than individual item for each item

**How?**

```
void someFunction(){  
    int a[5];
```

// stack allocate array

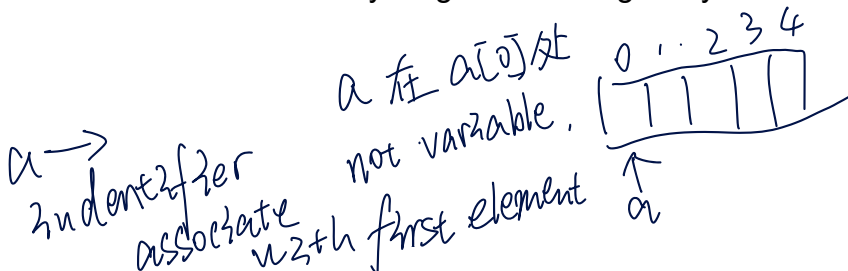
→ How many integer elements have been allocated memory? 5

→ Where in memory was the array allocation made? stack

→ Write the code that gives the element at index 1 a value of 11.

`a[1] = 11;`

→ Draw a basic memory diagram showing array a.



\* In C, the identifier for a stack allocated array (SAA) not variable.

\* A SAA identifier used as a source operand provide array address

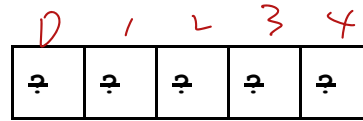
e.g., `printf("%p\n", a);`

\* A SAA identifier used as a destination operand result compile error.

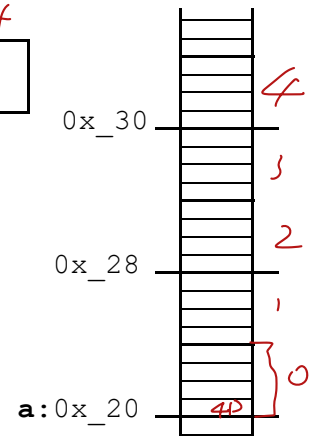
# 1D Arrays and Pointers

Given:

```
void someFunction() {
    int a[5]; //SAA
```



a  
 {int 4.  
 char 1  
 double 8



## Address Arithmetic

\*  $a[i] = *(a + i)$   
 (with 'int' written below 'a' in the original image)

1. compute the address

start at a's beginning 0x20.  
 add bytes offset to get to index i.

compile auto scale by 4  
 2. dereference the computed address to access the element

→ Write address arithmetic code to give the element at index 3 a value of 33.

$*(a + 3) = 33$

→ Write address arithmetic code equivalent to  $a[0] = 77$ ;

$*(a + 0) = 77$

## Using a Pointer

→ Write the code to create a pointer p having the address of array a above.

$\text{int } *p = a$ ; // &a

→ Write the code that uses p to give the element in a at index 4 a value of 44.

$*(p + 4) = 44$

\* In C, pointers and arrays

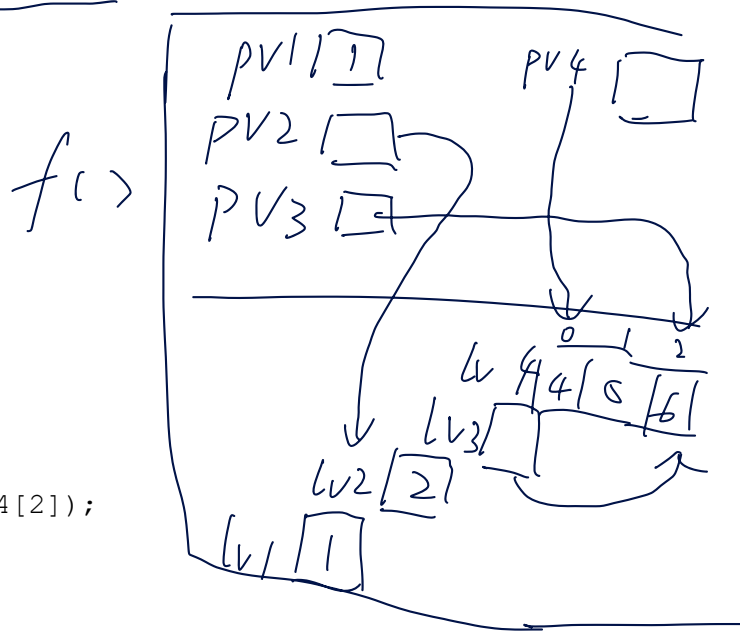
## Passing Addresses

### Recall Call Stack Tracing:

- ♦ manually trace for func call
- ♦ each func get a box (stack frame)
- ♦ top box is c r f ?

➤ What is output by the code below?

```
void f(int pv1, int *pv2, int *pv3, int pv4[]) {  
    int lv = pv1 + *pv2 + *pv3 + pv4[0];  
    pv1 = 11;  
    *pv2 = 22;  
    *pv3 = 33;  
    pv4[0] = lv;  
    pv4[1] = 44;  
}  
  
int main(void) {  
    int lv1 = 1, lv2 = 2;  
    int *lv3;  
    int lv4[] = {4, 5, 6};  
    lv3 = lv4 + 2;  
    > f(lv1, &lv2, lv3, lv4);  
    printf("%i,%i,%i\n", lv1, lv2, *lv3);  
    printf("%i,%i,%i\n", lv4[0], lv4[1], lv4[2]);  
    return 0;  
}
```



### Pass-by-Value

- ♦ scalars: param is a scalar variable that gets a copy of its scalar argument

- ♦ pointers: param is a ptr variable that get copy of addr

- ♦ arrays: param is a ptr variable that get copy of array address of [0]

\* Changing a callee's parameter

change callee's arg only

\* Passing an address

require caller trust callee.

∴ callee can change ptr points to

# 1D Arrays on the Heap

**What?** Two key memory segments used by a program are the

STACK

static (fixed in size) allocations

allocation size known during compile time

and HEAP

dynamic allocate.

**Why?** Heap memory enables

- ♦ access more memory than available at compile time
- ♦ having blocks of mem allocate and free when its running.

**How?**

func `void* malloc(size_in_bytes)`

reserve a block of heap memory of specific size.

func `void free(void* ptr)`

return a generic ptr that can assign to any ptr.  
free heap block ptr point to.

operate `sizeof(operand)` return size

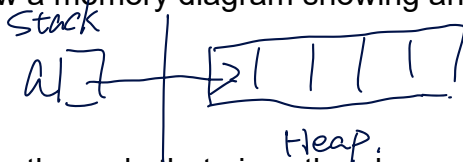
→ For IA-32 (x86), what value is returned by `sizeof(double)`? `sizeof(char)`? `sizeof(int)`?  
8 1 4

→ Write the code to dynamically allocate an integer array named `a` having 5 elements.

```
void someFunction() {
```

```
    int *a = malloc ( sizeof (int) x 5 );
```

→ Draw a memory diagram showing array `a`.



→ Write the code that gives the element at indexes 0, 1 and 2 a values of 0, 11 and 22 by using pointer dereferencing, indexing, and address arithmetic respectively.

```
*a = 0 // dereference
```

```
a[1] = 11 // indexing
```

```
*(a+2) = 22 // address arithmetic
```

→ Write the code that uses a pointer named `p` to give the element at index 3 a value of 33.

```
int *p = a;
```

```
*(p+3) = 33
```

→ Write the code that frees array `a`'s heap memory.

```
free ( a );
```

```
a = NULL
```

```
p = NULL.
```

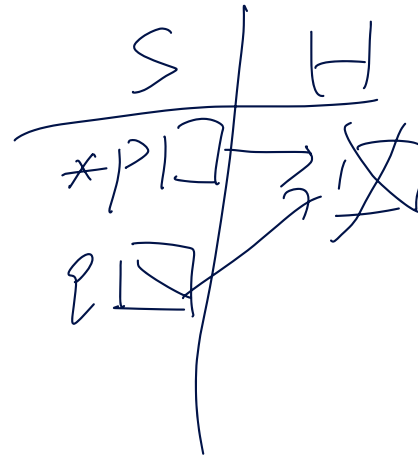


## Pointer Caveats

### \* Don't dereference uninitialized or NULL pointers!

```
int *p;
*p = 11; // intermittent error
```

```
int *q = NULL;
*q = 11;
```



### \* Don't dereference freed pointers!

```
int *p = malloc(sizeof(int));
int *q = p;
...
free(p);
```

```
...
*q = 11; // inter... error
```

dangling  
ptr

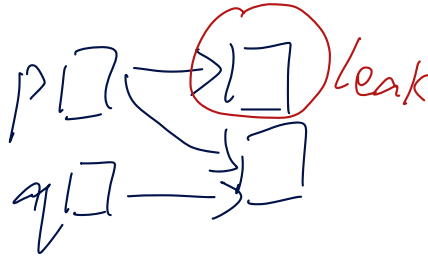
### dangling pointer:

a ptr var with a adress to heap memory that has been free.

### \* Watch out for heap memory leaks!

#### memory leak:

```
int *p = malloc(sizeof(int));
int *q = malloc(sizeof(int));
...
p = q;
```



### \* Be careful with testing for equality!

assume p and q are pointers

$p = q$

compares nothing because it's assignment

$p == q$

compares values in pointers

$*p == *q$

compares values in pointees

Compare int

### \* Don't return addresses of local variables!

```
int *ex1() {
    int i = 11; // local var
    return &i; // memory x available after function call end.
}
```

```
int *ex2(int size) {
    int a[size]; // SAA
    return a;
}
```

if in heap: malloc.  
// not able to return a address at stack.

## Meet C Strings

What? A string is

- a sequence of char terminated with null char
- a 1D array of char with string length + 1 '\0'

What? A string literal is

'CS 354'

- a construct space code string
- allocated prior to execution



size of C strings.

(Read only)

\* In most cases, a string literal used as a source operand

provide static addr

How? Initialization

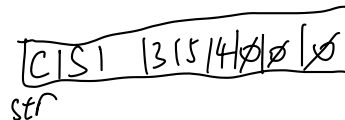
```
void someFunction() {  
    char *sptr = "CS 354";
```



→ Draw the memory diagram for sptr.

→ Draw the memory diagram for str below.

```
char str[9] = "CS 354";
```



→ During execution, where is str allocated?

SAA

How? Assignment

→ Given str and sptr declared in somefunction above, what happens with the following code?

```
sptr = "mumpsimus"; // OK
```

```
str = "folderol"; // Compiler Error
```

\* Caveat: Assignment cannot be used

## Meet `string.h`

What? `string.h` is

```
int strlen(const char *str)
```

Returns the length of string `str` up to but *not* including the null character.

```
int strcmp(const char *str1, const char *str2)
```

Compares the string pointed to by `str1` to the string pointed to by `str2`.

returns: < 0 (a negative) if `str1` comes before `str2`  
0 if `str1` is the same as `str2`  
> 0 (a positive) if `str1` comes after `str2`

```
char *strcpy(char *dest, const char *src)
```

Copies the string pointed to by `src` to the memory pointed to by `dest` and terminates with the null character.

```
char *strcat(char *dest, const char *src)
```

Appends the string pointed to by `src` to the end of the string pointed to by `dest` and terminates with the null character.

\* *Ensure the destination character array large enough including null char*

buffer overflow: exceed bounds of array.

How? `strcpy`

→ Given `str` and `sptr` as declared in some function on the previous page, what happens with the following code?

```
strcpy(str, "folderol");
```

```
strcpy(str, "formication");
```

```
strcpy(sptr, "vomitory");
```

*SPTR = "v..." // assignment.*

\* *Rather than assignment, `strcpy` (or `strncpy`) must be used to copy "c" string from one array to another*

\* *Caveat: Beware of*

*Buffer overflow / write to code segment.*