Digital System Design with HDL (I) Lecture 9

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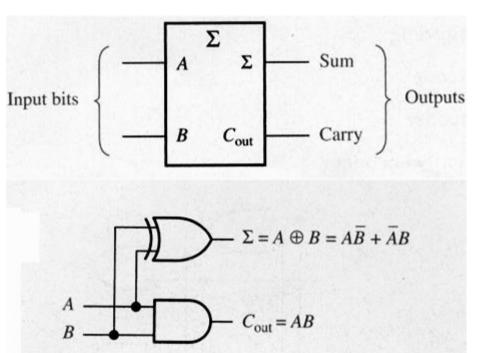
In This Session

- Verilog for Combinational Circuits
 - Adders
 - Arithmetic Overflow

Half-Adders

The half-adder does not add an input carry.

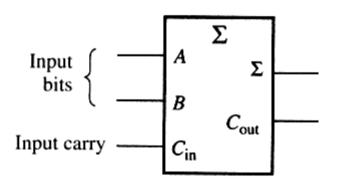
A	В	Cout	Σ
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0



- The sum is 1 only when the inputs are different

 — an XOR operation.
- The output carry is 1 only when both the inputs are 1 an AND operation.

A	В	C_{in}	Cout	Σ
0	0	0	0	0
0	0	1	0	1
0	· 1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1



$$\Sigma = (A \oplus B) \oplus C_{\rm in}$$

$$C_{out} = ABC_{in} + AB\overline{C}_{in} + AB\overline{C}_{in} + A\overline{B}C_{in}$$

$$= \left(ABC_{in} + AB\overline{C}_{in}\right) + \left(ABC_{in} + AB\overline{C}_{in}\right) + \left(ABC_{in} + \overline{A}BC_{in}\right)$$

$$= AB + AC_{in} + BC_{in}$$

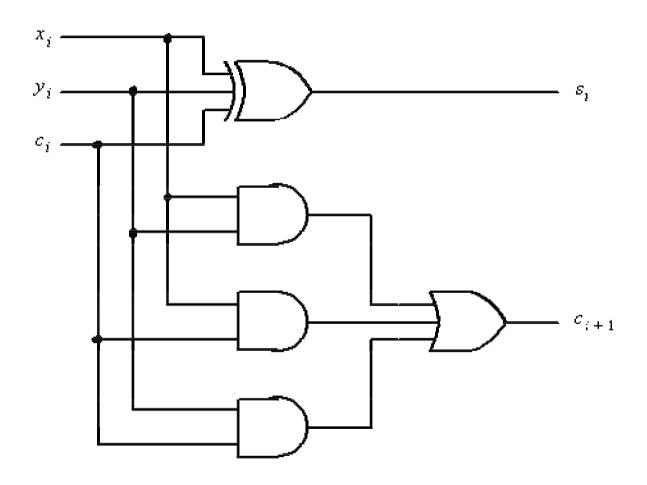
Exclusive-OR operation satisfies associative rule:

$$\begin{array}{rcl}
x \oplus (y \oplus z) & = & x \oplus (\overline{y}z + y\overline{z}) \\
& = & \overline{x}(\overline{y}z + y\overline{z}) + x(\overline{y} \cdot \overline{z} + yz) \\
& = & \overline{x} \cdot \overline{y}z + \overline{x}y\overline{z} + x\overline{y} \cdot \overline{z} + xyz
\end{array}$$

$$(x \oplus y) \oplus z = (\overline{x}y + x\overline{y}) \oplus z$$
$$= (\overline{x} \cdot \overline{y} + xy)z + (\overline{x}y + x\overline{y})\overline{z}$$
$$= \overline{x} \cdot \overline{y}z + xyz + \overline{x}y\overline{z} + x\overline{y} \cdot \overline{z}$$

We have

$$s = (x \oplus y) \oplus c = x \oplus y \oplus c$$



What is the output for an XOR gate with multiple inputs?

- If there are even 1s in the inputs, output a 0.
- If there are odd 1s in the input, output a 1.

Example:

 $0 \oplus 0 \oplus 0 = 0$

 $1 \oplus 0 \oplus 0 = 1$

 $1 \oplus 1 \oplus 0 = 0$

 $1 \oplus 1 \oplus 1 = 1$

Arithmetic Overflow

 Arithmetic overflow occurs when the outcome of addition does not fit within the magnitude bits used to represent the numbers.

Arithmetic Overflow

- It is important to be able to detect the overflow.
- In this example, Overflow = $c_3 \oplus c_4$
- For n-bit numbers we have

Overflow =
$$c_n \oplus c_{n-1} = c_n \oplus x_{n-1} \oplus y_{n-1} \oplus s_{n-1}$$

Since $s_k = x_k \oplus y_k \oplus c_k$, it follows that

$$x_k \oplus y_k \oplus s_k = (x_k \oplus y_k) \oplus (x_k \oplus y_k \oplus c_k)$$

$$= (x_k \oplus y_k) \oplus (x_k \oplus y_k) \oplus c_k$$

$$= 0 \oplus c_k$$

$$= c_k$$

Using continuous assignment

```
module fulladd (Cin, x, y, s, Cout);
input Cin, x, y;
output s, Cout;

assign s = x ^ y ^ Cin;
assign Cout = (x & y) | (x & Cin) | (y & Cin);
endmodule
```

A 4-bit adder using module instantiation

```
module adder4 (carryin, X, Y,S,carryout);
   input carryin;
   input [3:0] X, Y;
   output [3:0] S;
   output carryout;
   wire [3:1] C;
   fulladd stage0 (carryin, X[0], Y[0], S[0], C[1]);
   fulladd stage1 (C[1], X[1], Y[1], S[1], C[2]);
   fulladd stage2 (C[2], X[2], Y[2], S[2], C[3]);
   fulladd stage3 (C[3], X[3], Y[3], S[3], carryout);
endmodule
```

The **generate** Construct

- n instances of the fulladd subcircuit are required to build a n-bit ripple-carry adder.
- The **generate** construct (Verilog 2001) allows subcircuits to be instantiated in a loop.
- The loop index must be declared of type genvar a positive integer that appears only in generate blocks.

Syntax:

generate

[procedural statements]

[instantiation statements]

endgenerate

The **generate**Construct

 Each instance has an instance name addbit[i].stage.

```
module ripple_g (carryin, X, Y,S, carryout);
 parameter n = 4;
 input carryin;
 input [n-1:0] X, Y;
 output [n-1:0] S;
 output carryout;
 wire [n:0] C;
 genvar i;
 assign C[0] = carryin;
 assign carryout = C[n];
 generate
  for (i = 0; i \le n-1; i = i+1)
  begin:addbit
   fulladd stage (C[i], X[i], Y[i], S[i], C[i+1]);
  end
 endgenerate
endmodule
                                              13
```

An n-bit adder using procedural statements

```
module addern (carryin, X, Y, S, carryout);
  parameter n=32;
  input carryin;
  input [n-1:0] X, Y;
  output reg [n-1:0] S;
  output reg carryout;
  reg [n:0] C;
  integer k;
  always @(X, Y, carryin)
  begin
     C[0] = carryin;
     for (k = 0; k < n; k = k+1)
     begin
        S[k] = X[k] ^ Y[k] ^ C[k];
        C[k+1] = (X[k] \& Y[k]) | (X[k] \& C[k]) | (Y[k] \& C[k]);
     end
     carryout = C[n];
  end
endmodule
                                                              14
```

An alternative n-bit adder with the overflow signal

```
module addern (carryin, X, Y, S, carryout, overflow);
    parameter n = 32;
    input carryin;
    input [n-1:0] X, Y;
    output reg [n-1:0] S;
    output reg carryout, overflow;
    always @(X, Y, carryin)
    begin
        \{carryout, S\} = X + Y + carryin;
       overflow = carryout ^X[n-1] ^Y[n-1] ^S[n-1];
    end
endmodule
```

A one-digit BCD adder

```
module bcdadd(Cin, X, Y, S, Cout);
   input Cin;
  input [3:0] X, Y;
  output reg [3:0] S;
  output reg Cout;
  reg [4:0] Z;
  always@ (X, Y, Cin)
  begin
     Z = X + Y + Cin;
     if (Z < 10)
        {Cout, S} = Z;
     else
        \{Cout, S\} = Z + 6;
  end
endmodule
```