



CS 225

Data Structures

March 24– Graphs

G Carl Evans



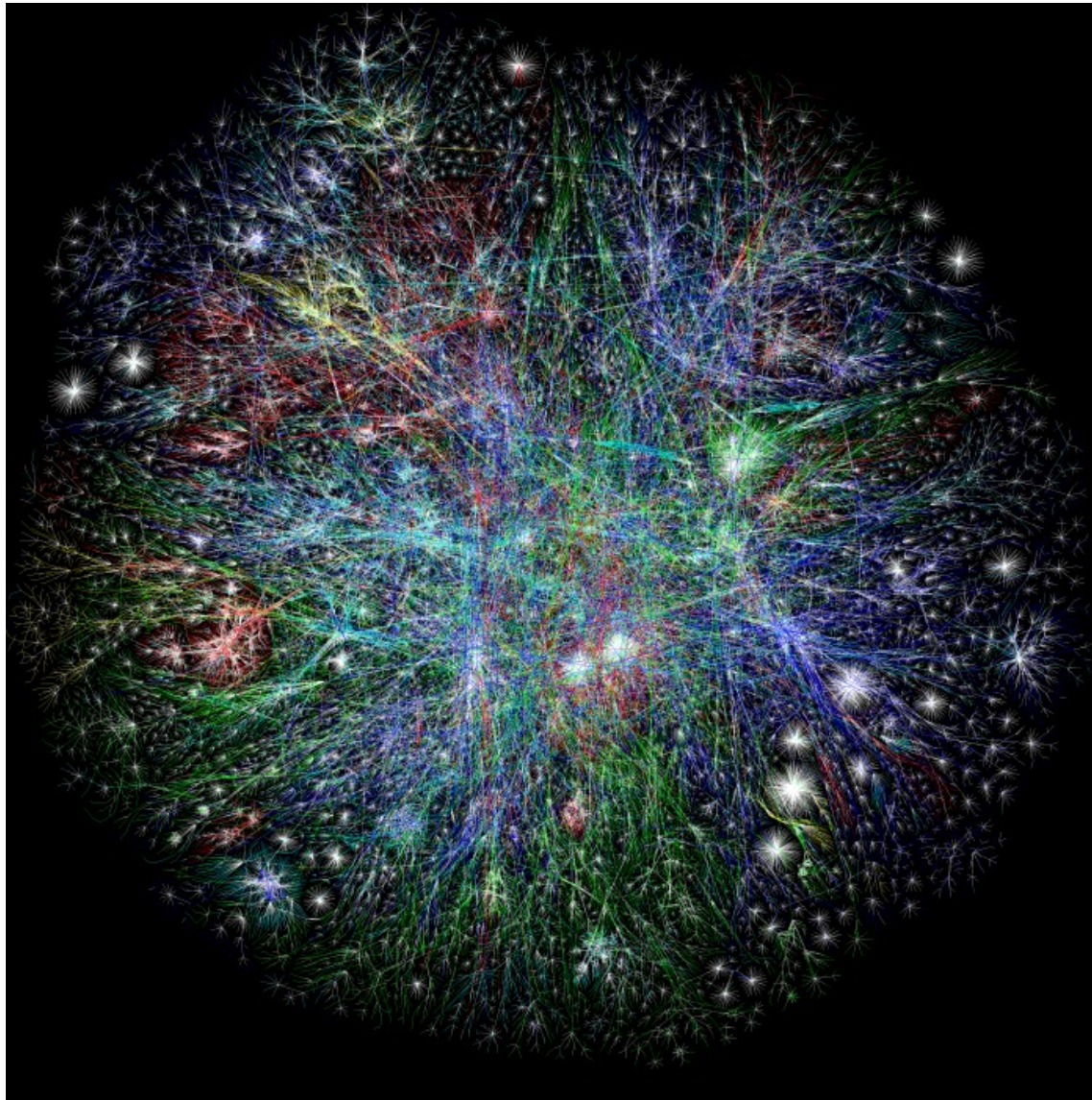
In Review: Data Structures

Array

- Sorted Array
- Unsorted Array
- Stacks
- Queues
- Priority Queues
 - Heaps
- Disjoint Sets
 - UpTrees

Linked

- Doubly Linked List
- Trees
 - BTree
 - Binary Tree
 - Huffman Encoding
 - kd-Tree
 - AVL Tree



The Internet 2003

The OPTE Project (2003)

Map of the entire internet; nodes are routers; edges are connections.

HeapifyUp BasicBlock Graph

```

heapifyUp(int*, unsigned int):
    push    rbp
    mov     rbp, rsp
    sub     rsp, 16
    mov     qword ptr [rbp - 8], rdi
    mov     dword ptr [rbp - 12], esi
    cmp     dword ptr [rbp - 12], 1
    jbe     .LBB0_4
    
```

```

heapifyUp(int*, unsigned int):@8
    mov     rax, qword ptr [rbp - 8]
    mov     ecx, dword ptr [rbp - 12]
    mov     edx, ecx
    mov     ecx, dword ptr [rax + 4*rdx]
    mov     rax, qword ptr [rbp - 8]
    mov     esi, dword ptr [rbp - 12]
    shr     esi, 1
    mov     esi, esi
    mov     edx, esi
    cmp     ecx, dword ptr [rax + 4*rdx]
    jge     .LBB0_3
    
```

```

heapifyUp(int*, unsigned int):@19
    mov     rax, qword ptr [rbp - 8]
    mov     ecx, dword ptr [rbp - 12]
    mov     edx, ecx
    mov     ecx, dword ptr [rax + 4*rdx]
    mov     dword ptr [rbp - 16], ecx
    mov     rax, qword ptr [rbp - 8]
    mov     ecx, dword ptr [rbp - 12]
    shr     ecx, 1
    mov     ecx, ecx
    mov     edx, ecx
    mov     ecx, dword ptr [rax + 4*rdx]
    mov     rax, qword ptr [rbp - 8]
    mov     esi, dword ptr [rbp - 12]
    mov     edx, esi
    mov     dword ptr [rax + 4*rdx], ecx
    mov     ecx, dword ptr [rbp - 16]
    mov     rax, qword ptr [rbp - 8]
    mov     esi, dword ptr [rbp - 12]
    shr     esi, 1
    mov     esi, esi
    mov     edx, esi
    mov     dword ptr [rax + 4*rdx], ecx
    mov     rdi, qword ptr [rbp - 8]
    mov     ecx, dword ptr [rbp - 12]
    shr     ecx, 1
    mov     esi, ecx
    call    heapifyUp(int*, unsigned int)
    
```

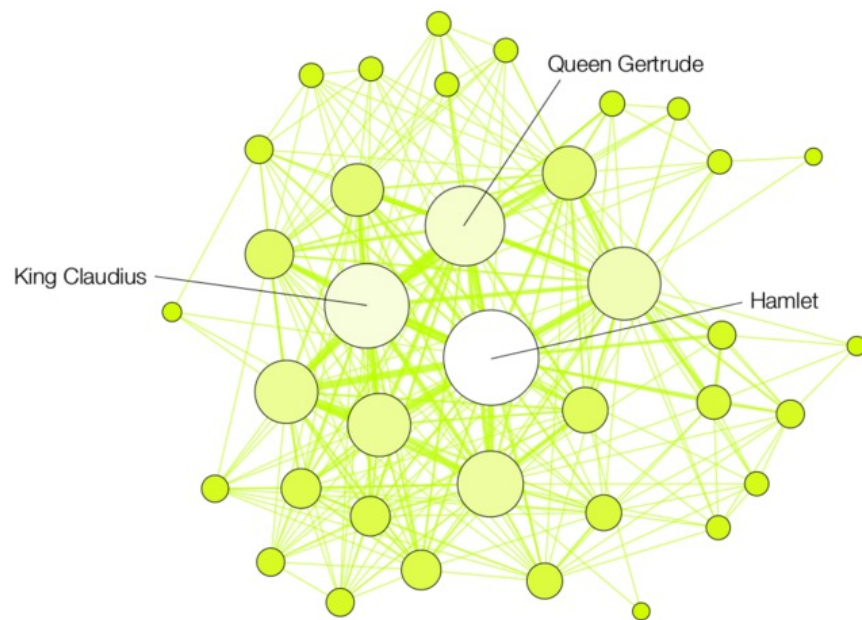
```

.LBB0_3:
    jmp     .LBB0_4
    
```

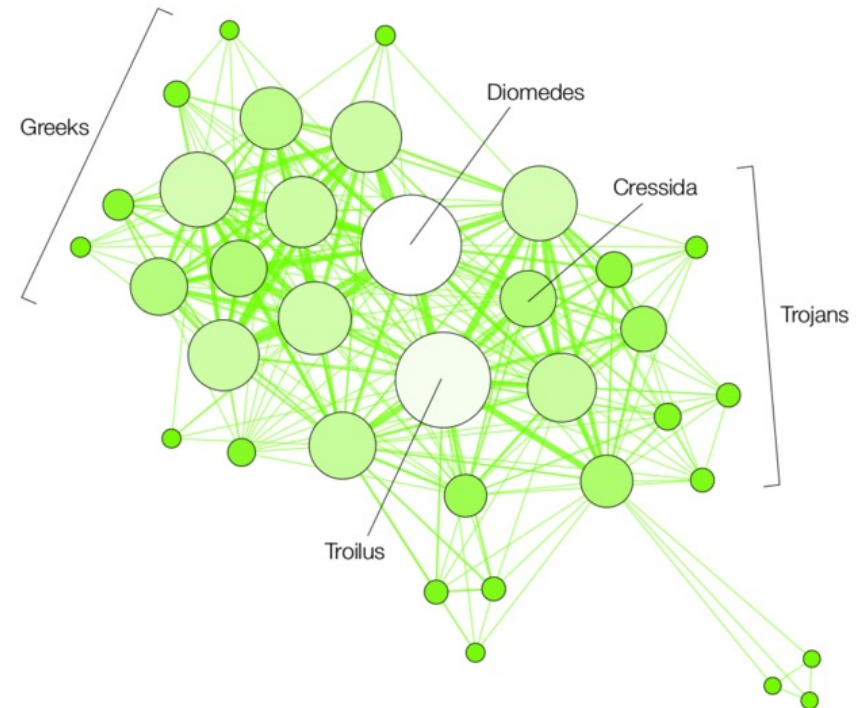
```

.LBB0_4:
    add     rsp, 16
    pop     rbp
    ret
    
```

Generated using tools at
<https://godbolt.org>



HAMLET

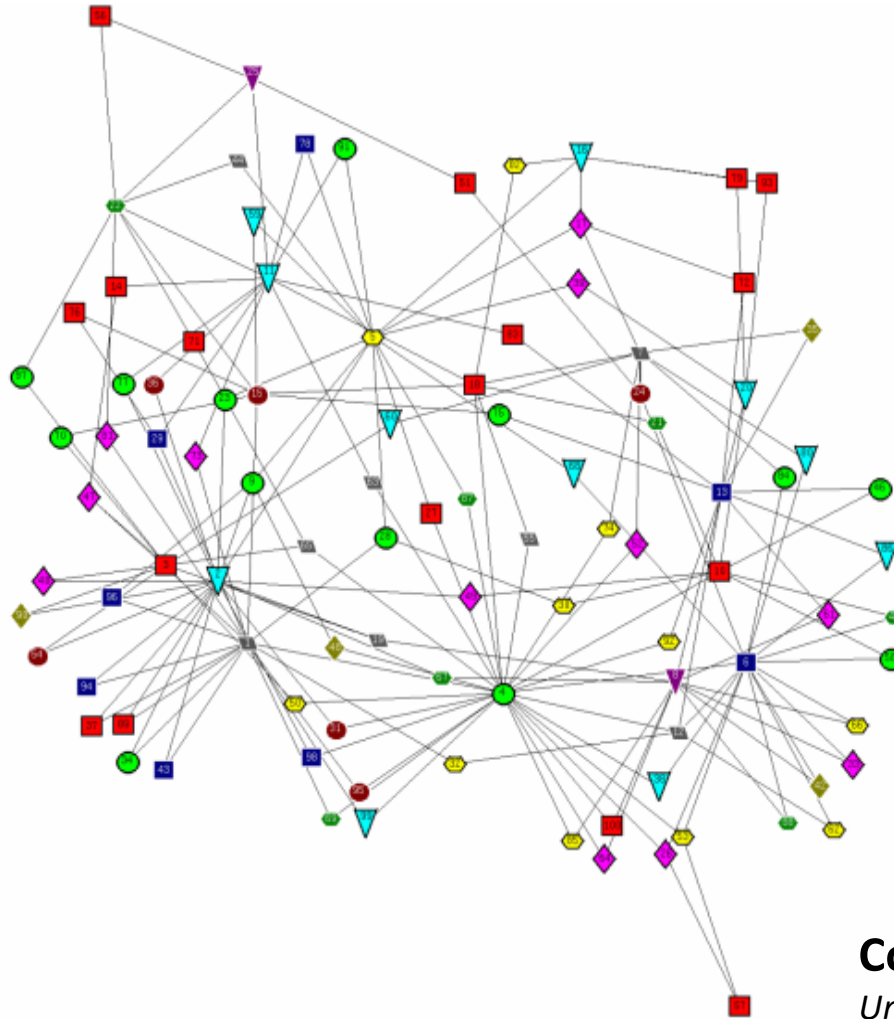


TROILUS AND CRESSIDA

Who's the real main character in Shakespearean tragedies?

Martin Grandjean (2016)

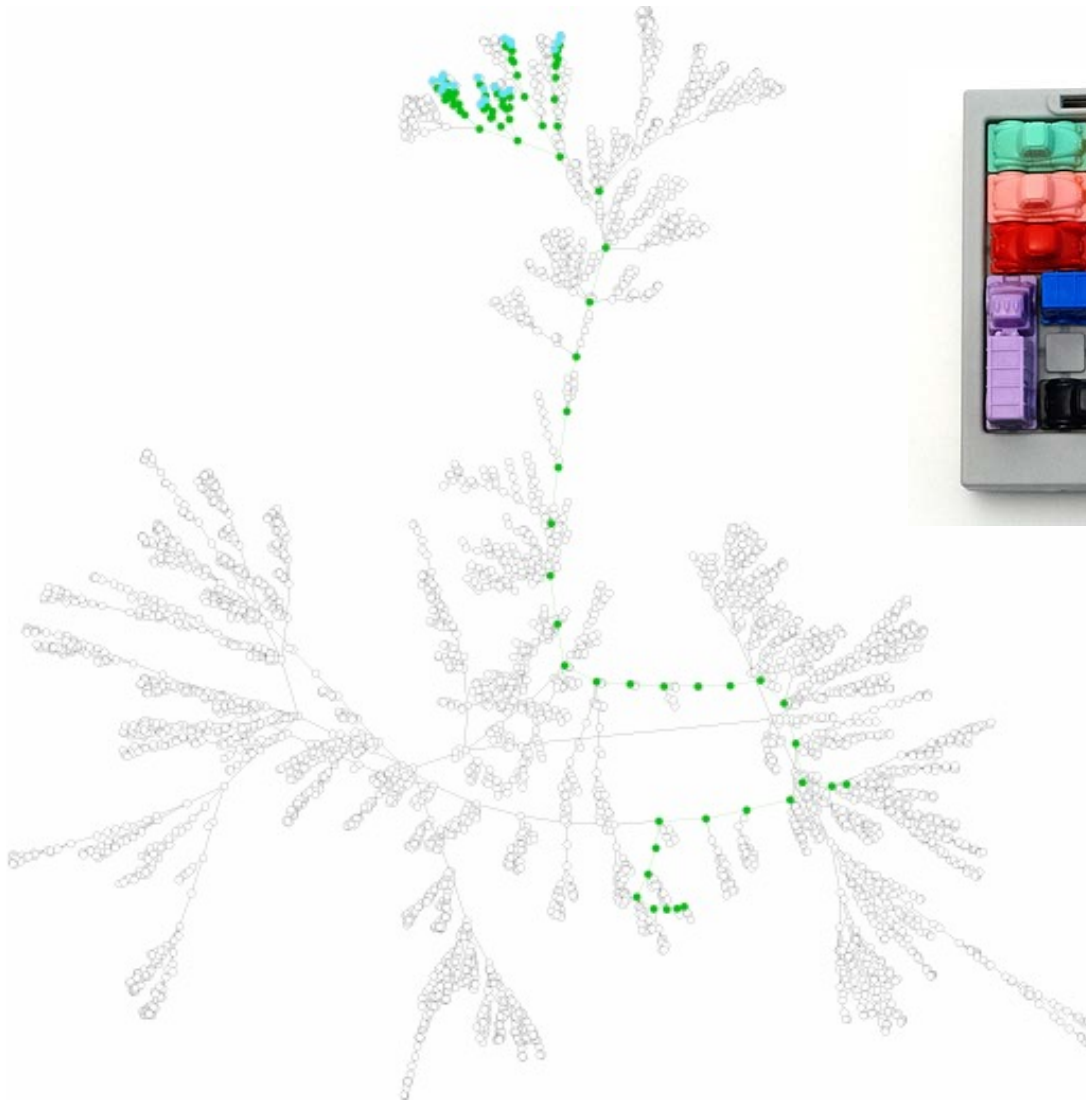
<https://www.pbs.org/newshour/arts/whos-the-real-main-character-in-shakespearean-tragedies-heres-what-the-data-say>



Conflict-Free Final Exam Scheduling Graph

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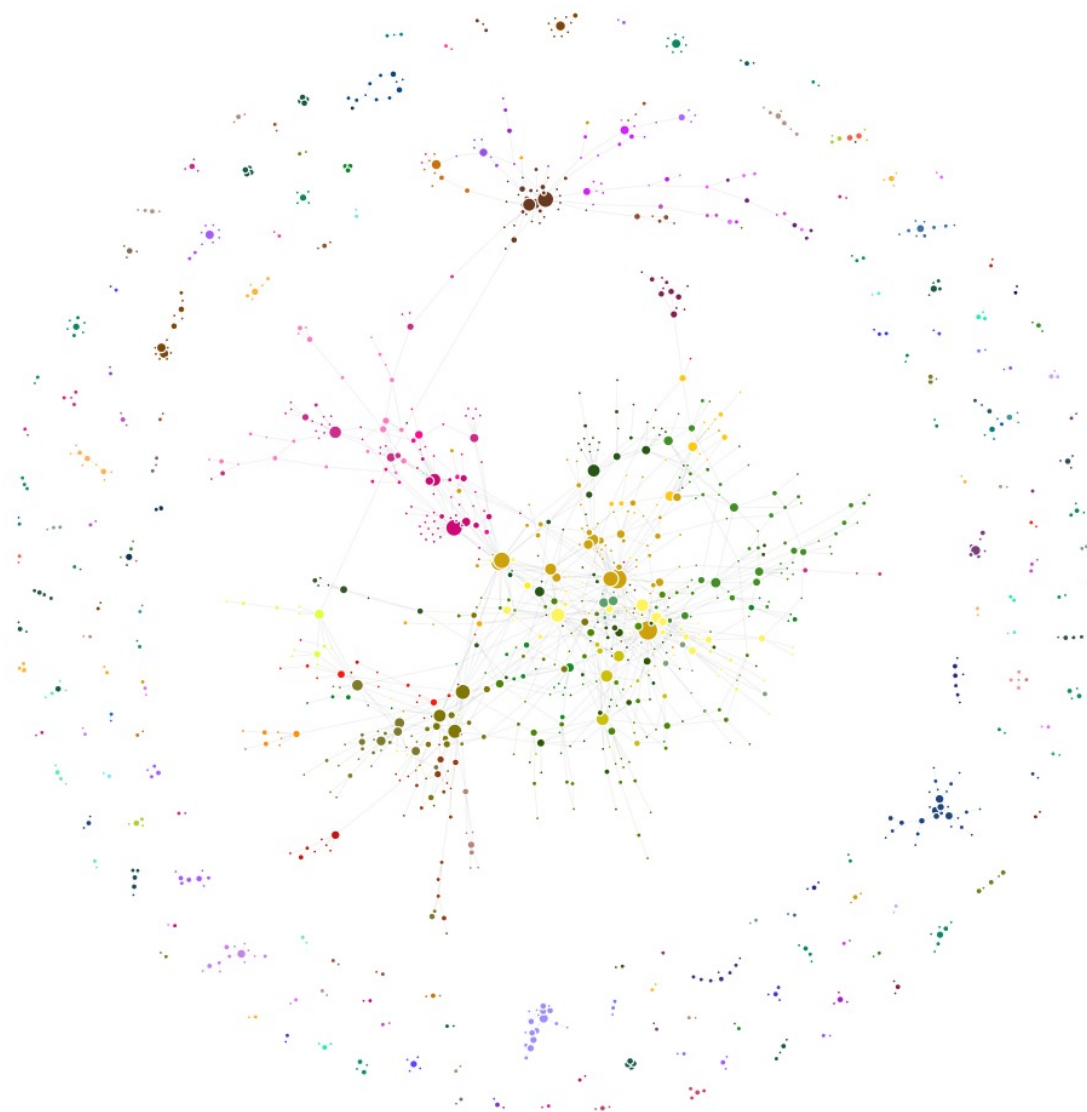
Presented by Cinda Heeren, 2016



“Rush Hour” Solution

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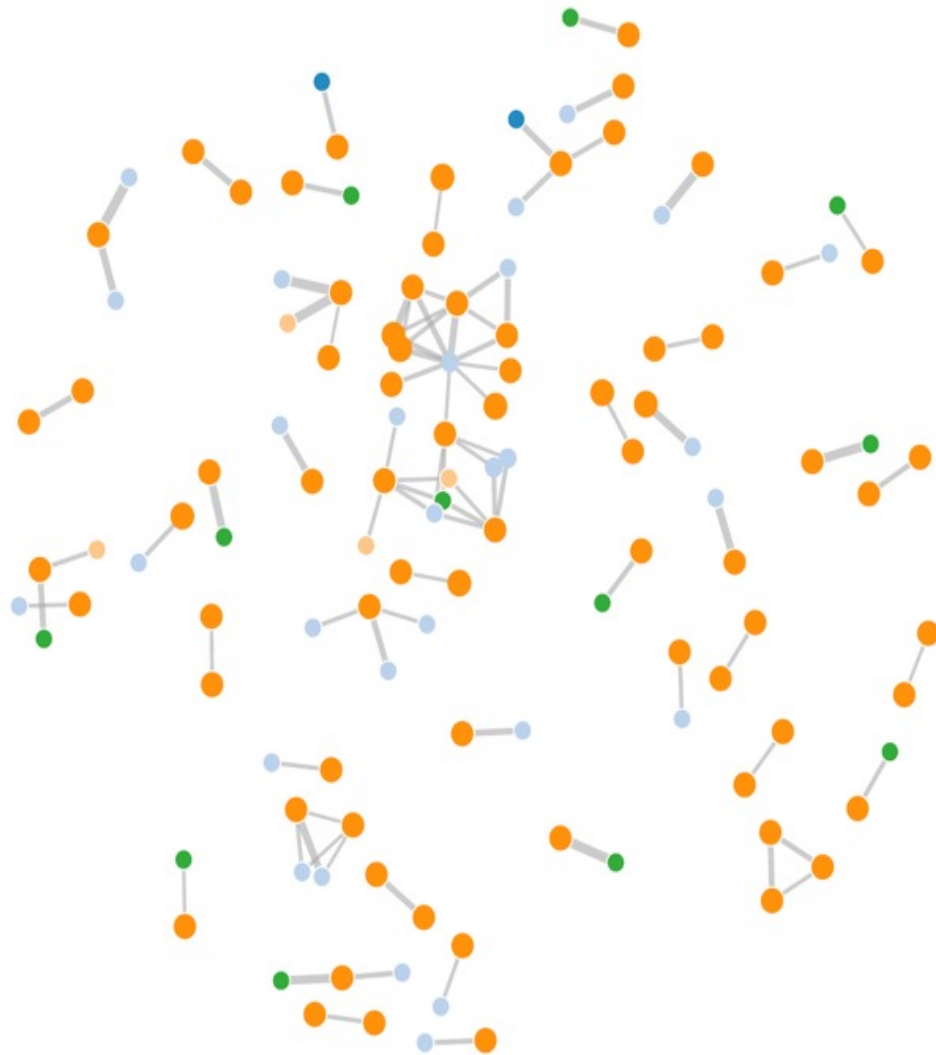


Class Hierarchy At University of Illinois Urbana-Champaign

A. Mori, W. Fagen-Ulmschneider, C. Heeren

Graph of every course at UIUC; nodes are courses, edges are prerequisites

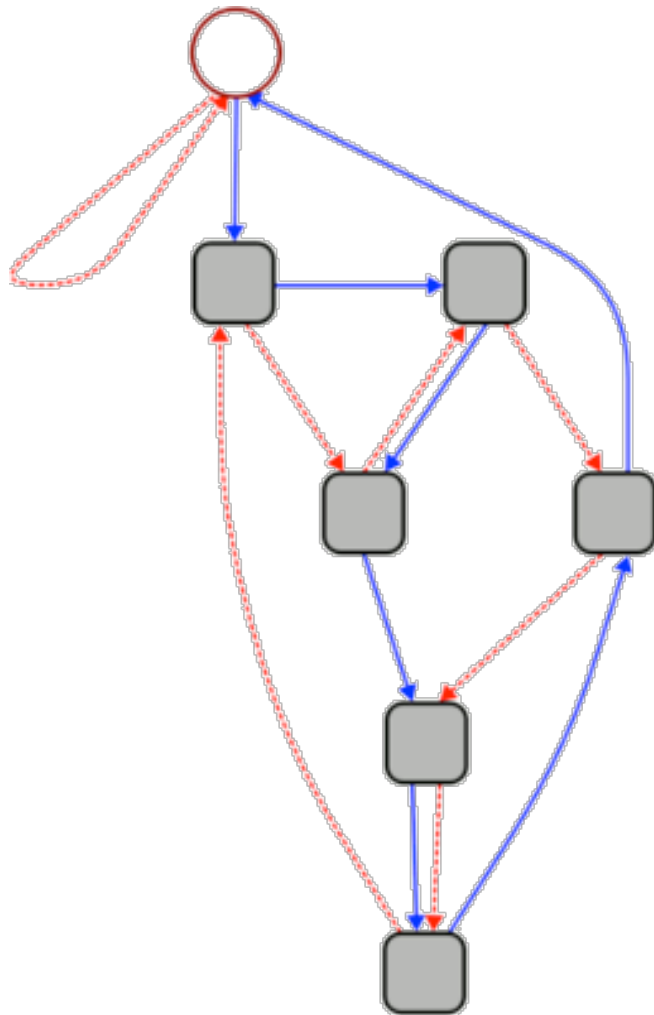
http://waf.cs.illinois.edu/discovery/class_hierarchy_at_illinois/



MP Collaborations in CS 225

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Presented by Cinda Heeren, 2016



This graph can be used to quickly calculate whether a given number is divisible by 7.

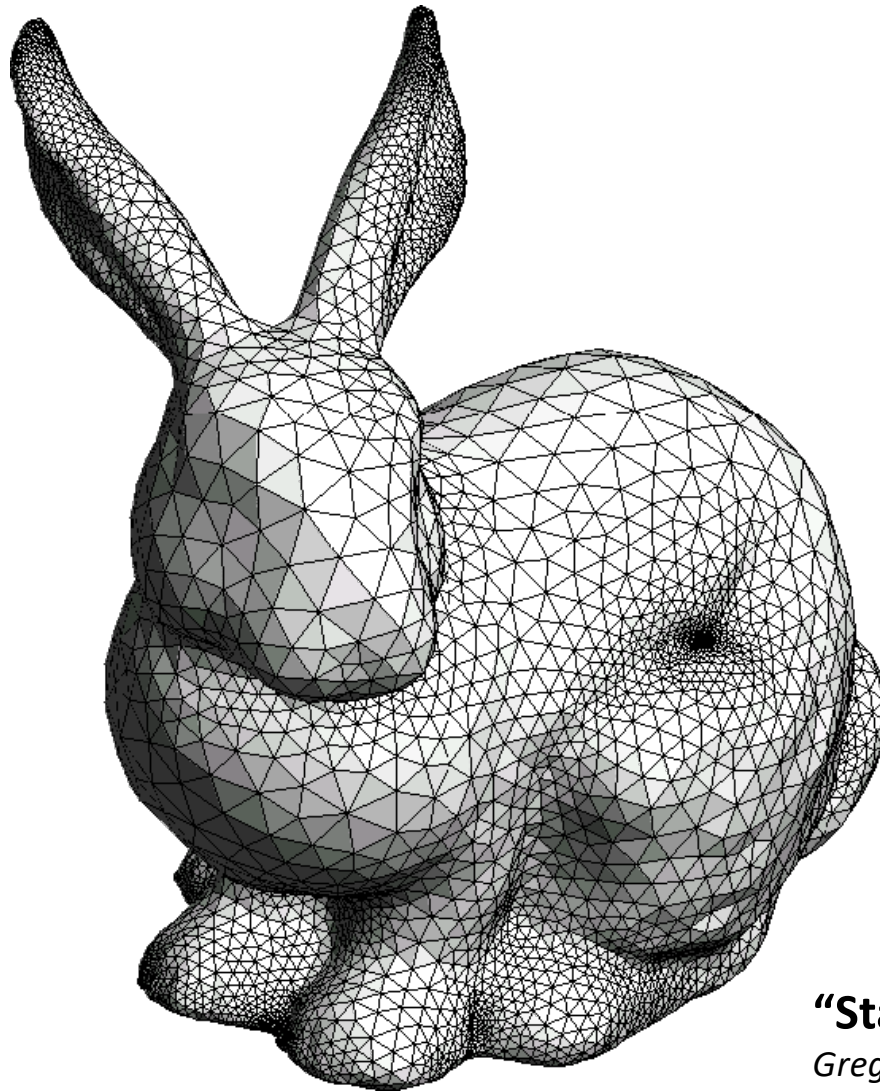
1. Start at the circle node at the top.
2. For each digit **d** in the given number, follow **d** blue (solid) edges in succession. As you move from one digit to the next, follow **1** red (dashed) edge.
3. If you end up back at the circle node, your number is divisible by 7.

3703

“Rule of 7”

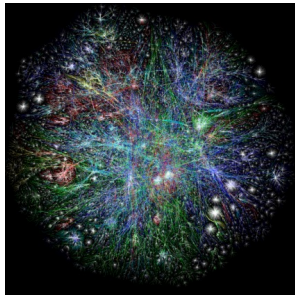
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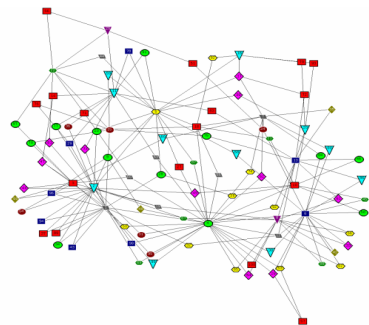
“Stanford Bunny”
Greg Turk and Mark Levoy (1994)

Graphs



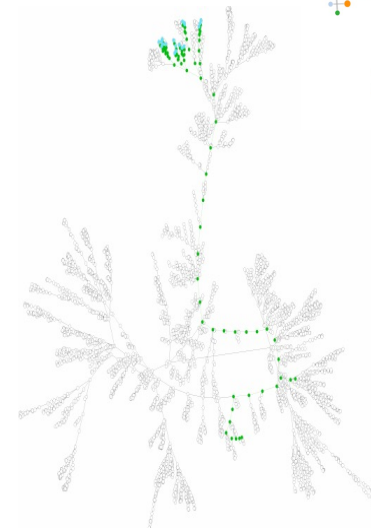
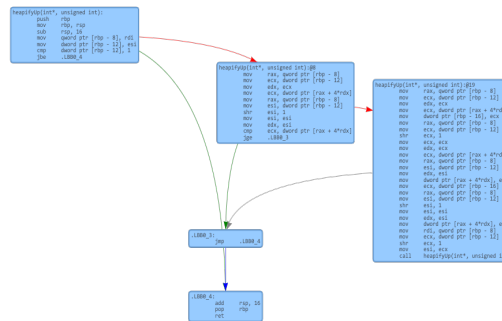
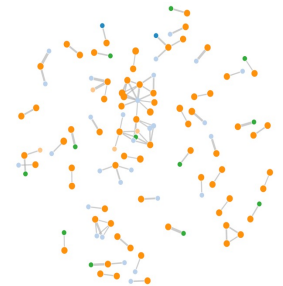
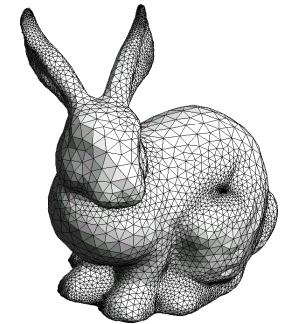
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TROILUS AND CRESSIDA



To study all of these structures:

1. A common vocabulary
2. Graph implementations
3. Graph traversals
4. Graph algorithms

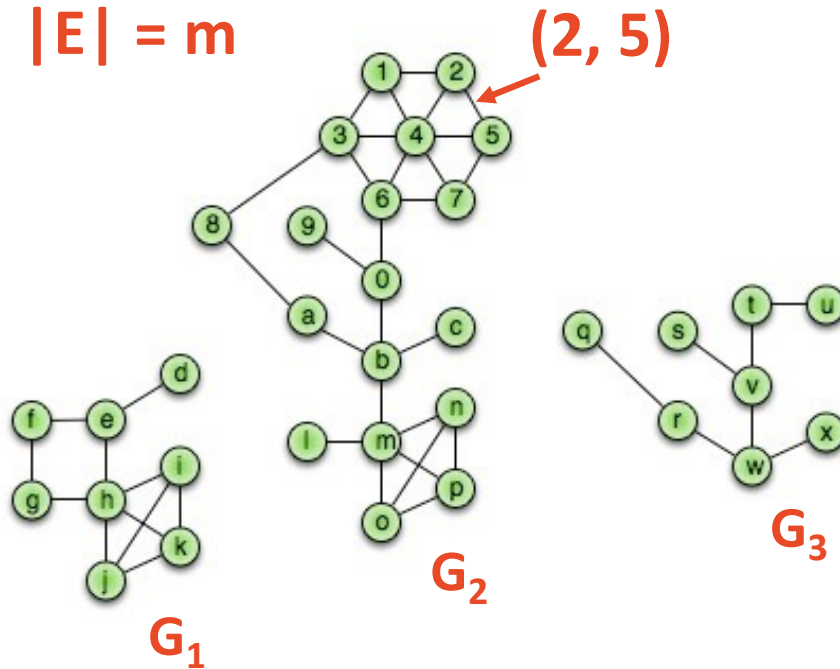


Graph Vocabulary

$$G = (V, E)$$

$$|V| = n$$

$$|E| = m$$



Incident Edges:

$$I(v) = \{ \{x, v\} \text{ in } E \}$$

Degree(v): $|I|$

Adjacent Vertices:

$$A(v) = \{ x : \{x, v\} \text{ in } E \}$$

Path or Walk: Sequence of vertices connected by edges

Path or Simple Path: Sequence of vertices connected by edges with no reuse

Cycle: A path with a common begin and end vertex with at least 3 vertices.

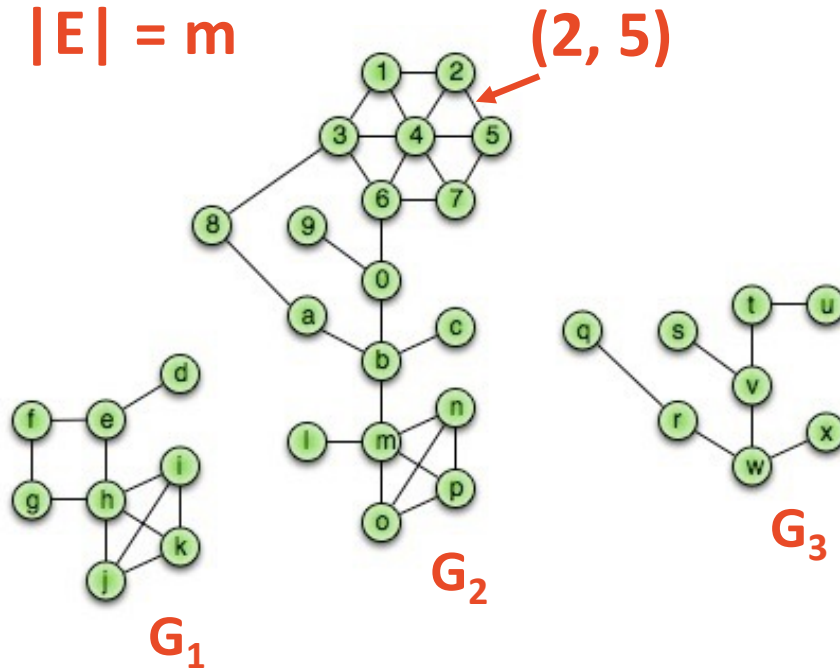
Simple Graph(G): A graph with no self loops or multi-edges.

Graph Vocabulary

$$G = (V, E)$$

$$|V| = n$$

$$|E| = m$$



Subgraph(G):

$$G' = (V', E'):$$

$V' \subseteq V, E' \subseteq E$, and

$$(u, v) \in E' \rightarrow u \in V', v \in V'$$

Complete subgraph(G)

Connected subgraph(G)

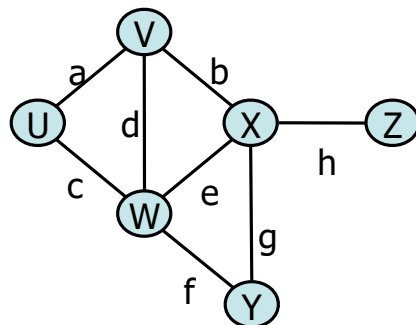
Connected component(G)

Acyclic subgraph(G)

Spanning tree(G)

Running times are often reported by **n**, the number of vertices, but often depend on **m**, the number of edges.

How many edges? **Minimum edges:**
Not Connected:



Connected*:

Maximum edges:
Simple:

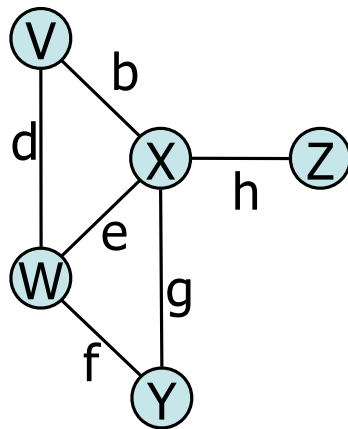
Not simple:

$$\sum_{v \in V} \deg(v) =$$

Graph ADT

Data:

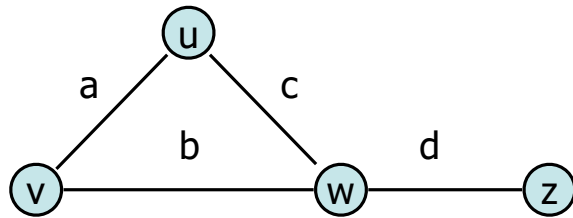
- Vertices
- Edges
- Some data structure maintaining the structure between vertices and edges.



Functions:

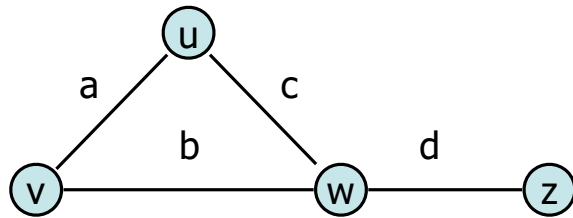
- insertVertex(K key);
- insertEdge(Vertex v1, Vertex v2, K key);
- removeVertex(Vertex v);
- removeEdge(Vertex v1, Vertex v2);
- incidentEdges(Vertex v);
- areAdjacent(Vertex v1, Vertex v2);
- origin(Edge e);
- destination(Edge e);

Graph Implementation Idea



Graph Implementation: Edge List

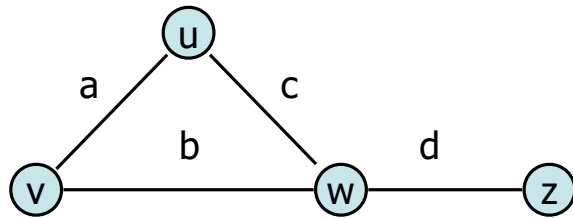
Vertex Collection:



u	u	v	a
v	v	w	b
w	u	w	c
z	w	z	d

Edge Collection:

Graph Implementation: Edge List

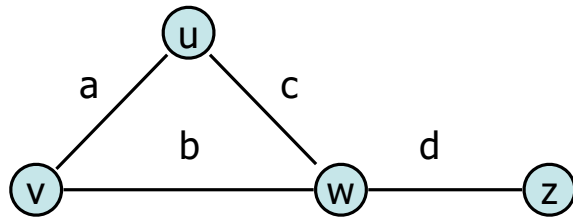


insertVertex(K key):

removeVertex(Vertex v):

u	u	v	a
v	v	w	b
w	u	w	c
z	w	z	d

Graph Implementation: Edge List



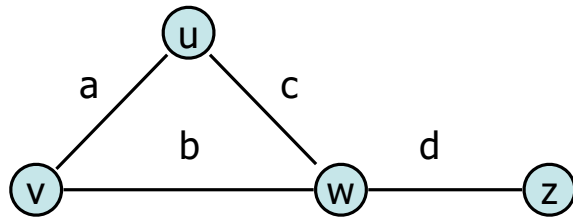
u	u	v	a
v	v	w	b
w	u	w	c
z	w	z	d

incidentEdges(Vertex v):

areAdjacent(Vertex v1, Vertex v2):

`G.incidentEdges(v1).contains(v2)`

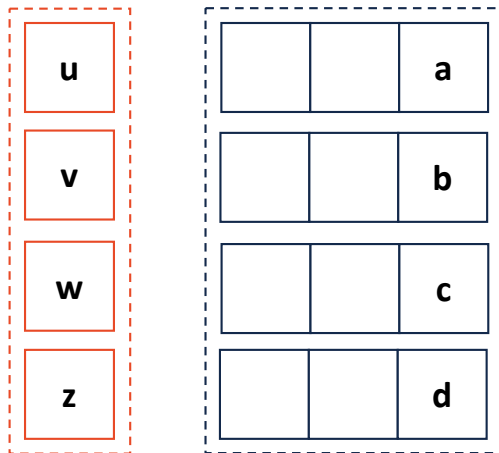
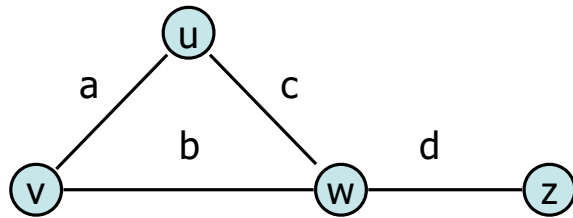
Graph Implementation: Edge List



insertEdge(Vertex v1, Vertex v2, K key):

u	u	v	a
v	v	w	b
w	u	w	c
z	w	z	d

Graph Implementation: Adjacency Matrix



insertVertex(K key);
removeVertex(Vertex v);
areAdjacent(Vertex v1, Vertex v2);
incidentEdges(Vertex v);

	u	v	w	z
u				
v				
w				
z				