Review Problem 3

- In assembly, compute the average of positive values X0, X1, X2, X3, and put into X10

```
ADD X4, X0, X1
ADD X5, X2, X3
ADD X5, X4, X5
LSR X10, X5, #2
```

Addressing Example

The address of the start of a character array is stored in X0. Write assembly to load the following characters

X2 = Array[0]

```
LDURB  X2,[X0, #0]
```

X3 = Array[1]

```
LDURB  X3, [X0, #1]
```

X4 = Array[2]

```
LDURB X4, [X0, #2]
```

X5 = Array[k] // Assume the value of k is in X1

MEM[X0 + X1^k]

```
ADD X5, X0, X1      // X5 = &(Array[k])
LDURB X5, [X5, #0]
```

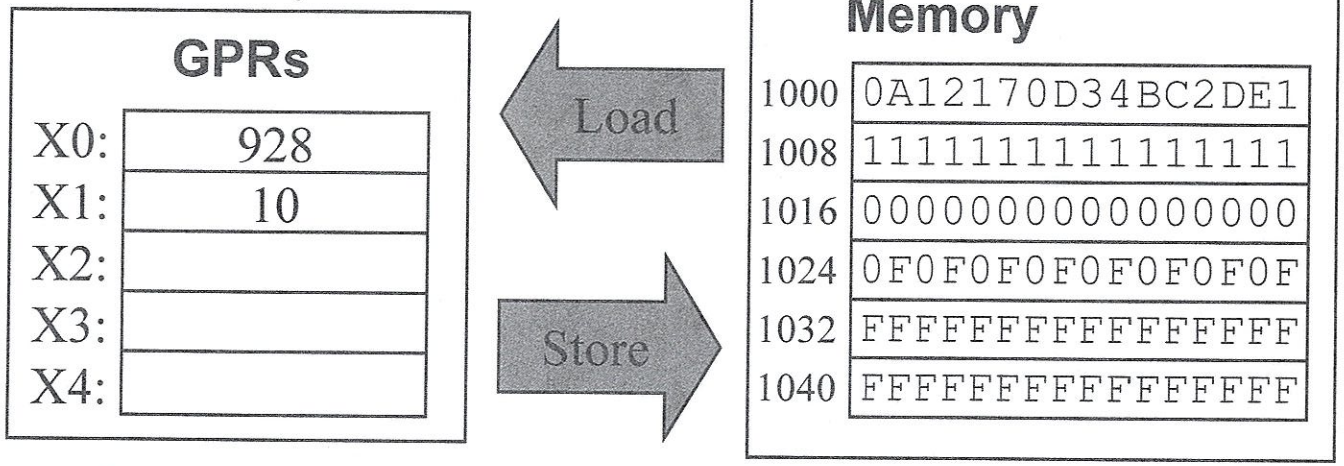
V[0] = MEM[X0] = mem[928] V[1] = MEM[X0+8] = MEM[9

V[K] = MEM[X0 + 8xk] V[K+1] = MEM[X0 + 8*K+8]

Array Example

```
/* Swap the kth and (k+1)th element of an array */
swap(int v[], int k) {
      int temp = v[k];
      v[k] = v[k+1];
      v[k+1] = temp;
}

// Assume v in X0, k in X1
```



SWAP:

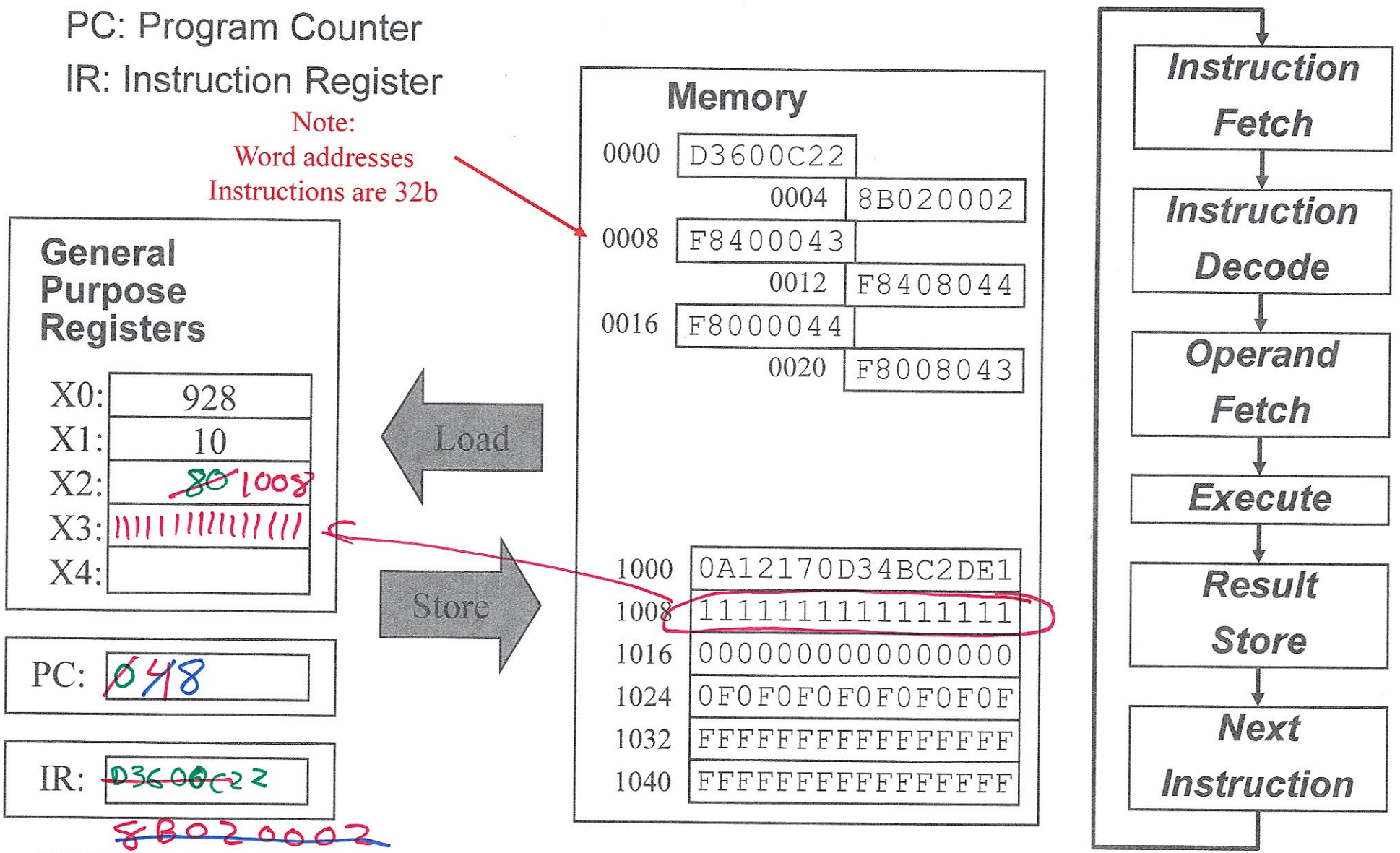
```
LSL  X2, X1, #3      // X2 = 8*K
ADD  X2, X0, X2      // X2 = &V[K]
LDUR X3, [X2, #0]    // get V[K]
LDUR X4, [X2, #8]    // get V[K+1]
STUR X4, [X2, #0]
STUR X3, [X2, #8]
```

Execution Cycle Example

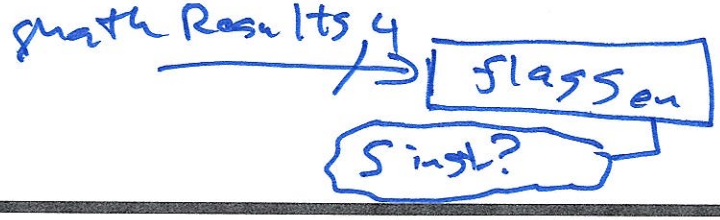
PC: Program Counter

IR: Instruction Register

Note:
Word addresses
Instructions are 32b



Flags/Condition Codes



Flag register holds information about result of recent math operation

- Negative: was result a negative number?
- Zero: was result 0?
- Overflow: was result magnitude too big to fit into 64-bit register?
- Carry: was the carry-out true?

Operations that set the flag register contents:

ADDS, ADDIS, ANDS, ANDIS, SUBS, SUBIS, some floating point.

Most commonly used are subtracts, so we have a synonym: CMP

CMP X0, X1 same as SUBS X31, X0, X1

CMPI X0, #15 same as SUBIS X31, X0, #15

Control Flow

Unconditional Branch – GOTO different next instruction

```
B START          // go to instruction labeled with "START" label
BR X30           // go to address in X30: PC = value of X30
```

Conditional Branches – GOTO different next instruction if condition is true

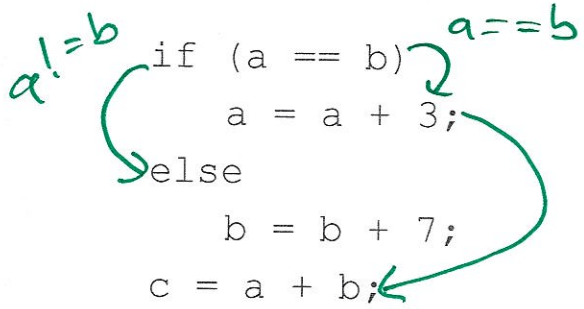
1 register: CBZ (==0), CBNZ (!= 0)

```
CBZ X0, FOO      // if X0 == 0 GOTO FOO: PC = Address of instr w/FOO label
```

2 register: B.LT (<), B.LE(<=), B.GE (>=), B.GT(>), B.EQ(==), B.NE(!=)

first compare (CMP X0, X1, CMPI X0, #12), then b.cond instruction

```
CMP X0, X1       // compare X0 with X1 - same as SUBS X31, X0, X1
B.EQ FOO         // if X0 == X1 GOTO FOO: PC = Address of instr w/FOO label
```




```
// X0 = a, X1 = b, X2 = c
    CMP X0, X1          // set flags
    B.NE ELSEIF         // branch if a!=b
    ADDI X0, X0, #3     // a = a + 3
    B DONE              // avoid else
ELSEIF:
    ADDI X1, X1, #7     // b = b + 7
DONE:
    ADD X2, X0, X1      // c = a + b
```