



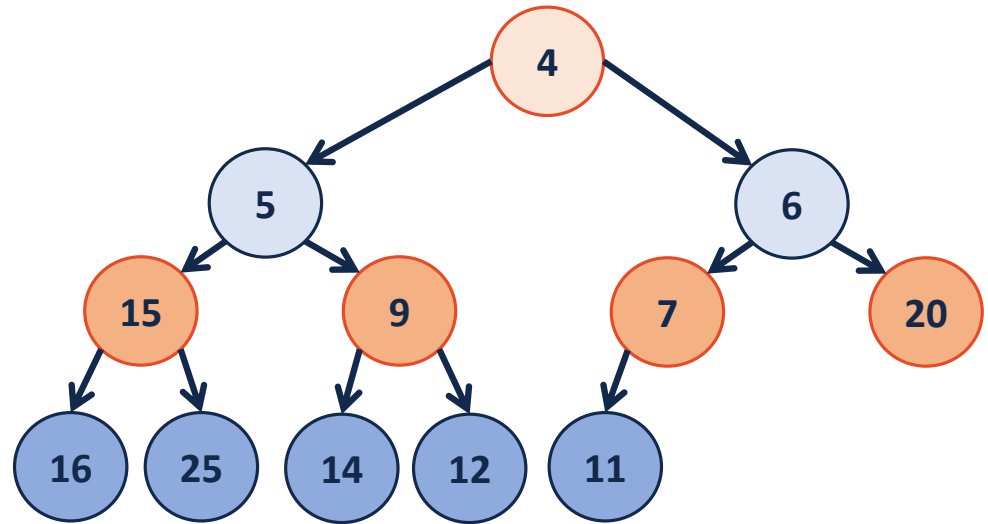
CS 225

Data Structures

March 10 – Heaps

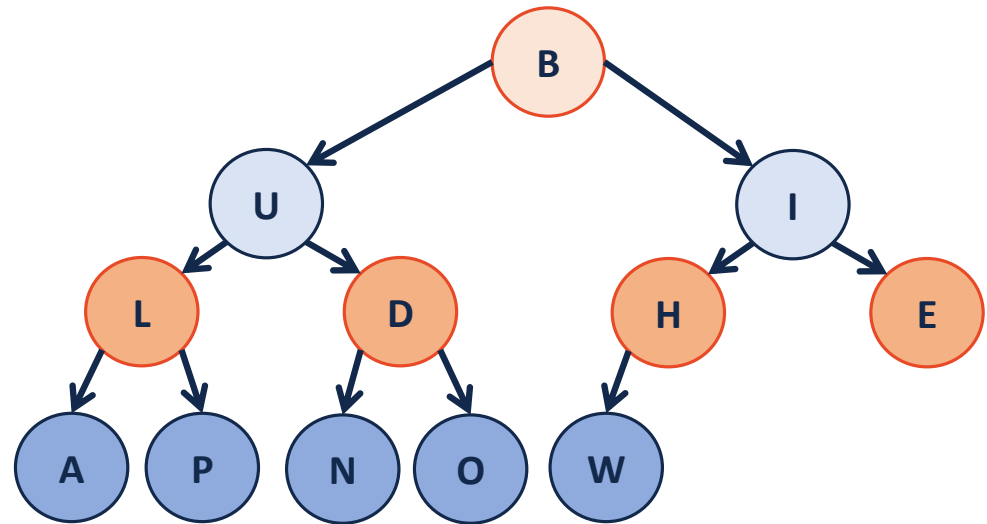
G Carl Evans

(min)Heap

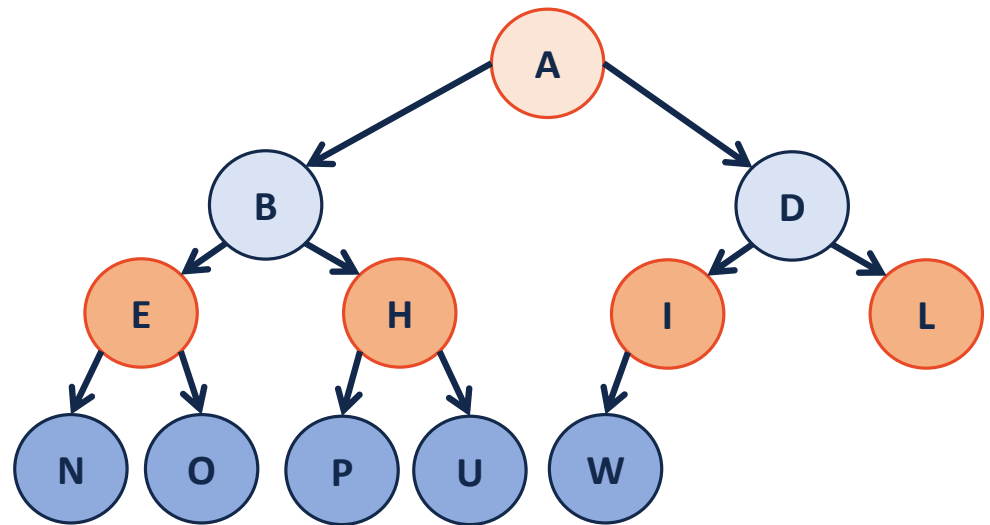


4	5	6	15	9	7	20	16	25	14	12	11			
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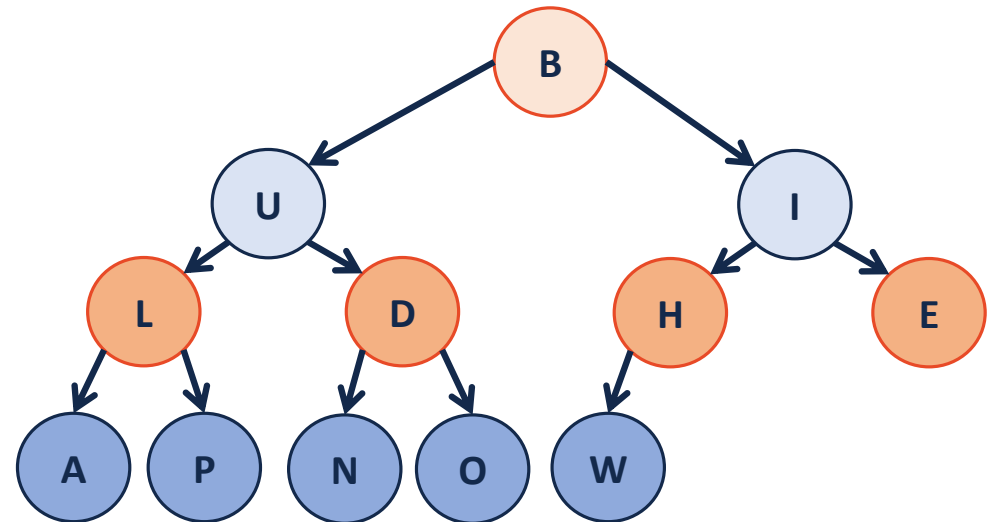
buildHeap



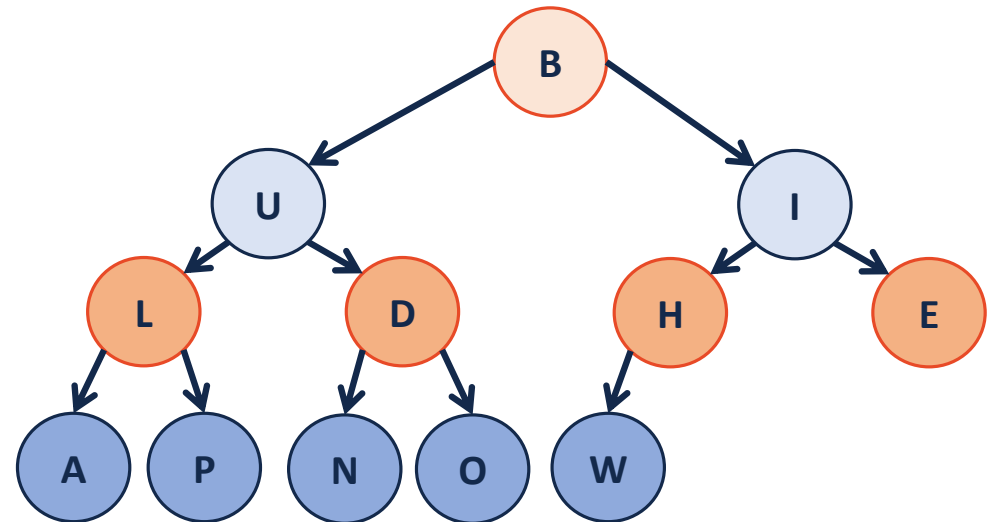
buildHeap – sorted array



buildHeap - heapifyUp



buildHeap - heapifyDown



buildHeap

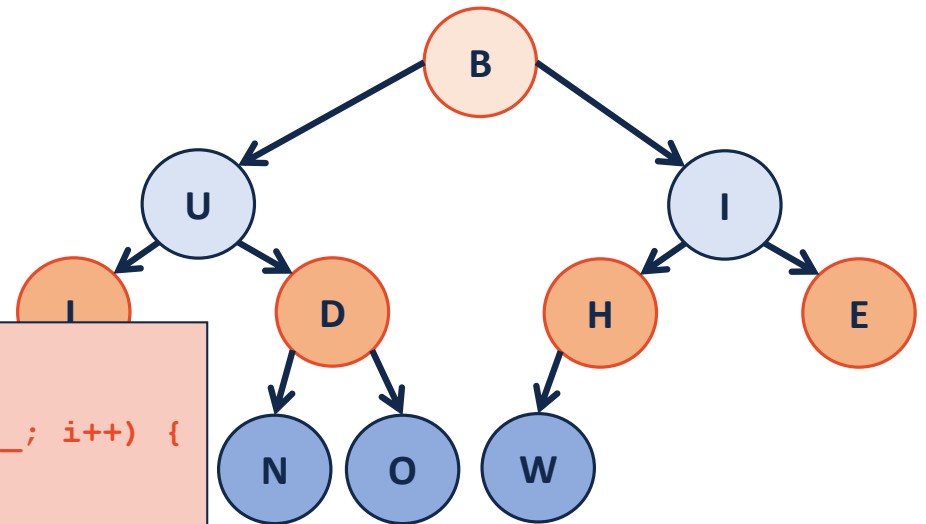
1. Sort the array – it's a heap!

2.

```
1 template <class T>
2 void Heap<T>::buildHeap() {
3     for (unsigned i = 2; i <= size_; i++) {
4         heapifyUp(i);
5     }
6 }
```

3.

```
1 template <class T>
2 void Heap<T>::buildHeap() {
3     for (unsigned i = parent(size); i > 0; i--) {
4         heapifyDown(i);
5     }
6 }
```





Proving buildHeap Running Time

Theorem: The running time of buildHeap on array of size n is: _____.

Strategy:

-

-

-

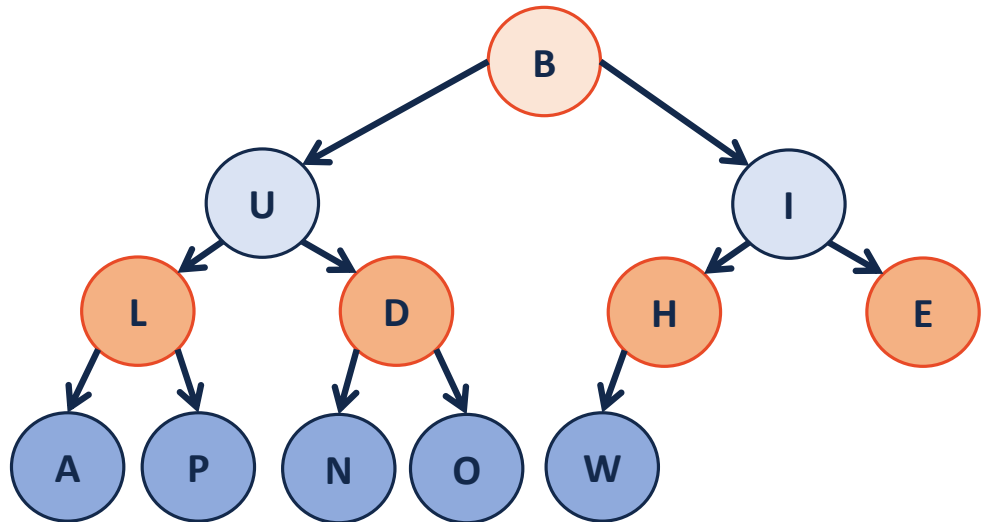
Proving buildHeap Running Time

$S(h)$: Sum of the heights of all nodes in a complete tree of height **h** .

$S(0)$ =

$S(1)$ =

$S(h)$ =





Proving buildHeap Running Time

Proof the recurrence:

Base Case:

General Case:



Proving buildHeap Running Time

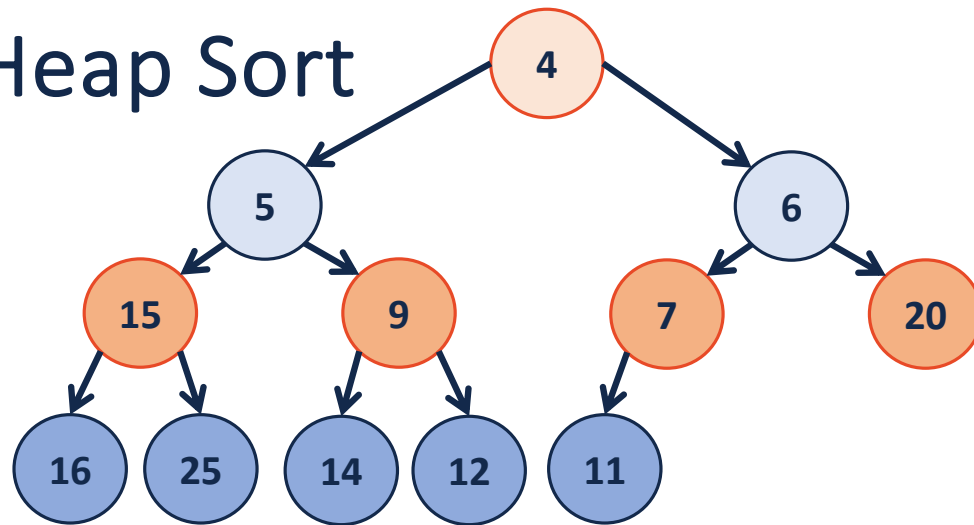
From $S(h)$ to RunningTime(n):

$S(h)$:

Since $h \leq \lg(n)$:

RunningTime(n) \leq

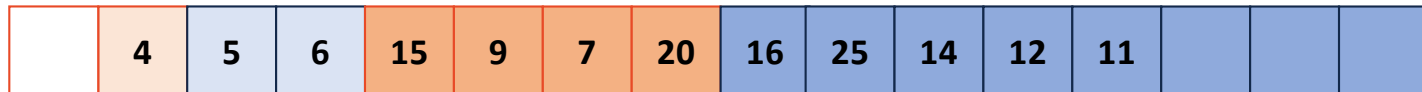
Heap Sort



1.

2.

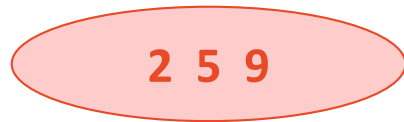
3.



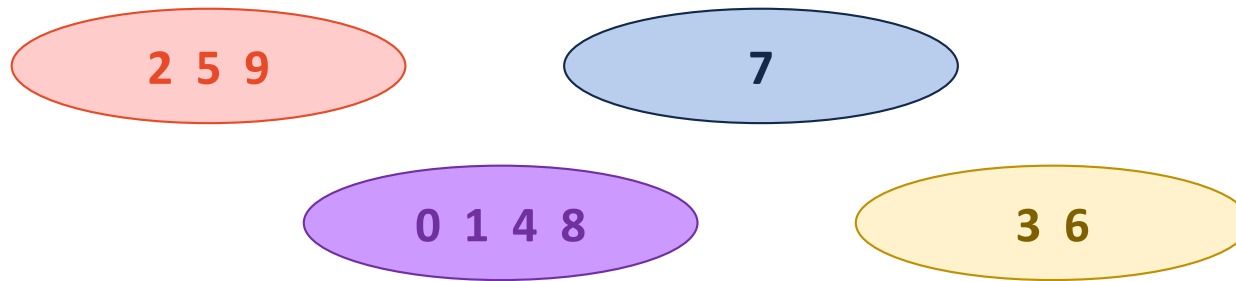
Running Time?

Why do we care about another sort?

Disjoint Sets

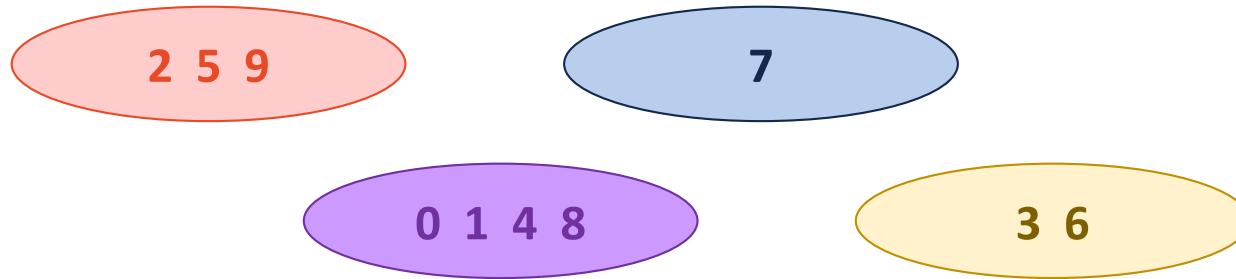


Disjoint Sets



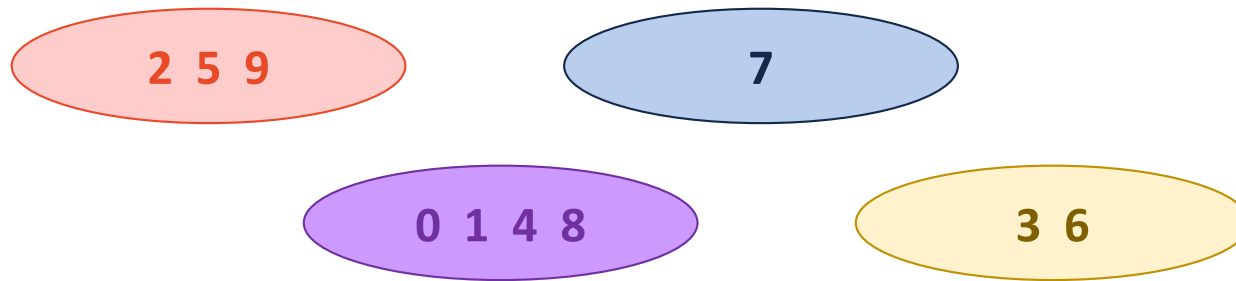
Operation: find(4)

Disjoint Sets



Operation: $\text{find}(4) == \text{find}(8)$

Disjoint Sets



Key Ideas:

- Each element exists in exactly one set.
- Every set is an equitant representation.
 - Mathematically: $4 \in [0]_R \rightarrow 8 \in [0]_R$
 - Programmatically: `find(4) == find(8)`

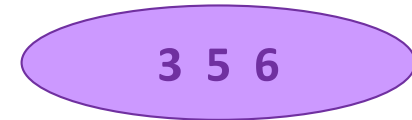
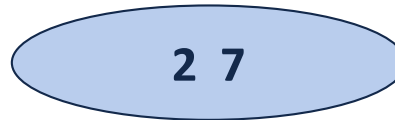


Disjoint Sets ADT

- Maintain a collection $S = \{s_0, s_1, \dots s_k\}$
- Each set has a representative member.
- API:

```
void makeSets(int number);  
void union(int k1, const int k2);  
int  find(int k);
```

Implementation #1



0	1	2	3	4	5	6	7

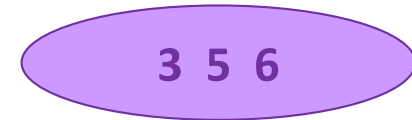
Find(k):

Union(k1, k2):

YOU EXPECTED A NEW DATA STRUCTURE

BUT IT WAS ME, TREE ALL ALONG

Implementation #2



0	1	2	3	4	5	6	7

Find(k):

Union(k1, k2):

Implementation #2

- We will continue to use an array where the index is the key
- The value of the array is:
 - **-1**, if we have found the representative element
 - **The index of the parent**, if we haven't found the rep. element
- We will call theses **UpTrees**:

