Tutorial
ELEC3506/9506
Communication Networks

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Tutorial 03 – Week 04

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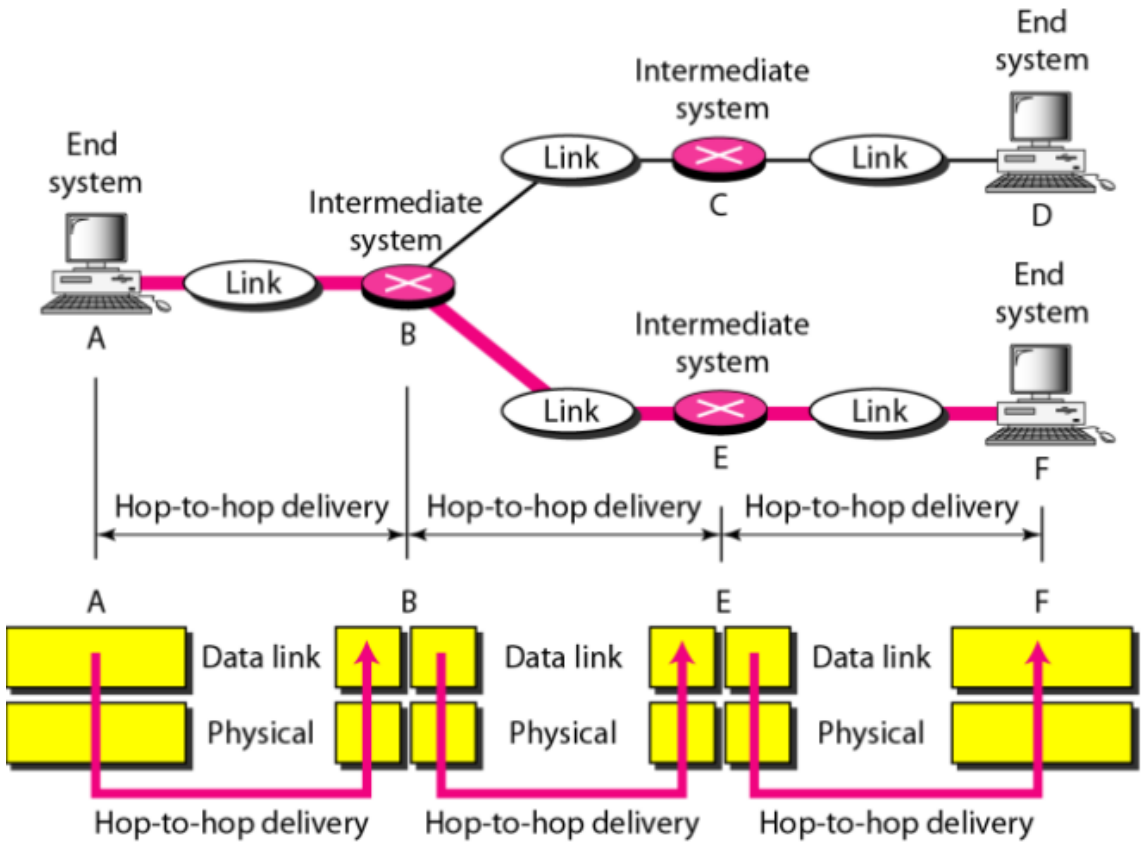
Q1. Briefly describe the services provided by the data link layer.

- **Framing Data Bits**: Divides the stream of bits received from the network layer into manageable frames.
- **Moving Frames**: Transfers frames from one hop (node) to the next in the network.
- **Physical Addressing**: Provides physical addresses (MAC addresses) for identifying the sender and receiver.
- **Flow Control**: Prevents the sender from overwhelming the receiver with too much data too quickly.
- **Error Control**: Adds reliability by detecting and correcting errors in the data transmitted over the physical layer.
- **Access Control**: Manages which device has control over the network link at any given time to avoid conflicts.



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Q2. Define framing and the reason for its need.

❑ Framing

A process in data link layer to divide the stream of bits received from the network layer into frames.

❑ Reasons for framing:

- Framing adds the physical sender and destination addresses. The destination address specifies where the frame should go; the sender address is used to send an acknowledgement (ACK) back to the sender.
- Framing in data link layer divides the message into appropriate-sized frames such that the flow control and error control become efficient. Smaller-sized frames make flow control and error control more efficient.

A frame has the following parts:

- ❑ **Frame header**: Contains the source and destination address
- ❑ **Payload field**: Contains the message to be delivered.
- ❑ **Trailer**: Contains error detection & correction bits.
- ❑ **Flag**: Marks the beginning and end of the frame.

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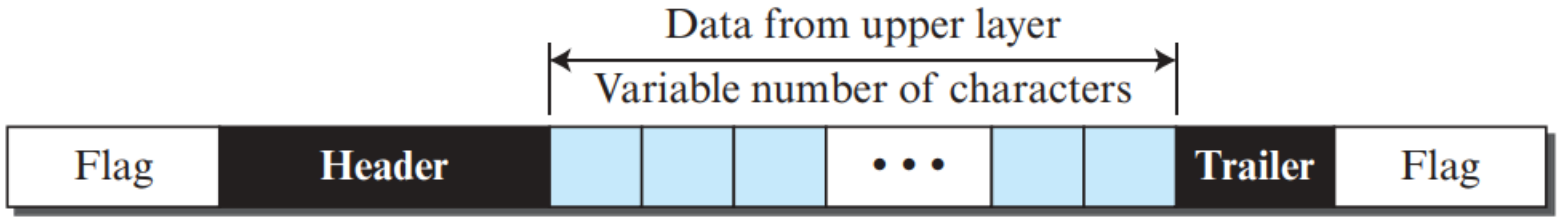
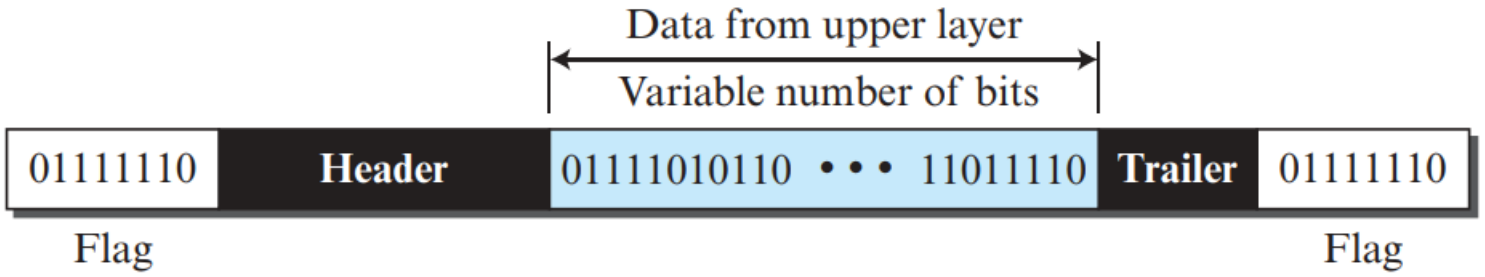
Q3. Compare and contrast byte-oriented and bit-oriented protocols. Which category has been popular in the past? Which category is popular now?

- **Byte-oriented (Character-oriented) protocol**
  - Data to be carried are 8-bit characters from a coding system, e.g., ASCII code
  - Character-oriented protocols were popular when only text was exchanged by the data link layers



- **Bit-oriented protocol:**
  - Data is transmitted as a stream of bits.
  - More popular today, because we need to send text, graphic, audio, and video, which can be better represented by a bit pattern rather than by using a sequence of characters.



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Q4. Compare and contrast the three most common error detection methods

- **Parity Checking (Vertical Redundancy Checking - VRC)**
  - One of the oldest and simplest methods
  - Adds 1 additional bit to each byte in the message.
  - Even parity causes the sum of all bits to be even.
  - Odd parity causes the sum of bits to be odd
- **Simple but only has about a 50% reliability rate.**

Example: Even Parity

Message	Codeword
000	0000
001	0011
010	0101
011	0110
100	1001
101	1010
110	1100
111	1111

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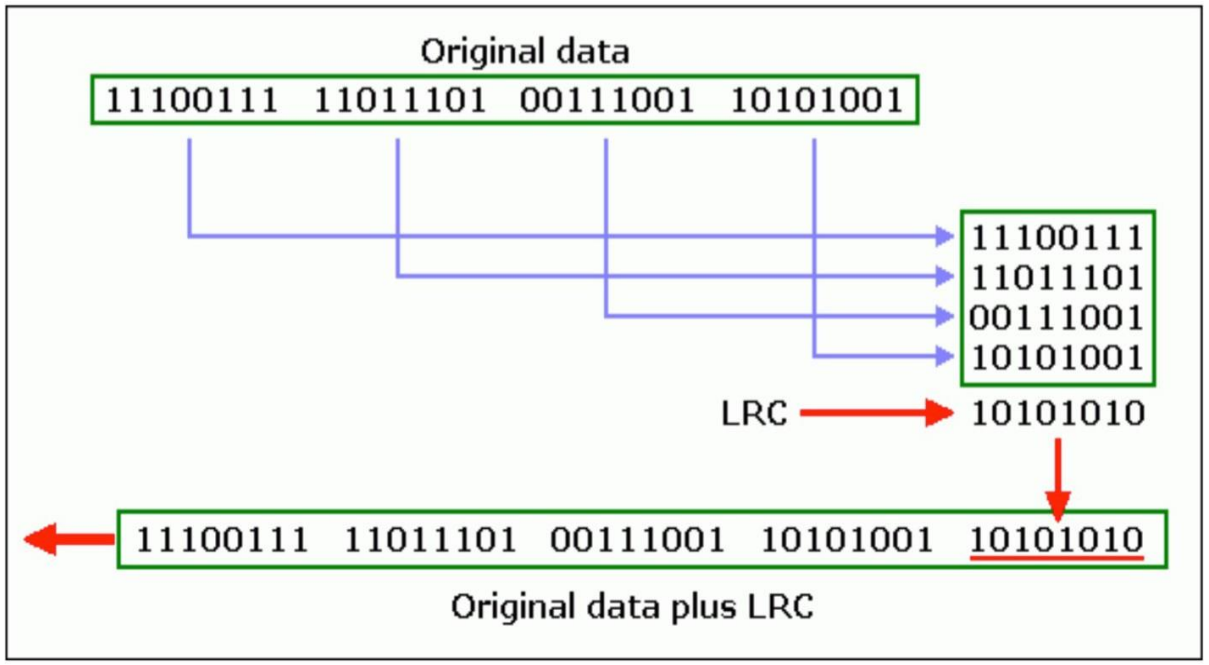
Q4. Compare and contrast the three most common error detection methods

- **Longitudinal Redundancy Checking (LRC)**
  - A block of bits is divided into rows and redundant row of bits is added to the whole block.



- At the receiver side, this redundant row is used to detect the error.
- The receiver calculates the LRC and compares it to the LRC appended to the transmitted data. If two LRCs match, no error is detected.
- Compared to VRC, LRC increases the likelihood of detecting burst errors.

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Q4. Compare and contrast the three most common error detection methods

- **Cyclic Redundancy Check Codes (CRC)**
  - Message length k, codeword length n
  - **Encode**
    - Multiply $u(x)$ with x^{n-k}
    - Divided $u(x)x^{n-k}$ by $g(x)$ and obtain the remainder $b(x)$
    - Forming the codeword

      $$\boldsymbol{v(x) = b(x) + u(x)x^{n-k}}$$

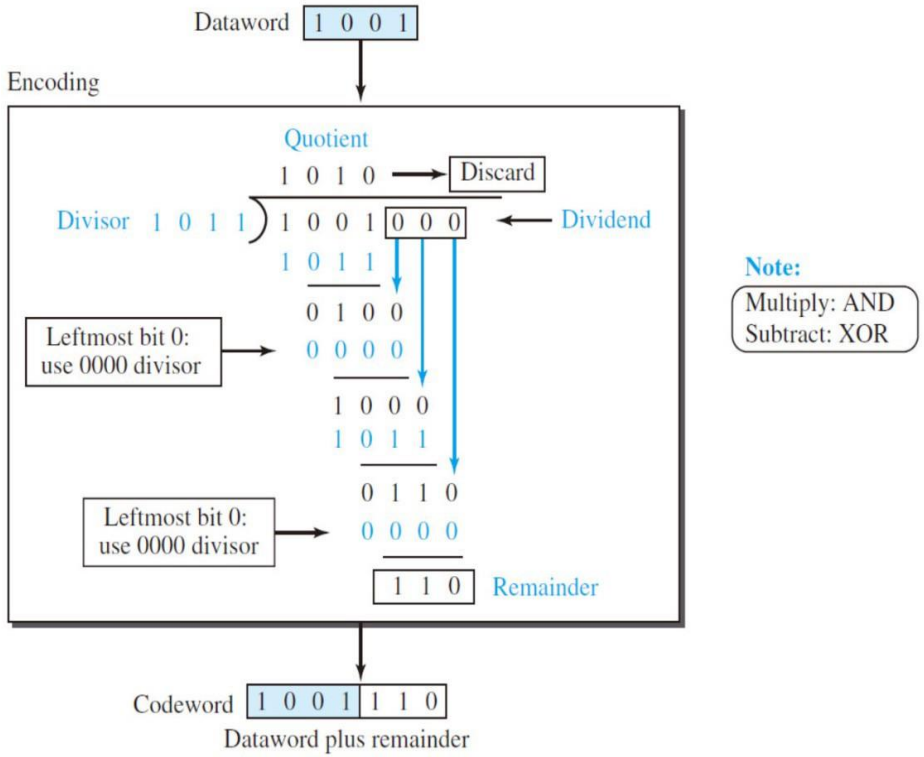
- Decode
  - Divide $v(x)$ by $g(x)$
  - If the remainder is zero, no error
  - Otherwise, error

  - CRC can detect more types of error than VRC and LRC.

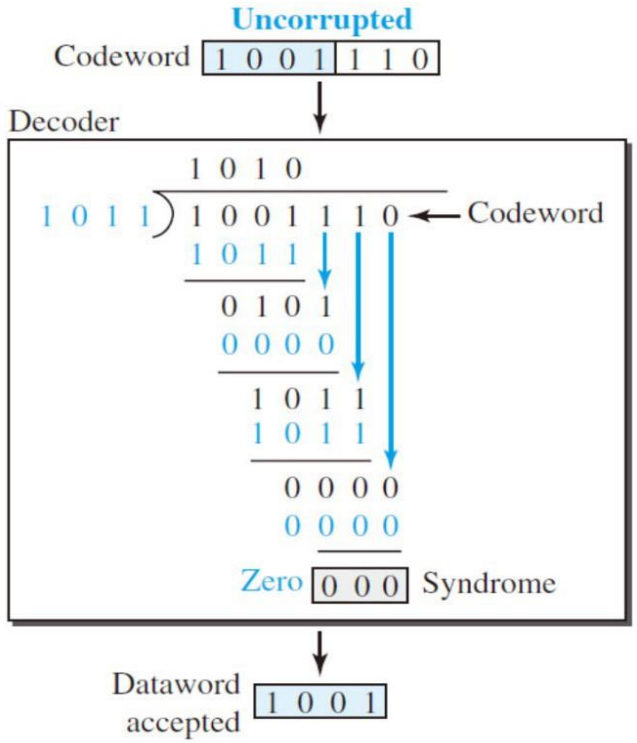
Q4. Compare and contrast the three most common error detection methods

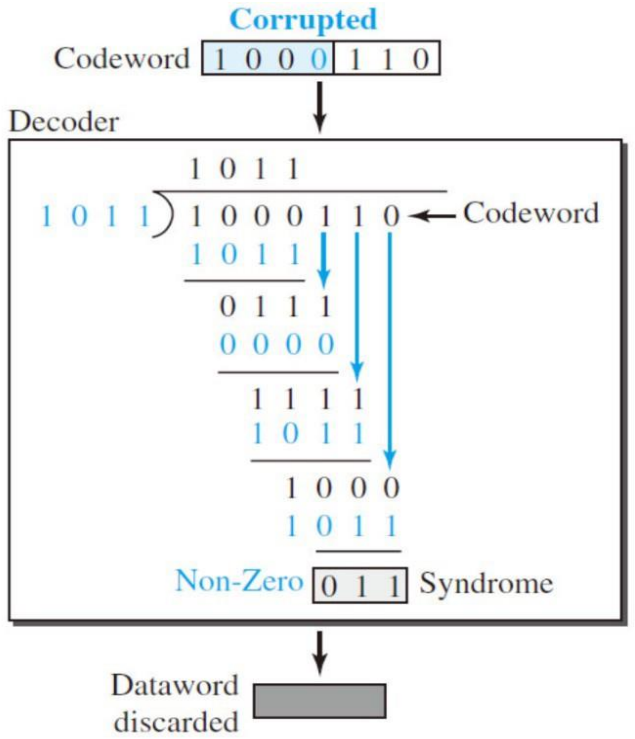
Cyclic Redundancy Check Codes (CRC)

Encoder



Note:
Multiply: AND
Subtract: XOR





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Q5. Compare and contrast flow control and error control. How does error control complement flow control?

- **Flow Control** refers to a set of procedures used to restrict the amount of data that sender can send before waiting for acknowledgement.
  - Any receiving device has a limited speed for processing data.
  - The processing rate of the receiver is usually slower than the transmission rate.
  - Each receiving device has a buffer for storing the incoming data before they are processed.
  - When the buffer begins to fill up, the receiver must tell the sender to slow down the transmission.

- **Error control** refers to a set of procedures used to detect and correct errors.
  - Both flow control and error control help in proper delivering of data to the receiver.
  - Flow control ensures that frames are not discarded due to buffer overflow. Error control complements flow control by ensuring that the frames received at the received side are error-free.

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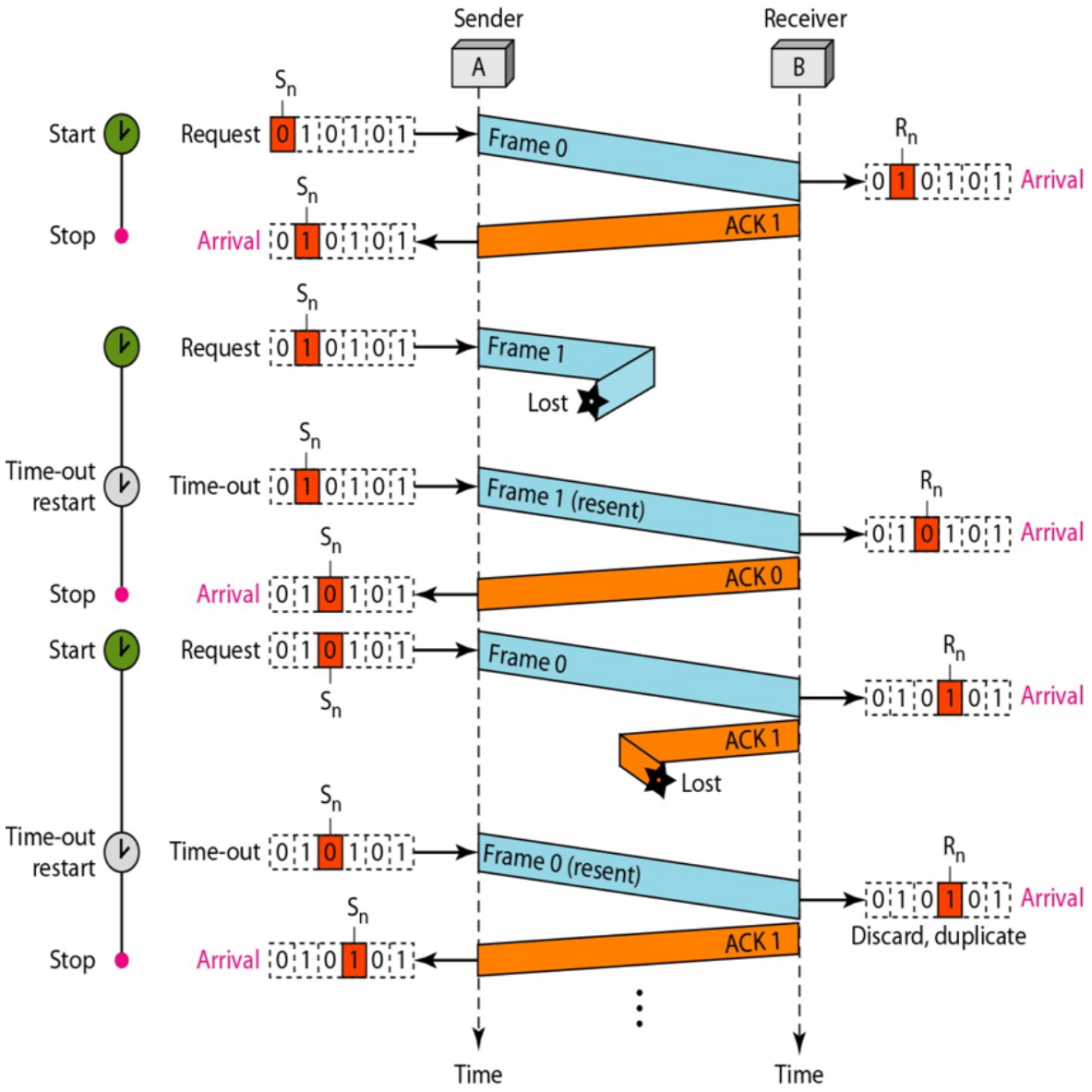
Q6. State three error control protocols for noisy channels (as discussed in the lecture).

- ❑ Stop-and-wait Automatic Repeat reQuest (ARQ) (Send one frame at a time)
- ❑ Go-back-N ARQ (Send several frames at a time)
- ❑ Selective-Repeat ARQ (Send several frames at a time)

Q7. Why is stop-and-wait ARQ inefficient?

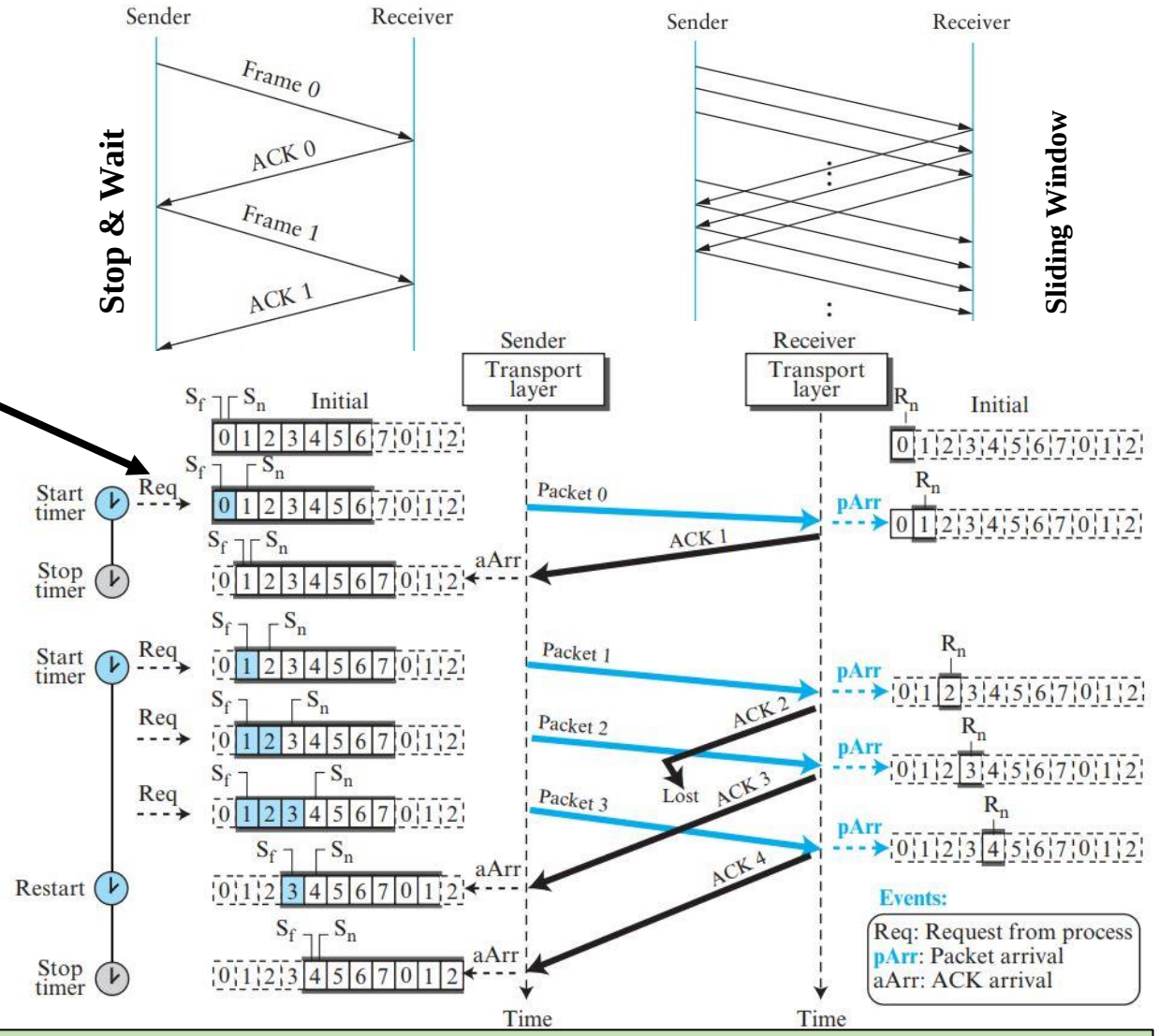


- Slow speed : The sender has to wait for a round trip time (RTT) before another frame can be sent.
- Low bandwidth ulitization : Only one frame can be sent at a time

Q8. What are the advantages and disadvantages of Go-Back-N ARQ?

Advantages

- **Multiple packet transmission**: the sender is allowed to transmit multiple packets (when available) without waiting for an acknowledgment.
- **Cumulative ACK**: Go-Back-N can acknowledge more than one frame with a single ACK, which reduces the time and bandwidth wasted for sending ACKs.
- **Complexity**: Less complicated than the Selective-Repeat ARQ. In case of Selective-Repeat ARQ, negative ACK (NACK) is required.
- **Buffering**: No buffering is needed at the receiver.
- **Sorting**: No packet sorting is needed at the receiver.

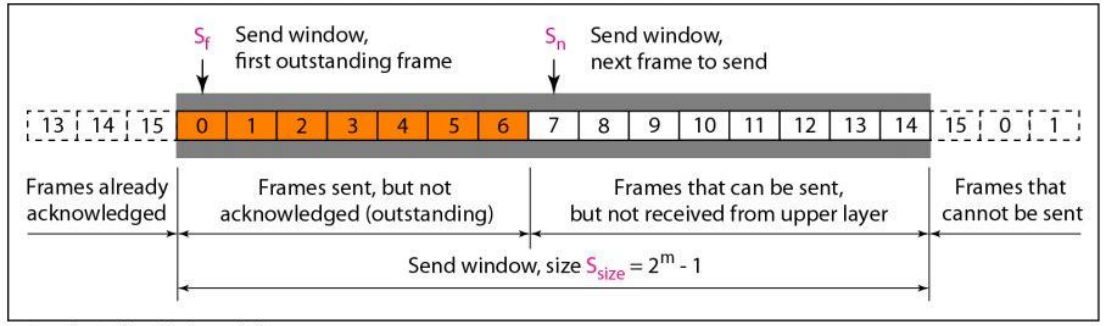


- Assume a sender sends 6 packets: packets 0, 1, 2, 3, 4, and 5. The sender receives an ACK with ackNo = 3. **That means?**
- If the system is using GBN, it means that packets 0, 1, and 2 have been received uncorrupted and the receiver is expecting packet 3.

Q8. What are the advantages and disadvantages of Go-Back-N ARQ?

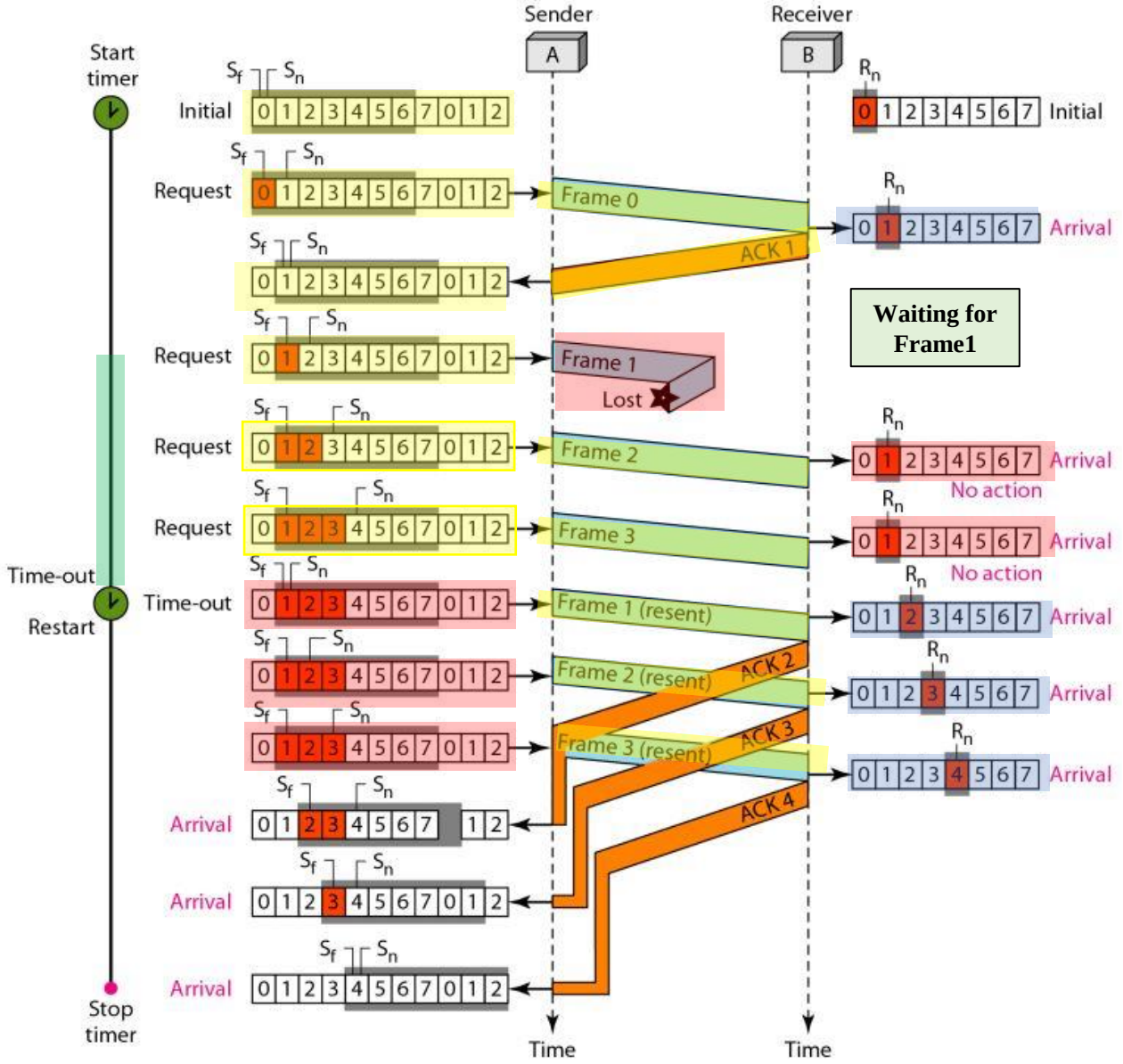
Disadvantages

- **Retransmission**: This protocol is very inefficient for a noisy link. Unnecessary resending uses up the bandwidth and slows down the transmission.
- **Transmitter Buffer**: Buffer requirement at the transmitter. Transmitter needs to store N packets.



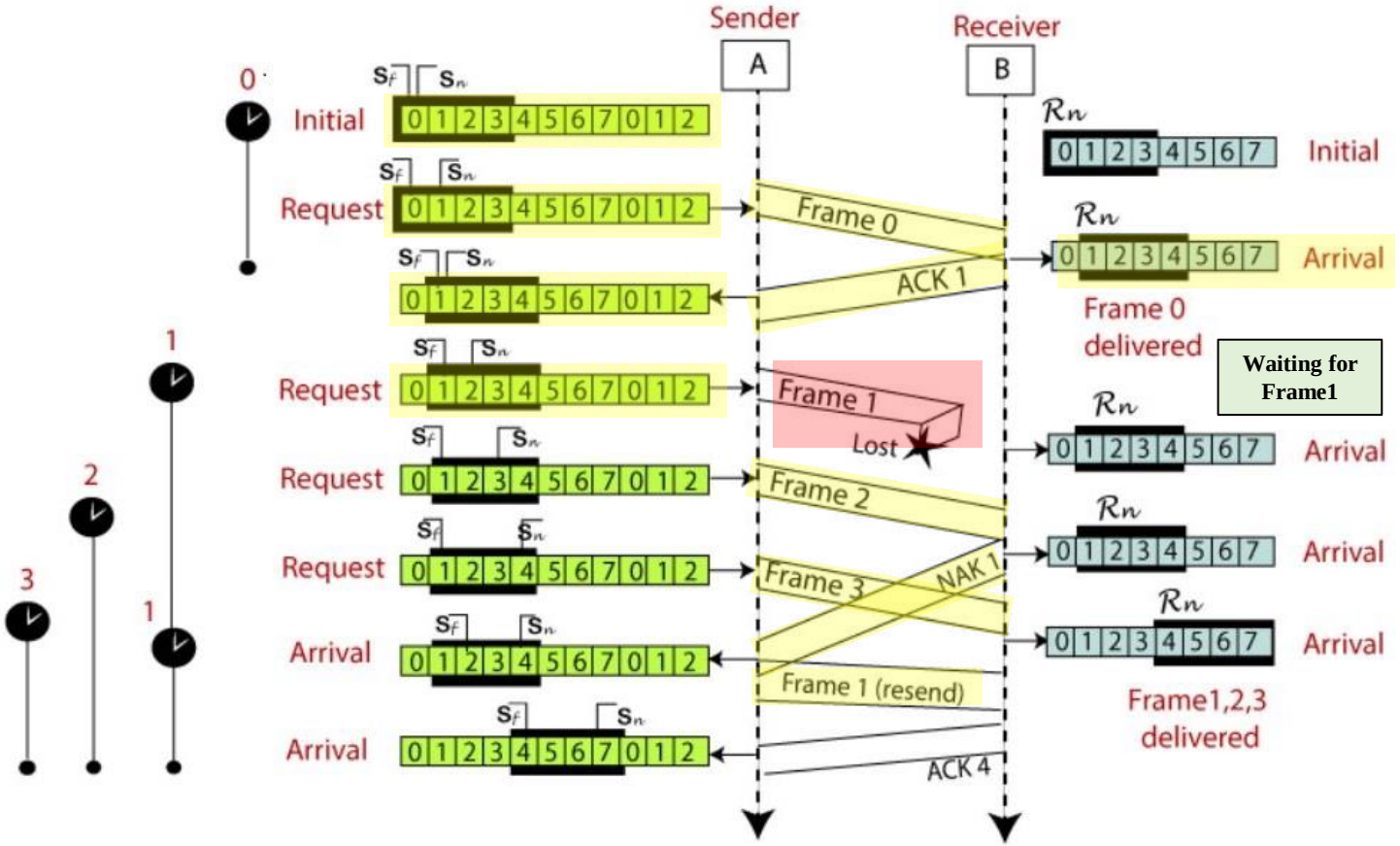
a. Send window before sliding

- **High bandwidth requirement**



Q9. How does Selective Repeat ARQ address the deficiencies of Go-Back-N ARQ? (Cont…)

- Unnecessary retransmission of the successfully received packets is one of the major problems in Go back N protocol.
- In **Selective-Repeat ARQ**, unnecessary transmission is avoided by sending only the frames that are corrupted or missing.
- **Negative Acknowledgement**: Selective-Repeat ARQ uses negative acknowledgement (NAK) which can speed up the retransmission of frames.
- Before the sender's timer expires for a lost frame, the receiver can send a NAK to let the sender know that a frame was lost or received incorrectly.



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Q10. Standard Ethernet defines several Physical layer implementations. Discuss the most common implementation categories?

❑ Standard Ethernet has defined several physical layer implementations, out of which the following four are commonly used.

Name	Media	Max Length	Line Encoding
10Base5	Thick coaxial cable	500 m	Manchester
10Base2	Thin coaxial cable	185 m	Manchester
10Base-T	2 UTP	100 m	Manchester
10Base-F	2 Fiber	2000 m	Manchester

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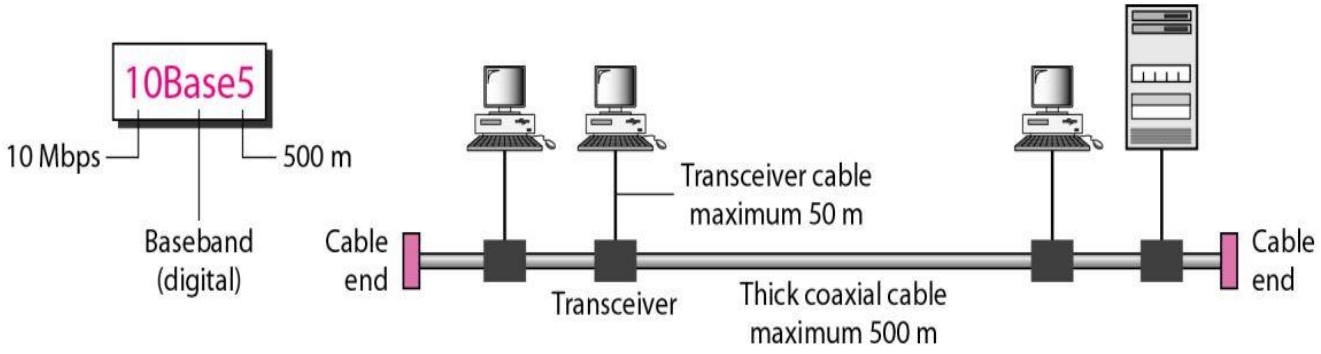
Q10. Standard Ethernet defines several Physical layer implementations. Discuss the most common implementation categories?

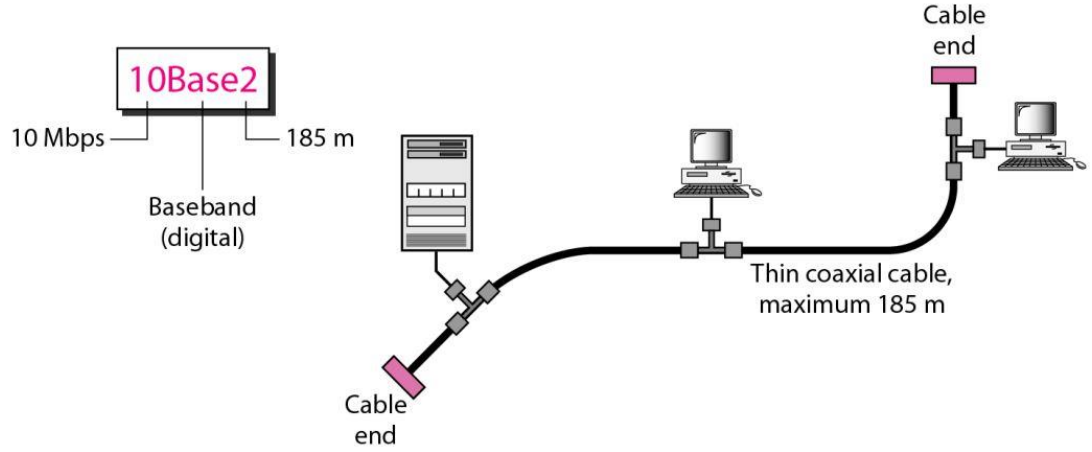
10Base5 (Thick Ethernet)

- 10Base5 is implemented in bus topology with an external transceiver connected to coaxial cable.
- The transceiver deals with transmitting, receiving and detecting collisions in the network.
- The 10Base5 Ethernet is also referred to by other names including thick **Ethernet** or **Thicknet**.
- The maximum length of the cable should not be more than 500 m.



10Base2 (Thin Ethernet)

- 10Base2 is also implemented in the bus topology but with a thinner and flexible coaxial cable.
- The cable can be bent to pass very close to the stations.
- The transceiver is normally part of the NIC card, which is installed inside the station.
- The 10Base2 Ethernet is also referred to as **Thin Ethernet** or **Cheapernet**.
- The maximum length of the cable should not be more than 185 m



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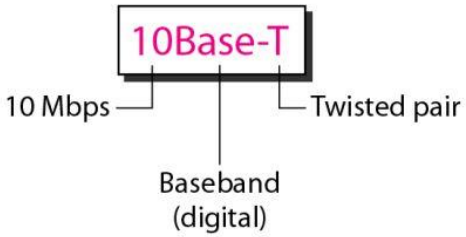
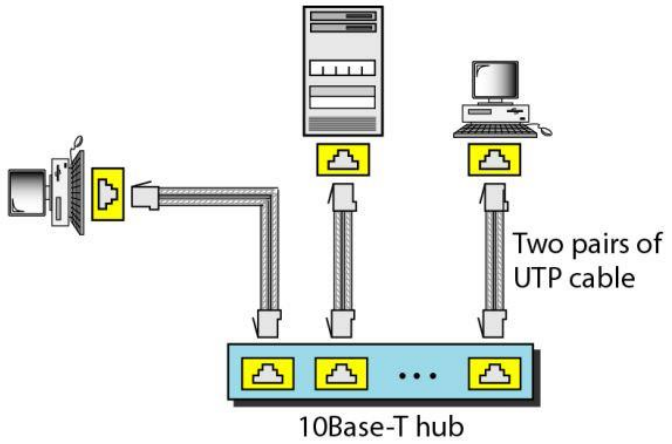
Q10. Standard Ethernet defines several Physical layer implementations. Discuss the most common implementation categories?

10Base-T (Twisted-Pair Ethernet)

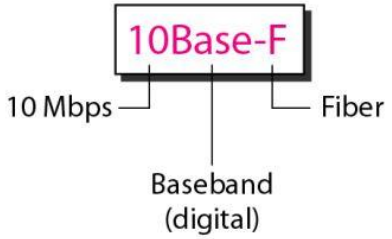
- The third implementation is called l0Base-T or twisted-pair Ethernet.
- 10Base-T uses two pairs of twisted cable and is implemented in star topology.
- All nodes are connected to the hub via two pairs of cable and thus, creating a separate path for sending and receiving the data.
- The maximum length of the cable should not be more than 100 m.

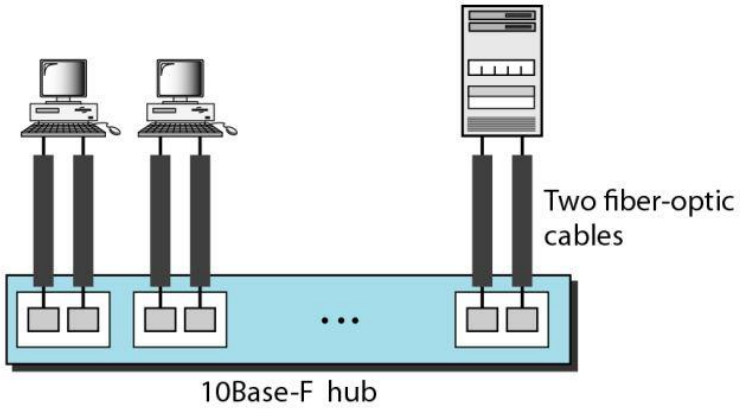




10Base-F (Fiber Ethernet)

- It is the most common 10-Mbps Ethernet that is implemented in star topology.
- It uses a pair of fiber optic cables to connect the nodes to the central hub.
- The transceiver is normally part of the NIC card, which is installed inside the station.
- The 10Base-F Ethernet is also referred to as **Fiber Ethernet**.
- The maximum length of the cable should not be more than 2000 m

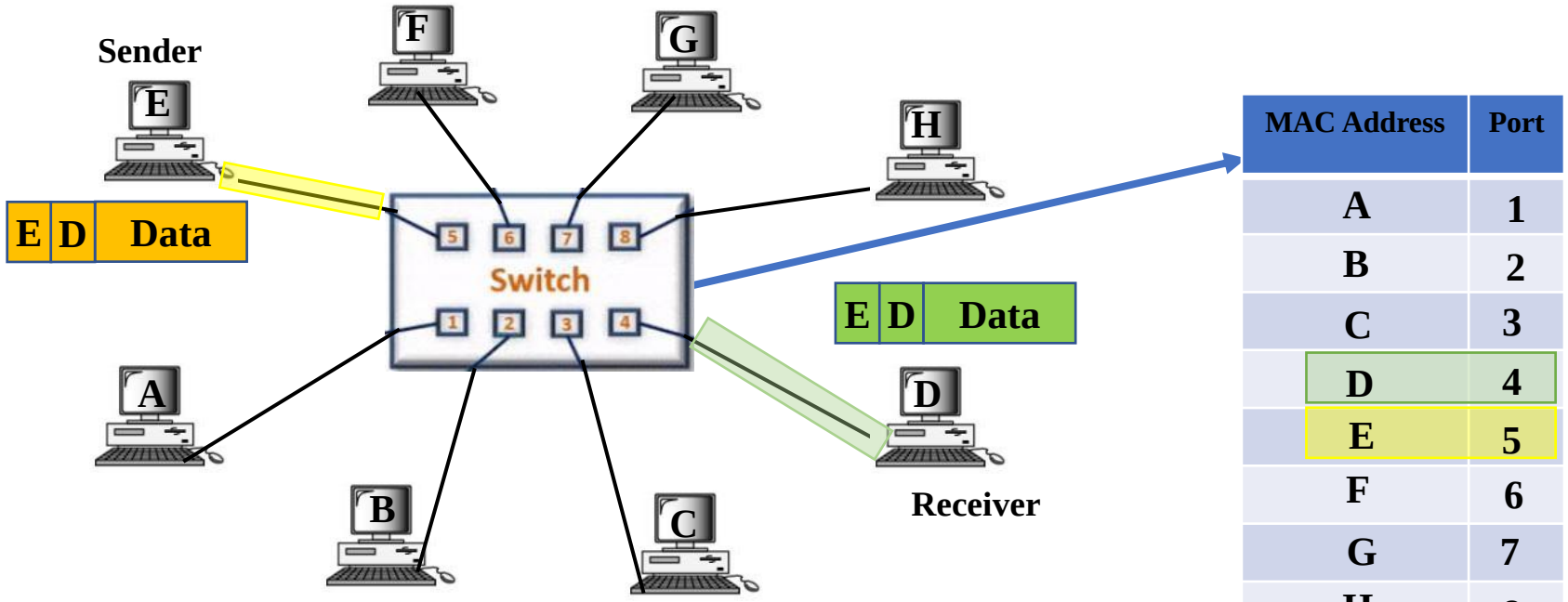




Q11. What is switched Ethernet?

- **Switch Ethernet** : It connects all the devices on same network, so that they can communicate with one another.

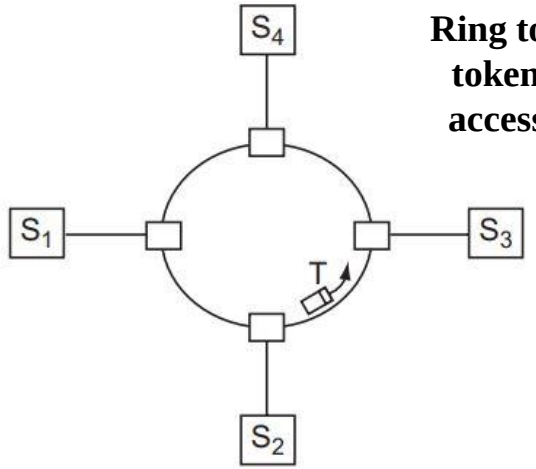
- When a port of the switch receives a frame, it checks the destination address in the frame and then sends the frame to the corresponding port for outgoing data.
- In switch Ethernet, MAC address Table (Forwarding Table) are used to forward the data.
- Bandwidth is shared only between the station and the switch.



MAC Address	Port
A	1
B	2
C	3
D	4
E	5
F	6
G	7
H	8

Q12. What are the advantages of token passing?

- **Token Passing**: A token is regularly passed from one node to another. And if a node has some information to pass on the network, the node releases the information.
- If the node does not have any data to release on the network, then it transfers the token to the next node.



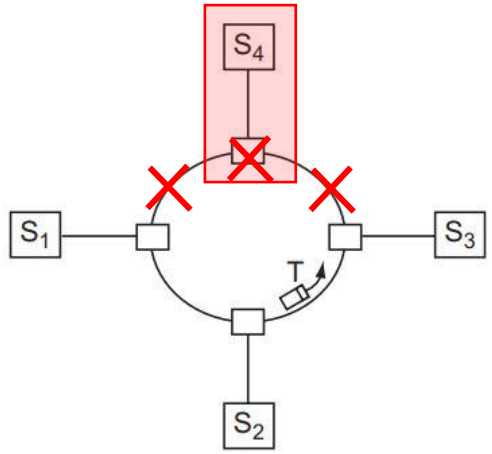
Ring topology in token-passing access method

Advantages of Token Passing

- Reduced chances of data collision as each node release a data packet after receiving the token.
- Token passing makes ring topology perform better than bus topology under heavy traffic.
- No need of server to control connectivity among the nodes.
- The channel bandwidth can be fully utilized when demand is heavy

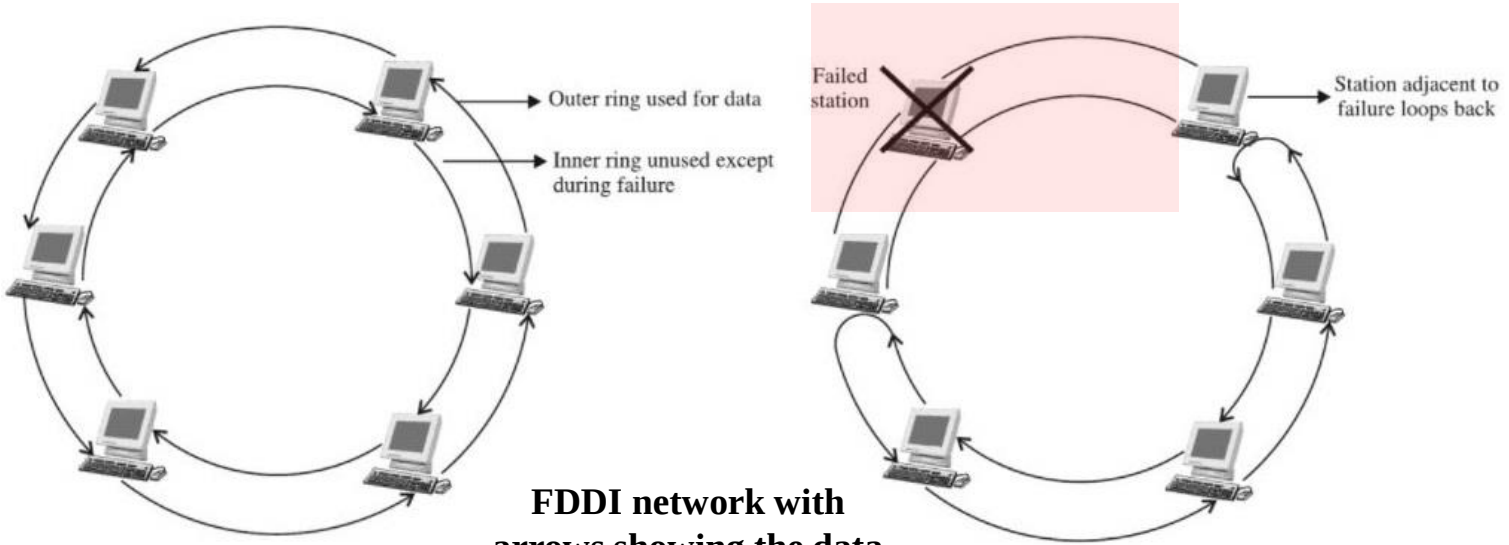
Q13. In the event of failure, how does the FDDI network automatically reconfigure?

- The problem with the **Token Ring** topology is that if one of the links-the medium between two adjacent stations fails, the whole system fails.
- To solve this problem, dual ring topology is used.



- **One of the dual ring token topology is Fiber Distributed Data Interface (FDDI).**
- **FDDI networks called counter rotating as direction of data flows between two rings are opposite.**

- The primary ring carries data, while the secondary ring is usually reserved as a backup in case a fault occurs on the primary ring.
- In the event of a failure on the primary ring, FDDI automatically reconfigures itself to use the secondary ring.
- This process of reconfiguring is called self-healing and FDDI is known as a self-healing network.



FDDI network with arrows showing the data flowing directions

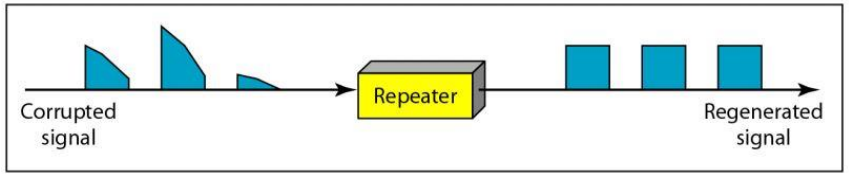
Q14. Compare and contrast between a hub (repeater) and a bridge.

Repeater

- Hub or Repeater operates at the physical layer of OSI model.
- Receives a signal and before it becomes too weak or corrupted, regenerates the original bit pattern.
- It is used to extend the physical length of a network.
- Repeater do not understand complete frame.
- Repeater cannot perform packet filtering.

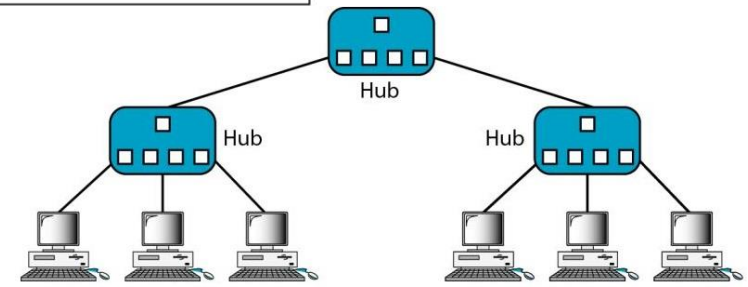
Hub

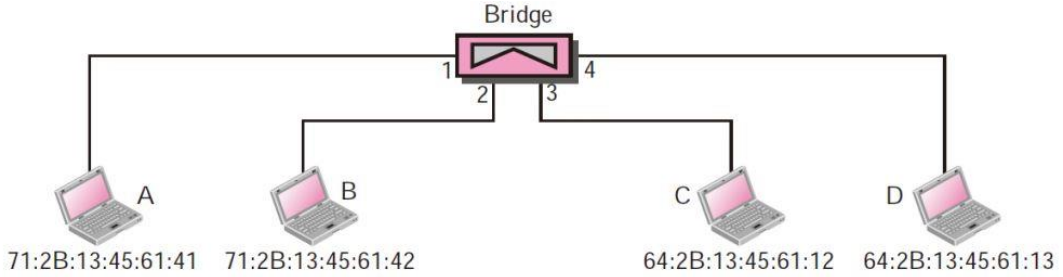
- A Hub is actually a multiport repeater.
- Broadcasts the message sent by one station to all other stations.
- It is most commonly used to connect multiple machines to the same LAN.



Bridge

- Bridge operates at the data link layer of OSI model.
- It can check the destination address of a frame and decide if the frame should be forwarded or discarded.
- It is mostly used in LANs.
- The complete frames is understood by bridge.
- Filtering of the packets in network is achieved by the bridge.





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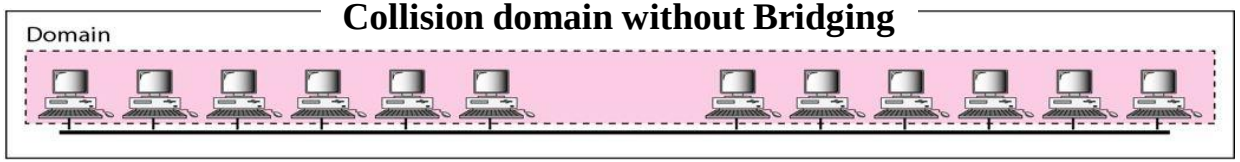
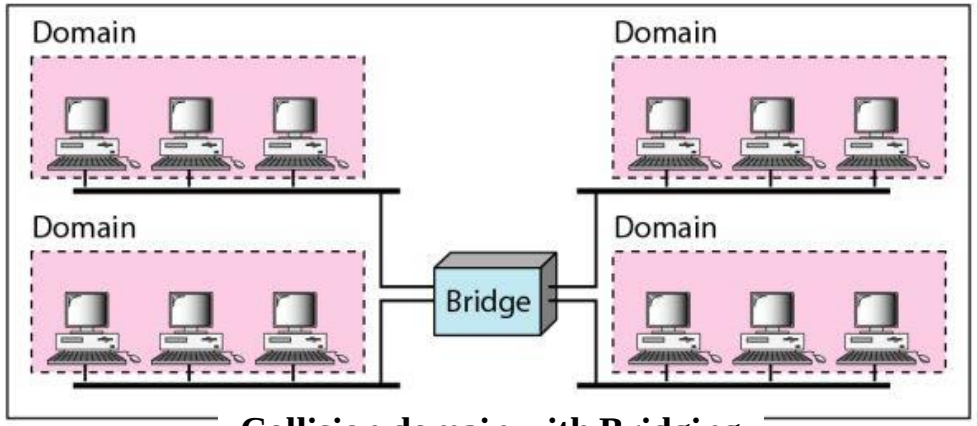
Q15. Compare and contrast between a bridge and a switch (Layer 2).

Bridge

- Bridge operated in the same LAN and performs relatively slow.
- Bridge has limited number of ports.
- Method of switching of a Bridge is store and forward.
- The collision domain becomes smaller

Collision domain without Bridging



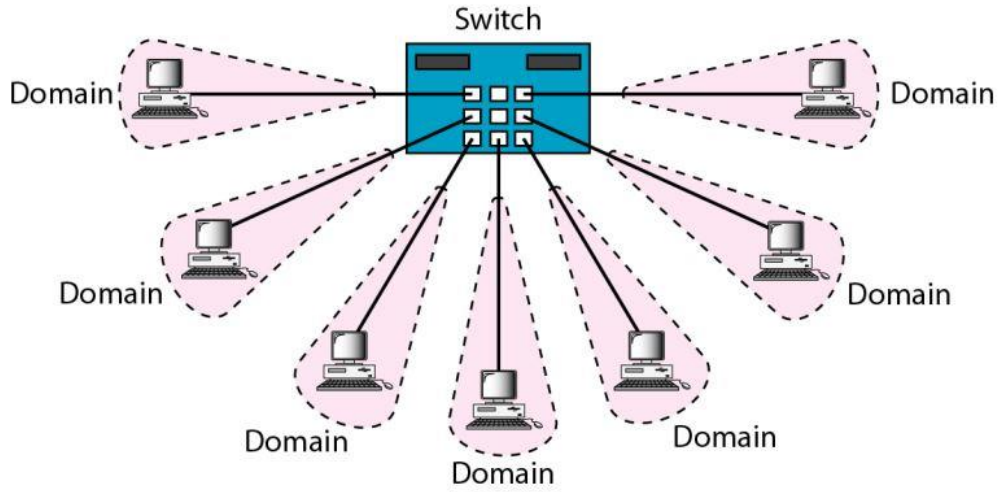


Collision domain with Bridging

Switch (Layer 2)

- A layer 2 switch is an N-port bridge with additional sophistication that allows faster handling of the packets.
- Switch can have many ports.
- Method of switching of a Switch can be store and forward, cut-through, or fragment-free.
- The probability of collision is reduced tremendously

Collision domain with Switch



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Q15. Compare and contrast between a bridge and a switch (Layer 2).

Bridge & Switch (Layer 2**)**

- In general, a switch is a multi-port bridge.
- Both operate in physical and data link layers .
- Both make filtering decisions based on the physical address of the frame they receive.
- Both Bridge and Switch have forwarding table of MAC address and port number.