- Was supposed to be an easy task.
 - Turned out: not so much.
- Linux, 32-bit, no NX, no PIE, no RELRO, stdin/stdout redirected.
- Simple echo logic in a loop:
 - Read uint32 (N) from input.
 - Read N bytes from input to static buffer.
 - Write the N bytes back.

```
void ReadBytes(char *buffer, unsigned int bytes_no) {
  unsigned int bytes_read = 0;

while (bytes_read != bytes_no) {
    ssize_t ret = read(STDIN_FILENO, &buffer[bytes_read], bytes_no - bytes_read);
    if (ret == 0) {
        exit(1);
    }

    bytes_read += ret;
}
```

- Typical overflow with controlled size, but there is nothing to overwrite after the statically located buffer - dead end.
- However, return value of read() is not fully checked.
 - May return -1 when an error occurs.
 - Example: when the syscall tries to write outside mapped memory.
 - When error code is returned, the internal position in the fd does not change; a subsequent read() will return the same data.
 - In our task, a return value of -1 results in shifting the pointer backwards (into the .got.plt section).

```
bytes_no = address of .data.buffer - .got.plt.write

data = dd(.got.plt.write + 4) + shellcode
```

- read() will read the data to &buffer[0], &buffer[-1], &buffer[-2] ... until it reaches the .got.plt.write function pointer.
- bytes_read == bytes_no is satisfied, the loop exits, overwritten write pointer is called.
- No PIE or NX, so you can just have shellcode directly in the static memory and point to it in the function pointer.

- Spawning /bin/sh directly doesn't work too well (probably due to buffered stdin input), it was easiest for us to just execute an open() + read() + write() shellcode.
- Flag:

```
DrgnS{Here's_an_easy_flag_for_an_easy_pwn}
```