

# OPERATOR PRECEDENCE PARSING

CSE 340 FALL 2022

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Notes based on the  
“Dragon Book”

# Parsing Operator Grammars

The grammar we have seen for expressions does not include the operator minus ('-').

This is not an oversight!

We can write the following grammar

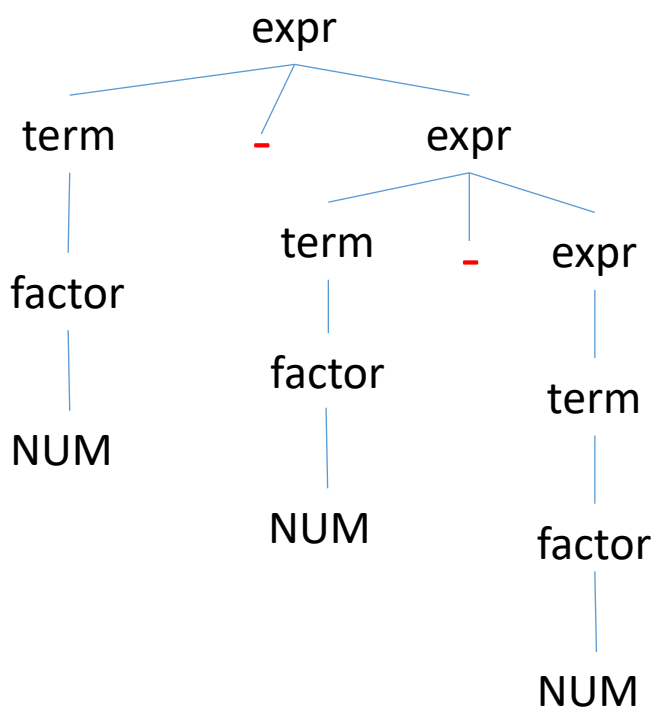
Expr  $\rightarrow$  term - Expr  
Expr  $\rightarrow$  term + Expr  
Expr  $\rightarrow$  term

but that would not work!

How do we parse the following?

1 - 2 - 3

According to the grammar above, we get



According to this tree,

1 - 2 - 3 = 2 !!!

# Parsing expressions with minus

The issue is that minus is left associative and the grammar treats minus as right-associative

Left associative grouping (**correct**)

$$1 - 2 - 3$$

$$(1 - 2) - 3$$

$$((1-2) - 3)$$

Right associative grouping (**wrong**)

$$1 - 2 - 3$$

$$1 - (2 - 3)$$

$$(1 - (2 - 3))$$

When we say that right associative grouping is wrong we mean that it is not according to the convention we adopted, not that there is something inherently wrong with it.

We could have decide the other way around, but, once decided, we should follow the adopted convention.

# Parsing expressions with minus

We can attempt to fix the problem by using the following grammar

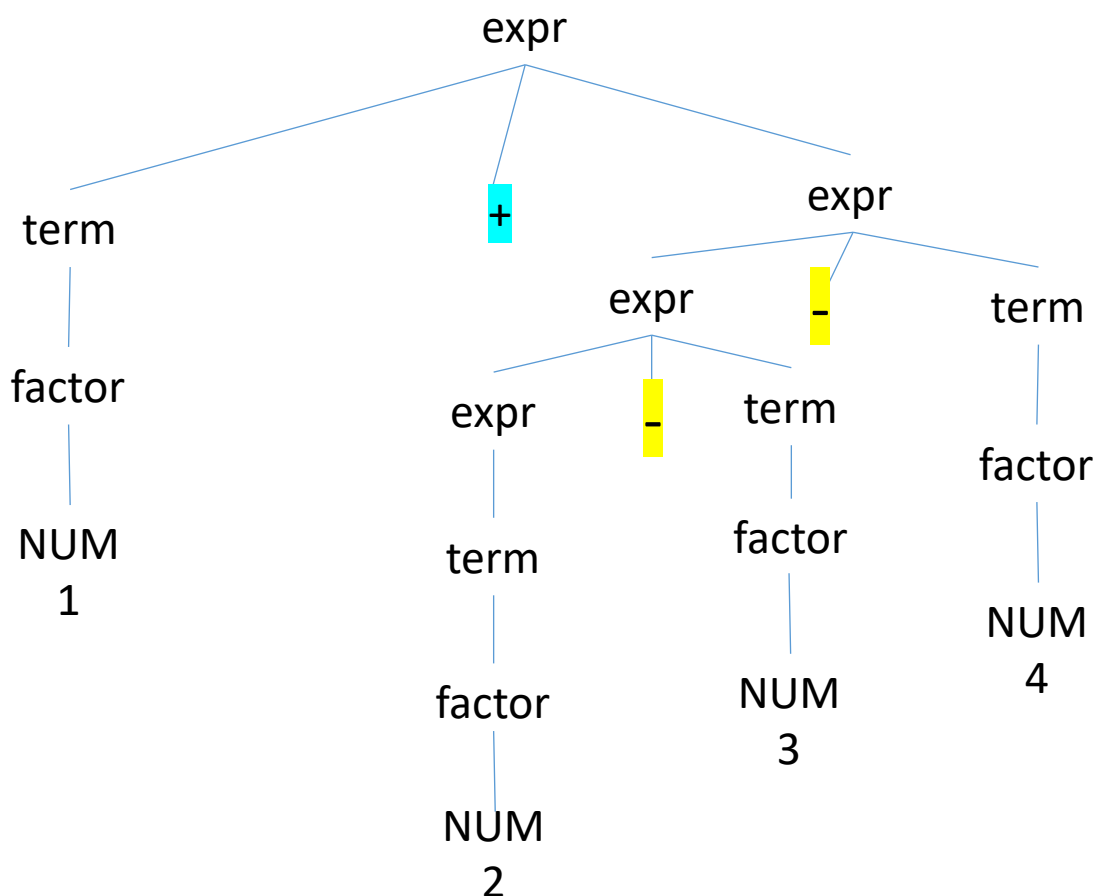
$\text{expr} \rightarrow \text{expr} - \text{term}$

$\text{expr} \rightarrow \text{term} + \text{expr}$

$\text{expr} \rightarrow \text{term}$

This grammar would give the following parsing for

1 + 2 - 3 - 4



# Parsing expressions with minus

We can attempt to fix the problem by using the following grammar

`expr -> expr - term`

`expr -> term + expr`

`expr -> term`

We cannot parse this grammar with a recursive descent parser!

```
parse_expr()
{
    // expr -> expr - term

    ....

    parse_expr() // infinite loop !!
}
```

We need another way to parse such expressions!

# A NEAT TRICK FROM FORTRAN COMPILER!

$a + b * c - d$

add ( ( ( at the beginning

replace every + with ) ) ) + ( ( (

replace every - with ) ) ) - ( ( (

replace every \* with ) ) \* ( (

replace every ^ with ) ^ (

add ) ) ) at the end

We get

(( ( a ) ) ) + ( ( ( b ) ) ) \* ( ( c ) ) ) - ( ( ( d ) ) )

(( ( a ) ) ) + ( ( ( b ) ) ) \* ( ( c ) ) ) - ( ( ( d ) ) )

Always works !

We can then parse with a simple parser that only has to worry about matching parentheses

# Operator Grammar

A grammar is called an operator grammar if

1. there is no righthand side of a rule which has two adjacent non-terminal
2. there is no rule of the form  $A \rightarrow \varepsilon$

**Example 1**      $E \rightarrow E A E \mid ( E ) \mid -E \mid ID$   
                   $A \rightarrow + \mid - \mid * \mid / \mid ^$

is not an operator grammar because of  $E A E$  has three adjacent non-terminals

**Example 2**      $E \rightarrow E + E \mid E - E \mid E * E \mid E / E$   
                   $\mid E ^ E \mid ( E ) \mid - E \mid ID$

is an operator grammar

# OPERATOR PRECEDENCE RELATIONSHIPS

To parse operator grammar, we first define parsing precedence relationships between the terminals of the grammar

We also introduce a new symbol \$ (end of input)

## parsing precedence relationships

- $<\cdot$  yields precedence to
- $\cdot>$  takes precedence over
- $\doteq$  has the same precedence as

These are not the same as the operator precedence levels.

These are used in guiding the parsing

You can think of  $<\cdot$   $\cdot>$  as matching parentheses that group what appears between them (this should become clearer with the examples)

There is a theory to determine these relationship for an unambiguous operator grammar. We will only look at heuristics for common expressions.

The parsing algorithm assumes that we already have a table that defines these relationships



## EXAMPLE

	+	-	*	/	^	id	(	)	\$
+	·	·	·	·	·	·	·	·	·
-	·	·	·	·	·	·	·	·	·
*	·	·	·	·	·	·	·	·	·
/	·	·	·	·	·	·	·	·	·
^	·	·	·	·	·	·	·	·	·
id	·	·	·	·	·			·	·
(	·	·	·	·	·	·	·	=	
)	·	·	·	·	·			·	·
\$	·	·	·	·	·	·	·		

# Parsing Algorithm

**Input**         $w \$$   
**Output**     abstract syntax tree of  $E$   
**Initially**    stack contains  $\$,$  scanning starts at the start of  $w$

**repeat**

```
if $ is on top of the stack and lexer.peek() = $ // EOF
    return;
else
{
    t = lexer.peek(); b = t.type;           // next token from w
    a = stack.terminalpeek().type;         // terminal at the top of stack
                                           // or just below if top is non-terminal

    if (table[a][b] == '<.') | ( table[a][b] = '=' ) // shift
        t = lexer.getToken();
        stack.push(t)
    else if (table[a][b] == '>') // reduce
    {
        RHS = an empty stack
        repeat
            s = stack.pop() // pop terminals and
                           // non-terminals

            if s is a terminal
                last_popped_term = s
            RHS.push(s)
        until ( ( is_a_terminal(stack.peek() ) and
                  ( table[stack.terminalpeek()][last_popped_term] == '<.' ) )

        if E -> RHS rule exists // RHS calculated above
        {
            reduce E -> RHS
            stack.push( E ) // we can think of E as the
                           // root of subtree for E -> RHS
        }
        else
            syntax_error()
    }
    else
        syntax_error();
}
```

**Note:**

- `stack.peek()` peeks at the symbol at the top of the stack, which could be a terminal or a non-terminal.
- `stack.terminalpeek()` peeks at the terminal closest to the top of the stack.



## EXAMPLE

## Stack

## Input

## Action

\$

$$a + b * (c + d) - a$$
 $\angle$ [illegible]

## EXAMPLE

## Stack

## Input

## Action

shift

\$

a

$$+ b * (c + d) - a$$
 $\angle$ [illegible]

## EXAMPLE

## Stack

\$

\$a

# Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$

## Action

shift

[illegible]

## EXAMPLE

## Stack

\$

\$a

## Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$

## Action

shift

[illegible]

## EXAMPLE

## Stack

\$

\$a

## Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$

## Action

shift

reduce  $E \rightarrow ID$

[illegible]



## EXAMPLE

## Stack

\$

\$a

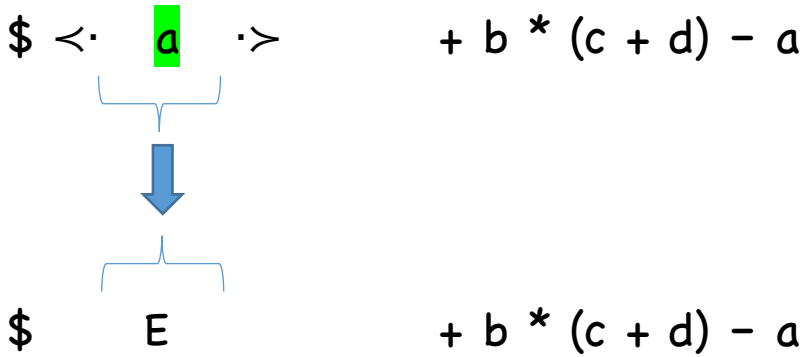
# Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$

## Action

shift

reduce  $E \rightarrow ID$

[illegible]

## EXAMPLE

## Stack

\$

\$a

**\$E**

id

## Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$

## Action

shift

reduce  $E \rightarrow ID$

	+	-	*	/	^	id	(	)	\$
+	·	·	·	·	·	·	·	·	·
-	·	·	·	·	·	·	·	·	·
*	·	·	·	·	·	·	·	·	·
/	·	·	·	·	·	·	·	·	·
^	·	·	·	·	·	·	·	·	·
id	·	·	·	·	·			·	·
(	·	·	·	·	·	·	·	=	
)	·	·	·	·	·			·	·
\$	·	·	·	·	·	·	·		

## EXAMPLE

## Stack

\$

\$a

**\$E**

id

## Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$

## Action

shift

reduce  $E \rightarrow ID$

[illegible]

## EXAMPLE

## Stack

\$

\$a

**\$E**

id

## Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$

## Action

shift

reduce  $E \rightarrow ID$

shift

[illegible]

## EXAMPLE

## Stack

\$

\$a

**\$E**

$\$E_+$

id

## Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$b * (c + d) - a$$

## Action

shift

reduce  $E \rightarrow ID$

shift

[illegible]

## EXAMPLE

## Stack

\$

\$a

**\$E**

**\$E+**

id

# Input

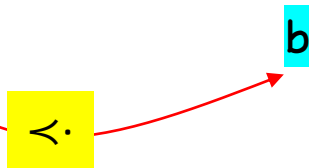
$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$b * (c + d) - a$$

## Action

shift

reduce  $E \rightarrow ID$

shift

[illegible]

## EXAMPLE

## Stack

\$

\$a

**\$E**

**\$E+**

id

# Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$b * (c + d) - a$$

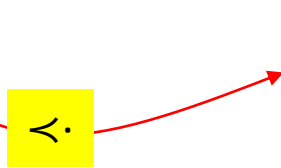
## Action

shift

reduce  $E \rightarrow ID$

shift

shift

[illegible]

## EXAMPLE

## Stack

\$

\$a

**\$E**

$\$E_+$

$\$E+b$

id

## Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$b * (c + d) - a$$
$$* (c + d) - a$$

## Action

shift

reduce  $E \rightarrow ID$

shift

shift

[illegible]



## EXAMPLE

## Stack

\$

\$a

**\$E**

$\$E_+$

$$E + b$$

id

•

## Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$b * (c + d) - a$$

\*  $(c + d) - a$

## Action

shift

reduce  $E \rightarrow ID$

shift

shift

[illegible]

## EXAMPLE

## Stack

\$

\$a

**\$E**

$\$E_+$

$$E + b$$

id

•

## Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$b * (c + d) - a$$

\*  $(c + d) - a$

## Action

shift

reduce  $E \rightarrow ID$

shift

shift

reduce E  $\rightarrow$  ID

[illegible]

## EXAMPLE

## Stack

\$

\$a

**\$E**

$\$E_+$

$\$E + b$

$\$E + E$

id id

## Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$b * (c + d) - a$$
$$* (c + d) - a$$
$$* (c + d) - a$$

## Action

shift

reduce  $E \rightarrow ID$

shift

shift

reduce  $E \rightarrow ID$

[illegible]

## EXAMPLE

## Stack

\$

\$a

**\$E**

$\$E_+$

$\$E + b$

$\$E + E$

id id

 $\prec$ 

# Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$b * (c + d) - a$$
$$* (c + d) - a$$

\*  $(c + d) - a$

## Action

shift

reduce  $E \rightarrow ID$

shift

shift

reduce  $E \rightarrow ID$

[illegible]

## EXAMPLE

## Stack

\$

\$a

**\$E**

$\$E_+$

$\$E+b$

$\$E + E$

 $\prec$ 

# Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$b * (c + d) - a$$
$$* (c + d) - a$$

\*  $(c + d) - a$



## Action

shift

reduce  $E \rightarrow ID$

shift

shift

reduce  $E \rightarrow ID$

shift

[illegible]

## EXAMPLE

## Stack

\$

\$a

**\$E**

$\$E_+$

$\$E + b$

$$\$E + E^*$$

id id

## Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$b * (c + d) - a$$
$$* (c + d) - a$$
$$(c + d) - a$$

## Action

shift

reduce  $E \rightarrow ID$

shift

shift

reduce  $E \rightarrow ID$

[illegible]

## EXAMPLE

## Stack

\$

\$a

**\$E**

$\$E_+$

$\$E + b$

$\$E + E^*$

id id

## Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$b * (c + d) - a$$
$$* (c + d) - a$$
$$(c + d) - a$$

## Action

shift

reduce  $E \rightarrow ID$

shift

shift

reduce  $E \rightarrow ID$

[illegible]

## EXAMPLE

## Stack

\$

\$a

**\$E**

$\$E_+$

$\$E+b$

$\$E + E^*$

id id

id id

 $\prec$ 

# Input

$$a + b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$+ b * (c + d) - a$$
$$b * (c + d) - a$$
$$* (c + d) - a$$
$$(c + d) - a$$

## Action

shift

reduce  $E \rightarrow ID$

shift

shift

reduce  $E \rightarrow ID$

shift

[illegible]



## EXAMPLE

## Stack

# Input

## Action

\$

$$a + b * (c + d) - a$$

shift

\$a

$$+ b * (c + d) - a$$

reduce  $E \rightarrow ID$

**\$E**

$$+ b * (c + d) - a$$

shift

$\$E_+$

$$b * (c + d) - a$$

shift

$\$E+b$

$$* (c + d) - a$$

reduce  $E \rightarrow ID$

$$E + E^*$$
$$(c + d) - a$$

shift

$$E + E^*$$
$$c + d) - a$$

id id

[illegible]

## EXAMPLE

[illegible][illegible]

## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E	$+ b * (c + d) - a$	shift
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift


Diagram illustrating the LR(0) item sets and transitions for the grammar:

- Item set:  $\{E \rightarrow ID\}$  (represented by the yellow box labeled  $< \cdot$ )
- Transitions:
  - Shift transition on  $id$  (represented by the red arrow pointing to the right).
  - Shift transition on  $($  (represented by the red arrow pointing to the right).

[illegible]

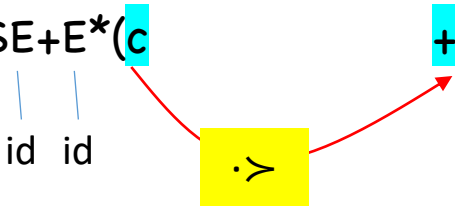
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\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
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\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	

[illegible]

## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
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\$E+	$b * (c + d) - a$	shift
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[illegible]

## EXAMPLE

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\$E	$+ b * (c + d) - a$	shift
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(c	$+ d) - a$	reduce $E \rightarrow ID$

[illegible]

## EXAMPLE

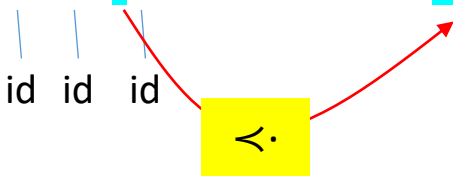
Stack	Input	Action
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\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E	$+ b * (c + d) - a$	shift
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(E	$+ d) - a$	

Diagram illustrating three vertical lines, each labeled 'id' below it.

	+	-	*	/	^	id	(	)	\$
+	→	→	→	→	→	→	→	→	→
-	→	→	→	→	→	→	→	→	→
*	→	→	→	→	→	→	→	→	→
/	→	→	→	→	→	→	→	→	→
^	→	→	→	→	→	→	→	→	→
id	→	→	→	→	→			→	→
(	→	→	→	→	→	→	→	≠	
)	→	→	→	→	→			→	→
\$	→	→	→	→	→	→	→		

## EXAMPLE

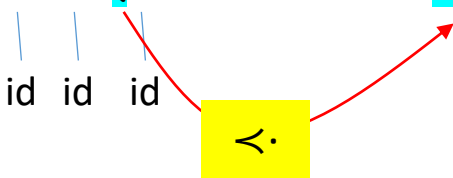
Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E	$+ b * (c + d) - a$	shift
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(E	$+ d) - a$	

[illegible]




## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E	$+ b * (c + d) - a$	shift
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(E	$+ d) - a$	shift

[illegible]

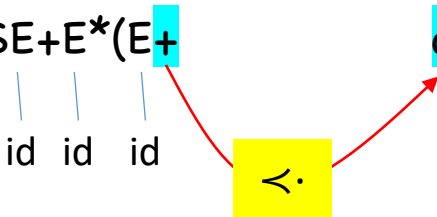
## EXAMPLE

Stack	Input	Action
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\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E	$+ b * (c + d) - a$	shift
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(E+	$d) - a$	

[illegible]

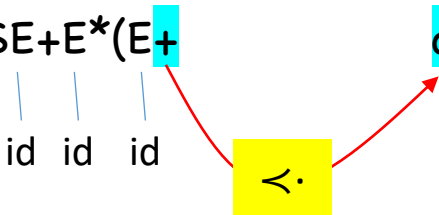
## EXAMPLE

Stack	Input	Action
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\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E	$+ b * (c + d) - a$	shift
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(E+	$d) - a$	

[illegible]


## EXAMPLE

Stack	Input	Action
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\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(E+	$d) - a$	shift

[illegible]

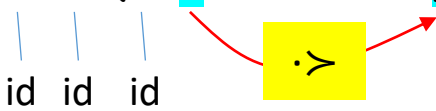
## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
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\$E	$+ b * (c + d) - a$	shift
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(E+d	$) - a$	

[illegible]

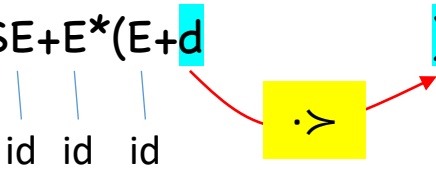
## EXAMPLE

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\$E	$+ b * (c + d) - a$	shift
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(E+d	$) - a$	

[illegible]

## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E	$+ b * (c + d) - a$	shift
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(E+d	$) - a$	reduce $E \rightarrow ID$

[illegible]

## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E	$+ b * (c + d) - a$	shift
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(E+E	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	

id

id

id

id

[illegible]



## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E	$+ b * (c + d) - a$	shift
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(E+E	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	

Diagram illustrating the LR(0) item sets and transitions for the grammar:

- Item sets are shown in boxes, with the current item set highlighted in yellow.
- Transitions are indicated by arrows:
  - Red arrows represent shift transitions (moving the input symbol to the stack).
  - Blue arrows represent reduce transitions (reducing the stack using a grammar rule).

[illegible]

## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E	$+ b * (c + d) - a$	shift
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(E+E	$) - a$	reduce $E \rightarrow ID$
\$E+E*(E+E	$) - a$	reduce $E \rightarrow E+E$

Diagram illustrating the LR(0) item sets and transitions for the grammar:

- Item sets are shown as states:  $S$ ,  $Sa$ ,  $SE$ ,  $SE+$ ,  $SE+b$ ,  $SE+E*$ ,  $SE+E*($ ,  $SE+E*(E+E$ , and  $SE+E*(E+E$ .
- Transitions are labeled with grammar symbols:  $a$ ,  $b$ ,  $($ ,  $)$ ,  $+$ ,  $*$ , and  $-$ .
- Red arrows indicate the transitions from the initial state  $S$  to  $Sa$  and from  $SE+E*(E+E$  to  $SE+E*(E+E$ .

[illegible]

# EXAMPLE

Stack

\$

\$a

\$E

\$E+

\$E+b

\$E+E\*

\$E+E\*(

\$E+E\*(E+E

\$E+E\*( E

id id

+  
id id

Input

a + b \* (c + d) - a

+ b \* (c + d) - a

+ b \* (c + d) - a

b \* (c + d) - a

\* (c + d) - a

(c + d) - a

c + d) - a

) - a

) - a

Action

shift

reduce E → ID

shift

shift

reduce E → ID

shift

shift

reduce E → ID

# When do we stop popping the stack?

\$ E + E \* ( E + E )

\$ <· + <· \* <· ( **E** + **E** )



when the E + E is popped, the following holds

1. the top of the stack is a terminal which is (
2. the last popped terminal is +
3. ( <· + so we stop

So, we keep on popping until

- (1) the top of stack is a terminal and
- (2) top\_of\_stack\_symbol <· last\_popped\_terminal

what is popped is between a pair <· and ·>:

<· RHS of reduction ·>

the **E** is also popped because we cannot stop with the **E** at the top of stack because it is not a terminal

# EXAMPLE

Stack

\$

\$a

\$E

\$E+

\$E+b

\$E+E\*

\$E+E\*(

\$E+E\*(E+E

\$E+E\*( E

id id

+  
id id

Input

a + b \* (c + d) - a

+ b \* (c + d) - a

+ b \* (c + d) - a

b \* (c + d) - a

\* (c + d) - a

(c + d) - a

c + d) - a

) - a

) - a

Action

shift

reduce E -> ID

shift

shift

reduce E -> ID

shift

shift

reduce E -> ID

# EXAMPLE

Stack

\$

\$a

\$E

\$E+

\$E+b

\$E+E\*

\$E+E\*(

\$E+E\*(E+E

\$E+E\*( E

id id

id id

Input

a + b \* (c + d) - a

+ b \* (c + d) - a

+ b \* (c + d) - a

b \* (c + d) - a

\* (c + d) - a

(c + d) - a

c + d) - a

) - a

) - a

Action

shift

reduce E -> ID

shift

shift

reduce E -> ID

shift

shift

reduce E -> ID



# EXAMPLE

Stack

\$

\$a

\$E

\$E+

\$E+b

\$E+E\*

\$E+E\*(

\$E+E\*(E+E

\$E+E\*( E

id id

id id

Input

a + b \* (c + d) - a

+ b \* (c + d) - a

+ b \* (c + d) - a

b \* (c + d) - a

\* (c + d) - a

(c + d) - a

c + d) - a

) - a

) - a

Action

shift

reduce E -> ID

shift

shift

reduce E -> ID

shift

shift

reduce E -> ID

shift



# EXAMPLE

Stack

\$

\$a

\$E

\$E+

\$E+b

\$E+E\*

\$E+E\*(

\$E+E\*(E+E

\$E+E\*( E

\$E+E\*( E )

id id

+

id id

Input

a + b \* (c + d) - a

+ b \* (c + d) - a

+ b \* (c + d) - a

b \* (c + d) - a

\* (c + d) - a

(c + d) - a

c + d) - a

) - a

) - a

- a

Action

shift

reduce E -> ID

shift

shift

reduce E -> ID

shift

shift

reduce E -> ID

shift



# EXAMPLE

Stack

\$

\$a

\$E

\$E+

\$E+b

\$E+E\*

\$E+E\*(

\$E+E\*(E+E

\$E+E\*( E

\$E+E\*( E )

id id

+

id id

Input

a + b \* (c + d) - a

+ b \* (c + d) - a

+ b \* (c + d) - a

b \* (c + d) - a

\* (c + d) - a

(c + d) - a

c + d) - a

) - a

) - a

- a

Action

shift

reduce E → ID

shift

shift

reduce E → ID

shift

shift

reduce E → ID

shift

· >

# EXAMPLE

Stack

\$

\$a

\$E

\$E+

\$E+b

\$E+E\*

\$E+E\*(

\$E+E\*(E+E

\$E+E\*( E

\$E+E\*( E )

id id

+

id id

Input

a + b \* (c + d) - a

+ b \* (c + d) - a

+ b \* (c + d) - a

b \* (c + d) - a

\* (c + d) - a

(c + d) - a

c + d) - a

) - a

) - a

- a

Action

shift

reduce E -> ID

shift

shift

reduce E -> ID

shift

shift

reduce E -> ID

shift

reduce E -> ( E )

.>

# EXAMPLE

Stack

Input

Action

\$

$a + b * (c + d) - a$

shift

\$a

$+ b * (c + d) - a$

reduce  $E \rightarrow ID$

\$E

$+ b * (c + d) - a$

shift

\$E+

$b * (c + d) - a$

shift

\$E+b

$* (c + d) - a$

reduce  $E \rightarrow ID$

\$E+E\*

$(c + d) - a$

shift

\$E+E\*(

$c + d) - a$

shift

\$E+E\*(E+E

$) - a$

reduce  $E \rightarrow ID$

\$E+E\*( E

$) - a$

shift

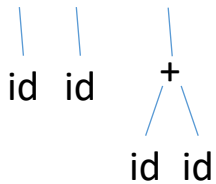
\$E+E\*( E )

$- a$

reduce  $E \rightarrow ( E )$

\$E+E \* E

$- a$



# EXAMPLE

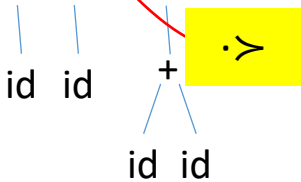
Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E	$+ b * (c + d) - a$	shift
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(E+E	$) - a$	reduce $E \rightarrow ID$
\$E+E*( E	$) - a$	shift
\$E+E*( E )	$- a$	reduce $E \rightarrow ( E )$
\$E+E * E	$- a$	

Diagram illustrating the state  $E + E * E$  and the next input token  $- a$ . A yellow box contains the LR item  $E \rightarrow E + E$ . A red arrow points from the  $*$  in the stack to the  $-$  in the input, indicating a shift action. Blue lines connect the  $id$  labels to the  $E$  non-terminals in the stack.

# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E	$+ b * (c + d) - a$	shift
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(E+E	$) - a$	reduce $E \rightarrow ID$
\$E+E*( E	$) - a$	shift
\$E+E*( E )	$- a$	reduce $E \rightarrow ( E )$
\$E+E * E	$- a$	reduce $E \rightarrow E * E$



# EXAMPLE

Stack

Input

Action

\$

a + b \* (c + d) - a

shift

\$a

+ b \* (c + d) - a

reduce E → ID

\$E

+ b \* (c + d) - a

shift

\$E+

b \* (c + d) - a

shift

\$E+b

\* (c + d) - a

reduce E → ID

\$E+E\*

(c + d) - a

shift

\$E+E\*(

c + d) - a

shift

\$E+E\*(E+E

) - a

reduce E → ID

\$E+E\*( E

) - a

shift

\$E+E\*( E )

- a

reduce E → ( E )

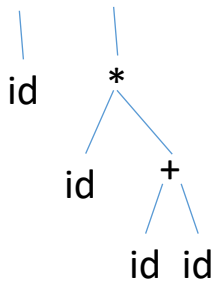
\$E+E \* E

- a

reduce E → E \* E

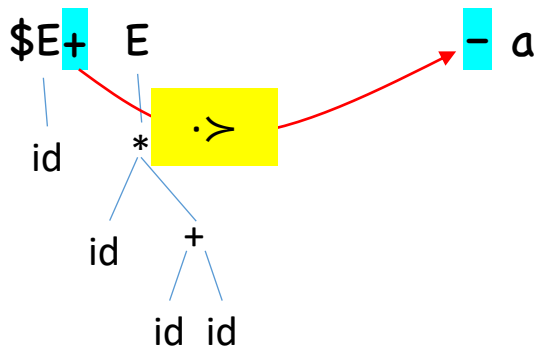
\$E+ E

- a



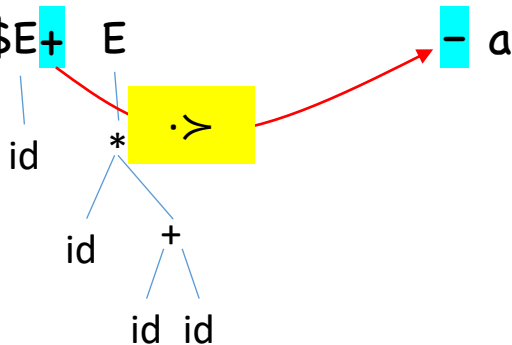
# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E	$+ b * (c + d) - a$	shift
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(E+E	$) - a$	reduce $E \rightarrow ID$
\$E+E*( E	$) - a$	shift
\$E+E*( E )	$- a$	reduce $E \rightarrow ( E )$
\$E+E * E	$- a$	reduce $E \rightarrow E * E$



# EXAMPLE

Stack	Input	Action
\$	a + b * (c + d) - a	shift
\$a	+ b * (c + d) - a	reduce E → ID
\$E	+ b * (c + d) - a	shift
\$E+	b * (c + d) - a	shift
\$E+b	* (c + d) - a	reduce E → ID
\$E+E*	(c + d) - a	shift
\$E+E*(	c + d) - a	shift
\$E+E*(E+E	) - a	reduce E → ID
\$E+E*( E	) - a	shift
\$E+E*( E )	- a	reduce E → ( E )
\$E+E * E	- a	reduce E → E * E
\$E+ E	- a	reduce E → E + E





# EXAMPLE

Stack

Input

Action

\$

$a + b * (c + d) - a$

shift

\$a

$+ b * (c + d) - a$

reduce  $E \rightarrow ID$

\$E

$+ b * (c + d) - a$

shift

\$E+

$b * (c + d) - a$

shift

\$E+b

$* (c + d) - a$

reduce  $E \rightarrow ID$

\$E+E\*

$(c + d) - a$

shift

\$E+E\*(

$c + d) - a$

shift

\$E+E\*(E+E

$) - a$

reduce  $E \rightarrow ID$

\$E+E\*( E

$) - a$

shift

\$E+E\*( E )

$- a$

reduce  $E \rightarrow ( E )$

\$E+E \* E

$- a$

reduce  $E \rightarrow E * E$

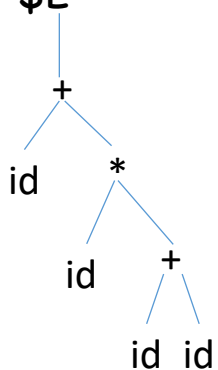
\$E+ E

$- a$

reduce  $E \rightarrow E + E$

\$E

$- a$



# EXAMPLE

Stack

Input

Action

\$

a + b \* (c + d) - a

shift

\$a

+ b \* (c + d) - a

reduce E → ID

\$E

+ b \* (c + d) - a

shift

\$E+

b \* (c + d) - a

shift

\$E+b

\* (c + d) - a

reduce E → ID

\$E+E\*

(c + d) - a

shift

\$E+E\*(

c + d) - a

shift

\$E+E\*(E+E

) - a

reduce E → ID

\$E+E\*( E

) - a

shift

\$E+E\*( E )

- a

reduce E → ( E )

\$E+E \* E

- a

reduce E → E \* E

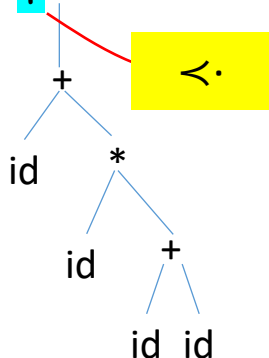
\$E+ E

- a

reduce E → E + E

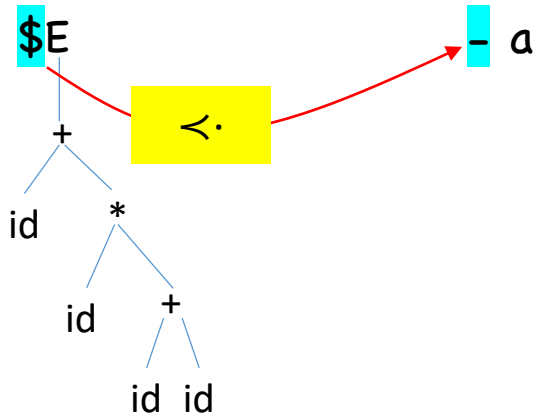
\$E

- a



# EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E	$+ b * (c + d) - a$	shift
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(E+E	$) - a$	reduce $E \rightarrow ID$
\$E+E*( E	$) - a$	shift
\$E+E*( E )	$- a$	reduce $E \rightarrow ( E )$
\$E+E * E	$- a$	reduce $E \rightarrow E * E$
\$E+ E	$- a$	reduce $E \rightarrow E + E$
\$E	$- a$	shift



# EXAMPLE

Stack

Input

Action

\$

a + b \* (c + d) - a

shift

\$a

+ b \* (c + d) - a

reduce E → ID

\$E

+ b \* (c + d) - a

shift

\$E+

b \* (c + d) - a

shift

\$E+b

\* (c + d) - a

reduce E → ID

\$E+E\*

(c + d) - a

shift

\$E+E\*(

c + d) - a

shift

\$E+E\*(E+E

) - a

reduce E → ID

\$E+E\*( E

) - a

shift

\$E+E\*( E )

- a

reduce E → ( E )

\$E+E \* E

- a

reduce E → E \* E

\$E+ E

- a

reduce E → E + E

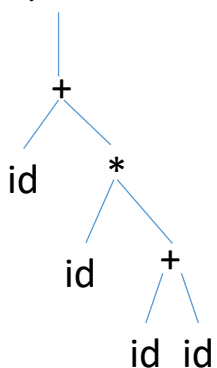
\$E

- a

shift

\$E-

a



# EXAMPLE

Stack

Input

Action

\$

a + b \* (c + d) - a

shift

\$a

+ b \* (c + d) - a

reduce E → ID

\$E

+ b \* (c + d) - a

shift

\$E+

b \* (c + d) - a

shift

\$E+b

\* (c + d) - a

reduce E → ID

\$E+E\*

(c + d) - a

shift

\$E+E\*(

c + d) - a

shift

\$E+E\*(E+E

) - a

reduce E → ID

\$E+E\*( E

) - a

shift

\$E+E\*( E )

- a

reduce E → ( E )

\$E+E \* E

- a

reduce E → E \* E

\$E+ E

- a

reduce E → E + E

\$E

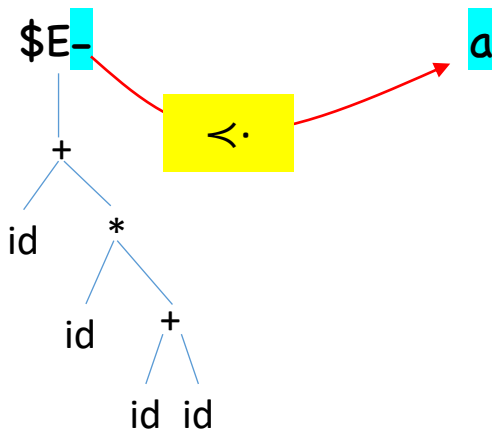
- a

shift

\$E-

a

<.



# EXAMPLE

Stack

Input

Action

\$

a + b \* (c + d) - a

shift

\$a

+ b \* (c + d) - a

reduce E → ID

\$E

+ b \* (c + d) - a

shift

\$E+

b \* (c + d) - a

shift

\$E+b

\* (c + d) - a

reduce E → ID

\$E+E\*

(c + d) - a

shift

\$E+E\*(

c + d) - a

shift

\$E+E\*(E+E

) - a

reduce E → ID

\$E+E\*( E

) - a

shift

\$E+E\*( E )

- a

reduce E → ( E )

\$E+E \* E

- a

reduce E → E \* E

\$E+ E

- a

reduce E → E + E

\$E

- a

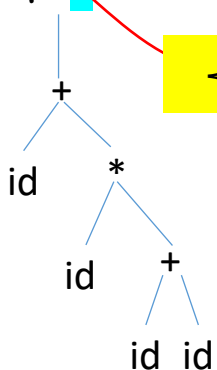
shift

\$E-

a

shift

<.



# EXAMPLE

Stack

Input

Action

\$

$a + b * (c + d) - a$

shift

\$a

$+ b * (c + d) - a$

reduce  $E \rightarrow ID$

\$E

$+ b * (c + d) - a$

shift

\$E+

$b * (c + d) - a$

shift

\$E+b

$* (c + d) - a$

reduce  $E \rightarrow ID$

\$E+E\*

$(c + d) - a$

shift

\$E+E\*(

$c + d) - a$

shift

\$E+E\*(E+E

$) - a$

reduce  $E \rightarrow ID$

\$E+E\*( E

$) - a$

shift

\$E+E\*( E )

$- a$

reduce  $E \rightarrow ( E )$

\$E+E \* E

$- a$

reduce  $E \rightarrow E * E$

\$E+ E

$- a$

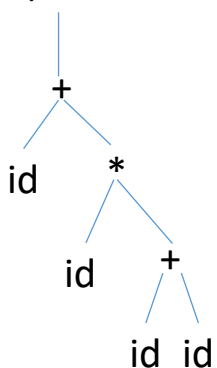
reduce  $E \rightarrow E + E$

\$E

$- a$

shift

\$E-a



# EXAMPLE

Stack

Input

Action

\$

a + b \* (c + d) - a

shift

\$a

+ b \* (c + d) - a

reduce E → ID

\$E

+ b \* (c + d) - a

shift

\$E+

b \* (c + d) - a

shift

\$E+b

\* (c + d) - a

reduce E → ID

\$E+E\*

(c + d) - a

shift

\$E+E\*(

c + d) - a

shift

\$E+E\*(E+E

) - a

reduce E → ID

\$E+E\*( E

) - a

shift

\$E+E\*( E )

- a

reduce E → ( E )

\$E+E \* E

- a

reduce E → E \* E

\$E+ E

- a

reduce E → E + E

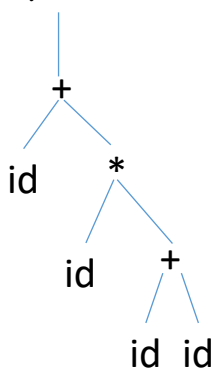
\$E

- a

shift

\$E-a

\$





# EXAMPLE

Stack

Input

Action

\$

a + b \* (c + d) - a

shift

\$a

+ b \* (c + d) - a

reduce E → ID

\$E

+ b \* (c + d) - a

shift

\$E+

b \* (c + d) - a

shift

\$E+b

\* (c + d) - a

reduce E → ID

\$E+E\*

(c + d) - a

shift

\$E+E\*(

c + d) - a

shift

\$E+E\*(E+E

) - a

reduce E → ID

\$E+E\*( E

) - a

shift

\$E+E\*( E )

- a

reduce E → ( E )

\$E+E \* E

- a

reduce E → E \* E

\$E+ E

- a

reduce E → E + E

\$E

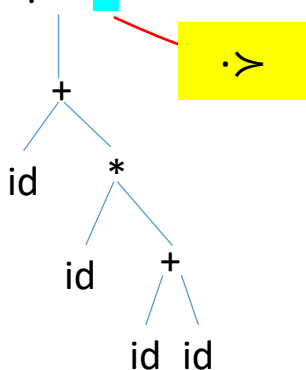
- a

shift

\$E-a

\$

.>



# EXAMPLE

Stack

Input

Action

\$

a + b \* (c + d) - a

shift

\$a

+ b \* (c + d) - a

reduce E → ID

\$E

+ b \* (c + d) - a

shift

\$E+

b \* (c + d) - a

shift

\$E+b

\* (c + d) - a

reduce E → ID

\$E+E\*

(c + d) - a

shift

\$E+E\*(

c + d) - a

shift

\$E+E\*(E+E

) - a

reduce E → ID

\$E+E\*( E

) - a

shift

\$E+E\*( E )

- a

reduce E → ( E )

\$E+E \* E

- a

reduce E → E \* E

\$E+ E

- a

reduce E → E + E

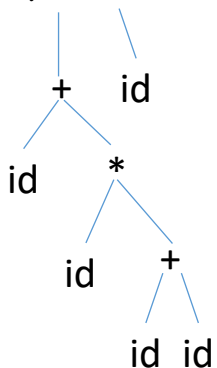
\$E

- a

shift

\$E-E

\$



# EXAMPLE

Stack

Input

Action

\$

a + b \* (c + d) - a

shift

\$a

+ b \* (c + d) - a

reduce E → ID

\$E

+ b \* (c + d) - a

shift

\$E+

b \* (c + d) - a

shift

\$E+b

\* (c + d) - a

reduce E → ID

\$E+E\*

(c + d) - a

shift

\$E+E\*(

c + d) - a

shift

\$E+E\*(E+E

) - a

reduce E → ID

\$E+E\*( E

) - a

shift

\$E+E\*( E )

- a

reduce E → ( E )

\$E+E \* E

- a

reduce E → E \* E

\$E+ E

- a

reduce E → E + E

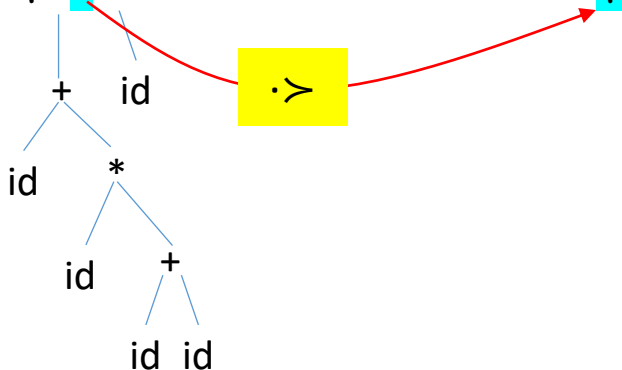
\$E

- a

shift

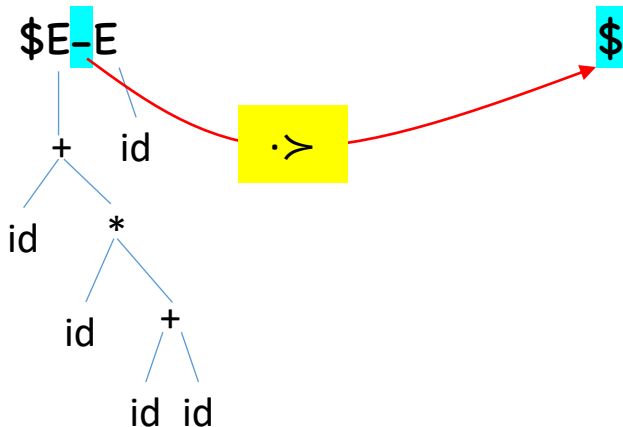
\$E-E

\$



## EXAMPLE

Stack	Input	Action
\$	$a + b * (c + d) - a$	shift
\$a	$+ b * (c + d) - a$	reduce $E \rightarrow ID$
\$E	$+ b * (c + d) - a$	shift
\$E+	$b * (c + d) - a$	shift
\$E+b	$* (c + d) - a$	reduce $E \rightarrow ID$
\$E+E*	$(c + d) - a$	shift
\$E+E*(	$c + d) - a$	shift
\$E+E*(E+E	$) - a$	reduce $E \rightarrow ID$
\$E+E*( E	$) - a$	shift
\$E+E*( E )	$- a$	reduce $E \rightarrow ( E )$
\$E+E * E	$- a$	reduce $E \rightarrow E * E$
\$E+ E	$- a$	reduce $E \rightarrow E + E$
\$E	$- a$	shift
\$E-E	\$	reduce $E \rightarrow E - E$



# EXAMPLE

Stack

Input

Action

\$

a + b \* (c + d) - a

shift

\$a

+ b \* (c + d) - a

reduce E → ID

\$E

+ b \* (c + d) - a

shift

\$E+

b \* (c + d) - a

shift

\$E+b

\* (c + d) - a

reduce E → ID

\$E+E\*

(c + d) - a

shift

\$E+E\*(

c + d) - a

shift

\$E+E\*(E+E

) - a

reduce E → ID

\$E+E\*( E

) - a

shift

\$E+E\*( E )

- a

reduce E → ( E )

\$E+E \* E

- a

reduce E → E \* E

\$E+ E

- a

reduce E → E + E

\$E

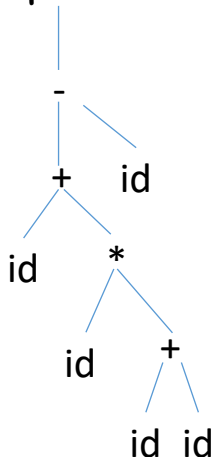
- a

shift

\$E

\$

reduce E → E - E



# EXAMPLE

Stack

Input

Action

\$

a + b \* (c + d) - a

shift

\$a

+ b \* (c + d) - a

reduce E → ID

\$E

+ b \* (c + d) - a

shift

\$E+

b \* (c + d) - a

shift

\$E+b

\* (c + d) - a

reduce E → ID

\$E+E\*

(c + d) - a

shift

\$E+E\*(

c + d) - a

shift

\$E+E\*(E+E

) - a

reduce E → ID

\$E+E\*( E

) - a

shift

\$E+E\*( E )

- a

reduce E → ( E )

\$E+E \* E

- a

reduce E → E \* E

\$E+ E

- a

reduce E → E + E

\$E

- a

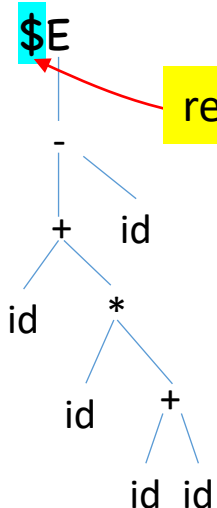
shift

\$E

\$

reduce E → E - E

return



# Dealing with non-terminals

We have the non-terminals on the stack as they are pushed when a reduction occurs

Given that the grammar is an operator grammar, we can have at most one non-terminal on the top of the stack. There is always a terminal on the top of the stack or just below

In the algorithm we assume that `stack.terminalpeek()` ignores non-terminals and returns the terminal symbol at the top of the stack or just below

# HEURISTIC FOR DETERMINING PRECEDENCE RELATIONSHIPS

We assume we have a set of operators with

- precedence levels
- associativity (left or right)
- operators at the same level have the same associativity

We assume the input is of the form  $w \$$

We have the following heuristics for determining  $\prec$ ,  $\succ$ , and  $\doteq$  relationships between operators and terminals

1. if  $op1$  has higher precedence level than  $op2$ , then

- $op1 \succ op2$
- $op2 \prec op1$

**Example:**  $* \succ +$

$+ \prec *$

$\wedge \succ +$

$+ \prec \wedge$



# HEURISTIC FOR DETERMINING PRECEDENCE RELATIONSHIPS

2. if op1 and op2 are operators of the same operator precedence, possibly the same operator, then

- If they are left associative:

- $op1 \cdot \succ op2$

- $op2 \cdot \succ op1$

**Example.** + and - are left associative, so we have

+  $\cdot \succ$  +

+  $\cdot \succ$  -

-  $\cdot \succ$  +

-  $\cdot \succ$  -

- If they are right associative:

- $op1 \prec \cdot op2$

- $op2 \prec \cdot op1$

**Example.** ^ is right associative, so we have

^  $\cdot \prec$  ^

^  $\prec \cdot$  ^

# HEURISTIC FOR DETERMINING PRECEDENCE RELATIONSHIPS

3. Also, we have the following

1.	op	$\prec\cdot$	ID
2.	ID	$\cdot\succ$	op
3.	op	$\prec\cdot$	(
4.	(	$\prec\cdot$	op
5.	op	$\cdot\succ$	)
6.	)	$\cdot\succ$	op
7.	\$	$\prec\cdot$	op
8.	op	$\cdot\succ$	\$
9.	(	$\doteq$	)
10.	\$	$\prec\cdot$	id
11.	id	$\cdot\succ$	\$
12.	\$	$\prec\cdot$	(
13.	(	$\prec\cdot$	(
14.	(	$\prec\cdot$	id
15.	)	$\cdot\succ$	\$
16.	)	$\cdot\succ$	)
17.	id	$\cdot\succ$	)

# Unary Operators (one operand)

If we have a unary operator `uop` **that is not a binary operator**, we can support it as follows

- `op <· uop` for every other operator `op`. `op` can be unary or binary
- `uop <· op` if `uop` has lower operator precedence level than `op`
- `uop ·> op` if `uop` has higher operator precedence level than `op`

If we have a unary operator **that is also a binary operator**, like MINUS, we cannot incorporate it in the scheme!

**Example** `id*-id` is not easily parsed

One solution is to have the `getToken()` function make the distinction by looking at the context in which the operator appears.

**Example** In FORTRAN a minus sign is unary if the previous token is an operator, LPAREN, COMMA, or EQUAL

It is better to handle this in the lexer than it is to make the parser more complicated

# Unary Operators (one operand)

If we have a unary operator uop that is not a binary operator, we can support it as follows

- $op \prec \cdot uop$  for every other operator  $op$ .  $op$  can be unary or binary
- $uop \prec \cdot op$  if  $uop$  has lower operator precedence level than  $op$
- $uop \succ \cdot op$  if  $uop$  has higher operator precedence level than  $op$

If we have a unary operator that is also a binary operator, like MINUS, we cannot incorporate it in the scheme!

**Example**  $id * -id$  is not easily parsed

One solution is to have the `getToken()` function make the distinction by looking at the context in which the operator appears.

**Example** In FORTRAN a minus sign is unary if the previous token is an operator, LPAREN, COMMA, or EQUAL

It is better to handle this in the lexer than it is to make the parser more complicated