

# CS323 Project Midterm Report

## C-like Compilers in Rust

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# **1 Software Stack and Project Goal**

# Software stack and Project Goal



Using `flex` and `bison` requires different tool chain and binding different files together, we are hoping to implement a software tool chain that can integrate Lexer, Parser, Semantic Analysis and the Intermediate Result Generation(Optimization) together.

## 1.a.a Why not Rust?

Rust is famous for its powerful compiler, by using Rust to build a compiler, we can dive more deeply into this amazing language. Also, it is easier to test using Rust by only a simple `cargo test`!

```
rust          1.82.0
logos         0.14.2
lalrpop       0.22.0
llvm_ir       0.11.1 (Support llvm-IR)
```

## 1.a.b Project Goal

We are trying to build a compiler frontend system to generate LLVM-IR result. And use the LLVM backend system, we can build a compiler thoroughly.

## **2 Specification and Core Feature**

# Lexer



By using `Logos`, a Rust-supported lexer, our language supports the following tokens. We are still planning to add more tokens when applying more features.

## Operators

>	OpGreaterThan	=	OpAssign	^	OpPow
<	OpLessThan	+	OpPlus	&&	OpAnd
<=	OpLessThanEqual	-	OpMinus		OpOr
>=	OpGreaterThanEqual	*	OpMul	!	OpNot
==	OpEqual	/	OpDiv	++	OpIncrement
!=	OpNotEqual	%	OpMod	--	OpDecrement

## Punctuation

.	Dot	[	LeftBracket	{	LeftBrace
,	Comma	]	RightBracket	}	RightBrace
:	Colon	(	LeftParen		
;	Semicolon	)	RightParen		

# Lexer (ii)



## Keywords

if	KeywordIf	for	KeywordFor	break	KeywordBreak
else	KeywordElse	return	KeywordReturn	continue	KeywordContinue
while	KeywordWhile				

## Declaration

enum	DeclarationEnum	struct	DeclarationStruct	fn	DeclarationFunction

## Type

bool	TypeBool	char	TypeChar	string	TypeString
int	TypeInt	float	TypeFloat	null	TypeNull

# Lexer (iii)



## Literals

<code>r"true false"</code>	LiteralBool: bool	<code>r"'\'' r"'\u{0-9a-fA-F}{1,6}'"</code>	LiteralChar: char	<code>r"-?[0-9]+" r"-?0[xX][0-9a-fA-F]+"</code>	LiteralInt: i32
<code>r"-?(?:0 [1-9]\d*)? \. \d+(?:[eE][+-])? \d+)"</code>	LiteralFloat: f32	<code>r#"([^\\" \\["\\bnfrt] \\x[0-9a-fA-F]{2} \\u[a-fA-F0-9]{1,6})*"#</code>	LiteralString: String		

## Identifiers

<code>r"[a-zA-Z_\$][a-zA-Z0-9_\$]*"</code>	Identifier: String

## Comment

<code>r"//[^\n]*\n?"</code>	Line Comment
<code>r"/\**"</code>	Block Comment



# Parser



By using `LALRPOP`, a Rust Parser Generator, our language will support a superset of SPL with the following syntax for now. Also, we tuned the original grammar to utilize `LALRPOP`'s features.

Term and Computation Expression (returns a value excluding bool or variable itself)

```
// Below are all left associative for now
CompExpr -> Term | // Priority 0
    CompExpr ^ CompExpr | CompExpr % CompExpr // Priority 1
    CompExpr * CompExpr | CompExpr / CompExpr // Priority 2
    CompExpr + CompExpr | CompExpr - CompExpr // Priority 3
    Identifier OpIncreament | Identifier OpDecreament

Term -> LiteralInt | LiteralFloat | LiteralChar | LiteralString | Identifier |
    LeftParen CompExpr RightParen

Specifier -> TypeInt | TypeFloat | TypeChar | TypeString | Null
// And struct...
```

## Parser (ii)



Conditional Expression (returns a boolean since if/loop only accept condition)

```
CondExpr -> CondTerm |
    CompExpr OpEqual CompExpr |
    CompExpr OpNotEqual CompExpr |
    CompExpr OpLessThan CompExpr |
    CompExpr OpGreaterThan CompExpr |
    CompExpr OpLessThanEqual CompExpr |
    CompExpr OpGreaterThanEqual CompExpr |
    CondExpr OpEqual CondExpr |
    CondExpr OpNotEqual CondExpr |
    CondExpr OpAnd CondExpr |
    CondExpr OpOr CondExpr |
    OpNot CondExpr

CondTerm -> LiteralBool | Identifier | LeftParen CondExpr RightParen
// Only Bool variable
```

## Parser (iii)



### Assignment Statement & Local Declaration

```
AssignStmt -> Identifier OpAssign CompExpr Semicolon |  
            Identifier OpAssign CondExpr Semicolon  
  
Defs -> Def Defs | $ // *  
Def -> Specifier Decs Semicolon  
  
Decs -> Dec | Dec Comma Decs // +  
Dec -> VarDec | VarDec OpAssign CompExpr | VarDec OpAssign CondExpr
```

### Statements (If/else, while/for)

```
Stmts -> Stmt Semicolon Stmts | $ // *  
  
ForInit -> Decs Comma ForInit | $ // *
```

## Parser (iv)



```
Stmt -> AssignStmt | Def | LeftBrace Stmts RightBrace |  
      | CondExpr | CompExpr | // In case someone writes this...  
      | KeywordReturn CompExpr Semicolon  
      | KeywordReturn CondExpr Semicolon // Or: Expr -> CompExpr | CondExpr  
      | KeywordIf LeftParen CondExpr RightParen Stmt  
      | KeywordIf LeftParen CondExpr RightParen Stmt KeywordElse Stmt  
      | KeywordWhile LeftParen CondExpr RightParen Stmt  
      | KeywordFor LeftParen ForInit Semicolon CondExpr Semicolon AssignStmt  
        Semicolon Stmt
```

### Function Declaration

```
FuncDec -> Type Identifier LeftParen ParaDecs RightParen  
         LeftBrace Stmts RightBrace
```

## Parser (v)



```
ParaDecs -> ParaDec Comma ParaDecs | ParaDec | $
```

```
ParaDec -> Specifier Identifier
```

### Global Definitions (Variable and Function)

```
ExtDefs -> ExtDef ExtDefs | $ // *
```

```
ExtDef -> FuncDec | Def
```

### The Program

```
Program -> ExtDefs
```

# Parser (vi)



## 2.b.a Error Recovery

**LALRPOP** provides error recovery by doing so.

```
Term: Box<tree::CompExpr> = {
    <n: "int"> => Box::new(tree::CompExpr::Value(tree::Value::Integer(n))),
    // ...
    "(" <CompExpr> ")",
    ! => { errors.push(<>); Box::new(tree::CompExpr::Error) },
}

#[test]
fn test_error_recovery() {
    assert_compexpr_parse("2 + * 5", "(2: i32 + (MissingTermError * 5: i32))");
    assert_compexpr_parse("2 + * 5 *",
        "(2: i32 + ((MissingTermError * 5: i32) * MissingTermError))");
}
```

# AST



Semantic action as building syntax tree. Every parsing result returns a node, represented by an object in Rust. We first define what a node looks like.

```
// ast.rs
pub enum CompExpr {
    Value(Value),
    Variable(Variable),
    UnaryOperation(UnaryOperator, Box<CompExpr>),
    BinaryOperation(Box<CompExpr>, BinaryOperator, Box<CompExpr>),
    Error
}
```

When facing an input, it will returns something in a recursive form, just like a tree.

```
// grammar.lalrpop
pub CompExpr: Box<tree::CompExpr> = {
```

## AST (ii)



```
Term,  
<lhs:CompExpr> "^" <rhs:CompExpr> => {  
    Box::new(tree::CompExpr::BinaryOperation (lhs, tree::BinaryOperator::Pow, rhs))  
},  
<lhs:CompExpr> "%" <rhs:CompExpr> => {  
    Box::new(tree::CompExpr::BinaryOperation (lhs, tree::BinaryOperator::Mod, rhs))  
},  
}  
  
#[test]  
fn test_expr() {  
    // Test if evaluation order is correct  
    assert_compexpr_parse("2 + 4 * 5", "(2: i32 + (4: i32 * 5: i32))");  
    // Test expression with bracket  
    assert_compexpr_parse("(2 + 4) * 5", "((2: i32 + 4: i32) * 5: i32)");  
}
```



## 3 The Design of the Compiler

## Design Principles



We try to implement our language from easy to hard. First, we are planning to implement C-basic features. Then, trying to dive deep into the principle of Rust compilers, we are trying to realize macro and lifetime features.

The **macro** will be parsed first in parser. To simplify the problem, we will only focus on macro translation instead of analyzing functional macros like Rust compiler does.

The **lifetime parameter** will be implemented using the natural feature `lalrpop` has. By locating the declaration of variable and function, we might be able to analyze the lifetime of variables and try to avoid some of the memory issues and de-referencing issues we mentioned in the former presentation. However, the notion of lifetime must come with ownership, therefore, our realization of lifetime checking might be a rather simple one.

## 4 Implementation Progress










# Progress



- ☒ Simple Lexer
- ☐ More complicated Lexer
- ☒ Simple Parser
- ☐ More complicated parser
- ☒ Simple AST structure
- ☐ AST with visitor mode
- ☐ Semantic Analysis
- ☐ Intermediate Result Generation
- ☐ Intermediate Result Optimization

# Schedule



-  Simple Lexer
-  12.1 More complicated Lexer
-  Simple AST structure
-  12.1 AST with visitor mode
-  Simple Parser
-  12.1 More complicated parser
-  12.1 Semantic Analysis
-  12.22 Intermediate Result Generation
-  12.22 Intermediate Result Optimization