

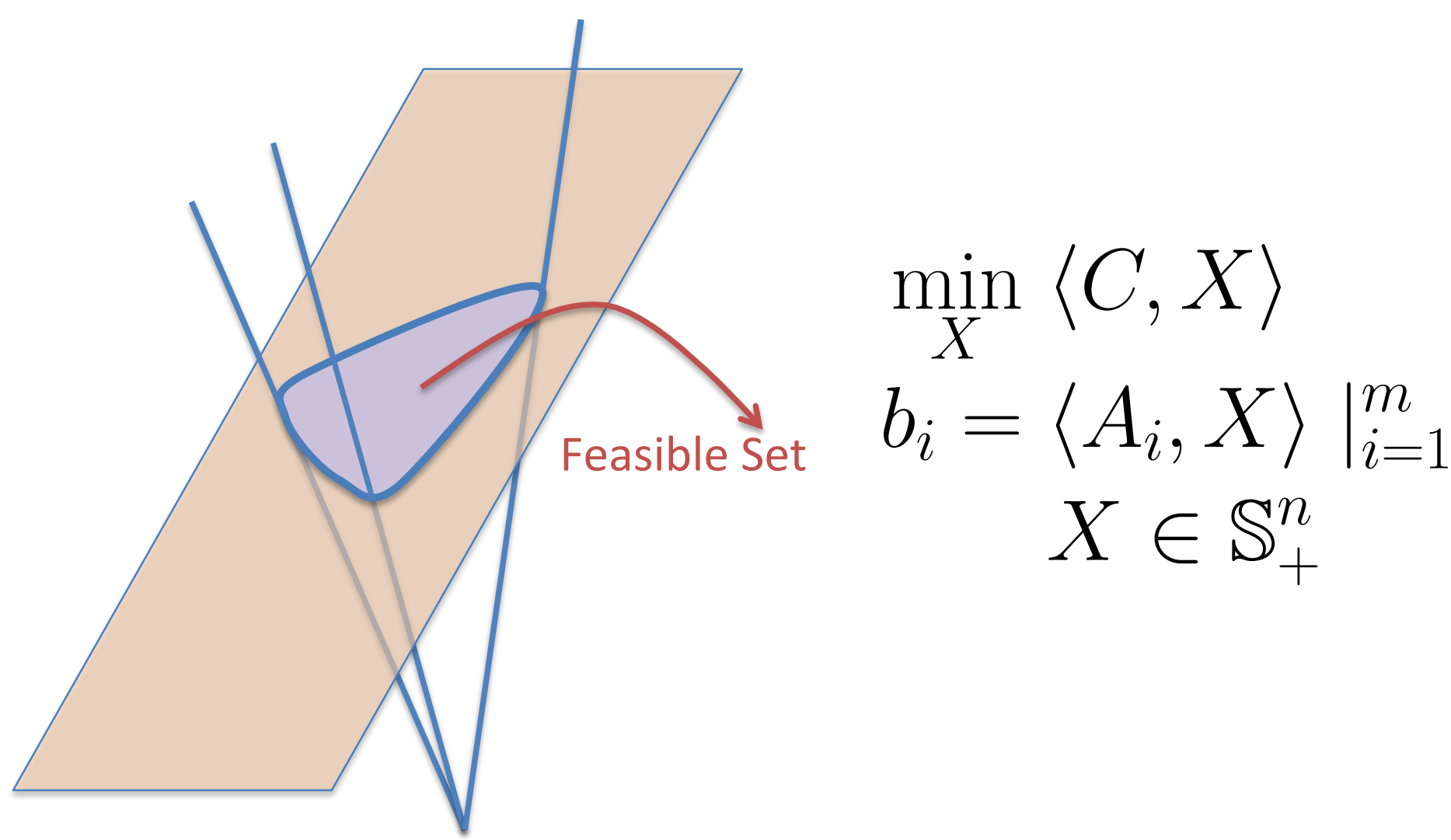
# Exploiting SDP Structure Yields Tighter Approximations

Jared Miller<sup>1</sup>, Yang Zheng<sup>2</sup>, Mario Sznaier<sup>1</sup>, Antonis Papachristodoulou<sup>3</sup>

<sup>1</sup>Northeastern University    <sup>2</sup>Harvard University    <sup>3</sup>University of Oxford

## Semidefinite Programs

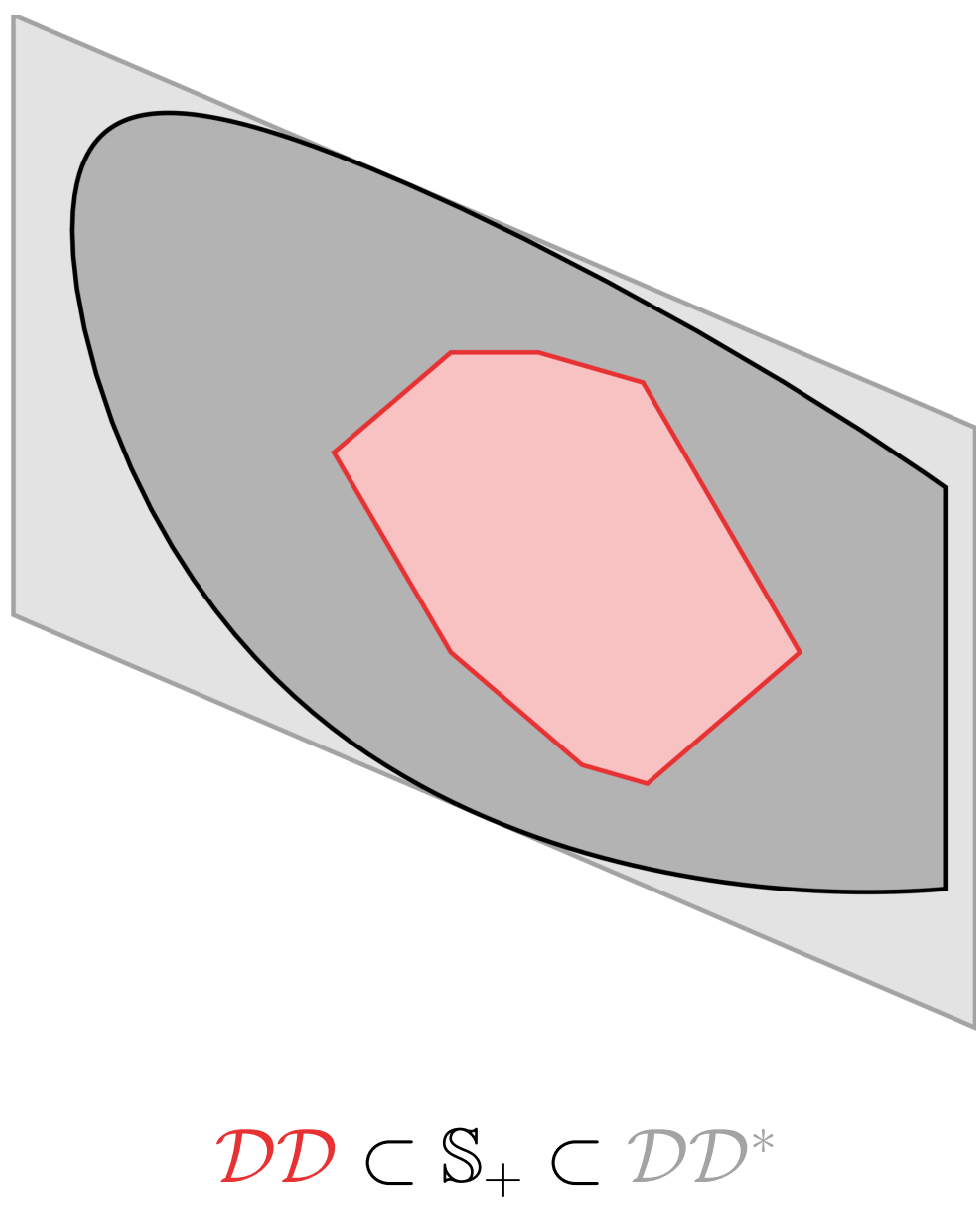
Modern control problems require solutions of large scale SDPs



Runtime scales as  $O(n^2m^2 + n^3m)$

## SDP Approximation

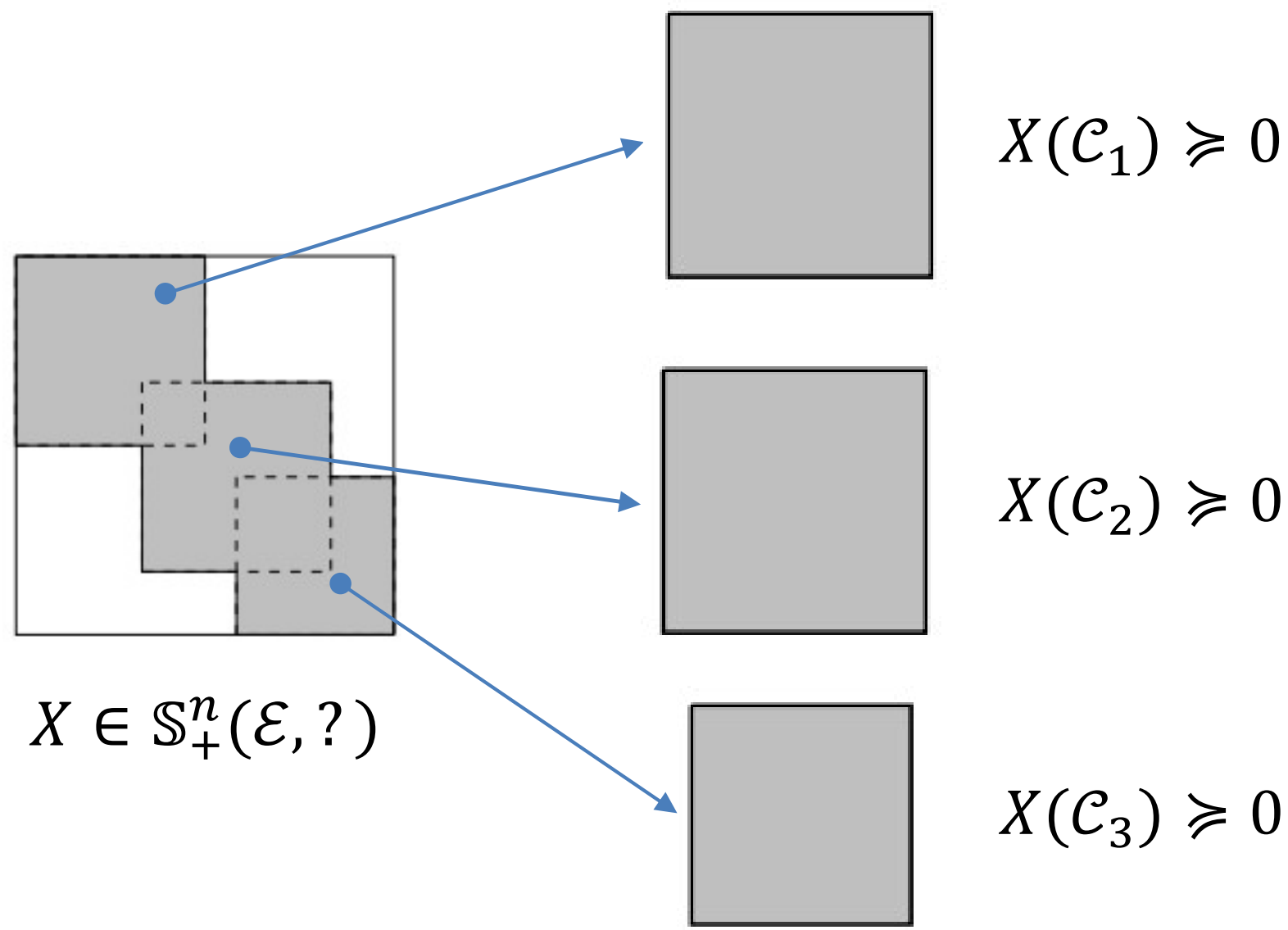
Structured (simpler) subset of PSD  
Diagonally Dominant:  $X_{ii} \geq \sum_{i \neq j} |X_{ij}|$



$DD$ : Upper and Lower bounds by LP

## SDP Structure

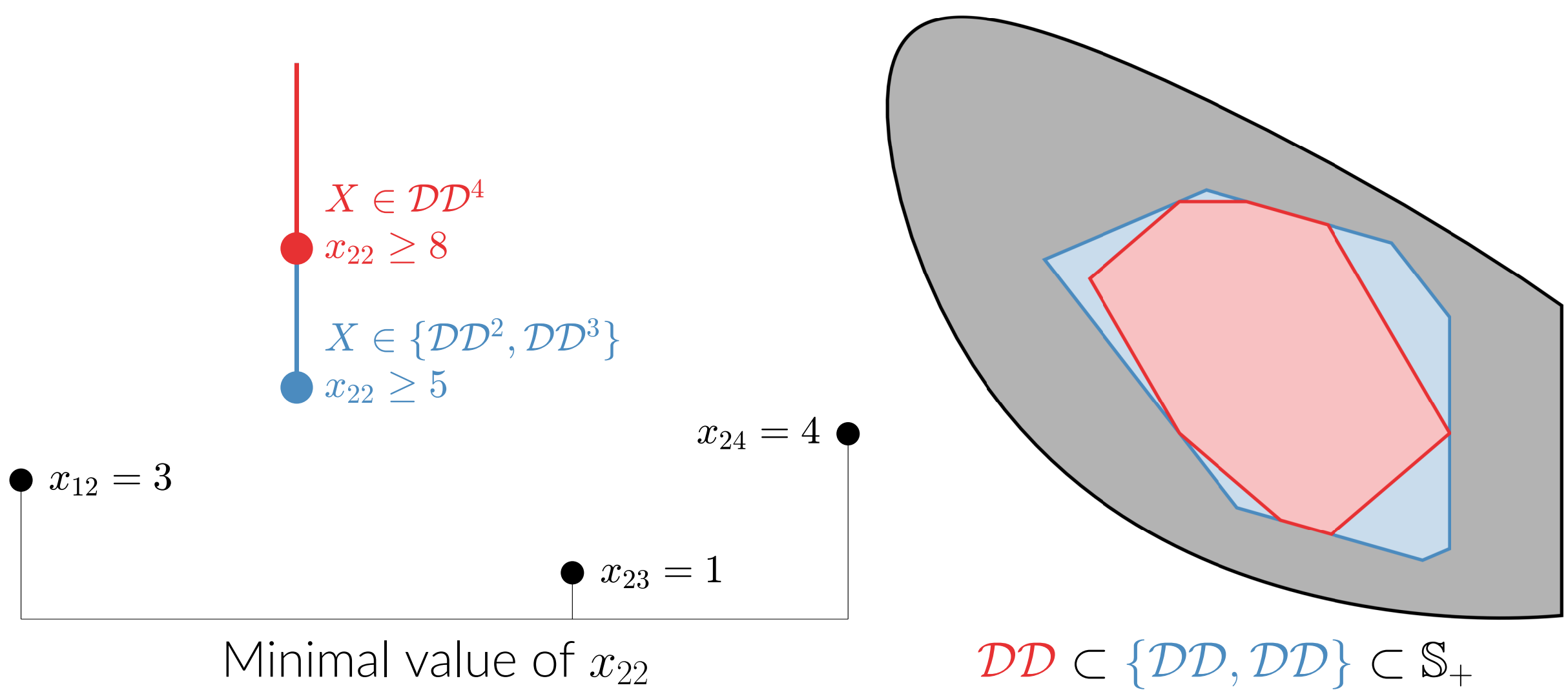
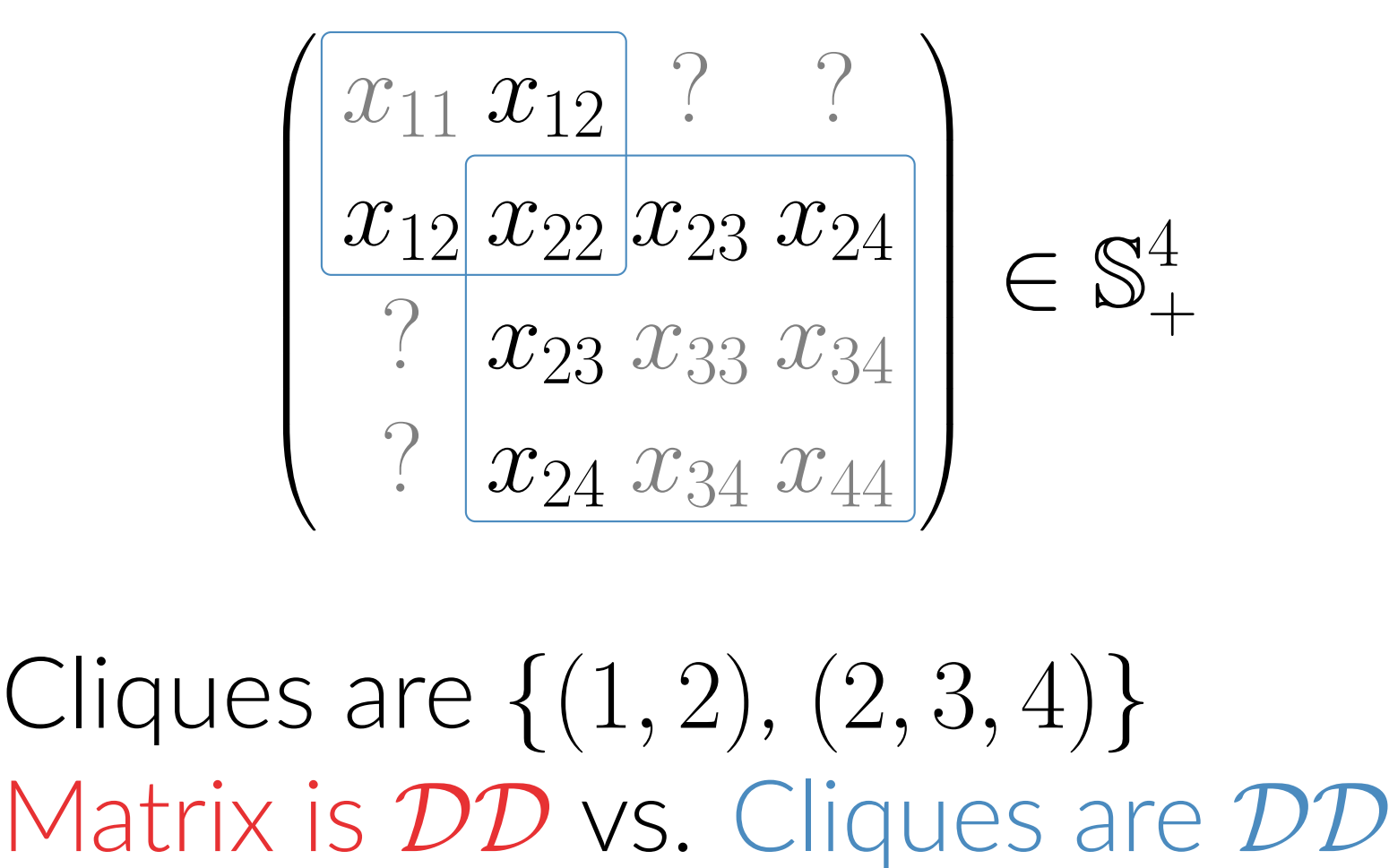
Improve runtime by reducing  $n$   
Decompose based on structure



Includes sparsity, symmetry, \*-algebra

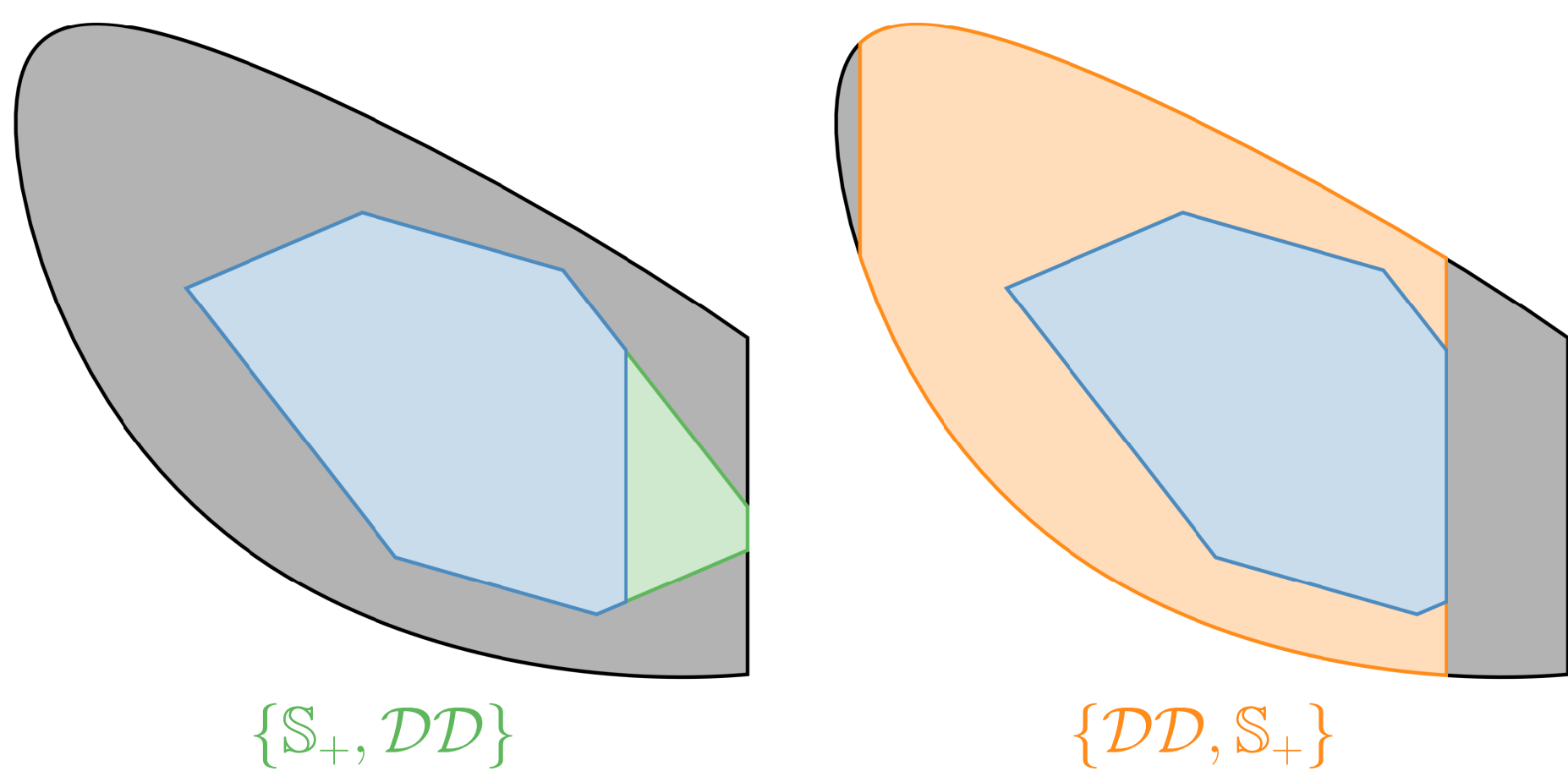
## Structure Broadens Feasible Regions

Approximations destroy structure, worse runtime and bounds



## Mixing Cones

Adds flexibility in optimization



Useful if problem has few large cliques

## Example: Polynomial Optimization

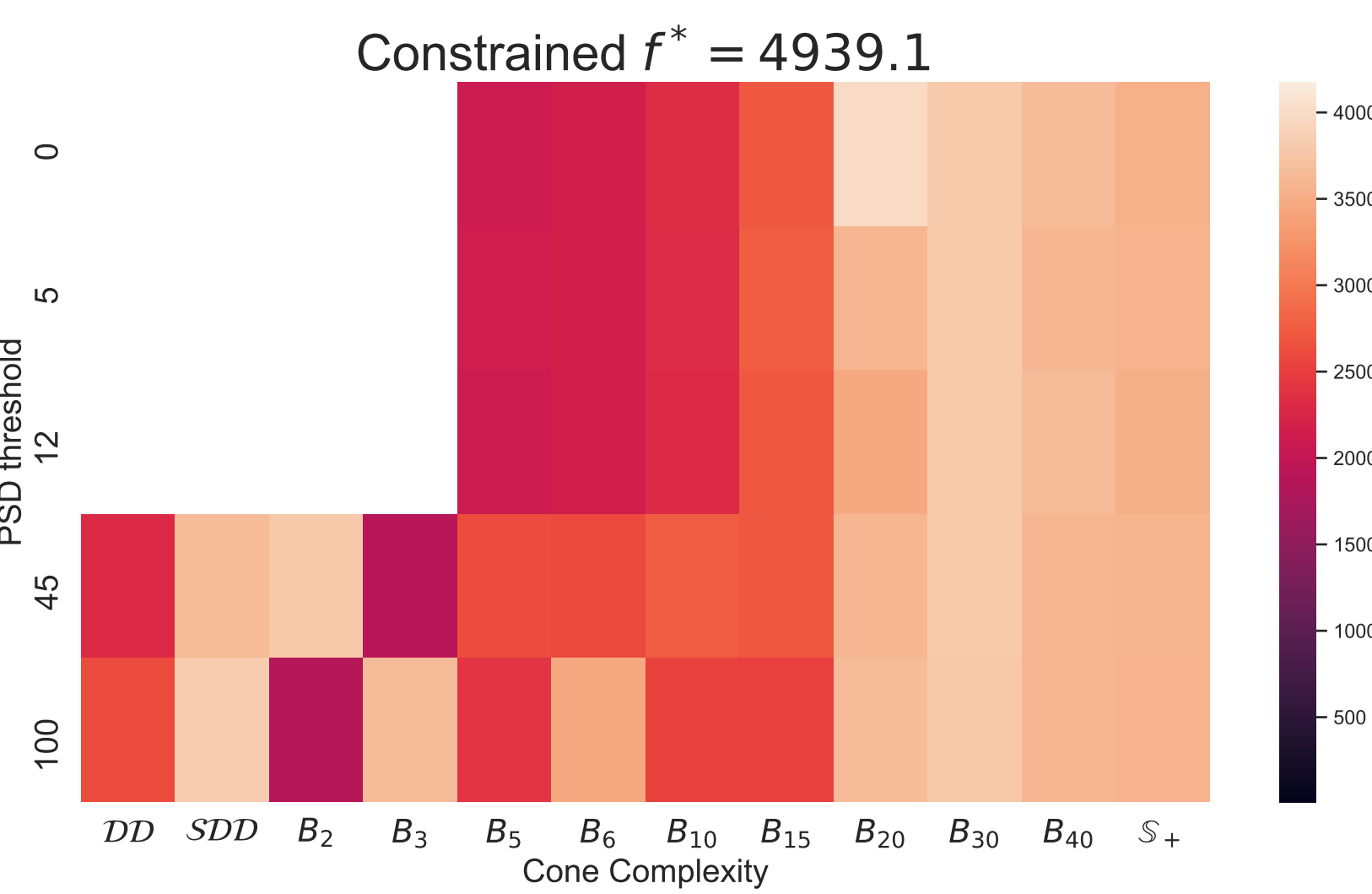
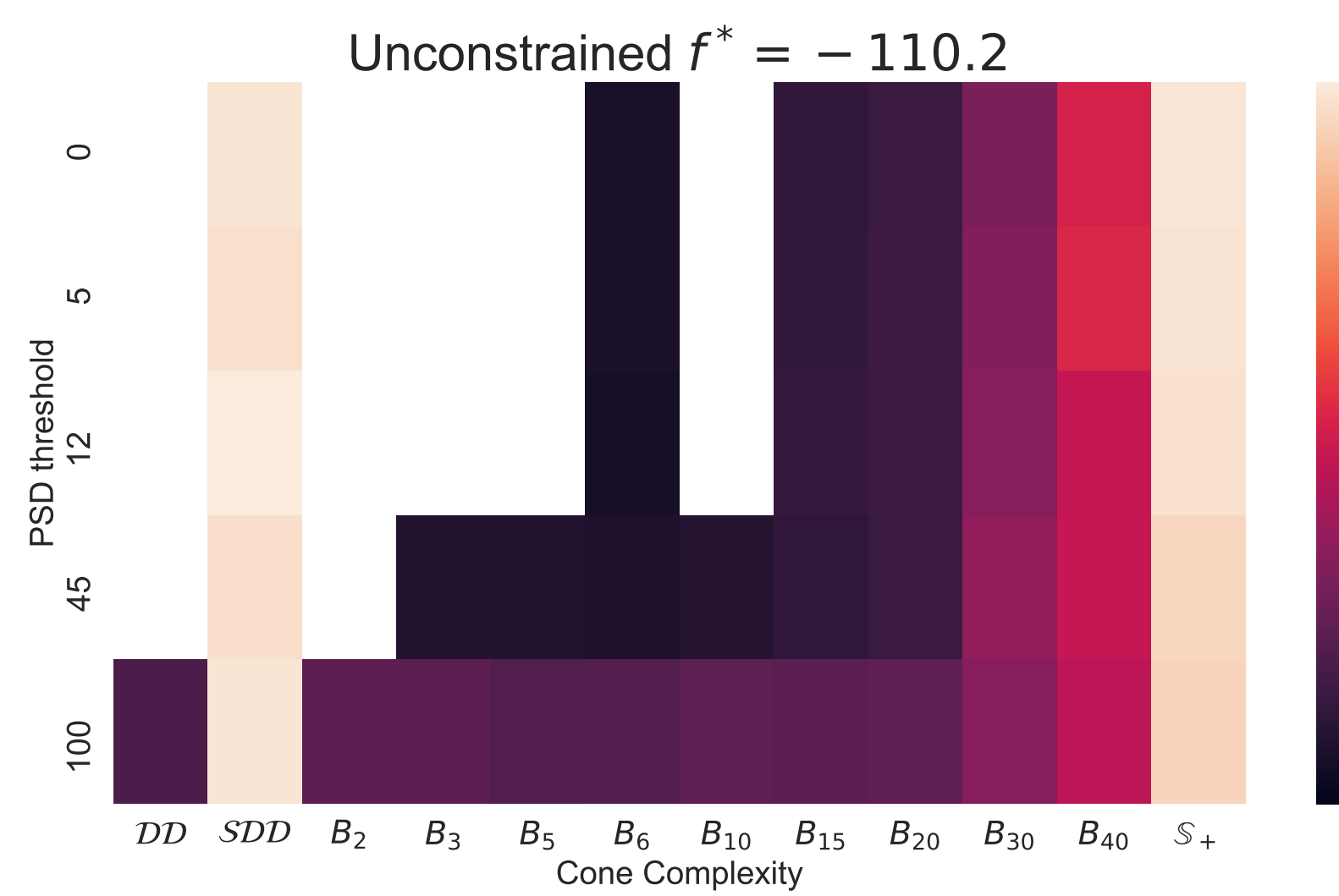
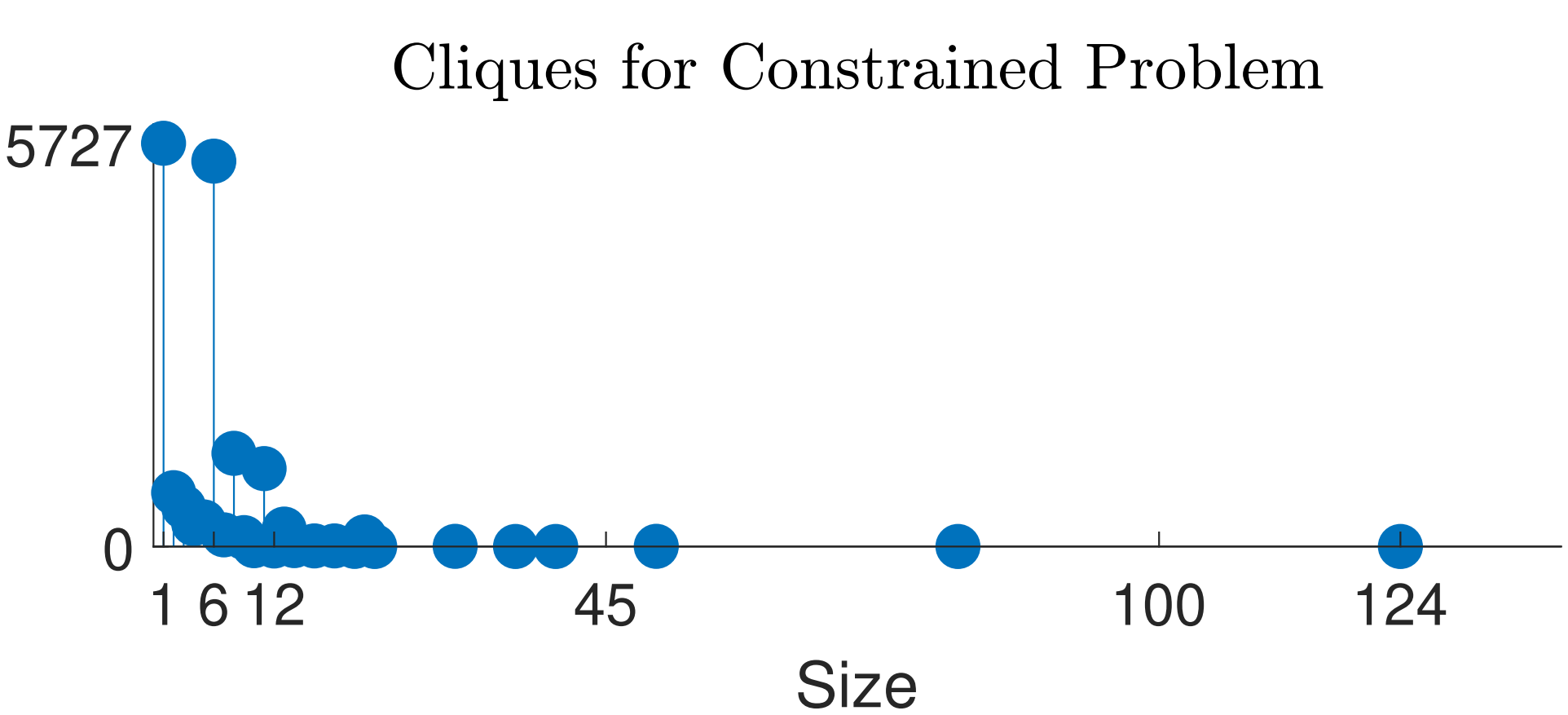
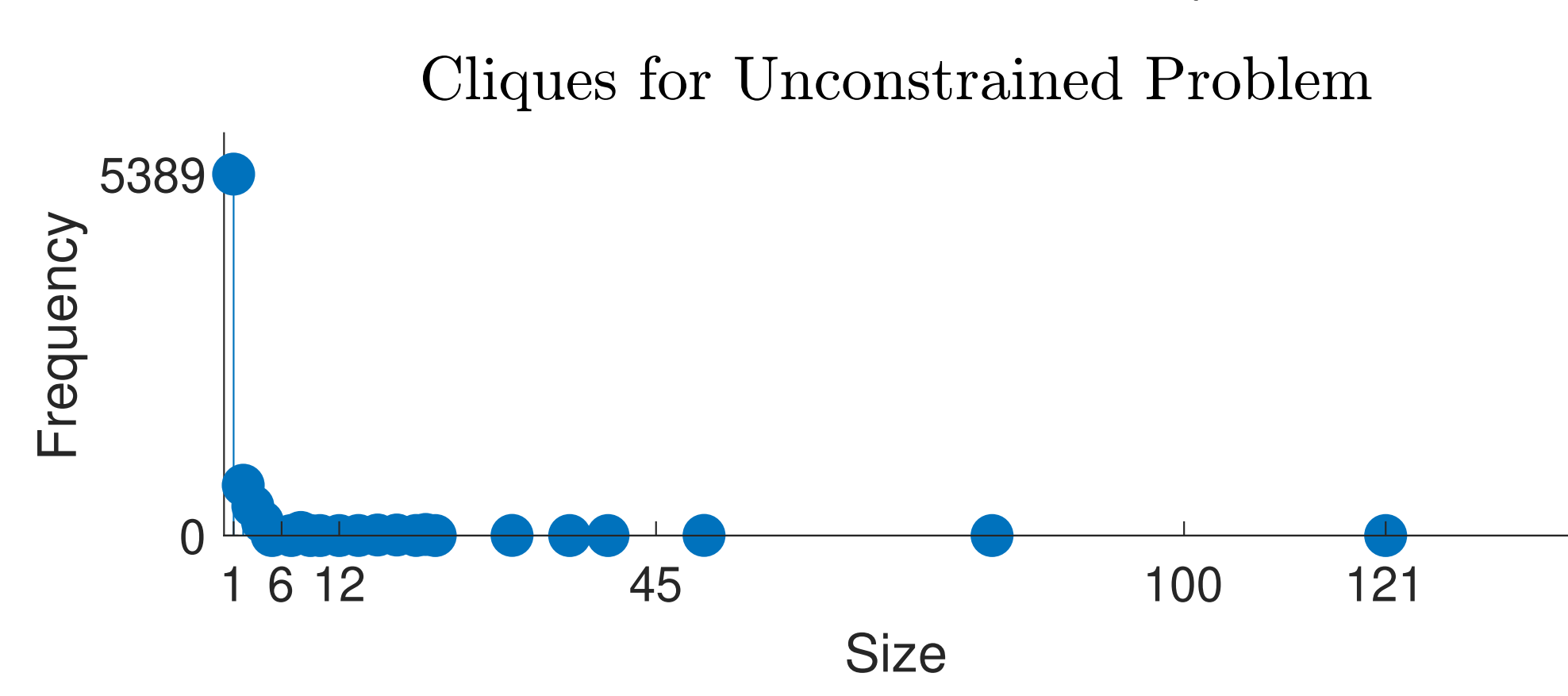
Lower bounds by 2<sup>nd</sup> order Moment-SOS, decomposition by term sparsity  
Sparse Quartic

$$f(x) = f_R(x) + f_Q(x)$$

Dense Quadratic

$$f^* = \min_x f(x)$$

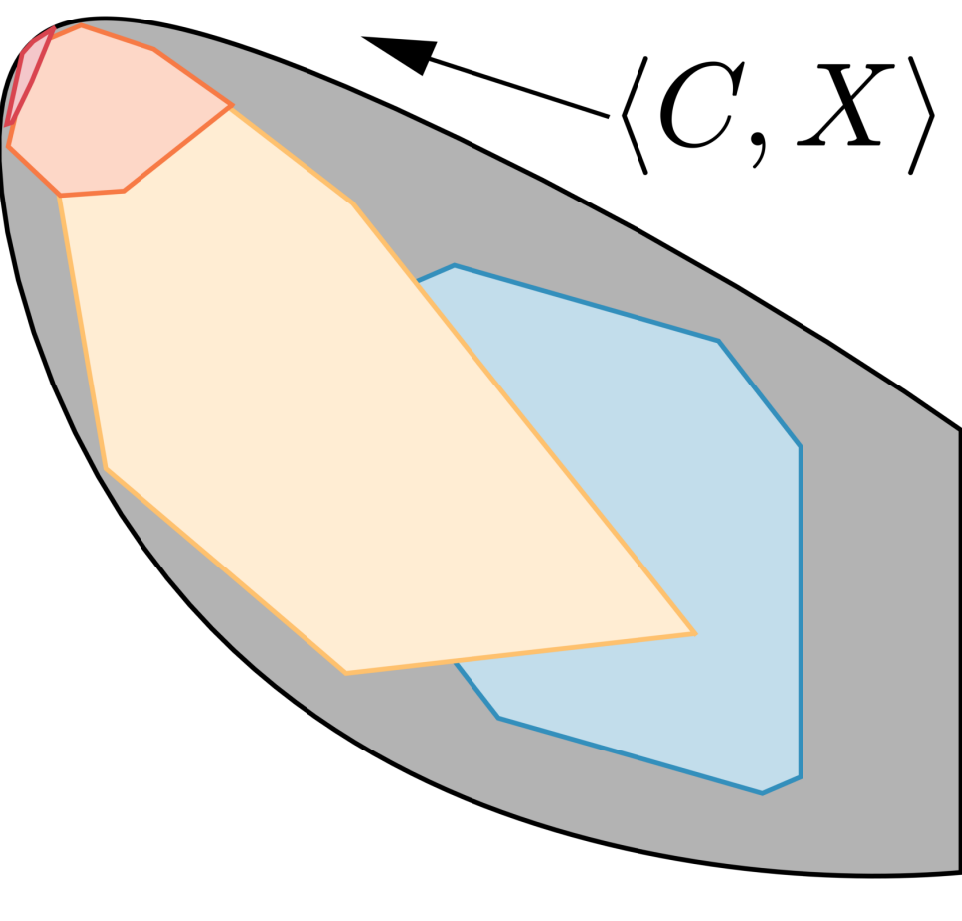
$x \in [1, 2]^{120}$  Box Constraint



Time to find SDP-matching lower bounds (seconds)

## Implications

Structure improves approximations  
Change of Basis: iterative refinement



Maximize cost:  $p_0 \leq p_1 \leq p_2 \leq p_3$

Future steps:

- Convergence to SDP optimum
- Optimal Power Flow
- $H_2/H_\infty$  Network Control



arXiv:1911.12859

github.com/zhengy09/SDPfW