Bounding the Distance to Unsafe Sets with Convex Optimization

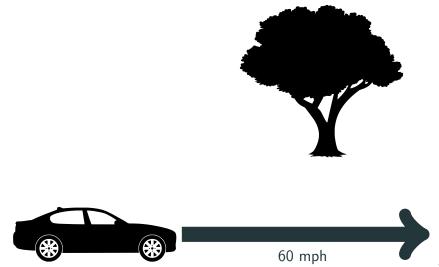
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DSCD Rising Stars (Robotics)



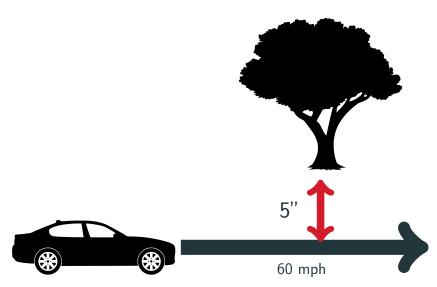
Safety Example



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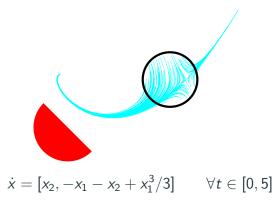
Main Ideas

Quantify safety of trajectories by distance to unsafe set

Create linear program to bound distance

Solve using Semidefinite Programming

Flow System Setting



$$X_0 = \{x \mid (x_1 - 1.5)^2 + x_2 \le 0.4^2\}$$

$$X_u = \{x \mid x_1^2 + (x_2 + 0.7)^2 \le 0.5^2,$$

$$\sqrt{2}/2(x_1 + x_2 - 0.7) \le 0\}$$

Distance Function

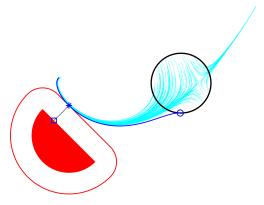
Metric space (X, c) satisfying $\forall x, y \in X$:

$$c(x,y) > 0$$
 $x \neq y$
 $c(x,x) = 0$
 $c(x,y) = c(y,x)$
 $c(x,y) \leq c(x,z) + c(z,y)$ $\forall z \in X$

Point-Unsafe Set distance: $c(x; X_u) = \min_{y \in X_u} c(x, y)$

Distance Estimation Problem (Nonconvex)

$$P^* = \min_{t, x_0 \in X_0} c(x(t \mid x_0); X_u)$$
$$\dot{x}(t) = f(t, x), \quad \forall t \in [0, T].$$

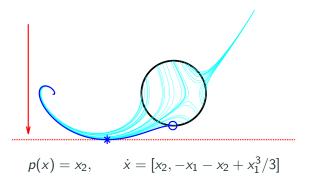


Peak Estimation

Peak Estimation Background

Find minimum value of p(x) along trajectories

$$P^* = \min_{t, x_0 \in X_0} p(x(t \mid x_0))$$
$$\dot{x}(t) = f(t, x(t)) \qquad t \in [0, T]$$



Peak Function Program

Infinite dimensional linear program (Fantuzzi, Goluskin, 2020) Uses auxiliary function v(t,x)

$$d^* = \max_{\gamma \in \mathbb{R}} \quad \gamma \tag{1a}$$

$$v(t,x) \le p(x)$$
 $\forall (t,x) \in [0,T] \times X$ (1b)

$$(\partial_t + f \cdot \nabla_x)v(t,x) \ge 0 \quad \forall (t,x) \in [0,T] \times X \quad (1c)$$

$$\gamma \le \nu(0, x) \qquad \forall x \in X_0 \tag{1d}$$

$$v \in C^1([0,T] \times X) \tag{1e}$$

 $P^* = d^*$ holds if $[0, T] \times X$ is compact, f Lipschitz

Sum-of-Squares Programming

Nonnegativity imposition $p(x) \ge 0 \ \forall x \in \mathbb{R}^n$

Sum-of-squares (SOS) $p(x) = \sum_i q_i(x)^2$ always nonnegative

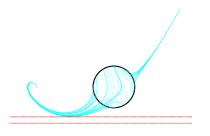
SDP-representable: $p(x) = m(x)^T Q m(x)$ for $Q \succeq 0$

Putinar Psatz: Positivity over $\{x \mid g_k(x) \geq 0\}$,

$$p(x) = \sigma_0(x) + \sum_i \sigma_i(x) g_i(x)$$

$$\exists \sigma_0(x) \in \Sigma[x], \quad \sigma(x) \in (\Sigma[x])^{N_g}$$

Peak Estimation Example Bounds

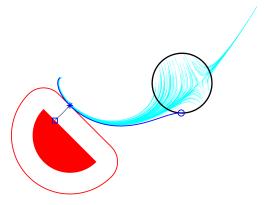


Converging bounds to min. $x_2 = -0.5734$ (moment-SOS) Box region X = [-2.5, 2.5], time $t \in [0, 5]$

Distance Program

Distance Estimation Problem (reprise)

$$P^* = \min_{t, x_0 \in X_0} c(x(t \mid x_0); X_u)$$
$$\dot{x}(t) = f(t, x), \quad \forall t \in [0, T].$$



L₂ bound of 0.2831

Connection to Peak Estimation

Specific form of problem

$$p(x) = c(x; X_u)$$

Moment-SOS hierarchy requires polynomial data

Function $c(x; X_u)$ generally non-polynomial

$$\min_{y \in [-1,1]} ||x - y||_2 = \begin{cases} 0 & x \in [-1,1] \\ |x - \operatorname{sign}(x)| & \text{else} \end{cases}$$

Distance Program (Functions)

Auxiliary v(t,x), point-set proxy $w(x) \le c(x; X_u)$:

$$d^* = \max_{\gamma \in \mathbb{R}} \quad \gamma \tag{2a}$$

$$v(0,x) \geq \gamma \qquad \forall x \in X_0 \tag{2b}$$

$$w(x) \geq v(t,x) \qquad \forall (t,x) \in [0,T] \times X \tag{2c}$$

$$c(x,y) \geq w(x) \qquad \forall (x,y) \in X \times X_u \tag{2d}$$

$$(\partial_t + f \cdot \nabla_x)v(t,x) \geq 0 \quad \forall (t,x) \in [0,T] \times X \tag{2e}$$

$$v \in C^1([0,T] \times X) \tag{2f}$$

$$w \in C(X)$$

Chain
$$\forall (t, x, y) \in [0, T] \times X \times X_u : c(x, y) \ge w(x) \ge v(t, x)$$

Computational Complexity

Use moment-SOS hierarchy

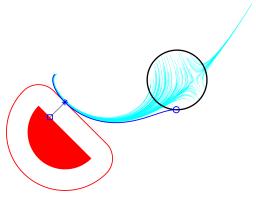
Degree d, dynamics degree $\widetilde{d} = d + \lceil \deg(f)/2 \rceil - 1$

Largest PSD matrix size
$$\max\left[\binom{1+n+\tilde{d}}{\tilde{d}},\binom{2n+d}{d}\right]$$

Timing scales approximately as max $\left[(1+n)^{6\tilde{d}},(2n)^{6d}\right]$

Approximation and Recovery

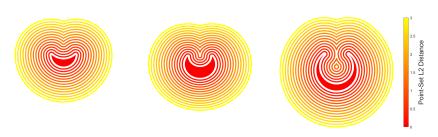
Attempt recovery if SDP (dual) solution has low rank Related to optima extraction in polynomial optimization



 L_2 bound of 0.2831

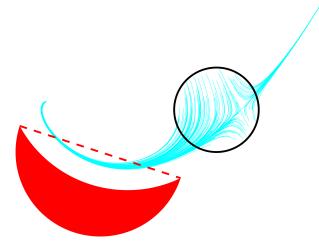
More Examples

Moon L2 Contours



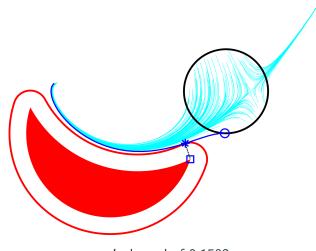
Inside one circle, outside another

Distance Example (Flow Moon)



Collision if X_u was a half-circle

Distance Example (Flow Moon)



 L_2 bound of 0.1592

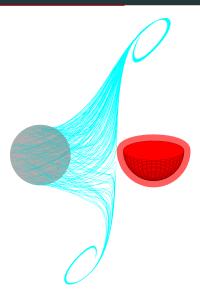
Distance Example (Twist)

'Twist' System,
$$T=5$$

$$\dot{x}_i = A_{ij}x_j - B_{ij}(4x_j^3 - 3x_j)/2$$

$$A = \begin{bmatrix} -1 & 1 & 1 \\ -1 & 0 & -1 \\ 0 & 1 & -2 \end{bmatrix}$$

$$B = \begin{vmatrix} -1 & 0 & -1 \\ 0 & 1 & 1 \\ 1 & 1 & 0 \end{vmatrix}$$



 L_2 bound of 0.0425

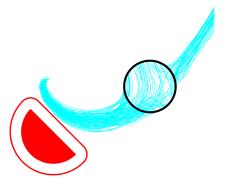
Distance Variations

Distance Uncertainty

Time dependent (bounded) uncertainty $w(t) \in W \ \forall t \in [0, T]$

Dynamics
$$\dot{x}(t) = f(t, x(t), w(t))$$

$$(\partial_t + f(t, x, w) \cdot \nabla_x) v(t, x) \ge 0, \quad \forall (t, x, w) \in [0, T] \times X \times W$$



 L_2 bound of 0.1691, $w(t) \in [-1, 1]$

Shapes along Trajectories

Orientation $\omega(t) \in \Omega$, shape S

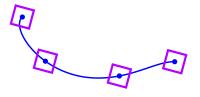
Body to global coordinate transformation A:

$$A: S \times \Omega \rightarrow X$$

$$(s,\omega)\mapsto A(s;\omega)$$

Angular Velocity = 0 rad/sec

Angular Velocity = 1 rad/sec



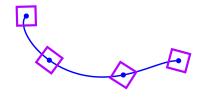
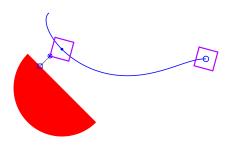


Figure 1: Shape translating and (possibly) rotating

Set-Set Distance Problem

Set-Set distance between $A(S; \omega(t))$ and X_u given t

$$P^* = \min_{t, \omega_0 \in \Omega_0, s \in S} c(A(s; \omega(t \mid \omega_0)); X_u)$$
$$\dot{\omega}(t) = f(t, \omega) \qquad \forall t \in [0, T]$$



L₂ bound of 0.1465

Take-aways

Conclusion

Distance as a measure for safety

Distance estimation with polynomial optimization

Approximate recovery if moment matrices are low-rank

Extend to uncertain and set-set scenarios

Future Work

- Distance-Maximizing Control
- Chance-constrained distance
- Further Sparsity
- Efficient Computation
- Other nonnegativity cones and proofs

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Thank you for your attention

arxiv:2110.14047

http://github.com/jarmill/distance

Graduating May 2023, looking for postdocs