

Exploiting SDP Structure Yields Tighter Approximations

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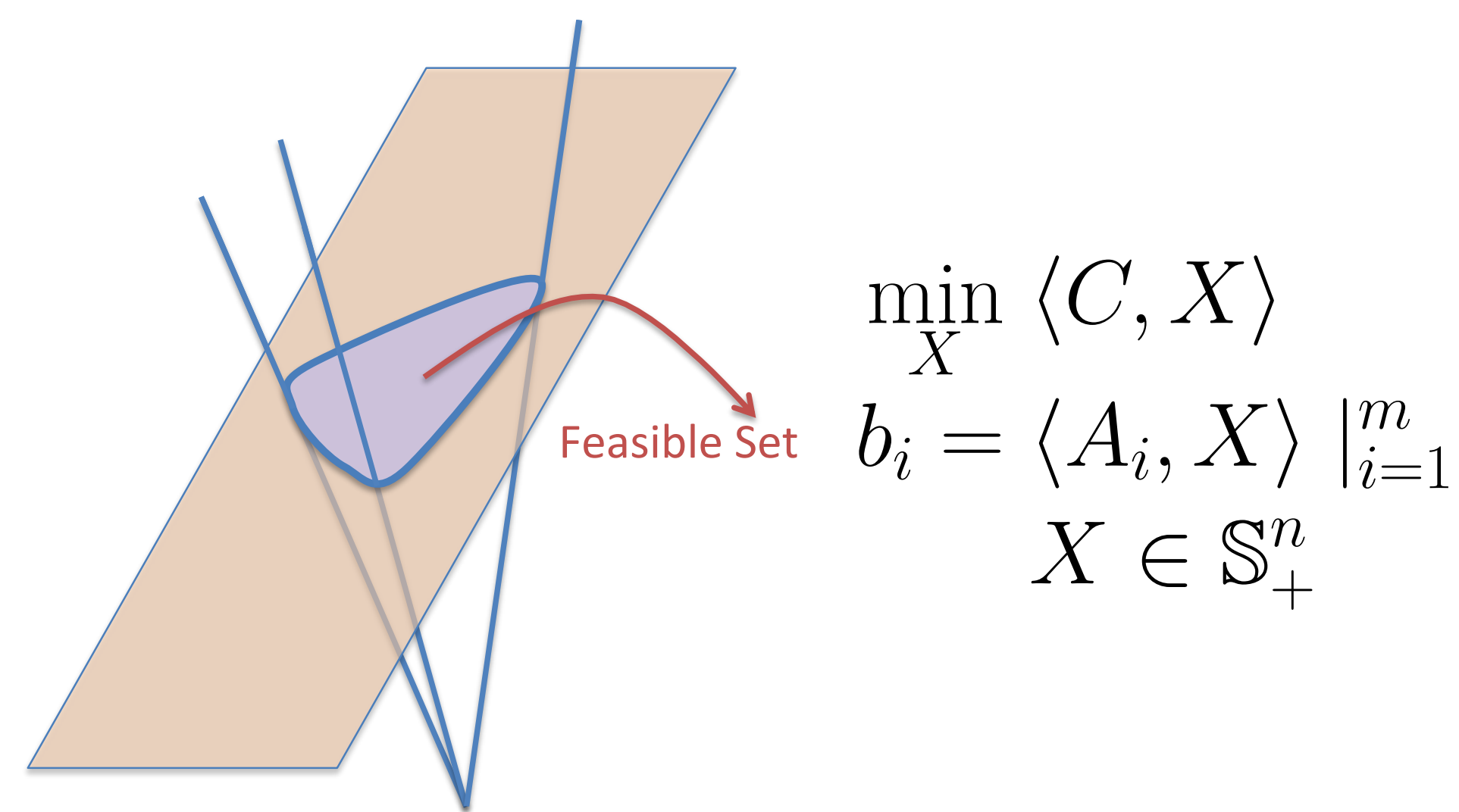
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Semidefinite Programs

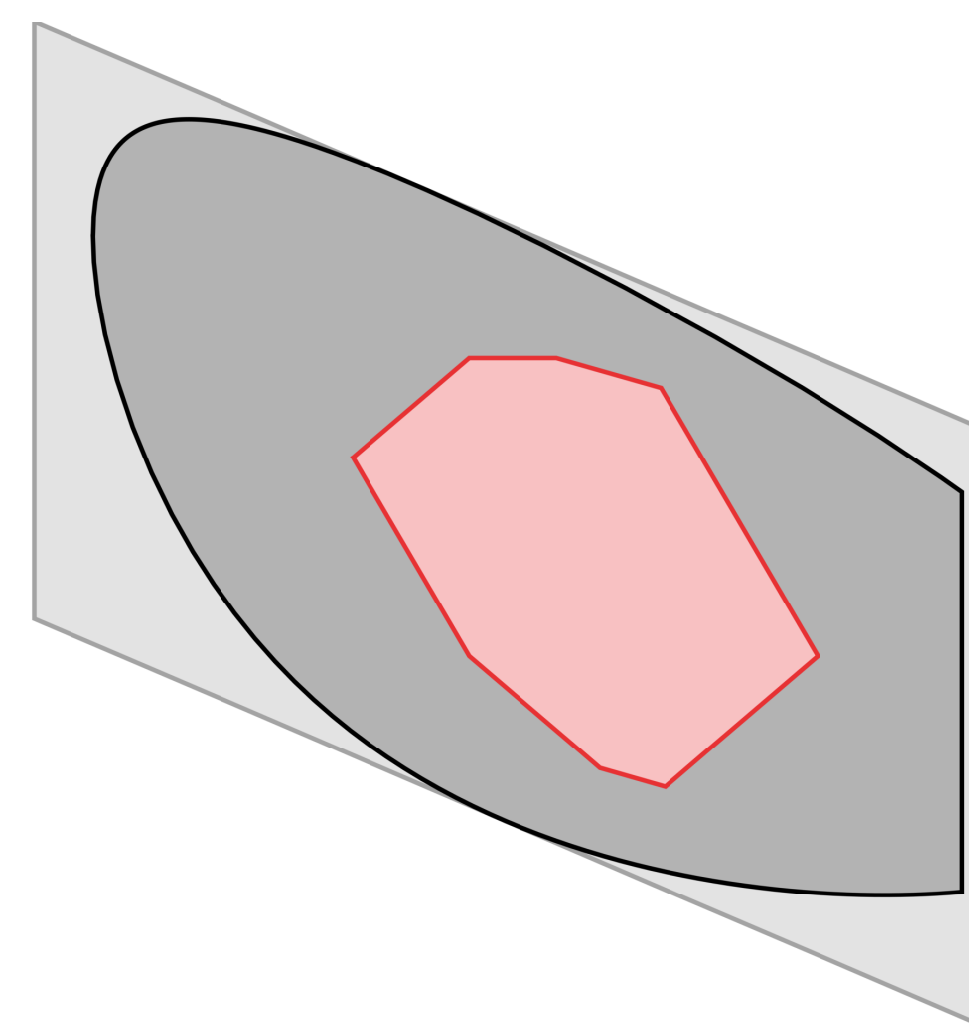
Modern control problems require solutions of large scale SDPs



Runtime scales as $O(n^2m^2 + n^3m)$

SDP Approximation

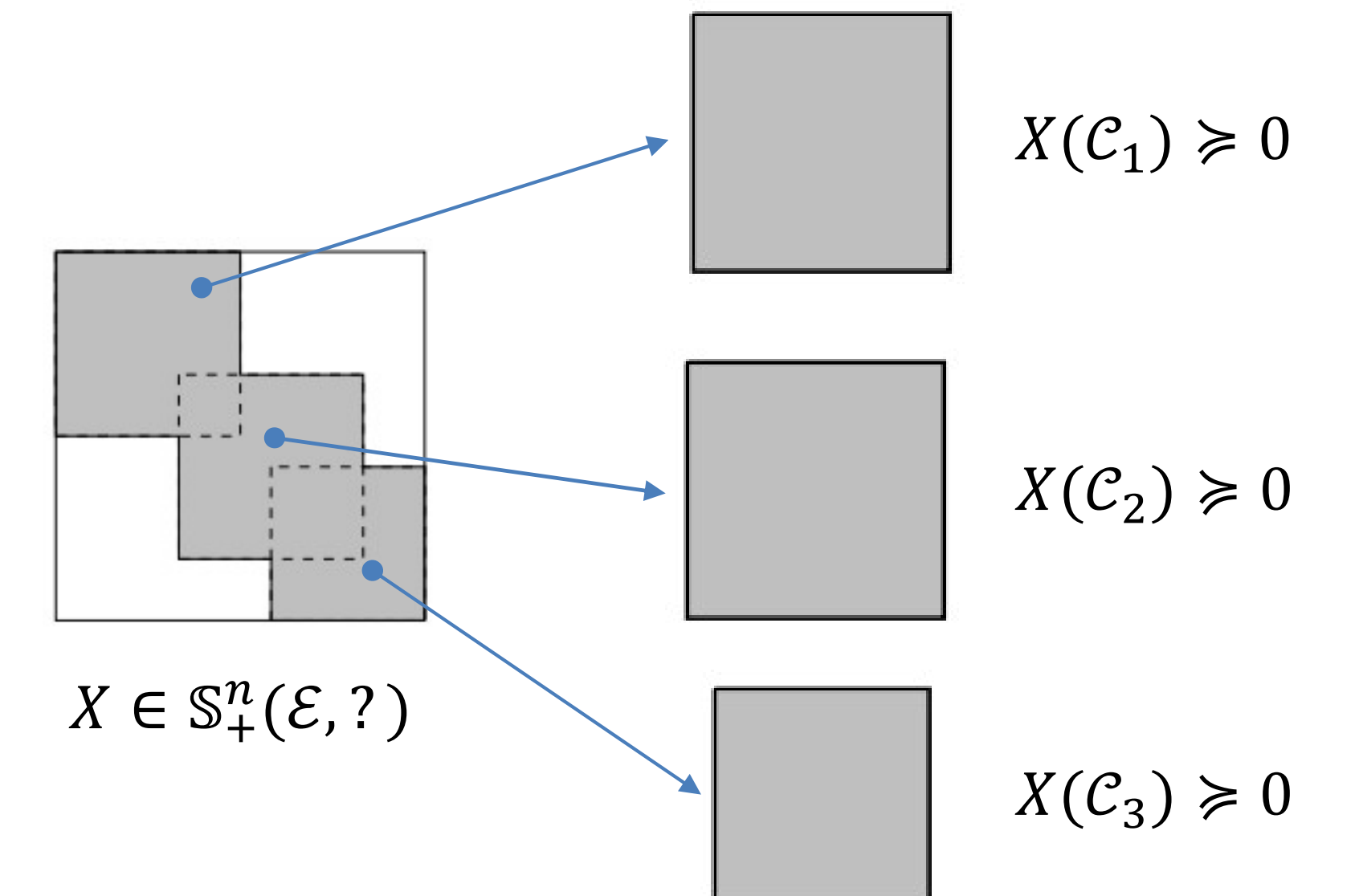
Structured (simpler) subset of PSD
 Diagonally Dominant: $X_{ii} \geq \sum_{i \neq j} |X_{ij}|$



\mathcal{DD} : Upper and Lower bounds by LP

SDP Structure

Improve runtime by reducing n
 Decompose based on structure



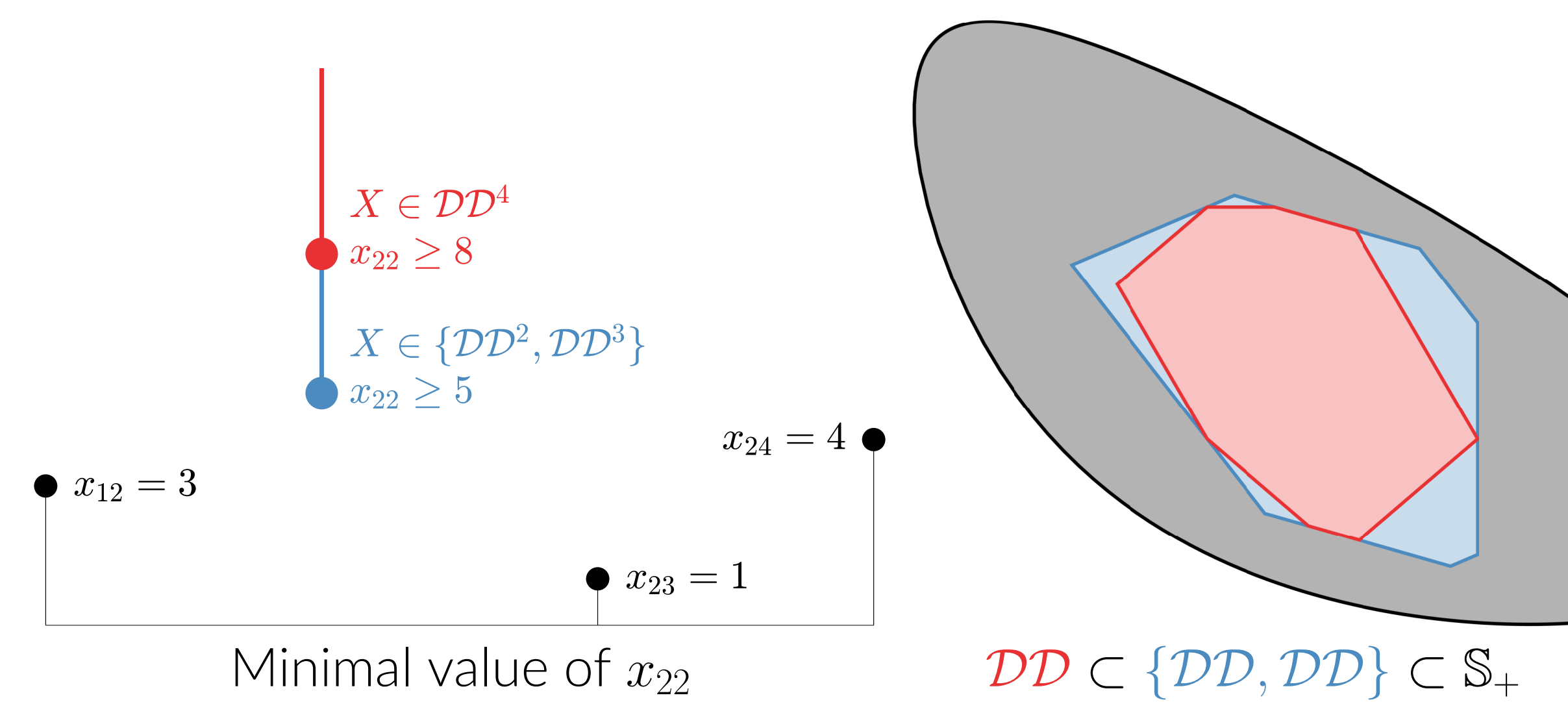
Includes sparsity, symmetry, *-algebra

Structure Broadens Feasible Regions

Approximations destroy structure, worse runtime and bounds

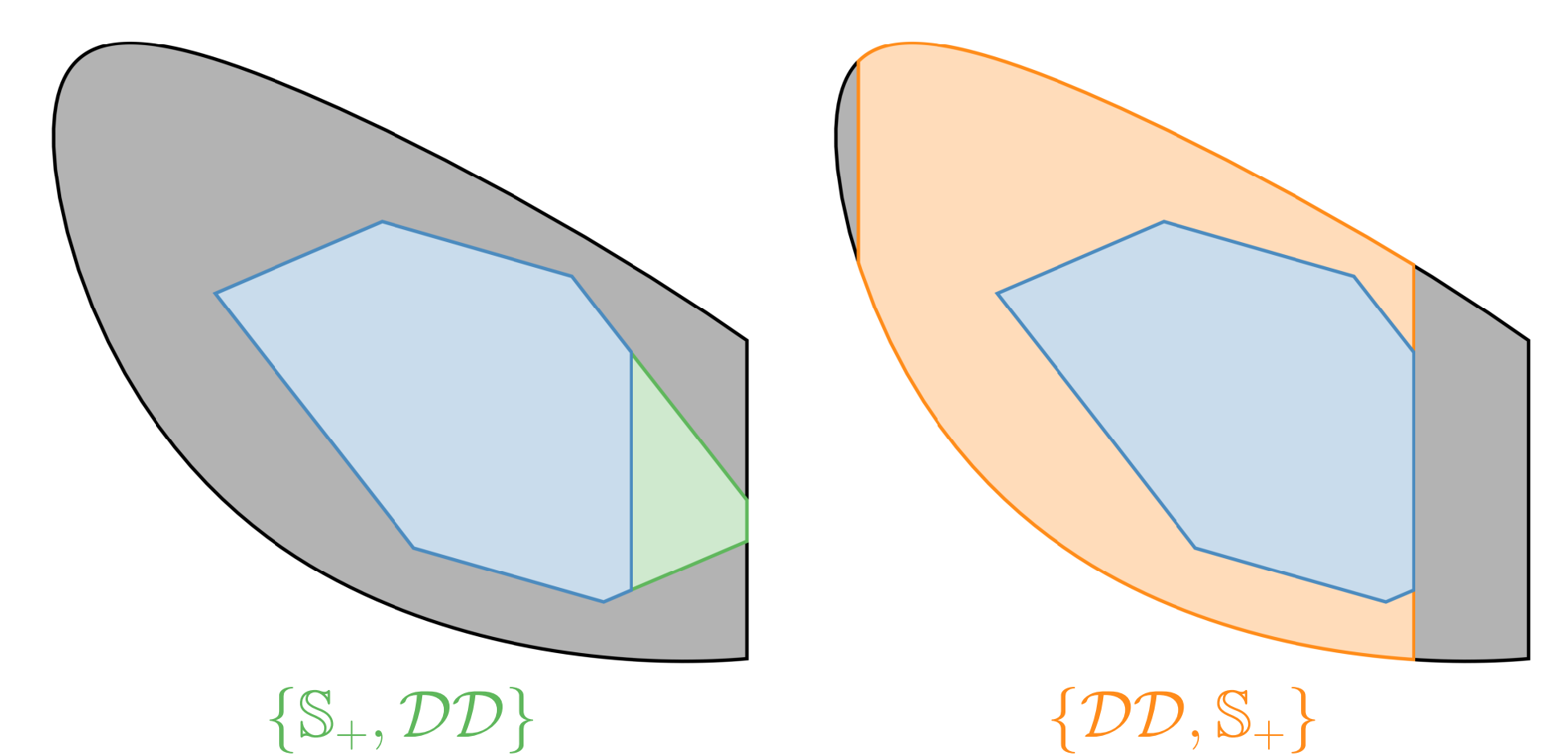
$$\begin{pmatrix} x_{11} & x_{12} & ? & ? \\ x_{12} & x_{22} & x_{23} & x_{24} \\ ? & x_{23} & x_{33} & x_{34} \\ ? & x_{24} & x_{34} & x_{44} \end{pmatrix} \in \mathbb{S}_+^4$$

Cliques are $\{(1, 2), (2, 3, 4)\}$
 Matrix is \mathcal{DD} vs. Cliques are \mathcal{DD}



Mixing Cones

Adds flexibility in optimization



Useful if problem has few large cliques

Example: Polynomial Optimization

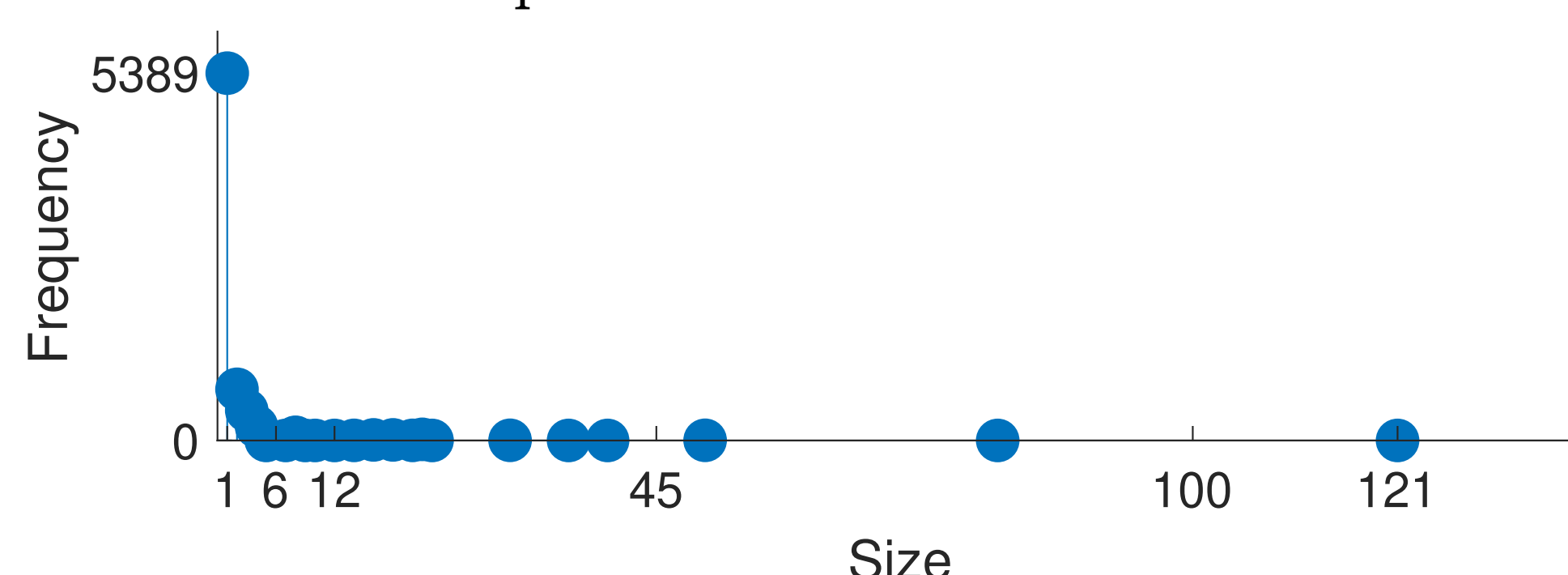
Lower bounds by 2nd order Moment-SOS, decomposition by term sparsity

Sparse Quartic

$$f(x) = f_R(x) + f_Q(x)$$

Dense Quadratic

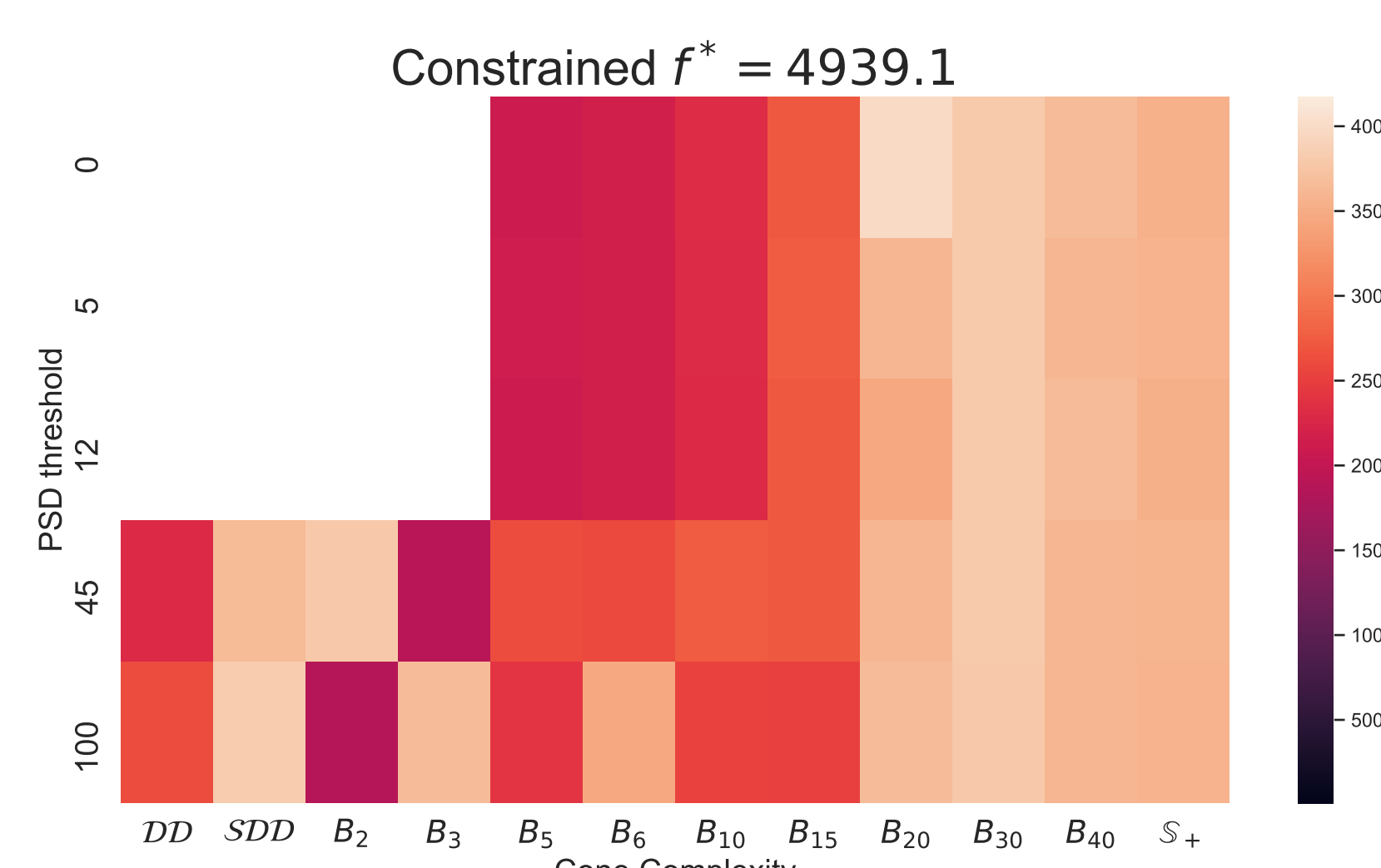
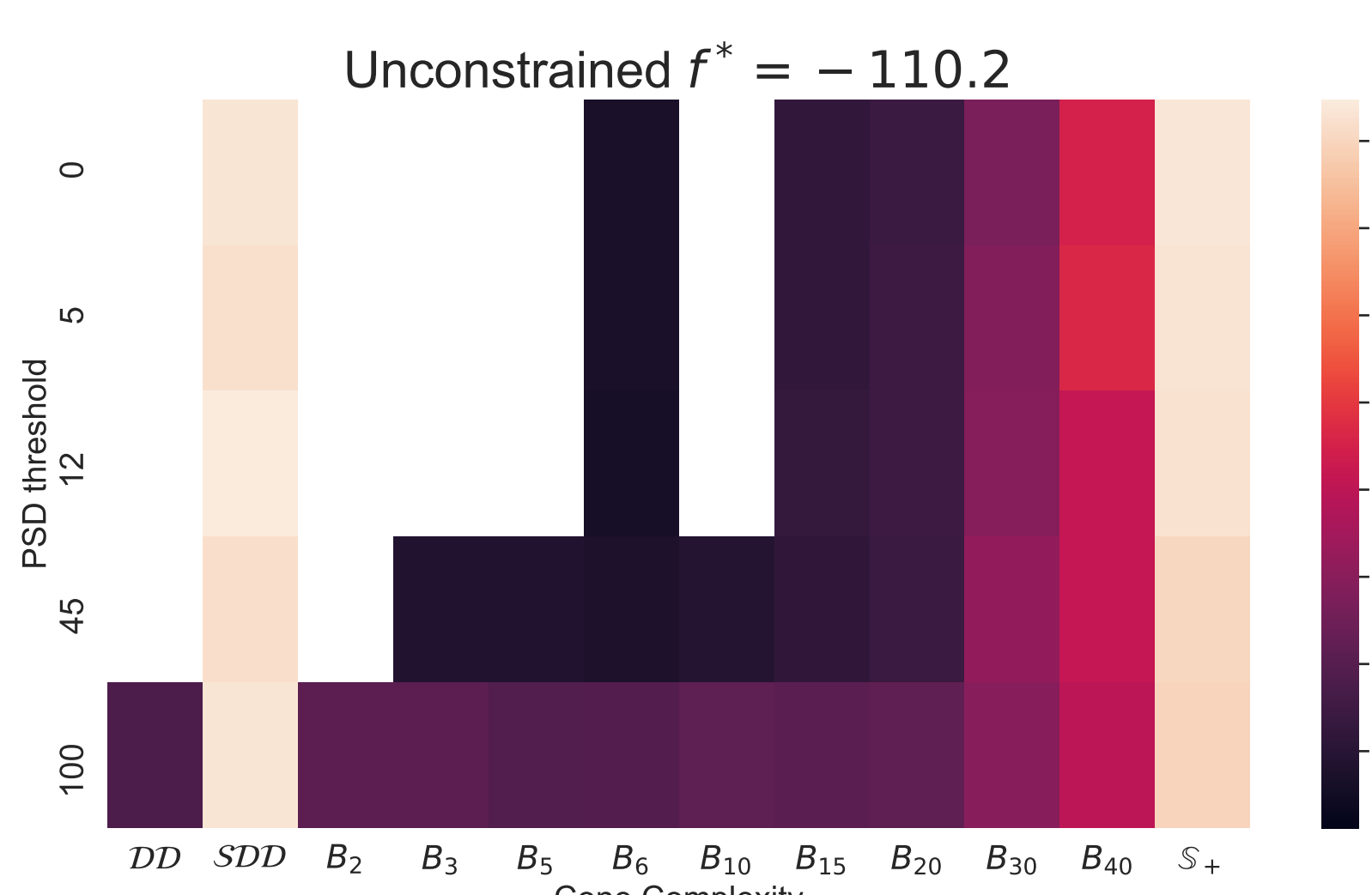
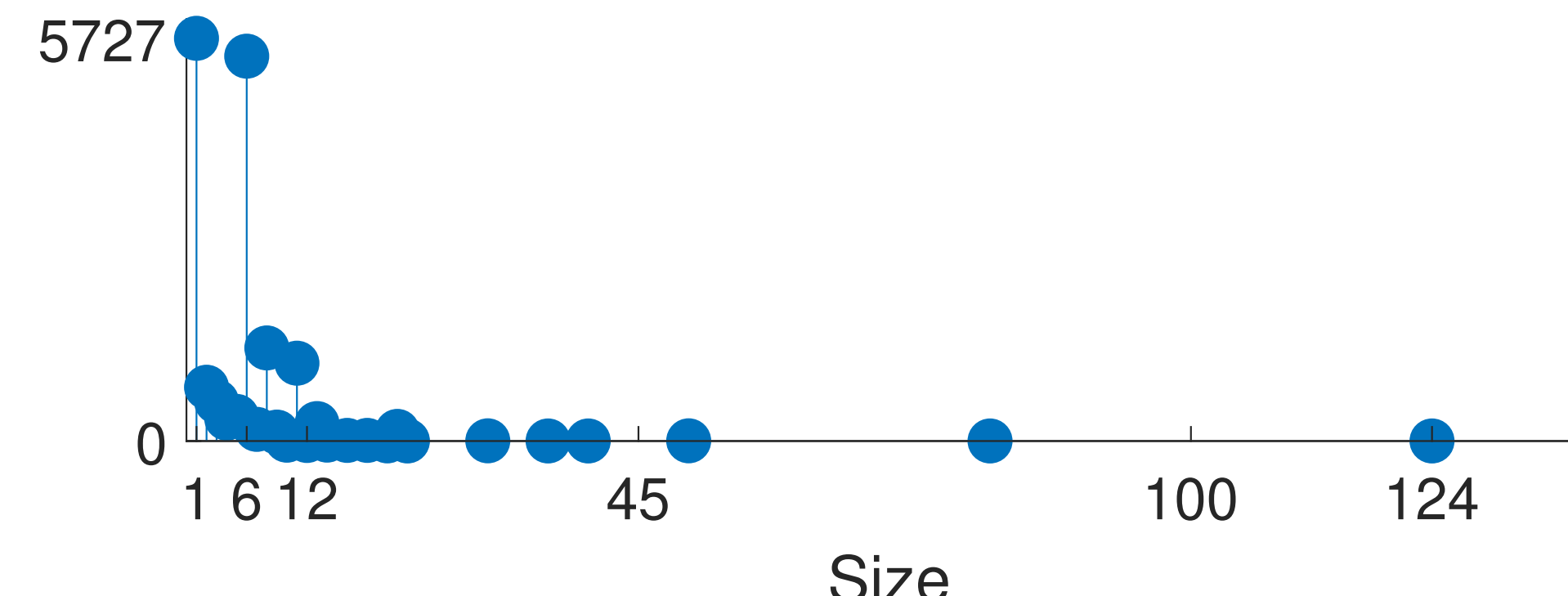
Cliques for Unconstrained Problem



$$f^* = \min_x f(x)$$

$x \in [1, 2]^{120}$ Box Constraint

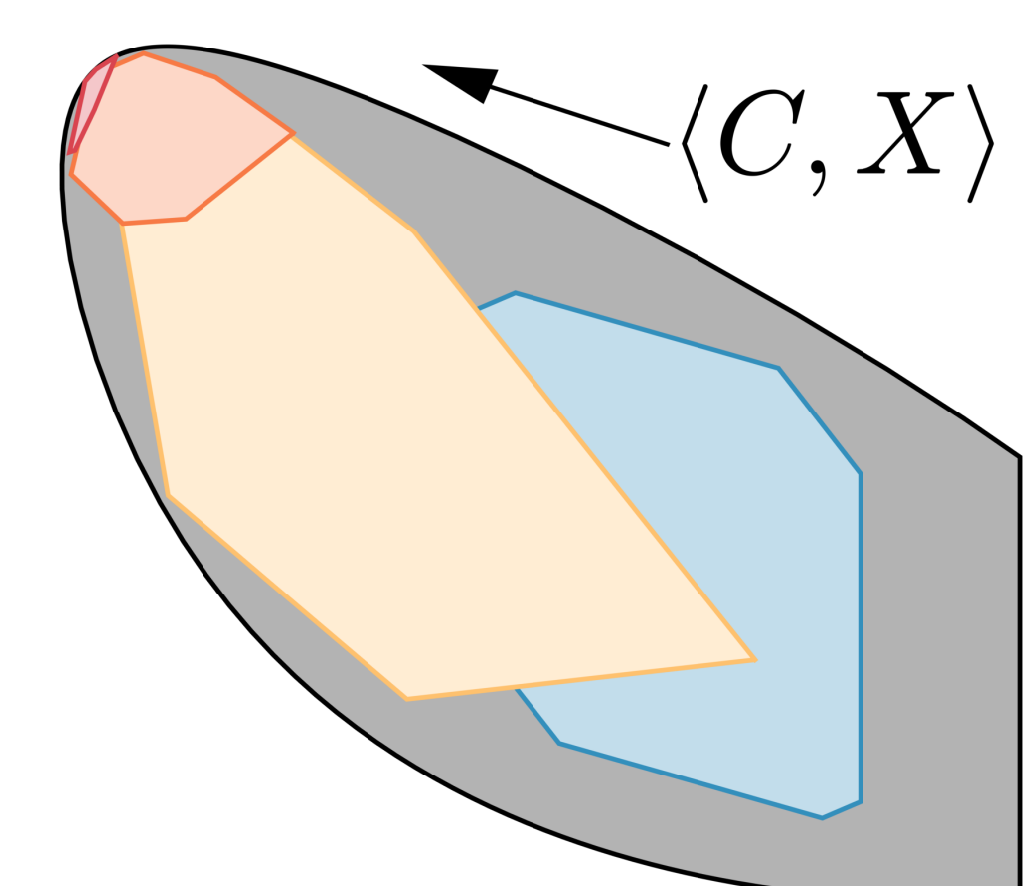
Cliques for Constrained Problem



Time to find SDP-matching lower bounds (seconds)

Implications

Structure improves approximations
 Change of Basis: iterative refinement



Maximize cost: $p_0 \leq p_1 \leq p_2 \leq p_3$

Future steps:

Convergence to SDP optimum

Optimal Power Flow

H_2/H_∞ Network Control



arXiv:1911.12859

github.com/zhengy09/SDPfw