

PROPOSITIONAL LOGIC 1

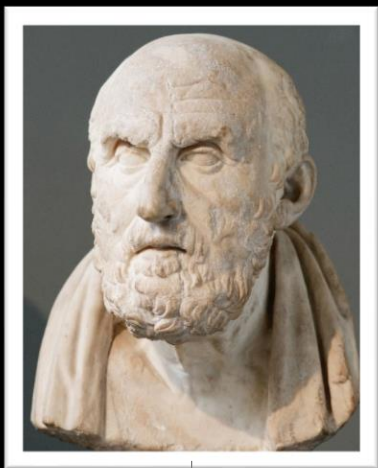
ARTIFICIAL INTELLIGENCE | COMP 131

TODAY ON AI

- What is Logic?
- Propositional Logic
- Questions?

SECTION 01

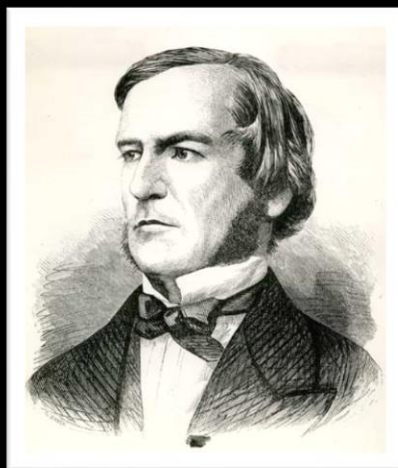
What is Logic?



500 BC



500



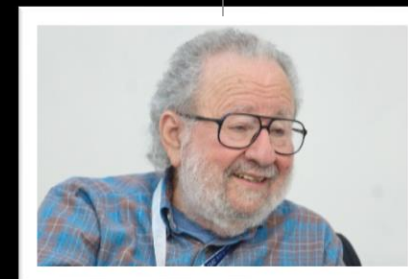
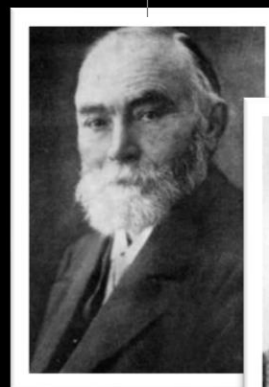
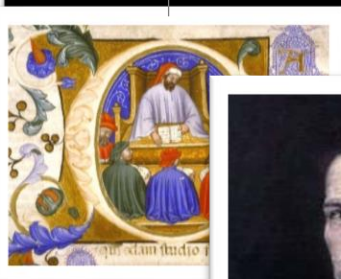
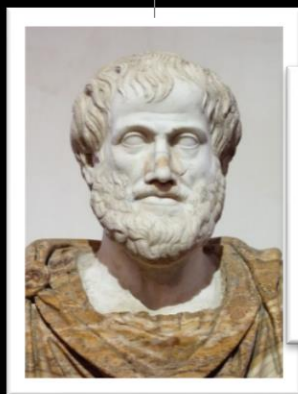
1850



1900



1950



Logic is a way of formally representing the **state of the world** and the world's **rules of operation** so that we can make **rational decisions** and learn **new knowledge** based on our **existing one**.

It allows to:

- **Express knowledge** using a formal language
- **To carry out** reasoning in that language

There are several types of logic. Each type is increasingly complex as it captures more advanced concepts:

LANGUAGE	ONTOLOGICAL COMMITMENT	EPISTEMOLOGICAL COMMITMENT
Propositional logic	Facts	True / False / Unknown
First-order logic	Facts, Objects, Relations	True / False / Unknown
Temporal logic	Facts, Objects, Relations, Times	True / False / Unknown
Probability theory	Facts	Degree of belief $\in [0, 1]$
Fuzzy logic	Facts with degree of truth $\in [0, 1]$	Known interval value
Markov logic	Facts, Objects, Relations	Degree of belief $\in [0, 1]$

- **Programming languages:**

- They are formal and not ambiguous
- Unfortunately they lack expressivity as they cannot accommodate partial information

- **Natural Language:**

- Very expressive but also ambiguous:
- Inference possible, but hard to automate

Flying planes can be dangerous.

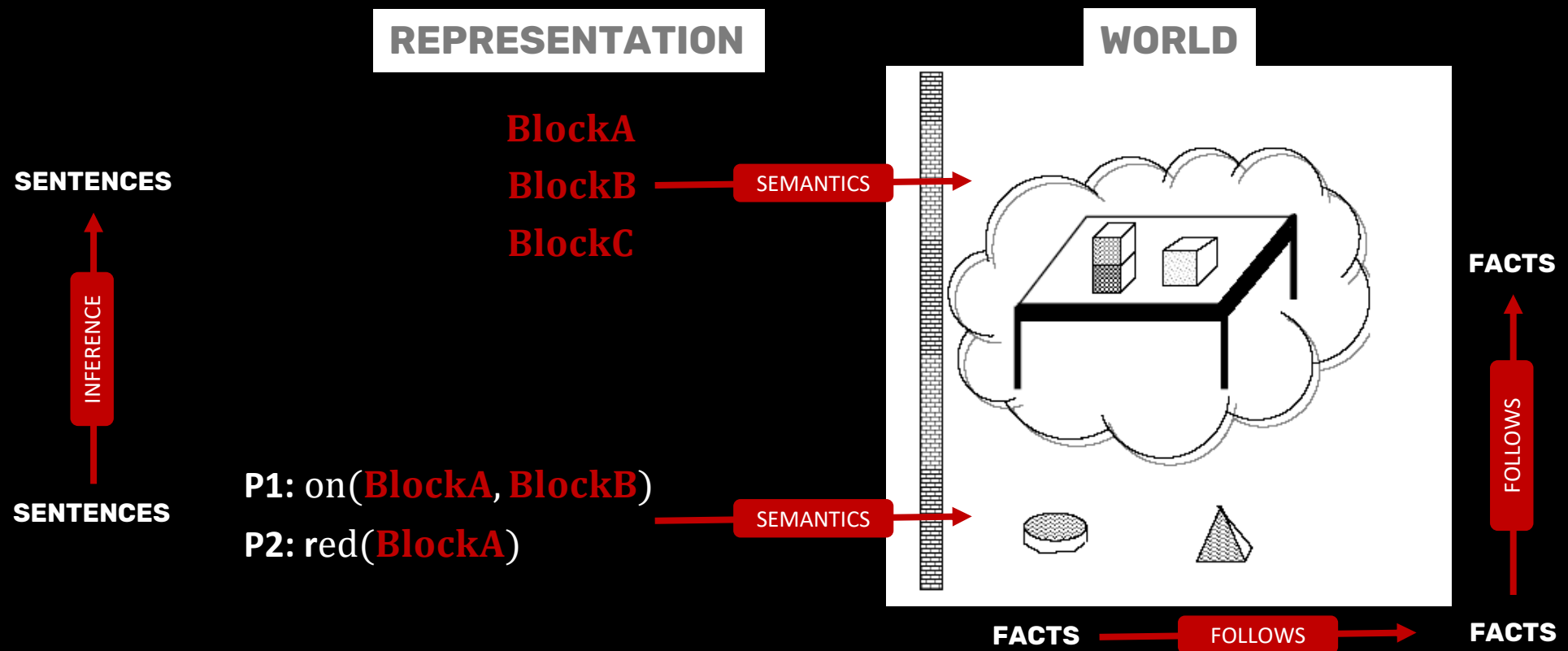
The teacher gave the boys an apple.

- A good **representation language** is:

- Both formal and can express partial information
- Can accommodate inference

The **fundamental elements** of a formal logic are:

- **Syntax**: it is the set of symbols and rules used to express knowledge
- **Semantics**: it specifies the way symbols and sentences relate to the world
- **Inference procedures**: they describe the rules for deriving new sentences (and therefore, new semantics) from existing sentences



SECTION 02

Propositional Logic

- A **sentence** expresses a possible condition of the world
- A sentence can either **true (T)** or **false (F)**
- Most basic sentences are called **simple**, **predicates** or **atomic sentences**
am_wet, is_raining, have_umbrella

- **Predicates:** Symbols or constants **true (T)** or **false (F)**
- **Symbols:** p, q, s, \dots
- Sentences are combined by **connectives** to produce other **sentences**:

\wedge	AND	Conjunction
\vee	OR	Disjunction
\rightarrow	IMPLIES	Implication / conditional
\leftrightarrow	IS EQUIVALENT	Biconditional
\neg	NOT	Negation

A **sentence** is a well-formed formula that can be defined recursively:

- A **symbol** is a proposition
 - If s is a sentence, then $\neg s$ is a sentence
 - If s is a sentence, then (s) is a sentence
 - If s and t are sentences, then $(s \wedge t)$, $(s \vee t)$, $(s \rightarrow t)$ and $(s \leftrightarrow t)$ are sentences
 - A sentence results from a finite number of iterations of the above rules
-
- **Operator precedence:** $\neg \wedge \vee \rightarrow \leftrightarrow$

PROPOSITIONAL LOGIC GRAMMAR

Sentence	\rightarrow	Predicate Proposition
Proposition	\rightarrow	(Sentence)
		Sentence Connective Sentence
		\neg Sentence
Predicate	\rightarrow	true (T) false (F) <i>Symbol</i>
Connective	\rightarrow	$\wedge, \vee, \rightarrow, \leftrightarrow$

- o means **"It is hot"**
- h means **"It is humid"**
- r means **"It is raining"**
- $(o \wedge h) \rightarrow r$ means **"IF it is hot AND humid, THEN it is raining"**
- $q \rightarrow p$ means **"IF it is humid, THEN it is hot"**

A better way to write sentences:

- Hot: **"It is hot"**
- Humid: **"It is humid"**
- Raining: **"It is raining"**

$$(o \wedge h) \rightarrow r \qquad \text{hot} \wedge \text{humid} \rightarrow \text{raining}$$

- The meaning or **semantics** of a sentence determines its **interpretation**
- Given the truth values of all symbols in a sentence, it can be **evaluated** to determine its **truth value** (true or false)
- A **knowledge base** is a list of sentences assumed to be **true**
- A **model** for a KB is a **possible world** (assignment of truth values to propositional symbols) in which each sentence in the KB is **true**
- A sentence p is **satisfiable** if it has some model:
 - $a \wedge b$ is satisfiable
 - $a \wedge \neg a$ is **not** satisfiable

- A **tautology** is a sentence that is **true** under all interpretations, no matter what the world is like or how the semantics are defined. Example: "It's raining or it's not raining."
- A **contradiction** is a sentence that is **false** under all interpretations. The world is never like what it describes, as in "It's raining and it's not raining."

p	$p \wedge \neg p$	$p \vee \neg p$
F	F	T
T	F	T

p	q	$\neg p$	$p \wedge q$	$p \vee q$	$p \rightarrow q$	$p \leftrightarrow q$
F	F	T	F	F	T	T
F	T	T	F	T	T	F
T	F	F	F	T	F	F
T	T	F	T	T	T	T

- $\neg p$ is a **negation** of p : $\neg \text{am_wet}$
- $p \wedge q$ is a **conjunction** of p and q : $\text{am_wet} \wedge \neg \text{is_raining}$
- $p \vee q$ is a **disjunction** of p or q : $\text{am_wet} \vee \neg \text{is_raining}$
- $p \rightarrow q$ is an **implication** of p (premise) implies q (conclusion): $\text{am_wet} \rightarrow \text{is_raining}$
- $p \leftrightarrow q$ is a **biconditional** of p if-and-only-if (iff) q : $\text{am_wet} \leftrightarrow \text{is_raining}$

Simple examples of knowledge base and models:

$\text{is_raining} \rightarrow \text{am_wet}$

model: {
 $\text{is_raining} = \text{true},$
 $\text{am_wet} = \text{true}$ }

$\text{is_raining} \rightarrow \text{am_wet}$

$\text{is_raining} \rightarrow \text{take_umbrella}$

model 1: {
 $\text{is_raining} = \text{false},$
 $\text{am_wet} = \text{true},$
 $\text{take_umbrella} = \text{true}$ }

model 2: {
 $\text{is_raining} = \text{true},$
 $\text{am_wet} = \text{true},$
 $\text{take_umbrella} = \text{true}$ }

$\text{is_raining} \rightarrow \text{am_wet}$

$\text{is_raining} \wedge \text{have_umbrella} \rightarrow \text{open_umbrella}$

model: {
 $\text{is_raining} = \text{true},$
 $\text{am_wet} = \text{true},$
 $\text{have_umbrella} = \text{true},$
 $\text{open_umbrella} = \text{true}$ }

FORM	EQUIVALENCE	NAME
$p \wedge \mathbf{T}, p \vee \mathbf{F}$	p	Identity laws
$p \vee \mathbf{T}$ $p \wedge \mathbf{F}$	\mathbf{T} \mathbf{F}	Domination laws
$p \vee p, p \wedge p$	p	Idempotent laws
$\neg(\neg p)$	p	Double negation law
$p \vee q$ $p \wedge q$	$q \vee p$ $q \wedge p$	Commutative laws
$(p \vee q) \vee r$ $(p \wedge q) \wedge r$	$p \vee (q \vee r)$ $p \wedge (q \wedge r)$	Associative laws
$p \vee (q \wedge r)$ $p \wedge (q \vee r)$	$(p \vee q) \wedge (p \vee r)$ $(p \wedge q) \vee (p \wedge r)$	Distributive laws
$\neg(p \wedge q)$ $\neg(p \vee q)$	$\neg p \vee \neg q$ $\neg p \wedge \neg q$	De Morgan's laws
$p \rightarrow q$	$\neg q \rightarrow \neg p$	Contrapositive equivalence
$p \vee \neg p$	\mathbf{T}	Excluded middle
$p \wedge \neg p$	\mathbf{F}	Negation creates opposite

Let's decide that:

- is_raining means **It's raining outside**
- have_umbrella means **I have an umbrella**
- am_wet means **I am wet**

We can condense propositions replacing them with symbols:

- r means: **It's raining outside**
- u means: **I have an umbrella**
- w means: **I am wet**

EVALUATION

Model: $\{w = \mathbf{T}, r = \mathbf{F}, u = \mathbf{T}\}$

Sentence: $(\neg w \wedge r) \wedge (\neg r \vee u)$

Result of the evaluation: **F**

SATISFACTION

Sentence: $(\neg w \wedge r) \wedge (\neg r \vee u)$

Model: $\{w = ?, r = ?, u = ?\}$

Result of satisfiability:

w	r	u	S
F	F	F	F
F	F	T	F
F	T	F	F
F	T	T	T
T	F	F	F
T	F	T	F
T	T	F	F
T	T	T	F

- KB **P entails Q**, written $p \models q$, means that whenever p is **true**, so is q . In other words, all models of p are also models of q .

q is entailed by p (a set of premises or assumptions) iff there is no logically possible world in which q is **false** while all the premises in p are **true**.

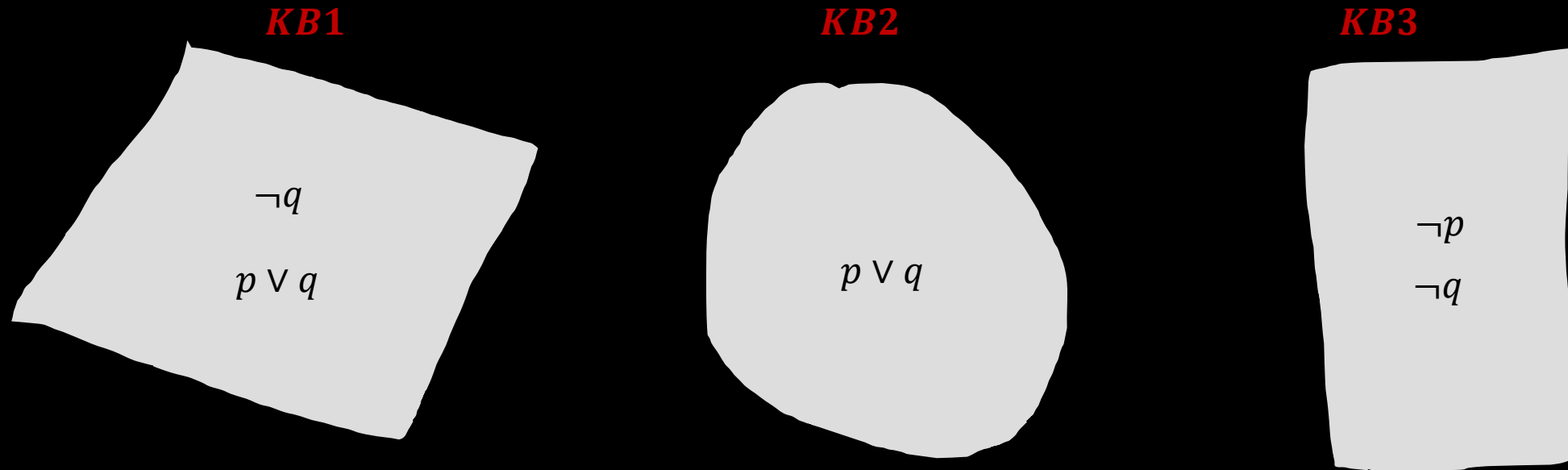
Example:

- RR1 found the victim \models RR1 looked for the victim
- BR2 was turned off \models BR2 is a robot

Entailment is **stronger** than implication because the former imposes a truth value to the latter. Implication does not do that.

If P is a sentence and K is a knowledge base, a fair question to ask is:
does K entail P ?

Or, if more than one knowledge base is available:



Proof by refutation is a complete inference procedure that is used to prove entailment. It tries to prove something demonstrating the opposite:

- It uses a single inference rule or **resolution**
- The knowledge base is represented in a **Conjunctive Normal Form** (or **CNF**)
- It **reduces** inference to the problem of checking satisfiability

A sentence is in **conjunctive normal form** (or CNF) if it is a **conjunction of disjunction terms**:

- Examples: $\boxed{q} \wedge \boxed{(r \vee s)}$
 \boxed{p}
 $\boxed{(p \vee \neg q)} \wedge \boxed{(r \vee s)}$
 $\boxed{(p \vee \neg q)}$
 $\boxed{(p \vee q \vee r)}$

- There is a way to convert a sentence into a clausal form:

SENTENCE	CLAUSAL FORM
$p \leftrightarrow q$	$p \rightarrow q$ $q \rightarrow p$
$p \rightarrow q$	$\neg p \vee q$
$\neg \neg p$	p
$\neg(p \vee q)$	$\neg p \wedge \neg q$
$\neg(p \wedge q)$	$\neg p \vee \neg q$

- The process is performed **conjoining** the *KB* with the negation of the query
- It's possible to say that *KB* entails the query iff **the truth table returns false in every row**

- Example:
Given the following *KB*:

1. $w \vee \neg u$
2. $\neg u \wedge r \rightarrow w$

S: $r \wedge \neg w$

Does ***KB* entail *S*** (or ***KB* \models *S***) ?

$$(w \vee \neg u) \wedge (\neg u \wedge r \rightarrow w) \wedge \neg(r \wedge \neg w)$$

$$(w \vee \neg u) \wedge ((\neg u \wedge r) \rightarrow w) \wedge (\neg r \vee \neg \neg w)$$

$$(w \vee \neg u) \wedge (\neg(\neg u \wedge r) \vee w) \wedge (\neg r \vee w)$$

$$(w \vee \neg u) \wedge (u \vee \neg r \vee w) \wedge (\neg r \vee w)$$

$$(w \vee \neg u) \wedge (u \vee \neg r \vee w) \wedge (\neg r \vee w)$$

<i>r</i>	<i>u</i>	<i>w</i>	<i>S</i>
F	F	F	T
F	F	T	T
F	T	F	F
F	T	T	T
T	F	F	F
T	F	T	T
T	T	F	F
T	T	T	T

PRACTICE

Exercises from the textbook (chapter 7):
7.1, 7.4, 7.5, 7.7, 7.10

QUESTIONS ?

ARTIFICIAL INTELLIGENCE COMP 131

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